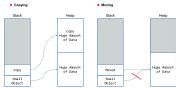
1 MOVE SEMANTICS



Sometimes, it is desirable to a - With a *copy*, the stack and h

- With a move, only the stack data is copied. The old heap pointer gets deleted and the new one is attached
 to the same heap data. Attention: The old stack data is still valid, but in a indeterminate state, risk or dangling pointers
- Ownership Transfer: Resources of expiring values can be transferred to a new owner. Might be more efficient than a (deep) copy and destroying the original. Might be feasible when copying is not (e.g. add::omique_pts Examples of such resources: Head Memory. Lock. (Fiel Handles)

1.1. EXAMPLE MOVE CONSTRUCTOR (TESTAT 1)

ntainerformigumjeu: q freefiglögisch() // Default constructor ~=fatd::make_unique<BigObject>()} {} // make_unique creates heap obj. & pointer

- erforBigObject(ContainerForBigObject consts other) // Copy constructo urce{std::make_unique-BigObject>(*other.resource)} {} utes a room of the bean object and a new mointer to the copy
- ntainerForBigObject(ContainerForBigObject&& other) // Move constructor
 resource(std::move(other.resource)) {}
 // Now pointer points to the same heap object, 'other' is valid, but :
- signment operator tor=(ContainerForBigObject const& other) -> ContainerForBigObject& { *sd::make_unique*BigObject>(*other.resource); // Same as copy cons this; // The "this" object gets returned
- / mwie assignment operator
 with operator-(ContainerPerligDbject& (
 wing std::map; // Search same in namespace scope, failback to std-imple
 swap(resource, other.resource); // Pointers get sampped
 //resource std::move(other.resource) is possible too, same as move come
 return = this:

1.2. LVALUE AND RVALUE REFERENCES

Ivalue: Everything that has an identity to name! The address can be taken. A lyalue reference can be used

return type (netwoed object must savvive, i.e. a reference parameter or vector.at ()). Beware of dangling references!

Realue: Disposable values without an address or a name. Either a literal, a temporary object or an explicitly converted Ivalue. If an rvalue is passed to a rvalue reference parameter, it gets a name and turns into an Ivalue. See Chapter 2.2. "Perfect forwarding with std::forward".

Binds to an rvalue (temporary objects, iterats)
T && T
Can extend the life-time of a temporary.
auto createT() → T; /** can be "stolen" (not needed anymore) auto consume(T&& t) → void { // manipulate t } // manipulate t } auto rvalueMeffexample() → void { consume(T&}); // new temporary T as param T&& t = createT(); // rarely used }
rvalue examples

The value which is the Ivalue/rvalue is marked as [value]			
// Has a name T value {}; std::cout « [value];	// Temporary object without a name int value{}; std::cout << [value + i];		
<pre>// Function calls returning a lvalue ref [std::cout << 23]; // returns 'std::cout &' [vec.front()]; // returns 'T &'</pre>	<pre>// Function call returning a value type std::abs(int n); // returns 'int'</pre>		
// Built-in prefix inc/dec expressions ++a; // returns 'T &'	// Built-in postfix inc/dec expressions a++; // returns 'T', the value without +1		
<pre>// lvalue ref, but has a name auto foo(T& param) → void { std::cout ≪ [param]; }</pre>	<pre>// Temporary T without a reference auto create() → T; [create()];</pre>		
<pre>// rvalue ref, but has a name auto print(T&& param) → void { std::cout ≪ [param]; }</pre>	<pre>// Transformation into rvalue with move T value{}; T o = [std::move(value)];</pre>		
// References have an address T& create(); [create()];	// xvalue, binds to rvalue references T&& create(); [create()];		
<pre>// String literals are always lvalues std::cout « ["Hello"];</pre>	Rule of thumb: Does element keep living? Vivalue (only copy), × rvalue (copy & move possible)		
// depends on the	implementation of +		

T value(): std: cout or [value + 1]:

1.3. VALUE CATEGORIES



A value is always either a Ivalue, xvalue or prvalue. Lvalue does not always mean "on the left side of an

Has Identitty?	Can be moved from?	Value category
Yes	No	Ivalue
Yes	Yes	xvalue (expiring solve)
No	No (Since C++17)	prvalue (pure notive)
No	Yes (Since C++17)	- (daes not exist anymane)

Yealue: Expiring Value, Address cannot be taken, cannot be used as left-hand operator of built-in assignment Conversion from prvalue through temporary materialization. Conversion from Ivalue through std::move(
Examples: Function call returning rvalue ref (i.e. ntd::move(x)), access of non-ref members of rvalue object: X x1fb, x2fb: consume(std::move(x1)): std::move(x2).member: Xfb.member:

1.3.1. Temporary Materialization

Getting from something imaginary to something you can point to (o volve getting am address). Transformation from proble to snable. Requires a destructor. Happens...

- when binding o reference to a proble pt.). when accessing a member of a prvalue (2),
when accessing an element of a prvalue array (int value = (int[])(10, 20)[1]; // value = 20],

- when converting a prvalue to a pointer (coset int &Gref = (5+3); coset int *ptr = &ref;)
 when initializing an std:: initializer_list<T> from a braced-init-list. (etd:: vector(1, 2, 3);)

const Ivalue reference: Binds Ivalues, xvalues and prvalues: auto f(T const &) → void CPLA | FS25 | Nina Grässli, Jannis Tschan, Matteo Gmür

- struct 6host { auto haunt() const \rightarrow void { std::cout \ll "boos!\n"; } // ~6host() = delete; would error
- ; auto evoke() \rightarrow 6host { return 6host{}; } auto main() \rightarrow int { 8host&sam = evoke(); /* (1) */ 6host{}.haunt(); /* (2) */ }

1.4. OVERLOAD RESOLUTION FOR ERFF FUNCTIONS

	f(s)	f(S &)	f(S const &)	f(S &&)
s{}; f(s);	>	✓ (preferred over const5)	~	×
const s{}; f(s);	×	×	~	×
[S{});	~	×	~	✓ (preferred over const
s{}; f(std::move(s));	~	×	~	✓ (preferred over const

The overload for value parameters imposes ambiguities. For deciding two Ivalue reference overloads, the const-ness of the argument is considered

5.	OVERLOAD RESOLUTION FOR MEMBER FUNCTIONS					
		S::m()	S::m()	const	S::eO &	

1.5. OVERLOAD RESOLUTION FOR MEMBER FUNCTIONS						
	S::m()	S::m() const	S::m() &	S::m() const &	S::m() &&	
	Value Members			Reference Members		
S s{}; f(s);	٧.	~	✓ (preferred over const&)	~	×	
S const s{}; s.=();	×	~	×	~	×	
\${}.=();	~	~	×	~	✓ (preferred over const&)	
S s{}; std::move(s).m();	~	~	×	~	✓ (preferred over const&)	

Reference and non-reference overloads cannot be mixed. The reference qualifier affects the this object and the overload resolution. const 86 is theoretically possible, but an artificial case.

1.6. SPECIAL MEMBER FUNCTIONS

- POSInference: The rule of three file faces
 Constructors: Default Constructor, Copy Constructor, Move Constructor (called an variable is
 Assignment Operators: Copy Assignment, Move Assignment (called on variable reassignment)
 Destructors: (Called automatically when variable goes out of scape to clean up the objects resources)

Assignment operators must be member functions. Move operations must not throw exceptions, thus aren't allowed to allocate memory. Use the default implementation of the special members whenever possible.

1 6 1 Moun Constructor 9 (9 68)

CPDReference: Move constructor

Takes the guts out of the argument and moves them to the constructed object. Leaves the argument in volid but indeterminate state. Don't use the argument after it has been moved from until you assign it are value

Default behavior: Initialize base classes and members with move initialization. (5 & v) := (set::move(v, w))

1.6.2. Conv/Move assignment auto operators (\$ const. \$) -> \$6 / auto operators (\$ \$6) -> \$6 CPDIsference: Copy, majornest, CPDIsference: Move majornest
Copies/Moves the argument into the this object. Executed when the variable to copy/move to has already
been nitriolized-Must be a member operator. Default behavior: Initializes base classes and members with

1.6.3. Destructor -S()

CPORference: Destructors

Deallocates resources held by the this object. No parameters. Must not throw (to seezespt). Must be implemented if a custom copy/move is defined. Default behavior: Calls destructor of base classes and members.

1.6.4. Copy-Swap-Idiom

Allows efficient implementation of the copy-assignment operator. Utilizes the copy constructor to create a temporary object, and exchanges its contents with itself using a non-throwing swap. Therefore, it swap old data with new data. The temporary object is then destructed automatically. It needs a copy-constructor a destructor and a swap function to work. Should be marked neexcept.

voing std::swap; // Fall back to std::swap if no user defined swap is available swap(member1, other.member1); // Calls std::swap or custom swap of member variable

// Move assignment
auto operator=(S&& s) → S& { // could be neexcept
 if (std::addressof(s) → this) {

} :: // Destructor is implicit

auto swap(Sū lhs, Sū rhs) noexcept → void { lhs.swap(rhs); // Calls S::swap() above

write/get	default ctor	destructor	copy ctor	copy assign	move ctor	move assign
nothing	defaulted	defaulted	defaulted	defaulted	defaulted	defaulted
any ctor	undeclared	defaulted	defaulted	defaulted	defaulted	defaulted
default ctor	user decl.	defaulted	defaulted	defaulted	defaulted	defaulted
destructor	defaulted	user decl.	defaulted!	defaulted!	undeclared	undeclared
copy ctor	undeclared	defaulted	user decl.	defaulted!	undeclared	undeclared
copy assign	defaulted	defaulted	defaulted!	user decl.	undeclared	undeclared
move ctor	undeclared	defaulted	deleted	deleted	user decl.	undeclared
move assign	defaulted	defaulted	deleted	deleted	undeclared	user decl.

1.7. COPY FUSION

some cases the compiler is required to elide (see) specific conv/move operations regardless of the side effects of the corresponding special member functions ("Mand tory elision"). The omitted operation

not exist. This happens...

— In initialization, when the initializer is a prvalue: S s = S{S{}}; (Only 2 constructor cell, 0 copy open When a function call returns a prvalue: (s) gets initialized directly in new_ne() instead of a auto_create() \rightarrow S { return S(}; } auto_main() \rightarrow int { S new_swfcreate()}; S * sp = new S{create()}; }

1.7.1. Further optional elisions

MANON found from two Copinions Return type is a value type, return expression is a local variable of the return type. The object is constructed in the location of the return value. The construction must rail? of the return type. The object is constructed in the location of the return value. The construction must rail? It is constructed in the return of the return type and type

1.7.2. Example

S { S s{}; std::cout « "\t --- create() ---\n"; return s; } std::cout e "\t -- S s{create()} ---\n"; S s{create()};
< "\t -- s = create() ---\n"; s = create();</pre>

x Move in z (), 1x Move in z = 2x Move in z (), 1x Move in z =		Ox Move in x (), Ox Move in x =		
S s{create()}	S s{create()}	S s{create()}		
Constructor S()	Constructor S()	Constructor S()		
create()	create()	create()		
Constructor S(S&&)	Constructor S(SSA)	// no return to create()		
Constructor S(S&&)	// no create() to s	// no create() to s		
s = create()	s = create()	s = create()		
Constructor S() create()	Constructor S()	Constructor S() create()		
Constructor S(S&&)	Constructor S(SSA)	// no return to create()		
operator =(S&&)	operator =(S&&)	operator =(S&&)		

1.8. LIFE-TIME EXTENSION

time ends at the end of the block. It is not transitive (Reference Return → Dangling Referen

struct Demon { /* ... */ }; suto summon() \rightarrow Demon { return Demon{}; } // Creates a demon suto countEyes(Demon const&) \rightarrow void { /* ... */ }

acto main() → int {
 summe();
 where ();
 // Deson disc at the end of the statement
 count(sec(Deson()); // Deson lives long enough for count, eyes to finish
 // Life-time can be extended by count S → Timpdom Lives until end of block
 // Life-time can also be extended by SS → Lizzoik Lives until end of block
 Deson(); Lizzoik = County ();
 // Life-time can also be extended by SS → Lizzoik Lives until end of block
 Deson(); Lizzoik = County ();
 // Rappion and Lizzoik dis here

2 TYPE DEDUCTION

2.1. FORWARDING REFERENCES AND TYPE DEDUCTION

2.1. TO ANNAHOUND IN CITE CHEMICAL SHAPE TO EXCOLATION TO THE WAY THE ANNAHOLD THE

template <typename T>
auto f(ParamType param) → void; // T and ParamType can be different types!
f(<expr>);

eduction of type. T depends on the structure of the type of the corresponding parameter PanamType: 2.1.1. Paramtype is a value type (T)

2.1.1. Paramtype is a value type (1)

(e.g. with / [f peans) → vat(f) Natherint (f *) are also value types

1. «sayer» is a reference type: (ignore the reference

2. Ignore the rightmost const of «sayer» (char canst * const → char canst *)

3. Pattern match «say» 's type against Paramtype to figure out f int x = 23; int const ex = x; int const & erx = x; char const + const ptr

f(x); auto f(int param) → void;
f(xx); auto f(int param) → void;
f(cxx); auto f(int param) → void;
f(crx); auto f(int param) → void;
f(ptr); auto f(char const * param)

2.1.2. Parantype is a reference (T&)

int x = 23; int const ex = x; int const & crx = x;

Examples for References: auto f(T & paren) -> void auto f(int& param) → void; auto f(int const& param) → void; auto f(int const& param) → void; auto f(int const& param) → void; Examples for Const References: outo f(T const &

2.1.3. Parantype is a forwarding reference (T&&)

1 severy is an Ivalue: T and ParanTyne become Ivalue reference int x = 23: int const ex = x: int const & erx = x:

// calls // instances f(x); auto f(int & param) → void; (ex); auto f(int const & param) → void; (erx); auto f(int const & param) → void; (27); auto f(int & param) → void; (*0SI*); auto f(char const (&) [4] param) → void;

When an initializer_list is used for type deduction, an error occurs: f({23}) For this to work, a separate template is needed

template <typename T>
auto f(std::initializer_list<T> param) -> void;
f(23): // T = int. ParamType = std::initialize 2.1.5. Type Deduction for auto

CPPReference: Placeholder type specifiers: auto/decl.tupe (auto Same deduction as above, auto takes the place of T:

autošá uref1 = x; autošá uref2 = cx; autošá uref3 = 23: // T&&: x = int (lvalue) → uref1 = int& // T&&: cx = int const (lvalue) → uref2 = int const& // T&&: 23 = inter (numlum) → unof1 = inter // Special cases auto init_listi = {23}; // std::initializer_listcint> auto init_list2{23}; // int, was std_initializer_listcin auto init_list3{23, 23}; // Error, requires one single argum

aute Return Type Deduction: auto can be used as return type and for parameter declarations. Body must be available to deduce the type. Multiple auto parameters are considered different types.

2.1.6. Type Deduction for decltype

description p =

2.1.7. Type Deduction in Lambdas

2.2. PERFECT FORWARDING WITH STD::FORWARD

template <typename T>
decltype(auto) forward(std::remove_reference_t<T>& param) {
 return static castsT&As(maram);

std:: forward is a conditional cost to an evalue reference. This allows arguments to be treated as what they originally were: halues remain (values and realues remain realues. If T is of value type: T && is an evalue reference in the return expression. (int → inter.)

If T is of Ivalue reference type, the resulting type is an i (e.g. T = int 4 → T 45 would mean "int 4 44" which can be col

complatectypename T>
tatic auto make_buffer(T && value) → BoundedBuffercvalue_type> {
BoundedBuffercvalue_type> new_buffer{};
new_buffer.jousi(ztdi::coreard<t>(value));
return new_buffer.joufie(ztdi::coreard<t>(value));

2.3. STD::MOVE

Does not actually move objects. It's just a unconditional cost to an evalue reference. This allows resolution of ryalue reference overloads (suto ((TGS ±)) and move-constructor/-assignment operators. Coution! Moving

To allow for moving, there should be no const member in a class. Reference Collapsing: "T& &", "T& &&" and "T&& &" become "T&", "T&& &&" becomes "T&&"

3. HEAP MEMORY MANAGEMENT

Lifetime on Stack: Deterministic, local variables get deleted automatically upon leaving their sco Lifetime on Heap: Creation and deletion happens: explicitly with new and delete (Insepresus, our auto -600) $\sim void$ { auto 1p = new fint(5); /* ... * */ delete 1p).Rules: Delete every object you allocated, do not delete an object twice or access a deleted object.

3.1. EXPLICIT LIFE-TIME MANAGEMENT

Global and local variables have life-time implicitly managed by the program flow. Some resources can be allocated and deallocated explicitly. This is error-prone. Guideline: Always wrap explicit resource management

3.2. POINTER SYNTAX

Heap Array Access: aute arr = new int[5]{}; // 5 needs to be compile-time-constant, auto → int
int v = arr[4]; // accessing element (warr + 4 pointer objects)

Direct Member Access (->)

ar() -> void { this--value = ...; } int value; // 'this': pointer to a instan

Pointer Parameters: Pointers can be used as parameters. Addresses can be taken with & ico.

suto foo(int* p) \rightarrow void { } suto bar() \rightarrow void { int* in = new int{5}; int local = 6; foo(in); foo(&local); }

Const Politers: const Politer can't be undiffed, but the object behind it may, const is on the right side of the . The declaration is read from right to left.

Int const + const + const is const politer to a power before the const politer to a const politer to a const politer to a const politer.

nullptr: Represents a null-Pointer. Is a *literal* (prusiw) and has type nullptr_t. Implicit conversion to any pointer type: T *. Prefer nullptr over 8 and NULL (no oversion emblasts, no implicit conversion).

Pointer vs. Reference are always bound to an object cannot be rebound to another obj can be changed (f not const) require dereferencing with "+" or "-> allow member access by "."

Use raw pointers only to explicitly model th roid (if (etr.) (...))) and for modelling borrowing only. Else, use smart pol

can be dangling (referencing

3.3. MEMORY ALLOCATION WITH NEW

new <type> <initializer>

Allocates memory for an instance of <type>. Returns a pointer to the object or array created on the heap of type <type> *. The arguments in the <intializer> are passed to the constructor of <type>. Memory Leak # not removed with delete. Anold manual allocation, use RAII intead. struct Point{ Point(int x, int y) : x{x}, y{y}{} int x, y; }; auto createPoint(int x, int y) → Point* { return new Point{x, y}; // constructor

auto createCorners(int x, int y) → Point* {
 return new Point[2]{{8, 8}, {x, y}};

3.3.1. Placement new

new (<location>) <type> <initializer> Used for placing elements on the heap in the location of a deleted Used for placing elements on the heap in the location of a deleted element. Does not allocate new memony (hasi-phase heaves)! The memory of <locations needs to be suitable for construction of a new object and any element there must be destroyed before. Calls the Constructor for creating the object at the given location and returns the memory location. Better use std::construct_at(). new (ptr) Point{7, 6}; delete ptr;

3.4. MEMORY DEALLOCATION WITH DELETE

delete «pointer: Deallocates the memory of a single object pointed to by the speinter>. Calls the Destructor of the destroyed type, delete nullptr does nothing. Deleting the same object twice is Undefined Behavior!

struct Point{ Point(int x, int y) : $x\{x\}$, $y\{y\}$ {} int x, y; }; auto funKithPoint(int x, int y) \rightarrow void { Point + pp = new Point{x,y}; delse pp; // calls destructor and releases memory

3.4.1. Placement delete

Does not exist, but a destructor can be called explicitly.

S * ptr = ...; ptr->-S(); Destroys the object, but does not free its memory. Called like any other member function. Better use

delete[] cointer-to-array; Deallocates the memory of an array pointed to by the <pointer-to-array>. Calls the Destructor of the

3.5. NON-DEFAULT CONSTRUCTIBLE TYPES
A type is non-default-constructible when there is *no explicit or implicit default constructo*To create arrays of NDC types, allocate the plain memory and initialize it later (skin page 26 auto mesory = std::sake_weique-cost;ericle[]ciaze@ficialty = 2 /* Array size */); auto location = reinterprat_castr@size.dischesory.ort(); // Address of first element std::construct_size.disches(), 2 // Equivalent to arr@() = Foint(); auto value = elementAt(sembor, ort(), 0); // Access value via helper function std::destruy_size.disches(), 2 // Access value via helper function

auto elementAt(std::byte+ memory, size_t index) \rightarrow Point& { // helper function

Don't use an element if it is uninitialized and destroy them before the memory is deallocate Use a std::byte array as memory for NDC Elements.

- Static:std::arraycstd::byte, no_of_bytes> values_memory; (on stack, size known of cor

Dynamic: std::unique_ptr<std::byte[]> values_memory; (on heap, size known at run-tim 3.6. CLASS-SPECIFIC OVERLOADING OF OPERATOR NEW/DELETE Overloading new and delate for a class can inhibit heap allocation. This can be used to provide efficient allocation, is useful with a memory pool for small instances or if thread-local pools are used. Can log or limit

jumber of heap-allocated instances. But in general, no truct not_on_heap { // Prevents heap sllocation of this class static auto operator new(statislize, t.g.) \rightarrow void * { three std:bad_slloc{};} static auto operator new(std:size, t.g.) \rightarrow void * { three std:bad_slloc{};} static auto operator new[(std:size, t.g.) \rightarrow void = { three std:bad_slloc{};} static auto operator deltar(void *ytr) \rightarrow void nexcept { /* do nothing \neq } static auto operator deltar[(void *ytr) \rightarrow void nexcept { /* do nothing \neq } }

3.7. READING DECLARATIONS Declarations are read starting by the declarator (name) First read to the right until a closing parenthesis is encountered Second read to the left until an opening parenthesis is encounte Third jump out of the parentheses and start over

Specifiers right to the declarator: Array Declarator ((1) and Function Parameter List ((sparameter

Specifies left to the declarator: References (ss, sl, Pointers (*) and Types (int) const: Applies to its left neighbor, if there is no left neighbor, it applies to its right neighbor. onst int 1; int const 1; Should always be written to the right of the type to a oid (* f)(int &. double)

tion that taker a reference to (ot and a double returning uply nt const * (* f [2][3]) [5]; ents of arrays of 3 elements of pointers to arrays of 5 elements of pointers to const

3.8. RESOURCE MANAGEMENT WITH RAII

Resource Acquisition is Initialization is an alternative to allocating and deallocating a reso Wraps allocation and deallocation in a class, uses regular constructor/destructor. Cleaned in a 3.8.1. std::unique_ptr and std::make_unique

CPDM/menor_itd_value_qtr_CPDM/menor_itd_male_unipue std::unique_ptr<char> cPtr = std::unique_ptr<char>('*'); Wraps a plain pointer, haz zero runtime overhead. A custom deleter could be supplied if required. Always use make unique for creation. Can create unbound arrays, but not fixed size array 3.8.2. Container Member Function emplace

or std:: vector<T.Allocator>::emplace.CPPNeference: std:: stack<T.Contains Constructs elements directly in a container, more efficient than moving them. Not available for std:: array.
std::stackPoint> vec{}; vec.emplace(3, 5); // std::vector requires position arguent

4. ITERATORS & TAGS 4.1. TAGS FOR DISPATCHING

ruct SpaceBriveTag{}; mplate<typename> struct SpaceShipTraits { using Drive = SpaceBriveTag; } truct SubspaceOriveTag : SpaceOriveTag{};
suplate<> struct SpaceShipTraits<BalaxyClassShip> { using Drive = SubspaceOriveTag; }; template <typename Spaceshipo auto travelToDispatched(Salaxy destination, Spaceship& ship, SpaceOriveTag) → void {

Late <typename Spaceship> travelTo(Galaxy destination, Spaceship\$ ship) → void { typename SpaceShipTraits<Spaceship>::Drive drive(); // get the Spaceship's drive tag travelToDispatched(destination, ship, drive): // call overloaded function

4.2. ITERATORS Different algorithms require different strengths of iterators. Iterators capabilities can be determined at

compile time with tag types.

Outputterator: Write results (to console, file etc.), without specifying an end (used on std::ostress)

Output distraction: With results in mouse, for mile, without specifying an end part on end-restroad, operated - returns mile black reference for subgenment of the value. Imput framework in the subgenment of the value of mile statement, imput framework reduced produced by the subgenment of the value of mile statement, postable - returns one follower depress, or reduce, Formward transferred. Easily first sequences, making pass jumps as statement, List shade fall, considerations and produced by the subgenment of making from control operator - returns leadure the subgenment of t

You need to implement the members required by your Stonaton too

struct Intiturator 4 // Provide these member types to align with STL iter using iterator_category = std:input_iterator_tag; // iterator_category = std:input_iterator_tag; // iterator_category = outling with _type = int; // type of elament to iterator iterators over using pointer = int = int // pointer type of the elements iterated over using pointer = int = int // pointer type of the elements iterated over using reference = int = int // reference type of the elaments iterated over

4.2.1. iterator traitso

CPPReference: etd:::iterator: traits STL algorithms often want to dete STL algorithms often want to determine the type of some specific thing related to an iterator. However, not all iterator types are actually classes. Default iterator_traits just pick the type aliases from those provided.

Specialization of iterator_traits also allows "naked pointers" to be used as iterators in algorithms template stynename. The struct iterator traits (Tee 4 /* Provide requiar ite

4.2.2. Problems with the Stream Input Iterator

Reference Member and Default Constructor When implementing a input iterator, we need to be able to create an E0F iterator. This dirty hack works, but the global variable to initialize the reference in an anonymous namespace is bad for multi-threading.

amespace { // global variable to initialize the reference in an empty namespace std::istringstream empty{}; // pseudo default |
IntInputter::IntInputter() : input { empty }
// guarantee the empty stream is !good()
input.clear(std::ios_base::eofbit); // mark empty stream as EOF

. Dereferencing and Equality

// Bereferencing with * reads the value from the input
wto IntInputter::operator*() → IntInputter::walue_type {
 value_type value;
 return value;
 fsteps tire arm:

Stream iter comparisons only make sense for testing if they can still be 1 intiputter::operator=(const IntImputter & other) const \rightarrow bool { storm !input.good() & | other.input.good() }

4.2.3. Custom Iterator Example (Testat 2)

utingvalue_type = BoundedBuffer::value_type;
uting value_type = BoundedBuffer::refreence;
uting oreference = BoundedBuffer::refreence;
uting const_reference = BoundedBuffer::const_reference;
uting citze_type = BoundedBuffer::tize_type;
uting difference_type = std::ptndiff_t;
uting iterator_category = std::rendem_access_iterator_tag;

auto operator==(iterator_base const & other) const -> bool { ... }
auto operator<=>(iterator_base const & other) const -> std::strong_ordering { ... }

auto operator*() const -> decltype(auto) { return Buffer->elementAt(Index); } auto operator->() const -> decltype(auto) { return &(Buffer->elementAt(Index)); } auto operator[](difference.type index) const -> decltype(auto) { ... } auto operator++() -> iterator_base & { Index++; return *this; }
auto operator++(int) -> iterator_base {
 auto const copy = *this;

auto operator--() -> iterator_base & { Index--; return *this; }
auto operator--(int) -> iterator_base {
 auto const copy = *this;

auto operator+(difference_type n) const -> iterator_base { auto copy = *thi copy += n; return copy; } auto operator-(difference_type n) const -> iterator_base {
 return this->operator+(-n); }

outo operator==(difference_type n) -> iterator_base & { Index += n; return *this; }
operator==(difference type n) -> iterator base & { return this->operator==(-n) auto operator-(iterator_base const & other) const -> difference_type { ... }
private: difference_type Index{}; BoundedBuffer* Buffer;

Boost would generate operator ++ (int), operator -- (int), operator + (diffe

using iterator = iterator_base<Container<T>>; using const_iterator = iterator_base<Container<T> co

difference_type n) with implementation shown above. Change signature to struct it boost::operators inpl::random access iterator helpersiterator base<*V. V> f}

5.1.1. Type traits

5. ADVANCED TEMPLATES Pros of static polymorphism

 Longer compile-times
 Template code has to be known when used (innext to be in IRP file)
 Larger binary size (copy of the used parts for each template instance) easier to optimize)
Type checks at compile-time A polymorphic call of a virtual function (whentance ownloading) requires lookup of the target function. Non-virtual calls (rempiar ownloading) directly call the target function. This is more efficient.

5.1. SFINAE (SUBSTITUTION FAILURE IS NOT AN ERROR)

CONFIGURATE ADMA*

Lived to eliminate overload candidates by substituting return type and parameters. During overload reco-lution the template parameters in a template declaration (u, v) are substituted with the deduced types (u + u). Lived the substituted with the deduced types (u + u). Lived the substitution of template parameter fails, that overload candidate is discurred.

decltype (<return-expr>) as return type checks if the overload candidate would work, but this approach is infeasible, because functions can be void and it's not elegant for complex bodie

CPSErference_Disc Tools (PSSerference indicates of content

#include <type_traits>
The standard library provides many predefined checks for type traits. A trait contains a boolean value. Usually available in a v (retur be boot onside and non- v variant (serves the interval constant) Example: std::is_class CPPErformer std::s class
std::is_class
:std::is_class_v<5>; // both true
std::is_class<int>::value; std::is_class_v<int>; // both fals

std::enable if/std::enable if t <u>Department: Rest enable_1f</u> t std::enable_1f_t takes an expression and a type. If the expression evaluates to true, std::enable_1f_t represents the given type, otherwise, it does not represent a type.

to main() → int {
std::emable_if_t<true, int> i; // int
std::emable_if_t<false, int> error; } // no type, compiler error std::enable_1f can be applied at different places (marked with "||", only one needs to be used)

template <typename T, [typename = enable_if_t<is_class_v<T>, void>]> auto increment([enable_if_t<is_class_v<T>, void>] value) // impairs t \rightarrow [enable_if_t<is_class_v<T>, T>] {

5.2. TEMPLATE PARAMETER CONSTRAINTS AND CONCEPTS

<u>CPFBeformer: Contraints and Concepts</u>

Provide a means to specify the characteristics of a type in template context. Better error messages, more

5.2.1. Keyword requires

ining template parameters, requires is followed by a compile-time constant boolean expres sion. Is either placed after the template parameter list or after the function template's declarato

template <typename T>
requires std::is_class_v<T> // either here...
auto function(T argument) -> void requires std::is_class_v<T> /* or here */ { ... }

requires Expression requires also starts an expression that evaluates to bool depending whether they can be compiled.

wires (\$parameter-list\$) { /* sequence of requirements */ } template <typename T>
requires requires (T comst v) { v.increment(); } // compiles if v has a increment function
auto increment(T value) \rightarrow T { return value.increment(); }

Type Requirements Check whether a type exists. Starts with typename keyword. Useful for nested types like in Rounded Ruffe

requires { typename StypeS }

requires { typename BoundedBuffer<int>::value type: }

Compound Requirements

Compound arquitements

Check whether an expression is valid and can check constraints on the expression's type. The return-type requirement is optional. Needs to be a valid type constraint, regular types can't be used.

requires (T v) { { \$expression\$ } → \$type-constraint\$; }

template <typename T^* // We can't use T as return type in requires, it is not a constrain requires $\{T$ const $v\}$ { v.increment() } \to std::same_as< T^* ; } auto-increment(T value) $\to T$ T error value increment(T).

template <typename T>
concept Incrementable = requires (T const v) { { v.increment() } → std::same_as<T> }; Named requirements can then be used in template parameter declarations or as part of a re

complate <Incrementable T> // either here...
requires Incrementable<T> // ...or here
nuto increment(T value) → T { return value.in:

auto can be used as parameter type instead of a template d

mplate <fnorementable T>
to increment(); }
its equivalent to Terse Syntax
to increment() ratue) → T { return value.increment(); }
its equivalent to Terse Syntax
to increment[Incrementable auto value) → T { return value.inc

CPPReference: Concept Abrox

The STL has predefined constraints: std::equality_comparable (can type be ==/== com

A. COMPTLE-TIME COMPUTATION

6.1. CONSTANT EXPRESSION CONTEXT hese expressions always need to be defined at compile-time.

Non-type template arguments (std::arraysElement, 5> arrs();)

constinit variables are non-const. They need to be initialized at compile-time, but can be changed at

ve local variables of literal type. The variables must be initialized before usage

Can use loops, recursion, arrays, references, branches

Can contain branches that rely on run-time only features, if branch is not executed during compile-time.

— Are usable in constexps and non-constexpr contexts (during runtime)

— Can allocate dynamic memory that is cleaned up by the end of the compilation (uncor C++20)

— Can be virtual member functions (uncor C++20)

constexpr auto factorial(unsigned n) \rightarrow unsigned { /* needs to have a body */ }

consteval auto factorial(unsigned n) { return result:

during compile-time. Instead, there will be a compilation error. If const

CPPReference: Named Requirements - Literatives

Literal types are built-in scalar types (like int, double, pointers, enum), Structs with some restrictions (must have matesper destructor and countexper / counterval constructor), Lambdas, References, Arrays of literal types and void e functions with side effects on literals)

6.3.2. Compile-Time Computation with Variable Templates

6.3.3. Captures as Literal Types

Capture types (the types returned by lambds expressions) are literal types as well. The can be used as types of consteapy variables and in consteapy functions.

5.2.2. Keyword concept

Specifies a named type requirement. Conjunctions (&&) and disjunctions (|||) can be used to combine constraints (&e. requires std::integral(-1) ||| std::flueting_paint(-1)|.

5.2.3. Abbreviated Function Templates

If there are two auto arguments, two template typenames T1, T2 get created.

5.2.4. STL Concepts

Array bounds (doubte matrix (ROWS) (cots) ():)

6.2. CONSTEXPR / CONSTINIT

6.3. CONSTEXPR / CONSTINIT

6.4. CONSTEXPR / CONSTINIT

6.5. CONSTEXPR / CONSTINIT

6.6. CONSTEXPR / CONSTINIT

6.7. CONSTEXPR / CONSTINIT

6.8. CONSTEXPR / CONSTINIT

6.9. CONSTINIT

6.9 CONSTINIT

6.9

insterne / const (a) t Variables are evaluated at compile-time. They are initialized by a constant expression

consteapy remiscrate variations are evaluated at compine-time. They are initiatized by a constant is and require a literal type (primitive data type whost knop allocation). They can be used in constant is contexts. Possible contexts are local scope, namespace scope and static data members. — consteapy variables are const, read only at run-time

run-time 6.2.1. constexpr Functions

6.2.2. consteval Functions Are usable in constever contexts only (only be collectuated or only and implicitly const

estavos auto fastopialOFE - fastopial(E): // works

signed input{};
(std::cin >> input) {
std::cout << factorial(input); // error, function cannot be used at</pre>

6.2 LITERAL TYPES

Literal Types can be used in constexpr functions, but only constexpr member functions

white: $\begin{array}{ll} \text{obstace} & \text{vector}(T \, x, \, T \, y, \, T \, z) : \text{values}(x, \, y, \, z) \cdot \{ \} \, / / \text{ constapp constructor} \\ \text{constapp auto length}() \, \text{const} \rightarrow T \, \{ \, / / \text{ constapp const member function} \\ \text{auto squares} & \times () + x() + y() + y() + z() + z(); \\ \text{return std:oper(squares)}. \end{array}$

constexpr auto create() → Vector<double> {
 Vector<double> v(1.0, 1.0, 1.0); v.x() = 2.0; return v; }
constexpr auto v = create();

template <size_t N>
constexpr size_t factorial = factorial<N-1> * N;

- Array doubnet (mode matrix(RME) (DLE)((1) - Case 42; /* /*))
- Enumerator initializers (mose Liph (Off + 0, 0s = 2); /* /*))
- Enumerator initializers (mose Liph (Off + 0, 0s = 2); /* /*)
- Sabtic Asserbins, (metric, mose (mode mose double of the control of the c

auto main() → int {
 static_assert(factorialOfS = 120); // works
 ussigned input{};

6.2.3. Undefined Behavior

6.3.1. Literal Class Type Example
template <typename T> // can be a template
class Vector { constexpr static size_t dimensions = 3; std::arrayeT, dimensions> walkers

onstexpr auto x() → T& { return values[0]; } // constexpr nor constexpr auto x() const → T const& { return values[0]; } // implicit default destructor is also constexpr

//v.x() = 1.8; // possible if v was constinit auto v2 = create(); v2.x() = 2.8; } // v2 is a regular variable, can be modified

(for Class Template see slides page 25)

Variable templates can be constexor/constinit and defined recursively. Usage: Template-ID

template o // Base case

Page 1

operator"" _UDLSuffix()"

Allows integer, flusting point, character, and string iteration produce objects of user-defined type by defining a user-defined type. By defining a user-defined soften. The suffice most start with an undercome, it allows to seld demotion, convention, extend in probable, defined to operated functions are consistent. Only contract functions are consistent, or convention operated functions are consistent. Only comparation processing functions like and a conversation operation in such as a format of the contraction of the contraction operator is needed. Ruler part overloaded UTL operators that belong together in a separate numeropace.

IMPopur.

amesgace velocity::literals {
 onsteep faline auto operator** _kph(unsigned long long value) → Speeds(ph freturn Speeds(ph-fasfelobookle(value)); } // user defined literal operator
 // called with Speeds = 80_kph;

mplate<typename T>
ncept arithmetic = std::is_arithmetic_v<T>; // allows ints & floats in + oper

static constexpr inline auto seferoBouble(long double value) → double {
 if (value > std::numeric_limitscdooble>::max()
 | value < std::numeric_limitscdooble>::min()) {
 throw std::invalid_argument{"Value must be within double range"};
 }
}

turn static castedoubles(value):

Truct space {
constexpr explicit Speed(double value) : value{value} {};
} // conversion to double
auto constexpr explicit operator double() const { return value; } // conversion to double
auto constexpr operator - (arithmetic autor orb;) - decltype(chs) { return value + rhs }
}

ivate: druble value{};

outo constexpr operator +(arithmetic auto lhs, Speed rhs) \rightarrow decltys { return rhs + lhs; } // Make + operator commutative (value + See

Signatures: unsigned long long for integral constants, long double for floating point constants, (char const *. size t len) → std::string for string literals, char const * for a raw UDL operator //

6.4.1. Template UDL Operator

template <char...>
auto operator""_suffix() → TYPE

Has an empty parameter list and a variadic template parameter. Characters of the literal are templa arguments. Often used for interpreting individual characters. Since C++20, the template UDL operator work with string literals as well (example and compile time steps in slides on page 42 - 44).

 $\textbf{Standard Literal Suffixes: Do not have leading underscore: std::string } (\mu), \ \text{std::complex } (\theta, \theta, \theta), \ \text{where} (\theta, \theta, \theta) = 0$

6.5. PREPROCESSOR hello.cop → preprocessor → hello.i → compiler → hello.o → linker → hello Object-like Macros

#define identifier replacement-list new-line

Identifier is a unique name, by convention in ALL_CAPS. Is valid until #undef NAME. Replacement-list is a

#define identifier (identifier-list?, ?...?) replacement-list new-line Features an optional parameter list, containing only names. Params with a #-prefix turn into string literals

Includes: Textual inclusion of another file. #include "outh" for including a header file from the same project or workspace, #include <peth> for external includes.

Conditional includes: Enable a section depending on a condition. (example and macros in skdes on page 52 - 62).

#ifdef identifer new-line

#if constant-expression new-line #elif constant-expression new-line #else new-line #endif new-line

7. MILLTT-THREADING & MILTEX

itd:: thread to explicitly run 7.1. API OF STD :: THREAD

auto main() → int { std::thread prector { [] { /*lambdax/ } }: prector.ioin(): }

A new thread is created and started automatically, Creates a new execution context, join() walts for the thread to finish. Besides lombels, functions or functor objects can also be executed in a thread. The return volue of the function is ignored. Threads are default-constructible and moveable. Courtion: Program terminates if thread gets destructed without calling join() before!

struct Functor { auto operator()() comst \rightarrow void { std::cout \ll "Functor" \ll std::endl; } to function() → void { std::cout « "Function" « std::endl; }

std::thread functionThread(function); std::thread functorThread(Functor{}); functorThread.join(); functionThread.join(); }

Streams: Using global streams does not create data races, but sequencing of characters could be mixed. std::this_thread helpers: get_id() (An ID of the underlying OS thread), sleep_for(durati

7.1.1. Passing arguments to a std::thread

template<class Function, class... Args> explicit thread(Function&& f, Args&&...args);

The std:: thread constructor takes a function/functor/lambda and arguments to forward. You should pass all arguments by value to avoid data races and dangling references. Capturing by reference in lambdas creates shared data as well (fusu how to use them, don't derin from an autobital

auto fibonacci(std::size_t n) \rightarrow std::size_t { /* ... */ } auto printFib(std::size_t n) \rightarrow void { auto fib = fibonacci(a); /* print... */ } auto main() \rightarrow int { std::thread function { printFib, 46 }; function.join(); }

Before the std::thread object is destroyed, you must join() (well until finished) or det thread and run in the background) the thread, otherwise you get a runtime error.

CPPReference: std:: ithread RAII wrapper that automatically calls: 1 CR.t.request_stop() sends the request, with

7.2 INTER-THREAD COMMUNICATION

CPPReference std::nutex CPPReference std::shared nutex Communication happens with mutable shared state

a Broblem: Data Pace Colution: Locking the charge

All mutexes provide the following operations - Acquire: lock() - blocking, try_lock() - non-blocking - Release: unlock() - non-blocking

Two properties specify the capabilities:

Recursive: Allow multiple nested acquire operations of the same thread (prevents self-decolock)

Timed: Also provide timed acquire operations (try_teok_pter(devotion), try_took_wetil(time))

Reading operations don't need exclusive access. Only concurrent writes need exclusive locking. Use

7.3.1. Acquiring / Releasing Mutexes

CPRyformors atd:liess, quark, CPRHyformors atd:scaped, lace, CPRHyformors atd:sunious_lock

CPRHyformors atd:sless, quark, CPRHyformors atd:scaped, lace, CPRHyformors atd:sunious_lock

Usually you use a lock that manages the mutex: ock_guard: RAll wrapper for a single mutex. Locks immediately when cons

destructed. a.
 ped_Lock: RAII wrapper for multiple mutexes. Locks immediately when constructed, unlocks

Standard Containers and Concurrency: There is no thread-safety wrapper for standard containers. Access to different individual elements from different threads is not a data race. Almost all other concurrent uses of containers are deapnersus. Stand-ay for copies to the same object can be used from different threads, but accessing the object itself can race if non-const (reference ensering atoms).

7.3.2. Thread-safe Guard Example [Testat 3]

Scoped Lock Pottern: Create a lock guard that (un)locks the mutex automatically. Every member function

is mutually exclusive because of scoped locking pattern. Strategized Lock Pottern: Template Parameter for

using guard = std::lock_guard<MUTEX>; using lock = std::unique_lock<MUTEX>; template_stynepuse_EngerTX

/ wait requires timed locking therefore unique_lock (lk, [this] return [a.empty():]): // checked once. no busy wai:

uto try_pop(T & t) \rightarrow bool {
quard lk/mx}; if (q.empty()) { return false; } t = q.front(); q.pop(); return true; / call container empty, not this→empty, would cause deadlock uto empty() const → bool { gward lk{mx}; return q.empty(); }

std::queue<T> q{}; mutable MUTEX mm{}: // mutable to unlock in const member functions ONDITION not_empty();

iable don't need to be swap()-ed. But notif 7.4. RETURNING RESULTS FROM THREADS
We can use shared state to "return" result, read the result. We cannot remainly a result result. We cannot remainly a result.

}): // Mutex is unlocked when guard goes out of scope std::Mis_thread::Risep,for(is); auto lock = std::widque_lock(motex); // Lock the mutex finished.wait(lock); // Release mutex, wait until thread unlocks mur std::cout « The ansers is: " « shared « "h"; thread.join(); } 7.4.1. std ··· future

PPReference: etd.: future future represents result

mait(). urocus umun available, mait_for(<timeout>): blocks until available or timeout elapsed, mait_until(<time>): blocks until available or the timepoint has been reached

Their destructor may wait for the result to become available

7.4.2. std::nrowise A resolver resolves the host names to endopints, as (a : connect () tries to establish a connection #include
Promises are the origin of futures. They allow us to obtain a future using get_future() and pub asio::ip::tcp::resolver resolver{context}; auto endpoints = resolver.resolve(domain, "88"); asio::connect(socket, endpoints)

auto main() -> int {

: {
 std::chrono_literals;
t> promise{}; auto result = promise.get_future(); Sing Ammapace Non-commandation of the promise (set_future(); uto thread = std::thread; [&]{std::this_thread::sleep_for(2s); promise.set_value(42); }

auto main() → int {
 auto the_amswer = std::async([] { /* calculate a while ... */ return 42; });
 std::cout ~ "The answer is: " ~ the answer.ost() ~ "\n": }

This function returns a std::future that will store the result. get() waits for the result to be available static static can take an argument of type stdr: faunth flower perigrid stdr. I stands it is sign. I stands as the stands are stdr. I stands it is sign. I stands are stdr. I stands it is sign. I stands are stdr. I stands it is sign. I stands are stdr. I stands it is sign. I stands are stdr. I stands it is sign. I stands are thread and executes it regardless of if we need the result or not, stdr. I stands it side even defers execution until the result is obtained from the star if structure (ayer variences), which comparise the result on the thread calling get(). The default policy is stdr: Launch: I sayn | stdr: Launch: I sayn | stdr: Launch: I deferred (minomener dependent).

of with this abstraction. Represents the "

of bits on the exchitecture defines, like 64-bit). Conflict: Two expression evaluations run in parallel. Both access the same Memory Location (et lea

Read/Writes in a single statement are "unsequenced": std:: cout « +1 « +1: // output u

uto outputWhenReady(std::atomic flag & flag, std::ostream & out) → void f

When creating your own atomic type with std::atomic<T>, the atomic member operations are

bool compare exchange strong(T & expected, T desired) (connet fol spuriously, but slower

Applying Memory Orders: All atomic operations take an additional argument to specify the me (type std::mmsrv_order).e.g. flag.clear(std::memory_order::seq_ost);

Sequential Consistency [seq_cst]: Global execution order of operations. Every thread obser order. This is the Default behavior. The latest modification will be available to a read operation Acquire (acquire): No reads or writes in the current thread can be recordered before this load other threads that release the some otomic are visible in the current thread. Not guarantees

Relaxed (relaxed): Does not give promises about sequencing. No data-races for atomic v

or move assignment. Object can only be accessed as a whole. No member access operator.

Release (release): No reads or writes in the current thread can be reordered after this store. All writes in

the current thread are visible in other threads that acquire the same atomic.

Acquire/Release (acq_rel): Guaranteed to work on the latest value unlike Acquire, used for Read-Modify-

can be inconsistent (parallel lead()/store()) (Lost Updates), but may be more efficient. Difficult to get right Release/Consume: Do not use! Data-dependency, hard to use. Better use acquire.

Memory Location: An object of scalar type (Arithmetic, pointer, enum, std::nullptr. These value

Data Race: The program contains two conflicting actions. Undefined Rehavior

when operations might be reordered. The Memory ordering

Sequentially-consistent (Same as code ordering and the default is

Acquire/Release (Weaker guaranters than sequentially-consistent)

Consume (Discouraged, slightly weaker than popular-release) and

Visibility of effects: (if one thread modifies a variable, under what con

// start critical section
while (flag.test_and_set() /* set flag to true and r
out << "Here is thread: " << get_id() << std::endl;
flag.clear(); // sets the flag to false
// end critical section

write (Half Fence), but ordering is consistent

8.1.2. Custom Types with std::atomic

Write operations (Full Fence) e.g. test_and_set(...).

// wa// to main() → int {
std: atomic_flag flag { };
std::thread { [Sflag] { outputNerReady(flag, std::cout); } };
outputNerReady(flag, std::cout); 't_jdin(); }
....h.nxd;:atomic=T, the atomic memb

8.1. MEMORY MODEL

std::this_thread::sleep_for(is);
std::cout « "The answer is: " « result.get() « "\n"; thread.join(); }

emplate<typename Function, typename ...Args> uto async(Function&& f, Args&... args) -> std::future</* implicitly from f */>;

Advanced Reading: asio:: read also allows to specify completion conditions.

- asio::transfer_all() (Definit behavior, transfer all date or areal buffer is full)

- asio::transfer_at_least(std::size_t bytes) fixed or least bytes number of asio::transfer_at_least(std::size_t bytes) fixed only bytes anable of asio::transfer_axeatty(std::size_t bytes) fixed easily system number of asio::transfer_axeatty(std::size_t bytes) fixed easily system number of a

asio::read_until allows to specify conditions on the data being read.

std::ostringstream request{};
request < "GET / HTTP/1.1/r\n";
request < "Host: " < domain << "\r\n"; request < "\r\n";
request < "Host: " < domain << \r\r\n"; request < "\r\n";
salo::wwifer(request.str())); // Slocks until data is sent</pre>

Simple matching of characters or strings (read until "x") or more complex matching using std :: regex
 Allows to specify a call able object (predicate → con be iterated over, returns sent data if predicate inse.

Close: shutdown() closes the read/write stream associated with the socket. The destructor cancels al

socket.shutdown(asio::ip::tep::socket::shutdown:both); // close read and write end socket.close();

9.2.2. Example: Synchronous TCP Server with ASIO Socket, Blind & Listen: An acceptor is a special socket responsible bound to a given local end point and starts listening automatically. The C++ Standard defines an abstract machine which describes how a program is executed. Platform to execute a C++ program. The abstract machine defines in what order initialization takes place and in who rder a program is executed. It defines what a thread is, what a memory location is, how threads interact

asia::io_context context{}; asia::ip::tep::endpoint localEndpoint{asia::ip::tep::v4(), port}; //uses an available IPv4 asia::ip::tep::acceptor acceptor(context, localEndpoint); Accept: accept() blocks until a client tries to establish a connection (with connect). It returns a new socket

asio::ip::tcp::endpoint peerEndpoint{}; // information about clier asio::ip::tcp::socket peerSocket = acceptor.accept(peerEndpoint);

9.2.3. Async communication more Model defines when the effect of an operation is visible to other threads and how and

8.2. VOLATILE

9. NETWORK & ASYNC

Close: Release the connection

9.2. ASIO LIBRARY

9.1. DATA SOURCES / BUFFERS

Volatile in C++ is different from volatile in Java and C#. Load and store operations of volatile variables mus

not be elided, even if the compiler cannot see any visible side-effects within the same thread. Prevents the compiler from reordering within the same thread (but the hardware might receder instructions onyway). Useful when

Interrupts are events originating from underlying system which interrupt the normal execution flow of the program. Depending on the platform, they can be suppressed. When an interrupt occurs, a previously registered function is called "never-six rever sources" calls. Should be short and must run to completion. After the interrupt was handled, execution of the program resumes.

5004

Socket: All ASIO Operations require an I/O context. Create a TCP Socket using the context. There are

Connect: If the IP address is known an endooint can be constructed easily socket connect() tries to

Write: asio::write() sends data to the peer the socket is connected to. It returns when all data is sent or

Read: asio::read() receives data sent by the peer the socket is connected to. It returns when the read-buffer is full, when an error occurred or when the stream is closed. The error code is set if a problem occurs,

constexpr size_t bufferSize = 1824; std::array-char, bufferSize> repty{}; asia::errer_code errerCode{}; auto::mercy_code errerCode{}; auto:readingsth = asia::meadisocket, asia::buffer(repty.data(), bufferSize), errerCode)

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ation buffers. The ASIO library doesn't manage memory

Data shared between an ISR and the normal program execution needs to be prote atomic modifications need to become visible, volotile helps because it suppress

nterrupts may need to be disabled temporarily to guarantee atomicity.

Sockets are an abstraction of endpoints for communication.

— TCP Sockets are reliable, stream-oriented and require a connection setup

— UDP Sockets are unreliable, datagram-oriented and do not require a connection.

std::string and std::vector.

Streambuf Buffers: asio::streambuf. Works with std::istream and std::os

9.2.1. Example: Synchronous TCP Client Connection with ASIO

asynchronous and synchronous functions to communicate with sockets

auto address = asio::ip::make_address("127.8.8.1"); auto endosint = asio::io::to::endosint(address. 88): socket.conne

Accept: Block caller until a connection request arrive

elve: Receive some data over the connection

ect: Actively attempt to establish a connection

idling multiple requests simultaneously: Using synchronous operations blocks the current thread. Ass ations allow further processing of other requests while the async operation is executed. Most 10.1.3. Member Functions that should not throw OS support asynchronous IO operations.

. The program lives an async operation on an I/O object (seeker) & passes

The I/O object delegates the operation and the callback to list so_context.

The OS performs the asynchronous operation.

The OS signois the 5 so_context that the operation has been completed.

When the program calls so_context ir un() the remaining asynchronous

for the result of the operating system).

6. Still inside the io_context::run() the completion handler is called to handle the result of the asynchro

Asymc_read_operations: asio::asymc_read_asio::asymc_read_until, asio::read_at(ord Asymc_world operations: asio::asymc_world, asio::asymc_world_at They return immediately. The operation is processed by the executor associated with the stream's associate context. A completion handler is called when the operation is done.

/ or an error occurs. Then calls the completion handler.

"uto writeCompletionNandler = [] (asio::error_code ec, std::size_tlen_of_mrite) {/*...*/}
sic::ssyme_mriteGooket, buffer, writeCompletinaN+ndmail.

9.2.5. Asynchronous Acceptor Create an accept handler that is called when an incoming connection has been established. The second parameter is the socket of the newly connected client. A session object is created on the heap to handle all communication with the client. accept() is called to confluence new inbound connection ottempts. The accept handler is registered to handle the next accept asynchronously.

autops insinited to regulate the factor of institution restances and port. It calls accept() for registering the accept handler and does not block. To use it, create an io_context and the server. The executor will run until no async operation is left thanks to async_accept(). It is important that the server lies as long as saync operations on it are processions on it are

ipt() → void {
rectHandler = [this](asio::error_code ec, tcp::socket peer) { if (lec) { // start new session if no error has occurred session = std::make_shared<Session>c(std::movesession>ctart();

struct Server { using tcp = asio::ip::tcp;
 Server(asio::io_context &c, unsigned short port)
 : acceptor(c, tcp::endpoint{tcp::v4(), port}} { accept(); }

cept(); // call accept again to continue accepting new connect

uctor stores the socket with the client connection, stort() initializes the first async read. head() invokes async reading. write() invokes async writing. The fields store the data of the nable_shared_free_this is needed because the session object would die at the end of the accept in the control of the second of the accept in the control of the second of the accept in the control of the second of the accept in the control of the second of the accept in the control of the second of therefore it needs to be allocated on the heap. The handlers need to keep the object alive by pointing on it on the heap. If there is no pointer to the object left, it gets deleted

ivate: suto read() → void; auto write(std::string data) → void; usio::streambuf buffer{}; std::istream input{&buffer}; asio::ip::tcp::socket sock

antO - world / mandO: N

de in Accept handler
session = std::make_shared<Session>(std::move(peer)); session->start();

code in Read handler
Session:read() → void { auto readComplHandl = [self = shared_from_this()] /*...*/}

/ Code in Write handler
uto Session::write(std::string input) → void {
 auto data = std::make_shared<std::string>(input)
 auto writeCompletionHandler = [self = shared_fr

9.2.7. Async Operation without Calibacks

wrations can work "without" callbacks. Specify "special" objects as callbacks.

use_future: Returns a std:: future<1>. Errors are communicated via excep

es in the detected lignores the result of the operation assis:: secaches: ignores the result of the operation asis:: use_awaitable: Returns a std:: awaitable<T> that can be awaited in a coroutine. Complicated!

CEPPs(renner: Standard Identy-header coutenate). #Include <estional>
Most OS support signals. Signals provide a synchronous notifications. They are used to gracefully terminate a program, communicate errors, notify about traps ("ir's a report") Complete Name - Poor - - Coop STERM (Termination requested), SIBSEBV (Insalid memory access), SIBINT (User Interrupt), SIBILL (Ille

9.3.1. Signal Handling in ASIO asio::signal_set defines a set of signals to walt for. Handlers can be set up with signal_set

Placed Size Buffers: asia:::buffee(). Must provide at least as much memory as will be read. Can use several standard containers as a backend. Pointer + Size combinations are available.

Dynamically Sized Buffers: asio::bufmanic_buffer() Use if you do not know the required space and with signal_set_async_wait() that take a lambda. The signal handler receives the signal that occurred and an error if the wait was aborted: [8](auto error, auto sig){...} Useful to cleanly stop server applications.

9.4. ACCESSING SHARED DATA
Multiple ayex operations can be in flight. All completion handlers are dispatched through as to:
and run on a thread executing 10_context.run(). Multiple threads can call run() or
asio::ia_context.This results in a possible data race.

9.3. SIGNAL HANDLING

Strands are a mechanism to en

trands are a mechanism to ensure sequential execution of handlers.

Implicit strands: I implicit strands:

o results = std::vector<int> { }: auto strand = asio::make strand(context)

'in connection class

initiation:

| initiation | initiat

18. ADVANCED LIBRARY DESIGN 10.1. EXCEPTION SAFETY

There is code that handles exceptions code that throws exception, and exc does not cotch, just forwards exceptions). Exception safety is important in generic code that manages resources or data structures (might coil user-defined operations, must not grable its data structures and must not leak resources)

The deterministic lifetime model of C++ requires exception safety. When an exception is thrown unwinding" ends the lifetime of temporary and local objects. Throwing an exception while another ex is "in flight" in the same thread causes the program to terminate.

10.1.1. Safety Levels (from highest to lowest)

sometimes even impossible, e.g. memory allocation. (Examples: Swap, Move Constructor, Move Assign

add::weter
**Discoser
**Discoser
**Comparison
**Compari

auto push(value_type const & elem) \rightarrow void {
 value_type val{elem}; // might threw here due to copy operation failing
 tail_ = (tail_ + 1) % (capacity() + 1); elements_++; /*elements not yet changed+/}

No gowranter: Often unintentional, but happens. Invalid or corrupted data when an exception is thrown

Invariant OK

CPTheleasure nonnext secretar nearcest secretary nearcest belongs to the function signature. You cannot overload on nearcept. can be used to determine the "nearceptiness" of an expression, v nearcept(nearcept(re))) means "Outer function is nearcept, if function f() is all

The compiler might optimize a call of a neexcept function better. But if you throw an exception from a

Destructors must not throw when used during stack unwinding

Move construction and move assignment better not throw (This is why it often uses away is swap should not throw (std:: swap requires non-throwing move op Copyling might throw, when memory needs to be allocated

10.1.4 Standard Library Helners

nevent-move like std : move 16 negreent instead if negreent then move

There are other helpers like is nothrow copy_assignable, destructible, copy_constructible, saspable assignable, move_assignable, copy_assignable, destructible, copy_constructible, saspable

10.1.5. Wide and narrow contracts

- Wide Contract: A function that can handle all argument values of the given param It cannot fail and should be specified as neaccept(true). this, global and exter your Contract: A function that has preconditions on its parameters, e.g. (of parameter must not be negative. Even if not checked and no exception is thrown, those fun

allows later checking and throwing. 10.2. PIMPLIDIOM (POINTER TO IMPLEMENTATION IDIOM) CEPReference: Pixel

Opaque/Incomplete Types: Name known (declared), but not the content. Introduced by a forward declaration.

struct S(\frac{1}{2}; // Definition auto main() \rightarrow int { S s(\frac{1}{2}; foo(s); /* s can now be used */ } Problem: Internal changes in a class' definition require clients to re-compile (e.g. changing a type of a remetale). Solution: Create a "Compilation Firewall": Allow changes to implementation without the need to recompile users. It can be used to shield client code from implementation changes, meaning you must not

auto foo(S & s) → void { foo(s); /* S s{}; is invalid */}

Put in the "exported" header file a class consisting of a "Pointer to Implementation" plus all public members. Basically place all public definition in the header file.

/ wizard.hoo: Minimal Header // WizardImpl.cpp: Full impleme lass Wizard {
 std::shared_ptr<class WizardImpl> pImpl;

ublic: #izard(std::string name = "Wincewind"); plmplstd::mizard(std::string name) plmplstd::make_shareddizardIng auto ddhagic(std::string wish) auto Wizard::doMagic(std::string wish)
 → std::string {
 return pImpl→doMagic(wish);

e streetess. Tento we need to define the destructor of Wisserd after the definition of WizardInpl. The compiler can't move the destructor by himsel

// WizardImpl.cpp rlass WizardImpl { /* ... */ }; std::unique_ptr<class WizardImpl> pImpl;

// ... Wizard::~Wizard() = default: Because the default deleter of std::unique_ptr can't delete an incomplete type, we need to define the destructor explicitly.

10.2.1. Design Decisions with Pimpi Idio

Thems should objects be copied?

No copyling - only moving: std::unique_ptr-class Impl> (device des - shouldow copyling - only moving: std::shney_freclass Impl> (dening me implement)

Deep copyling: std::unique_ptr-class Impl> (with OV copy constructor o 12.1. BUILD AUTOMATION Good for reproducibility, productivity, maintainability and shareability. There are many IDEs which help build our projects. But sometimes, you don't want to rely on an IDE (Nie one build server, when sharing or reproducing

Never do o Inol = null otr, and do not inherit from Pimpi class.

11 HOUDGLASS THTEDEACES 11.1. APPLICATION BINARY INTERFACES (ABI)

ABIs define how programs interact on a binary level (Names of structures and functions, calling conventions, instruction sets). C++ does not define any specific ABI, because they are tightly coupled to the platform. They change between OSes, compiler versions, library versions, etc. Different STL implementations are (usually) incompa



11.1.1. Example: Wizard Class

t Mizard {
 recreated the service of the s

Background C API: Abstract data types can be represented by pointers. Ultimate abstract pointer: void s map to functions taking the abstract data type pointer as first argument. Requires and disposal functions to manage object lifetime. Strings can only be represented by char *. Make sure to not return pointers to temporary objects. Exceptions do not work across a C API, use a Error struct.

struct Wizard * wizard; // Wizard can only be accessed thre struct Wizard const * cwizard;

truct Error + error t; // Stores exception messages, needs to be cleaned up eateMizard(char comst + name, error_t + out_error); // Factory func to mrap oseMizard(wizard toDispose); // Factory function to move

sposefiliare/(extand toDispose); // Factory function to wrop district
sort *error.massog(error.t error); // Alloates sorre assage or
ror.dispose(error.t error); // Alloates sorre assage into the hosp
error.dispose(error.t error); // Alloates error assage from the hosp
error.dispose(error.t error); // Alloates error.er

Parts of C++ that can be used in an extern "C" interface: nctions, but not templates. No overlo

- Transcribe, due not emphases, not openiousnig)
- Cyristiffice Types (see, not, not, see, see)
- Pointers, including function pointers
- Ponward declared drancts
- Unicoped framms withboard class of base types
- Unicoped framms withboard class of base types
- Unicoped framms withboard class of base types
- United Types (you must embrace at with extern "C" when compiling it with C++ (extern "C" (i))

Implementing the Onaque Wizard Type: Wizard class must be implemented. To allow full C++ including ne" class. It wraps the actual Wizard im // Wizard.cpp
extern "C" {
struct Wizard { // C linkage trampoline
Wizard(char const * name) : wiz{name} {}
unseen::Wizard wiz;

Dealing with Exceptions: You can't use references in C API, you must use pointers to pointers. In case of an error, allocate error value on the heap. You must provide a disposal function to clean up. Internally, you can use C++ types, but you should return charconst +, because the caller owns the object providing memory.

Creating Error Messages from Exceptions: Call the function body and catch expressions. Map them to

or object, set the pointer pointed to by out_error. Passed out out_error must not be nullptr root_error = new Errors*("liknown internal errors*) Frors Handling at Client Side: Clientside C+- usage requires mapping error codes Unfortunately, the exception type doesn't map through. But you can use a generic Internetial error. There is a declared RAII class for disposal. You could also use a ten throwing destructor, but this is tricky because of possible leaking.

// C++ Client API (Header Only) // Header continued struct ErrorRAII { struct ThrowOnError nError() = default; OnError() noexcept(false) { -inrowOnError() moexcept(false) {
 if (error.opaque) { throw
 std::runtime_error[error_message(error.opa pperator error_t*() { return Serror.opaque; } private:
 ErrorRAII error{nullptr}; error t opaque:

To complete WizardClient.cpp, call the C functions from WizardClient.hpp from global namespace (i.e. :: do_mogác()). Delete the copy constructor & assignment to prevent a double fre

auto doMagic(std::string, const &wish) → std::string {
 return ::do_magic(wiz, wish.c_str(), ThrowUnError{}); private:
 Wizard(Wizard const &) = delete;
 Wizard & operator =(Wizard const &) = delete;
 wizard wiz;

11.2. JAVA NATIVE ACCESS (JNA) les. Consists of a single JAR file and is cross-platfo JNA provides a simple interface to 0 Type mappings see slides page 31.

blic interface Cplatib extends Library {
Cplatib INSTANCE = (Cplatin) Native load("cpla", Cplatib.class);
void printInt(int number); // matches the function in the next of Calling the loaded library handle THETANCE is only by convention. The leader rearcher for a cuitable library

Plain non-opaque atreet types must inherit from Structure. You must override getFieldUrder() and you can use the tag-interface Structure. ByValue. You can access pointers to such types using getPainter().

first in the path specified by jna. Library. path, otherwise in the system default library search path. Fallback is the class path.

Function names and parameter types must motch. However, the types are not validated. Parameter name 11.2.3. Interfacing with Plain structs

12.2. BUILD TOOLS tools: GNU make, Scons, Ninia, CMake, autotools. Features of Build Tools: Incremental builds, parallel builds, automatic dependency resolution, package management, automatic test execution, platform independence, additional processing of build products. Different Classes of Build Automation Software: ifferent Classes of Build Automation Software:

Make-style Build Tools: Run build scripts, produce your final products, often verbose, use a language

Do not write your own scripts for this process, because then every source file gets built every time, the

Managing lifetime is not trivial. Using dispose ... () API functions in finalizers is not recommended. Either provide a dispose method on your Java type or implement AutoClosable and use your objects with try-

Use Intibule foreigness are supported by the size of the buffer. Requires that the API supports it. getByteArray() copies the data from the buffer. Make sure to free the buffer either using an API free ...() functions or

agnostic configuration language Build Script Generators: Generate configurations for Make-style Build Systems or Build Scripts, configuration is independent of actual build tool, often have advanced features like download dependencies.

11.2.4. Working with Raw Byte Arrays

12. BUILD SYSTEMS

ands tend to be platf

Well-known tool to build all kinds of projects. Many IDEs "understand" make projects. The workflow descrip tion is in Makefile via "Target" rules. Each target may have one or more prerequisites and execute one or more commands to generate one or more results. Targets are then executed "top-down". A target is only executed if required.

target: prereq_target
prereq_target: prereq_file other_target
command_to_generate_output

Pros: Very generic, powerful pattern matching mechanism, builds only what is needed, when its
 Cons: Often platform-specific commands, need to specify how to do thins

12.3. BUILD SCRIPT GENERATORS int to achieve, not how to do it. Work on a higher level, let the tool create the actual build

configurations. Platform independent build specification, tool indep 12.3.1. CMake # CMakeLists.txt

#Project("my_app" LANGUAGES CXX)

#:cout << "Hello There!\n";} add_executablef"## ann" "mai- --

midir build # create separate build directory for the build output

fred(...) sets the minimum required CMake version. This implicitly defines the and defines the name of the project and which language we us

reference. J command defines the name of the proper are whose which imaging features are used by the target i.e. the C+ interpet. Companies. J command the command of the c

(es (...) defines additional target propertie

add_Library(...): Defaults to static libraries, can be overridden at configuration time. All features, include paths and dependencies should be PUBLIC. get_link_libraries is used to define libraries required by a target. Can be PUBLIC or PRIVATE, applies

PUBLIC features/dependencies/includes of the library Variables can be defined using set (VAR_MAME VALUE). They are referenced using \${VAR_MAME}\$. This global variables like PROJECT_MAME, PROJECT_SOURCE_DIR or PROJECT_BINARY_DIR. Can be used in p concrete values: add_executable (\${PROJECT_MAME}} "source1.cpp" "source2.cpp" ...).

12.3.2. Testing with CMake
CMake includes CTest. Enable it with enable_testing(). Create a "Test Ru

12.4. PROJECT LAYOUT

Project Layout (Genera

late Laufe Contains headers. Add subfolders for separate subsystems if needed were Contains implementations. Subfolders inpost should match six-Level folder List history, party; Contains external resources like librarie starts: Contains texts, and above folders as necessary fault orange fifes should be in the project root them creating, all been, introduce another layer of nesting to avoid filename clashes in clients.

truct Example {
Example(int a) { std::cout < "Ctr\n"; }
Example(int a) { std::cout < "Copy Ctr\n"; }
Example(Example(int) { std::cout < "Copy Ctr\n"; }
Example(Example(int) { std::cout < "Move Ctr\n"; }
-*Example(j std::cout < "Workn"; }</pre>

13. MOVE SEMANTICS DUTPUT

1. Creating new object with a still calls Constructor.

Example a = 5; // "Ctr"
Example b = std::move(a); // "Move Ctr"
Example c = Kexmple(std::move(f)); // "Move Ctr"
Example const d = c; // "Copy Ctr"
Example const d = c; // "Copy Ctr"
Example = std::move(a); // "Copy Ctr"
(Moving const no

2. Constructor elides creation of temporaries (Optimization suto defaultExample() → Example { return Example{5}; }
ixample f{Example{5}}; // "Ctr"

3 Temporary Ivalue in methods auto noveExample(Example e) { auto h = std::forward<Example>(e); } // "Move Ctr" "Dto:
moveExample(g); // "Copy Ctr" "Dtor"

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→ std::string;