

# EECS3311 – Software Design

## Visitor Pattern

## A CFG for Arithmetic Expressions

An example grammar that generates strings representing arithmetic expressions with the four operators +, -, \*, /, and numbers as operands is:

1. <expression> --> number
2. <expression> --> ( <expression> )
3. <expression> --> <expression> + <expression>
4. <expression> --> <expression> - <expression>
5. <expression> --> <expression> \* <expression>
6. <expression> --> <expression> / <expression>

The only nonterminal symbol in this grammar is <expression>, which is also the start symbol. The terminal symbols are {+,-,\*/, (,),number}. (We will interpret "number" to represent any valid number.)

The first rule (or production) states that an <expression> can be rewritten as (or replaced by) a number. In other words, a number is a valid expression.

The second rule says that an <expression> enclosed in parentheses is also an <expression>. Note that this rule defines an expression in terms of expressions, an example of the use of recursion in the definition of context-free grammars.

The remaining rules say that the sum, difference, product, or division of two <expression>s is also an expression.

# A simple grammar

 $v ::= v\delta | \delta$  $\delta ::= 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9$ 

**V: Num → Int**

$$V[v\delta] = 10 \times V[v] + V[\delta]$$

$$V[0] = 0 \quad V[1] = 1$$

$$V[2] = 2 \quad V[3] = 3$$

$$V[4] = 4 \quad V[5] = 5$$

$$V[6] = 6 \quad V[7] = 7$$

$$V[8] = 8 \quad V[9] = 9$$

$$\begin{aligned} V[123] &= 10 \times V[12] + 3 \\ &= 10 \times (10 \times V[1] + 2) + 3 \\ &= 123 \end{aligned}$$

# A BNF grammar is a way to express a Context Free Grammar

```
expression = term | expression, "+" , term;
term       = factor | term, "*" , factor;
factor     = constant | variable | "(" , expression , ")";
variable   = "x" | "y" | "z";
constant   = digit , {digit};
digit      = "0" | "1" | "2" | "3" | "4" | "5" | "6" | "7" | "8" | "9";
```

- Each rule is called a production
- Green is non-terminal
- In quotes are terminals  
(Wikipedia)

FYI: not needed for the course

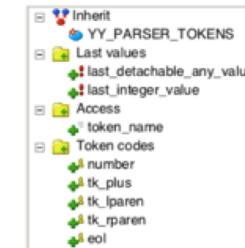
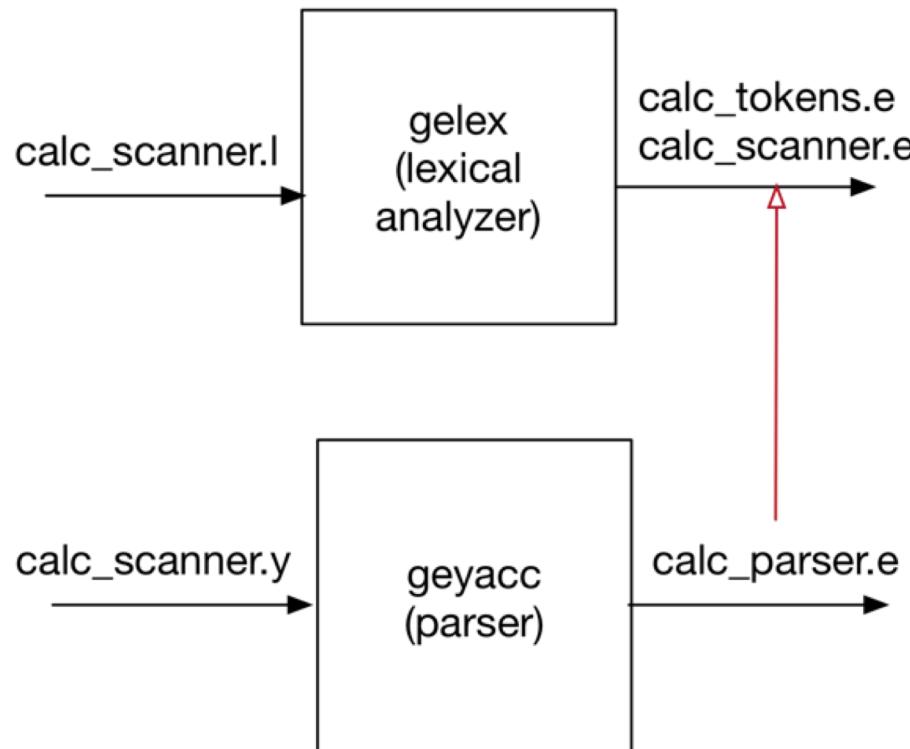
## Theory of Computation

### Lexer and Parser

Tutorial Documentation (login with `anonymous` and empty password).

The basic idea is as follows:

`text:STRING -- last token read`  
`last_integer_value:INTEGER`



# Syntax and Semantics

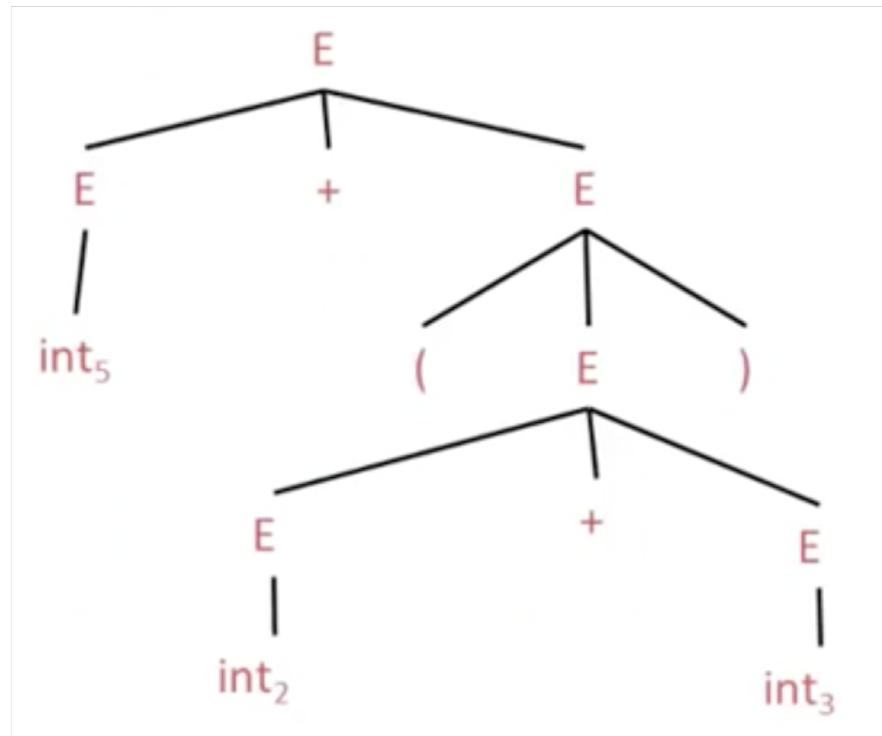
- **Syntax** is the **grammar**. It describes the way to construct a correct sentence.
  - For example, *this water is triangular* is syntactically correct (<subject><predicate>)
  - A program that is syntactically correct will compile and run.
- **Semantics** relates to the **meaning**.
  - *this water is triangular* does not mean anything, though the grammar is ok.
  - A program that is semantically correct will actually do what you as the programmer intended it to do (it will also not crash or have bugs/errors)

# Grammar (for expression E)

- $E \rightarrow \text{Int} \mid (E) \mid E + E$

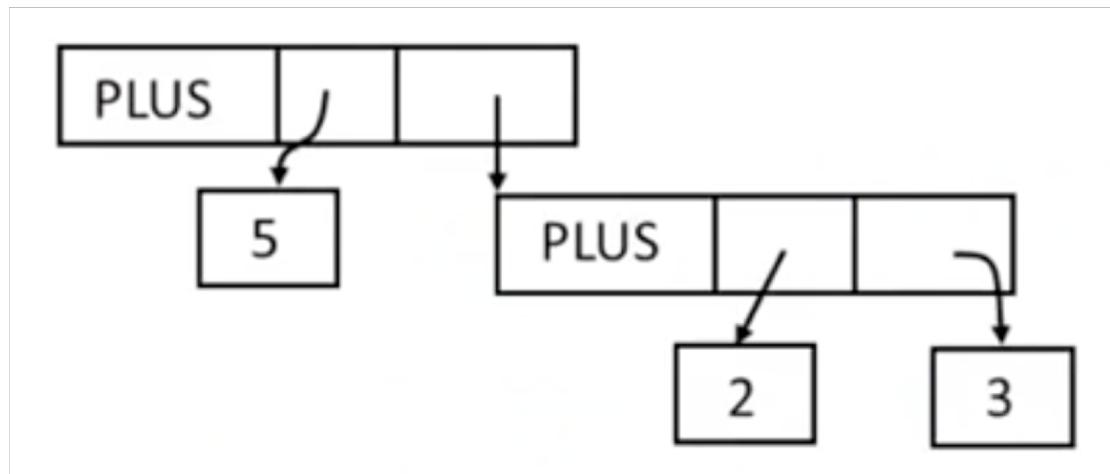
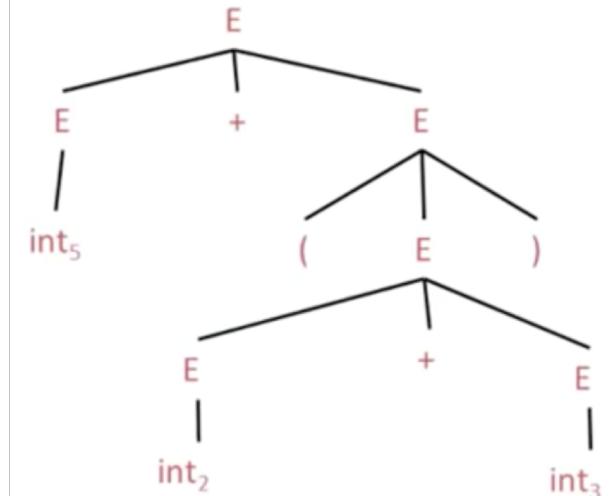
# Parse Tree

- $E \rightarrow \text{Int} \mid (E) \mid E + E$
- $5 + (2 + 3)$

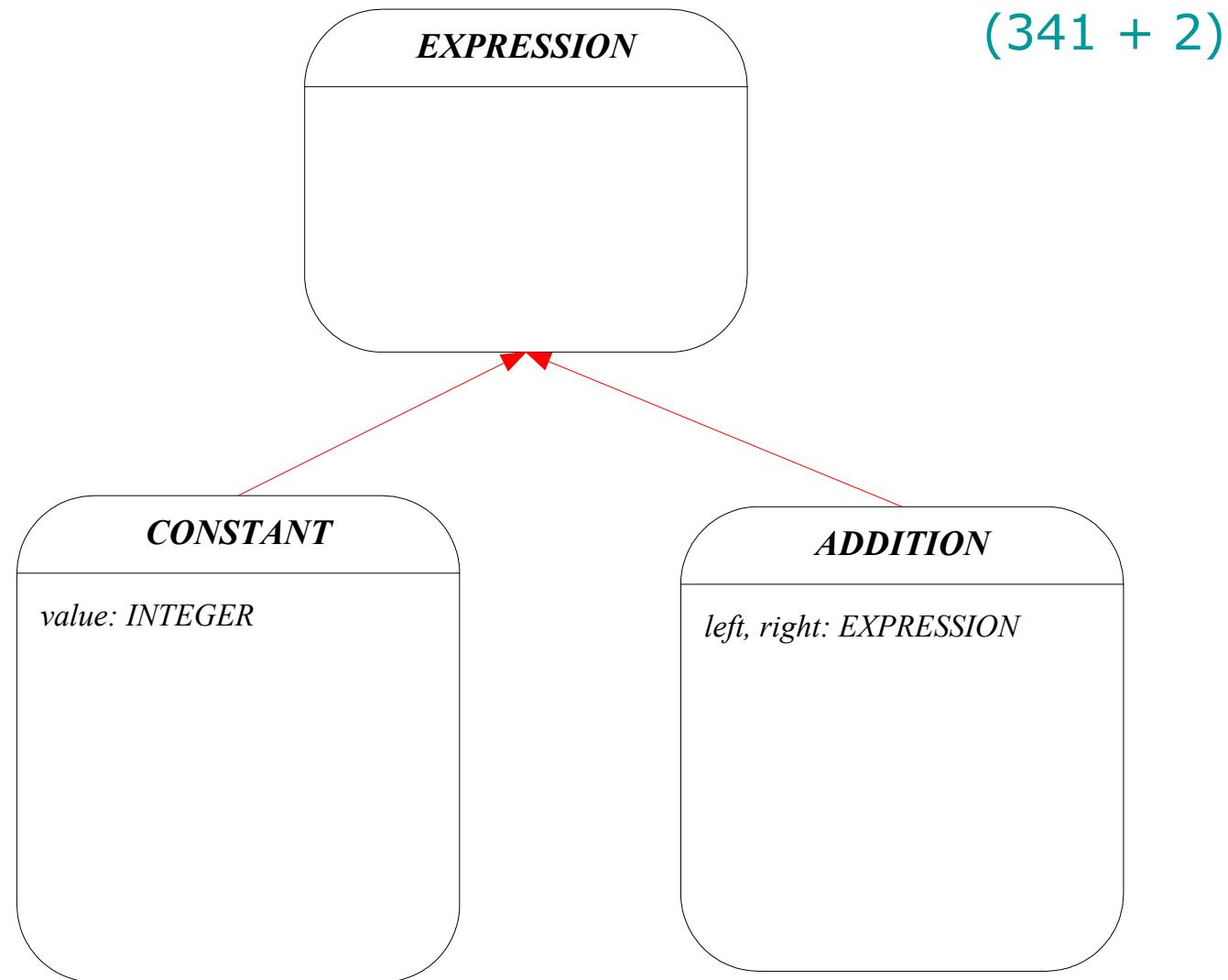


# AST (Abstract Syntax Tree)

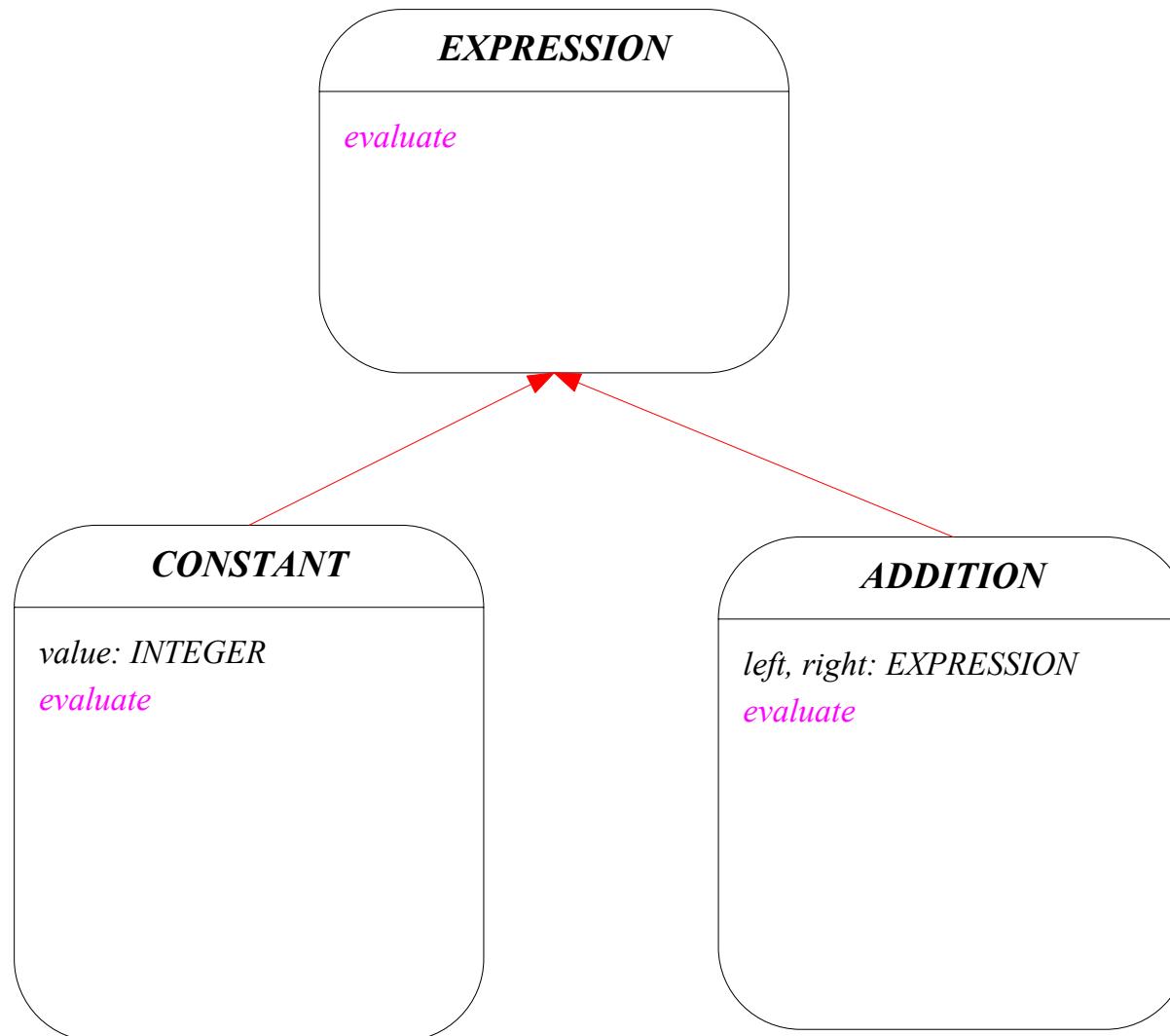
- $E \rightarrow \text{Int} \mid (E) \mid E + E$
- $5 + (2 + 3)$



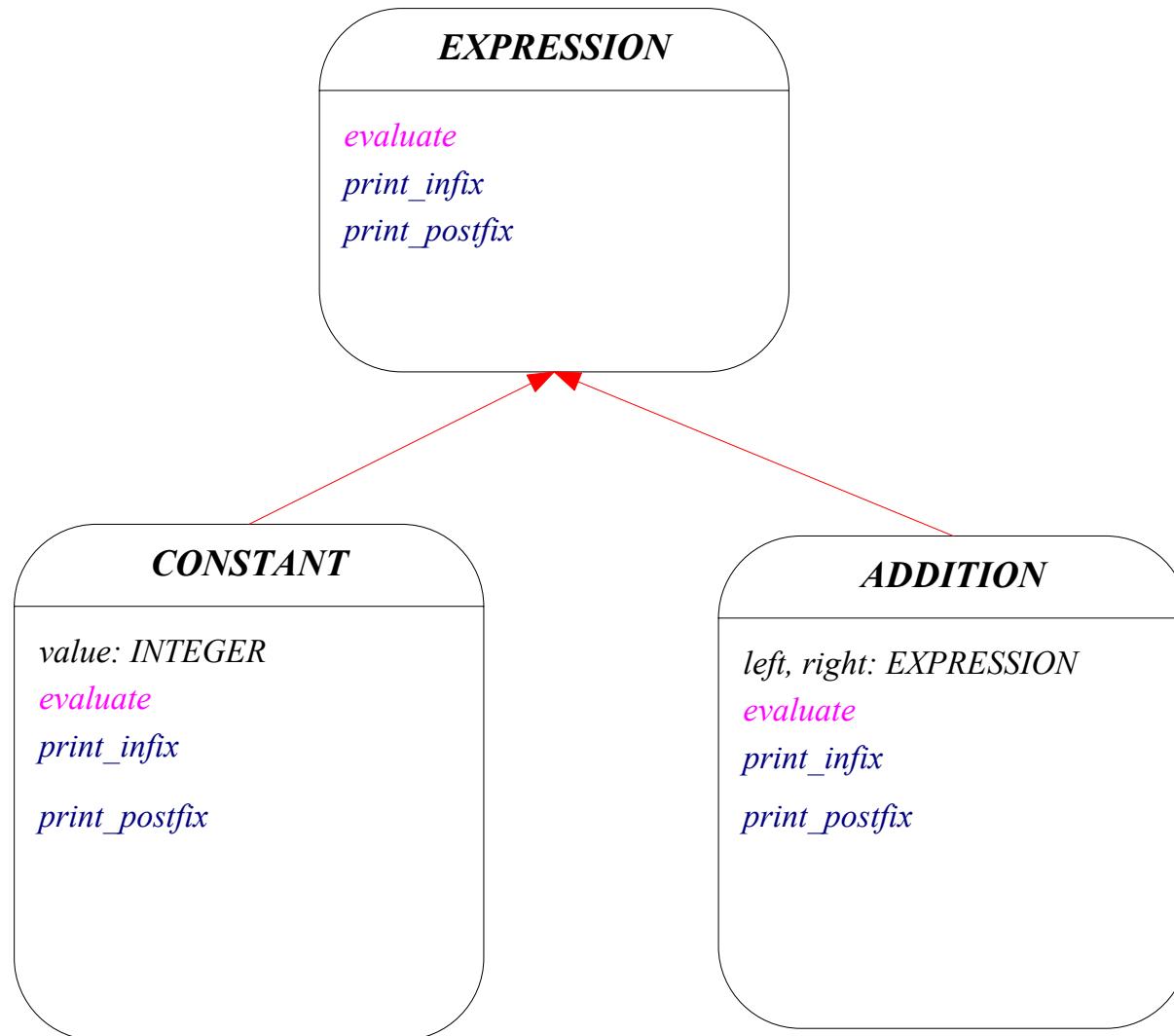
# A parsing problem – Integer expressions



$$\text{Evaluate} - (341 + 2) = 343$$



# More functions?



# Problems

- Distributing the various operations for varieties of semantic analysis across the nodes of the abstract syntax tree
  - leads to a system that is hard to understand
  - hard to maintain
  - hard to change
- It is confusing to have evaluation code mixed with pretty printing or type checking code.

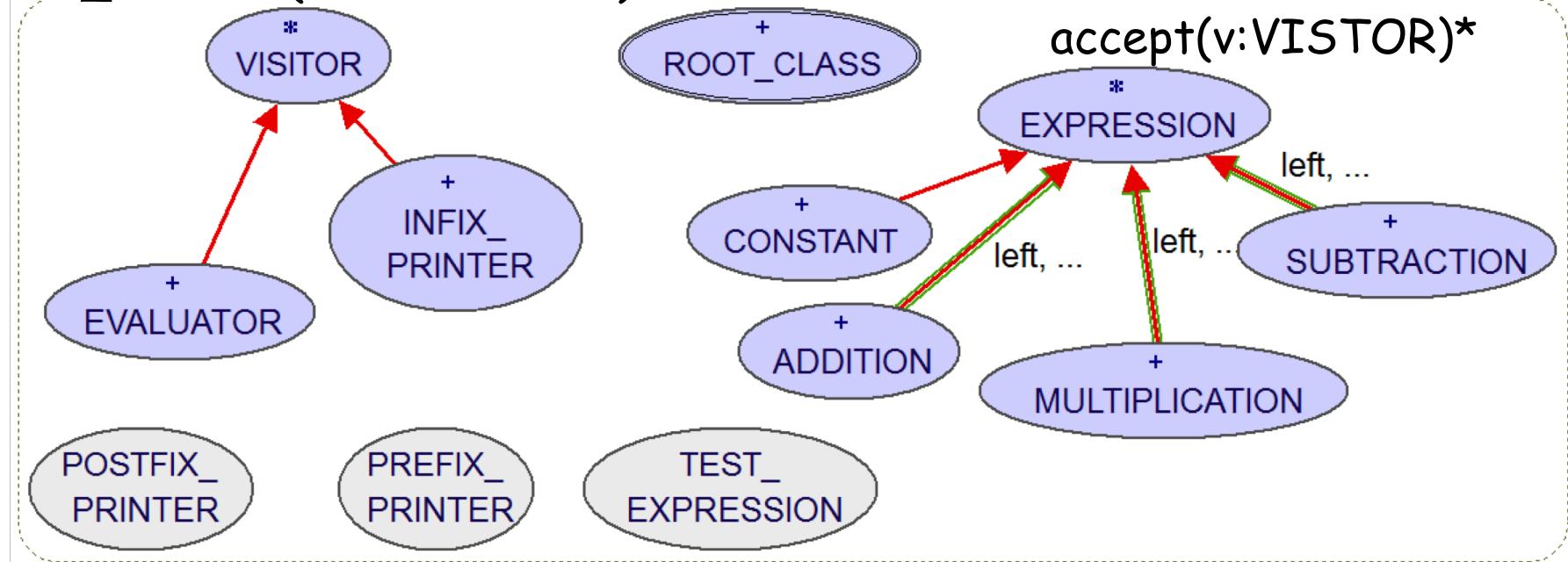
# Visitor Pattern

- An object structure contains many classes (expressions) with the need for unrelated operations (e.g. *evaluate*, *pretty print*, etc) to be performed on the classes.
- We want to avoid “polluting” the classes with these various unrelated operations.
- Define a general interface (*accept*) to apply various computations to a set of classes (e.g. expressions) without modifying them each time.
- “Call back” type of a mechanism

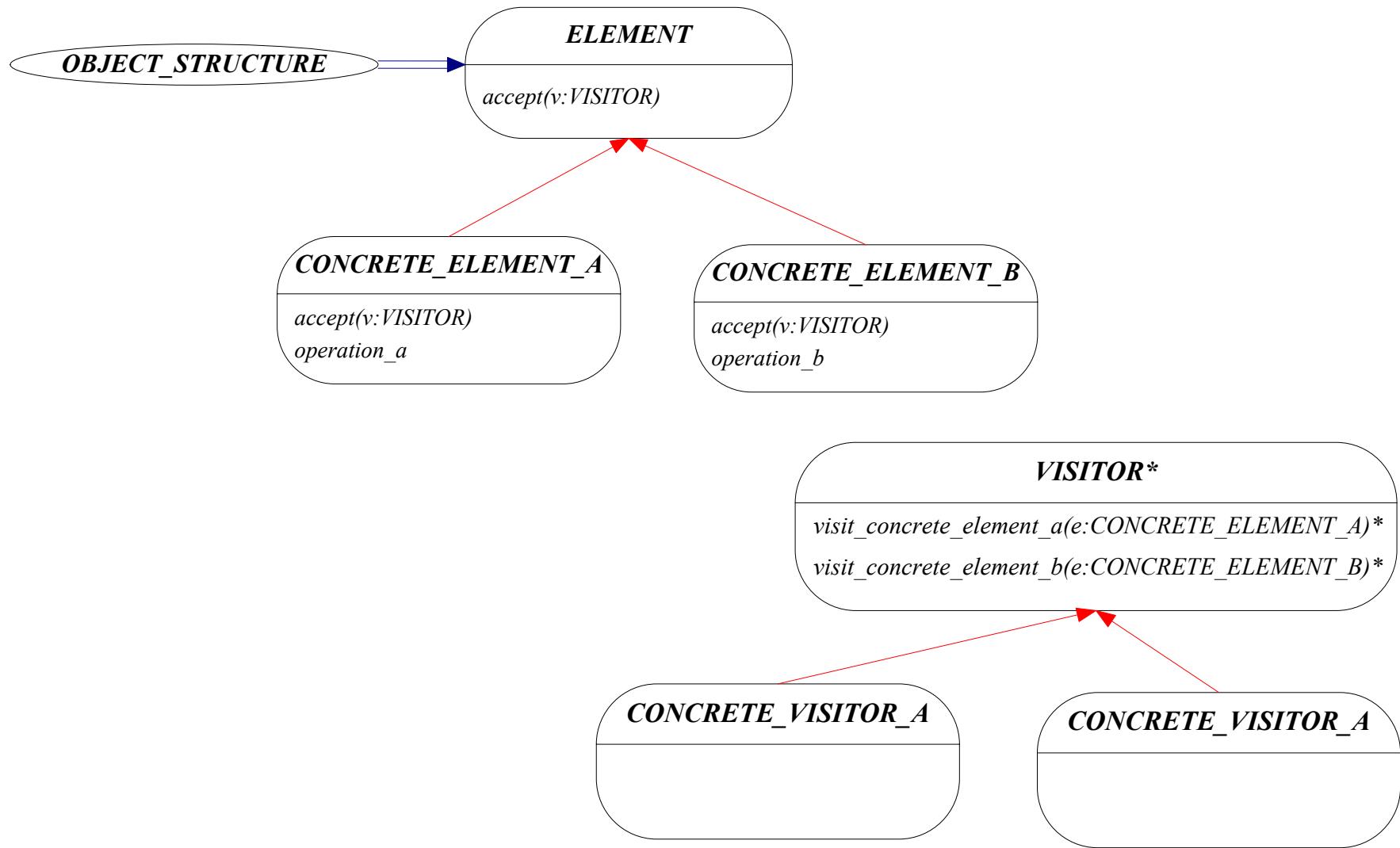
```
visit_constant(a: CONSTANT)*  
visit_addition(a: ADDITION)*
```

root\_cluster

```
accept(v:VISITOR)*
```



# Static Diagram

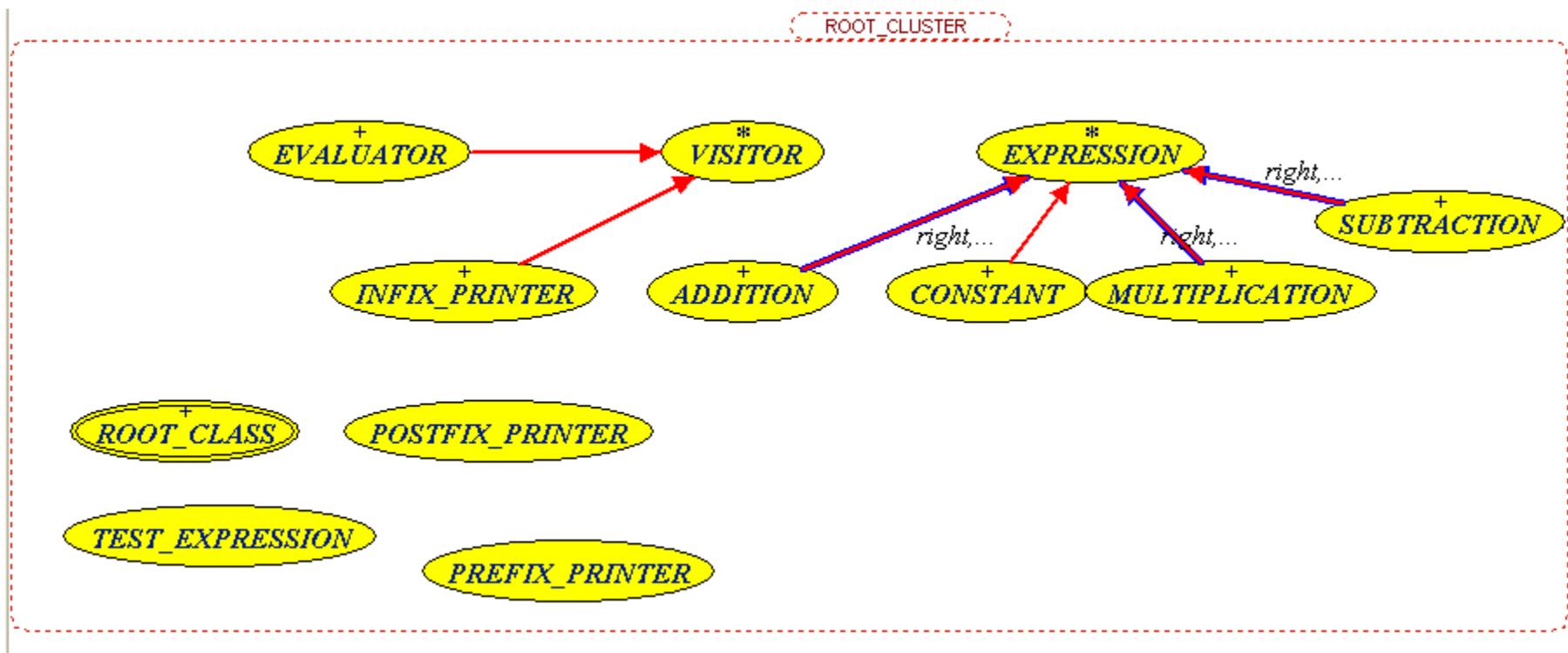


# Key to visitor pattern

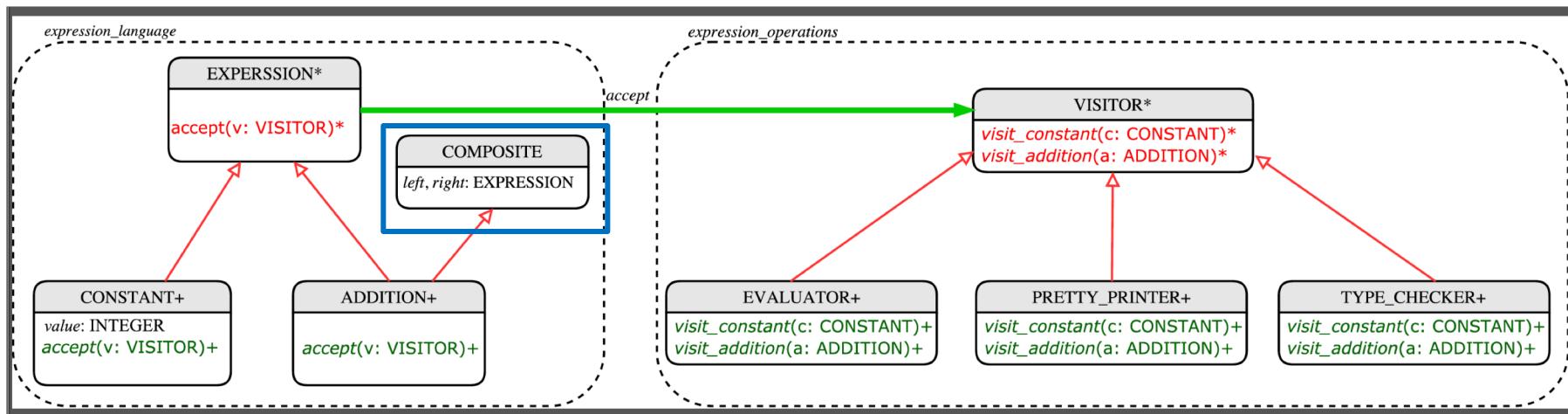
- Call back mode (double dispatch)
- The operation that gets executed depends on both the type of VISITOR and the type of ELEMENT it visits.

# Parsing expression

- See visitor-code.zip (in SVN)
- Sequence diagram on next few slides/with debugger



# Top Level View Draw.io Template



# Debugger/UML Sequence Diagram

- Follow in debugger using an ESpec unit test
- During usage, a Visitor visits each element in the structure via recursive descent
  - Top Down
  - Left to Right
- See Prof. Wang's Video  
[https://www.youtube.com/watch?v=z0bjLzs\\_2-s](https://www.youtube.com/watch?v=z0bjLzs_2-s)
- [https://www.eecs.yorku.ca/~jackie/teaching/lectures/index.html#EECS3311\\_F17](https://www.eecs.yorku.ca/~jackie/teaching/lectures/index.html#EECS3311_F17)
  - For code

# Test: $(341 + 2) = 343$

```
test_evaluator_with_plus: BOOLEAN
  local
    constant341, constant2: CONSTANT
    plus: ADDITION
    evaluator: EVALUATOR
  do
    comment ("test_evaluator_with_plus")
    create constant341.make (341)
    create constant2.make (2)
    create plus.make (constant341, constant2)
    create evaluator
    plus.accept (evaluator)
    Result := evaluator.value = 343
  end
```

# Test: $(341 + 2) = 343$

```
test_evaluator_with_plus: BOOLEAN
  local
    constant341, constant2: CONSTANT
    plus: ADDITION
    evaluator: EVALUATOR
  do
    comment ("test_evaluator_with_plus")
    create constant341.make (341)
    create constant2.make (2)
    create plus.make (constant341, constant2)
    create evaluator
    plus.accept (evaluator)
    Result := evaluator.value = 343
  end
```

Test:  $(341 + 2) = 343$

```
test_evaluator_with_plus: BOOLEAN
local
    constant341, constant2: CONSTANT
    plus: ADDITION
    evaluator: EVALUATOR
do
    comment ("test_evaluator_with_plus")
    create constant341.make (341)
    create constant2.make (2)
    create plus.make (constant341, constant2)
    create evaluator
    plus.accept (evaluator)
    Result := evaluator.value = 343
end
```

Objects			
Name	Value	Type	Address
+ Current object	<0x102C27FD0>	ROOT_CLASS	0x102C27FD0
- Locals			
constant2	<0x102C27FF8>	CONSTANT	0x102C27FF8
! value	2	INTEGER_32	
Once routine			
constant341	<0x102C28000>	CONSTANT	0x102C28000
! value	341	INTEGER_32	
Once routine			
evaluator	<0x102C27FE8>	EVALUATOR	0x102C27FE8
! value	0	INTEGER_32	
Once routine			
plus	<0x102C27FF0>	ADDITION	0x102C27FF0
left	<0x102C28000>	CONSTANT	0x102C28000
! value	341	INTEGER_32	
Once routine			
right	<0x102C27FF8>	CONSTANT	0x102C27FF8
! value	2	INTEGER_32	
Once routine			
Once routine			
Result			
! Result	False	BOOLEAN	

# Test: $(341 + 2) = 343$

test\_evaluator\_with\_plus: BOOLEAN

```

local
  constant341, constant2: CONSTANT
  plus: ADDITION
  evaluator: EVALUATOR
do
  comment ("test_evaluator_with_plus")
  create constant341.make (341)
  create constant2.make (2)
  create plus.make (constant341, constant2)
  create evaluator
  plus.accept (evaluator)
  plus.value = 343
end

```

Flat view of feature `accept' of class ADDITION

```

accept (visitor: VISITOR)
  -- Traverse with `visitor'.
  do
    visitor.visit_addition (Current)
  end

```

Objects

Name	Value	Type	Address
Current object	<0x102C27FF0>	ADDITION	0x102C27FF0
left	<0x102C28000>	CONSTANT	0x102C28000
value	341	INTEGER_32	
Once routines			
right	<0x102C27FF8>	CONSTANT	0x102C27FF8
value	2	INTEGER_32	
Once routines			
Arguments			
visitor	<0x102C27FE8>	EVALUATOR	0x102C27FE8
value	0	INTEGER_32	
Once routines			

# Test: $(341 + 2) = 343$

test\_evaluator\_with\_plus: BOOLEAN

```

local
    constant341, constant2: CONSTANT
    plus: ADDITION
    evaluator: EVALUATOR
do
    comment ("test_evaluator_with_plus")
    create constant341.make (341)
    create constant2.make (2)
    create plus.make (constant341, constant2)
    create evaluator
    plus.accept (evaluator)
plus.value = 343

```

Flat view of feature `visit\_addition' of class EVALUATOR

```

visit_addition (a: ADDITION)
    -- Process an addition expression.
local
    left_evaluator, right_evaluator: EVALUATOR
do
    create left_evaluator
    create right_evaluator
    a.left.accept (left_evaluator)
    a.right.accept (right_evaluator)
    value := left_evaluator.value + right_evaluator.value
end

```

Objects (EVALUATOR)

Name	Value	Type	Address
Current object	<0x102C27FE8>	EVALUATOR	0x102C27FE8
value	0	INTEGER_32	
Once routines			
Arguments			
a	<0x102C27FF0>	ADDITION	0x102C27FF0
left	<0x102C28000>	CONSTANT	0x102C28000
value	341	INTEGER_32	
right	<0x102C27FF8>	CONSTANT	0x102C27FF8
value	2	INTEGER_32	

# Test: $(341 + 2) = 343$

Flat view of feature `visit\_addition' of class EVALUATOR

```

test_evaluator_with_plus: BOOLEAN
local
  constant341, constant2: CONSTANT
  plus: ADDITION
  evaluator: EVALUATOR
do
  th_plus" )
  1)

41, constant2)
343

```

```

visit_addition (a: ADDITION)
  -- Process an addition expression.
local
  left_evaluator, right_evaluator: EVALUATOR
do
  create left_evaluator
  create right_evaluator
  a.left.accept (left_evaluator)
  a.right.accept (right_evaluator)
  value := left_evaluator.value + right_evaluator.value
end

```

Feature Class

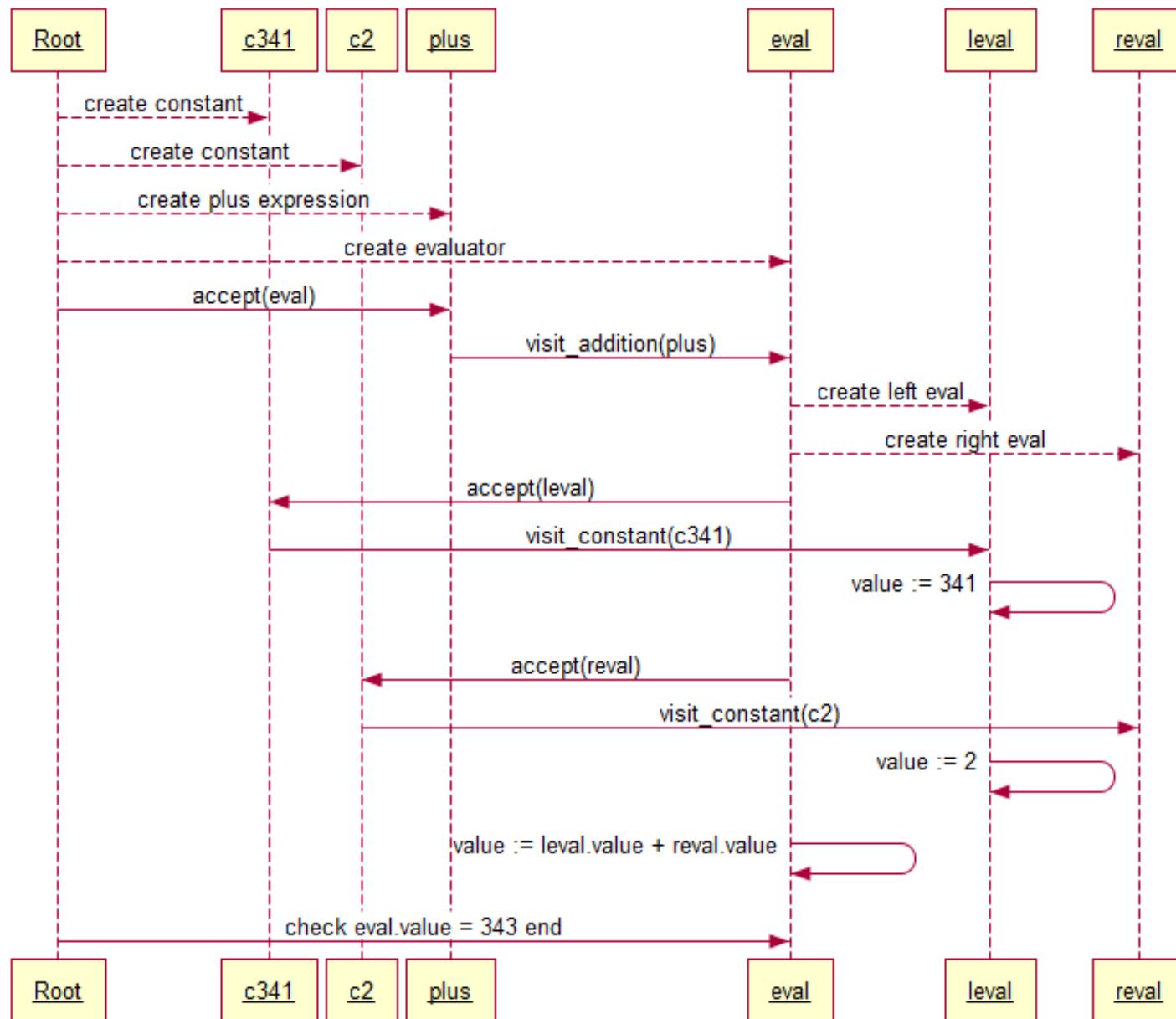
Objects			
Name	Value	Type	Address
Current object	<0x102C27FE8>	EVALUATOR	0x102C27FE8
value	343	INTEGER_32	
Once routines			
Arguments			
a	<0x102C27FF0>	ADDITION	0x102C27FF0
left	<0x102C28000>	CONSTANT	0x102C28000
value	341	INTEGER_32	
Once ro...			
right	<0x102C27FF8>	CONSTANT	0x102C27FF8
value	2	INTEGER_32	
Once ro...			
Once routi...			
Locals			
left_evaluator	<0x102C28030>	EVALUATOR	0x102C28030
value	341	INTEGER_32	
Once routi...			
right_evaluator	<0x102C28028>	EVALUATOR	0x102C28028
value	2	INTEGER_32	
Once routi...			

# Test: $(341 + 2) = 343$

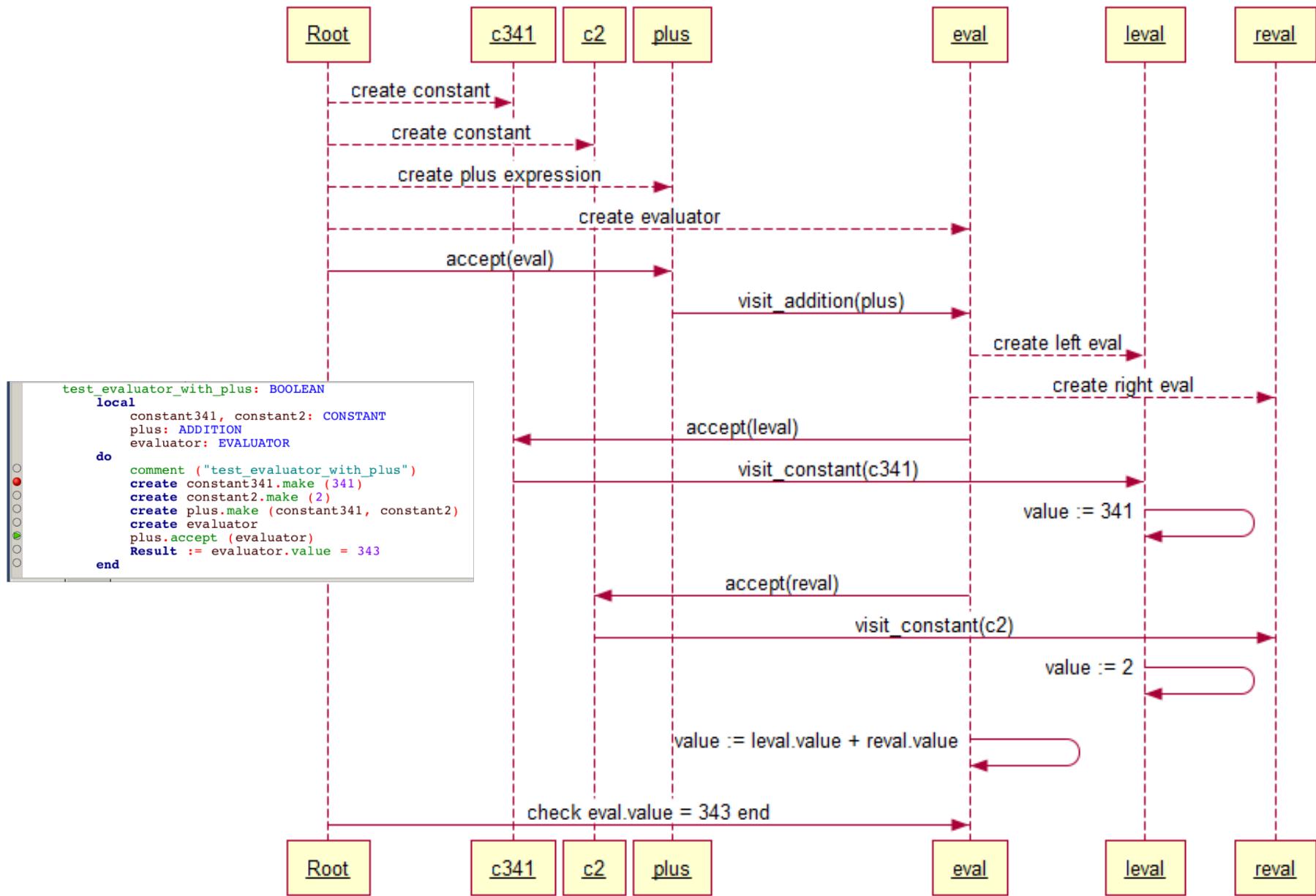
Flat view of feature `visit_addition` of class EVALUATOR			
{EVALUATOR}.visit_addition			
visit_addition (a: ADDITION) -- Process an addition expression.			
local left_evaluator, right_evaluator: EVALUATOR			
do create left_evaluator create right_evaluator a.left.accept (left_evaluator) a.right.accept (right_evaluator) value := left_evaluator.value + right_evaluator.value			
end			
Feature  Class			
Objects			
Name	Value	Type	Address
Current object	<0x102C27FE8>	EVALUATOR	0x102C27FE8
value	343	INTEGER_32	
Once routines			
Arguments			
a	<0x102C27FF0>	ADDITION	0x102C27FF0
left	<0x102C28000>	CONSTANT	0x102C28000
value	341	INTEGER_32	
Once routine			
right	<0x102C27FF8>	CONSTANT	0x102C27FF8
value	2	INTEGER_32	
Once routine			
Locals			
left_evaluator	<0x102C28030>	EVALUATOR	0x102C28030
value	341	INTEGER_32	
Once routine			
right_evaluator	<0x102C28028>	EVALUATOR	0x102C28028
value	2	INTEGER_32	
Once routine			

```
Flat view of feature `test_evaluator_with_plus` of class
test_evaluator_with_plus
local
constant341, d
plus: ADDITION
evaluator: EVALUATOR
do
comment ("test_evaluator_with_plus")
create constant341.make (341)
create constant2.make (2)
create plus.make (constant341, constant2)
create evaluator
plus.accept (evaluator)
Result := evaluator.value = 343
end
```

## Visitor Sequence Diagram: $341 + 2 = 343$



## Visitor Sequence Diagram: $341 + 2 = 343$



# Problems with visitor pattern

- Adding a new CONCRETE\_ELEMENT\_C class is hard
  - Every concrete visitor must be provided with a *visit\_concrete\_element\_c* operation
- Depends on what extent ELEMENT structure will change.
  - Use visitor when the algorithms applied to the object structure is likely to change
  - Do not use when the classes of elements are likely to change.

# Visitors Exercise

- Consider a grammar for expressions like  
 $(1.6*2 = 3.2) \wedge (2.8*4.6-9.7 <= 3.1)$
- Arithmetic expressions that may be embedded in Boolean expressions
- Assume that the input syntax is correct, but now you must type check before evaluation. For example,  $(2>3) + (1<3)$  may be syntactically correct according to a grammar, but is not type correct. So the type is either Arithmetic, Boolean or Unknown.
- Design two visitors, one to type check, and one to evaluate. The precondition for evaluation is that it should be type correct (either Arithmetic or Boolean).