

Range Query Interview Questions

Q)

0	1	2	3	4	5	6	7	8
1	3	5	2	7	6	1	4	8

Sum b/w index (2, 6)

Ans : Run a for loop from index 2 to 6.

Worst case : Time Complexity = $O(N)$

→ Do better.

Algorithm : Square Root Decomposition.

* Used to optimize queries

* Reduces time complexities to $O(\sqrt{N})$

Steps :-

① Divide the array into blocks of size \sqrt{N} .

② Compute ans for every block

③ Given a question or query for left to right, combine the blocks that contain this

0	1	2	3	4	5	6	7	8
1	3	5	2	7	6	1	4	8

$$\sqrt{N} = \sqrt{9} = 3$$



1, 3, 5	2, 7, 6	1, 4, 8
---------	---------	---------

0	1	2	3	4	5	6	7	8	9
1	3	5	2	7	6	3	1	4	8



0	1	2	3
1, 3, 5	2, 7, 6	3, 1, 4	8, 0, 0



0	1	2	3
9	15	8	8

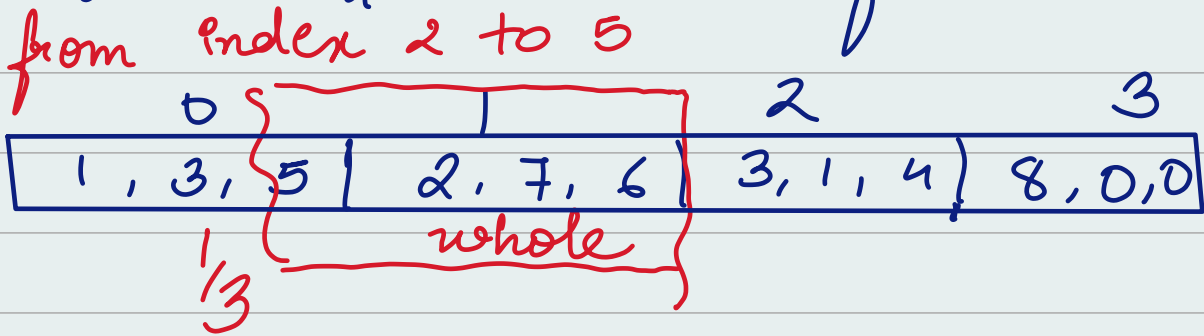
(Preprocessing)

0	1	2	3	4	5	6	7	8	9
1	3	5	2	7	6	3	1	4	8

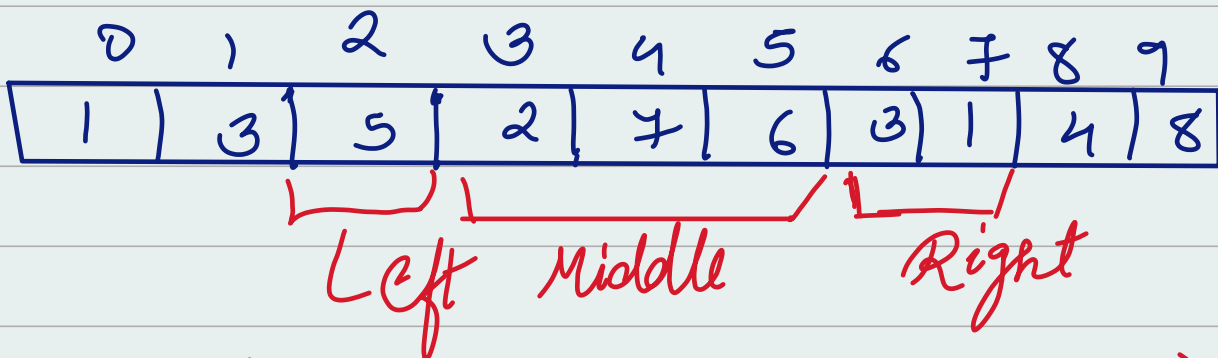
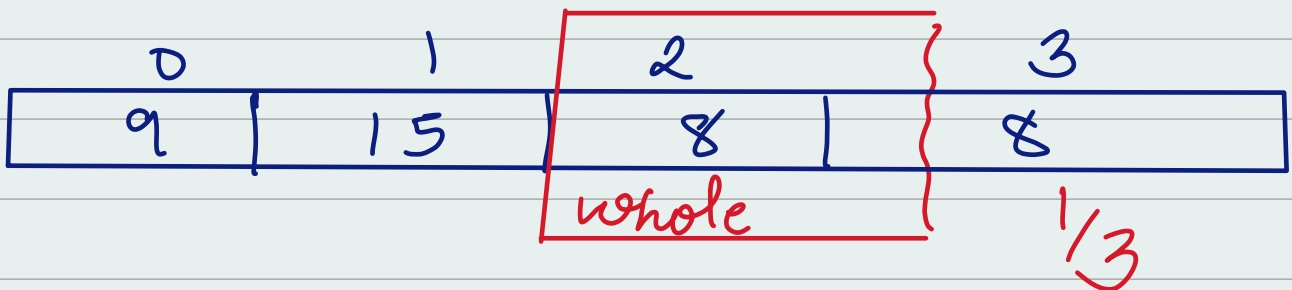
i-case

i) First case: taking whole blocks.
like indices from 3 till 8

ii) Second Case: taking $\frac{1}{3}$ from first block and the whole from other block



iii) Third case: Similarly taking last block as $\frac{1}{3}$ and the other block as whole From index 6 till 9.

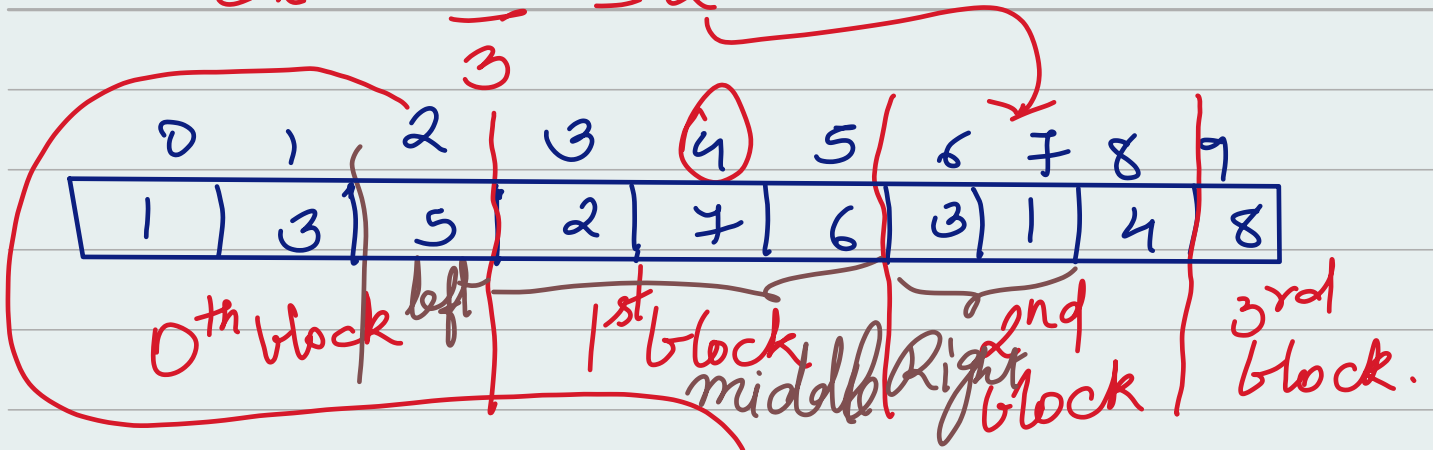


$$(2, 7) = (2, 2) + (3, 5) + (6, 7)$$

Here $\boxed{9} = \frac{9}{\sqrt{N}} = \frac{9}{3} = 3 \rightarrow \text{index} = 1$

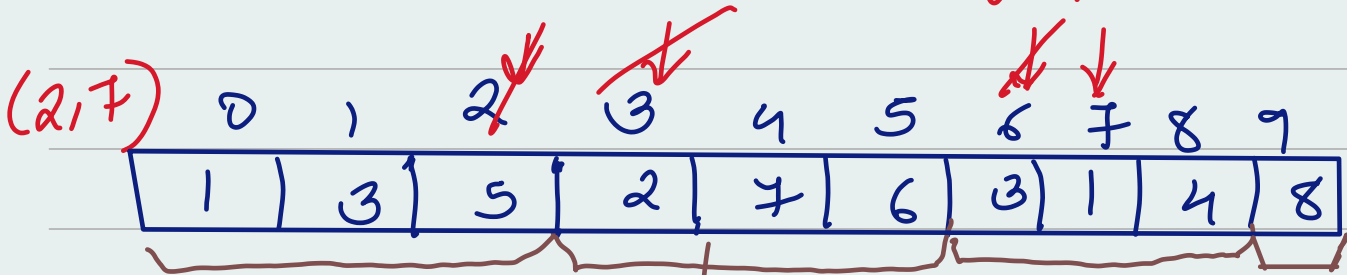
which block does index no 7 lies in? Ans block 2.

Ex: $\frac{7}{3} = 2$



$i = \frac{2}{3} = 0$

(jump left 2 times)



Here chunks are size of \sqrt{N}



l	n	sum
2	7	5
3		5+15=20
6		20+3=23
7		23+1=24

left ✓

$l \% \sqrt{N} = 2 \% 3 = 2$

l++

$l \% \sqrt{N} = 3 \% 3 = 0$

left become 0 come out of the loop left part is done, now for middle part.

while (middle ?) ✓

$$\frac{1}{8} (3) + \sqrt{N} \leq 9$$
$$\frac{1}{8} (3) + (3) \leq 7 (\text{table})$$

$$\text{Sum} = \text{sum} + \text{block} \left(\frac{1}{\sqrt{N}} \right)$$

$$= \text{sum} + \left(\frac{3}{3} \right)$$

$$= \text{sum} + \text{index}(1)$$

$$= 5 + 15 = 20$$

Here we don't have to calculate the sum of each & every block In the above array it is

already preprocessed calculated. Take the whole block sum as it is

Now jump the left index by \sqrt{N} times

Now, right part:

while ($l \leq r$)

$$\left(\begin{array}{l} 6 \leq 7 \rightarrow 3 \\ 7 \leq 7 \rightarrow 1 \end{array} \right) \rightarrow 4$$

Now left = Right
break.

Done

Suppose we want to update the index number 4 value 7 to 8

arr[4] = 8 update $i = l / \sqrt{n}$
 $i = 4 / 3 = 1$
Block[1]

then

$$\text{Block}[1] = 15 - 7 + 8$$

$$= 15 + (8 - 7)$$

$$\text{Block}[i] + (\text{value} - \text{arr}[i])$$

0	1	2	3	4	5	6	7	8	9
1	3	5	2	8	6	3	1	4	8

0	1	2	3
9	16	8	8

updated.

Code:

```
public class Main {
```

```
    public static void main(String[] args) {
```

```
        int[] arr = {1, 3, 5, 2, 7, 6, 3, 1, 4, 8};
```

```
int n = arr.length;
```

```
// build a block array.
```

```
int sqrt = (int) Math.sqrt(n);
```

```
int block_id = -1;
```

```
int blocks = new int[sqrt + 1];
```

```
for (int i = 0; i < n; i++) {
```

```
    // new block is starting
```

```
    if (i % sqrt == 0) {  
        block_id++;
```

```
}
```

```
blocks[block_id] = blocks  
[block_id] + arr[i];
```

```
}
```

```
System.out.println(query(blocks,  
arr, 2, 7, 3));
```

```
}
```

```
private static int query(int[] blocks,  
    int[] arr, int l, int r, int segt)  
    int sum = 0;  
    // left part
```

```
while (l % segt != 0 && l < r &&  
    l != 0) {
```

```
    sum = sum + arr[l]  
    l++;
```

```
}
```

```
// middle part
```

```
while (l + segt <= r) {  
    sum = sum + blocks[l / segt];  
    l += segt;
```

```
}
```

```
// right part
```

```
while (l <= r) {  
    sum = sum + arr[l];  
    l++;
```

```
}
```


return sum;

}

```
public void update (int [] blocks  
int [] arr, int i, int val, int i) {
```

```
int block_id = i / 4;
```

```
blocks [block_id] = block  
[block_id] + (val - arr[i]);
```

}

}

