Computer Architecture Tutorial 4

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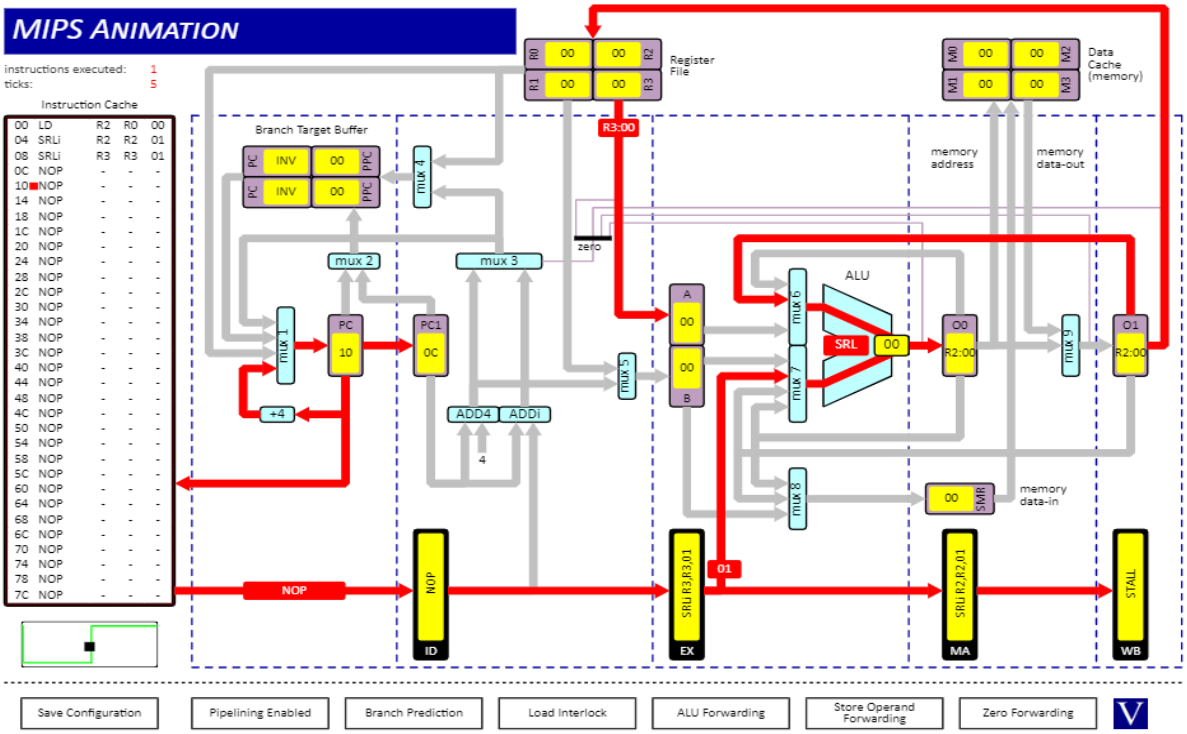
Q1.

a)

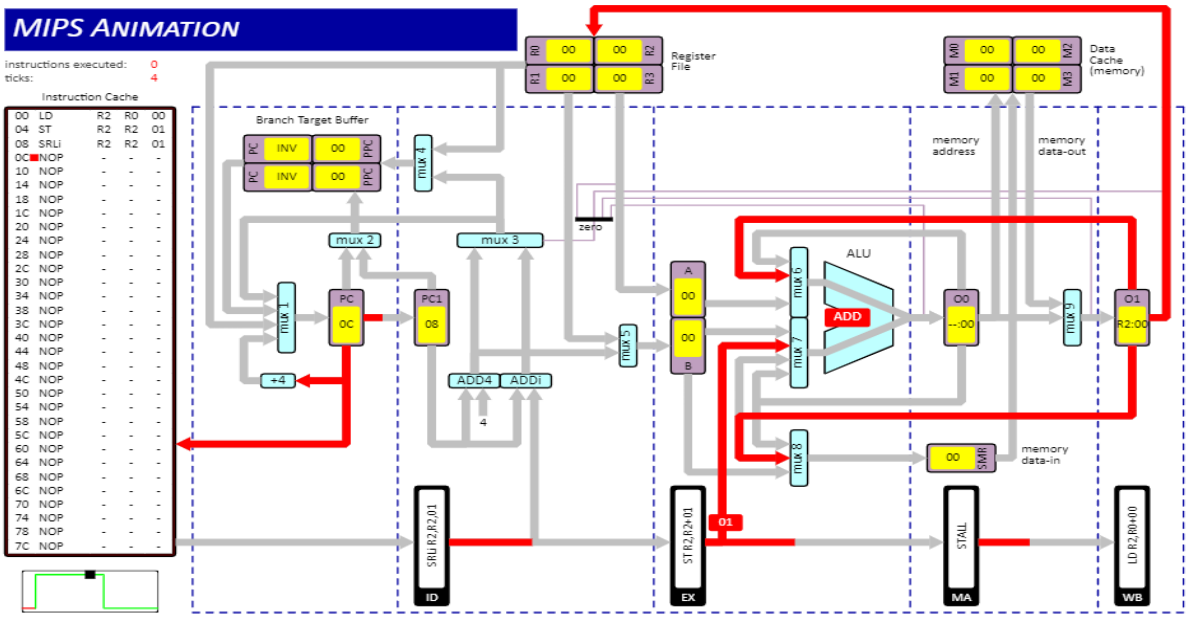


LD R2, R0, 00

SRLi R2, R2, 01

SRLi R3, R3, 01

b)

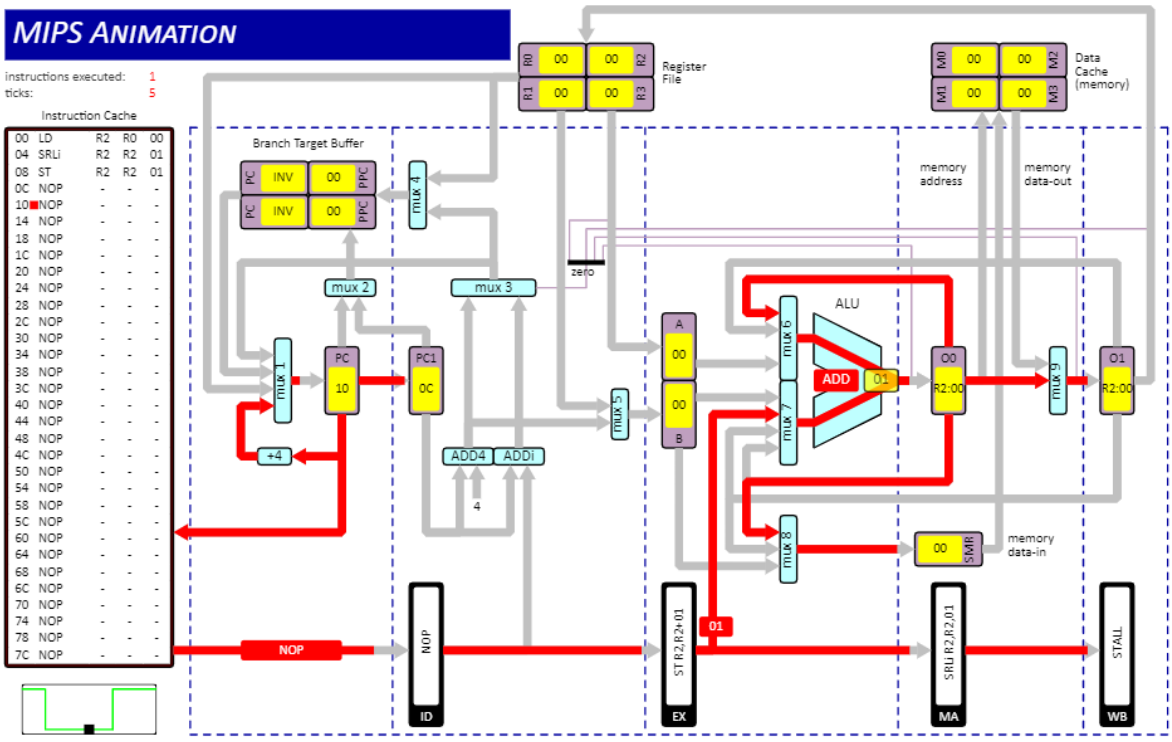


LD R2, R0, 00

ST R2, R2, 01

SRLi R2, R2, 01

c)

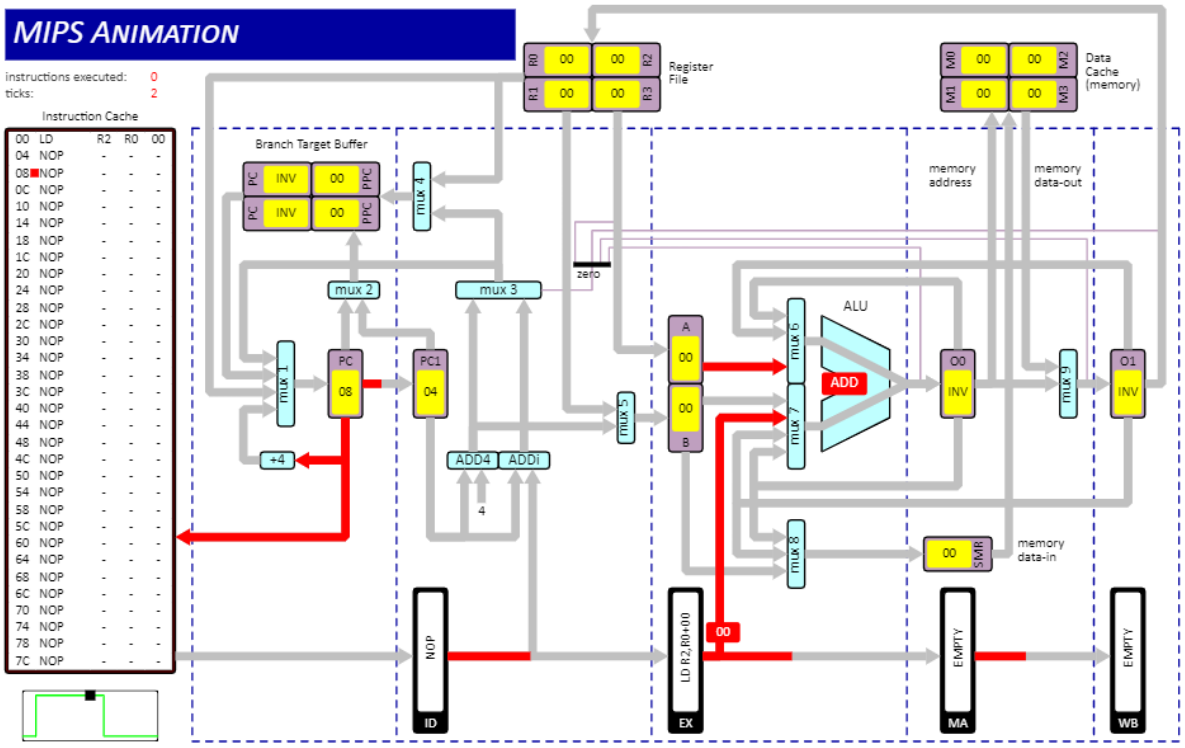


LD R2, R0, 00

SRLi R2, R2, 01

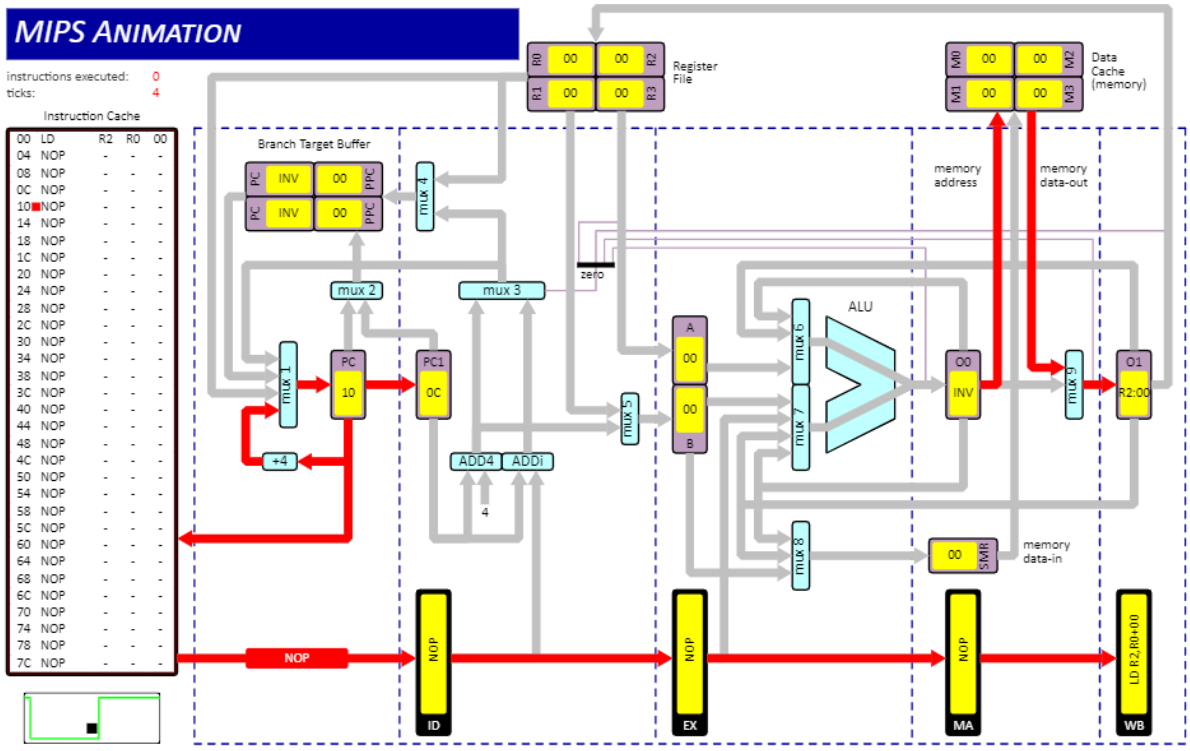
ST R2, R2, 01

d)



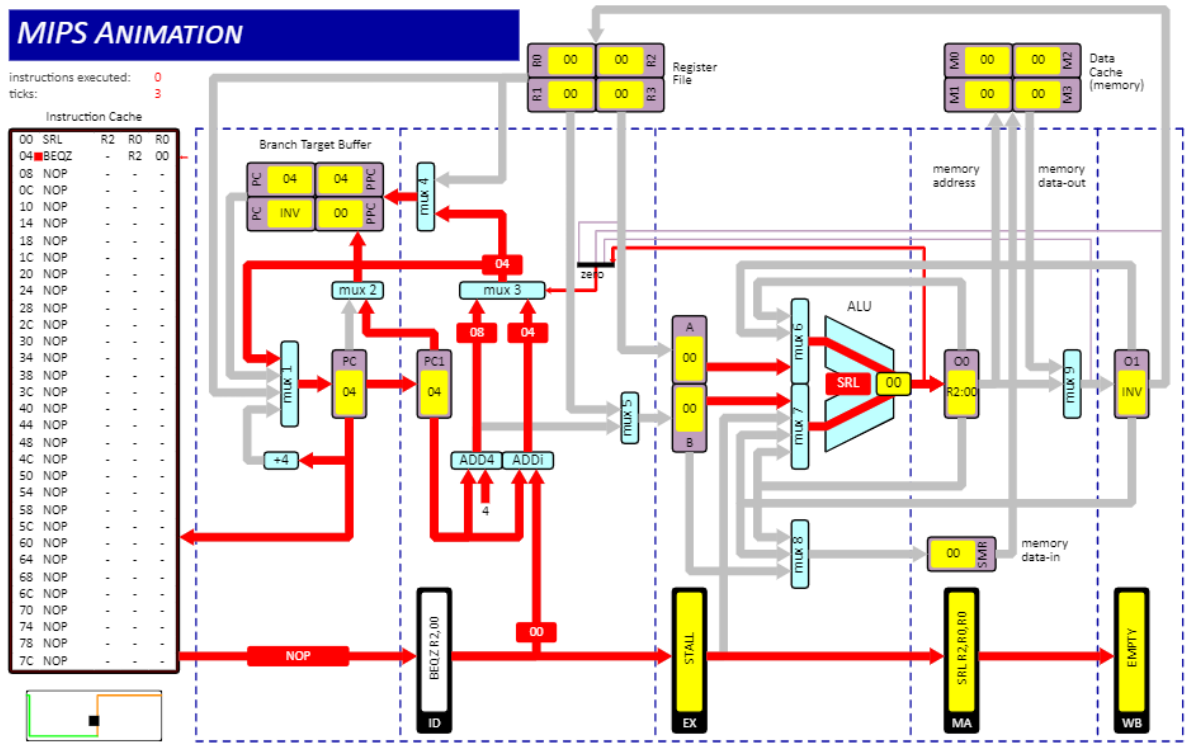
LD R2, R0, 00

e)



LD R2, R0, 00

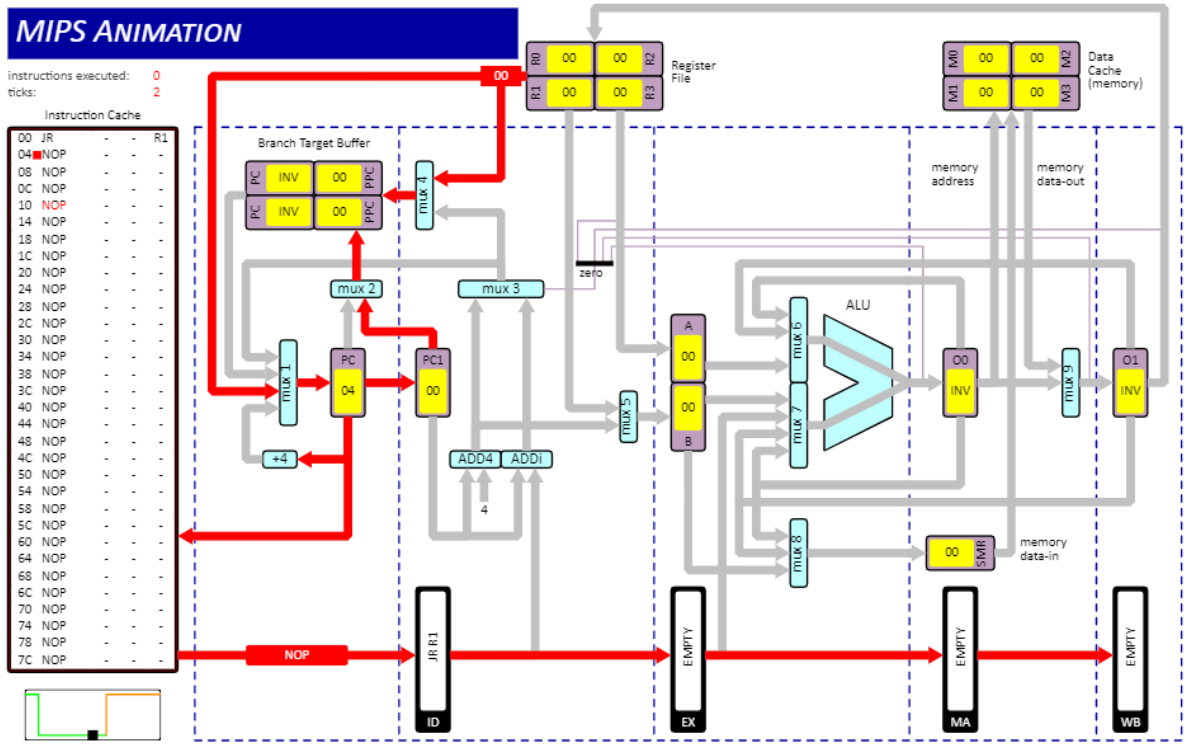
f)



SRL R2, R0, R0

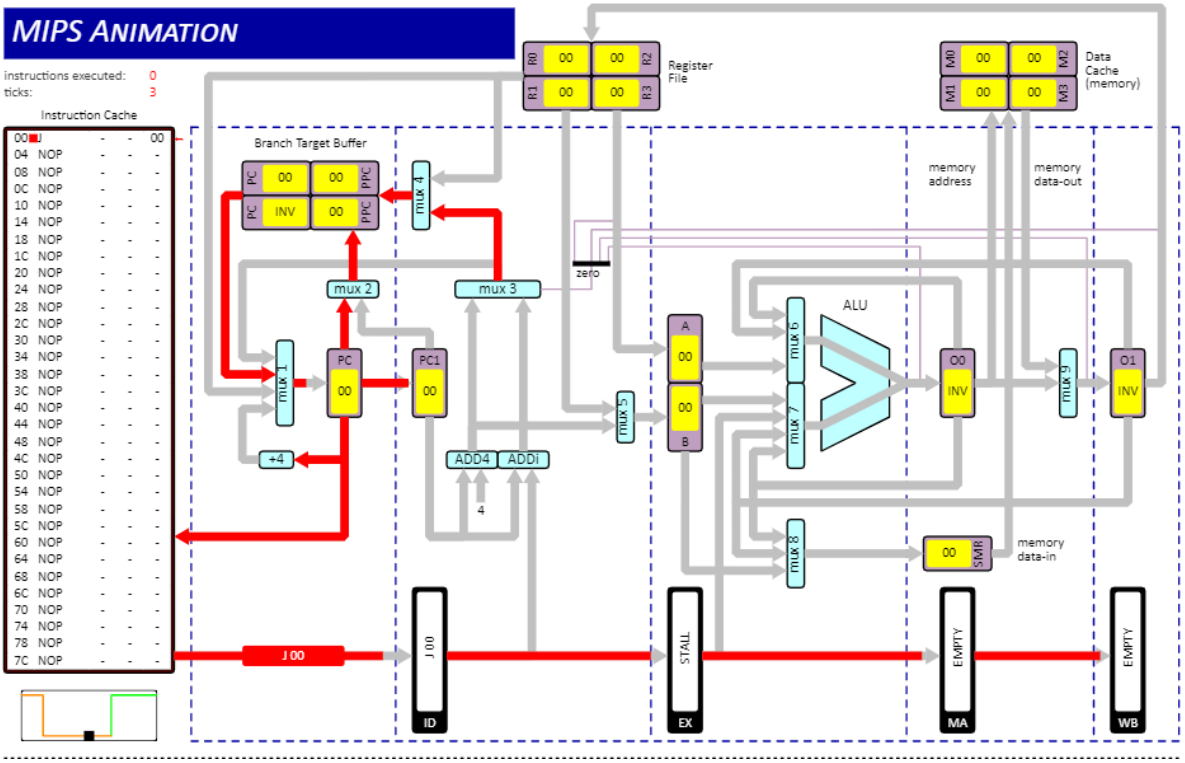
BEQZ R2, 00

g)



JR R1

h)



J 00

Q2.

(i)

R1 result = 0x15

Clock cycles = 10

(ii)

R1 result = 0x15

Clock cycles = 18

(iii)

R1 result = 0x06

Clock cycles = 10

In (i) the instruction is executed and the intermediate result is passed back to the ALU through the immediate registers for the next instruction while also being sent to the register file. In (ii) pipeline stalls occur because the next instruction is waiting for the answer from the previous instruction due to data dependency. In (iii) there is no data dependency or immediate registers to store results from previous instructions so the instructions occur as independent operations. The instructions keep reading values from the registers that are not yet updated.

Q3.

(i)

Instructions executed: 38

Clock cycles: 50

The pipeline takes the first four clock cycles to execute the 1st instruction. After this, 8 pipeline stalls occur. These stalls occur mainly in relation to all calls of the LD instruction (3 times) as this is a memory access instruction. Stalls also occur for SLLi (2 times), ANDi (2 times) and J (1 time).

(ii)

Instructions executed: 38

Clock cycles: 53

The pipeline takes the first four clock cycles to execute the 1st instruction. 10 pipeline stalls occur. These stalls occur mainly in relation to all calls of the LD instruction (4 times) as this is a memory access instruction. Stalls also occur for all calls of SLLi (4 times), one call of ANDi and one call of J. The lack of branch prediction is the cause for an increase in pipeline stalls, especially in relation to ANDi.

(iii)

Instructions executed: 38

Clock cycles: 46

The pipeline takes the first four clock cycles to execute the 1st instruction. 4 pipeline stalls occur. The stalls occur mainly in relation to calls of ANDi (2 times), SRLi (1 time) and J (1 time). The stalls for ANDi relate to the following instruction BEQZ checking the value of R2 immediately after ANDi operates on R2, mean BEQZ must wait for the result to be stored in R2. The stalls are fewer because having SLL occur first allows for the value of R2 to be updated for SRL as SLL operates on R3.