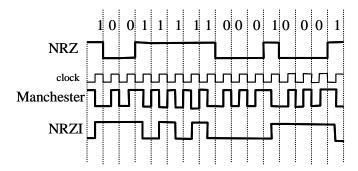
## ECE 4380/6380: Computer Communications Problem Set 2 Solutions

Each problem is worth 10 points

1. **Encoding**. Encoding for sequence 1001 1111 0001 0001



## 2. **Encoding** Chapter 2, number 4.

One can list all 5-bit sequences and count, but here is another approach: there are  $2^3$  sequences that start with 00, and  $2^3$  that end with 00. There are two sequences, 00000 and 00100, that do both. Thus, the number that do either is 8 + 8 - 2 = 14, and finally the number that do neither is 32 - 14 = 18. Thus there would have been enough 5-bit codes meeting the stronger requirement; however, additional codes are needed for control sequences.

## 3. Framing.

A 0 is stuffed in at the two places marked with the  $\underline{\mathbf{0}}$ .

- 4. **Framing** ..., DLE, DLE, DLE, ETX, ETX
- 5. Error detection with CRC. Show that (x + 1) is not a factor of C(x).

Note that the standard CRC-12 polynomial is  $x^{12} + x^{11} + x^3 + x^2 + x + 1$ , and as can be easily verified, (x + 1) is a factor.

6. **Error detection with CRC**. Calculate the CRC value of the bit sequence 0011 1010 1011, and list the message that should be transmitted

$$\begin{array}{r}
100000111 \overline{\smash) 001110101011.000000000} \\
\underline{100000111} \\
110100101 \\
\underline{100000111} \\
10100010 \\
0
\end{array}$$

The transmitted sequence: 0011 1010 1011 0110 0111

7. **Error detection with CRC**. Find a 13-bit burst error polynomial  $E(x) = x^{12} + ... + 1$  that cannot be detected by a CRC check (with CRC-8)...

An error polynomial cannot be detected if it is a multiple of C(x). We are looking for a polynomial P(x), such that C(x)P(x) = E(x). One simple approach is to select a P(x) that has the terms  $x^4$  and 1, so that we get the terms  $x^{12}$  and 1 in E(x).  $P(x) = x^4 + 1$  is one such polynomial, and

$$E(x) = (x^8 + x^2 + x + 1)(x^4 + 1)$$
  
=  $x^{12} + x^8 + x^6 + x^5 + x^4 + x^2 + x + 1$