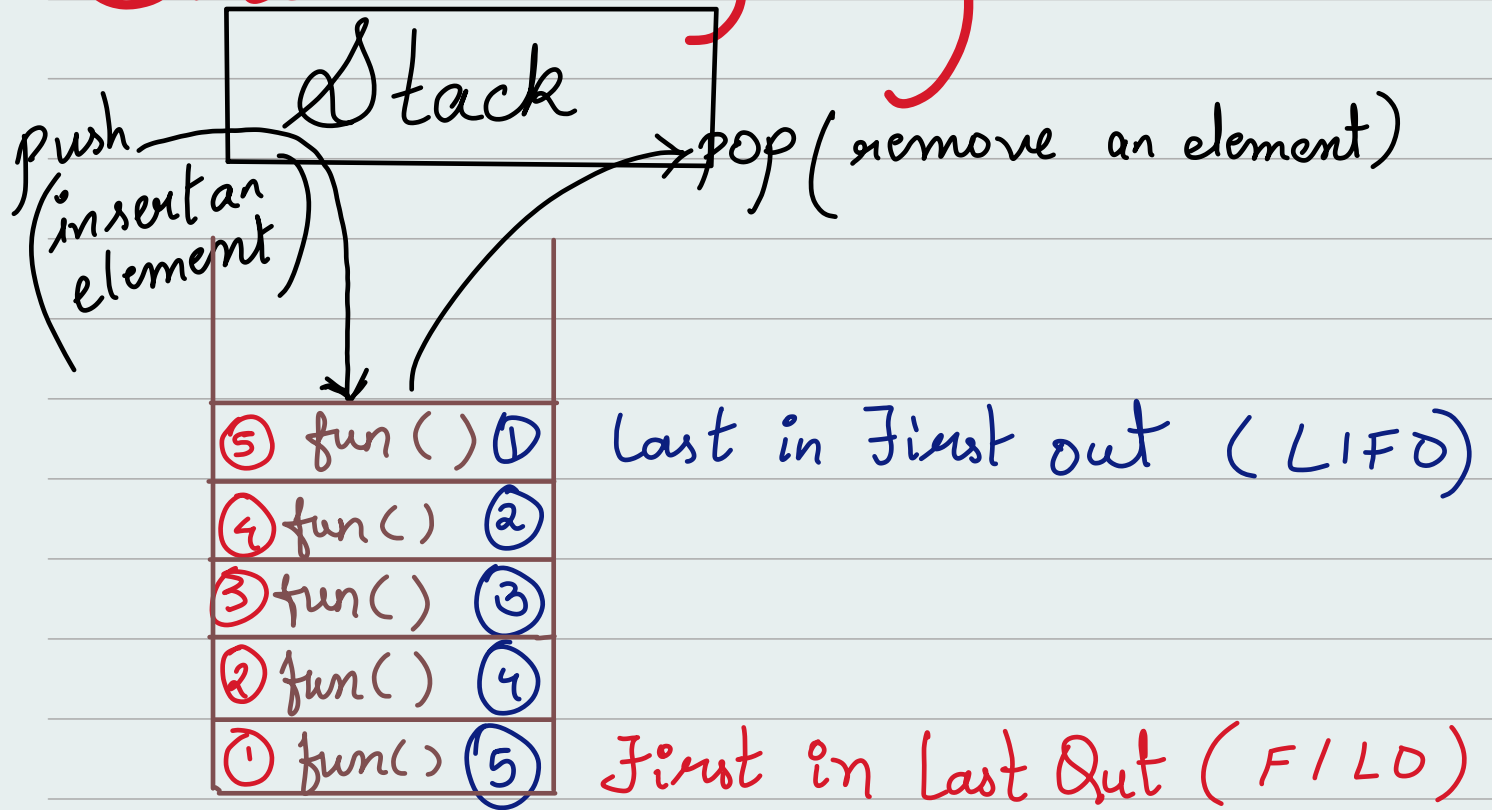
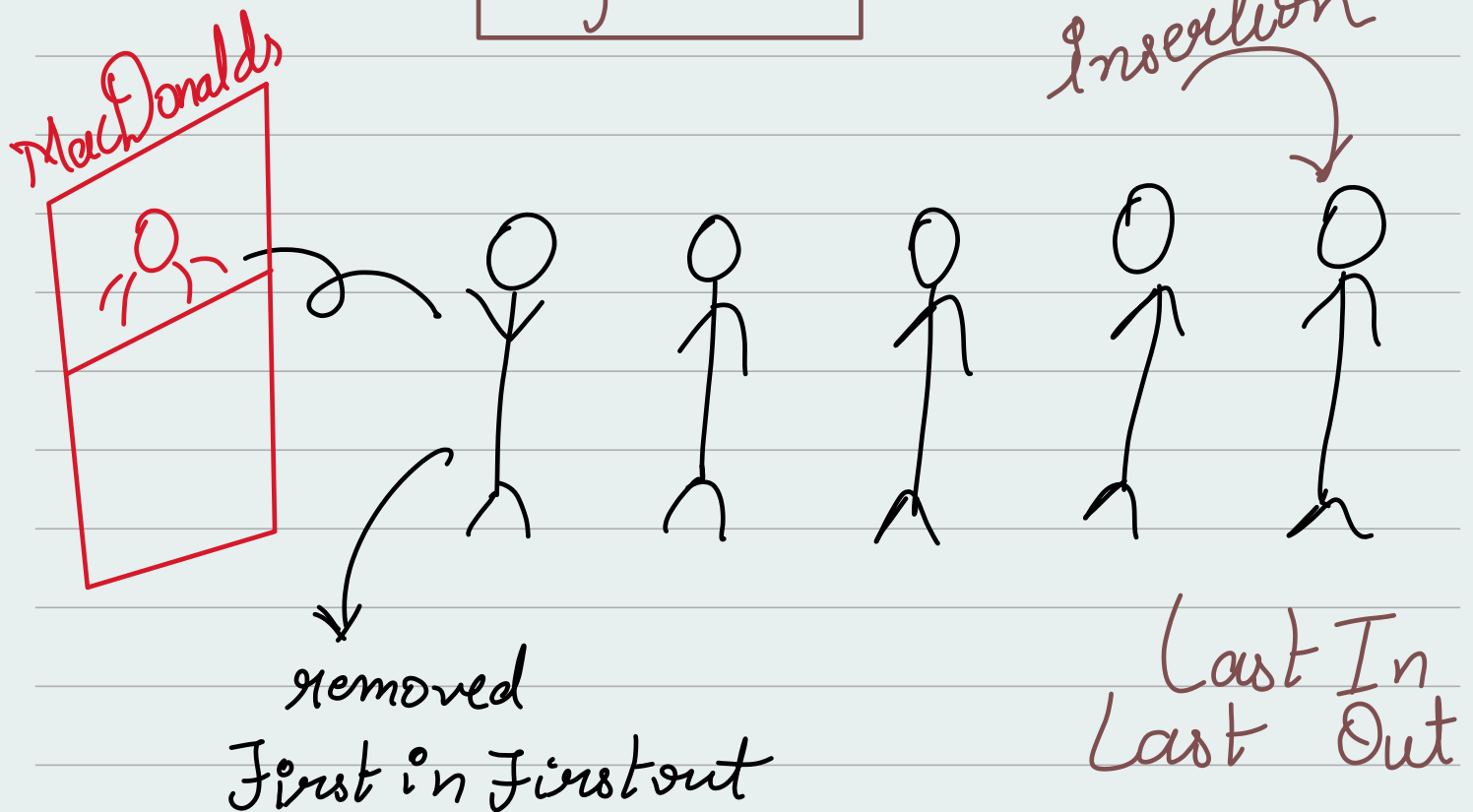


# Stack & Queues ~



## Queues

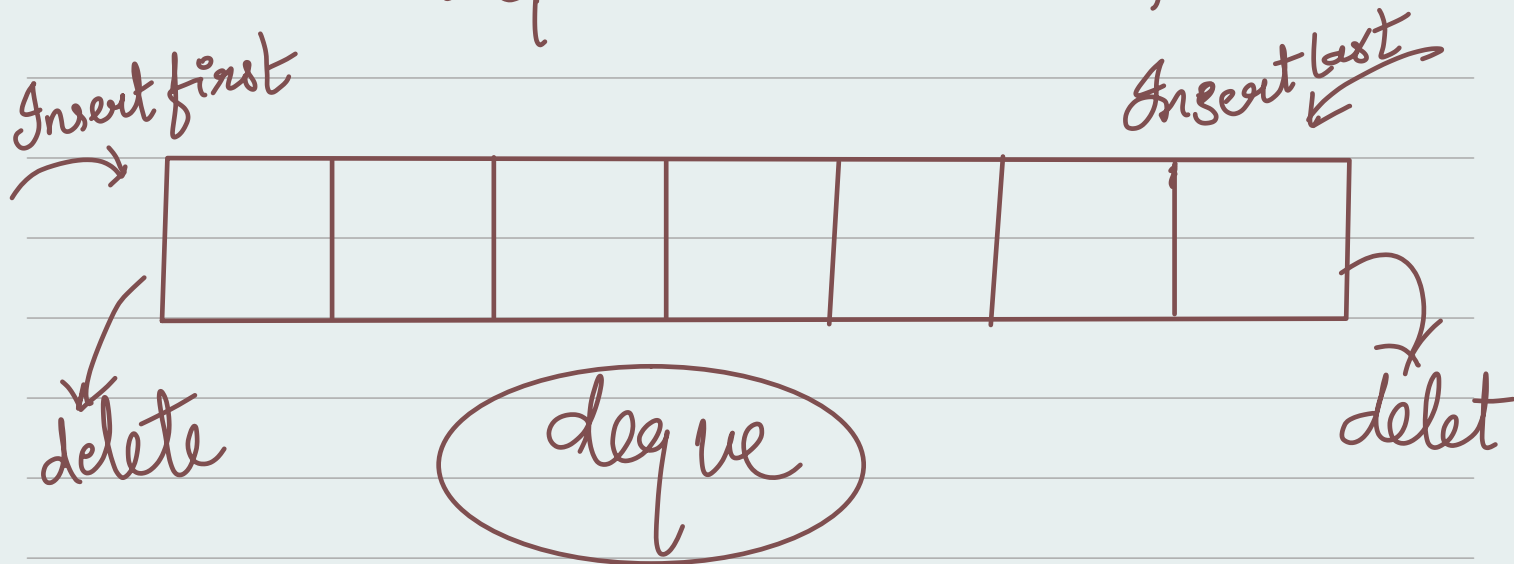


```
Queue<Integer> queue = new LinkedList<>();
```

Deque : Counter on both the sides we can insert & remove elements from both the sides.

```
Deque<Integer> deque = new ArrayDeque<>();
```

```
deque.add(10);  
deque.addLast(78);  
deque.removeFirst();
```



\* Custom Stack.

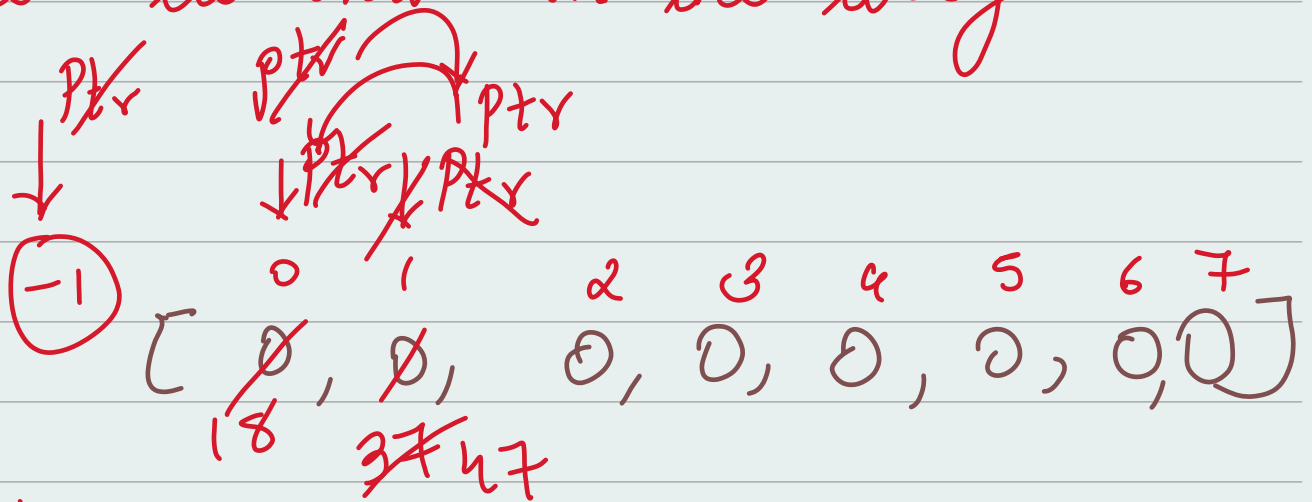
ptr  $\rightarrow$   $\textcircled{-1}$

0	1	2	3	4	5	6	7
0	0	0	0	0	0	0	0

ptr pointed to -1 index value.

So if we add or insert a item in this array we have to maintain

a pointer that would add or remove the elements in the array



$ptr$   
 0 add an element <sup>(18)</sup> increment  $-1 + 1 = 0$   
 1 add an element 37  $0 + 1 = 1$   
 0 delete an element 37  $1 - 1 = 0$   
 1 add an element 47  $0 + 1 = 1$

Using this pointer to go left and right to add or delete

When array is full  $\{ ptr == data\ length - 1 \}$

Custom Queue:

data =  $[ 34, 0, 0, 0, 0, 0 ]$   
 indices 0 1 2 3 4 5  
 45

end  
~~0~~  
 1

when data reach index 5 it's full.

2

data[i] data = [3, 9, 14, 18, 77]

0 1 2 3 4

remove

shift

data[i-1] [9, 14, 18, 77, 0]

0 1 2 3 4

→ Adding complexity is  $O(1)$  because we are adding an element

→ But removing an element changes the complexity to  $O(N)$  cause it's shifting left to  $O(N-1)$  which is  $O(N)$

```
public int remove() throws Exception {
```

```
    if (isEmpty()) {
```

```
        throw new Exception("Queue is Empty");
```

```
    }
```

```
    int removed = data[0]
```

```
    // shift the elements to left,
```

```

for (int i = 1; i < end; i++) {
    data[i-1] = data[i]
}
end--;
return removed;
}

```

```

public int front () throws Exception {
    if (isEmpty()) {

```

```

        throw new Exception("Queue is Empty");
    }

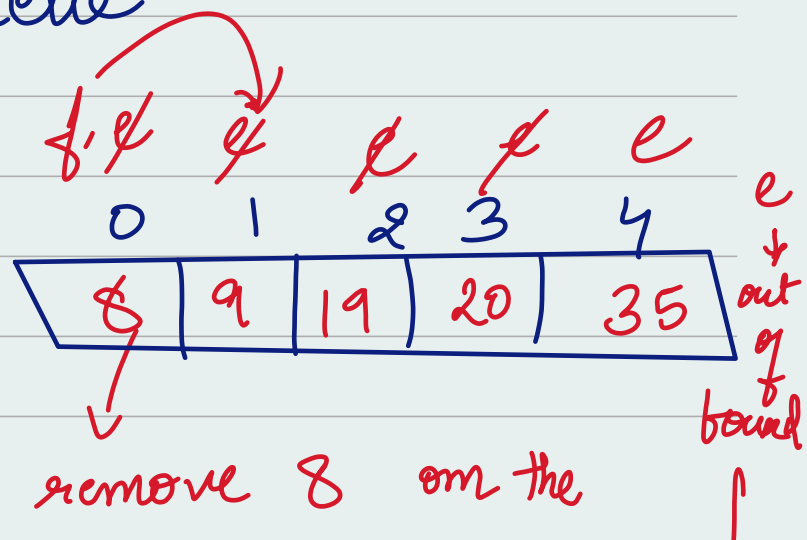
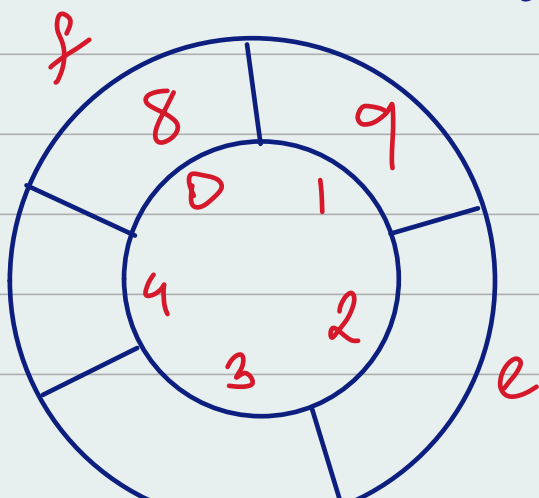
```

```

    return data[0];
}

```

## \* Circular Queue



index number zero.  
front item is getting removed.

→ Use remainder theorem

$e \% \text{Size of the array.}$

$$5 \% 5 = 0$$

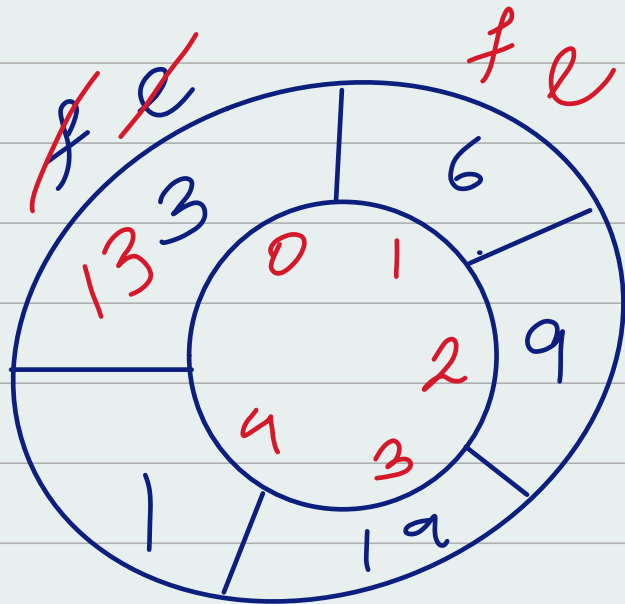
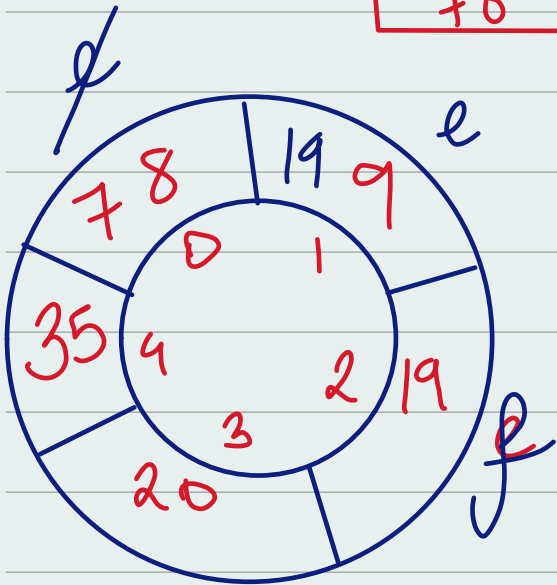
end is at 0<sup>th</sup> index  
again.

$$6 \% 5 = 1$$

Now,

~~e~~ → e f

78	199	19	20	35
----	-----	----	----	----



(Remove an item)

(Insert an item)

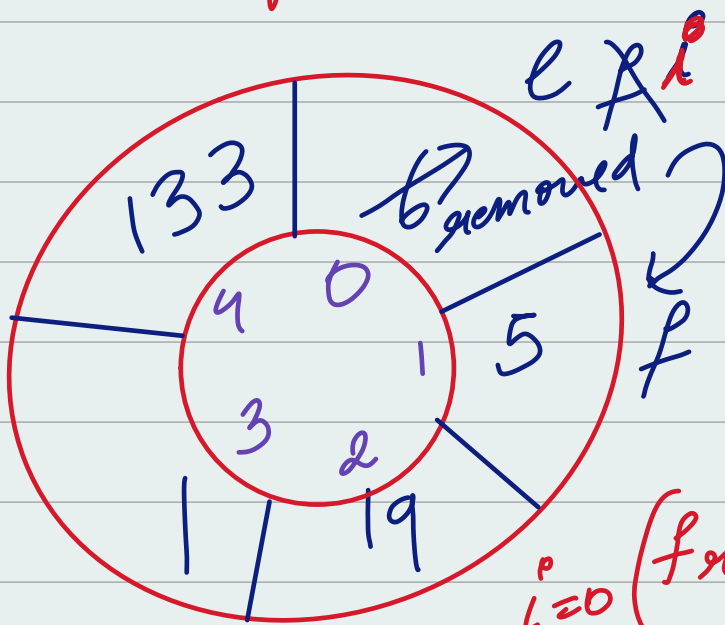
# Remove 3 insert 133.

taking i pointer till its not equal to end using do while loop

from 1 to 5 index.

when elements are copied dynamically

After for loop cond  
 $\text{temp}[i] = \text{data}[(\text{front} + i) \% \text{data.length}];$



front = 2  
i = 0

$i = 0 \quad (\text{front} + i) \% \text{data.length}$

2nd index value	$i = 0$	$2 \% 5 = 2$	It will go data-length times
3rd index value	$i = 1$	$3 \% 5 = 3$	
4th index value	$i = 2$	$4 \% 5 = 4$	
	$i = 3$	$5 \% 5 = 0$	
	$i = 4$	$6 \% 5 = 1$	

$f \downarrow$	$\downarrow$	$\downarrow$			$l \downarrow$				
0	1	2	3	4	5	6	7	8	9
5	19	1	133	6					

$i = 5$   $i \neq \text{data length}$   
 terminate loop

so, front = 0  
 & end = data length



