Finding the inverse of a polynomial in the finite field $\mathbb{Z}_p/(f(x))$:

Recall the algorithm from your earlier project or from the slides 8-12 of sec 2.3 to find gcd of two integers.

We write a table with 7 columns. In the first row, the first 4 numbers are always 1, 0, 0, 1 (in that order), while the 5th and 6th numbers are the two numbers for which one seeks the gcd, and the 7th is always zero. Now we choose 5th column a prime number (here 271) and so gcd will be 1.

u_1	v_1	u_2	v_2	u_3	v_3	q
1	0	0	1	271	192	0

Then we add a new row using the following rules:

- The new q is the greatest integer less than or equal to the quotient of the old u_3 and v_3 .
- The new u_i is the old v_i .
- The new $v_i = \text{old } u_i (\text{ current } q)(\text{ old } v_i)$
- We do this multiple times, until we produce a row where $v_3 = 0$

We get **Division Algorithm**:

u_1	v_1	u_2	v_2	u_3	v_3	q
1	0	0	1	271	192	0
0	1	1	-1	192	79	1
1	-2	-1	3	79	34	2
-2	5	3	-7	34	11	2
5	-17	-7	24	11	1	3
-17	*	24	*	1	0	11

Now
$$\gcd(271, 192) = 1$$
 and $(271)(-17) + (192)(24) = 1 \implies (192)^{-1} = 24$ in mod 271.

In order to find the inverse in a finite field, we need to repeat the above algorithm.

For example let us consider the finite filed $\mathbb{Z}_2/(f(x))$ where $f(x) = x^4 + x + 1$.

Suppose that $g(x) = x^3 + x^2 + 1$ which corresponds to the 4-bit 1101. In order to find the inverse of g(x) in this field, we simply repeat above division algorithm starting with f(x) and g(x), but keep in mind that we are dealing with \mathbb{Z}_2 and so $-x^n$ is same as x^n and $2x^n = 0$.

Division Algorithm for Polynomials:

u_1	v_1	u_2	v_2	u_3	v_3	q
1	0	0	1	$x^4 + x + 1$	$x^3 + x^2 + 1$	0
0	1	1	x+1	$x^3 + x^2 + 1$	x^2	x+1
1	x+1	x+1	x^2	x^2	1	x+1
*	*	x^2	*	1	0	x^2

Therefore, the inverse of g(x) is x^2 . In other words, $(1101)^{-1} = 0100$.

Now let us consider the finite filed $\mathbb{Z}_2/(f(x))$ where $f(x)=x^8+x^4+x^3+x+1$.

Suppose that $g(x) = x^5 + x^4 + x + 1$. In order to find the inverse of g(x) in this field, we again repeat above division algorithm starting with f(x) and g(x), but keep in mind that we are dealing with \mathbb{Z}_2 and so $-x^n$ is same as x^n and $2x^n = 0$.

Division Algorithm for Polynomials:

u_1	v_1	u_2	v_2	u_3	v_3	q
1	0	0	1	$x^8 + x^4 + x^3 + x + 1$	$x^5 + x^4 + x + 1$	0
0	1	1	$x^3 + x^2 + x + 1$	$x^5 + x^4 + x + 1$	$x^4 + x^3 + x$	$x^3 + x^2 + x + 1$
1	x	$x^3 + x^2 + x + 1$	$x^4 + x^3 + x^2 + x + 1$	$x^4 + x^3 + x$	$x^2 + x + 1$	x
x	$x^3 + x + 1$	$x^4 + x^3 + x^2 + x + 1$	$x^6 + x^5 + x^3 + x^2$	$x^2 + x + 1$	1	$x^2 + 1$
*	*	$x^6 + x^5 + x^3 + x^2$	*	1	0	$x^2 + x + 1$

For example,

3rd row
$$v_2$$
 = 2nd row u_2 – (3rd row q)(2nd row v_2)
= 1 – $(x)(x^3 + x^2 + x + 1)$
= 1 – x^4 – x^3 – x^2 – x
= x^4 + x^3 + x^2 + x + 1

Ultimately u_3 will become 1 and the inverse is the u_2 in the corresponding row. So

$$(x^5 + x^4 + x + 1)^{-1} = x^6 + x^5 + x^3 + x^2$$

Exercise: Show that in the above field:

$$(x^6 + x^4 + x + 1)^{-1} = x^7 + x^6 + x^3 + x$$

If you prefer to write a program to find the inverse, see next page.

If you would like to write a program to find the inverse, here is an algorithm:

Extended Euclidean Algorithm to find the inverse of a polynomial a(x) in the finite field $\mathbb{Z}_p/(f(x))$:

INPUT: A nonzero binary polynomial a of degree at most m-1.

OUTPUT: $a^{-1} \mod f$.

- 1. $u \leftarrow a, v \leftarrow f$.
- 2. $g_1 \leftarrow 1, g_2 \leftarrow 0.$
- 3. While $u \neq 1$ do
- 3.1. $j \leftarrow \deg(u) \deg(v)$.
- 3.2. If j < 0 then: $u \leftrightarrow v$, $g_1 \leftrightarrow g_2$, $j \leftarrow -j$.
- 3.3. $u \leftarrow u + z^j v$.
- 3.4. $g_1 \leftarrow g_1 + z^j g_2$.
- 4. Return (g_1) .

For Python or Perl version of algorithm see section 7.11 of this

lecture note.