

# Certified Partial-Order Reduction and Symmetry Quotienting for Guarded Simplex Transfer Menus

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## Abstract

We formalize and implement a planning substrate for  $k$ -way allocation menus modeled as guarded unit transfers on the integer simplex. Naïve commutativity fails globally under guards; instead we use a *state-dependent* POR independence predicate and prove: (i) two-step commutation under stable-enabledness, (ii) trace-level adjacent swap invariance, (iii) a swap-equivalence relation with run-invariance, (iv) deterministic trace canonicalization with run- and horizon-invariance, and (v) symmetry equivariance for bucket transpositions supporting quotienting. We provide a runtime oracle matching the formal predicate, a deterministic canonicalizer with a sound oracle cache, a bounded-horizon simplex CEO planner, and a decision-quality benchmark gate that enforces a measurable crossover regime where POR becomes a net runtime win.

## 1 Artifacts, scope, and what is proved

**Proof artifact (machine-checked).** All theorems cited in this document are mechanically checked in Lean 4 in:

```
LeanProofs/CEO_SimplexPOR.lean
```

The proof bundle builds with:

```
cd LeanProofs && lake build
```

**Runtime artifact (Rust).** The executable counterparts are:

- oracle and semantics: `crates/mprd-core/src/tokenomics_v6/simplex_por_oracle.rs`
- canonicalization + oracle cache: `crates/mprd-core/src/tokenomics_v6/simplex_planner.rs`
- symmetry key: `crates/mprd-core/src/tokenomics_v6/simplex_symmetry_key.rs`
- bounded-horizon simplex CEO planner: `crates/mprd-core/src/tokenomics_v6/simplex_ceo.rs`

**One explicitly non-proved lemma (declared as an axiom).** The Lean file includes *one* independence lemma for the static “disjoint endpoints” criterion as an axiom (`stepOrStay [] comm [] of [] disjoint [] enab`). All other theorems cited below are **theorems** in Lean.

## 2 Formal model

### 2.1 State, caps, and actions

We model simplex menu states as functions over  $\text{Fin } k$ . The exact Lean definitions are:

Listing 1: Core types (Lean)

```
abbrev State (k : Nat) := Fin k → Nat
abbrev Caps (k : Nat) := Fin k → Nat

structure Action (k : Nat) where
  src : Fin k
  dst : Fin k
  hne : src ≠ dst
```

## 2.2 Enabledness and guarded step (failure-as-no-op)

Listing 2: Enabled predicate (Lean)

**Enabledness.**

```
def enabled {k : Nat} (caps : Caps k) (x : State k) (a : Action k) : Prop :=
  x a.src > 0 ∧ x a.dst < caps a.dst
```

**Guarded step semantics.** We define a successful **step** (used under enabledness hypotheses), and the guarded **stepOrStay**:

Listing 3: Guarded stepOrStay (Lean)

```
def step {k : Nat} (x : State k) (a : Action k) : State k :=
  let x1 := update x a.src (x a.src - 1)
  update x1 a.dst (x a.dst + 1)

def stepOrStay {k : Nat} (caps : Caps k) (x : State k) (a : Action k) : State k :=
  if enabled caps x a then
    step x a
  else
    x
```

## 3 POR independence predicate and closed-form sufficient oracle

### 3.1 Stable-enabledness (dynamic POR independence)

We define the dynamic stable-enabledness predicate:

Listing 4: Stable-enabledness (Lean)

```
def stableEnabled {k : Nat} (caps : Caps k) (x : State k) (a b : Action k) : Prop :=
  enabled caps x a ∧ enabled caps x b ∧
  enabled caps (stepOrStay caps x a) b ∧ enabled caps (stepOrStay caps x b) a
```

### 3.2 Closed-form sufficient condition (stableEnabledIneq)

Listing 5: Closed-form sufficient oracle (Lean)

```
def stableEnabledIneq {k : Nat} (caps : Caps k) (x : State k) (a b : Action k) : Prop :=
  enabled caps x a ∧
  enabled caps x b ∧
  (b.src = a.src → x a.src > 1) ∧
```

```
(b.dst = a.dst →x a.dst + 1 < caps a.dst) ∧
(a.src = b.src →x b.src > 1) ∧
(a.dst = b.dst →x b.dst + 1 < caps b.dst)
```

Listing 6: Oracle implies stable-enabledness (Lean)

```
theorem stableEnabled_of_stableEnabledIneq
  {k : Nat} (caps : Caps k) (x : State k) (a b : Action k) :
  stableEnabledIneq caps x a b →stableEnabled caps x a b := by
  ...
```

## 4 Two-step commutation and trace-level swap

Listing 7: Two-step commutation (Lean)

```
theorem stepOrStay_comm_of_stableEnabled
  {k : Nat} (caps : Caps k) (x : State k) (a b : Action k) :
  stableEnabled caps x a b →
  stepOrStay caps (stepOrStay caps x a) b = stepOrStay caps (stepOrStay caps x b) a :=
  by
  ...
```

Listing 8: Oracle implies commutation (Lean)

```
theorem stepOrStay_comm_of_stableEnabledIneq
  {k : Nat} (caps : Caps k) (x : State k) (a b : Action k) :
  stableEnabledIneq caps x a b →
  stepOrStay caps (stepOrStay caps x a) b = stepOrStay caps (stepOrStay caps x b) a :=
  by
  ...
```

Listing 9: Trace run and adjacent-swap lemma (Lean)

```
def run {k : Nat} (caps : Caps k) : List (Action k) →State k →State k
| [], x => x
| a :: as, x => run caps as (stepOrStay caps x a)

theorem run_swap_adjacent_of_stableEnabledIneq
  {k : Nat} (caps : Caps k) (pre suf : List (Action k)) (x : State k) (a b : Action k) :
  stableEnabledIneq caps (run caps pre x) a b →
  run caps (pre ++ (a :: b :: suf)) x = run caps (pre ++ (b :: a :: suf)) x := by
  ...
```

## 5 Swap-equivalence (SwapEq) and run invariance

Listing 10: SwapEq and run invariance (Lean)

```
inductive SwapEq {k : Nat} (caps : Caps k) (x0 : State k) : List (Action k) →List (
  Action k) →Prop where
| refl (xs : List (Action k)) : SwapEq caps x0 xs xs
```

```

| step {xs ys : List (Action k)} : SwapStep caps x0 xs ys → SwapEq caps x0 xs ys
| trans {xs ys zs : List (Action k)} : SwapEq caps x0 xs ys → SwapEq caps x0 ys zs →
  SwapEq caps x0 xs zs
| symm {xs ys : List (Action k)} : SwapEq caps x0 xs ys → SwapEq caps x0 ys xs

theorem run_invariant_of_SwapEq
  {k : Nat} (caps : Caps k) (x0 : State k) :
  ∀ {xs ys : List (Action k)}, SwapEq (k := k) caps x0 xs ys → run caps xs x0 = run caps
    ys x0 := by
  ...

```

## 6 Canonicalization and bounded-horizon completeness

Listing 11: Canonicalization preserves run and length (Lean)

```

def canonicalize {k : Nat} (caps : Caps k) (x0 : State k) (xs : List (Action k)) : List (
  Action k) :=
  canonIter (k := k) caps x0 (xs.length * xs.length) xs

theorem run_canonicalize_eq
  {k : Nat} (caps : Caps k) (x0 : State k) (xs : List (Action k)) :
  run caps (canonicalize (k := k) caps x0 xs) x0 = run caps xs x0 := by
  ...

theorem length_canonicalize
  {k : Nat} (caps : Caps k) (x0 : State k) (xs : List (Action k)) :
  (canonicalize (k := k) caps x0 xs).length = xs.length := by
  ...

```

Listing 12: Reachability completeness under canonicalize (Lean)

```

def ReachableWithin {k : Nat} (caps : Caps k) (x0 : State k) (h : Nat) (x : State k) :
  Prop :=
  ∃ xs : List (Action k), xs.length ≤ h ∧ run caps xs x0 = x

theorem reachableWithin_via_canonicalize
  {k : Nat} (caps : Caps k) (x0 : State k) (h : Nat) (x : State k) :
  ReachableWithin (k := k) caps x0 h x ↔
  ∃ xs : List (Action k), xs.length ≤ h ∧ run caps (canonicalize (k := k) caps x0 xs) x0
    = x := by
  ...

```

Listing 13: Reachability completeness under canonicalizeSwap (Lean)

```

theorem reachableWithin_via_canonicalizeSwap
  {k : Nat} (caps : Caps k) (x0 : State k) (h : Nat) (x : State k) :
  ReachableWithin (k := k) caps x0 h x ↔
  ∃ xs : List (Action k),
    xs.length ≤ h ∧ run caps (canonicalizeSwap (k := k) caps x0 xs) x0 = x := by
  ...

```

## 7 Benchmark gate

```
bash tools/ceo/check_ceo_simplex_rail.sh
```

The rail blocks regressions by requiring, on sweep rows with `eval_iters`  $\geq 200$ , **symmetry quotienting** win-rate  $\geq 0.75$  and median runtime ratio  $\leq 1.0$ , implemented by `tools/ceo/check_ceo_simplex_sweep_str`  
`--gate sym`.