# **Assignment 2B**

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Problem Description:  $\exists$  two sets  $a, b \in \mathbb{Z} \{1..n\} \mid \exists$  a set  $c \in \mathbb{Z} \{1..n\} \mid |c| \le |a|, |b| \mid c = a \cap b$ 

## Algorithm: Array Intersection

input: 1-Based arrays arr\_a and arr\_b of positive integers. denote the size of arr\_a by |arr\_a| and similarly for b

output: 1-Based array arr\_c of positive integers, such that arr\_c contains exactly one copy of each integer that appears somewhere in both arr\_a arr\_b. Thus,  $|arr_c| \le min(|arr_a|, |arr_b|)$ .

input: 1-Based array arr temp and a integer num

output: true if arr\_temp contains num and false if arr\_temp does not contain num

```
def include ← arr_temp, num
    if |arr_temp| < 1 then ↑ false
    for k ← 1..|arr_temp|
        if arr_temp[k] = num then ↑ true
    ↑ false</pre>
```

#### Correctness

• Returning the Correct Value

The algorithm checks all the elements of arr\_b for arr\_a and if the elements are the same and not already in the array that holds onto elements shared in arr\_a and arr\_b then that value is added to the array that holds onto elements shared in arr\_a and arr\_b. These requirements are all met and arr\_c is not appended additional copies of times when  $arr_a[i] == arr_b[j]$  because of the include function.

Halting

the loops are the same size as arr\_a and arr\_b or smaller for arr\_c therefore the loops

will end resulting in a return statement which guarantees that as long as arr\_a and arr\_b are finite the algorithm will halt.

#### Worst Case

 $n = |arr_a|$ ,  $m = |arr_b|$ , o = the size of a growing array that doesn't exceed the length of n or m, whichever is smaller.

This simplifies to  $O(mn^2)$  or  $O(nm^2)$ , for whichever of n or m is smaller.

O(n) because the outer loop goes n times, and O(m) because the inner loop goes m times and O(n) or O(m) because that loop runs through an array no larger than the smaller of n or m.

# Generalization to allow for any number of arrays

```
\forall sets \in \mathbb{Z} \{1..n\} \exists a number n \in \mathbb{Z} \{1..n\} of sets m \in \mathbb{Z} \{1..n\} \mid \exists a set c \in \mathbb{Z} \{1..n\} \mid \forall sets m, |c| \leq |\text{the smallest array in } m \mid |c| = |\text{the intersection of sets in } m
```

### Algorithm 2 for Generalization

input: an array of arrays of integers

outout: an array of integers shared by the arrays in the input such that no integer appears more than once in the output array.

input: 1-Based arrays arr\_a and arr\_b of positive integers. denote the size of arr\_a by |arr\_a| and similarly for b

output: 1-Based array arr\_c of positive integers, such that arr\_c contains exactly one copy of each integer that appears somewhere in both arr\_a arr\_b. Thus,  $|arr_c| \le min(|arr_a|, |arr_b|)$ .

input: 1-Based array arr\_temp and a integer num

output: true if arr\_temp contains num and false if arr\_temp does not contain num

```
def include \( \text{ arr_temp, num} \)
    if |arr_temp| < 1 then \( \text{ false} \)
    for \( k \lefta 1.. |arr_temp| \)
        if \( \arr_temp[k] = \text{ num then } \) true

\( \text{ false} \)
</pre>
```

## Correctness Algorithm 2

• Returning the Correct Value

I built the second algorithm to use the previous one, so if the previous one is assumed to be correct then those parts in the n array intersect are still correct. The n array intersect uses a similar strategy and does an n² number of calls to the previous algorithm with an additional n. The algorithm for n array intersect returns the correct value as it checks all possible sets of different array combinations from the array of arrays and only adds to the array intersect array when that array does not contain the value shared by the current arr\_all[i] and arr\_all[j]. therefore all values in the returned array are unique and all possible combinations of arrays are checked, thus the intersection of the arrays is correctly computed.

Halting

Like the previous algorithm this one operates on a finite array of finite arrays and so the loops will end and the return will be called halting the algorithm.

# Implementation

Implementation in Ruby (for both algorithms) with Testing

```
#! /usr/bin/env ruby
# Author: Isaac Archer
# Description:
# test for an algorithm to find the array
```

```
# intersect of two arrays of length n and m
# input: two arrays of integers
# output: array of integers of the
# intersect of the arrays
def array_intersect(arr_a, arr_b)
      0.upto(arr_a.length - 1) do |i|
            0.upto(arr_b.length - 1) do |j|
                  if arr a[i] == arr b[j]
                         if arr c.include? arr a[i]
                         else
                  end
            end
      end
      return arr c
end
# input: array of arrays of integers
# output array of integers that is the
# intersect of the arrays in the input
def narray_intersect(arr_all)
      0.upto(len all) do |i|
            O.upto(len all) do |j|
                  if i == j
                  else
                         temp_intersect.each do |k|
                               if arr_c.include? k
                                     next
                               else
                         end
            end
      end
      return arr c
end
def test arr all int(num arr, arr len)
      0.upto(num_arr - 1) do |1|
             \hbox{\tt 0.upto(arr\_len - 1) \{ |m| temp\_arr.push(rand(arr\_len) + 1) } \} \\
      end
```

```
ll.each αο |a|
            print "ARRAY #{i}: "
            puts a.to_s
      end
      print "INTERSECT ARRAY: "
      puts arr_c.to_s
end
def test arr int(len a, len b)
       0. upto(len_a - 1) \ \{ \ |i| \ arr_a.push(rand(len_a) + 1) \ \} 
       0. upto(len_b - 1) \ \{ \ |i| \ arr_b.push(rand(len_b) + 1) \ \} 
      print "ARRAY A: "
      puts arr a.to s
      print "ARRAY B: "
      puts arr_b.to_s
      print "INTERSECT ARRAY: "
      puts arr_c.to_s
end
def main()
      if ARGV.length != 4
           puts "invalid arguments"
      end
end
```

The algorithm slows down considerably as the input size increases above small values like 100 or 1000.

- n and m size 100 completed in 0.044s
- n and m size 1000 completed in 0.204s

- n and m size 10000 completed in 17.774s
- n and m size 100000 completed in a really long time (significantly over an hour)