

Jeffrey
Lonsford

0.1 (a) $f(n) = n - 100 \quad g(n) = n - 200$

$$\lim_{n \rightarrow \infty} \frac{n - 100}{n - 200} = 1$$

$$\lim_{n \rightarrow \infty} \frac{n - 200}{n - 100} = 1$$

$$f = O(g)$$

(b) $f(n) = n^{1/2} \quad g(n) = n^{4/3}$

$$\lim_{n \rightarrow \infty} \frac{n^{1/2}}{n^{4/3}} = \frac{\cancel{n}^{1/2}}{\cancel{n}^{4/3}} = n^{-1/2} = n^{-1/6} = 0$$

$$\lim_{n \rightarrow \infty} \frac{n^{4/3}}{n^{1/2}} = \frac{\cancel{n}^{4/3}}{\cancel{n}^{1/2}} = n^{11/6} = n^{11/6} = \infty$$

$$f = O(g)$$

(c) $f(n) = 100n + \log n \quad g(n) = n + (\log n)^2$

$$\lim_{n \rightarrow \infty} \frac{100n + \log n}{n + (\log n)^2} = 100$$

$$\lim_{n \rightarrow \infty} \frac{n + (\log n)^2}{100n + \log n} = 1/100$$

$$f = O(g)$$

$$(d) f(n) : n \log n \quad g(n) : 10n \log 10n$$

$$\lim_{n \rightarrow \infty} \frac{n \log n}{10n \log 10n} = \frac{1}{10} + \log n$$

$$\frac{1}{10} + \log n$$

$$\lim_{n \rightarrow \infty} \frac{\log n}{10 + 10 \log 10n} = \frac{\frac{1}{n}}{\frac{10}{n} + \frac{10 \log 10n}{n}} = \frac{1}{10}$$

$$\lim_{n \rightarrow \infty} \frac{10n \log 10n}{n \log n} = \frac{10n + 10 \log 10n}{n + \log n}$$

$$\lim_{n \rightarrow \infty} \frac{10 + 10 \log 10n}{\log n} = \frac{\frac{10}{n}}{\frac{\log n}{n}} = 10$$

$f = \Theta(g)$

$$(e) f(n) : \log 2n \quad g(n) : \log 3n$$

$$\lim_{n \rightarrow \infty} \frac{\log 2n}{\log 3n} = \frac{\frac{1}{n}}{\frac{\log 3n}{n}} = 1$$

$$\lim_{n \rightarrow \infty} \frac{\log 3n}{\log 2n} = \frac{\frac{\log 3n}{n}}{\frac{\log 2n}{n}} = 1$$

$f = \Theta(g)$

$$(f) f(n) = 10 \log n \quad g(n) = \log(n^2)$$

$$\lim_{n \rightarrow \infty} \frac{10 \log n}{\log(n^2)} = \frac{\frac{10}{n}}{\frac{2}{n}} = 5$$

$$\lim_{n \rightarrow \infty} \frac{\log(n^2)}{10 \log n} = \frac{\frac{2}{n}}{\frac{10}{n}} = \frac{1}{5}$$

$$f = \Theta(g)$$

$$(g) f(n) = n \quad g(n) = n \log^2 n$$

$$\lim_{n \rightarrow \infty} \frac{n}{n \log^2 n} = \frac{n^{0.01}}{\log^2 n} = \frac{0.01n^{-0.99}}{2 \log n \cdot \frac{1}{n}} = \frac{0.01n^{0.01}}{2 \log n}$$

$$= \frac{0.001n^{-0.99}}{2} = \frac{0.001n^{0.01}}{2} = 0$$

$$\lim_{n \rightarrow \infty} \frac{n \log^2 n}{n} = \frac{\log^2 n}{n^{0.01}} = \frac{2 \log n \cdot \frac{1}{n}}{0.01n^{-0.99}}$$

$$= \frac{2 \log n}{0.01n^{0.01}} = \frac{2}{0.01n^{-0.99}} = \frac{2}{0.001n^{0.01}} = 0$$

$$f = \Omega(g)$$

(h) $f(n) = n^2/\log n$ $g(n) = n(\log^2 n)$

$$\lim_{n \rightarrow \infty} \frac{n^2/\log n}{n(\log^2 n)} = \frac{n(2\log(n)-1)}{\log^2(n)} \cdot \frac{1}{\log(n)(\log(n)+2)}$$

$$\lim_{n \rightarrow \infty} \frac{n(2\log(n)-1)}{\log^3(n)(\log(n)+2)} = \infty$$

$$\lim_{n \rightarrow \infty} \frac{n^2/\log n}{n^2/\log n} = \frac{\log(n)(\log(n)+2)}{n(2\log(n)-1)} \cdot \frac{1}{\log^2(n)}$$

$$= \frac{\log^3(n)(\log(n)+2)}{n(2\log(n)-1)} = 0$$

$$\boxed{f = \Omega(g)}$$

$$(i) f(n) = n^{\alpha} \quad g(n) = (\log n)^{\beta}$$

$$\lim_{n \rightarrow \infty} \frac{n^{\alpha}}{(\log n)^{\beta}} = \frac{\alpha n^{\alpha-1}}{\beta (\log n)^{\beta-1}} = \frac{\alpha n^{\alpha-1}}{\beta (\log n)^{\beta-1}}$$

$$= \frac{\alpha \cdot \log n}{\beta \cdot \log n} = \frac{\alpha}{\beta} = \frac{n^{\alpha}}{(\log n)^{\beta}}$$

$$\frac{\alpha n^{\alpha-1}}{\beta \log n} = \frac{\alpha n^{\alpha-1}}{\beta} = \frac{\alpha n^{\alpha-1}}{n} = \alpha n^{\alpha-1} = \infty$$

$$\lim_{n \rightarrow \infty} \frac{(\log n)^{\beta}}{n^{\alpha}} = \dots \frac{1/n}{\alpha n^{\alpha-1}} = \frac{1}{\alpha n^{\alpha-1}} = 0$$

$f = \Omega(g)$

$$(j) f(n) = (\log n)^{\log n} \quad g(n) = n / \log n$$

With the Common Sense Rule, an exponent will grow faster than a polynomial. Therefore $f = \Omega(g)$ because g will be the lower bound of f .

$$(K) f(n) = \sqrt{n} \quad g(n) = (\log n)^3$$

$$\lim_{n \rightarrow \infty} \frac{n^{1/2}}{(\log n)^3} = \frac{\frac{1}{2}n^{-1/2}}{3(\log n)^2} = \frac{\frac{1}{2}n^{-1/2}}{3(\log n)^2} =$$

$$\frac{\frac{1}{2}n^{-1/2}}{6(\log n)^2} = \frac{\frac{1}{2}n^{-1/2}}{6(\log n)} = \frac{\frac{1}{2}n^{-1/2}}{6/n} =$$

$$\frac{\frac{1}{2}n^{-1/2}}{6} = \infty$$

$$\lim_{n \rightarrow \infty} \frac{n^{1/2}}{(\log n)^3} = \dots \frac{6}{8n^{1/2}} = 0$$

$$f = \Omega(g)$$

$$(D) f(n) = n^{1/2} \quad g(n) = 5^{\frac{\log_2 n}{\log_2 5}}$$

$$= n^{\frac{\log_2 5}{2}}$$

$$= n^{2.32}$$

Since $\frac{1}{2} < 2.32$, $n^{2.32}$ dominates $n^{1/2}$
 therefore by the Common Sense Rules,
 $f = O(g)$ and g is the upper bound
 of f

$$(m) f(n) = n2^n \quad g(n) = 3^n$$

By the Common Sense Rule, $2^n < 3^n$, so
 3^n dominates becoming the upper bound
 $f \in O(g)$

$$(n) f(n) = 2^n \quad g(n) = 2^{n+1}$$
$$= 2^n \cdot 2^1$$

Both function have the exponent, so
no function dominates making both
the limits 2 and 1/2, therefore $f \in O(g)$

$$(o) f(n) = n! \quad g(n) = 2^n$$

Factorial grows faster than exponential
therefore $n!$ dominates 2^n , so g will
be the lower bound of f
 $f \in \Omega(g)$

$$(p) f(n) = (\log n)^{\log n} \quad g(n) = 2^{\log n}$$
$$\in n^{\log(\log n)} \quad = 2^{\log n / \log 2}$$

Exponent with grows faster than
any polynomial, so $f \in O(g)$ from
the Common sense rules.

$$(g) f(n) = \sum_{i=1}^n c^{k+1} = n^{k+1} \quad g(n) = n^{k+1} = n^k \cdot n$$

Since both are polynomials to the k ,
 they will both be $f = O(g)$ and $g = \Omega(g)$,
 therefore $|f = \Theta(g)|$

$$0.2 \quad g(n) = 1 + c + c^2 + \dots + c^n = \frac{c^{n+1} - 1}{c - 1}$$

$$(a) \lim_{n \rightarrow \infty} \frac{c^{n+1} - 1}{c - 1} = \frac{0 - 1}{c - 1} = \frac{1}{1-c} = g(n)$$

$$\lim_{n \rightarrow \infty} c^{n+1} = 0$$

$$g = O(1) = \lim_{n \rightarrow \infty} \frac{g(n)}{1} = \frac{1}{1-c}$$

$$\lim_{n \rightarrow \infty} \frac{g(n)}{1} = 1-c$$

$$(b) \quad G(n) \text{ if } c=1 \quad g(n) = n$$

$$\lim_{n \rightarrow \infty} \frac{g(n)}{n} = \frac{n}{n} = 1 \quad g = \Theta(n) \text{ if } c=1$$

$$\lim_{n \rightarrow \infty} \frac{n}{g(n)} = \frac{n}{n} = 1$$

$$\text{d) } O(c^n) \text{ if } c > 1 \quad \frac{c^n - 1}{c - 1} = g(n)$$

$$\lim_{n \rightarrow \infty} \frac{g(n)}{c^n} = \frac{c^{n-1}}{c - 1} = \frac{c^n}{c^n} = 1$$

$$\lim_{n \rightarrow \infty} \frac{c^n}{g(n)} = \frac{c^n}{c^{n+1} - 1} = \frac{c^n}{c^n} = 1$$

$\boxed{g = O(c^n)}$

0.3

$$\text{(a) } F_n \approx 2^{0.5n} \quad n \geq 6$$

$$F_6 = F_5 + F_4 \quad F_7 = F_{n-1} + F_{n-2}, \quad F_0 = 0, \quad F_1 = 1$$

$$F_0 = 0 \quad F_1 = 0 + 1 = 1 \quad F_2 = 1 + 1 = 2$$

$$F_3 = 1 + 2 = 3 \quad F_4 = 2 + 3 = 5$$

$$F_5 = 5 + 3 = 8$$

$$2^{\lceil \log_2 8 \rceil} = 2^3 = 8$$

$$8 = 8, \quad n = 6$$

$$n = 7$$

$$F_7 = 8 + 5 = 13$$

$$2^{\lceil \log_2 13 \rceil} = 2^3.5 \approx 11.3$$

$$13 > 11.3, \quad n = 7$$

c_{n-2c}

(b) $f_n \leq 2^{cn}$

$$f_{n-1} + f_{n-2} \leq 2^{c(n-1)} + 2^{c(n-2)}$$
$$2^{cn} = 2^{cn} (2^{-c} + 2^{-c})$$

$$1 = 2^c + 2^{-c} \quad x = 2^c$$

$$1 = x + x^{-1}$$

$$0 = x^2 + x - 1$$

$$\frac{-1 \pm \sqrt{1-4(1)(-1)}}{2c} = \frac{-1 \pm \sqrt{5}}{2}$$

$$2^c = \frac{-1 \pm \sqrt{5}}{2} \quad \text{let } n=6$$

$$-c = \log_2 \left(\frac{-1 \pm \sqrt{5}}{2} \right) \quad b_6 = 8$$

$$C = -\log_2 \left(\frac{-1 \pm \sqrt{5}}{2} \right) \quad 8 \leq 2^{\frac{-c}{6}(6)} \left(\frac{-1}{2} + \frac{\sqrt{5}}{2} \right)$$

$$C \approx -\log_2 \left(\frac{-1 \pm \sqrt{5}}{2} \right) \quad \approx 18.4 (.62 + .39)$$

$$C \approx .6942 \approx .7 \quad 8 \leq 18.4, \text{ let } n=6$$

(c) So we found that $2^{4n} > f_n$, then
that means that 2^{4n} grows faster than
 f_n and it dominates it for all n

Then we when take the limit of f_n and
 2^{4n} , we get:

$$\lim_{n \rightarrow \infty} \frac{f_n}{2^{4n}} = 0$$

$$\lim_{n \rightarrow \infty} \frac{2^{4n}}{f_n} = \infty$$

Therefore $f_n = \mathcal{O}(2^{4n})$.