Warp Scheduling and Divergence

Soumyajit Dey, Assistant Professor, CSE, IIT Kharagpur

February 13, 2020



GPU can be viewed as an array of Streaming Multiprocessors (SMs) Each SM has the following elements

- ► Registers that can be partitioned among threads of execution
- ► Several Caches: Shared memory, Constant, Texture, L1 etc
- ► Warp Schedulers (More on this later)
- Scalar Processors(SPs) for integer and floating-point operations
- ► Special Function Units (SFUs) for single-precision floating-point transcendental functions

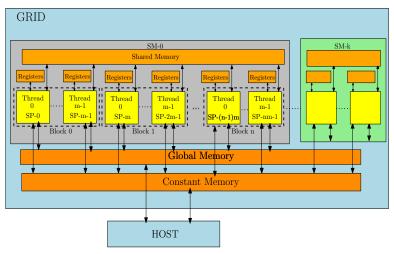


Table: CUDA Device Memory Types and Scopes

Variables Declaration	Memory	Scope	Lifetime
Automatic Variables other than arrays	Register	Thread	Kernel
Automatic array variables	Local	Thread	Kernel
deviceshared int SharedVar	Shared	Block	Kernel
device int GlobalVar	Global	Grid	Application
deviceconstant int ConstVar	Constant	Grid	Application



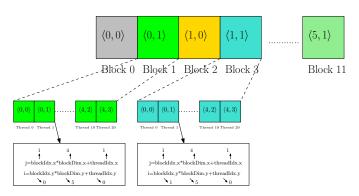
Mapping to Hardware







Example: CUDA Thread and Block Definition





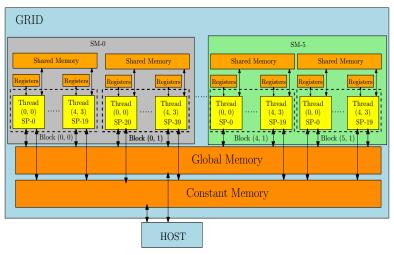


Generalized Mapping Scenario

- ► Let us consider a scenario for the grid and block dimensions specified above.
- gridDim = <6,2> and blockDim = <5,4>
- ► #SMs = 6 #SPs per SM = 40
- ► Two Blocks are mapped to one SM at a time.
- ► Hardware resources are completely utilized.



Mapping to Hardware





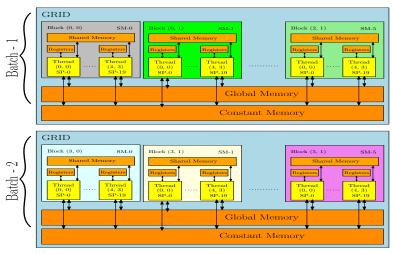


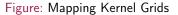
Mapping in a resource constrained setting

- ► Consider a scenario where the resources of the architecture are limited.
- gridDim = < 6, 2 > and blockDim = < 5, 4 >
- ▶ #SMs = 6 #SPs per SM = 20
- Thread Blocks are launched in batches sequentially.
- ▶ Execution is serialized to some extent.



Mapping to Hardware







SM, SP, Block and thread

- ▶ thread block max size : 1024 (modern archs 2048)
- ► SM can store max 1024 "thread contexts"
- ► can have much less than 1024 SPs
- ► GTX 970 : 13 SMs : 13 X 1024 thread contexts in parallel
- ► GTX 970 : 128 SP per SM



SM, SP, Block and thread

- ► One block in one SM
- ► One SM can have multiple blocks

If SM can store max 1024 "thread contexts", and block size is 256, we have 4 blocks per SM.

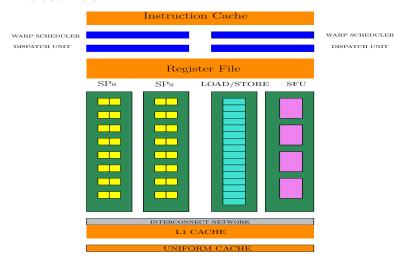


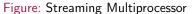
GPU HW scheduler

- ► The hw scheduler decided which threads to map to a collection of SPs in SIMD fashion :: SIMT model of execution
- ► This collection is physically guaranteed to execute in parallel
- ► The unit of such collections is "warp"



SM: A closer look





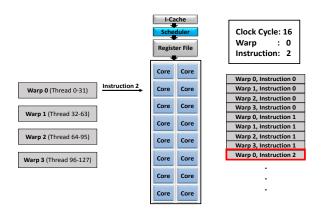


Warps

- ► Warp is a unit of thread Scheduling in SMs.
- ► Warp size is implementation specific (typically 32 threads)
- Warps are executed in an SIMD fashion i.e. the warp scheduler launches warps of threads and each warp typically executes one instruction across parallel threads.

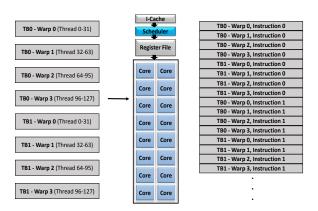
Ex : If a SM has 128 SPs, it can execute 4 Warps at a given time (one Warp has 32 Threads)





Ref : Henk Corporaal, Gert-Jan van den Braak - "Introduction to GPGPU Architectures"





► Thread block scheduler (TBS) is *believed to use* round robin policy to schedule thread blocks - implementation dependent



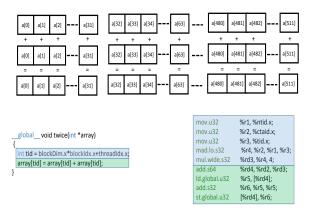


Figure: Simple CUDA Kernel



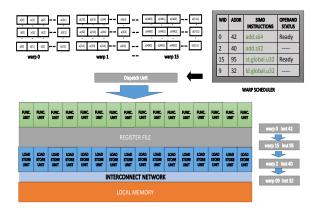


Figure: Warp Scheduler



- ► Issue one "ready-to-go" warp instruction/cycle
- Use operand score-boarding to prevent hazards
- ▶ Issue selection based on round-robin/age of warp
- Score-boarding determines if a thread is ready to execute?
- Scoreboard is a HW implemented table that tracks instrs fetched, resource availability for fetched instrs (FU and operand), register file modifications by instrs.



Latency Tolerance

- ► When threads in one warp execute a long-latency operation (read from global memory), the warp scheduler will dispatch and execute other warps until that operation is finished.
- ▶ Other long latency operations : FP units, Branch instructions
- ► After all, all threads in the same control-flow execute same instruction sequence on different data points!
- ► A common practice is to launch thread blocks of a size that is a multiple of the warp size to maximally utilize threads.
- ► Slow global memory accesses by threads in a warp may be optimized using coalescing (more on this later)



Efficient use of thread blocks

Target System Constraints

- ► A maximum of 8 blocks and 1024 threads per SM
- ► A maximum of 512 threads per block

Table: Solutions for various block scenarios

Input Block Size	Blocks per SM	Threads per Block	Remarks
8 * 8	12	64	SM execution resources will be underutilized
16*16	4	256	Achieves full thread capacity in SMs
32*32	1	1024	Exceeds the limit of 512 threads per block



CUDA API provides constructs for obtaining properties of the target GPU.

- cudaGetDeviceCount(): Obtains the number of devices in the system.
- cudaGetDeviceProperties(): Returns the property values of a particular device



```
int main()
{
    int devCount;
    cudaGetDeviceCount(&devCount);
    for (int i = 0; i < devCount; ++i)
    {
        cudaDeviceProp devp;
        cudaGetDeviceProperties(&devp, i);
        printDevProp(devp);
    }
    return 0;
}</pre>
```



```
void printDevProp(cudaDeviceProp devProp)
printf("Major revision number: %d\n",devProp.major);
 printf("Minor revision number: %d\n",devProp.minor);
 printf("Name: %s\n",devProp.name);
 printf("Total global memory: u\n",devProp.totalGlobalMem);
 printf("Total shared memory per block: %u\n", devProp.
     sharedMemPerBlock):
 printf("Total registers per block: %d\n", devProp.
    regsPerBlock);
 printf("Warp size: %d\n",devProp.warpSize);
 printf("Maximum memory pitch: %u\n",devProp.memPitch);
 printf("Maximum threads per block: %d\n",devProp.
    maxThreadsPerBlock):
 for (int i = 0; i < 3; ++i)
 printf("Maximum dimension %d of block: %d\n",i,devProp.
      maxThreadsDim[i]):
 for (int i = 0; i < 3; ++i)
 printf("Maximum dimension %d of grid: %d\n", i, devProp.
      maxGridSize[i]):
```





Example: Tesla K40m Characteristics

```
Major revision number: 3
Minor revision number: 5
Name: Tesla K40m
Total global memory: 3405643776
Total shared memory per block:49152
Total registers per block: 65536
Warp size: 32
Maximum memory pitch: 2147483647
Maximum threads per block: 1024
Maximum dimension 0 of block: 1024
Maximum dimension 1 of block: 1024
Maximum dimension 2 of block: 64
Maximum dimension 0 of grid: 2147483647
Maximum dimension 1 of grid: 65535
Maximum dimension 2 of grid: 65535
Clock rate: 745000
Total constant memory:65536
Texture alignment: 512
Concurrent copy and execution: Yes
Number of multiprocessors: 15
```

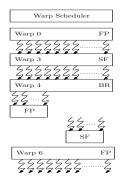


Control Flow Divergence

- ► Threads inside a warp execute the same instruction.
- ► How does a warp handle if statements / branch instructions?
- ► The GPU is not capable of running both the if else blocks at the same time.



Warp Scheduling



Threads inside a warp executing same instruction - efficient

Threads inside a warp executing different instruction - in efficient $\,$

Figure: Warp Divergence



Consider the following kernel code

Half the threads of a warp execute the addition instruction while the other half execute the subtraction instruction.

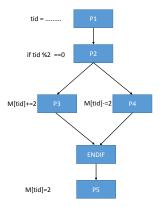


The Hardware's Job

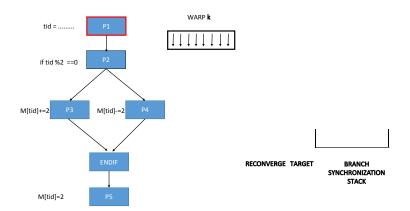
The GPU has hardware support for handling divergent branch instructions in code.

- ► The PTX Assembler maintains internal masks, a branch synchronization stack and special markers
- ► The PTX Assembler sets a branch synchronization marker first for the divergent if statement that pushes the active mask on a stack inside each SIMD thread
- Depending on the value of the mask relevant threads execute instructions,
- ► Once the instructions in the if block are finished, the active mask is popped from the stack, flipped and pushed back.

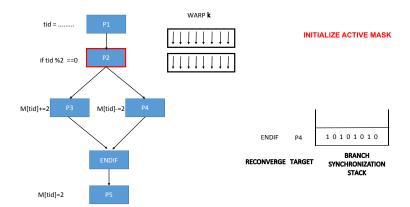




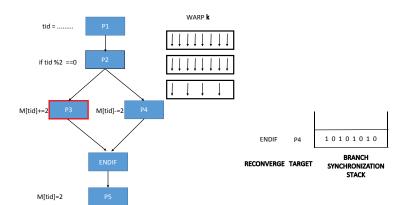




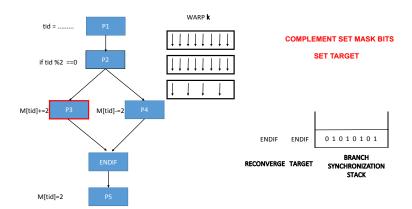




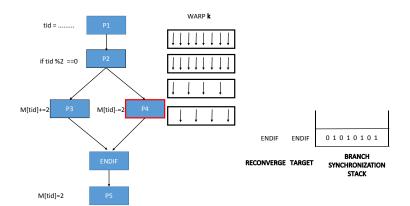




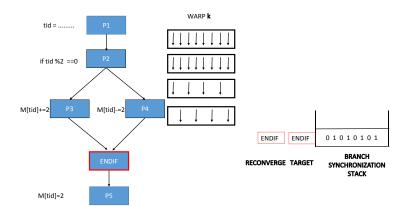




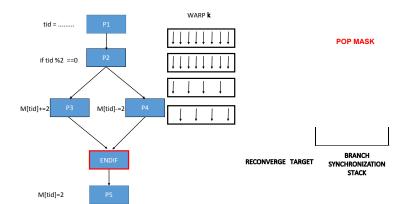




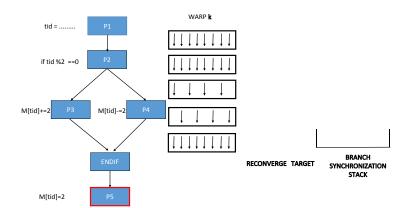














Observations

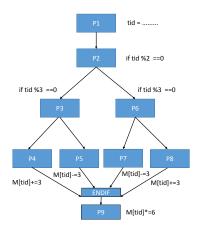
- ► The target value represents the address of the instruction to be executed by the warp, once the current conditional block of instructions has finished.
- ► The reconvergence value represents the address of the convergence statement, once both the conditional blocks associated with an if-else statement has finished execution.
- ► The mask is popped from the stack once the conditional block is executed both way (which is known from the target and reconvergence marker).



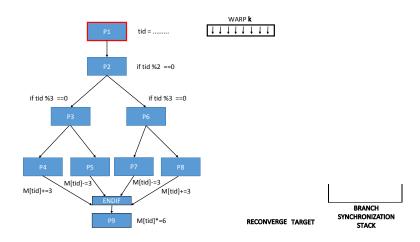
Let us consider an example that has nested if/else statements.

```
__global__
void divergence(float *M)
/*P1*/ int tid=blockIdx.x*blockDim.x+threadIdx.x;
         if (tid%2==0)
/*P2*/
/*P3*/ if (tid%3==0)
             M[tid]+=3;
/*P4*/
           else
/*P5*/
             M[tid] = 3;
          else
/*P6*/
          if(tid%3==0)
/*P7*/
             M[tid]-=3;
           else
/*P8*/
             M[tid] += 3;
          M[tid]*=6;
/*P9*/
```

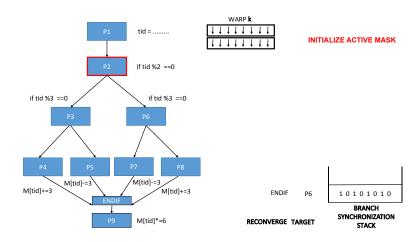




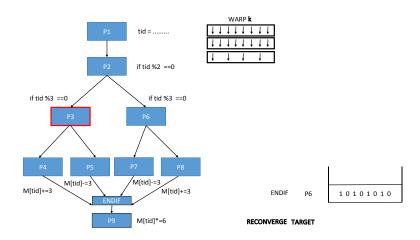




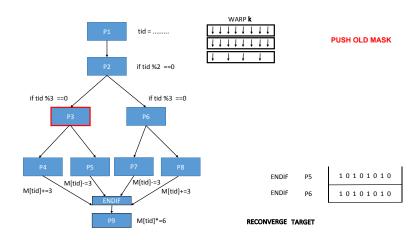




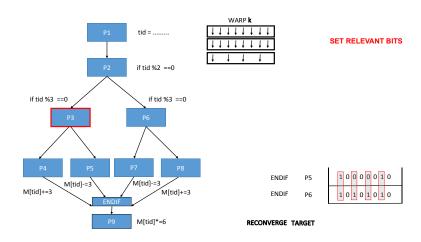




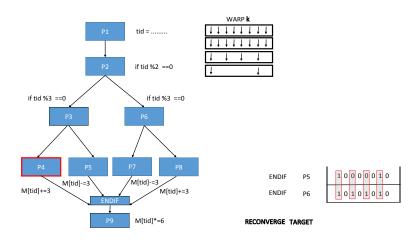




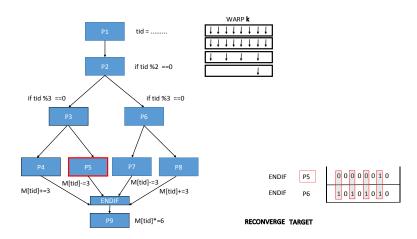




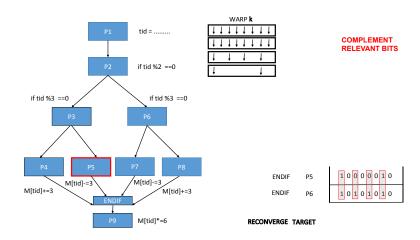




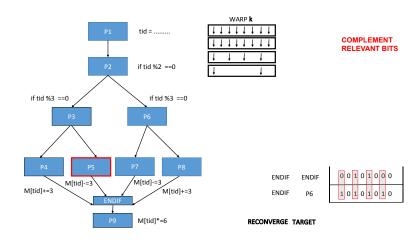




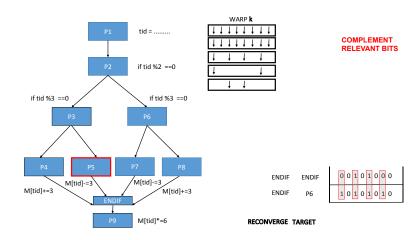




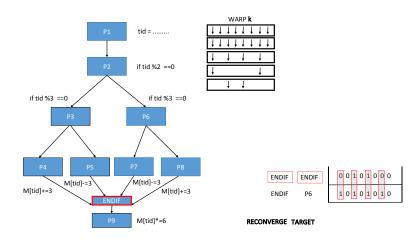




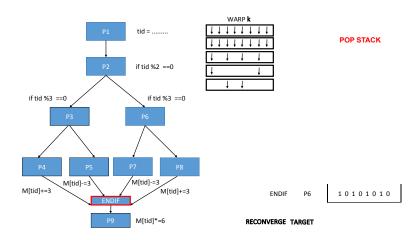




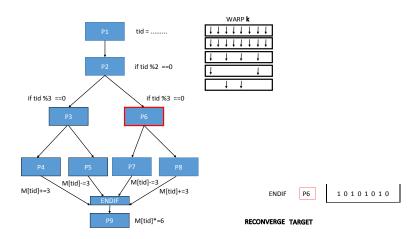




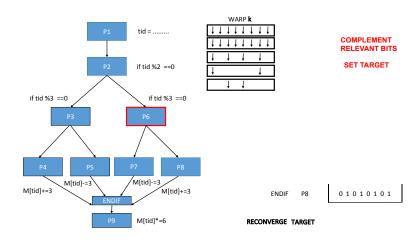




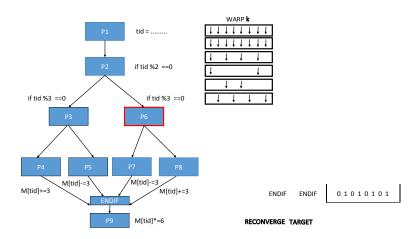




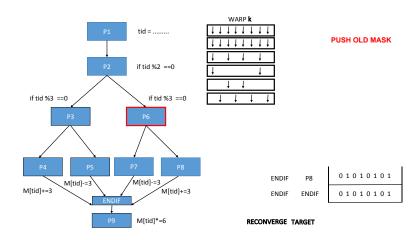




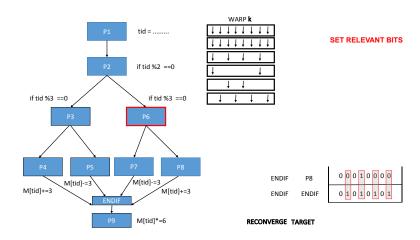




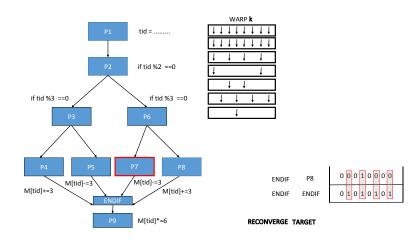




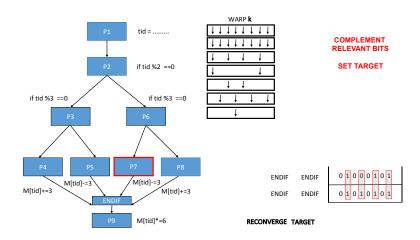




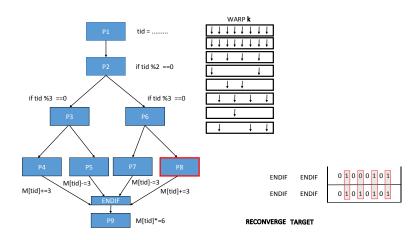




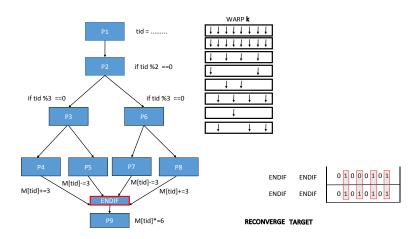




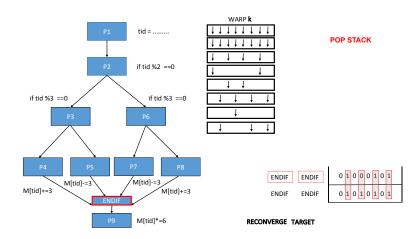




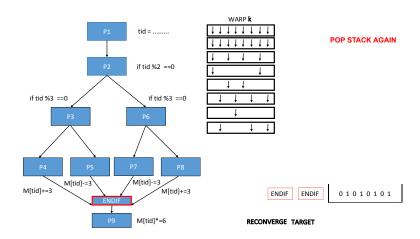




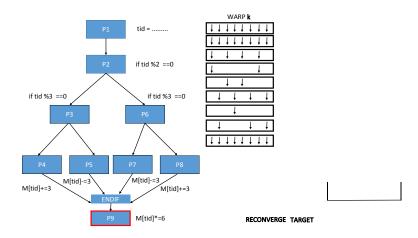














Programming tips

- GPU programmer has to be aware of hardware imposed restrictions - threads/SM, blocks/SM, threads/blocks, threads/warps
- ► The only safe way to synchronize threads from different blocks is to terminate kernel and make a fresh launch at the target synchronization point

