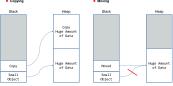
Zusammenfassung 1 MOVE SEMANTICS

Copying



With a move, only the stack data is copied. The old heap pointer gets deleted and the new one attached to the same heap data. Attention: The old stack data is still valid but in a undefined state, risk of dangling

Ownership Transfer: Resources of expiring values can be transferred to a new owner. Might be more efficient than (deep) copy and destroying original. Might be feasible when copying is not (e.g. std::::mique.ptr) Examples of such resources: Head Memory. Look. (fiel Handlet)

1.1. EXAMPLE MOVE CONSTRUCTOR (TESTAT 1)

ttainerforstyoupe... {
rForBigObject() // default constructor
reefstd::make_unique<BigObject>()} {} // make_unique creates points

untainerForBigObject(ContainerForBigObject constS other) // copy constructor
resource(std::make_unique<BigObject>(rether.resource)} {}
// creates per mointer to joid Heap object >> copy.

ontainerForBigObject(ContainerForBigObject88 other) // move constri : resource(std::move(other.resource)} {} // pointer cets changed

/ copy assignment operator
upto operator
upto operator = (ContainerForBigübject constå other) -> ContainerForBigübjectå {
 resource = std::make_unique=Bigübjectv(=other.resource);
 return =this; // the 'this' object gets returned, same as copy constructor

"(ContainerForRigDhiert&& other) → ContainerForRigDhiert& {

std::unique_ptr<BigObject> resource;

1.2 IVALUE AND DVALUE DESERVES

Ivalue references	rvalue references		
Binds to an Ivalue (everything with a name)	Binds to an rvalue (temporary objects, literals)		
T &	T &&		
The original must exist as long as it is referred to.	Can extend the life-time of a temporary.		
auto modify(T& t) → void { // manipular t } suct lvalueMerExample() → void { Tt = 5; modify(C); // modify(5) would not /// work because 5 is not a reference } t i r = t; // rarely used }	auto createI() → T; // Object can be "stolen" auto consume(T&E t) → void { // manipulate t } auto rvalueMefExample() → void { consume(T(E); // new T as param T&E t = createI(); // rarely used }		
Ivalue examples	rvalue examples		
The value that is the Ivalue,	/rvolue is in marked as [volue]		
// has name	// temporary object without page		

Ivalue examples	rvalue examples			
The value that is the Ivalue/rvalue is in marked as [value]				
// has name T value {}; std::cout << [value];	<pre>// temporary object without name int value{}; std::cout « [value + 1];</pre>			
// Function call returning lvalue ref [std::cout << 23]; // returns 'cout &' [vec.front()]; // returns 'T &'	<pre>// Function call returning value type std::abs(int n); // returns int</pre>			
// Built-in prefix inc/dec expressions	// Built-in postfix inc/dec expression a++; // returns value of a (w/out +1)			
<pre>// lvalue reference, but has a name auto foo(T& param) -> void { std::cout << [param]; }</pre>	<pre>// without reference auto create() → T; [create()];</pre>			
<pre>// rvalue reference, but has name auto print(T&& param) -> void { std::cout << [param]; }</pre>	<pre>// transformation into rvalue T value(); T o = [std::move(value)];</pre>			
// Reference has adress T& create(); [create()];	// xvalue T&& create(); [create()];			
// String literals are always lvalues std::cout << ["Hello"];	Rule of thumb: Does element keep living? I value (only copy), × rvalue (copy & move possible)			

// depends on the implementation of T value(): std::cout << [value + 1]:

1.3. VALUE CATEGORIES



Yes	No	Ivalue			
Yes	Yes	xvalue (expiring value)			
No	No (Since C++17)	prvalue (pure notue)			
No	Yes (Since C++17)	- (does not exist anymore)			
Xvalue: Expiring Value - Address cannot be taken - cannot be used as left-hand operator of built-in assignm					

erialization Conversion from Ivalue through ste Examples: Function call returning rvalue ref (ptd::nove(x)), access of non-ref members of rvalue object X x1{}, x2{}; consume(std::move(x1)); std::move(x2).member; X{}.member;

1.3.1. Temporary Materialization

Getting from something imaginary to something you can point to (e velue) prvalue to xvalue. Requires a destructor. Happens...

- when binding a reference to a prvalue (1)
- when accessing an element of a privative array (int value = (int[])(10, 20)[1]; // value = 20), when converting a privative to privative from the converting a privative to a pointer (cost int fairer = (2-3); cost int value (inter-fairer); when initializing an std::initializer_list*xi> from a broad-initializer_list*vector(x, 2);)

Property is a successful to the second of t

const Ivalue reference: Binds Ivalues, xvalues and prvalues: auto f(T const &) → void

1.4. OVERLOAD RESOLUTION FOR ERFF FUNCTIONS

	7(2)	7(5 a)	f(s const a)	7(3 88)
S s{}; f(s);	*	✓ (preferred over const&)	~	×
S const s{}; f(s);	~	×	~	×
f(S{});	~	×	~	✓ (preferred over const&)
S s{}; f(std::move(s));	~	×	-	✓ (preferred over const&)

2 TYPE DEDUCTION

template <typename T> auto f(ParamType param) → void;

2.1.1. Parantype is a value type (T)

2.1.2. Paramtype is a reference (T&)
(e.g. outo f(T & param) → void or outo f(T or

CPPReference: Templote argument deduction
2.1. FORWARDING REFERENCES AND TYPE DEDUCTION

Deduction of type T depends on the structure of the type of the corr

(e.g. stor) (r) percay — (obs) note: extern or answers per 1. <expre> is a reference type: (gnore the reference 2. (gnore the rightmost const of <expre> (char cosst + cosst — char cosst +) 3. Pattern match <expre> 's type against ParamType to figure out T

int x = 23; int const ex = x; int const & erx = x; char const

«expr» is a reference type: ignore the reference
 Pattern match <expr»'s type against ParamType to figure out 1

Examples for Const References: outp ((7 court & cores) -) with

// calls // instances
f(x); auto f(int const& param) → void;
f(cx); auto f(int const& param) → void;
f(crx); auto f(int const& param) → void;

2.1.3. Paramtype is a forwarding reference (T&&)

2.1.4. Type Deduction and Initializer Lists
When an initializer_list is used for type

template <typename T> auto f(std::initializer_list<T> param) → void; f({23}); //T = int, ParamTyre = ctd:for(+calf-c-

<u>CPPReference: Placeholder type specifiers</u>
Same deduction as above, auto takes the place of T.

Left x=23; where x=x is the first x=x is the first x=x; x=x;

template <typename T>
decltype(auto) forward(std::remove_reference_t<T>& param) {
 return static cast<T&&/(param):

auto Return Type Deduction: auto can be used as return type and for parameter declarations. Body must be

CPPReference decitive securities
Represents the declared type of a name expression, decitype (auto) allows deduction of (inline) function

return types, but does not strip references like out o. Can take an expression for specifying trailing ret, types

std:: forward is a conditional cost to an replue reference. This allows arguments to be treated as what they

- If T is of value type, T && is an avalue reference in the return expression. ($int \rightarrow int E$)
- If T is of ivalue reference type, the resulting type is an avalue reference to an ivalue re ($int = 1 \text{ int } E \rightarrow T \text{ &} E \text{ would mean 'int } E \text{ &} E \text{ 'which can be collapsed into 'int } E'$).

plate<typename D fitc outc make_buffer(T && value) → BoundedBuffer-value_type> { BoundedBuffer-value_type> new_buffer(f); new_buffer.poin(dfd:ferward<t>(value)); return new_buffer; }

auto x = 23; auto const cx = x; auto& rx = x;

2.1.6 Tune Deduction for deal two

2.1.7. Type Deduction in Lambda

2.2. PERFECT FORWARDING WITH STD::FORWARD

CPRinfermon: atd:: farward; Slides Page 25 anward

// calls // instances // deduced Ts f(x); auto f(int& param) → void; T = int f(cx); auto f(int const& param) → void; T = int const f(cw): auto f(int const& param) → void; T = int const

«expr» is an Ivalue: T and ParanType become Ivalue references (rint
 Otherwise (if couprs is an evalue): Rules for references apply (test exemple

int y = 23: int const ey = y: int const & cey = y:

int x = 23; int const ex = x; int const & erx = x;

pe deduction applies for type T (Danger) heaves can also hind to TSS as a TSJ. Only for Method Template emplates, and only for "T SS" not "std:: vector<T> SS" or "T const SS".

f(<expr>);

The overload for value parameters imposes ambiguities. For deciding two Ivalue reference overloads, the constness of the argument is considered.

OVEREDAD RESOLUTION FOR MEMBER FORCHORS						
	S::m()	S::m() const	S::=() &	S::m() const &	S::=() &&	
	Va	ue Members	Reference Members			
s(); (s);	*	~	(preferred over const5)	~	×	
const s{}; =();	×	· ·	×	~	×	
0:=O;	~	~	×	~	(preferred over const&)	
s{}; td::move(s).m();	~	~	×	~	(preferred over consts)	

1.6. SPECIAL MEMBER FUNCTIONS

Normally not necessary to implement ("rule of zero"), but if one is needed, likely all special m should be implemented ("rule of five"). Copy Constructor/Assignment usually const. Assignment operators must be member functions. Move operations must not throw exception

1.6.1. Move Constructor S(S &&)

1.6.2. Copy/Move assignment auto operator=(S const &) → S&/auto operator=(S &&) → S&

CDMsferror: Copy any comment. CPMsferrors: Many analysis and Copies/Moves the argument into the this object. Executed when the variable to copy/move to has afreed these initialized Must be member operator. Default helpsalor: Initialized has a classes and members with cross

CPDReference: Denterotors

Considerates resources held by the this object. No parameters. Must not throw (in resempt) Must be implemented if a custom copy/move is defined. Default behavior: Calls destructor of base classes & member:

Allow implementing the copy-assignment operator. Utilizes the copy constructor to create a tempological, and exchanges its contents with itself using a non-throwing swap. Therefore, it swaps the dates with new date. The temporary object is then destructed automatically. It needs a copy-construct destructor and a swap function to work.

/ copy assignment - usually not needed
uto operator=(S const& s) → S& {
 if (std::addressof(s) → this) { // avoids ut
 S copy = s; // can throw exception, this-ol
 sump(copy);

turn *this;

/ move assignment uto operator-(S&& s) \rightarrow S& { // could be measurept if (std::addressof(s) \Rightarrow this) { swap(s): // s contains value of this-object, but is okay because (

1.6.5. What you write vs what you get

makes the implicit default explicitly <u>user decl.</u>							
write/get	default ctor	destructor	copy ctor	copy assig	move ctor	move assig	
nothing	defaulted	defaulted	defaulted	defaulted	defaulted	defaulted	
any ctor	undeclared	defaulted	defaulted	defaulted	defaulted	defaulted	
default ctor	user decl.	defaulted	defaulted	defaulted	defaulted	defaulted	
destructor	defaulted	user decl.	defaulted!	defaulted!	undeclared	undeclared	
copy ctor	undeclared	defaulted	user decl.	defaulted!	undeclared	undeclared	
copy assig	defaulted	defaulted	defaulted!	user decl.	undeclared	undeclared	
move ctor	undeclared	defaulted	deleted	deleted	user decl.	undeclared	
move assig	defaulted	defaulted	deleted	deleted	undeclared	user decl.	

1.7. COPY ELISION

Disabled elision (C++14)

2x Move in x(), 1x Move in x

1.8. LIFE-TIME EXTENSION

Demonia laznik = summon();

The life-time of a temporary can be extended time ends at the end of the block. It is not train

-- S s{create()} Constructor S()

 In initialization, when the initializer is a privalue: S s = S{S{}} (Only 1 ctor cell, 0 copy c
 When a function call returns a privalue: (S() gets initialized directly in new_set} instead of in c auto create() → S { return S{} }
auto main() → int f S new swicreate()}: S * sp = new Sicreate()}: }

7.1. Make a fation take optimization, fation type, is also type, attent approach is a local stable MNOV intermed attent to helpful commonder in the action of the intermed use, the contraction must all relate of the intermed to the contraction must all relate of the action of the properties MNOV.

There is present, fation expression is a local variable from the innermost surrounding by block. The color contracted in the location where it would be moved or copied.

Only mandatory elision (C++17)

--- S s{create()} Constructor S()

and only Ω int I when then at the end of the statement Ω when Ω is a support of the statement Ω is the statement Ω is the statement Ω is the same three long enough for count, γ is the statement Ω is the statement Ω is support Ω in Ω in Ω is the statement Ω is support Ω in Ω in Ω is the statement Ω is support Ω in Ω in Ω in Ω is the statement Ω is support Ω in Ω in Ω in Ω is support Ω in Ω

Ox Move in s (), Ox Move in s

- S s{create()} -

Does not actually move objects. It's just a unconditional cost to an evalve reference. This allows resolution of auto create() \rightarrow S { S s{}; std::cout \ll "\t --- create() --auto main() \rightarrow int { std::cout \ll "\t --- S s{create()} ---\n"; S s{create()}; std::cout \ll "\t --- s = create() ---\n"; s = create();

2.3. STD::MOVE

Reference Collapsing: "T& &", "T& &&" and "T&& &" become "T&", "T&& &&" becomes "T&&"

3. HEAD MEMORY MANAGEMENT

Ifetime on Stack: Deterministic, local variables get deleted automatically upon leaving their scop

Ifetime on Heap: Creation and deletion happens explicitly with new and delete (Desperous, exelution for $() \rightarrow void \{ auto ip = new int{S}; /* ... */ delete ip; \}$ Rules: Delete every object you allocated, do not delete an object twice or access a deleted object

3.1. EXPLICIT LIFE-TIME MANAGEMENT

Global and local variables have life-time implicitly managed by the program flow. Some resources can be allocated and deallocated explicitly. This is error-prone. Guideline: always wrap explicit resource management to an object which has implicit life-time management (RAII).

3.2. POINTER SYNTAX

auto arr = new int[5]{}; // 5 needs to be complile-time-constant, auto = int
int v = arr[4]: // accessing element (* arr + 4 pointer objects)

Direct Member Access (->) struct S {
 auto member() \rightarrow void { this->value = ...; } int value; // this is a pointer to instance

Pointer Parameters: Pointers can be used as parameters. Addresses can be taken with (&) can be overriden

auto foo(int* p) \rightarrow void { } auto bar() \rightarrow void { int* ip = new int{S}; int local = 6; foo(ip); foo(&local); }

Const Pointers: Pointers can be const. const is on the right side of the *.

int const + const + const icpeppe = &icpep;

"kpcppc is a const pointer to a pointer to a const pointer to a const."

nullptr: Represents a null-Pointer. Is a literal (produe) and has type nullptr_t. Implicit conversion to any pointer type: T *. Prefer nullate over 8 and NULL (no overload emblaulty, no in ointer vs. Reference

are always bound to an object can be changed (f not o

teuto for(int* a) → void (if (a) (...))) and for modelling borrowing only. Else, use

3.3. MEMORY ALLOCATION WITH NEW

new <type> <initializer>

Allocates memory for an instance of <type>. Returns a pointer to the object or array created on the heap of type <type> *. The arguments in the <initializer> are passed to the constructor of <type>. Memory Lea wed with delete. Avoid manual allocation, use RAII instead.

struct Point{ Point{int x, int y} : x{x}, y{y}{} int x, y; }; auto createPoint{int x, int y} \rightarrow Point* { return new Point{x, y}; // constructor uto createCorners(int x, int y) → Point* {
return new Point[2]{{0, 0}, {x, y}};

3.3.1. Placement new

new (<location>) <type> <initializer>

Used for placing elements on the heap in the location of a deleted element. Does not allocate new memory place-free fairneed. The memory of Aceas those needs to be suitable for construction of a new object and any element there must be destroyed before. Calls the Constructor for creating the object at the given location and returns the memory location. Better use at all construct, at (ptr). If element there is not the construction of the cons new (ptr) Point{7, 6}; delete ptr;

3.4. MEMORY DEALLOCATION WITH DELETE

Deallocates the memory of a single object pointed to by the cpointer>. Calls the Destructor of the destroyer
type. delate nullptr does nothing. Deleting the same object twice is Undefined Behavior!

truct Point{ Point(int x, int y) : $x\{x\}$, $y\{y\}$ { int x, y; }; uto *umitinPoint(int x, int y) \rightarrow void { Point x pp = new Point $\{x,y\}$; delta destructor and releases memory

3.4.1. Placement delete

S * ptr = ...; ptr->-S(); Destroys the object, but does not free its memory. Better use std::destroy_at(ptr). Use this for non-

3.4.2. Array Memory Deallocation with delete[]

delete[] <pointer-to-array>

Deallocates the memory of an array pointed to by the <pointer-to-array>. Calls the Destructor of the destroyed objects. Also deletes multidimensional arrays. Not necessary to know exact amount of elements in the array. Undefined behavior on non-array types.

3.5. NON-DEFAULT CONSTRUCTIBLE TYPES
A type is non-default-constructible when there is no explicit or implicit defauldicate the plain memory and initialize it later (stder page 36 convert). auto memory = std::make unique<std::byte[]>(sizeof(Point) + 2 /*Array sizex/)

unto location - esistepret_cartPoint*(sizeo(Point) + 2 /%errsy sizestTo location - esistepret_cartPoint*(seeser,point)); std::construct_at(Location, 1, 2); // Equivalent to arr[0] - Point[1, 2); autro value - sizesstif(seeser,point), 0); // Access value via helper funct std::destroy_at(Location);

auto elementAt(std::byte+ memory, size_t index) -> Point& { // helper function return reinterpret_cast<Point>>(memoru))forfact|

Do not use an element if it is uninitialized and destroy them before the memory is deallocated Use a std::byte array as memory for NDC Elements.

- Static: std::array<std::byte, no_of_bytes> values_memory; (on stock, size known at compile-time)

- Dynamic: std::unique_ptr<std::byte[]> values_memory; (on heap, size known at run-time 3.6. CLASS-SPECIFIC OVERLOADING OF OPERATOR NEW/DELETE

Overloading new and delete for a class can inhibit heap allocation. This can be used to provide efficient allocation, is useful with a memory pool for small instances or if thread-local pools are used. Can log or limit number of heap-allocated instances. But in general, not advised.

struct not_om_heap { // Prevents heap allocation of this class static auto operator new[statisize_t zz] \rightarrow woid * { three statished_allocf[:] static auto operator new[statisize_t zz] \rightarrow woid * { three statished_allocf[:] static auto operator deltafe(wid *ptr) \rightarrow woid newcept { /* do nothing w } tatic auto operator deltafe(yidid**ptr) \rightarrow woid newcept { /* do nothing w } tatic auto operator deltafe(yidid**ptr) \rightarrow woid newcept { /* do nothing w } }

3.7. READING DECLARATIONS

3.7. READING DECLARATIONS

Declarations, are read storting by the declarator (name)

First read to the right until a closing parenthesis is encountered

Second read to the left until an opening parenthesis is encountered

Third jump out of the parentheses and start over

Specifiers right to the declarator: Array Declarator (()) and Function Parameter List ((<po Specifiers left to the declarator: References (65, 6), Pointers (+) and Types (4nt onst: applies to its left neighbor, if there is no left neighbor, it applies to its right neighbor. Should always

be written to the right of the type to avoid surprises.

int const * (* f [2][3]) [5]:

Fix an array of 2 elements of arrays of 3 elements of pointers to arrays of 5 elements of pointers to const. Int. nt (* f(int(*)(int))) (int); //7: int (3: * 1: f (5: int(2: *)(4: int))) (6: int) (1) is a function that takes a pointer to a function as argument (2 pointing to 4) and returns a pointer to a function (s pointing to 5) which takes an int as argument (4) and returns an int (5). The function takes an int as argument (4) and returns an int (7). Always use type aliases on cases like this!

Resource Acquisition is initialization is an alternative to allocating and dealloca allocation and deallocation in a class, use constructor for allocation, destructo

3.8.1. std::unique_ptr and std::make_unique

CPReference infrusions at a CPReference infrance unique charge (**); \$10::unique_ptr<chary < Ptr = std::make_unique<chary(**); Wagsa plain pointer, has zero runtime overhead. A custom deleter could be supplied if required. Always use nake_unique for creation. Can create unbound arrays, but not fixed size arrays. 3.8.2. Container Member Function emplace

CPPReference: std:: vector<7. Allocator :: seplace. CPPReference: std:: stack<7. Cont Constructs elements directly in a container, better than moving. Not availal std:: stack<Point> vec{}; vec.emplace(3, 5); // std:: vector requ A TTERATORS & TAGS

4.1. TAGS FOR DISPATCHING

<u>CPPReference states</u>... <u>Iterator too</u>

If the <u>same operation</u> can be implemented more/less efficiently depending on the capabilities of the argument, togs can be used to find the "best" imp Tags are used to mark capabilities of associated types. They do not contain any r

provides travelThroughSpace, the base fu

 $\label{toy-continuous} \begin{tabular}{ll} template < typename Spaceships \\ auto travelToDispatched(Galaxy destination, Spaceships ship, SpaceOriveTag) \to void { } \\ \end{tabular}$ 5.2. TEMPLATE PARAMETER CONSTRAINTS AND CONCEPTS

uto travelToDis
ship.travelThr

Allow constraining template parameters, requires is followed by a compile-time constant boolean expression. Is either placed after the template parameter list or after the function template's declarator. plate <typename Spaceship> n travelTo(Galaxy destination, Spaceship& ship) → void { template <typename T>
requires std::is_class_v<T> // either here...
auto-function(T argument) -> void requires std::is_class_v<T> /* or here */ {...} ename SpaceShipTraits<SpaceShip>::Orive drive{}; // insta /elToDispatched(destination, ship, drive); // call overla

4.2. ITERATORS Different algorithms require different strengths of iterators. Iterators capabilities can be determined at

compile time with tag types.

Outputiterator: Write results, without specifying an end (used on std:: ostreon).

Objective result, without people grapher of the Value of the Control of the Contr

CommunicationComm

You need to implement the members required by your iterator_tag

4.2.1. iterator traits o

CPPReference: atd::iterator traits

STL algorithms often want to determine the type of some specific thing related to an iterator. However, not all iterator types are actually classes. Default iterator—traits just pick the type aliases from those provided. Specialization of terator—traits also allows "naded pointers" to be used as iterators in algorithms.

4.2.2. Problems with the Stream Input Iterator

4.2.2. Problems with the survenil representation of the following the fo

namespace { //global variable to initialize the reference in an empty std::istringstream empty{}; // pseudo default

IntInputter::IntInputter() : input { empty }

Dereferencing and Equality

value_type operator=(); auto operator=(Intinputter const & other) const // just input

/ just input
wto Intinputer::operator*() → Intinputer::value_type {
 value_type value{};
 input > value;
 return value;

// and compare. Only equal if both are !good()
auto IntInputer::operator==(const IntInputer & other) const → bool {
 return !input.good() &&b !other.input.good();

2.2. Custom Herator Example [Fasted 2]
temp Lister(spossus Py
temp L

auto operator==(iterator_base const & other) const -> bool { ... }
auto operator<=>(iterator_base const & other) const -> std::strong_ordering { ... }

auto operator*() const -> decltype(auto) { ... }
auto operator->() const -> decltype(auto) { ... } auto operator[](difference_type index)

auto operator++() -> iterator_base & { ... }
auto operator++(int) -> iterator_base {
 auto comst copy = *this;
 +-(*this);
 return copy; }

auto operator--() -> iterator_base & { ...
auto operator--(int) -> iterator_base {
 auto coest copy = *this;
 --(*this);
 return copy; } auto operator+(difference_type n) comst -> iterator_base {
 auto copy = *this;
 copy += n;
 return copy; }

auto operator-(difference_type n) const -> iterator_base {
 return this->coerator+(-n): }

auto operator == (difference_type n) -> iterator_base & { ... }
auto operator == (difference_type n) -> iterator_base & { return

Boost if used, would generate operator+(int), operator--(int), operator-(difference_type n), operator-(difference_type n) with implementation shown above. Change signature to street interator_base : boost:operator_imple:random_access_timent_helperitator_basevb.

5. ADVANCED TEMPLATES

Pros of static polymorphism

Hannens at compile time

Cons of static polymorphis Longer compile-times
 Template code has to be known when used (Needs)

 Type checks at compile-time Larger binary size (copy of the used ports for each

A polymorphic call of a virtual function (whentence ownloading) requires lookup of the target function. Non-virtual calls (semplose overloading) directly call the target function. This is more efficient. 5.1. SFINAE (SUBSTITUTION FAILURE IS NOT AN ERROR)

nate overload candidates. During overload resolution the template parameters in a template is used to entirely expended calciducates. You ling oversided in constant of the empirical parameters in a comparison declaration are substituted with the deduced types. This may result in template parameters in a comparison compiled (i.e. ceiting a member/junction on a volue type). If the substitution of template parameter fails, that overfload candidate is discorreded.

5.1.1. Type traits

Other forces. Do Taste. Children and interest contact.

The standard library provides many predefined checks for type traits. A trait contains a boolean value. Usually available in a _v fertures the text result and non-_v variant fertures the integral contains.

Example: std :: is class

std::is_class<>::value; std::is_class_v<8>; // both true std::is_class<int>::value; std::is_class_v<int>; // both false std::enable_if/std::enable_if_t CPPReference add::enable_if

std::enable_if_t takes an expression and a type. If the expression evaluates to true, std::enable_if_t
represents the given type, otherwise, it does not represent a type.

auto main() \rightarrow int {
 std::enable_if_t<true, int> i; // int
 std::enable_if_t<false, int> error; // no type

std::enable_if can be applied at different places (marked with "||", only one needs to be used):

requires Expression requires also starts an expression that evaluates to bool depending whether they can be compiled.

<u>CPFBeformer: Contraints and Concepts</u>

Provide a means to specify the characteristics of a type in template context. Better error messages, more

quires (\$parameter-list\$) { /* sequence of requirements */ }

template <typename T> requires requires (T const v) { v.increment(); } auto increment(T value) → T { return value.increment(); } Type Requirements

5.2.1. Keyword requires

template <typename T>
requires requires (T const v) { v.increment() } \rightarrow std::same_as<T>; }
auto increment(T value) \rightarrow T { return value.increment(); }

5.2.2. Keyword concept

Specifies a named type requirement. Conjunctions (&&) and disjunctions ((()) can be used to combine constraints (&e. requires std::integral<f) // std::fleeting.point<f).

template <Incrementable T> // either here...
requires Incrementable<T> // ...or here
auto increment(T value) \rightarrow T { return value.ii

5.2.3. Abbreviated Function Templates

If there are two auto arguments, two template typenames T1, T2 get created.

6.1. CONSTANT EXPRESSION CONTEXT Need to be defined at compile-time.

Non-type template arguments (std::arrayclissent, 5> arr();)

Array bounds (double matrix[ROWS](GLIS]();) - Static Assertions (static_assert(order = 66);)
- constexpr / constinit variables (constexpr unsigned pi = 3;)

constexpr if-statements (if constexpr (nize > 0) ()) noexcept expressions (stob (Stab 26) noexcept (true);)

static const variables in namespace scope of built-in types initialized with a constant express put into ROMAble memory, if at all. Allowed in constant expression content. No complicated no guarantee to be done at comple

- constexpr variables are const, read only at run-time

6.2.1. constexor Functions

computations (e.g. three. Compile

consteval auto factorial(unsigned n) {

consteaps auto factorialOf5 = factorial(5); // works auto auin() → int { translof5 = 120); // works works workspecify: unsigned import(); if (ratician > input) { sticcore * (ratician > input) { sticcore * (ratician > input) { sticcore * (ratician > input) { raticion * (ratician > input) { workspecify: sticcore * (ratician > input) { sticcore * (ratician > input) { workspecify: s

Literal Types can be used in constexor functions, but only constexor member functions can be called or

6.3.1. Literal Class Type Example constexpr static size_t dimensions = 3; std::array<T. dimensions> valuesfb:

while: constaypr Vactor(T x, T y, T z) : values(x, y, z), $\{\}$ // constaypr constant auto length() const \rightarrow T $\{$ // constaypr const number function auto squares \sim 2) \sim 4() \sim 9() \sim 10 \times 2(); extern satisfayer(equivars).

:onstexpr auto v = create();
suto main() \rightarrow int { //v.x() = 1.8; // possible with constinitauto v2 = create(); v2.x() = 2.8;}

6.3.2. Compile-Time Computation with Variable Templates (for Class Template see slides page 25) Variable templates can be constexpr/constinit and defined rec

6.3.3. Captures as Literal Types

Page 1

CPLA | FS25 | Nina Grässli, Jannis Tschan, Matteo Gmür

Check whether a type exists. Starts with typename keyword. Useful for nested types like in Rounded Ruffe

requires { typename StypeS }

requires { typename BoundedBuffer<int>::value type: } Compound Requirements

Compound requirements

Check whether an expression is valid an can check constraints on the expression's type. The return requirement is optional.

requires (T v) { { \$expression\$ } → \$type-constraint\$; }

concept Incrementable = requires (T const v) { { v.increment() } \rightarrow std::same_as<T> }

Can be used in template parameter declarations or as part of a requires claus

toplate (Incrementable T)
auto increment(T value) → T { return value.increment(); }
// is equivalent to Tera Syntax
auto increment(Incrementable auto value) → T { return value.increment() }:

6. COMPILE-TIME COMPUTATION

- Case expressions (mitch(value) { case 42: /* ... s/ })
- Enumerator initializers (new Light { Off = 0, on = 1 };)

constexpr / constinit Variables are evaluated at compile-time. They are initialized by a constant expr and require a literal type (primitive date type). They can be used in contents are local scope, namespace scope and static data members.

constinit variables are non-const. They need to be initialized at compile-time, but can be changed at

Can have local variables of "literal" type. The variables must be initialized before usage.

Can use loops, recursion, arrays, references

Can cute loops, recursion, arrays, references.

Can contain branches that refy on run-time only features, if branch is not executed during compile-time

Computations (sp. roves, comparement or your an arches who company, /
- Can only cold const expr functions.
- Are usable in const expr and non-const expr context
- allocate dynamic memory that is cleaned up by the end of the compilation (since c+20)
- be virtual member functions (since c+20)

stexpr auto factorial(unsigned n) \rightarrow unsigned { /* needs to have a body */ } 6.2.2. consteval Functions

Are usable in constexpr contexts only and implicitly of

There is no undefined behavior during compile-time. Instead, there will be a compilation error. If constaxy evaluation does not reach an invalid statement, the code is still valid. 6.3. LITERAL TYPES

}
constexpr auto x() \rightarrow T& { return values[8]; } // constexpr non
constexpr auto x() const \rightarrow T const& { return values[8]; }
} // implicit default destructor is also renotewne

emplate <size_t N>
onstexpr size_t factorial = factorial<N-i> + N; template o // Base case

ure types (the types returned by lambdo expressions) are literal types as well. The can be used as types of texpr variables and in constexor functions

operator"" _UDLSuffix()"

Allows integer, flusting point, character, and string iteration produce objects of user-defined type by defining a user-defined type. By defining a user-defined soften. The suffice most start with an undercome, it allows to seld demotion, convention, extend in probable, defined to operated functions are consistent. Our probable defined to operate functions are consistent, or consistent in probable defined to operate functions are consistent. Our probable defined to probable defined on a finishment of the control of the control operation operatio

amespace velocity::literals {
onstexpr inline auto operator*"_kph(unsigned long long value) → Speed<kph
⟨ return Speed<kph>⟨safeToDouble(value)}; } // user defined literal operat

uplate<typename T>
ncept arithmetic = std::is_arithmetic_v<T>; // allows ints & floats in + ope

static constexpr inline auto safeToBouble(long double value) → double {
 if (value > std::nomeric_limits<double>::max() ||
 value < std::nomeric_limits<double>::min()) { throw std::invalid_argument{"Value must be within double rance"}

unn static castedoubles(value):

Truct space {
constexpr explicit Speed(double value) : value(value) {};
constexpr explicit operator double() const { return value; } // conversion operator
auto constexpr operator(arithmetic auto rhs) → decltype(rhs) { return value * rhs } ivate: druble value{};

uto constexpr operator*(arithmetic auto lhs, Speed rhs) -> decltype(lhs)
{ return rhs * lhs; } // Make * operator commutative (value * Speed calls Speed * val

Signatures: unsigned long long for integral constants, long double for floating point constants, (char const *. size t len) → std::string for string literals, char const * for a raw UDL operator //

6.4.1. Template UDL Operator

template <char...>
auto operator""_suffix() → TYPE

Has an empty parameter list and a variadic template parameter. Characters of the literal are templat arguments. Since C++20, the template UDL operator works with string literals as well (example and compile-ter

Standard Literal Suffixes: Do not have leading underscore: std::string μ_0 , std::complex μ_0 ψ_0

Object-like Macros

#define identifier replacement-list new-line

Identifier is a unique name in ALL CAPS. Is valid until #undef NAME. Replacement-list is a possible empty

Function-like Macros

6.5. PREPROCESSOR

#define identifier (identifier-list?, ?...?) replacement-list new-line Features an optional parameter list, containing only names. Params with # prefix turn into string:

Includes: Textual inclusion of another file. #Include "poth" for including a header file from the same of

or workspace, #Include <path> for external includes.

Conditional includes: Enable a section depending on a condition. (example and macros in skides on page 52 - 62). #ifdef identifer new-line

#if constant-expression new-line #elif constant-expression new-line #else new-line #endif new-line

7. MILLTT-THREADING & MILTEX

itd:: nutex and co. to facilita 7.1. API OF STD :: THREAD

A new thread is created and started automatically. Creates a new execution context, join() waits for the thread to finish. Besides (nombdos, functions or functor objects can also be executed in a thread. The return value of the function is (gnored. Threads are defquite.constructible and moveable. Courtion: Program terminations). nates if thread gets destructed without calling join() before

struct Functor {
 auto operator()() const -> void { std::cout « "Functor" « std::endl; }

function() → void { std::cout ≪ "Function" ≪ std::endl; }
main() → int d t {
 ctionThread(function);

std::thread functorThread{Functor{}}; functorThread.join(); functionThread

Streams: Using global streams does not create data races, but sequencing of characters could be mixed std::this_thread helpers: get_id() (An id of the underlying OS thread), sleep_for(durat

7.1.1. Passing arguments to a std::thread templatecclass Function, class... Args> explicit thread(Function&& f, Args&&...args);

The std :: thread constructor takes a function/functor/lambda and arguments to forward. You should pass all arguments by violue to avoid data races and dangling references. Copturing by reference in humbdashard data as well!

auto fibonacci(std::size_t n) \rightarrow std::size_t { /* ... */ } auto printfib(std::size_t n) \rightarrow void { auto fib = fibonacci(n); /* print... */ } auto main() \rightarrow int { std::thread function { printfib, 46 }; function.jcin(); }

Before the std::thread object is destroyed, you must join() (work until file

7.2. STD::JTHREAD CPPReference: std:: ithre RAII wrapper that aut ically calls 10in(). Also supports external stop requests (t. request atea(), similar

Concellation Token in CA). 7.3. INTER-THREAD COMMUNICATION

mor. std.: sutes. CPPR/mmor std.: shared sutes inication happens with mutable shared stat access or make access atomic.

All mutexes provide the following operation

Two properties specify the capabilities

Recursive: Allow multiple nested acquire operations of the same thread (prevents self-decided)
 Timed: Also provide timed acquire operations (try_tock_fer(devotion), try_tock_until(time))

Reading operations don't need exclusive access. Only concurrent writes need exclusive locking. Uss std::shared_mxtex/std::shared_timed_mxtex with lock, shared() to read with multiple threads. Shared mutex also have the exclusive lock methods from a regular mutex.

7.3.1. Acquiring / Releasing Mutexes

eck overd: RAII wrapper for a single mutex. Locks immediately when constructed, unlocks whe

destructed ed_lock: RAII wrapper for multiple mutexes. Locks immediately when constructed, unlocks when destructed. Acquires multiple locks in the constructor, avoids deadlocks by relying on internal

sequence. Blocks until all locks could be acquired. (atd::acoped_teck both(mx, other.mx);)
std::unique_teck: Mutex wrapper that allows defered and timed locking. Similar interface to timed

mutex, allows explicit locking/unlocking. Unlocks when destructed. std::shared_lock: Wrapper for shared mutexes. Allows explicit locking/unlocking, un

Standard Containers and Concurrency: There is no thread-safety wrapper for standard containers. Access

to different individual elements from different threads is not a data race. Almost all other concurrent uses of containers are dangerous, shared_ptr copies to the same object can be used from different threads, but accessing the object itself can race if non-const (reference counter is atom)

7.3.2. Threadsafe Guard Example (Testat 3)

Scoped Lock Pottern: Create a lock guard that (un)locks the mutex automatically. Every member function

is mutually exclusive because of scoped locking pattern. Strategized Lock Pottern: Template Parameter for

using guard = std::lock_guard<MUTEX>; using lock = std::unique_lock<MUTEX>; template_stynepuse_EngerTX

emplate <typename ForwardType>
uto push(ForwardType &&t) → void {
 guard lk{mx}; q.push(std::forward<F

/ wait requires timed locking therefore unique_lock (lk, [this] return [a.empty():]): // checked once. no busy wai:

uto try_pop(T & t) \rightarrow bool {
quard lk-{mx}; if (q.empty()) { return false; } t = q.front(); q.pop(); return true; // call container empty, not this→empty, would cause deadlock
nuto emoty() const → bool { guard lk{mx}; return q.empty(); }

std::queue<T> q{};
mutable MUTEX mxf}: // mutable to unlock in const member functions

CONDITION not_empty{}; er, don't need to be swap()-ed. But notify t 7.4. RETURNING RESULTS FROM THREADS
We can use shared state to "return" results. Acquire lock in prod result, read the result was result.

o main() ~ inf {
to make , *stiristice(f); auto finished = std:condition_variable(f); auto shared = 8;
tot thread = std:thread(f(d);
stotithics_thread intege_for(Ta);
auto guard = std::locs_quard(purtar);
sufto guard = std::locs_quard(purtar);

*inished.wait(lock);
std::enut or "The answer is: " or shared or "\n": thread ininf): } 7.4.1. std::future nos: atd:: future

available,
= mit():blocks until available,
= mit():blocks until available or timeout elapsed,
= mit_for(-timeout-): blocks until available or the timeout near the cash
= mit_mit(<time>): blocks until available or the timeopoint has been reached
and then get the result
= get() - blocks until available and returns the result value or throws).

The destructor may wait for the result to become available

7.4.2. std::promise

Promises are the origin of futures. They allow us to obtain a future using get_future() and publish results

promise.intb promise(); auto result = promise.get_future();
thread = std::thread{
}{std::this_thread::sleep_for(2s); promise.set_value(42); }

}
std::this_thread::sleep_for(is);
std::this_thread::sleep_for(is);
std::saut ~ "The answer is: " < result.get() < "\n"; thread.join(); }

nouting asynchronously. It allows us to return our result from our computation

function. Additionally, it cotches all exceptions and propagates template<typename Function, typename ...Args> auto async(Funtion&& f, Args&&... args) \rightarrow std::future</* implicitly from f */>;

main() → int {
co the_answer = std::async([] { /* calculate a while ... */ return 42; });
:::cout < "The answer is: " << the_answer.get() << "\n"; }</pre>

The function centures a std:://stw.th.th.th.ts.tent heresult.yev.ly.iv.ir.ir.

**The function centures a std:://stw.th.th.th.ts.tent heresult.yev.ly.iv.it.for the result to be available.

**std::async:can take an argument of type std::/launch/anyopaus/.This is an anna with neumerators a

**std::async:can take an argument of type std::/launch/anyopaus/.This is an anna with enumerators a

**std::async:can take an argument of type std::launch/anyopaus/.This is an anna with enumerators a

**std::async:can take an argument of type std::launch/anyopaus/.This is an anna with enumerators of if we need

**result or no., std::launch::async:can take an argument of its enumerators of its enumerators

The C++ Standard defines an abstract machine which describes how a program is executed. Platform specifics are no longer relevant with this abstraction. Represents the "minimal viable computer" required to execute a C++ program. The abstract machine defines in what order initialization takes place and in what order a program is executed. It defines what a thread is, what a memory location is, how threads interact and what constitutes a data roce.

Memory Location: An object of scalar type (Ant bits as the architecture defines, like 64-bit).

Conflict: Two expression evaluations run in parallel. Both access the same Memory Location (relience).

Data Race: The program contains two conflicting actions. Undefined Behavior!

8.1. MEMORY MODEL + Memory Model defines when the effect of an operation is visible to other threads and how and

The C++ Memory Model defines when the effect of an operation is visible to other threads a when operations might be recordered. The Memory ordering define when effects become visibl - Sequentially-consistent flame a cost ordering and the definal behavior.), - Acquire/Release (Worlder purmaters than sequentially-consistent), - Consume (Discoveped, sliptly wester than sequentially-consistent),

Visibility of effects:

Read/Writes in a single statement are "unsequenced": std::cout « +1 « +1:

8.1.1. Atomics

Template class to create attomic types. Attomics are guaranteed to be don-once free. There are several specializations in the standard library. The most basic type std::atomic_flag is lock-free. utWhenReady(std::atomic_flag & flag, std::ostream & out) → void {

hile (flag.test_and_set() /*set flag to true and returns old valuex/) Here is thread: " < get_id() << std::endl;
ar(): // sets the flag to false

wto main() → int {
 std::atomic_flag flag { };
 std::thread t { [&flag] { outputWhenReady(flag, std::cout); } ;
 outputWhenReady(flag, std::cout); t.join(); }

When creating your own atomic type with std::atomic<T>, the atomic member operations are:

roid store(T) (set the new value)

Load() (get the current value)

exchange(T) (set the new value and retur

bool_compare_exchange_weak(T & expected, T desired) (compare expect
the current value with desired, otherwise replace expected with current value. May spuriously
bool_compare_exchange_strong(T & expected, T desired) (connet fell as

Applying Memory Orders: All atomic operations take an additional argument to specify the memory order std::memory_order).e.g.flag.clear(std::memory_order::seq_cst);

reguental Considery (see, exit) Global execution order of operations. Every thread observes the same reder. This is the Default behavior. The latest modification will be available to a read. equire (sequire) No reads or writes in the current thread can be reordered before this load. All writes in which through that relieise the same othering are visible in the current thread. You grownsteed to see latest

other timeast that revesse the same atomic are visible in the current timead, not guaranteed to see lottest white [Half Fench, but ordering is consistent.

Release [Pelasea]: No reads or writes in the current thread can be reordered offer this store. All writes in the current thread are visible in other threads that orquire the same atomic.

Acquire/Release [sec_rex]: Guaranteed works on the letter value untils Acquire, used for Read-Modify-

Acquire/Release (e.g., rel.): Guaranteed works on the latest value unlike Acquire, used for Read-Modify-Write operations (rull Frence) e.g., tark_mal_set(...). Relaxed (rel.awel): Does not give promises about sequencing. No data-races for atomic variables. Order can be inconsistent (journel exal/pizerse()) (Lost Updates), but may be more efficient. Difficult to get right. Releaze/Consume: Do not use! Data-dependency, hard to use. Better use acquire. 8.1.2. Custom Types with std.: atomic

Custom types need to be trivially copyable. You cannot have a custom copy ctor, move ctor, copy assignment or move assignment. Object can only be accessed as a whole. No member access operator.

8.2. VOLATILE

Volatile in C++ is different from volatile in Java and C#. Load and store operations of volatile variables mus not be elided, even if the compiler cannot see any visible side-effects within the same thread. Prevents the compiler from reordering within the same thread (but the hardwore might receder instructions onyway). Useful when

Interrupts are events originating from underlying system which interrupt the normal execution flow of the program. Depending on the platform, they can be suppressed. When an interrupt occurs, a previously registered function is called "hierowysterics inconsine" cross, Should be short and must run to completion. After the interrupt was handled, execution of the program resumes

Data shared between an ISR and the normal program execution needs to be prob atomic modifications need to become visible. Volatile helps because it suppres nterrupts may need to be disabled temporarily to guarantee atomicity.

9. NETWORK & ASYNC

Sockets are an abstraction of endpoints for communication.

— TCP Sockets are reliable, stream-oriented and require a connection setup

— UDP Sockets are unreliable, datagram-oriented and do not require a connection.

Accept: Block caller until a connection request arrive ect: Actively attempt to establish a connection

elve: Receive some data over the connection

Close: Release the connection 9.1. DATA SOURCES / BUFFERS

tion buffers. The ASIO library doesn't manage memory Fixed Size Buffers: a size: buffer (). Must provide at least as much memory as will be read. Can use several standard containers as a backen-di-Pointer + Size combinations are available.

Dynamically Sized Buffers: as is city-dynamic_put-fer(). Use if you do not know the required space and with

std::string and std::vector. Streambuf Buffers: asio::streambuf. Works with std::istream and std::os 9.2. ASIO LIBRARY

9.2.1. Example: Synchronous TCP Client Connection with ASIO
Socket: All ASIO Operations require an I/O context. Create a TCP Socket using the context.

asio::io_context context{}; // multiple contexes possible asio::ip::tep::socket socket{context}; // multiple sockets per context possible

Connect: If the IP address is known, an endpoint can be constructed easily, socket.connect() tries to auto address = asio::ip::make_address("127.8.8.1"); auto endpoint = asio::ip::tep::endpoint(address, 88); socket.connect(endpoint);

A resolver resolves the host names to endooints, as io :: connect() tries to establish a connection b::ip::tcp::resolver resolver{context};
o endocints = resolver.resolve(domain, "88"); asio::connect(socket, endpoints);

Write: asio :: write() sends data to the peer the socket is connected to. It returns when all data is sent or

std::ostringstream request(); request « "GET / HTTP/1.1\n\n"; request « "GET / HTTP/1.1\n\n"; asio::write(socket, asio::bwffer(request.str()));

Read: asio::read() receives data sent by the peer the socket is connected to. It returns when the read-buffer is full, when an error occurred or when the stream is closed. The error code is set if a problem occurs, or the stream by been closed in securior.

constexpr size_t bufferSize = 1824; std::arraycchar, bufferSize> reply(t); asio::error_code errorCode(t); asio::buffer(reply.dsta(), bufferSize), error unto readingth = asio::read(socket, asio::buffer(reply.dsta(), bufferSize), error

Advanced Reading: as 10:: read also allows to specify completion conditions. asio::transfer_all() (Default behavior, transfer all or until buffer is full)
asio::transfer_at_least(std::size_t bytes) (Read of least bytes asio::transfer_exactly(std::size_t bytes) (Read exactly bytes number of bytes

asio::read_until allows to specify conditions on the data being read.

- Simple matching of characters or strings (medy unit "a") or more complex x

- Allows to specify a callable object (predictor -- cen be iterated over, returns sent

Dipsets add::psir/citerator, best- specific (iterator- begin, iterator- end)

Close: shutdown() closes the read/write stream associated with the socket. The destructor cancels all nending operations and destroys the object socket.shwtdown(asio::ip::tep::socket::shwtdown:both); // close read and write end

9.2.2. Example: Synchronous TCP Server with ASIO Socket, Bind & Listen: An acceptor is a special socket responsible bound to a given local end point and starts listening automatically.

asio::io_context context{}; asio::ip::tep::endpoint localEndpoint{asio::ip::tep::v4(), port}; //uses an available IPv4 asio::ip::tep::acceptor acceptor(context, localEndpoint); Accept: accept() blocks until a client tries to establish a connection (with connect). It returns a new sorker

asio::ip::tcp::endpoint peerEndpoint{}; // information about cl: asio::ip::tcp::socket peerSocket = acceptor.accept(peerEndpoint)

Handling multiple requests simultaneously: Using synchronous operations blocks the current thread. Asyr ronous Operations allow further processing of other requests while the async operation is executed. Most OS support asynchronous IO operations.

The program invokes on async operations.

1. The program invokes on async operation on an I/O object (necket) & passes.

2. The I/O object delegates the operation and the callback to its io_context.

3. The OS performs the asynchronous operation.

3. The OS signois the Lo_context that the operation has been completed.

3. When the program calls io_context:::vun() the remaining asynchronous.

for the result of the operating system).

6. Still inside the 10 context::run() the completion handler is called to handle the result of the asynchro-

9.2.4. Asynchronous Read/Write on Sockets

They return immediately. The operation is processed by the executor associated with the stream's asio:io_context. A completion handler is called when the operation is done.

of the state of th

The constructor creates the server. It initializes its acceptor with the given 10 context and port. It calls ccept() for registering the accept handler and does not block. To use it, create an io_context and the erver. The executor will run until no async operation is left thanks to async_accept(). It is important that

the server lives as long as async operations on it are processes struct Server { using tcp = asio::ip::tcp;
Server(asio::io_context &c, unsigned short port)
 : acceptor{c, tcp::endpoint{tcp::v4(). aort)} {

cept() → void {
acceptHandler = [this](asio::error_code ec, tcp::socket peer) acceptHandler = [this](asio::error_code ec, tcp::socket p
 (!ec) {
 auto session = std::make_shared<Session>(std::move(peer));
 session=>start();

9.2.6. Session with Asynchronous IO

Los. Jesulon water Asymmetrous are the constructor stores the socket with the client connection. start() initializes the first async read, lead() invokes async reading, write() invokes async writing. The fields store the data of the session belte where, from this is needed because the session object would die at the end of the accept handler, therefore it needs to be allocated on the heap. The handlers need to keep the object alive by pointing on it on the heap. If there is no pointer to the object left, it gets deleted

Otto - until 4 mark () to lvate:
auto read() → void; auto write(std::string data) → void;
asio::streambuf buffer{}; std::istream input{&buffer}; asio::ip::tep::socket sock

de in Accept handler
session = std::make_shared<Session>(std::move(peer)); session->start();

code in Read handler
Session:read() → void { auto readComplHandl = [self = shared_from_this()] /*...*/}

/ Code in Write handler
uto Session::write(std::string input) → void {
 auto data = std::make_shared<std::string>(input)
 auto writeCompletionHandler = [self = shared_fr

9.2.7. Async Operation without Calibacks wrations can work "without" callbacks. Specify "special" objects as callbacks.

use_future: Returns a std:: future<1>. Errors are communicated via excep

asso:: secaches: ignores the result of the operation asso:: use_awaitable: Returns a std:: awaitable<T> that can be awaited in a coroutine. Complicated!

9.3. SIGNAL HANDLING

Most OS support signals. Signals provide asynchronous notifications. They are used to gracefully terminate a program, communicate errors, notify about traps ("it's a traps!") SIGTERM (Terminodian requested), SIGSEGV (tracitic memory access), SIGINT (User Interrupt), SIGILL (II

9.3.1. Signal Hamelling in ASIO
assion: stgnal_sate definer a set of signals to wait for Handler can be set up with signal_sate
signal_sate.signal_sate signal_sate.signal shandler receives the signal that occurre
error if the wait was aborted. Useful to cleanly stop server applications.

9.4 ACCESSING SHAPED DATA

rands are a mechanism to ensure sequential execution of handlers.

Implicit stronds: If only one thread calls is_context.run() or program logic ensures only one operation is in progress at a time.

Explicit stronds: For multiple threads on the same io_context. Objects of type asia::strand<...>

globally accessible to results = std::vector<int> { }; auto strand = asio::make_strand(context). In connection class
is::async_read(scoket, asio::buffer(buffer),
asio::bind_executor(strand, [8][auto err, auto bytes) { // wrap access in bind_executo
asio::bind_executor(strand, [8][auto err, auto bytes) { // wrap access in bind_executo
asio result = parse(buffer): results.oush back(result): /* bve bve data race */ H):

18 ADVANCED LIBRARY DESTON

10.1. EXCEPTION SAFETY does not cotch, just forwards exceptions). Exception safety is important in generic code that manages resources or data structures (might coil user-defined operations, must not grable its data structures and must not leak resources)

10.1.1. Safety Levels (from highest to lowest)

nexcept disc no-thres: Will never throw an exception (and a always accessful. Very hard to achieve, sometimes even impossible, e.g. memory allocation. (Rumples: Swap, Move Contractor, Move Assignment Operator, std.): wester-Oricities; Original States (Swap, Move Contractor, Move Assignment Operator, std.): strong exception safety: Operation succeeds and doesn't throw, or nothing happens but an exception is

Basic exception safety: Does not leak resources or garble internal data structures in case of an exception (quarantees invariance), but operations might be only half-done (i.e. if copy throws, but in

uto push(value_type const & elem) \(\to \text{void} \) {
\text{value_type val(elem); } // \text{sight throw due to copy} \)
\text{tail_= (tail_+ 1) \(\text{x}(capacity() + 1); \text{elements_++; } /* \text{internal data not yet changed*/} \) No guarantee: Often unintentional, but hannens invalid or corrunted data when an exception is thrown

A function can only be as exception-safe as the weakest sub-function it calls!

Basic Guarantee

10.1.2. noexcept Ker ot (noexcept(f())) means "Outer function is noexcept, if function f() is also not

The compiler might optimize a call of a nexcept function better. But if you throw an except nexcept function, std::terminate() will be called.

10.1.3. Member Functions that should not throw rs must not throw when used during stack unwinding Move construction and move assignment better not throw (This is why it often uses as

copying might throw, when memory needs to be allocated

10.1.4. Standard Libary Helpers <u>PPReference std::nove_if_neexcept</u>
It may be hard for a container to implement its move operations if the element type does not support

noexcept-move. Use std::move_if_noexcept instead: If noexcept, then mov

There are other helpers like is nothrow . constructible, move_constructible, default_constructible, assignable, move_assignable, copy_assignable, destructible, copy_constructible, samppable

10.1.5. Wide and narrow contracts roct: A function that can handle all argument values of the given parameter types successfully.

It cannot fall and should be specified as noexcept(true). this, global and external resources are also possible parameters.

Narrow Contract: A function that has preconditions on its parameters, e.g. int parameter negative. Even if not checked and no exception is thrown, those functions should not be not allows later checking and throwing.

10.2. PIMPLIDIOM (POINTER TO IMPLEMENTATION IDIOM) ruct S; // Forward Declaration to foo(S & s) \rightarrow void { foo(s); /* S s{}; is invalid */}

GEDBritanse: Zhous

Opaque Types: Name known, but not the content. Introduced by a forward declaration. Can be used for cointers and references, but it cannot dereference values or access members without a later definition.

struct S{}; // Definition auto main() → int { S s{}; foo(s); } Problem: Internal changes in a class' definition require clients to re-compile (s.g. changing appendip protein member varieties). Solution: Create a "Compilation Firewall". Allow changes to implementation without the need to recompile users. It can be used to shield client code from implementation changes, meaning you must not change header files your client relies upon.

Put in the "exported" header file a class consisting of a "Pointer to Implementation" + all public m With std :: shared intercal ass Tana > we can use a minimal header and hide all details in the implementation WizardImpl. The Wizard class called from the header delegates all calls to WizardImpl

std::shared_ptr<class WizardImpl> pImpl;

subtrainer_invectors maximum to pump. // impresentation or mixed subtrained in the mixed subtrained subtrain

enterclass. Ten1> we need to define the destructor of Wilserd after the definition of 12 RIITI D SVSTEMS WizardImpl. The compiler can't move the destructor by himsel

std::unique_ptr<class WizardImpl> pImpl; // ... Wizard::~Wizard() = default: bblic: Bizard(std::string name); -Wizard(); // define explicitly auto deMagic(std::string wish) -d::string; };

Because the default deleter of std::unique_ptr can't delete an incomplete type, we need to define the destructor explicitly.

New should objects be copied?

- No copyling - only movings std::unique_ptr<class Impl> (declore destructor on - Sholllow copylings: std::shared_ptr<class Impl> (having the implementation)

- Deep copyling: std::shared_ptr<class Impl> (with DPC copy constructor, dylout for

Never do o Inpl = nullate, and do not inherit from PIMPL class.

11 HOUDGLASS THTEDEACES

11.1. APPLICATION BINARY INTERFACES (ABI)

ABIs define how programs interact on a binary level (Names of structures and functions, calling conventions, instruction sets). C++ does not define any specific ABI, because they are tightly coupled to the platform. They change between OSes, compiler versions, library versions, etc. Different STL implement. are (usually) incompatible Use C(89) as an "intermediate" layer, as "C frontend for our C++ code". This is a extremely stable ABI. No



11.1.1. Example: Wizard Class

me = "Rincewind") : name{name}, wand{} {} ing const & wish) \rightarrow char const *;

o learnSpell(std::string const & newspell) → void; o mixAndStorePotion(std::string const & potion) → void; o getName() const → char const + { return name.c_str(); } g wizard_client::Wizard; Wizard const magician();
ASSERT_EQUAL("Rincewind", magician.getName());

rekground C APE Abstract data types can be represented by pointers. Ultimate abstract pointer void *
ember functions map to functions taking the abstract data type pointer as first argument. Requires factory and disposal functions to manage object lifetime. Strings can only be represented by chan *. Make sure to
t return pointers to temporary objects. Exceptions do not work across a C API, use a Error struct.

struct Wizard * wizard; // Wizard can only be accessed through pointers

THREE THREE CONTROL OF CHARMES OF THE THREE CONTROL OF CHARMES OF

Parts of C++ that can be used in an extern "C" interface: Functions, but not templates. No overloa
 C primitive Types (char, int, double, vaid)

Pointers, including function pointers

Dealing with Exceptions: You can't use references in C API, you must use pointers to pointers. In case of an renor, allocate error value on the heap. You must provide a disposal function to clean up. Internally, you cause C++ types, but you should return char-const +, because the caller owns the object providing memory. Creating Error Messages from Exceptions: Call the function body and catch expressions. May be then to

Implementing the Opaque Wizard Type: Wizard class must be implemented. To allow full C++ including templates, we need to use a "trampolin" class. It wraps the actual Wizard implementation.

or object, set the pointer pointed to by out_error. Passed out out_error must not be nullptr

From Handing at Client Side: Client-side C++ usage requires mapping error codes back to exceptions.

Unfortunately, exception type doesn't map through. But you can use a generic standard exception (strict-restate, exception). There is a dedicated Rail casts for disposal. You could also use a temporary object with throwing destructor, but this is tricky because of possible leaking. ct ThrowOnError(rowOnError() = default; hrowOnError() moexcept(false) { if(error.opaque) { throw perator error_t+() { return &error.opaque; } rror_t opaque; private: ErrorRAII error(nullotr):

struct Wizard {
Wizard(std::string const & who = "Rincewing twizard(std::string const & who = "Rincewing twizard(who.c_str(), Throw())

ovides a simple interface to Clibraries. Consists of a single JAR file and is cross-platform. For C++/Java Type mappings see slides page 31. 11.2.1. Loading Libraries

Calling the loaded library handle INSTANCE is only by convention. The loader searches for a suitable library, first in the path specified by jna.library.path, otherwise in the system default library search path. Fallback is the class path

11.2.2. Interfacing with Functions extern "C" { void printInt(int number): } Function names and parameter types must motch. However, the types are not validated. Parameter names

constructor & assignment to prevent a double free

11.2. JAVA NATIVE ACCESS (JNA)

11.2.3. Interfacing with Plain structs (see exemple on sides storing of page 34)

Plain non-opaque struct types must inherit from Structure. You must of can use the tag-interface Structure. SyVelue. You can access pointers to

11.2.4. Working with Raw Byte Arrays

// WizardImpl.cpp class WizardImpl { /* ... */ };

target: prereq_target
prereq_target: prereq_file other_target
command_to_generate_output

\$ cmake .. # configure build environme \$ cmake --build . # build the project \$ rfset --output-on-failure # run ctest

Project Layout (General) Project Layout (Libraries)

Example a = 5; // "Ctr"
Example b = std::move(a); // "Hove Ctr"
Example c = Example(std::move(f)); // "Move Ctr"
Example castd d = c; // "Copy Ctr"
Example e = std::move(d); // "Copy Ctr"

3. Temporary Ivalue in methods

12.1. BUILD AUTOMATION urtivity maintainahility and shareahility. There are many IDEs which help build our projects. But sometime, you don't want to rely on an IDE (like o

agnostic configuration language

Build Script Generators: Generate configurations for Make-style Build Systems or Build Scripts, configurations for Make-style Build Systems or Build Scripts, configurations in dependent of actual build tool, often have advanced features like download dependencies.

only executed if required

12.3. BUILD SCRIPT GENERATORS

Has support for many languages and is platform independent

target_caspit_petrors_...) defines which language features are used by the target i.e. the C+Sandard or specific features. Peter petrojne standard starber than specific features!

Language features are used to define the behalds examply part of the target. Cashe is non-specific example. The cashe is no specific example features are used to define the behalds example part of the target. Cashe is non-specific example features are used to the cashe is no specific example features. In a petroportion of the target properties are discussed in the cashe is the cash

global variables like PROJECT NAME. PROJECT SOURCE DIR or PROJECT BINARY DIR. Can be used in place of

12.3.2. Testing with CMake cludes CTest. Enable it with enable testing(). Create a "Test Runner" executable

- include: Contains headers. Add subfolders for separate sub

test: Contains tests. Add above folders as necessary Build conflig files should be in the project root

When creating a library, introduce another layer of nesting to avoid filename clashes in clients

Managing lifetime is not trivial. Using dispose . . . () API functions in finalizers is not recommended. Eith provide a dispose method on your Java type or implement AutoClosable and use your objects with try

Do not write your own scripts for this process, because then every source file gets built every time, the commands tend to be platform specific, build order must be managed manually and scripts tend to become

12.2. BUILD TOOLS

12.2. BULD TOOS.
There are plenty of build-foots: GNU make, Score, Ninja, CMake, autotools, ...
Features of Build Tools: Incremental builds, parallel builds, automatic dependency resolution, package management, automatic test execution, placetom independency, additional processing of build products.

Different Classes of Build Automation Software:

- Make-style Build Pools: Run build scripts, produce your final products, often verbose, use a language

to make tool to build all kinds of projects. Many IDEs "understand" make projects. The Workflow in Mokefile via "Target" rules. Each target may have one or more prerequisites and execute commands to senerate one or more results. Targets are then executed "too-down". A target is

\$ mkdir build # create seperate build directory for the build outputs
\$ cd build

Variables can be defined using set(VAR_NAME_VALUE). They are referenced using \${VAR_NAME}\$. There are

enable_testing()
add_executable("test_runner" "Test.cpp")

12.4. PROJECT LAYOUT

see: Contains implementations. Subfolder layout should match include folder Lib/third_party: Contains external resources like libraries

13. MOVE SEMANTICS DUTPUT

3 Constructor elider creation of temporaries (On auto defaultExample()

Example { return Example(5); }

Example { fexample(5)}; // "Ctr"

Example { feraultExample()}; // "Ctr"

Opaque struct types should inherit from Pointer and provide a constructor using the create ... () function

Use IntByReference to retrieve the size of the buffer. Requires that the API supports it. getByteArrey() copies the data from the buffer. Make sure to free the buffer either using an API free...() functions or

add_executable(...) defines binaries target_compile_features(...) defines which language features are used by the target i.e. the C++

concrete values: add_executable(\${PROJECT_NAME} "source1.cpp" "source2.cpp" ...).



set Example {
rample(int a) { std::cout « "Ctr\n"; }
rample(Example const&) { std::cout « "Copy
rample(Example&) { std::cout « "Move Ctr\n
Example() { std::cout « "Utor\n"; }