

UNIVERSITY OF ST ANDREWS

CS4099

ILNP Routing for IoT

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1 NOTES

1.1 Fixing overhead

<https://ieeexplore.ieee.org/document/7993954> to fix route reply storm issue: add jitter before responding with route reply and listen for other replies

On receiving route request, spawn cache reply thread and return CacheReplyThread: wait jitter time, check if reply recently seen in queue, discard if seen, otherwise send. Need RecentRouteReplies: Fixed size FIFO buffer of (dest, last_{hop})

Stopping duplication of packets on forwarding though?

Abstract

Declaration

I declare that the material submitted for assessment is my own work except where credit is explicitly given to others by citation or acknowledgement. This work was performed during the current academic year except where otherwise stated. The main text of this project report is #TODO NN,NNN words long, including project specification and plan. In submitting this project report to the University of St Andrews, I give permission for it to be made available for use in accordance with the regulations of the University Library. I also give permission for the title and abstract to be published and for copies of the report to be made and supplied at cost to any bona fide library or research worker, and to be made available on the World Wide Web. I retain the copyright in this work.

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2 Introduction

Despite the imminent exhaustion of IPv4 addresses [1], IPv6 is still being adopted slowly [2]. Brittle solutions such as NAT are being used to temporarily expand the IP address space, and to avoid the transition costs involved in upgrading to IPv6. Whilst IPv6 does expand the address space greatly and introduces functionality such as multicast, the internet protocol itself suffers from many issues.

2.1 Issues with IP

IP addresses are used both to identify a system and to determine its topological location. [3] lists several of the downsides to this overloading of IP addresses, and why the protocol was still used despite these concerns.

The separation of concerns that should be achieved by a layered model is not possible, since the IP address is used by each layer in some way. IP addresses can be used in the application layer, and are bound to physical network interfaces, which goes against the end-to-end argument where each layer should provide an opaque abstraction to those above it.

The issues with IP are not just semantic. Due to the overloading of the IP address and the rapid increase in internet connected devices [4], the scalability of the system is being challenged. Implementations of multipath routing with the intention of balancing load are improving network performance for the operators that use them, but with IP it places greater stress on the default-free zone (DFZ) routing information base (RIB). Multihoming is also being used to improve reliability, but with IP this requires routing entries to store multiple addresses for one host. An IAB workshop [5] detailed how the DFZ RIB databases are growing in size exponentially due to the increasing number of devices and an inability to aggregate address prefixes. With IPv6 allowing for an even larger address space, this problem will only get worse.

Due to the growing number of Internet of Things (IoT) devices, mobility is also a necessary feature for a networking protocol. Mobile IP currently requires another entity (a home agent) to track and proxy packets to the mobile host as it moves from network to network. This mobility is also problematic for IPSec, which requires that the end system addresses remain fixed.

Given the difficulty involved in simply migrating from IPv4 to IPv6, it is

very doubtful that introducing an entirely different protocol for the internet would be successful. A backwards compatible solution would likely be the only solution that would be adopted within a reasonable time frame.

2.2 ILNP

Both multihoming and mobility are far simpler to implement and maintain if the identity and topological locator of a host are separated, and this is how the Identifier-Locator Network Protocol functions. [6] proposes ILNPv6, which implement ILNP with the same address space as IPv6 and the same packet structure as IPv6, but with different semantics for interpreting the addresses. ILNPv6 splits the original 128-bits used for an IPv6 address into two 64-bit fields: the upper bits representing the locator and the lower 64 bits representing the identifier. The version field in the IP header is used to differentiate between ILNPv6 and IPv6 packets, and routers that don't support ILNP can interpret the packets as IP without any issue.

The locator value identifies the subnetwork that a host belongs to, and a host can have multiple locator values, providing multihoming with smaller memory requirements for the RIB. The ID part of the address is unique to the host, and provides a fixed address which can be used at the transport layer.

Though ILNPv6 is very backwards-compatible, there are still some difficult challenges involved in its deployment. [7] describes how the tight coupling of the C socket API and IP addresses could cause issues in some legacy applications. ILNPv6 also requires some additions to the Domain Name System (DNS) to properly support multihoming.

2.3 Goal

The benefits of ILNP are clear for the entire internet infrastructure. The native support for multihoming and multipath routing is especially beneficial to typically mobile internet of things (IoT) devices. This project aims to exemplify these benefits for agricultural sensor networks in particular.

These devices are often restricted by limited battery life, memory, and computational capabilities. Most routing protocols focus on finding the shortest route between a source and destination, which often results in a small number of paths being heavily used and so some nodes are especially drained due to processing and forwarding of packets. This can result in a

network partition once crucial nodes fail (due to loss of battery), rendering a section of still operational nodes useless. By reducing the networking overhead and attempting to balance traffic across several paths, the network can remain operational for longer.

3 Context Survey

3.1 Internet Routing and Addressing Architecture

ILNP was listed alongside several other solutions to the current challenges facing the internet infrastructure in RFC6115 [8]. Locator-Identifier Split Protocol (LISP) also uses locator-identifier semantics and has already been deployed in 60 sites over 10 countries [9]. This ease of deployment is likely due to the fact that it requires no changes to hosts, and instead only required updates to core routers to be operational.

Another possible alternative is Routing Architecture for the Next Generation Internet (RANGI) [10]. RANGI again uses the identifier-locator split, but also has a cryptographic identifier which provides sender identification.

3.2 Ad Hoc Sensor Networks

3.3 Energy Efficient Routing Protocols

4 Requirements Specification

1. Describe requirements of resulting python library

5 Design

1. Component structure (socket interface, router/dsrservice/forwardingtable, raw sockets)
2. Runtime behaviour (packet parsing, routing, and forwarding)
3. Use figures to visualise project structure and workflow

6 Experiment

1. Discuss aim of experiment (to measure efficiency of the used routing protocol with ILNP, and compare to IP).
2. Explain case study, with reference to source (i.e. agricultural sensor setup)
3. Use visuals to show locators to real life position and sensor radi
4. Discuss experiment configuration (how machines were chosen, results collected, battery life simulated, etc)
5. discuss choice of metrics, justification and how to compare results.

7 Results and Discussion

1. Show heat map of results
2. Explain features of heat map
3. Describe the behaviour if IP was used instead through analysis
4. Discuss weaknesses with experiment

8 Conclusions

1. was the goal met, and if so how well?
2. future work with ILNP, possible suggestions of better alternatives to the routing protocol used.

9 Appendix

1. Instructions on installing, and executing and using the python module, and how to configure the experiments.

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