

# Zip-Tries: Simple Dynamic Data Structures for Strings

David Eppstein, *Ofek Gila*, Michael T. Goodrich, and Ryuto Kitagawa

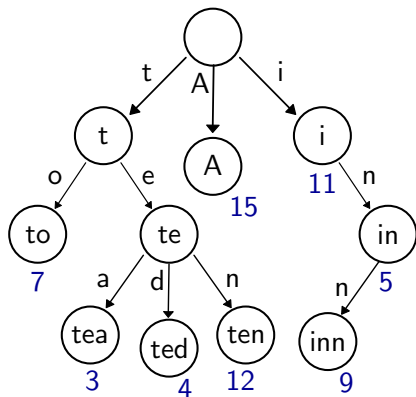
University of California, Irvine

ACDA, 2025



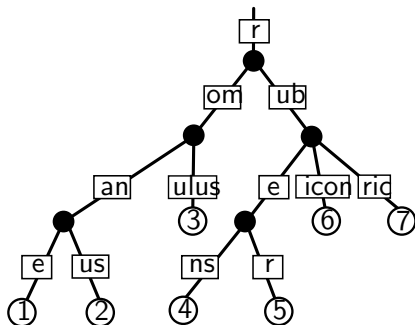
# Selected History I

- 1912 / 1959 / 1960 – Trie (Thue, de la Braindais, Fredkin) [3, 5]
  - Query time:  $\Theta(k)$



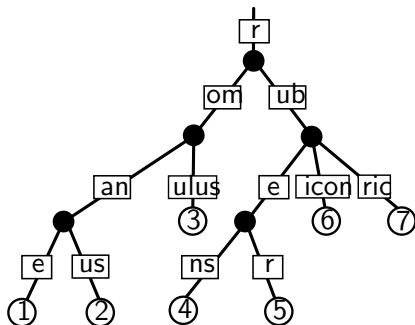
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- 1968 – Compressed trie / radix tree (Morrison, Gwehenberger) [8]
  - Query time:  $\mathcal{O}(k)$



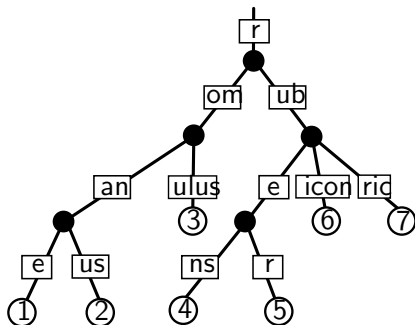
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- 2010+ – Dynamic z-fast tries, packed compact tries, etc. [1, 10]
  - Query time:  $\mathcal{O}(k/w + \log k)$



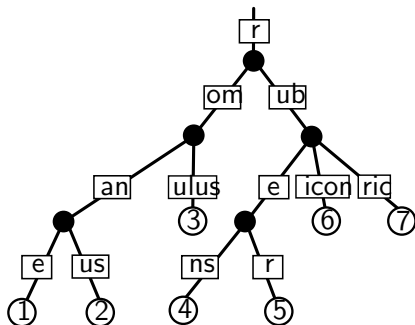
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  - Very fast, very complex



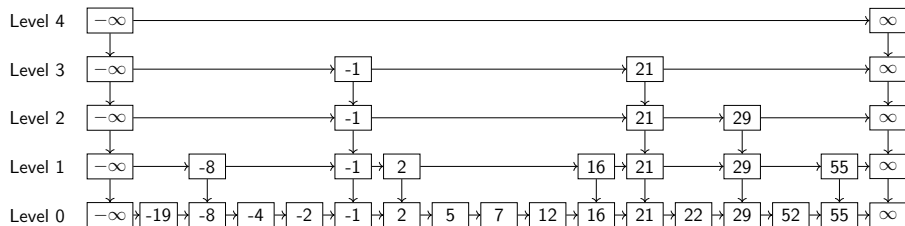
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  - Query time:  $\mathcal{O}(k/w + \log k)$
  - Very fast, very complex
  - Lots of branching,  $\mathcal{O}(k/w)$



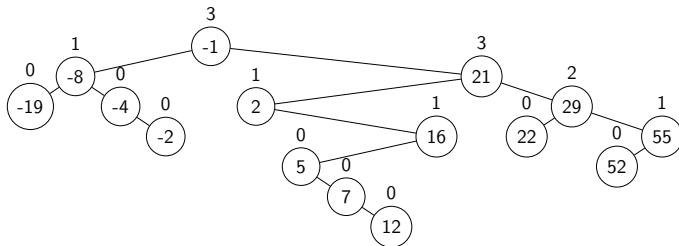
# Selected History II

- Meanwhile...
- 1989 – Skip list (Pugh) [9]



# Selected History II

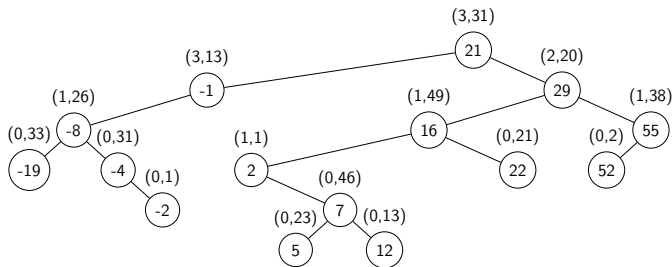
- Meanwhile...
- 1989 – Skip list (Pugh) [9]
- 2018 – Zip tree (Tarjan, Levy, Timmel) [11]
  - Flat-out better than skip lists
  - Only  $\mathcal{O}(\log \log n)$  bits metadata per node





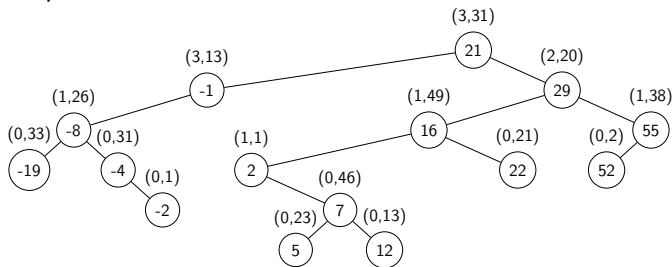
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- 2023 – Zip-zip tree (Gila, Goodrich, Tarjan) [6]
  - Flat-out better than zip trees



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    - Flat-out better than zip trees
- ✓ Simple, practical, and efficient

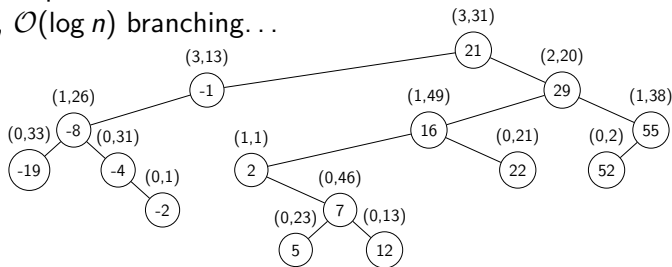


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✓ Simple, practical, and efficient

✓ Low,  $\mathcal{O}(\log n)$  branching...



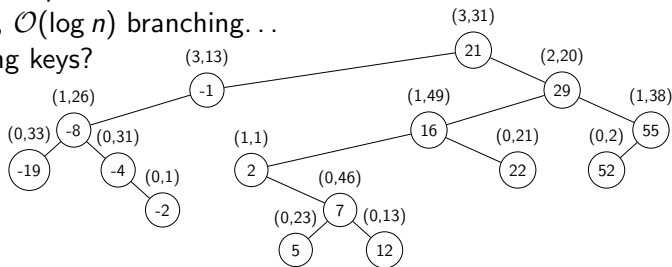
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- Meanwhile...
- 1989 – Skip list (Pugh) [9]
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✓ Simple, practical, and efficient

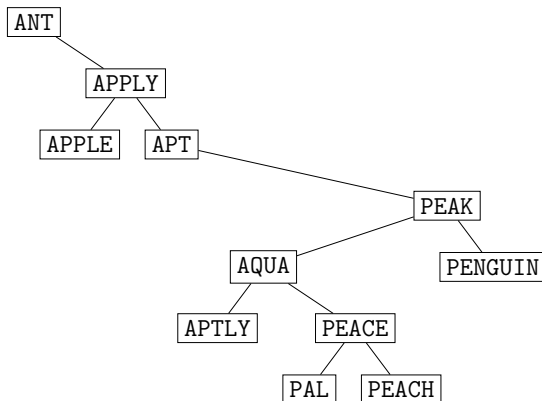
✓ Low,  $\mathcal{O}(\log n)$  branching...

? String keys?



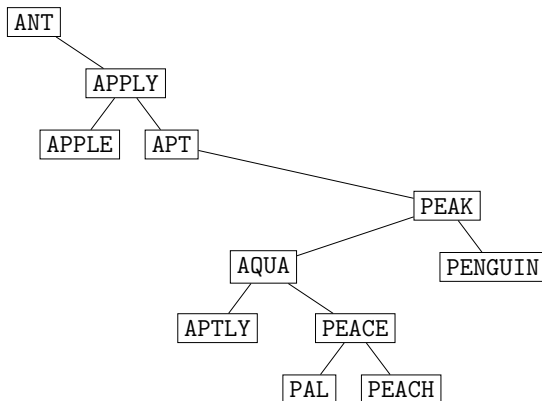
# Selected History – Bookend Search I

- Many structures for 1D keys...



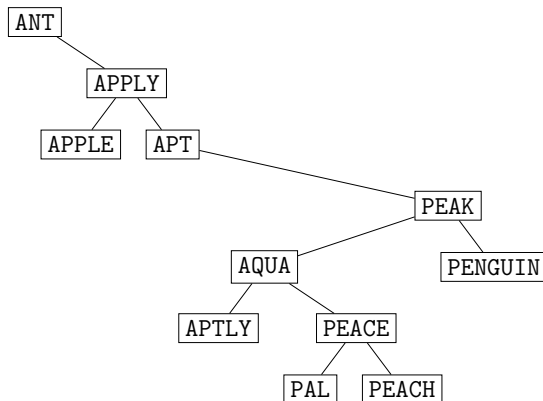
# Selected History – Bookend Search I

- Many structures for 1D keys. . .  $k$ -D keys?



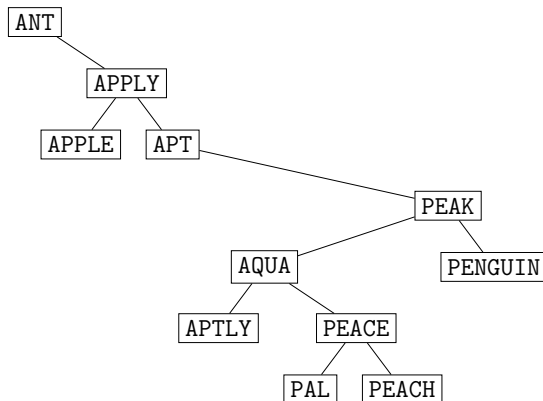
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- Many structures for 1D keys. . .  $k$ -D keys?
- Naively:  $\mathcal{O}(\mathcal{B}(n))$  search  $\rightarrow \mathcal{O}(k \times \mathcal{B}(n))$



# Selected History – Bookend Search I

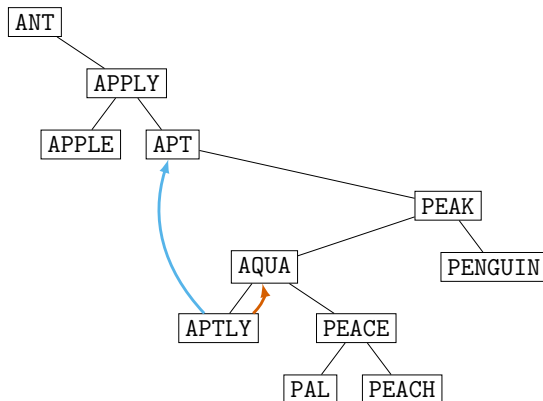
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- Grossi and Italiano [7] (2002),  $\mathcal{O}(k + \mathcal{B}(n))$





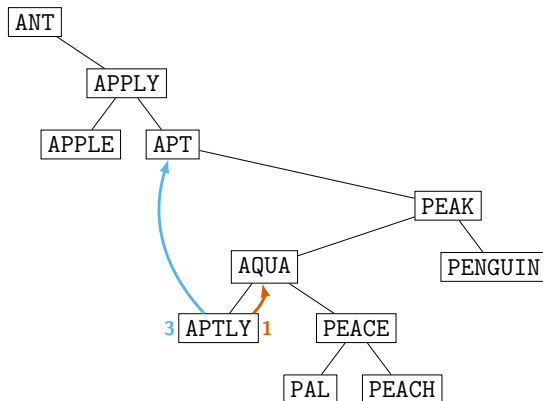
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- Grossi and Italiano [7] (2002),  $\mathcal{O}(k + \mathcal{B}(n))$ 
  - Pointers to lexical predecessor and successor



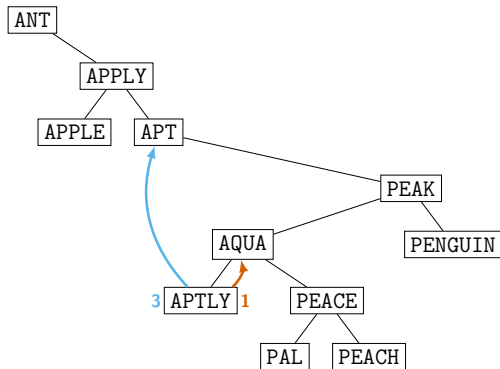
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- Many structures for 1D keys. . .  $k$ -D keys?
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- Grossi and Italiano [7] (2002),  $\mathcal{O}(k + \mathcal{B}(n))$ 
  - Pointers to lexical predecessor and successor
  - *Longest Common Prefix* (LCP)



# Selected History – Bookend Search II

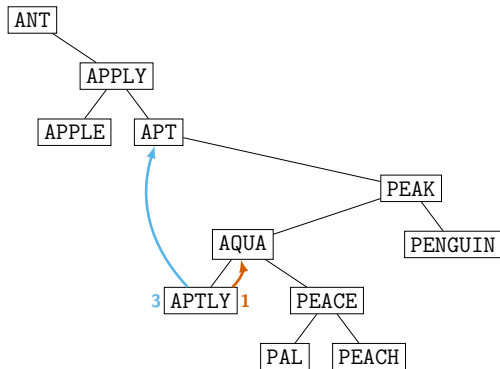
- Avoid unnecessary comparisons,  $\mathcal{O}(k + \mathcal{B}(n))$



# Selected History – Bookend Search II

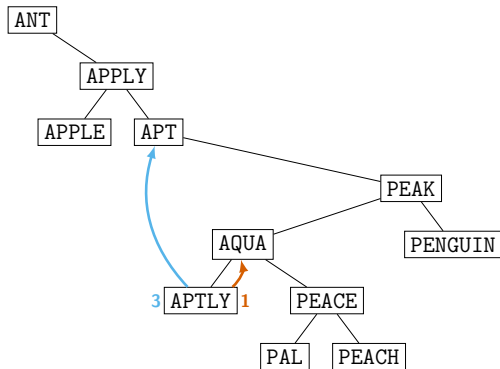
- Avoid unnecessary comparisons,  $\mathcal{O}(k + \mathcal{B}(n))$

✗ Not 'best' in theory



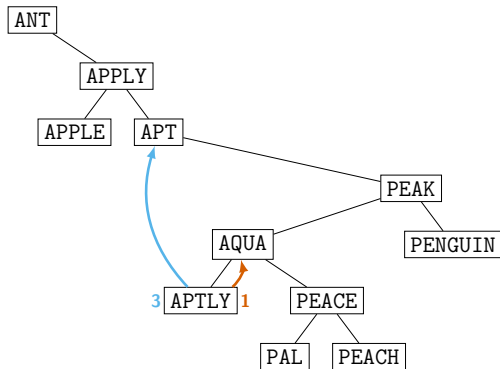
# Selected History – Bookend Search II

- Avoid unnecessary comparisons,  $\mathcal{O}(k + \mathcal{B}(n))$
- ✗ Not 'best' in theory
- ✓ Simple, practical, & efficient [2]



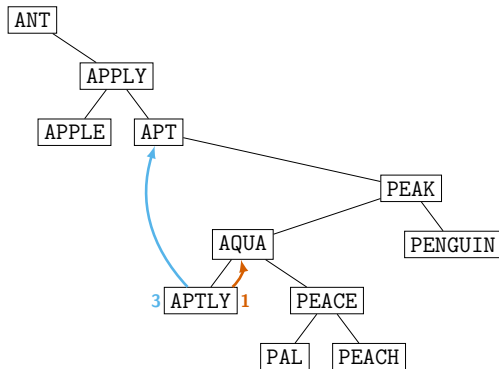
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- Avoid unnecessary comparisons,  $\mathcal{O}(k + \mathcal{B}(n))$
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- ✓ Simple, practical, & efficient [2]
- ✓ Branching can be low,  $\mathcal{O}(\log n)$  or  $\mathcal{O}\left(\frac{\log n}{\log \log n}\right)$  [4]



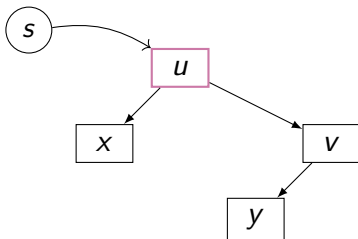
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- ❓ Parallel comparisons?



# Sequential Longest Common Prefix

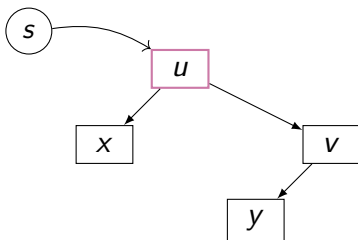
- Best way to find LCP?
- $k = |s| = 38$
- # computations: 0
- # branching: 0





# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 0
- # branching: 0

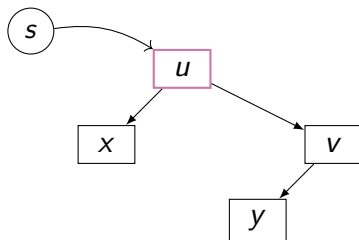


$s = \text{GGCTAGCTACAACGAGCTGGGGCAGCAGCTAGTAGTAGG}$

$u = \text{GGCTAACTACAACGAGCTGGCGTTGTGAGCTAC}$

# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 1
- # branching: 0

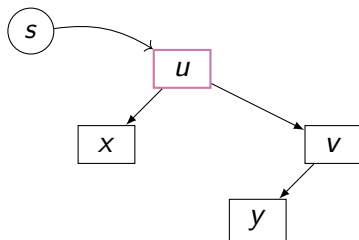


$s = \text{GGCTAGCTACAACGAGCTGGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTAACTACAACGAGCTGGCGTTGTGAGCTAC}$

GG

# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 2
- # branching: 0

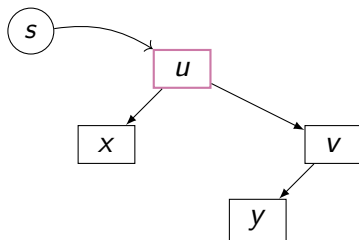


$s = \text{GGCTAGCTACAACGAGCTGGGCGAGCTAGTAGTAGG}$   
 $u = \text{GGCTAACTACAACGAGCTGGCGTTGTGAGCTAC}$

$\text{GG}$

# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 3
- # branching: 0

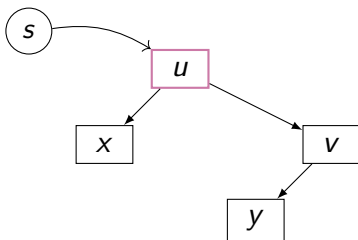


$s = \text{GGCTAGCTACAACGAGCTGGGCGAGCTAGTAGTAGG}$   
 $u = \text{GGCTAACTACAACGAGCTGGCGTTGTGAGCTAC}$

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# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 4
- # branching: 0

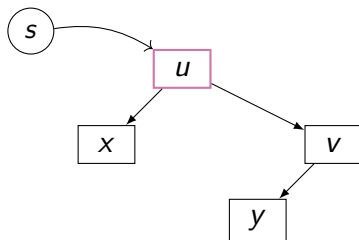


$s = \text{GGCTAGCTACAACGAGCTGGGCGAGCTAGTAGTAGG}$   
 $u = \text{GGCTAACTACAACGAGCTGGCGTTGTGAGCTAC}$

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# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 5
- # branching: 0

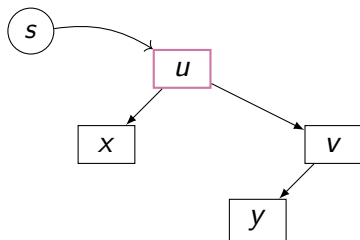


$s = \text{GGCTAGCTACAACGAGCTGGGCGAGCTAGTAGTAGG}$   
 $u = \text{GGCTACTACAACGAGCTGGCGTTGTGAGCTAC}$

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# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 6
- # branching: 0

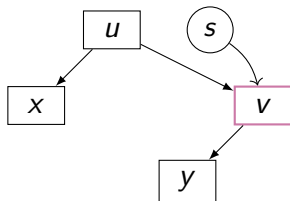


$s = \text{GGCTAGCTACAACGAGCTGGGCGAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$

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- Best way to find LCP?
- $k = |s| = 38$
- # computations: 6
- # branching: 1

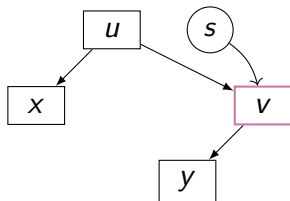


$s = \text{GGCTA} \text{GCTACAACGAGCTGGGCGAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTA} \text{GCTACAACGAGCTGGTCAGTAACGAGCTGGG}$



# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 7
- # branching: 1

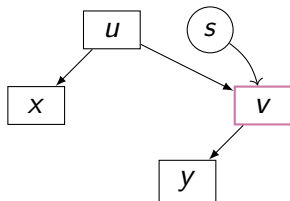


$s = \text{GGCTAGCTACAACGAGCTGGGCGAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGGTCAGTAACGAGCTGGG}$

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# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 8
- # branching: 1

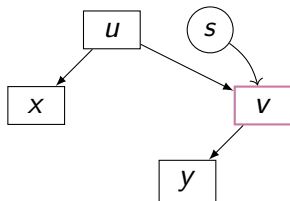


$s = \text{GGCTAGCTACAAACGAGCTGGGCGAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \textcolor{brown}{\text{A}} \text{CTACAAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAAACGAGCTGGTCAGTAACGAGCTGGG}$

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# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 9
- # branching: 1

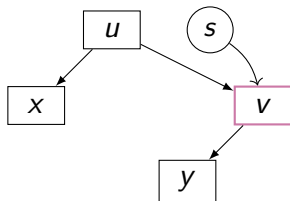


$s = \text{GGCTAGCTACAAACGAGCTGGGCGAGCTAGTAGTAGG}$   
 $u = \text{GGCTA}^{\text{orange}}\text{CTACAAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAAACGAGCTGGTCAGTAACGAGCTGGG}$

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# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 10
- # branching: 1

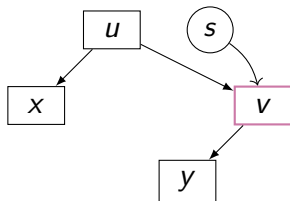


$s = \text{GGCTAGCTA} \text{CAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTA} \text{CAACGAGCTGGTCAGTAACGAGCTGGG}$

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# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 11
- # branching: 1

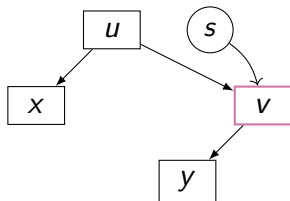


$s = \text{GGCTAGCTAC} \text{AACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTAC} \text{AACGAGCTGGTCAGTAACGAGCTGGG}$

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# Sequential Longest Common Prefix

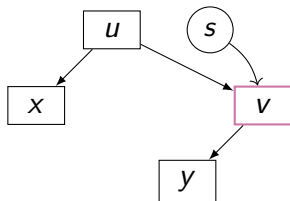
- Best way to find LCP?
- $k = |s| = 38$
- # computations: 12
- # branching: 1



$s = \text{GGCTAGCTAC} \text{ACGAGCTGGGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTAC} \text{ACGAGCTGGTCAGTAACGAGCTGGG}$

# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 13
- # branching: 1

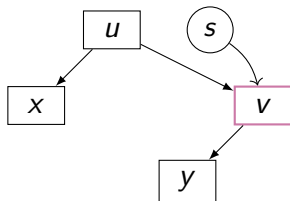


$s = \text{GGCTAGCTACAA} \text{CGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAA} \text{CGAGCTGGTCAGTAACGAGCTGGG}$

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# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 14
- # branching: 1



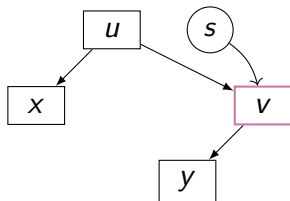
$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA}^{\text{orange}}\text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGGTCAGTAACGAGCTGGG}$

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# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 15
- # branching: 1

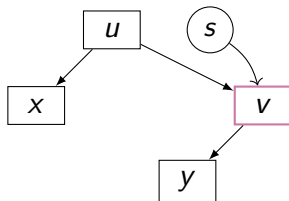


$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGGTCAGTAACGAGCTGGG}$

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# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 16
- # branching: 1

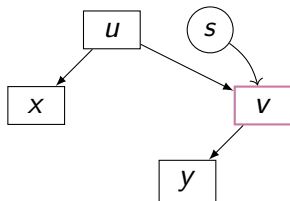


$s = \text{GGCTAGCTACAACGA} \text{GCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGA} \text{GCTGGTCAGTAACGAGCTGGG}$

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# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 17
- # branching: 1

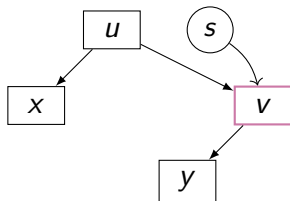


$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGGTCAGTAACGAGCTGGG}$

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# Sequential Longest Common Prefix

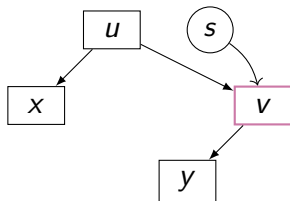
- Best way to find LCP?
- $k = |s| = 38$
- # computations: 18
- # branching: 1



$s = \text{GGCTAGCTACAACGAGCTGGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGGTCAGTAACGAGCTGGG}$

# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 19
- # branching: 1

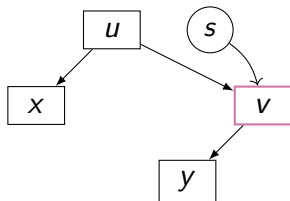


$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGGTCAGTAACGAGCTGGG}$

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# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 20
- # branching: 1

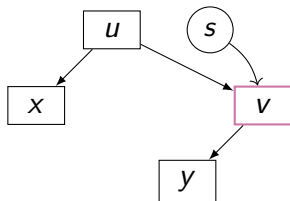


$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGGTCAGTAACGAGCTGGG}$

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# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 21
- # branching: 1

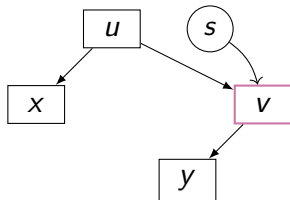


$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGG} \text{TCAGTAACGAGCTGGG}$

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# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 23
- # branching: 1



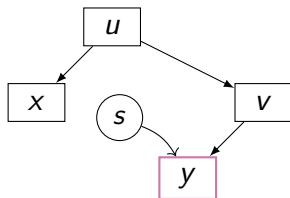
$s = \text{GGCTAGCTACAACGAGCTGGG} \text{CAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGG} \text{CAGTAACGAGCTGGG}$

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# Sequential Longest Common Prefix

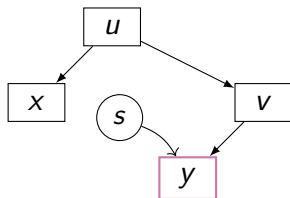
- Best way to find LCP?
- $k = |s| = 38$
- # computations: 23
- # branching: 2



$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGG} \text{CAGTAACGAGCTGGG}$   
 $y = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$

# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 23
- # branching: 2

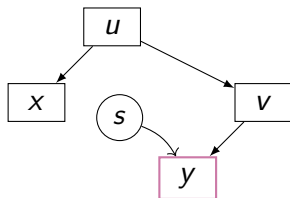


$s = \text{GGCTAGCTACAACGAGCTGGG} \text{CAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGG} \text{CAGTAACGAGCTGGG}$   
 $y = \text{GGCTAGCTACAACGAGCTGGG} \text{CAGCAGCTAGTAGTAGG}$

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# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 24
- # branching: 2

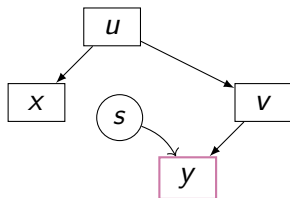


$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGG} \text{CAGTAACGAGCTGGG}$   
 $y = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$

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# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 25
- # branching: 2

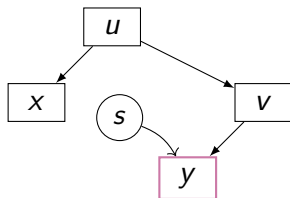


$s = \text{GGCTAGCTACAACGAGCTGGGCA} \text{GCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGG} \text{CAGTAACGAGCTGGG}$   
 $y = \text{GGCTAGCTACAACGAGCTGGGCA} \text{GCAGCTAGTAGTAGG}$

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# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 26
- # branching: 2

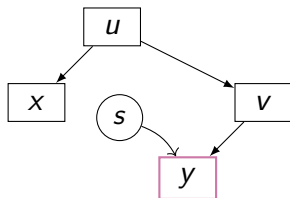


$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA}^{\text{A}}\text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGG}^{\text{T}}\text{CAGTAACGAGCTGGG}$   
 $y = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$

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# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 27
- # branching: 2

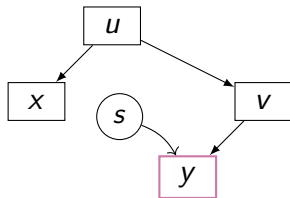


$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGG} \text{CAGTAACGAGCTGGG}$   
 $y = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$

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# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 28
- # branching: 2

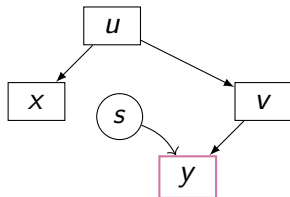


$s = \text{GGCTAGCTACAACGAGCTGGGCAGCA} \text{GCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGG} \text{CAGTAACGAGCTGGG}$   
 $y = \text{GGCTAGCTACAACGAGCTGGGCAGCA} \text{GCTAGTAGTAGG}$

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# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 29
- # branching: 2



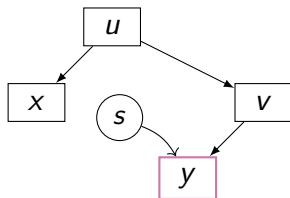
$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGG} \text{CAGTAACGAGCTGGG}$   
 $y = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$

└─┘



# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 30
- # branching: 2

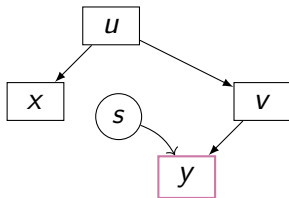


$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \textcolor{brown}{\text{A}} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGG} \textcolor{brown}{\text{T}} \text{CAGTAACGAGCTGGG}$   
 $y = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$

$\vdash$

# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 31
- # branching: 2



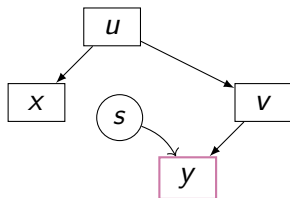
$s = \text{GGCTAGCTACAACGAGCTGGGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \textcolor{brown}{A} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGG} \textcolor{brown}{T} \text{CAGTAACGAGCTGGG}$   
 $y = \text{GGCTAGCTACAACGAGCTGGGCAGCTAGTAGTAGG}$

$\vdash$



# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 33
- # branching: 2

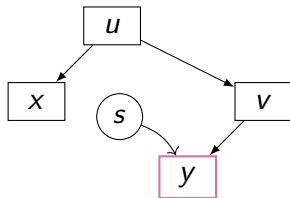


$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA}^{\text{A}}\text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGG}^{\text{T}}\text{CAGTAACGAGCTGGG}$   
 $y = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$

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# Sequential Longest Common Prefix

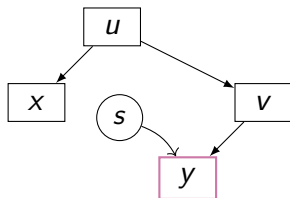
- Best way to find LCP?
- $k = |s| = 38$
- # computations: 34
- # branching: 2



$s = \text{GGCTAGCTACAACGAGCTGGGCAGCTAGT}$ AGTAGG  
 $u = \text{GGCTA}$ **A**CTACAACGAGCTGGCGTTGTGAGCTAC  
 $v = \text{GGCTAGCTACAACGAGCTGG}$ **T**CAGTAACGAGCTGGG  
 $y = \text{GGCTAGCTACAACGAGCTGGGCAGCTAGT}$ AGTAGG  
└─┘

# Sequential Longest Common Prefix

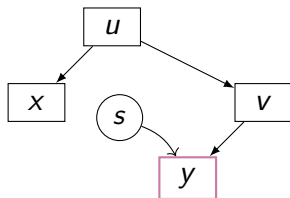
- Best way to find LCP?
- $k = |s| = 38$
- # computations: 35
- # branching: 2



$s = \text{GGCTAGCTACAACGAGCTGGGCAGCTAGTA} \text{GTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGG} \text{CAGTAACGAGCTGGG}$   
 $y = \text{GGCTAGCTACAACGAGCTGGGCAGCTAGTA} \text{GTAGG}$

# Sequential Longest Common Prefix

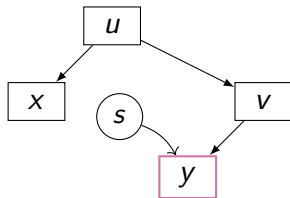
- Best way to find LCP?
- $k = |s| = 38$
- # computations: 36
- # branching: 2



$s = \text{GGCTAGCTACAACGAGCTGGGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGG} \text{CAGTAACGAGCTGGG}$   
 $y = \text{GGCTAGCTACAACGAGCTGGGCAGCTAGTAGTAGG}$

# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 37
- # branching: 2

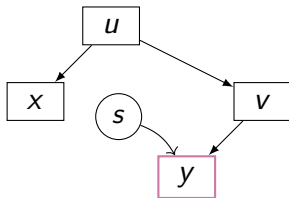


$s = \text{GGCTAGCTACAACGAGCTGGGCAGCTAGTAGT} \text{AGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGG} \text{CAGTAACGAGCTGGG}$   
 $y = \text{GGCTAGCTACAACGAGCTGGGCAGCTAGTAGT} \text{AGG}$



# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 38
- # branching: 2



$s = \text{GGCTAGCTACAACGAGCTGGGCAGCTAGTAGTA}\text{GG}$

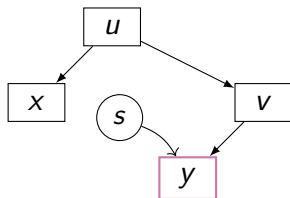
$u = \text{GGCTA}\text{ACTACAACGAGCTGGCGTTGTGAGCTAC}$

$v = \text{GGCTAGCTACAACGAGCTGG}\text{CAGTAACGAGCTGGG}$

$y = \text{GGCTAGCTACAACGAGCTGGGCAGCTAGTAGTA}\text{GG}$

# Sequential Longest Common Prefix

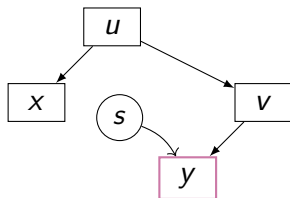
- Best way to find LCP?
- $k = |s| = 38$
- # computations: 39
- # branching: 2



$s = \text{GGCTAGCTACAACGAGCTGGGCAGCTAGTAGG}$   
 $u = \text{GGCTA}^{\text{A}}\text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGG}^{\text{T}}\text{CAGTAACGAGCTGGG}$   
 $y = \text{GGCTAGCTACAACGAGCTGGGCAGCTAGTAGG}$

# Sequential Longest Common Prefix

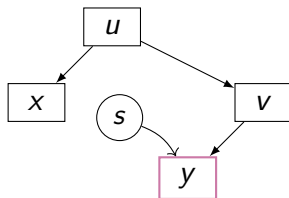
- Best way to find LCP?
- $k = |s| = 38$
- # computations: 40
- # branching: 2



$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGG}$   
 $u = \text{GGCTA}^{\text{A}}\text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGG}^{\text{T}}\text{CAGTAACGAGCTGGG}$   
 $y = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGG}$

# Sequential Longest Common Prefix

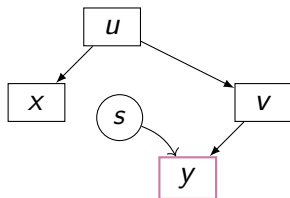
- Best way to find LCP?
- $k = |s| = 38$
- # computations: 40
- # branching: 2
- # computations =  $k + \text{\# branching}$



$s = \text{GGCTAGCTACAACGAGCTGGGCAGCTAGTAGG}$   
 $u = \text{GGCTA} \textcolor{brown}{\text{A}} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGG} \textcolor{brown}{\text{T}} \text{CAGTAACGAGCTGGG}$   
 $y = \text{GGCTAGCTACAACGAGCTGGGCAGCTAGTAGG}$

# Sequential Longest Common Prefix

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 40
- # branching: 2
- # computations =  $k + \# \text{ branching}$
- ✓ # computations =  $\mathcal{O}(k + \# \text{ branching})$



$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$

$u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$

$v = \text{GGCTAGCTACAACGAGCTGG} \text{CAGTAACGAGCTGGG}$

$y = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$

⊥

# Naive Parallel LCP

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 0
- # branching: 0
- # steps: 0

$s = \text{GGCTAGCTACAACGAGCTGGGCGAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTAACTACAACGAGCTGGCGTTGTGAGCTAC}$

# Naive Parallel LCP

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 33
- # branching: 0
- # steps: 1

$s =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG  
 $u =$  GGCTAACTACAACGAGCTGGCGTTGTGAGCTAC



33

# Naive Parallel LCP

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 33
- # branching: 1
- # steps: 1

$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \textcolor{brown}{\text{ACTACAACGAGCTGG}} \textcolor{brown}{\text{CGTTGTGAGCTAC}}$   
 $v = \text{GGCTAGCTACAACGAGCTGGTCAGTAACGAGCTGGG}$



# Naive Parallel LCP

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 64
- # branching: 1
- # steps: 2

$s =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG  
 $u =$  GGCTAACTACAACGAGCTGGCGTTGTGAGCTAC  
 $v =$  GGCTAGCTACAACGAGCTGGTCAGTAACGAGCTGGG

31

# Naive Parallel LCP

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 64
- # branching: 2
- # steps: 2

$s =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG  
 $u =$  GGCTA**ACT**ACAACGAGCTGG**CGTTGTGAGCTAC**  
 $v =$  GGCTAGCTACAACGAGCTGG**TCAGTAACGAGCTGGG**  
 $y =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG

# Naive Parallel LCP

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 82
- # branching: 2
- # steps: 3

$s =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG  
 $u =$  GGCTA**ACT**ACAACGAGCTGG**CGTTGTGAGCTAC**  
 $v =$  GGCTAGCTACAACGAGCTGG**TCAGTAACGAGCTGGG**  
 $y =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG

18

# Naive Parallel LCP

- Best way to find LCP?
  - $k = |s| = 38$
  - # computations: 82
  - # branching: 2
  - # steps: 3
- ✓ # steps = # branching + 1 = # comparisons

$s =$  GGCTAGCTACAACGAGCTGGGCGAGCAGCTAGTAGTAGG  
 $u =$  GGCTA**ACT**ACAACGAGCTGG**CGTTGTGAGCTAC**  
 $v =$  GGCTAGCTACAACGAGCTGG**TCAGTAACGAGCTGGG**  
 $y =$  GGCTAGCTACAACGAGCTGGGCGAGCAGCTAGTAGTAGG

18

# Naive Parallel LCP

- Best way to find LCP?
  - $k = |s| = 38$
  - # computations: 82
  - # branching: 2
  - # steps: 3
- ☒ # steps = # branching + 1 = # comparisons
- ☐ # computations?

$s =$  GGCTAGCTACAACGAGCTGGGCGAGCAGCTAGTAGTAGG  
 $u =$  GGCTA**ACT**ACAACGAGCTGG**CGTTGTGAGCTAC**  
 $v =$  GGCTAGCTACAACGAGCTGG**TCAGTAACGAGCTGGG**  
 $y =$  GGCTAGCTACAACGAGCTGGGCGAGCAGCTAGTAGTAGG

18

# Naive Parallel LCP

- Best way to find LCP?
  - $k = |s| = 38$
  - # computations: 82
  - # branching: 2
  - # steps: 3
- ✓ # steps = # branching + 1 = # comparisons
- ✗ # computations? –  $\mathcal{O}(k \times \# \text{ comparisons}) = \mathcal{O}(k \times \# \text{ branching})$

$s =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG  
 $u =$  GGCTA**ACT**ACAACGAGCTGG**CGTTGTGAGCTAC**  
 $v =$  GGCTAGCTACAACGAGCTGG**TCAGTAACGAGCTGGG**  
 $y =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG

18

# Parallel LCP

- Best way to find LCP?
- $k = |s| = 38$
- Only compare  $\frac{k}{\# \text{ comparisons}} = 13$  characters?
- # computations: 0
- # branching: 0
- # steps: 0

$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTAACTACAACGAGCTGGCGTTGTGAGCTAC}$

# Parallel LCP

- Best way to find LCP?
- $k = |s| = 38$
- Only compare  $\frac{k}{\# \text{ comparisons}} = 13$  characters?
- # computations: 13
- # branching: 0
- # steps: 1

$s = \text{GGCTA}\text{GCTACAAC}\text{GAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA}\text{ACTACAAC}\text{GAGCTGGCGTTGTGAGCTAC}$

$\underbrace{\hspace{10em}}_{13}$



# Parallel LCP

- Best way to find LCP?
- $k = |s| = 38$
- Only compare  $\frac{k}{\# \text{ comparisons}} = 13$  characters?
- # computations: 13
- # branching: 1
- # steps: 1

$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGGTCAGTAACGAGCTGGG}$

# Parallel LCP

- Best way to find LCP?
- $k = |s| = 38$
- Only compare  $\frac{k}{\# \text{ comparisons}} = 13$  characters?
- # computations: 26
- # branching: 1
- # steps: 2

$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA}^{\text{orange}}\text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGGTCAGTAACGAGCTGGG}$

13

# Parallel LCP

- Best way to find LCP?
- $k = |s| = 38$
- Only compare  $\frac{k}{\# \text{ comparisons}} = 13$  characters?
- # computations: 39
- # branching: 1
- # steps: 3

$s =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG  
 $u =$  GGCTA**A**CTACAACGAGCTGGCGTTGTGAGCTAC  
 $v =$  GGCTAGCTACAACGAGCTGG**T**CAG**T**A**A**C**G**AGCTGGG

|-----|  
13

# Parallel LCP

- Best way to find LCP?
- $k = |s| = 38$
- Only compare  $\frac{k}{\# \text{ comparisons}} = 13$  characters?
- # computations: 39
- # branching: 2
- # steps: 3

$s =$  **GGCTAGCTACAACGAGCTGG**GCAGCAGCTAGTAGTAGG  
 $u =$  **GGCTA****ACTACAAC**GAGCTGGCGTTGTGAGCTAC  
 $v =$  **GGCTAGCTACAACGAGCTGG****TCAGTAACGAG**CTGGG  
 $y =$  **GGCTAGCTACAACGAGCTGG**GCAGCAGCTAGTAGTAGG

# Parallel LCP

- Best way to find LCP?
- $k = |s| = 38$
- Only compare  $\frac{k}{\# \text{ comparisons}} = 13$  characters?
- # computations: 52
- # branching: 2
- # steps: 4

$s =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG  
 $u =$  GGCTA**ACT**ACAACGAGCTGGCGTTGTGAGCTAC  
 $v =$  GGCTAGCTACAACGAGCTGG**TCAGTAAC**AGCTGGG  
 $y =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG

13

# Parallel LCP

- Best way to find LCP?
- $k = |s| = 38$
- Only compare  $\frac{k}{\# \text{ comparisons}} = 13$  characters?
- # computations: 57
- # branching: 2
- # steps: 5

$s =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG  
 $u =$  GGCTA**A**CTACAACGAGCTGGCGTTGTGAGCTAC  
 $v =$  GGCTAGCTACAACGAGCTGG**T**CAG**T**A**A**C**G**AGCTGGG  
 $y =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG

|-----|  
5

# Parallel LCP

- Best way to find LCP?
  - $k = |s| = 38$
  - Only compare  $\frac{k}{\# \text{ comparisons}} = 13$  characters?
  - # steps?
- # computations: 57
  - # branching: 2
  - # steps: 5

$s =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG  
 $u =$  GGCTA**A**CTACAACGAGCTGGCGTTGTGAGCTAC  
 $v =$  GGCTAGCTACAACGAGCTGG**T**CAG**T**A**A**C**G**AGCTGGG  
 $y =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG

|-----|  
5

# Parallel LCP

- Best way to find LCP?
- $k = |s| = 38$
- Only compare  $\frac{k}{\# \text{ comparisons}} = 13$  characters?
- ✓ # steps?  $2 \times \# \text{ comparisons} = \mathcal{O}(\# \text{ branching})$
- # computations: 57
- # branching: 2
- # steps: 5

$s =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG  
 $u =$  GGCTA**A**CTACAACGAGCTGGCGTTGTGAGCTAC  
 $v =$  GGCTAGCTACAACGAGCTGG**T**CAG**T**A**A**C**G**AGCTGGG  
 $y =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG

5



# Parallel LCP

- Best way to find LCP?
  - $k = |s| = 38$
  - Only compare  $\frac{k}{\# \text{ comparisons}} = 13$  characters?
  - # computations: 57
  - # branching: 2
  - # steps: 5
- ☒ # steps?  $2 \times \# \text{ comparisons} = \mathcal{O}(\# \text{ branching})$
- ☐ # computations?

$s =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG  
 $u =$  GGCTA**A**CTACAACGAGCTGGCGTTGTGAGCTAC  
 $v =$  GGCTAGCTACAACGAGCTGG**T**CAG**T**A**A**C**G**AGCTGGG  
 $y =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG

5

# Parallel LCP

- Best way to find LCP?
  - $k = |s| = 38$
  - Only compare  $\frac{k}{\# \text{ comparisons}} = 13$  characters?
  - # computations: 57
  - # branching: 2
  - # steps: 5
- ✓ # steps?  $2 \times \# \text{ comparisons} = \mathcal{O}(\# \text{ branching})$
- ✓ # computations?  $\frac{k}{\# \text{ comparisons}} \times \# \text{ steps} = \mathcal{O}(k)$

$s =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGG  
 $u =$  GGCTA**A**CTACAACGAGCTGGCGTTGTGAGCTAC  
 $v =$  GGCTAGCTACAACGAGCTGG**T**CAG**T**A**A**C**G**AGCTGGG  
 $y =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGG

5

# Parallel LCP

- Best way to find LCP?
- $k = |s| = 38$
- # computations: 57
- # branching: 2
- # steps: 5
- Only compare  $\frac{k}{\# \text{ comparisons}} = 13$  characters?
- ✓ # steps?  $2 \times \# \text{ comparisons} = \mathcal{O}(\# \text{ branching})$
- ✓ # computations?  $\frac{k}{\# \text{ comparisons}} \times \# \text{ steps} = \mathcal{O}(k)$
- What if  $\text{LCP} \ll k$ ?

$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \color{brown}{\text{ACTACAACGAGCTGGCGTTGTGAGCTAC}}$   
 $v = \text{GGCTAGCTACAACGAGCTGG} \color{brown}{\text{TCAGT}} \color{brown}{\text{A}} \color{brown}{\text{C}} \color{brown}{\text{AGCTGGG}}$   
 $y = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$

5

# Parallel LCP-Aware LCP

- Best way to find LCP?
- $k = |s| = 38$
- Should be *LCP-aware*
- Double and halve # compared characters
- # computations: 0
- # branching: 0
- # steps: 0

$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTAACTACAACGAGCTGGCGTTGTGAGCTAC}$

# Parallel LCP-Aware LCP


- Best way to find LCP?
- $k = |s| = 38$
- Should be *LCP-aware*
- Double and halve # compared characters
- # computations: 1
- # branching: 0
- # steps: 1

$s =$  GGCTAGCTACAACGAGCTGGGGCAGCAGCTAGTAGG  
 $u =$  GGCTAACTACAACGAGCTGGCGTTGTGAGCTAC  
1

# Parallel LCP-Aware LCP

- Best way to find LCP?
- $k = |s| = 38$
- Should be *LCP-aware*
- Double and halve # compared characters
- # computations: 3
- # branching: 0
- # steps: 2


$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTAACTACAACGAGCTGGCGTTGTGAGCTAC}$



# Parallel LCP-Aware LCP

- Best way to find LCP?
- $k = |s| = 38$
- Should be *LCP-aware*
- Double and halve # compared characters
- # computations: 7
- # branching: 1
- # steps: 3

$s = \text{GGCTAGCTACAAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTAAC TACAAACGAGCTGGCGTTGTGAGCTAC}$



# Parallel LCP-Aware LCP

- Best way to find LCP?
- $k = |s| = 38$
- Should be *LCP-aware*
- Double and halve # compared characters
- # computations: 7
- # branching: 1
- # steps: 3


$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGGTCAGTAACGAGCTGGG}$



# Parallel LCP-Aware LCP

- Best way to find LCP?
- $k = |s| = 38$
- Should be *LCP-aware*
- Double and halve # compared characters
- # computations: 9
- # branching: 1
- # steps: 4


$s = \text{GGCTAGCTACAAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \textcolor{brown}{\text{AC}} \text{TACAAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAAACGAGCTGGTCAGTAACGAGCTGGG}$



# Parallel LCP-Aware LCP

- Best way to find LCP?
- $k = |s| = 38$
- Should be *LCP-aware*
- Double and halve # compared characters
- # computations: 13
- # branching: 1
- # steps: 5

$s = \text{GGCTAGCTAC} \text{ACGAGCTGGGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTAC} \text{ACGAGCTGGTCAGTAACGAGCTGGG}$



# Parallel LCP-Aware LCP

- Best way to find LCP?
- $k = |s| = 38$
- Should be *LCP-aware*
- Double and halve # compared characters
- # computations: 21
- # branching: 1
- # steps: 6

$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \color{brown}{\text{CT}} \text{ACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTG} \text{GTCAGTAACGAGCTGGG}$

|-----|  
8

# Parallel LCP-Aware LCP

- Best way to find LCP?
- $k = |s| = 38$
- Should be *LCP-aware*
- Double and halve # compared characters
- # computations: 37
- # branching: 2
- # steps: 7

$s =$  GGCTAGCTACAACGAGCTGGCAGCAGCTAGTAGTAGG  
 $u =$  GGCTA<sup>A</sup>CTACAACGAGCTGGCGTTGTGAGCTAC  
 $v =$  GGCTAGCTACAACGAGCTGG<sup>T</sup>CAG<sup>T</sup>A<sup>A</sup>C<sup>A</sup>G<sup>A</sup>CTGGG

16

# Parallel LCP-Aware LCP

- Best way to find LCP?
- $k = |s| = 38$
- Should be *LCP-aware*
- Double and halve # compared characters
- # computations: 37
- # branching: 2
- # steps: 7

$s =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG  
 $u =$  GGCTA<sup>orange</sup>CTACAACGAGCTGGCGTTGTGAGCTAC  
 $v =$  GGCTAGCTACAACGAGCTGG<sup>orange</sup>TCAG<sup>orange</sup>T<sup>orange</sup>A<sup>orange</sup>C<sup>orange</sup>GAGCTGGG  
 $y =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG

# Parallel LCP-Aware LCP

- Best way to find LCP?
- $k = |s| = 38$
- Should be *LCP-aware*
- Double and halve # compared characters
- # computations: 45
- # branching: 2
- # steps: 8

$s =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG  
 $u =$  GGCTA**AC**TACAACGAGCTGGCGTTGTGAGCTAC  
 $v =$  GGCTAGCTACAACGAGCTGG**TCAGTAAC**AGCTGGG  
 $y =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG

|-----|  
8

# Parallel LCP-Aware LCP

- Best way to find LCP?
- $k = |s| = 38$
- Should be *LCP-aware*
- Double and halve # compared characters
- # computations: 55
- # branching: 3
- # steps: 9

$s =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG  
 $u =$  GGCTA**ACT**ACAACGAGCTGGCGTTGTGAGCTAC  
 $v =$  GGCTAGCTACAACGAGCTGG**TCAGTAAC**AGCTGGG  
 $y =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG

10

# Parallel LCP-Aware LCP

- Best way to find LCP?
  - $k = |s| = 38$
  - Should be *LCP-aware*
  - Double and halve # compared characters
  - # computations: 55
  - # branching: 3
  - # steps: 9
- # steps?

$s =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG  
 $u =$  GGCTA**ACT**ACAACGAGCTGGCGTTGTGAGCTAC  
 $v =$  GGCTAGCTACAACGAGCTGG**TCAGTAAC**AGCTGGG  
 $y =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG



# Parallel LCP-Aware LCP

- Best way to find LCP?
  - $k = |s| = 38$
  - Should be *LCP-aware*
  - Double and halve # compared characters
  - # computations: 55
  - # branching: 3
  - # steps: 9
- ✓ # steps?  $\log \ell + 2 \times \# \text{ branching} = \log \ell + \mathcal{O}(\# \text{ branching})$

$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \color{brown}{\text{A}} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGGT} \color{brown}{\text{C}} \color{brown}{\text{AGT}} \color{brown}{\text{A}} \color{brown}{\text{C}} \color{brown}{\text{AG}} \text{CTGGG}$   
 $y = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$

# Parallel LCP-Aware LCP

- Best way to find LCP?
  - $k = |s| = 38$
  - Should be *LCP-aware*
  - Double and halve # compared characters
  - # computations: 55
  - # branching: 3
  - # steps: 9
- ☒ # steps?  $\log \ell + 2 \times \# \text{ branching} = \log \ell + \mathcal{O}(\# \text{ branching})$
- ☐ # computations?

$s = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$   
 $u = \text{GGCTA} \color{brown}{\text{A}} \text{CTACAACGAGCTGGCGTTGTGAGCTAC}$   
 $v = \text{GGCTAGCTACAACGAGCTGGT} \color{brown}{\text{C}} \color{brown}{\text{AGT}} \color{brown}{\text{A}} \color{brown}{\text{C}} \color{brown}{\text{AG}} \text{CTGGG}$   
 $y = \text{GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG}$

# Parallel LCP-Aware LCP

- Best way to find LCP?
  - $k = |s| = 38$
  - Should be *LCP-aware*
  - Double and halve # compared characters
  - # computations: 55
  - # branching: 3
  - # steps: 9
- ✓ # steps?  $\log \ell + 2 \times \# \text{ branching} = \log \ell + \mathcal{O}(\# \text{ branching})$
- ✓ # computations?  $2\ell + \# \text{ branching} = \mathcal{O}(\ell + \# \text{ branching})$

$s =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG  
 $u =$  GGCTA**A**CTACAACGAGCTGGCGTTGTGAGCTAC  
 $v =$  GGCTAGCTACAACGAGCTGG**T**CAG**T**A**A**C**G**AGCTGGG  
 $y =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG

# LCP Overview

- Subsequent LCP calculations

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- **Sequential:**
  - ✓ # computations (work):  $\mathcal{O}(\ell + \# \text{ branching})$
  - ✗ # steps (time):  $\mathcal{O}(\ell + \# \text{ branching})$

- Subsequent LCP calculations
- **Sequential:**
  - ✓ # computations (work):  $\mathcal{O}(\ell + \# \text{ branching})$
  - ✗ # steps (time):  $\mathcal{O}(\ell + \# \text{ branching})$
- **Parallel:**
  - ? # computations (work):  $\mathcal{O}(k + \# \text{ branching})$
  - ✓ # steps (time):  $\mathcal{O}(\# \text{ branching})$

- Subsequent LCP calculations
- **Sequential:**
  - ✓ # computations (work):  $\mathcal{O}(\ell + \# \text{ branching})$
  - ✗ # steps (time):  $\mathcal{O}(\ell + \# \text{ branching})$
- **Parallel:**
  - ? # computations (work):  $\mathcal{O}(k + \# \text{ branching})$
  - ✓ # steps (time):  $\mathcal{O}(\# \text{ branching})$
- **LCP-aware Parallel:**
  - ✓ # computations (work):  $\mathcal{O}(\ell + \# \text{ branching})$
  - ? # steps (time):  $\log \ell + \mathcal{O}(\# \text{ branching})$

- Subsequent LCP calculations

- **Sequential:**

- ☒ # computations (work):  $\mathcal{O}(\ell + \# \text{ branching})$
- ☒ # steps (time):  $\mathcal{O}(\ell + \# \text{ branching})$

- **Parallel:**

- ☐ # computations (work):  $\mathcal{O}(k + \# \text{ branching})$
- ☒ # steps (time):  $\mathcal{O}(\# \text{ branching})$

- **LCP-aware Parallel:**

- ☒ # computations (work):  $\mathcal{O}(\ell + \# \text{ branching})$
- ☐ # steps (time):  $\log \ell + \mathcal{O}(\# \text{ branching})$

- Zip-trie:  $\mathcal{O}(\log n)$  branching

- Parallel String B-Tree:  $\mathcal{O}(\log_B n) = \mathcal{O}\left(\frac{\log n}{\log \log n}\right)$  branching



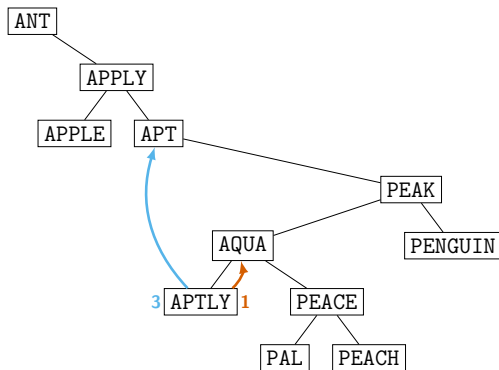
# LCP Requirement

- Requirement: LCP length *never* decreases
- How?

$s =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG  
 $u =$  GGCTA|ACTACAACGAGCTGGCGTTGTGAGCTAC  
 $v =$  GGCTAGCTACAACGAGCTGG|TCAGTAACGAGCTGGG  
 $y =$  GGCTAGCTACAACGAGCTGG|GCAGCAGCTAGTAGTAGG|

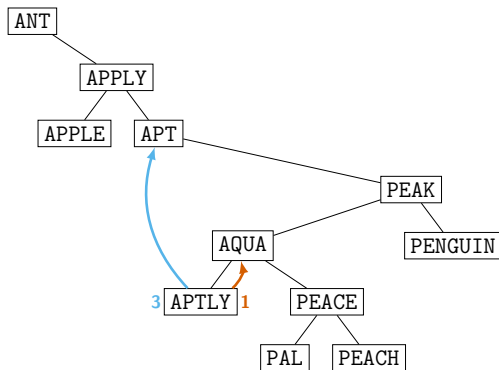
# LCP Requirement

- Requirement: LCP length *never* decreases
- How? – LCP length metadata & pointers



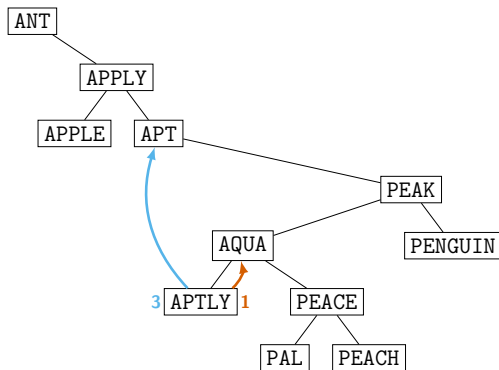
# LCP Requirement

- Requirement: LCP length *never* decreases
- How? – LCP length metadata & pointers
- Space usage?



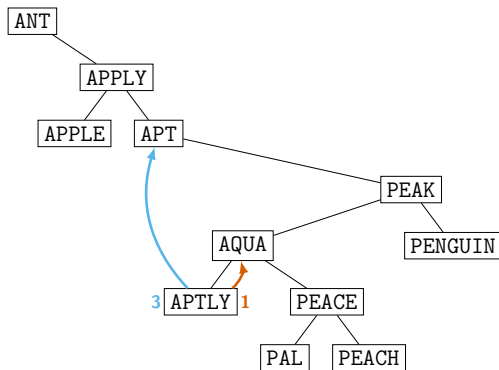
# LCP Requirement

- Requirement: LCP length *never* decreases
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- Space usage? –  $\mathcal{O}(\log \max k) + \text{pointers}$



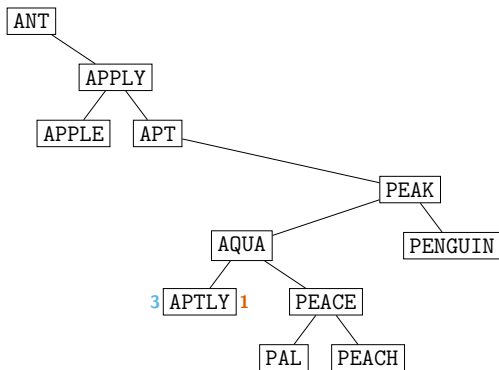
# LCP Requirement

- Requirement: LCP length *never* decreases
- How? – LCP length metadata & pointers
- Space usage? –  $\mathcal{O}(\log \max k) + \text{pointers}$
- Zip trie:  $\log \log$  bits of metadata



# LCP Requirement

- Requirement: LCP length *never* decreases
- How? – LCP length metadata ~~& pointers~~
- Space usage? –  $\mathcal{O}(\log \max k)$  ~~+ pointers~~
- Zip trie: log log bits of metadata



# LCP Problem

- Zip tries require LCP length metadata
- Max LCP length is  $\max k \leftrightarrow \log(\max k)$  bits

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- We can only represents  $\mathcal{O}(\log(\max k))$  values!

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- Which values to pick?

# LCP Problem

- Zip tries require LCP length metadata
- Max LCP length is  $\max k \leftrightarrow \log(\max k)$  bits
- **Goal:** Use only  $\mathcal{O}(\log \log(\max k))$  bits of metadata
- We can only represent  $\mathcal{O}(\log(\max k))$  values!
- Which values to pick? – [see paper!](#)

$s =$  GGCTAGCTACAACGAGCTGGGCAGCAGCTAGTAGTAGG  
 $u =$  GGCTAACTACAACGAGCTGGCGTTGTGAGCTAC

## Lemma

*A data structure,  $\mathcal{D}_S$ , spending  $\mathcal{O}(\ell + A(n))$  time on comparisons using exact LCP lengths will spend at most  $\mathcal{O}(\ell)$  additional time using approximate LCP lengths.*

## Theorem

*Let  $\ell$  be the length of the longest common prefix between a string key  $x$  of length  $k$ , and the stored keys in a parallel zip-trie,  $T$ .  $T$  can perform prefix search, predecessor/successor queries, and update operations on  $x$  in  $\mathcal{O}(\log n)$  span and  $\mathcal{O}(\frac{k}{\alpha} + \log n)$  work under the practical PRAM model, or alternatively in  $\mathcal{O}(\log \frac{\ell}{\alpha} + \log n)$  span and  $\mathcal{O}(\frac{\ell}{\alpha} + \log n)$  work.*

- <https://github.com/ofekih/ZipAndSkipTries>

# Results – Parallel String $B$ -Tree

## Theorem





*By setting  $B = \log n$ , a parallel string  $B$ -tree can perform prefix search in  $\mathcal{O}(\frac{\log n}{\log \log n})$  span and  $\mathcal{O}(\frac{k}{\alpha} + \log^2 n)$  work in the practical CRCW PRAM model. Operations that return  $m$  keys can be done in the same span and in  $\mathcal{O}(m)$  additional work.*

## Theorem






*A parallel string  $B$ -tree can perform prefix search in  $\mathcal{O}(\log_B n)$  I/O span and  $\mathcal{O}(\frac{k}{\alpha B} + \log_B n)$  I/O work in the practical CRCW PEM model. Operations that return  $m$  keys can be done in the same span and in  $\mathcal{O}(m/B)$  additional I/O work.*

- Concurrent updates
- Compression techniques, e.g., DNA-specific
- Update operations in the string B-tree
- Reduce gap between RAM and PRAM work for string B-tree

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