

## Exercise Sheet 6

### Task 6.1 Propositional Logic

Recall the proof system of natural deduction defined by the four inference rules. Here is a proof that inconsistent assumptions  $p$  and  $\neg p$  imply any  $q$ .

$p \rightarrow \perp \in \{\dots\}$	$\frac{}{p, p \rightarrow \perp, q \rightarrow \perp \vdash p \rightarrow \perp}$	ASSUMP	$p \in \{\dots\}$	$\frac{}{p, p \rightarrow \perp, q \rightarrow \perp \vdash p}$	ASSUMP
	$p, p \rightarrow \perp, q \rightarrow \perp \vdash \perp$			$p, p \rightarrow \perp, q \rightarrow \perp \vdash \perp$	IMPELIM
	$p, p \rightarrow \perp \vdash \perp$			$p, p \rightarrow \perp \vdash \perp$	DOUBLENEG
	$p \vdash (p \rightarrow \perp) \rightarrow q$			$p \vdash (p \rightarrow \perp) \rightarrow q$	IMPINTRO
	$\vdash p \rightarrow (p \rightarrow \perp) \rightarrow q$			$\vdash p \rightarrow (p \rightarrow \perp) \rightarrow q$	IMPINTRO

- a) Use the rules to construct proof trees for the following formulas. Prove them, and also  $p \rightarrow (p \rightarrow \perp) \rightarrow q$ , in Coq using tactics. Finally, prove them by giving explicit proof terms. Use the given template file for the Coq proofs.

- (i)  $p \rightarrow (p \rightarrow \perp) \rightarrow \perp$
- (ii)  $(p \rightarrow q) \rightarrow (q \rightarrow r) \rightarrow (p \rightarrow r)$

- b) Extend the syntax of propositional logic by conjunction and disjunction. Extend the semantics (evaluation function). Extend the inference rules of the proof system. *Hint:* For both operators, you will need introduction and elimination rule(s), analogously to the two rules IMPINTRO, IMPELIM. If you get stuck, you may get inspiration from the internet.
- c) Prove the following formulas in your extended proof system. Then prove them in Coq both by using tactics and by giving a proof term (use again the template file).

- (i)  $(p \wedge q) \rightarrow (q \wedge p)$  “ $\wedge$  is commutative”
- (ii)  $p \wedge (q \vee r) \rightarrow (p \wedge q) \vee (p \wedge r)$  “ $\wedge$  is distributive over  $\vee$ ”

### Task 6.2 Programming in Coq

Consider the following `option` data type that provides two kinds of values, either `none` or `some v` for some value `v`. Like `list`, it is parameterized over a type `A`.

```
Inductive option (A: Type) : Type :=
| some : A → option A
| none : option A.
```

Write the following programs, using the provided template file.

- (i) `pred_opt: nat → option nat`  
This function returns the predecessor of a natural number, or `none` if it does not exist.
- (ii) `pred: nat → nat`  
This function returns the predecessor of a natural number, or 0 if it does not exist.

- (iii) `minus: nat → nat → option nat`  
This function computes  $x - y$  for two natural numbers. If the result is not a natural number, `none` is returned.
- (iv) `is_none: ∀(X: Type), option X → bool`  
This polymorphic function takes an optional value and returns `true` if the value is `none` and `false` otherwise.
- (v) `head: ∀(X: Type), list X → option X`  
This function returns the first element of a list, or `none` if the list is empty.

### Task 6.3 Proving in Coq

Prove the following propositions in Coq, using the provided template file.

- (i)  $\forall (X: \text{Type}) (x: X) (xs: \text{list } X), \text{head } (\text{cons } x \text{ } xs) = \text{some } x$
- (ii)  $\forall (n: \text{nat}), S (\text{pred } n) = n \vee n = 0$
- (iii)  $\forall (x \ y \ z: \text{nat}), (x + y) + z = x + (y + z)$
- (iv)  $\forall (m \ n: \text{nat}), \text{minus } m \ n = \text{none} \vee \text{minus } n \ m = \text{none} \vee m = n$