

Controllability and open-loop control



Lecture 10

Jerome Jouffroy, Professor
jerome@sdu.dk

Today's lecture

- **Controllability concept and criterion**
- **Basics of open-loop control for linear systems**
- **Polynomial trajectories**
- **Open-loop control for general linear systems**

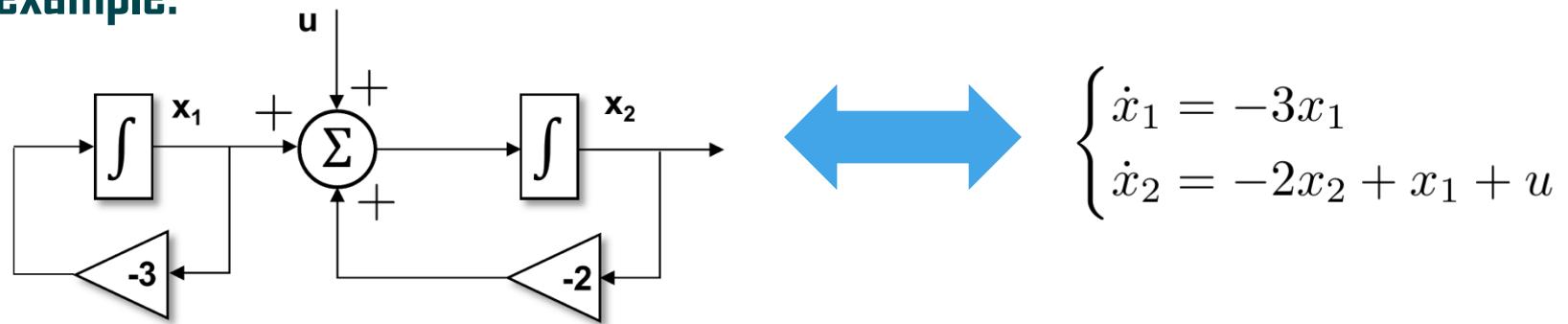
Controllability in a nutshell (1/2)

Basic controllability notion: Starting from the origin, can we use $u(t)$ to bring the state to any target state x_T ?

→ $u(t)$ needs to be able to influence all state components
(directly or indirectly)

Graphical interpretation with block diagrams:

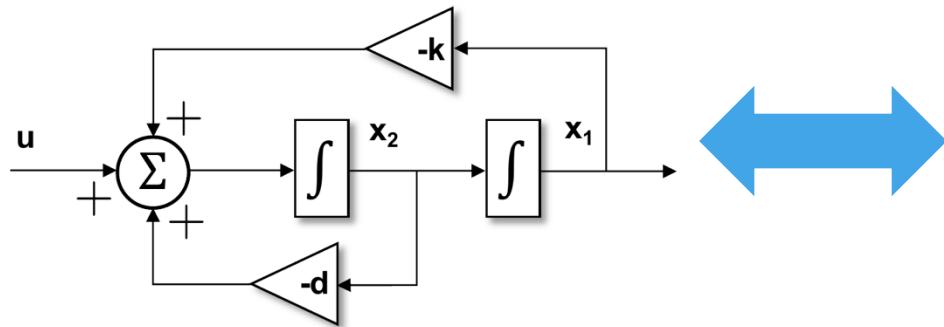
First example:



follow the flow of the arrows: u influences x_2 but not x_1

Controllability in a nutshell (2/2)

Second example:



$$\begin{bmatrix} \dot{x}_1(t) \\ \dot{x}_2(t) \end{bmatrix} = \begin{bmatrix} 0 & 1 \\ -\frac{k}{m} & -\frac{d}{m} \end{bmatrix} \begin{bmatrix} x_1(t) \\ x_2(t) \end{bmatrix} + \begin{bmatrix} 0 \\ \frac{1}{m} \end{bmatrix} u(t)$$

u influences directly x_2 and indirectly x_1 (through x_2)

→ this system is controllable

Definition: controllability

The system represented by: $\dot{\mathbf{x}} = \mathbf{A}\mathbf{x} + \mathbf{B}u$ is controllable if, for $\mathbf{x}(0) = 0$ and for any point \mathbf{x}_T of the state-space, there exists a finite time $T > 0$ and a control input $u(t)$, $t \in [0, T]$, such that $\mathbf{x}(T) = \mathbf{x}_T$. \square

The Kalman criterion for checking controllability

Systematic way to check controllability: Kalman criterion

consider the controllability matrix defined as

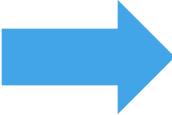
$$\mathbf{W}_c := [\mathbf{B}, \mathbf{AB}, \mathbf{A}^2\mathbf{B}, \dots, \mathbf{A}^{n-1}\mathbf{B}]$$

Main result

The system $\dot{\mathbf{x}} = \mathbf{Ax} + \mathbf{Bu}$ is controllable if, and only if the controllability matrix \mathbf{W}_c as given by equation (4.16) is **invertible!** □

Example:

$$\dot{\mathbf{x}} = \begin{bmatrix} -1 & 1 & 0 \\ 1 & -1 & 1 \\ 0 & 1 & -1 \end{bmatrix} \mathbf{x} + \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} u$$

compute 

$$\mathbf{W}_c = [\mathbf{B}, \mathbf{AB}, \mathbf{A}^2\mathbf{B}] = \begin{bmatrix} 0 & 1 & -2 \\ 1 & -1 & 3 \\ 0 & 1 & -2 \end{bmatrix}$$
$$\Rightarrow \det(\mathbf{W}_c) = 0$$



Controllability and the Controllability Canonical Form (CCF)

Throughout this course, we have often seen this form:

$$\dot{\mathbf{x}}(t) = \begin{bmatrix} 0 & 1 & 0 & 0 & \dots & 0 \\ 0 & 0 & 1 & 0 & \dots & 0 \\ 0 & 0 & 0 & 1 & \dots & 0 \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & 0 & \dots & 1 \\ -a_0 & -a_1 & -a_2 & -a_3 & \dots & -a_{n-1} \end{bmatrix} \mathbf{x}(t) + \begin{bmatrix} 0 \\ 0 \\ 0 \\ \vdots \\ 0 \\ b \end{bmatrix} u(t)$$

It is called a **Controllability Canonical Form (CCF)**

why?

$$\mathbf{B} = \begin{bmatrix} 0 \\ 0 \\ 0 \\ \vdots \\ 0 \\ b \end{bmatrix} \quad \mathbf{AB} = \begin{bmatrix} 0 \\ 0 \\ \vdots \\ 0 \\ b \\ * \end{bmatrix} \quad \mathbf{A}^2\mathbf{B} = \begin{bmatrix} 0 \\ \vdots \\ 0 \\ b \\ * \\ * \end{bmatrix} \quad \text{and so on until} \quad \mathbf{A}^{n-1}\mathbf{B} = \begin{bmatrix} b \\ * \\ * \\ \vdots \\ * \\ * \end{bmatrix}$$



\mathbf{W}_c shows that this system is always controllable as long as $b \neq 0$

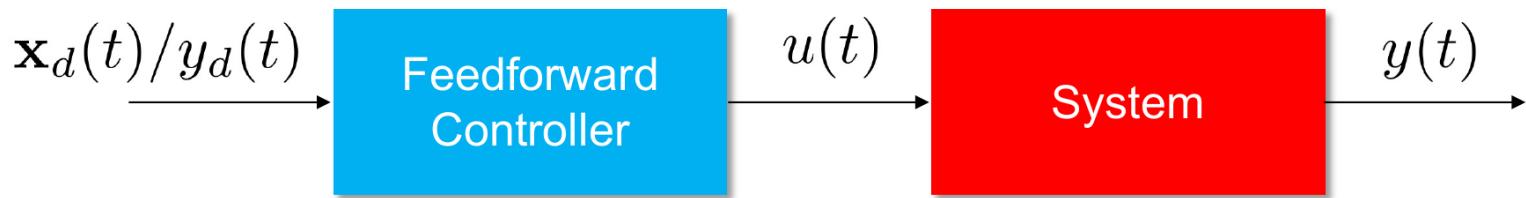
The open-loop control problem

control is not always only about feedback

open-loop control objective:

Make the system follow a
pre-defined trajectory $y_d(t)/x_{1d}(t)$

by computing the
corresponding $u(t) = u_d(t)$



advantages:

- no sensor is needed to do real-time control (lower budget on sensing)
- computationally more effective (everything is pre-computed, no need to sensing-computing-actuation loop)

alternative names: feedforward control, motion planning

Open-loop control for CCFs (1/2)

Take the CCF

Assume that $y_d(t)$ is a polynomial desired trajectory. If $y(t)$ follows $y_d(t)$, then

$$\begin{cases} \dot{x}_1 = x_2 \\ \dot{x}_2 = x_3 \\ \dot{x}_3 = x_4 \\ \vdots \\ \dot{x}_{n-1} = x_n \\ \dot{x}_n = -a_0x_1 - a_1x_2 - \dots - a_{n-1}x_n + bu \end{cases} \quad \text{with} \quad y = x_1$$

→ $x_1(t) = x_{1d}(t) = y_d(t)$ is a polynomial

→ $x_2(t) = \dot{x}_1(t) = \dot{x}_{1d}(t) = \dot{y}_d(t)$ is a polynomial (ie derivative of polynomial)

→ $x_3(t) = \dot{x}_2(t) = \ddot{x}_1(t) = \ddot{x}_{1d}(t) = \ddot{y}_d(t)$ is a polynomial

→ $x_n(t) = \dot{x}_{n-1}(t) = x_1^{(n-1)}(t) = y_d^{(n-1)}(t)$ is a polynomial

→ $\dot{x}_n(t) = x_1^{(n)}(t) = y_d^{(n)}(t)$ is a polynomial

finally, isolating $u(t)$ in the CCF:

$$u(t) = u_d(t) = \frac{1}{b} \left[a_0 y_d(t) + a_1 \dot{y}_d(t) + \dots + a_{n-1} y_d^{(n-1)}(t) + y_d^{(n)}(t) \right]$$

Open-loop control: the MSD example

Example:
$$\begin{bmatrix} \dot{x}_1(t) \\ \dot{x}_2(t) \end{bmatrix} = \begin{bmatrix} 0 & 1 \\ -k/m & -d/m \end{bmatrix} \begin{bmatrix} x_1(t) \\ x_2(t) \end{bmatrix} + \begin{bmatrix} 0 \\ 1/m \end{bmatrix} u(t)$$
 with $y(t) = x_1(t)$

Take the polynomial desired trajectory $y_d(t)$

If the output follows the desired trajectory, then

$x_1(t) = y_d(t)$ is a polynomial

$x_2(t) = \dot{x}_1(t) = \dot{x}_{1d}(t) = \dot{y}_d(t)$ is a polynomial

$\dot{x}_2(t) = \ddot{x}_1(t) = \ddot{y}_d(t)$ is a polynomial

So that we have the open-loop controller

$$u_d(t) = k y_d(t) + d \dot{y}_d(t) + m \ddot{y}_d(t)$$
 which will make the output follow the desired trajectory

Polynomial trajectories (1/2)

Let us see how to construct a polynomial trajectory (2nd order case).

we want to go from $\mathbf{x}(0) = \begin{bmatrix} x_{10} \\ x_{20} \end{bmatrix} =: \mathbf{x}_0$ to $\mathbf{x}(T) = \begin{bmatrix} x_1(T) \\ x_2(T) \end{bmatrix} = \begin{bmatrix} x_{1T} \\ x_{2T} \end{bmatrix} =: \mathbf{x}_T$
($\mathbf{x}(t)$ is the state of a CCF) in time T

a polynomial trajectory is described by

$$y_d(t) = \alpha_0 + \alpha_1 t + \alpha_2 t^2 + \alpha_3 t^3 + \dots$$

2 initial and 2 final conditions  we need 4 coef. = poly of order 3

we have $\begin{cases} x_1(t) = y_d(t) = \alpha_0 + \alpha_1 t + \alpha_2 t^2 + \alpha_3 t^3 & (1) \\ x_2(t) = \dot{x}_1(t) = \dot{y}_d(t) = \alpha_1 + 2\alpha_2 t + 3\alpha_3 t^2 & (2) \end{cases}$
and its derivative

at $t=0$, (1) gives $x_1(0) = y_d(0) = x_{10} = \alpha_0$  we have the first
2 coefficients!
and (2) gives $x_2(0) = \dot{y}_d(0) = x_{20} = \dot{x}_1(0) = \alpha_1$

Polynomial trajectories (2/2)

$$\begin{cases} x_1(t) = y_d(t) = \alpha_0 + \alpha_1 t + \alpha_2 t^2 + \alpha_3 t^3 \\ x_2(t) = \dot{x}_1(t) = \dot{y}_d(t) = \alpha_1 + 2\alpha_2 t + 3\alpha_3 t^2 \end{cases} \quad (1) \quad (2)$$

At time $t = T$

(1) and (2) give $\begin{cases} x_1(T) = y_d(T) = \alpha_0 + \alpha_1 T + \alpha_2 T^2 + \alpha_3 T^3 = x_{1T} \\ x_2(T) = \dot{y}_d(T) = \alpha_1 + 2\alpha_2 T + 3\alpha_3 T^2 = x_{2T} \end{cases}$

separate known from unknown terms

$$\begin{cases} x_{1T} - (x_{10} + x_{20}T) = \alpha_2 T^2 + \alpha_3 T^3 \\ x_{2T} - x_{20} = 2\alpha_2 T + 3\alpha_3 T^2 \end{cases}$$

rewrite in vectorial form

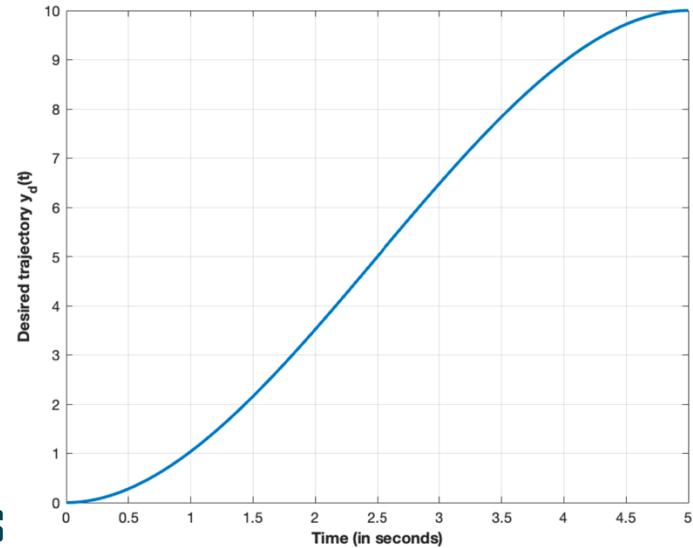
$$\begin{bmatrix} x_{1T} - (x_{10} + x_{20}T) \\ x_{2T} - x_{20} \end{bmatrix} = \begin{bmatrix} T^2 & T^3 \\ 2T & 3T^2 \end{bmatrix} \begin{bmatrix} \alpha_2 \\ \alpha_3 \end{bmatrix}$$

invert the time matrix to get the last two coefficients

$$\begin{bmatrix} \alpha_2 \\ \alpha_3 \end{bmatrix} = \begin{bmatrix} T^2 & T^3 \\ 2T & 3T^2 \end{bmatrix}^{-1} \begin{bmatrix} x_{1T} - (x_{10} + x_{20}T) \\ x_{2T} - x_{20} \end{bmatrix}$$

(implementation trick:)

$$\begin{bmatrix} x_{1T} - (x_{10} + x_{20}T) \\ (x_{2T} - x_{20})T \end{bmatrix} = \begin{bmatrix} T^2 & T^3 \\ 2T^2 & 3T^3 \end{bmatrix} \begin{bmatrix} \alpha_2 \\ \alpha_3 \end{bmatrix}$$



$$y_d(t) = \alpha_0 + \alpha_1 t + \alpha_2 t^2 + \alpha_3 t^3$$

example:

$$y_d(0) = 0 \text{ to } y_d(T) = 10$$

in $T = 5$ seconds

$$\text{with } \dot{y}_d(0) = \dot{y}_d(T) = 0$$

Back to the MSD

Example: Previously (slide 9), we obtained the open-loop controller

$$u_d(t) = ky_d(t) + d\dot{y}_d(t) + m\ddot{y}_d(t)$$

since $y_d(t) = \alpha_0 + \alpha_1 t + \alpha_2 t^2 + \alpha_3 t^3$

we have $\dot{y}_d(t) = \alpha_1 + 2\alpha_2 t + 3\alpha_3 t^2$

and $\ddot{y}_d(t) = 2\alpha_2 + 6\alpha_3 t$

so that we have the polynomial expression

$$u_d(t) = k(\alpha_0 + \alpha_1 t + \alpha_2 t^2 + \alpha_3 t^3) + d(\alpha_1 + 2\alpha_2 t + 3\alpha_3 t^2) + m(2\alpha_2 + 6\alpha_3 t)$$

Remark: do not implement things like this! This way of programming is too tedious ☺

State-coordinate transformation

Our DL algorithm works for a CCF. What about more general forms?

→ trick: transform a SS rep. back into a CCF to do DL control

Relation between 2 different SS rep.?

start with (1) $\begin{cases} \dot{x} = Ax + Bu \\ y = Cx \end{cases}$ and define the invertible change
of coordinates $z(t) = Tx(t)$
implies $x = T^{-1}z$ and $\dot{z} = Tx$ with constant $T \in \mathbb{R}^{n \times n}$

so that (1) is transformed into $\begin{cases} \dot{z} = T(Ax + Bu) \\ y = CT^{-1}z \end{cases}$



$$\begin{cases} \dot{z} = TAT^{-1}z + TBu \\ y = CT^{-1}z \end{cases}$$



$$\boxed{\begin{cases} \dot{z} = \tilde{A}z + \tilde{B}u \\ y = \tilde{C}z \end{cases}}$$

A second-order example

Example: start with ODE $\ddot{y} + 3\dot{y} = 2y + u$ (ODE)

a CCF of (ODE) is $\dot{x} = \begin{bmatrix} 0 & 1 \\ 2 & -3 \end{bmatrix} x + \begin{bmatrix} 0 \\ 1 \end{bmatrix} u$

Define a change of coordinate $z = T x$ where $T = \begin{bmatrix} 3 & 1 \\ 1 & 0 \end{bmatrix}$

so that we have $\dot{z} = \tilde{A}z + \tilde{B}u$ (SS rep. after transform, in the z-space)

with $\tilde{A} = T A T^{-1} = \begin{bmatrix} 3 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 0 & 1 \\ 2 & -3 \end{bmatrix} \begin{bmatrix} 0 & 1 \\ 1 & -3 \end{bmatrix} = \begin{bmatrix} 0 & 2 \\ 1 & -3 \end{bmatrix}$

and $\tilde{B} = T B = \begin{bmatrix} 3 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$

ie we have the new SS $\dot{z} = \begin{bmatrix} 0 & 2 \\ 1 & -3 \end{bmatrix} z + \begin{bmatrix} 1 \\ 0 \end{bmatrix} u$ (not a CCF!)

Transforming a SS-rep. into a CCF (1/2)

Objective: we want to find the change of coordinates $\mathbf{z} = \mathbf{T}\mathbf{x}$

so that $\dot{\mathbf{z}} = \tilde{\mathbf{A}}\mathbf{z} + \tilde{\mathbf{B}}\mathbf{u}$ is a CCF

Recall that

$$\tilde{\mathbf{A}} = \mathbf{T}\mathbf{A}\mathbf{T}^{-1}$$



$$\tilde{\mathbf{A}}\mathbf{T} = \mathbf{T}\mathbf{A}$$

and

$$\tilde{\mathbf{B}} = \mathbf{T}\mathbf{B}$$

Let

$$\mathbf{T} = \begin{bmatrix} \mathbf{T}_1 \\ \mathbf{T}_2 \\ \vdots \\ \mathbf{T}_n \end{bmatrix}$$

(we split \mathbf{T} into n row vectors)

$$\begin{bmatrix} 0 & 1 & \dots & 0 \\ \vdots & 0 & \ddots & \vdots \\ 0 & \dots & 0 & 1 \\ -a_0 & -a_1 & \dots & -a_{n-1} \end{bmatrix}$$

$$\begin{bmatrix} \mathbf{T}_1 \\ \mathbf{T}_2 \\ \vdots \\ \mathbf{T}_n \end{bmatrix} = \begin{bmatrix} \mathbf{T}_1\mathbf{A} \\ \mathbf{T}_2\mathbf{A} \\ \vdots \\ \mathbf{T}_n\mathbf{A} \end{bmatrix}$$

SDU $\begin{bmatrix} \mathbf{T}_2 \\ \mathbf{T}_3 \\ \vdots \\ -a_0\mathbf{T}_1 - a_1\mathbf{T}_2 - \dots - a_{n-1}\mathbf{T}_n \end{bmatrix} = \begin{bmatrix} \mathbf{T}_1\mathbf{A} \\ \mathbf{T}_2\mathbf{A} \\ \vdots \\ \mathbf{T}_n\mathbf{A} \end{bmatrix}$ ie $\boxed{\mathbf{T}_2 = \mathbf{T}_1\mathbf{A}}$
 $\boxed{\mathbf{T}_3 = \mathbf{T}_2\mathbf{A}}$ etc...

if we have \mathbf{T}_1 ,
we have all the other \mathbf{T}_i 's!

Transforming a SS-rep. into a CCF (2/2)

Remaining question: we need to find T_1



let us use

$$\tilde{\mathbf{B}} = \mathbf{T}\mathbf{B}$$

with $\tilde{\mathbf{B}}^T = [0 \dots 0 1]$

(CCF with $b=1$)

this gives

$$\tilde{\mathbf{B}} = \begin{bmatrix} \mathbf{T}_1\mathbf{B} \\ \mathbf{T}_2\mathbf{B} \\ \mathbf{T}_3\mathbf{B} \\ \vdots \\ \mathbf{T}_n\mathbf{B} \end{bmatrix} = \begin{bmatrix} \mathbf{T}_1\mathbf{B} \\ \mathbf{T}_1\mathbf{A}\mathbf{B} \\ \mathbf{T}_1\mathbf{A}^2\mathbf{B} \\ \vdots \\ \mathbf{T}_1\mathbf{A}^{n-1}\mathbf{B} \end{bmatrix} \xrightarrow{\text{transpose}} \tilde{\mathbf{B}}^T = [0 \dots 0 1] = [\mathbf{T}_1\mathbf{B} \quad \mathbf{T}_1\mathbf{A}\mathbf{B} \quad \mathbf{T}_1\mathbf{A}^2\mathbf{B} \quad \dots \quad \mathbf{T}_1\mathbf{A}^{n-1}\mathbf{B}]$$

so that we have $\tilde{\mathbf{B}}^T = \mathbf{T}_1\mathbf{W}_c$



inverting \mathbf{W}_c , we get \mathbf{T}_1 as

$$\mathbf{T}_1 = \tilde{\mathbf{B}}^T \mathbf{W}_c^{-1}$$

Back to the second-order example

Example: start with $\dot{x} = \begin{bmatrix} 0 & 2 \\ 1 & -3 \end{bmatrix} x + \begin{bmatrix} 1 \\ 0 \end{bmatrix} u$ (not a CCF)

Transform this into a CCF?

compute the controllability matrix $\mathbf{W}_c = [\mathbf{B} \quad \mathbf{AB}] = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$
and set $\tilde{\mathbf{B}}^T = [0 \quad 1]$

Calculate $\mathbf{T}_1 = \tilde{\mathbf{B}}^T \mathbf{W}_c^{-1} = [0 \quad 1] \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} = [0 \quad 1]$

which gives $\mathbf{T}_2 = \mathbf{T}_1 \mathbf{A} = [0 \quad 1] \begin{bmatrix} 0 & 2 \\ 1 & -3 \end{bmatrix} = [1 \quad -3]$

so that we have the coordinate transform $\mathbf{T} = \begin{bmatrix} 0 & 1 \\ 1 & -3 \end{bmatrix}$

giving the new SS rep.

$$\dot{\mathbf{z}} = \begin{bmatrix} 0 & 1 \\ 2 & -3 \end{bmatrix} \mathbf{z} + \begin{bmatrix} 0 \\ 1 \end{bmatrix} u \quad (\text{CCF})$$

Open-control for general linear systems

Algorithm - Open-Loop control for general LTI systems

Given the following state-space representation

$$\begin{cases} \dot{\mathbf{x}} = \mathbf{Ax} + \mathbf{Bu} \\ y = \mathbf{Cx} \end{cases}$$

- After checking the controllability of the system, compute the transform \mathbf{T} and obtain a CCF.
- We want the state $\mathbf{x}(t)$ to go from \mathbf{x}_0 to \mathbf{x}_T . Translate this in the z -space: we want the state $\mathbf{z}(t)$ to go from $\mathbf{z}_0 = \mathbf{T}\mathbf{x}_0$ to $\mathbf{z}_T = \mathbf{T}\mathbf{x}_T$.
- Construct polynomial trajectory $y_d(t)$ using \mathbf{z}_0 and \mathbf{z}_T as boundary conditions (ie we want to have $y_d(0) = z_{1,0}$, $\dot{y}_d(0) = z_{2,0}$, ..., for the initial conditions and $y_d(T) = z_{1,T}$, $\dot{y}_d(T) = z_{2,T}$, ..., for the final conditions).
- Implement feedforward controller

$$u_d(t) = a_0 y_d(t) + a_1 \dot{y}_d(t) + \dots + a_{n-1} y_d^{(n-1)}(t) + y_d^{(n)}(t)$$

Application



Beyond linear systems (1/2)

Beyond linear systems (2/2)

