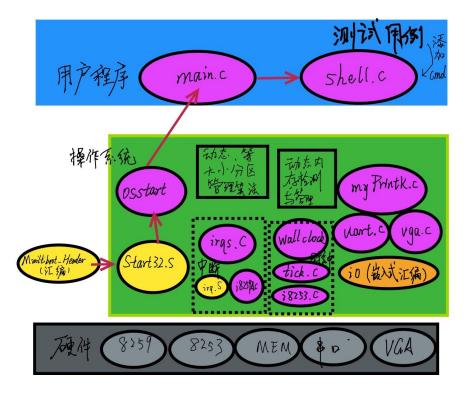
# 操作系统实验报告

实验四 Memory Management

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### 一、 软件框图

紫色为C程序,黄色为汇编程序,橙色为嵌入式汇编



概述:软件大体上可以分为两个层次:用户程序和操作系统,其中操作系统又可以分为与用户程序的接口、I0设备的驱动程序、中断控制程序、时钟功能程序和内存管理程序等各功能模块,在每块中又可以进一步划分更细的层次(越上方的层次越高),如图所示。

# 二、 主流程说明

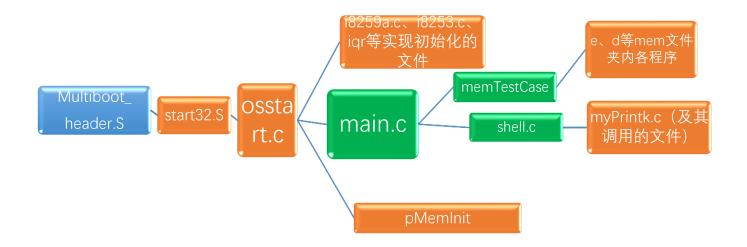
- 1. qemu启动header;
- 2. header通过调用myOS提供的\_start, 跳转到汇编文件start32.S; (进入myOS)
- 3. 从start32调用C语言入口, 进入到c程序osstart.c;
- 4. osstart.c调用i8259A和i8253函数进行初始化,调用enable\_interrupt函数启用中断;调用 pMemInit接口进行内存检测;并调用用户程序的接口,即myMain函数,从而执行用户程序 main.c;

(进入UserApp)

- 5. main.c调用memtestinit接口以及shell接口;
- 6. memtest添加新的指令,并链接各指令对应的内存分配/管理函数
- 7. shell使用myprintk.c中的函数实现交互;

#### 流程图:

(橙色表示myOS内的程序,绿色表示用户程序)



## 三、 主要功能模块说明&源代码说明

### 1. 内存检测模块

内存检测的主要函数:

Memtest: 从start开始,以grainSize为步长,进行内存检测。由于start是无符号长整型 (4byte),此处使用的检测方法为把start作为int型的指针,检测头尾的4个字节。

```
□void memTest(unsigned long start, unsigned long grainSize){
    /*本函数需要实现!!!*/
    /* ... */
    //最后,输出可用的内存的起始地只和大小,别忘记赋值给上面的全局变量
    if (start < 0x100000) {</pre>
        myPrintk(0x7, "the start location should be no less than <math>1M\n");
        return;
    if (grainSize < 0x400) {</pre>
        myPrintk(0x7, "the grainsize should be no less than 1K\n");
 *p = 0xaa55aa55;
 if (*p != 0xaa55aa55)
     flag = 1;//结束标志
 *p = 0x55aa55aa;
 if (*p != 0x55aa55aa)
     flag = 1;
 *p = i;
 if (flag)break;
```

### 2. 等大小分区管理算法模块

结构体:

```
//eFPartition是表示整个内存的数据结构
// 一个EEB表示一个空闲可用的Block

struct EEB {
    unsigned long next_start;
    };
    //eFPartition是表示整个内存的数据结构
    unsigned long totalN;
    unsigned long perSize; //unit: byte
    unsigned long firstFree;
};
```

这部分的主要函数有:

a) eFPartitionTotalSize, 用于计算A的合理尺寸:

b) eFPartitionInit,用于对A进行划分和管理:

c) eFPartitionAlloc和eFPartitionFree, 用于按需求分配和释放:

```
□unsigned long eFPartitionAlloc(unsigned long EFPHandler){
     //本函数需要实现!!!
     /*本函数分配一个空闲块的内存并返回相应的地址,EFPHandler表示整个内存的首地址
     struct eFPartition* efp = (struct eFPartition*) EFPHandler;
     struct EEB* eeb = (struct EEB*) efp->firstFree;
     if (efp->firstFree > EFPHandler + eFPartitionTotalSize(efp->perSize, efp->totalN)) //failed
        return 0;
     efp->firstFree = eeb->next_start;//success
     return (unsigned long)eeb;
|unsigned long eFPartitionFree(unsigned long EFPHandler, unsigned long mbStart) {
    //本函数需要实现!!!
    /* ... */
    struct eFPartition* efp = (struct eFPartition*) EFPHandler;
    if (mbStart == 0) mbStart = EFPHandler + eFPartitionTotalSize(efp->perSize, efp->totalN);
    efp->firstFree = EFPHandler + sizeof(struct eFPartition);
    struct EEB* eeb = (struct EEB*) efp->firstFree;
    int i = 0:
    while ((unsigned long)eeb < mbStart) {</pre>
        eeb->next start = (unsigned long)eeb + efp->perSize;
        eeb = (struct EEB*) ((unsigned long)eeb + efp->perSize);
    eeb = (struct EEB*) ((unsigned long)eeb - efp->perSize);
    if (i == efp->totalN)
        eeb->next_start = 0;
    return 1;
```

### 3. 动态分区管理算法模块 结构体:

```
//dPartition 是整个动态分区内存的数据结构
∃struct dPartition{
      unsigned long size;
      unsigned long firstFreeStart;
 };
 // EMB每一个block的数据结构, userdata可以暂时不用管。
□struct EMB{
    unsigned long size;
   union {
                          // if free: pointer to next block
       unsigned long nextStart;
                           // if allocated, blongs to user
       unsigned long userData;
a) dPartitionInit: 对B进行初始化
unsigned long dPartitionInit(unsigned long start, unsigned long totalSize) {
   //本函数需要实现!!!
   if (totalSize < sizeof(struct EMB) + sizeof(struct dPartition) )
      return 0:
   struct dPartition* dp = (struct dPartition*) start;
   dp->size = totalSize;
   dp->firstFreeStart = start + 8;
   struct EMB* emb = (struct EMB*)(dp->firstFreeStart);
   emb->size = totalSize - 16:
   emb->nextStart = 0;
   return start:
b) dPartitionAlloc和dPartitionFree: 按需求分配和回收
unsigned long dPartitionAlloc(unsigned long dp, unsigned long size) {
   return dPartitionAllocFirstFit(dp, size);
unsigned long dPartitionFree(unsigned long dp, unsigned long start){
   return dPartitionFreeFirstFit(dp, start);
调用的两个F-F算法函数功能如下: (实现代码较长,此处略)
unsigned long dPartitionAllocFirstFit(unsigned long dp, unsigned long size) [
     //本函数需要实现!!!
     使用firstfit的算法分配空间,当然也可以使用其他fit,不限制。
     最后,成功分配返回首地址,不成功返回0
□unsigned long dPartitionFreeFirstFit(unsigned long dp, unsigned long start) {
     //本函数需要实现!!!
     /*按照对应的fit的算法释放空间
     注意检查要释放的start~end这个范围是否在dp有效分配范围内
     返回1 没问题
     返回0 error
流程图:
                                 dPartitionAlloc
```

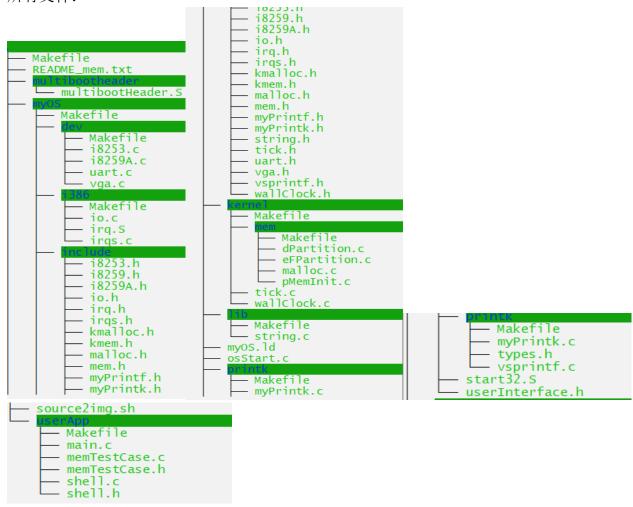
# dPartitionAlloc dPartitionAllocFF (k)malloc/(k)free dPartitionFreeFF

### 4. shell模块:新增如下功能

```
void addNewCmd( unsigned char *cmd,
        int (*func) (int argc, unsigned char **argv),
        void (*help_func) (void),
        unsigned char* description) {
   struct cmd* tcmd = (struct cmd*) malloc(sizeof(struct cmd));
   strcpy(cmd, tcmd->cmd);
   tcmd->func = func;
   tcmd->help_func = help_func;
   strcpy(description, tcmd->description);
   tcmd->nextCmd = NULL;
   struct cmd* tmpCmd = ourCmds;
   if (tmpCmd == NULL)
        ourCmds = tcmd:
        return:
   while (tmpCmd->nextCmd != NULL) {
       tmpCmd = tmpCmd->nextCmd;
    tmpCmd->nextCmd = tcmd;
```

### 四、 目录组织

所有文件:



### Makefile组织:

见上图中的各Makefile文件

# 五、 代码布局

由myOS.ld的代码可知, myOS.elf文件中有三个 section:

1. 第一个section为.text,位置从1M处开始,在.text内的分布为8字节对齐,前12字节为魔术,从第16字节开始是代码部分;

代码结束后16位对齐;

- 2. 第二个section为.data,位置从.text结束并对齐后开始; 末尾16位对齐;
- 3. 第三个section为.bss,位置从.data结末尾对齐后开始; 末尾16位对齐;
- 4. .bss结束后是 end, 此处是我们可以操作的内存空间的开始, 512位对齐;

### 六、 编译过程说明

由makefile可知,编译过程有以下两步:

- 1. 编译汇编代码(header.S和start32.S)和C代码(osstart.c等)生成.o文件;
- 2. 根据myOS.ld的部署要求,把上述.o文件链接成myOS.elf文件 如下图所示:

ld -n -T myOS/myOS.ld output/multibootheader/multibootHeader.o output/myOS/start32.o output/myOS/osStart.o output/myOS/dev/uart.o output/myOS/dev/vga.o output/myOS/dev/i8253.o output/myOS/dev/i8259A.o output/myOS/i386/io.o output/myOS/i386/irq.o output/myOS/i386/irqs.o output/myOS/printk/myPrintk.o output/myOS/lib/string.o output/myOS/kernel/tick.o output/myOS/ernel/wallClock.o output/myOS/kernel/mem/pMemInit.o output/myOS/kernel/mem/dPartition.o output/myOS/kernel/mem/eFPartition.o output/myOS/kernel/mem/malloc.o output/userApp/main.o output/userApp/shell.o output/userApp/memTestCase.o -o output/myOS.elf make succeed

### 七、 运行和运行结果说明

运行指令qemu-system-i386 -kernel output/myOS.elf -serial stdio 运行指令sudo screen /dev/pts/0 运行结果如下如所示:

1. 内存检测

```
G pm@DESKTOP.75F7V98: -/workspace/ab4

MemStart: 0x1000000
lemd: 0x7f000000
lemd: 1057d0
dPartiion(start=0x1057d0, size=0x7efa830, firstFreeStart=0x105848)
EMB(start=0x105848, size=0x7efa87b0, nextStart=0x0)
dPartiion(start=0x1057d0, size=0x7efa830, firstFreeStart=0x1057d8)
EMB(start=0x1057d8, size=0x7efa820, nextStart=0x0)
START RUNNINS....
IJX->cnd
cmd
list all registered commands:
command name: description
cmd: list all registered commands
help: help [cmd]
tostMalloc1: Malloc, write and read.
nextMalloc2: Malloc, write and read.
nextMalloc3: India dPartiion(size=0x100). A:B:C: ==> -B:C: ==> -C: ==> testdP3: Init a dPartiion(size=0x100). A:B:C: ==> -B:C: ==> -C: ==> testdP3: Init a dPartiion(size=0x100). A:B:C: ==> -B:C: ==> -C: ==> testdP3: Init a dPartiion(size=0x100). A:B:C: ==> -B:C: ==> -C: ==> testdP3: Init a dPartiion(size=0x100). A:B:C: ==> -B:C: ==> -C: ==> testdP3: Init a dPartiion(size=0x100). A:B:C: ==> -B:C: ==> -C: ==> testdP3: Init a dPartiion(size=0x100). A:B:C: ==> -B:C: ==> -C: ==> testdP3: Init a dPartiion(size=0x100). A:B:C: ==> -B:C: ==> -C: ==> testdP3: Init a dPartiion(size=0x100). A:B:C: ==> -B:C: ==> -C: ==> testdP3: Init a dPartiion(size=0x100). A:B:C: ==> -B:C: ==> -B
```

2. 功能函数:

```
We had successfully malloc() a small memBlock (size=0x100, addr=0x105f50)
  It is initialized as a very small dPartition;
 EMB(start=0x105f58, size=0xf0, nextStart=0x0)

Alloc a memBlock with size 0x10, success(addr=0x105f58)!....Relaesed;

Alloc a memBlock with size 0x20, success(addr=0x105f58)!....Relaesed;

Alloc a memBlock with size 0x40, success(addr=0x105f58)!....Relaesed;

Alloc a memBlock with size 0x80, success(addr=0x105f58)!....Relaesed;

Alloc a memBlock with size 0x80, success(addr=0x105f58)!....Relaesed;

Alloc a memBlock with size 0x100, failed!
  Now, converse the sequence.
  Alloc a memBlock with size 0×100, failed!
 Alloc a memBlock with size 0x80, success(addr=0x105f58)!.....Relaesed; Alloc a memBlock with size 0x40, success(addr=0x105f58)!.....Relaesed; Alloc a memBlock with size 0x20, success(addr=0x105f58)!.....Relaesed;
  Alloc a memBlock with size 0x10, success(addr=0x105f58).....Relaesed;
   in@DESKTOP-7SE7V9B: ~/workspace/lab4
                                                                                                                                                                     П
 Alloc a memBlock with size 0x10, success(addr=0x105d38)!.....Relaesed;
TJX->testdP2
 We had successfully malloc() a small memBlock (size=0x100, addr=0x105e40);
 It is initialized as a very small dPartition;
dPartition(start=0x105e40, size=0x100, firstFreeStart=0x105e48)
dPartition(start=0x105e40, size=0x100, firstFreeStart=0x105e48)
EMB(start=0x105e48, size=0xf0, nextStart=0x0)
Now, A:B:C:- ==> -:B:C:- ==> -:C- ==> -
Alloc memBlock A with size 0x10: success(addr=0x105e48)!
dPartition(start=0x105e40, size=0x100, firstFreeStart=0x105e68)
EMB(start=0x105e68, size=0xd0, nextStart=0x0)
Alloc memBlock B with size 0x20: success(addr=0x105e68)!
dPartition(start=0x105e40, size=0x100, firstFreeStart=0x105e98)
EMB(start=0x105e98, size=0x30, nextStart=0x0)
Alloc memBlock C with size 0x30: success(addr=0x105e98)!
dPartition(start=0x105e40, size=0x100, firstFreeStart=0x105e48)
EMB(start=0x105e48, size=0x300, firstFreeStart=0x105e48)
EMB(start=0x105e48, size=0x60, nextStart=0x0)
 EMB(start=0x105ed8, size=0x60, nextStart=0x0)
Now. release A.
   jin@DESKTOP-7SF7V9B: ~/workspace/lab4
Alloc a memBlock with size 0x10, success(addr=0x105f58)!.....Relaesed;
 TJX->testdP3
 testdP3
 We had successfully malloc() a small memBlock (size=0x100, addr=0x106060
we had successfully malloc() a small memblock (size=0x100, addriti in initialized as a very small dPartition;
dPartition(start=0x106060, size=0x100, firstFreeStart=0x106068)
EMB(start=0x106068, size=0xf0, nextStart=0x0)
Now, A:B:C:- ==> A:B:- ==> A:- ==> -
Alloc memBlock A with size 0x10: success(addr=0x106068)!
dPartition(start=0x106068, size=0x100, firstFreeStart=0x106088)
FWB(start=0x106088, size=0xd0, nextStart=0x0)
 EMB(start=0x106088, size=0xd0, nextStart=0x0)
Alloc memBlock B with size 0x20: success(addr=0x106088)!
dPartition(start=0x106060, size=0x100, firstFreeStart=0x1060b8)
at at tribinstat to 100000, size=0x10, firstFreeStart=0x1000008/
EMB(start=0x106008, size=0xa0, nextStart=0x0)
Alloc memBlock C with size 0x30: success(addr=0x106008)!
dPartition(start=0x106060, size=0x100, firstFreeStart=0x1060f8)
 EMB(start=0x1060f8, size=0x60, nextStart=0x0)
Now, release C.
 ₫ jin@DESKTOP-7SF7V9B: ~/workspace/lab4
  EEB(start=0x105d99, next=0x0)
EEB(start=0x105d99, next=0x0)
Alloc memBlock C, start = 0x105d7a: 0xccccccc
eFPartition(start=0x105d30, totalN=0x4, perSize=0x20, firstFree=0x105d99)
EEB(start=0x105d99, next=0x0)
Alloc memBlock D, start = 0x105d99: 0xdddddddd
eFPartition(start=0x105d30, totalN=0x4, perSize=0x20, firstFree=0x0)

[Alloc memBlock E, failed!
eFPartition(start=0x105d30, totalN=0x4, perSize=0x20, firstFree=0x6000ff53)
EEB(start=0xf000ff53, next=0x0)
Nnww release A.
  Now, release A.
  eFPartition(start=0x105d30, tota1N=0x4, perSize=0x20, firstFree=0x105d3c)
  EEB(start=0x105d3c, next=0xaaaaaaaa)
  EEB(start=0xaaaaaaaa, next=0x0)
 Now, release B.
 eFFartition(start=0x105d30, tota1N=0x4, perSize=0x20, firstFree=0x105d3c)
EEB(start=0x105d3c, next=0x105d5c)
EEB(start=0x105d5c, next=0xbbbbbb)
  Now, release C.
  eFPartition(start=0x105d30, tota1N=0x4, perSize=0x20, firstFree=0x105d3c)
 EEB(start=0x105d3c, next=0x105d5c)
EEB(start=0x105d5c, next=0x105d7c)
  EEB(start=0x105d7c, next=0xcccc)
  EEB(start=0xcccc, next=0x0)
  Now, release D.
```

# 八、 遇到的问题和解决办法

问题: mentest使用short指针读写失败;

解决:询问了朱同学,利用int指针解决了这一问题。