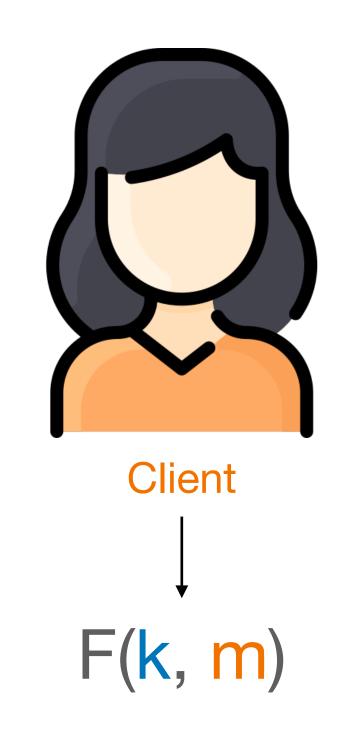
# A Post-Quantum Oblivious PRF from Isogenies

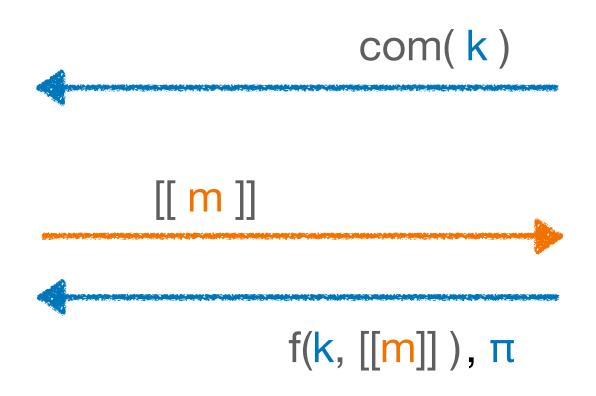
**Andrea Basso** 





# Oblivious PRF

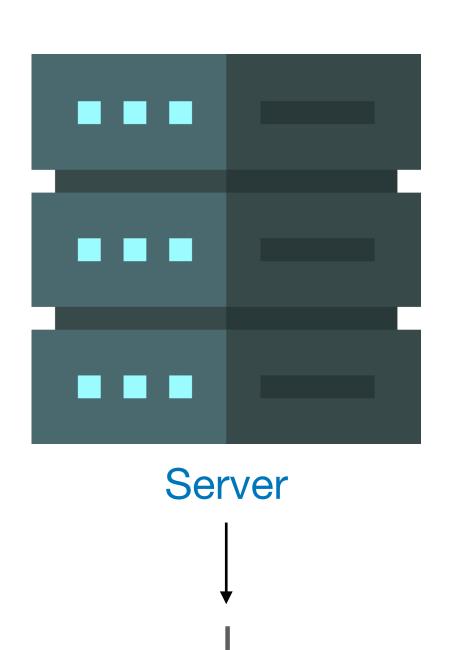




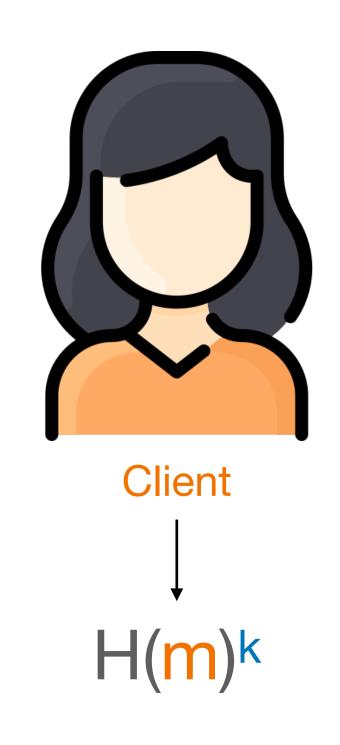
- Password-checking in Microsoft Edge
  - •OPAQUE

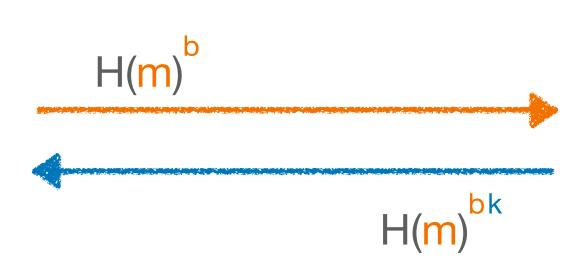
• ,,,,,

- Privacy pass
- Private-set intersection
  Adaptive OT

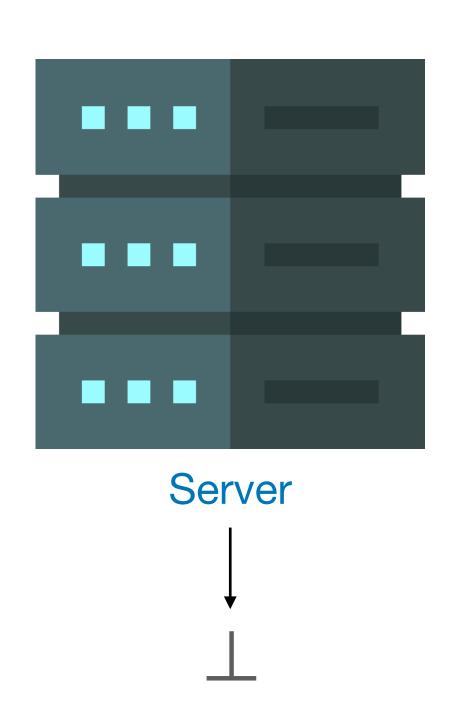


# HashDH OPRF





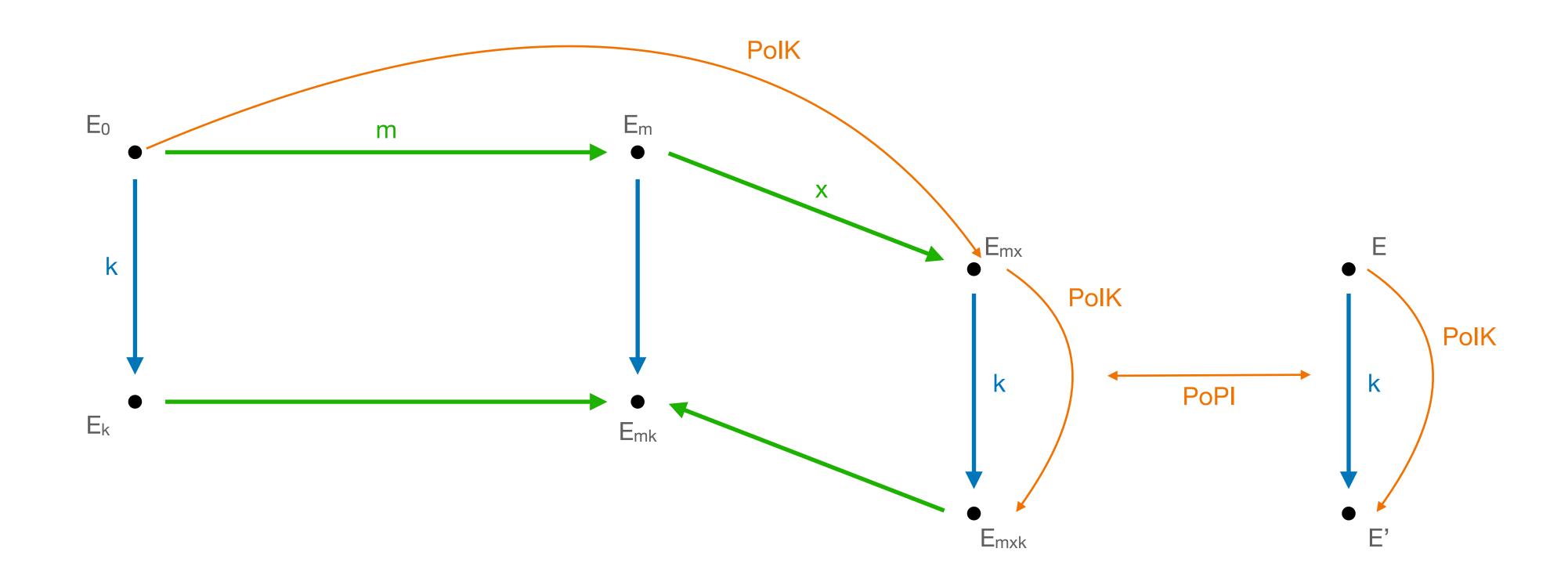
- Server doesn't learn
  anything
  Output is
  deterministic
  deterministic
- Client only learns one output -



# Post-quantum OPRFs

 Generic MPC techniques many rounds (≥ 5) round optimal VOPRF based on lattices [ADDS19] • feasibility result (> 2<sup>40</sup> bits of comms) six rounds VOPRF based on SIDH [BKW20] broken by attack on PR and on SIDH three rounds (OT required) OPRF based on CSIDH [BKW20] CSIDH parameters?

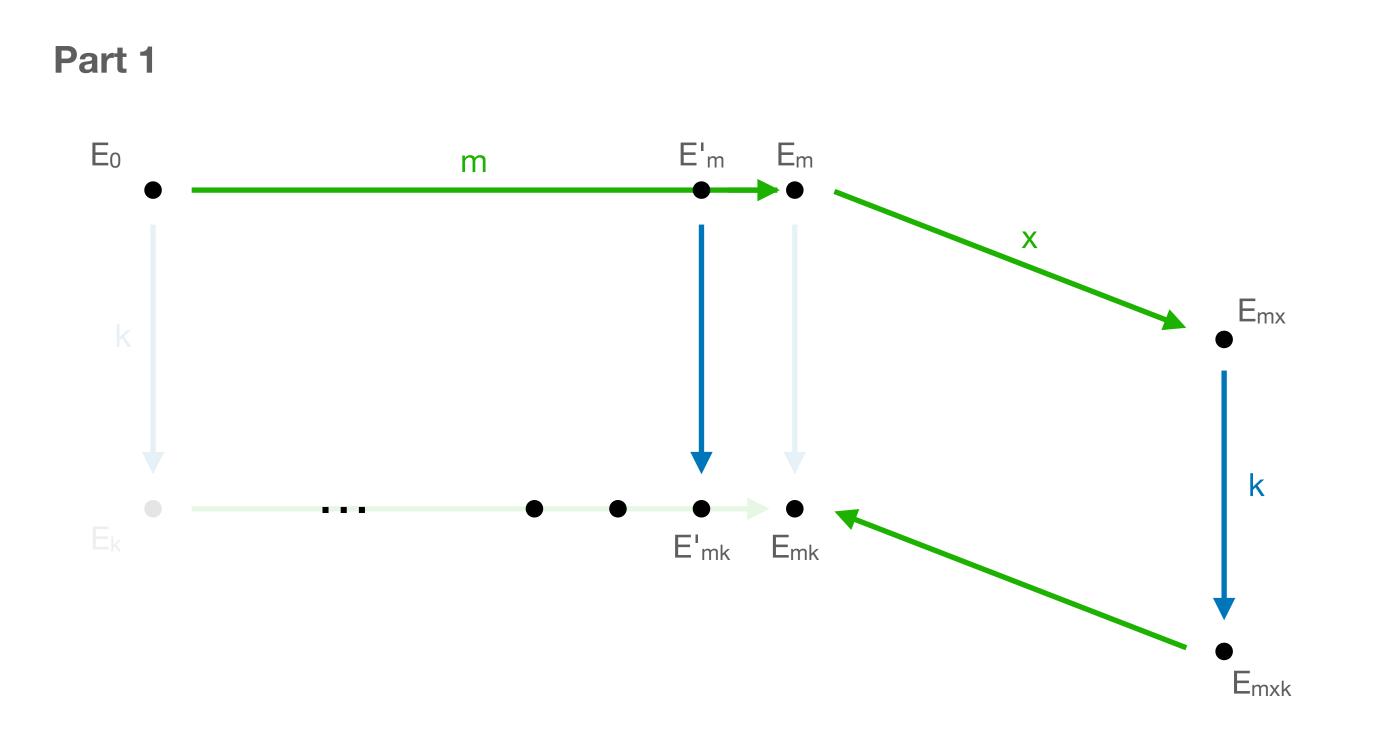
# The original OPRF [BKW20]



$$F(k, m) = H(m, j_{mk}, E')$$

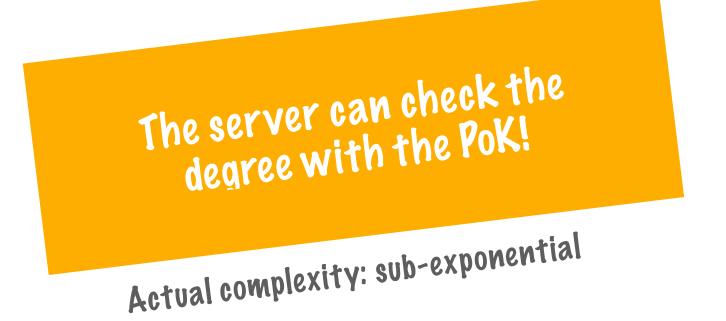
# Breaking pseudorandomness [BKMPS21]

Pseudorandomness: after n interactions, an attacker cannot generate n+1 PRF outputs



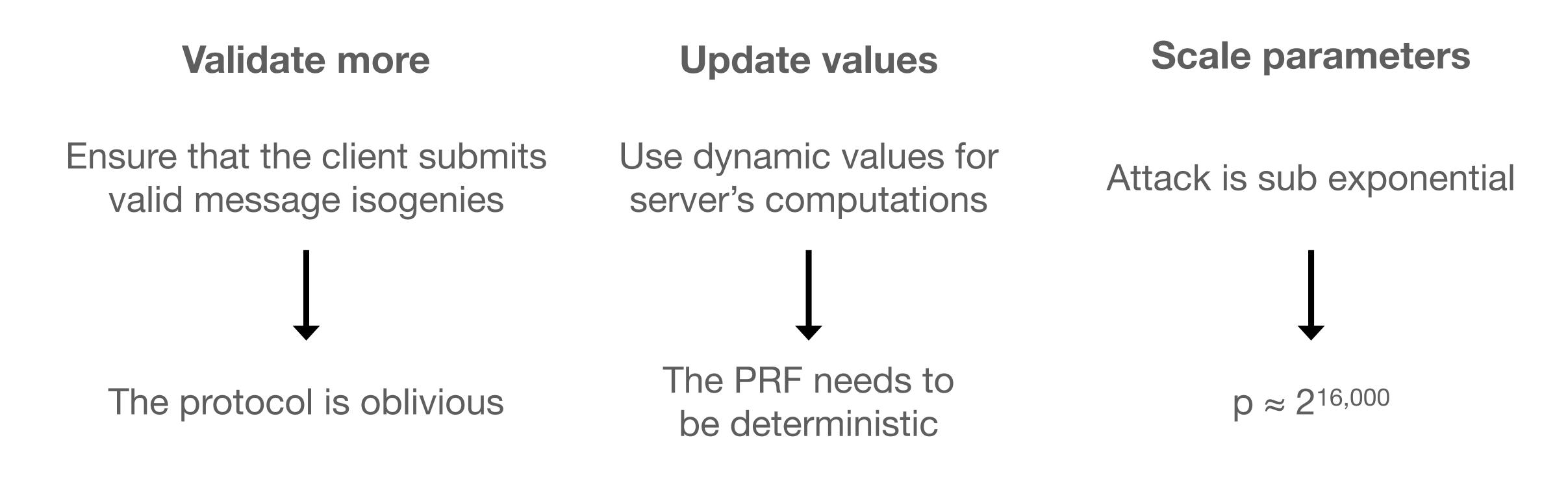
#### Part 2

- Repeat the attack 3 times
- Find a basis on E<sub>k</sub>
- Evaluate the PRF on any message



## Countermeasures?

It seems hard to prevent an attacker from recovering a basis on Ek



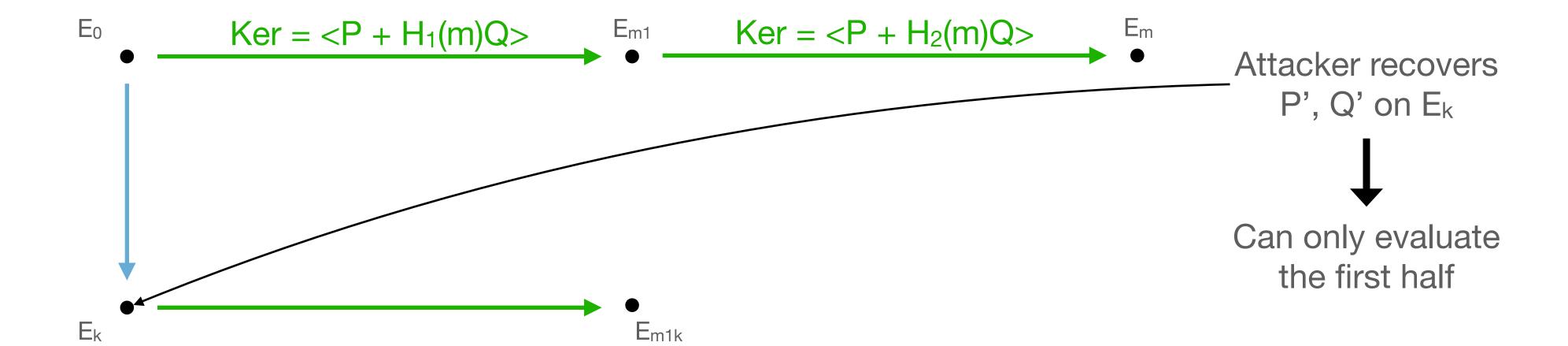
Idea: make the basis on Ek not enough for an attack

## An efficient countermeasure

[BKM20]

 $Ker = \langle P + H(m)Q \rangle$   $E_m$ Attacker recovers  $P', Q' \text{ on } E_k$  Can evaluate the PRFon any message

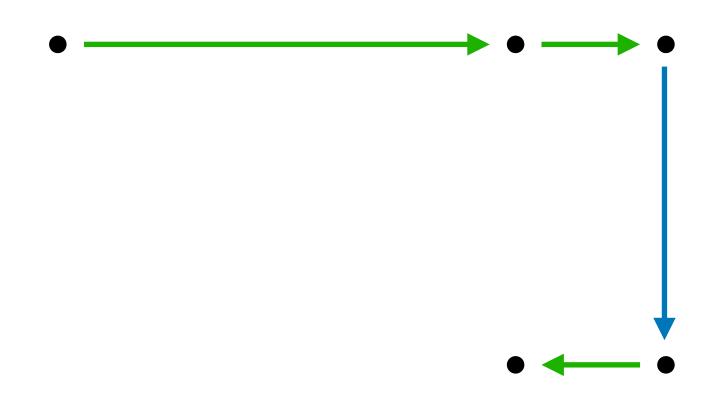
Our countermeasure



# Preventing the SIDH attacks

#### First attempt

SIDH attacks requires N torsion to recover a N<sup>2</sup> - degree isogeny



not really an OPRF

#### **SIDH** countermeasures

only works for one party

• Longer isogenies

Masked-degree isogenies [Mor22]
 hard to build proofs

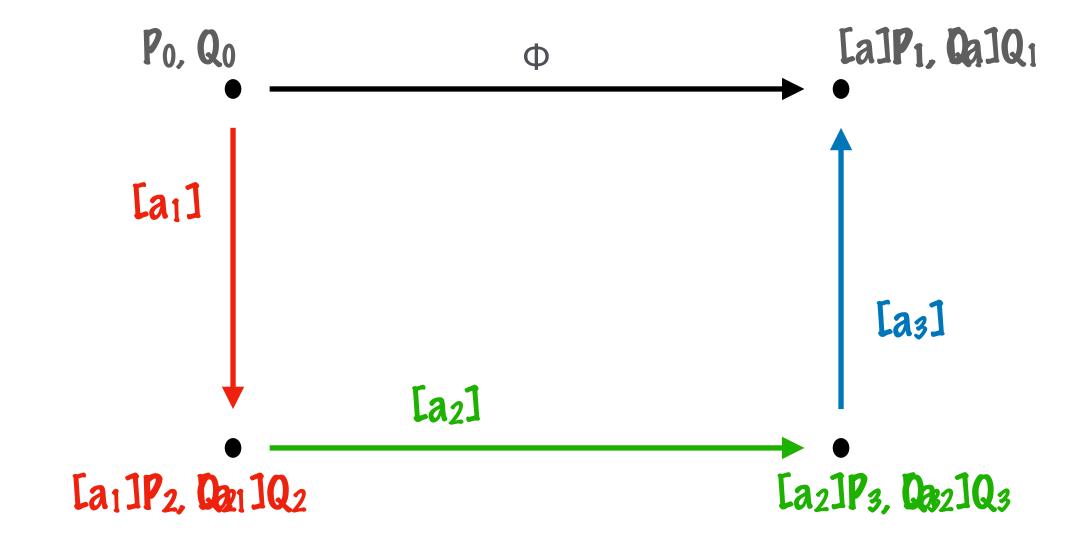
Masked torsion points [Fou22]

it works

large prime

needs new PolK

## PolK with masked torsion



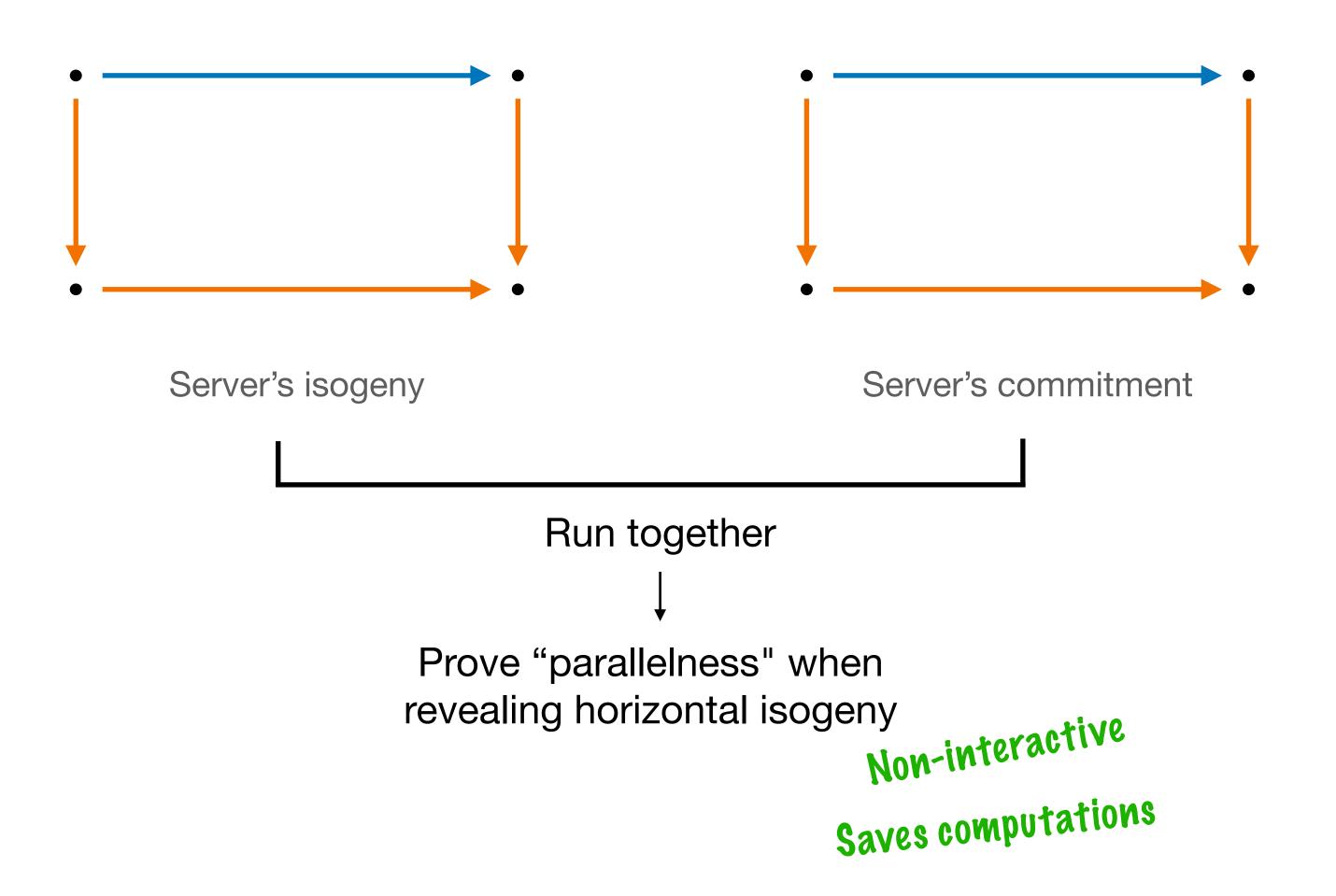
challenges from {-1, 0, 1}

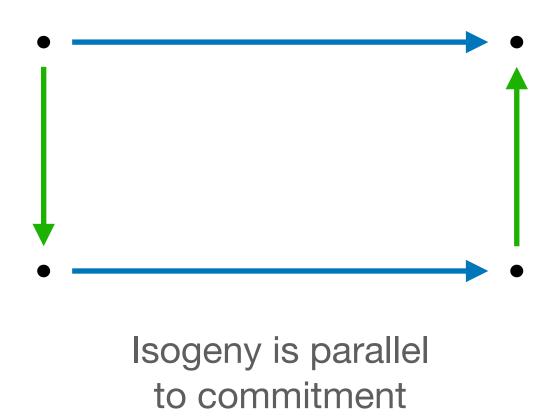
soundness error = 2/3 $\rightarrow$  need  $1.7\lambda$  repetitions

 $a = a_1 \times a_2 \times a_3$ 

# Verifiability

[BKW20] uses 3 proofs:





Interactive (5 rounds)

# Putting it all together

- Pseudorandomness countermeasure more efficient than original
- SIDH counteremeasures would require p ≈ 26000
- New SIDH proof requires a larger prime requires p ≈ 29000
- New Popi more efficient than original round optimal

Protocol	Rounds	Bandwidth (avg.)	Verifiable	Secure
[ADDS21] (LWE)	2	>128 GB	✓	✓
[BKW20] (SIDH)	6	1.4 MB		X
[BKW20] (CSIDH)	3	$424~\mathrm{kB}$	X	
[This work]	2	1.9 MB		