Requisites - functions, & tack memory, heap memory

2) vosite a function that puints a message

- A function is calling another function

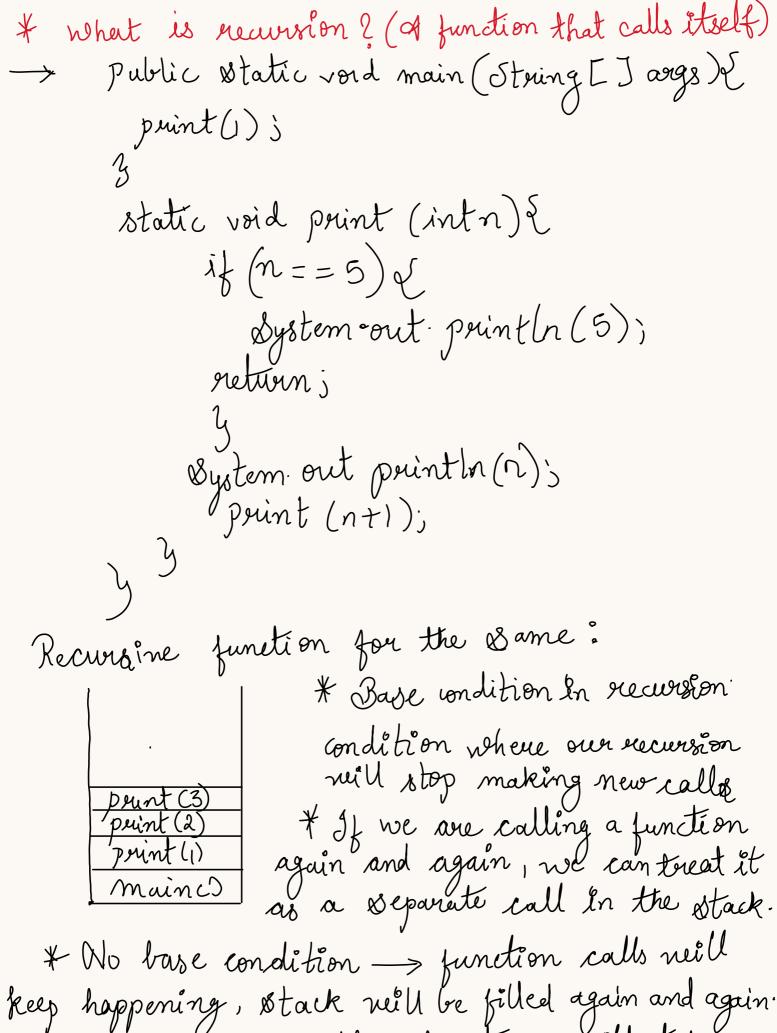
- main is called in the stack memory trow function calls works in languages.

		output
-	print 5(5)	1 2
-	print 4 (4)	3
-	print 2 (2)	4
ļ	puint 1 (1)	5
	main ()	

while the function is not finished executing, it mill remain in stack

point to be noted: when a function finishes executing it is removed from the stack and the flow of program is restorted to where that function near called

> After printing all the numbers the print function will be removed from the stack.



* No base condition -> junction calls neill keep happening, stack neill be filled again and again. We know that every call of function will take some memory.

* Memouy of computer will exceed the limit. This will give us Stack Overflow evoror

* Why Recursion!

- Junction calling itself. It helps us in solving bigger and complex problems in a simple way.

- you can convert recursion solutions into a tenation and like versa.

- Space complexity is not constant because of necursine calls.

- It helps in breaking down bigger problems into smaller problems.

* Visualizing recursions: VVIP

main () I programower

Print (1) / Recursion

Jace

Point (2)

point (3)

point (4)

Point (5)

3) Fibonacci numbers: Find nth fibonacci numbers and 3rd 4th 5th 6th th 8th 9th index 7 oth 1st 2nd 3rd 4th 5th 6th th 8th 9th gloment 30, 1, 1, 2, 3, 5, 8, 13, 21, 34, Fibonaci (N) = Fibonaci (N-1) + Fibonaci (NZ) > Example: 3 rdindex (N-1) (N-2) Similarly for (N-1) Fibo it would be (N-2) + (N-3) index elements Flbo(N-1) = Flbo(N-2) + Flbo(N-3 It will keep going on so if we by to build recursine tree for fibos.

Fibo (5) (Jibo3) 7ibo (4) +

+ Fibo(2) Fibo W_ Fibo(1)+Fibo(0) Jibo(2) + Jibo(1) Fibo (0) -> Wimilarly for] (4) 4) would be (Recursive recursions? 2 Ans

Break it down into smaller problems.

The abone formula is known recurrence relation.

* The base condition is represented

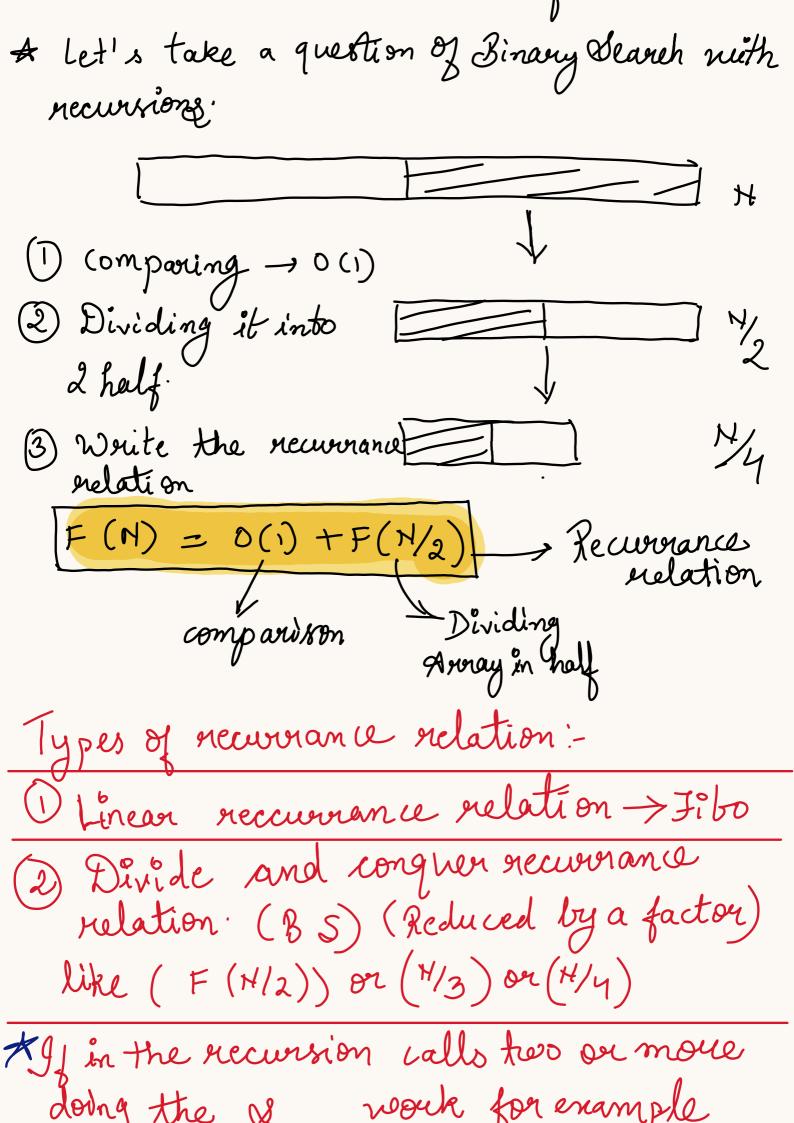
A The base condition is represented by answers we already have In the above case we know that F(0) = 0and F(1) = 1 This is base condition

P& vm (String [] wgs) { int ans = fibo (4)

System.out println (ans); static int fibo (int n) ¿ // base condition if (n < 2) f z return n; 3 return fibo (n-2) + fibo (n-1);

How to un our and and appro a publish

break down problem into Or Jodentify if you can Vmaller problèms Don'te the necuvernce relation if needed. 3 - Draw the recursive tree. (9) About the tree *See the flow of functions, how they are getting in stack * Identify and focus on left tree calls and thenright tree calls and pointer again and and paper Draw the tree again using pen A Use a debugger to see the flow of 5) -> See now the values are returned at each Step. Der where the function neill come out of. In the end it will come out of the main function. (1) Auguments (vill go in the next fut call) A Variables 2 Return Type VVIP (3) (Rody of the function How to use variables in the functi ?



F(50) - 50th fibonaci number en that
F(2), F(1) and F(0) neill be called
F(2), F(1) and F(0) neill be called again and again F(2) and this
veill create a repeatation in the function And the sans will take much longer to compute.
well create a repeatation in the function
And the sans will take much longer to
compute.
How can we solve this issue ? The ans
is Dynammic perogramming (If In the recurrison
calls two or more recursion calls are
is Dynammic perogramming (If in the recurrision calls are loing the same work don't compute it again and cagain
and cagain
Tip. Do not overthink
Make some to return the result of
Make some to return the result of a function call of the return type.
Code!
PUM (String [], args)

PSVM (String[], orgs) \(\)
int[] are = \(\ 2, 4, 6, 8, 10, 12 \ \ \);

int target = 12; System out println (binary dewich (

vu, torget, 0, over-length-1)); Static ent binaryssearch (int []avor, int target, int start, int end) & if (start > end) & return -1; int mid = stout + (end-stout)/2, if (arr [mid] == target) { return mid; if (tanget < aun [mid]) { retwin binary Dearch (aun, target, start, mid-1); return binary sewich (avr, target, mid+1, end);

AU of these weill be in stack!

main ¥3 midt < 6 6 return 6th index