

Fundamentals of Mathematical Logic by Hinman

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Contents

1 Propositional Logic and Other Fundamentals

1.1 The Propositional Language

Problem 10

Give a precise definition of the set of sentences using infix notation and prove unique readability for your definition.

Solution:

Definition of infix sentences:

1. S_0 := the set of L -atomic sentences.
2. For each $n \in \omega$,

$$\begin{aligned} S_{n+1} := & S_n \cup \{\neg\phi : \phi \in S_0\} \cup \{(\neg\phi) : \phi \in S_n \sim S_0\} \\ & \cup \{(\phi) \cdot (\psi) : \phi, \psi \in S_n \sim S_0, \bullet \in \{\wedge, \vee, \rightarrow, \leftrightarrow\}\} \\ & \cup \{(\phi) \cdot \psi : \phi \in S_n \sim S_0, \psi \in S_0, \bullet \in \{\wedge, \vee, \rightarrow, \leftrightarrow\}\} \\ & \cup \{\phi \cdot (\psi) : \phi \in S_0, \psi \in S_n \sim S_0, \bullet \in \{\wedge, \vee, \rightarrow, \leftrightarrow\}\} \\ & \cup \{\phi \cdot \psi : \phi, \psi \in S_0, \bullet \in \{\wedge, \vee, \rightarrow, \leftrightarrow\}\} \end{aligned}$$

3. Let $S_f := \bigcup_{n \in \omega} S_n$.

Definition 1 (Parenthesization of binary connectives). *Let ϕ and ψ be sentences, and let \bullet be any binary connective. Then the four possible parenthesizations of $\phi \bullet \psi$ are:*

1. $(\phi) \cdot (\psi)$
2. $(\phi) \cdot \psi$
3. $\phi \cdot (\psi)$
4. $\phi \cdot \psi$

Proposition 1 (Unique readability for propositional sentences). *For any sentence $\theta \in S_f$, exactly one of the following holds:*

1. θ is atomic;
2. for some sentence ϕ , $\theta = \neg\phi$;
3. θ is a **conjunction** $\phi \wedge \psi$ for some sentences ϕ and ψ , with one of the four parenthesizations from Definition ??;

4. θ is a **disjunction** $\phi \vee \psi$ for some sentences ϕ and ψ , with one of the four parenthesizations from Definition ??;
5. θ is an **implication** $\phi \rightarrow \psi$ for some sentences ϕ and ψ , with one of the four parenthesizations from Definition ??;
6. θ is a **bi-implication** $\phi \leftrightarrow \psi$ for some sentences ϕ and ψ , with one of the four parenthesizations from Definition ??.

Proof of unique readability:

Proof.

■

Problem 11

For any expression $\phi = s_0 \dots s_k$, a **proper initial segment** of ϕ is any sequence of symbols $s_0 \dots s_l$ with $l < k$. Prove that no proper initial segment of a sentence is a sentence, and show that this can be used as an alternative to (4) as a technical lemma for the proof of Proposition 1.1.5.

Problem 12

Give a careful proof of Proposition 1.1.9, and show how Theorem 1.1.7 is an application of it.