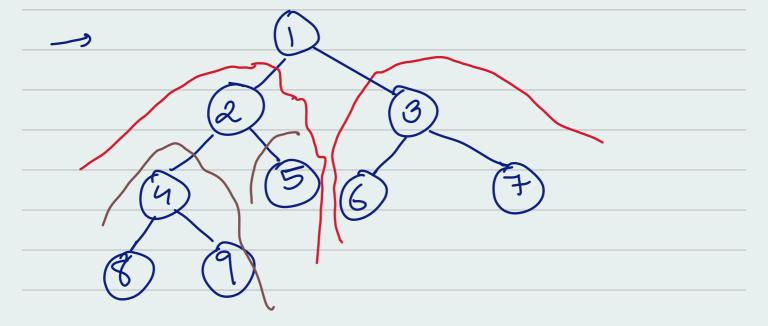
Advanced Inel Questions:

8) Construct a binavy tree from Preorder and InOrder Transmal



pre Orden:

(1), (2), 4,8,9,5,3,6,F

8, 4, 9, 2, 5, 1), 6, 3, 7

Suring these two exects the above binary there.

> using (Hash map) we can store the Constant O(1) Home Pone: 1, 2, 4, 8, 9, 5, 3, 6, 7

In '8, 4, 9, 2, 5, 1, 6, 3, 7

In '8, 4, 9, 2, 5, 1, 6, 3, 7

In similarly for 6 putit

In R. H 5014. node left = Recursion

node night = Recursion

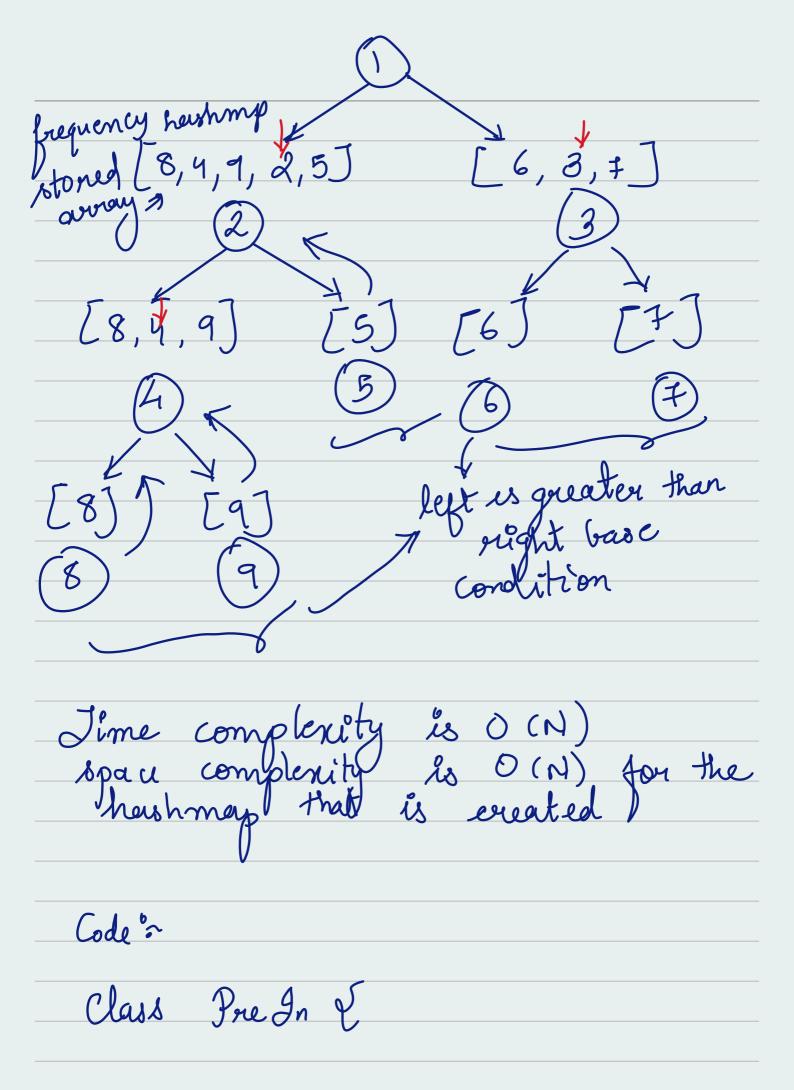
start (left will not dange)

node left = l, (i-1)

Right will change

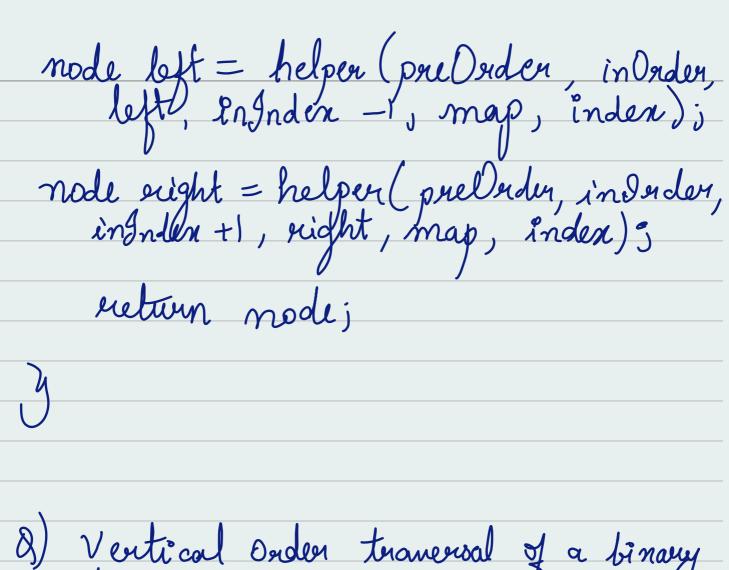
node sight = (i+1), r to index -1.

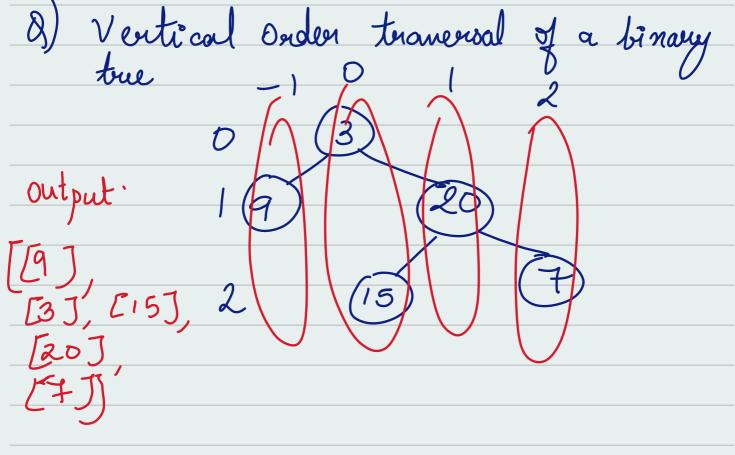
Left will change to i+1 Jelou is a recurrire diagram.



Dublic dree Node build Juie
jublic due Node build due (int [] preOrder, int [] in Order) E
Hashmap < Integer, Integer > map = new Hashmap < > ();
Jon (int i=0; iz inOrden length -); i+t)
map-put (înOnden [i], r),
snt index = 403 ;
retrun helper (preOrder, in Order, O pre Order-length -1, map, index),
J. J
public Treexode helper (int [] pre Orden,
Hashman < Integer, Integer > map & int
public TreeNode helper (int[] preOrder, ent[] inOrder, int left, int eight, Hashmap < Integer, Integer > map, ent [] index)
Ef (left > right) E
El (left > sight) E return null 3

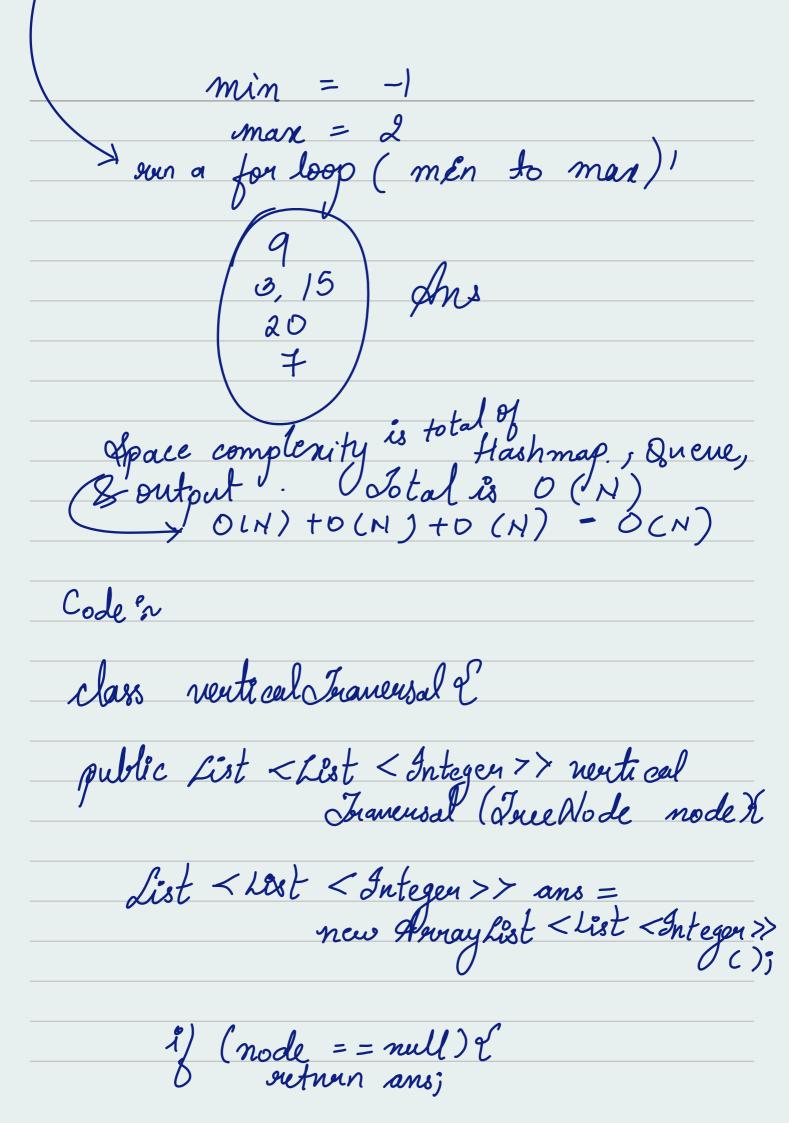
I This check whether the current subtree is valid or not of the left is greater
is valid on not et left is greater
than night it means indices have crossed each other and there are no elements so it will return null:
crossed each other and there are no
elements soit will return null.
8 → Here left is 0 and index is index-1 [leaf node] boundary crossed 80 return mull.
index-1 bleaf node J boundary
crossed so return mill'
ent euwent = preOrder index[[0]];
Index [0]++;
Jue Node mode = new Jue Node
Cowvent);
if (left = = right) E return mode 3
Jettivin mode z
int inIndex = map get (current);
1/This code in Index obtained from the
map to connectly divide the in Orde
int inIndex = map get (current); 1/This code in Index obtained from the map to connectly divide the inOrder overay into left and right subtrees for the recursive calls
Jor V the recursive call





We mill use Breadth First Search by uring greve (3-0) (9-1) (2011) null null. Desplay 9 first and then 1-1 s/+1 and then rust of et neutical level by level Store it en a hashmap The key will be col, nalue is nodes

column		mn	nodes	
	0		[3,15]	
	-1		[9]	
	J		[20]	
	2		[7]	
1		1 11		
(904	t t	iese keys	by maintainin
W	u'n	and	d max	



int col = 0; Queue < Map. Entry Tree Node Integer >>

queue = new Away Beque <>>); Map < Integer, Annay List < Integer >>
map = new Hash Map (); queue. offen (new Abstract Map. Simple Entry < \((node, col)\); notile (!queue : is Empty ()) d Mag-Entry < Jule Wode, Integer > removed = queue poll (); node = sumoned getkey (); col = sumoned get value (); 3) (mode = = mull) {

i) (! map · eentalns Key Ccol)) {

map· put (col, new Array fist

Snteger > ();

map. get (col). add (node nal); min = Math·men (men, col); max = Math·max (max, col); queul effen (new Abstract Map. Simple Entry <7 (node left, col-10); queue offen (new Abstract Map semple Entery <> (node oright, col+1)); fou (Int i = min, ix= max; i+t) {

ans. add (map. get (i));

g Jetwin ans;

Word Ladden	- Hard!
words ~	hit—> cog
dot, hot, dog	, løt, sog, log,
Using set cause aus	e it is easy to
	<u>v</u>
2 bet bot dot	lot dog løg
Here every 2	voud is compared
with the	re envoue set
	imes the word
com pouls e	m x the character
companison is 0 (N	2 × m) Time Complexity
	is space complexity.

