Introduction to Information Retrieval

Lectures 7: Dictionaries and tolerant retrieval

Dictionary

- Given a query and an inverted index, first tasks:
 - Determine whether each query term exists in the vocabulary
 - If yes, identify the pointers to the corresponding postings lists

 The vocabulary lookup operation uses a data structure called dictionary.

Dictionary data structures for inverted indexes

The dictionary data structure stores the term vocabulary, document frequency, pointers to each postings list ... in what data structure?

```
31
                                                  173
 Brutus
                         2
                              4
                                   11
                                             45
                                                        174
 Caesar
                         2
                                    5
                                                        132
                                         6
                                             16
                                                   57
                              4
Calpurnia
                        31
                             54
                                  101
```

i

dictionary

postings

A naïve dictionary

An array of struct:

term	document	pointer to
	frequency	postings list
а	656,265	\longrightarrow
aachen	65	\longrightarrow
zulu	221	\longrightarrow

char[20] int Postings *
20 bytes 4/8 bytes 4/8 bytes

- How do we store a dictionary in memory efficiently?
- How do we quickly look up elements at query time?

Dictionary data structures

- Two main choices:
 - Hashtables
 - Trees
- Some IR systems use hashtables, some trees

Hashtables

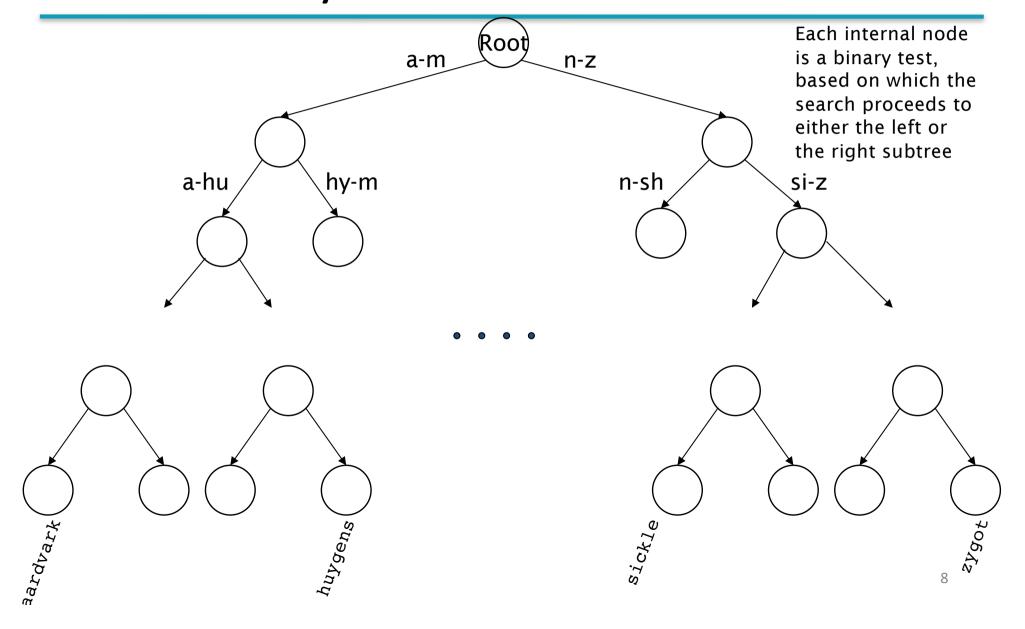
- Each vocabulary term is hashed to an integer over a large enough space that hash collisions are unlikely
 - (Assumed: you are familiar with hashtables)
- Pros:
 - Lookup is faster than for a tree: O(1)
- Cons:
 - No easy way to find minor variants (judgment/judgement) since they could be hashed to very different indexes:
 - No easy way to handle wild-card queries / prefix search
 - If vocabulary keeps growing, need to occasionally do the expensive operation of rehashing everything

Trees

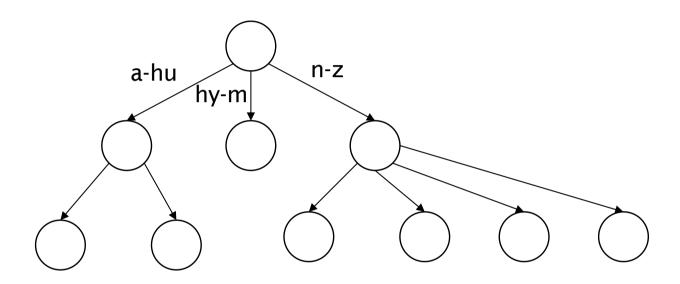
Simplest: binary tree

More usual: B-trees

Tree: binary tree



Tree: B-tree



- Definition: Every internal nodel has a number of children in the interval [a,b] where a, b are appropriate positive integers, e.g., [2,4].
- Each branch under an internal node represents a test for a range of character sequences (to guide the search)

Trees

- Trees require a standard ordering of characters and hence strings ... but we typically have one for all languages
- Pros:
 - Solves the prefix problem (terms starting with hyp)
- Cons:
 - Slower than hashtables: O(log M) [and this requires balanced tree]
 - Rebalancing binary trees is expensive
 - But B-trees mitigate the rebalancing problem

WILD-CARD QUERIES

Users may give wild-card queries

- mon*: find docs having any word beginning with "mon"
- Why wild-card queries?
 - Users may not be sure of a spelling
 - Users may want to match multiple variants of a term, e.g., query judicia* for matching both judicial and judiciary
- Types of wild-card queries
 - Trailing wild-card queries, e.g., mon*
 - Leading wild-card queries, e.g., *mon
 - General wild-card queries, e.g., s*dney, fas*in*te
- Key challenge: which dictionary terms match a wild-card query?

Trailing wild-card queries

- Trailing wild-card queries (mon*) easier
 - Use a binary tree (or B-tree) over the dictionary:
 - Retrieve all dictionary words in range: mon ≤ w < moo</p>
 - Then process postings lists of all such words

Leading wild-card queries

- *mon: find words ending in "mon": slightly harder than trailing wild-card queries
- Maintain an additional B-tree for terms backwards (reverse B-tree)
 - Each root-to-leaf path corresponds to a term in the dictionary written backwards
 - Can retrieve all words in range: nom ≤ w < non.</p>

Query processing

- At this point, we have an enumeration of all terms in the dictionary that match the wild-card query.
- We still have to look up the postings for each enumerated term, merge the postings, etc.

Exercise: from what we have discussed, how can we enumerate all terms matching the query *pro*cent*?

B-trees handle *'s at the end of a query term relatively efficiently

- How can we handle *'s in the middle of query term?
 - co*tion

- One solution: look up co* in a B-tree and *tion in a reverse B-tree, and then intersect the two term sets
 - Expensive
- Better solution: transform wild-card queries so that the *'s occur at the end
 - This gives rise to the Permuterm Index

Permuterm index

- Introduce \$ as a special symbol to mark the end of a term (a symbol that does not appear in the text)
- For term *hello*, index under:
 - hello\$, ello\$h, llo\$he, lo\$hel, o\$hell, \$hello
 - Various rotations of each term (augmented with \$) all link to the original vocabulary term
- Permuterm vocabulary: set of all rotated terms in the permuterm index

Handling wild-card queries with Permuterm index

- Given a wild-card query
 - Rotate the query so that the * symbol appears at the end of the string
 - Look up the rotated query in the permuterm index
- Queries:

```
X lookup on X$ X* lookup on $X*
```

- *X lookup on X\$* *X* lookup on X*
- X*Y lookup on Y\$X*
 X*Y*Z
 ??? Exercise!

```
Query = hel*o
X=hel, Y=o
Lookup o$hel*
```

Queries having multiple wild-cards

- E.g., fi*mo*er
- First enumerate all dictionary terms that are in the permuterm index of er\$fi*
- Not all such terms will have "mo" in the middle need to filter out mismatched terms exhaustively

Permuterm query processing

- Rotate query wild-card to the right
- Now use B-tree lookup as before
- Once the permuterm index enables us to identify the original vocabulary terms matching a wild-card query, we can look up these terms in the usual way
- Permuterm problem: ≈ quadruples lexicon size

Empirical observation for English.

k-gram (e.g., Bigram) indexes

- Enumerate all k-grams (sequence of k chars) occurring in any term
- Use \$ as a special character to denote the beginning and end of each term
- e.g., from text "April is the cruelest month" we get the 2-grams (bigrams)

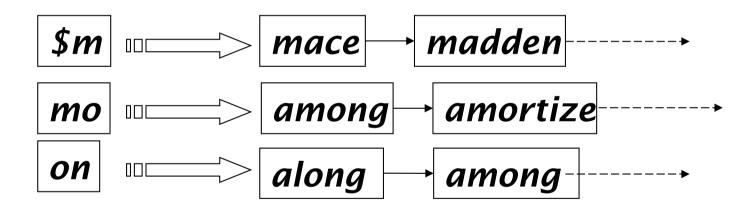
\$a,ap,pr,ri,il,l\$,\$i,is,s\$,\$t,th,he,e\$,\$c,cr,ru,ue,el,le,es,st,t\$, \$m,mo,on,nt,h\$

K-gram index

- Dictionary contains all k-grams that occur in any term in the vocabulary
- Maintain a <u>second</u> inverted index <u>from k-grams to</u> <u>dictionary terms</u> that match each k-gram
 - Each postings list points from a k-gram to all vocabulary terms containing that k-gram
- Example with k=2 shown on next slide

K-gram index (k=2) example

The k-gram index finds terms based on a query consisting of k-grams (here k=2).



Processing wild-cards with k-gram index

- Query mon* can now be run as
 - \$m AND mo AND on
- Gets terms that match AND version of our wildcard query.
- But we'd enumerate moon.
- Must post-filter these terms against query.
- Surviving enumerated terms are then looked up in the term-document inverted index.
- Fast, space efficient (compared to permuterm).

Processing wild-card queries

- As before, we must execute a Boolean query for each enumerated, filtered term.
- Wild-cards can result in expensive query execution (very large disjunctions...)
 - pyth* AND prog*
- If you encourage "laziness" people will respond!

Type your search terms, use '*' if you need to.
E.g., Alex* will match Alexander.

Recap till now

- Dictionary Data Structures:
 - Hashtables
 - B-trees

- Tolerant Retrieval
 - Wildcard queries
 - Permuterm index
 - k-gram index

SPELLING CORRECTION

Spell correction

- Two principal uses
 - Correcting document(s) being indexed
 - Correcting user queries to retrieve "right" answers
- Two main flavors:
 - Isolated word
 - Check each word on its own for misspelling: jacson
 - Will not catch typos resulting in correctly spelled words
 - e.g., from → form
 - Context-sensitive
 - Look at surrounding words,
 - e.g., I flew form Heathrow to Narita.

Document correction

- Needed for OCR'ed documents or for correcting typing errors
 - Correction algorithms are tuned for this: rn/m
 - Can use domain-specific knowledge
 - OCR confuses O and D more often than it would confuse O and I
 - O and I are adjacent on the QWERTY keyboard, so more likely interchanged in typing
- Often we don't prefer to change the documents; instead fix the query-document mapping

Query mis-spellings

- Our principal focus here
 - E.g., the query "IIT Khagarpur"
- We can either
 - Retrieve documents indexed by the correct spelling (with a declaration of the changed query), OR
 - Return several suggested alternative queries with the correct spelling
 - Did you mean ... ?

Isolated word correction

- Fundamental premise there is a lexicon from which the correct spellings come
- Two basic choices for this
 - A standard lexicon such as
 - Webster's English Dictionary
 - An "industry-specific" lexicon hand-maintained
 - The lexicon of the indexed corpus
 - E.g., all words on the web
 - All names, acronyms etc.
 - (Including the mis-spellings)

Isolated word correction

- Given a lexicon and a character sequence Q, return the words in the lexicon closest to Q
- What's "closest"?
- We'll study several alternatives
 - 1. Edit distance (Levenshtein distance)
 - 2. Weighted edit distance
 - *3. n*-gram overlap

Edit distance

- Given two strings S_1 and S_2 , the minimum number of operations to convert one to the other
- Operations are typically character-level
 - Insert, Delete, Replace, (Transposition)
- E.g., the edit distance from dof to dog is 1
 - From cat to act is 2 (Just 1 with transpose.)
 - from *cat* to *dog* is 3.
- Generally found by dynamic programming.

Weighted edit distance

- Similar to above, but the weight of an operation depends on the character(s) involved
 - Meant to capture OCR or keyboard errors
 Example: m more likely to be mis-typed as n than as q
 - Therefore, replacing m by n is a smaller edit distance than replacing m by q
 - This may be formulated as a probability model
- Requires weight matrix as input
- The dynamic programming approach can be modified to handle weights

Using edit distances for correction

- Given query, first enumerate all character sequences within a preset (may be weighted) edit distance (e.g., 2)
- Intersect this set with list of "correct" words
- Show terms you found to user as suggestions
- Alternatively,
 - Look up all possible corrections in our inverted index and return all docs ... slow
 - Retrieve with a single most likely correction
- The alternatives disempower the user, but save a round of interaction with the user

n-gram overlap

- Enumerate all the n-grams in the query string as well as in the lexicon
- Use the n-gram index (recall what we discussed in wild-card search) to retrieve all lexicon terms matching any of the query n-grams

- Variations
 - Can threshold by number of matching *n*-grams
 - Weight by keyboard layout, common typos, etc.

Example with trigrams

- Suppose the text is *november*
 - Trigrams are nov, ove, vem, emb, mbe, ber
- The query is december
 - Trigrams are dec, ece, cem, emb, mbe, ber
- So 3 trigrams overlap (of 6 in each term)
- How can we turn this into a normalized measure of overlap?

One option – Jaccard coefficient

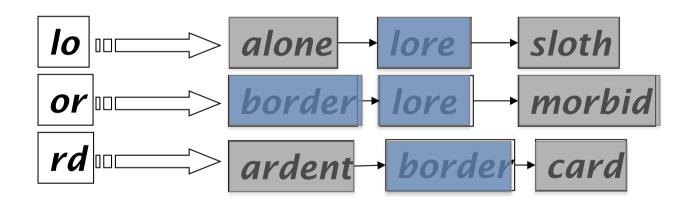
- A commonly-used measure of overlap
- Let X and Y be two sets; then the J.C. is

$$|X \cap Y|/|X \cup Y|$$

- Equals 1 when X and Y have the same elements and zero when they are disjoint
- X and Y don't have to be of the same size
- Always assigns a number between 0 and 1
 - Now threshold to decide if you have a match
 - E.g., if J.C. > 0.8, declare a match

Matching trigrams

 Consider the query *lord* – we wish to identify words matching 2 of its 3 bigrams (*lo, or, rd*)



Standard postings "merge" will enumerate ...

Adapt this to using Jaccard (or another) measure.

Context-sensitive spell correction

- Consider the phrase query "flew form Heathrow"
- We'd like the IR system to respond: Did you mean "flew from Heathrow"? because no docs matched the query phrase.
- Need surrounding context to catch this error

Context-sensitive correction

- Query: flew form Heathrow
 - Note: we do not know which word(s) is/are in error
- First idea: retrieve dictionary terms close (e.g., in terms of weighted edit distance) to each query term
- Now try all possible resulting phrases with one word "fixed" at a time
 - flew from heathrow
 - fled form heathrow
 - flea form heathrow
 - .. and so on
- Hit-based spelling correction: Suggest alternative(s) that have lots of hits (from query logs / corpus)

Exercise

 Suppose that for "flew form Heathrow" we have 7 alternatives for flew, 19 for form and 3 for heathrow.

How many "corrected" phrases will we enumerate in this scheme?

Another approach

- Break phrase query into a conjunction of biwords (we discussed in Lecture 2).
- Look for biwords that need only one term corrected.
- Enumerate only phrases containing "common" biwords.

General issues in spell correction

- We enumerate multiple alternatives for "Did you mean?"
- Need to figure out which to present to the user
 - The alternative hitting most docs
 - Query log analysis
- More generally, rank alternatives probabilistically argmax_{corr} P(corr | query)
 - From Bayes rule, this is equivalent to argmax_{corr} P(query | corr) * P(corr)



Language model

SOUNDEX

Soundex

- Class of heuristics to expand a query into phonetic equivalents
 - Language specific mainly for names :
 - E.g., chebyshev → tchebycheff
- Invented for the U.S. census ... in 1918

Soundex – typical algorithm

- Turn every token to be indexed into a 4-character reduced form
- Do the same with query terms
- Build and search an index on the reduced forms
 - (when the query calls for a soundex match)
- http://www.creativyst.com/Doc/Articles/SoundEx1/SoundEx1.htm#Top

Soundex – typical algorithm

- 1. Retain the first letter of the word.
- Change all occurrences of the following letters to '0' (zero):

- 3. Change letters to digits as follows:
- B, F, P, V \rightarrow 1
- C, G, J, K, Q, S, X, $Z \rightarrow 2$
- D,T \rightarrow 3
- $L \rightarrow 4$
- M, N \rightarrow 5
- $R \rightarrow 6$

Soundex continued

- 4. Remove all pairs of consecutive digits.
- 5. Remove all zeros from the resulting string.
- 6. Pad the resulting string with trailing zeros and return the first four positions, which will be of the form <uppercase letter> <digit> <digit> <digit>.

E.g., *Herman* becomes H655.

Will *hermann* generate the same code?

Soundex

- Soundex is the classic algorithm, provided by most databases (Oracle, Microsoft, ...)
- How useful is soundex? Not very for information retrieval
- Okay for "high recall" tasks (e.g., Interpol), though biased to names of certain nationalities
- Zobel and Dart (1996) show that other algorithms for phonetic matching perform much better in the context of IR

What queries can we process?

- We have
 - Positional inverted index with skip pointers
 - Wild-card index
 - Spell-correction
 - Soundex
- Queries such as

(SPELL(moriset) /3 toron*to) OR SOUNDEX(chaikofski)