# Análise de Redes

Slide 03 - Walks, Paths,

**Distances, Centrality** 

Prof. Patrick Terrematte



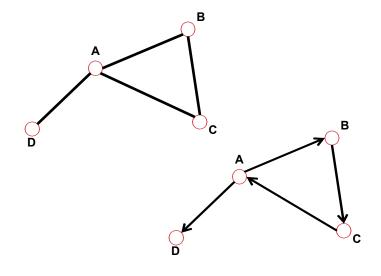
## **Análise de Redes**

- Caminhada / Walk
  - Sequência de vértices consecutivos conectados por arestas.
- Caminho / Path
  - Sequência de vértices consecutivos conectados por arestas, sem ciclos.
- Distância
  - Menor caminho entre dois vértices.



#### Caminho

- Caminho: sequência de vértices consecutivos conectados por arestas <s, u, v, ..., t>.
- Em um grafo direcionado, o caminho segue o sentido da aresta. AB ≠ BA.
- Distância (caminho mínimo, caminho geodésico): o menor caminho entre dois vértices.



#### Grafo não-direcionado

<B, C, A, D> é caminho de comprimento 3.

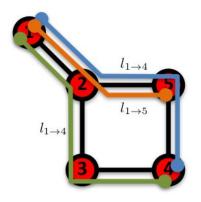
<B, A, D, C> não é caminho.

#### Grafo direcionado

<A, B, C> é caminho.

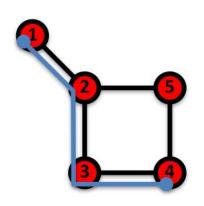
<A, C, B> não é caminho.

#### **Caminhos**



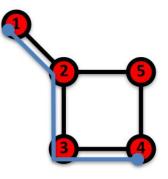
$$l_{1\to 4} = 3$$
  $l_{1\to 5} = 2$ 

**Distância**: menor comprimento entre 2 vértices (caminho mínimo).



$$l_{1\to 4}=3$$

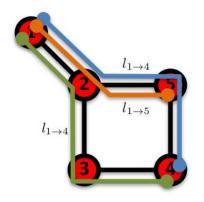
**Diâmetro**: maior distância entre quaisquer 2 vértices (maior caminho mínimo).



$$(l_{1\to 2} + l_{1\to 3} + l_{1\to 4} + l_{1\to 5} + l_{2\to 3} + l_{2\to 4} + l_{2\to 5} + l_{3\to 4} + l_{3\to 5} + l_{4\to 5})/10 = 1.6$$

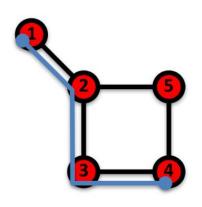
**Caminho médio**: média das distâncias entre todos os pares de vértices.

#### **Caminhos**



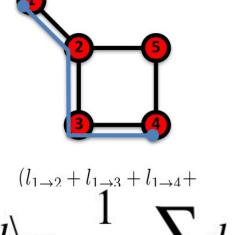
$$l_{1\to 4} = 3$$
  $l_{1\to 5} = 2$ 

**Distância**: menor comprimento entre 2 vértices (caminho mínimo).



$$l_{1\to 4}=3$$

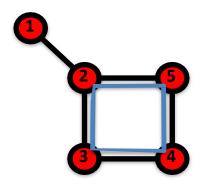
**Diâmetro**: maior distância entre quaisquer 2 vértices (maior caminho mínimo).



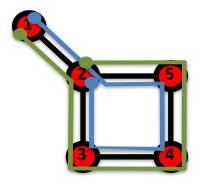
**Caminho médio**: média das distâncias entre todos os pares de vértices.

 $\max i, j > i$ 

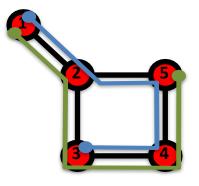
## **Caminhos**



**Ciclo**: caminho que começa e termina no mesmo vértice.

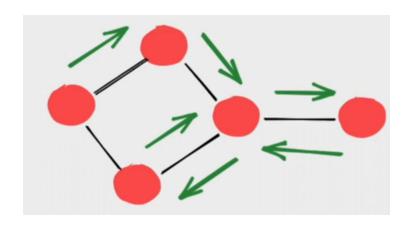


Caminho Euleriano: caminho que passa por cada aresta uma única vez.

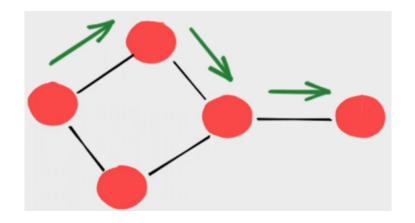


Caminho Hamiltoniano: caminho que passa por cada vértice uma única vez.

## Walk x Path



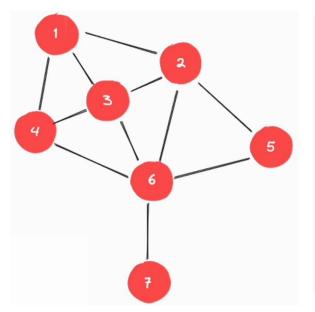
Caminhada de comprimento 6

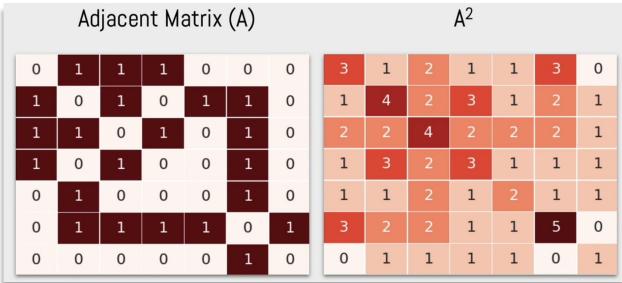


Caminho de comprimento 3



## Walk

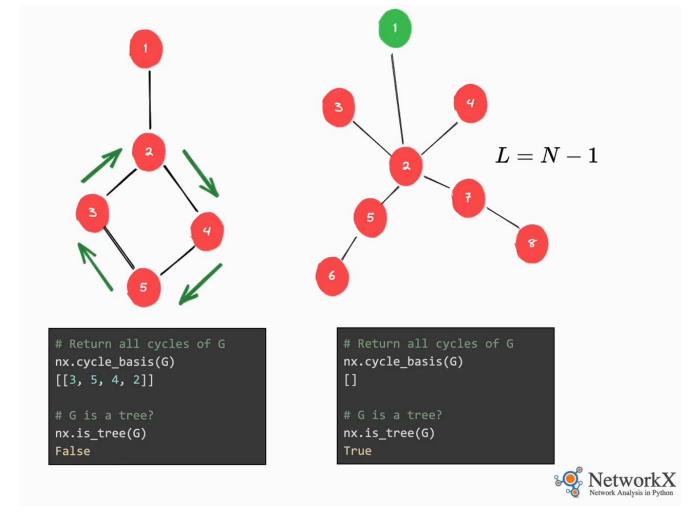




Os caminhos de comprimento 2 da rede. A diagonal representa os graus de cada nó.

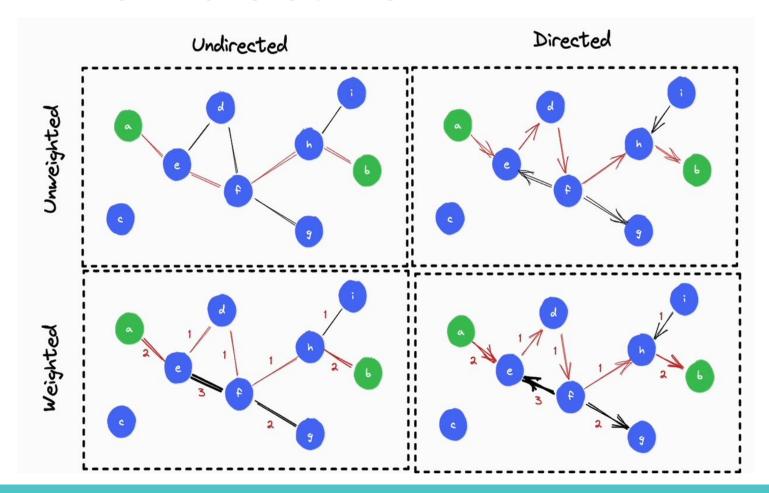


## Ciclos



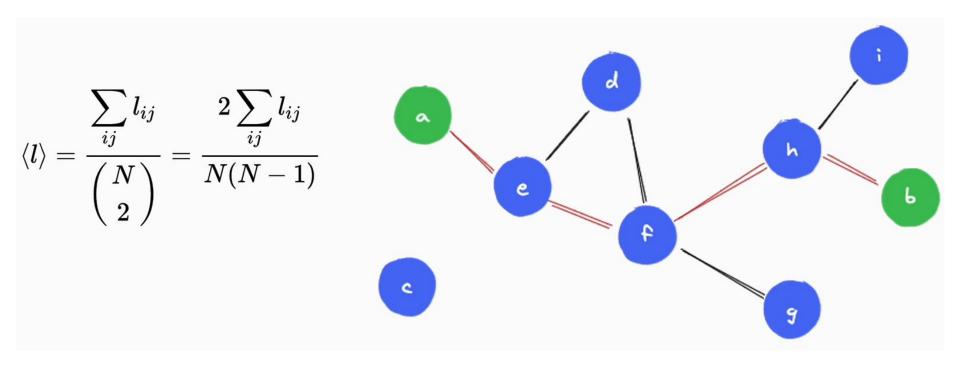


## Caminho mais curto



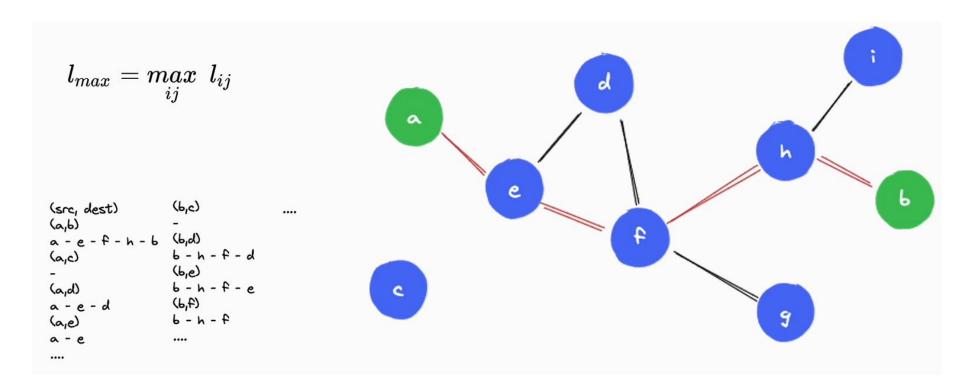


## Caminho mais curto médio





## Diametro da rede





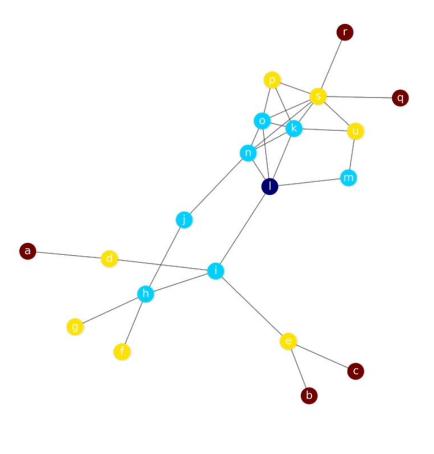
#### Diameter

The diameter of a network is the maximum eccentricity.

```
nx.eccentricity(g)
{'a': 6, 'b': 6, 'c': 6, 'd': 5,
  'e': 5, 'f': 5, 'g': 5, 'h': 4,
  'i': 4, 'j': 4, 'k': 4, 'l': 3,
  'm': 4, 'n': 4, 'o': 4, 'p': 5,
  'q': 6, 'r': 6, 's': 5, 'u': 5}

nx.diameter(g)
6

[k for k,v in nx.eccentricity(g).items()
if v == nx.diameter(g)]
['a', 'b', 'c', 'q', 'r']
```

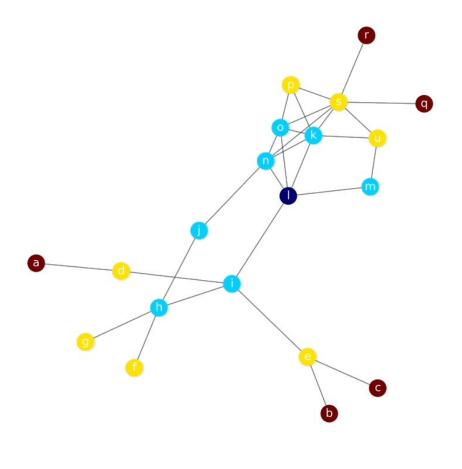




#### **Eccentricity**

It is the maximum distance from a node to all other node in the network.

```
nx.eccentricity(g)
{'a': 6, 'b': 6, 'c': 6, 'd': 5,
  'e': 5, 'f': 5, 'g': 5, 'h': 4,
  'i': 4, 'j': 4, 'k': 4, 'l': 3,
  'm': 4, 'n': 4, 'o': 4, 'p': 5,
  'q': 6, 'r': 6, 's': 5, 'u': 5}
```





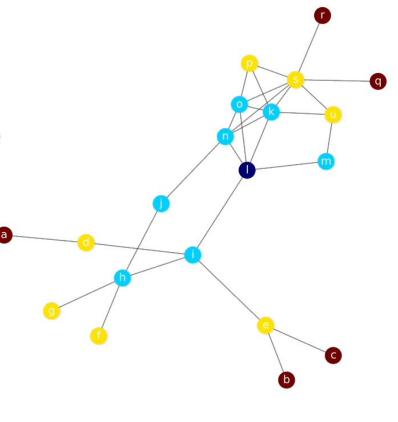
#### Periphery

The periphery of a network is a set of all nodes whose eccentricity equals the diameter.

```
nx.eccentricity(g)
{'a': 6, 'b': 6, 'c': 6, 'd': 5,
  'e': 5, 'f': 5, 'g': 5, 'h': 4,
  'i': 4, 'j': 4, 'k': 4, 'l': 3,
  'm': 4, 'n': 4, 'o': 4, 'p': 5,
  'q': 6, 'r': 6, 's': 5, 'u': 5}

nx.diameter(g)
6

nx.periphery(g)
['a', 'b', 'c', 'q', 'r']
```



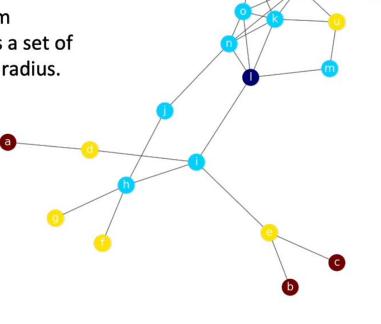


#### **Radius & Center**

The radius of a network is the minimum eccentricity. The center of a network is a set of all nodes whose eccentricity equal the radius.

```
nx.eccentricity(g)
{'a': 6, 'b': 6, 'c': 6, 'd': 5,
  'e': 5, 'f': 5, 'g': 5, 'h': 4,
  'i': 4, 'j': 4, 'k': 4, 'l': 3,
  'm': 4, 'n': 4, 'o': 4, 'p': 5,
  'q': 6, 'r': 6, 's': 5, 'u': 5}

nx.radius(g)
3
[k for k,v in nx.eccentricity(g).items()
if v == nx.radius(g)]
['l']
nx.center(g)
['l']
```



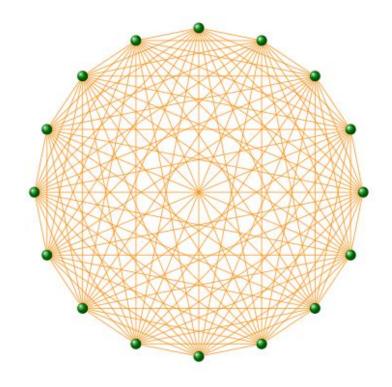


## **Grafos Completos**

- Grafo com tamanho  $L = L_{max}$  e grau médio  $\langle k \rangle = N-1$ .
- O maior número de arestas de em um grafo de ordem N:

$$L_{\text{max}} = \begin{pmatrix} N \\ 2 \end{pmatrix} = \frac{N!}{(N-2)!2!} = \frac{N(N-1)}{2}$$

- Densidade: número de arestas L em relação ao grafo completo  $L_{max}$ .
- Dado um grafo de ordem N e tamanho L.
  - Grafo esparço: L ~ N.
  - $\circ$  Grafo denso: L  $\sim$  N<sup>2</sup>.



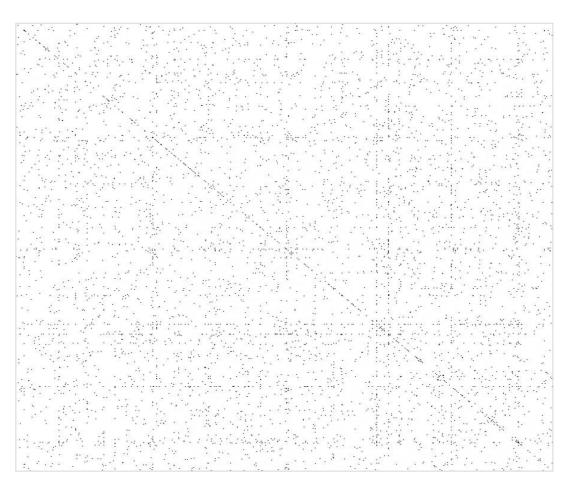
# Redes Reais são Esparsas

$$L \ll L_{max}$$
 ou  $\langle k \rangle \ll N-1$ 

	Nodes	Links	Lmax	<k></k>
WWW (sample)	325,729	1.4 x10 <sup>6</sup>	10 <sup>12</sup>	4.51
Proteina (S. cerevisiae)	1,870	4,470	10 <sup>7</sup>	2.39
Colaboração (math)	70,975	2 x10 <sup>5</sup>	3 x10 <sup>10</sup>	3.9
Atores em um filme	212,250	6 x10 <sup>6</sup>	1.8 x10 <sup>13</sup>	28.78

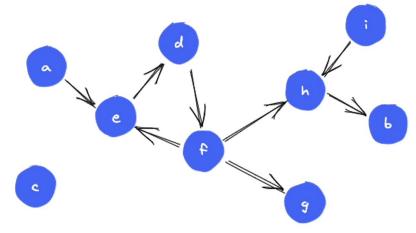
## Redes Reais são Esparsas

A matriz de adjacência da rede de interação proteína-proteína de levedura, constituída por 2.018 vértices, cada um representando uma proteína.



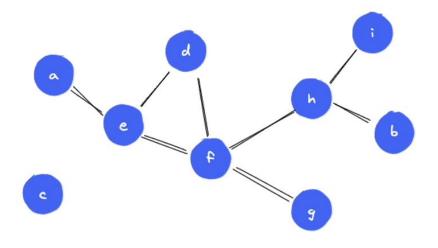
# **Componentes Conectados**

- Se dois nós não podem ser conectados por uma caminhada, então estão em diferentes componentes conectados (subgrafos).
- Componentes conectados são subgrafos com nós que podem ser atingidos pelas arestas/links da rede.
- Giant Connected Components (GCC) é um subgrafo que se destaca dos outros subgrafos pelo número de nós.





```
# G is connected or not?
nx.is_connected(G)
False
# interact under all connected component of G
for component in nx.connected_components(G):
       print(component)
{'a', 'i', 'e', 'g', 'h', 'd', 'b', 'f'}
{'c'}
# how many connected components has G?
nx.number_connected_components(G)
2
# which connected component is a node N?
nx.node_connected_component(G,"a")
{'a', 'b', 'd', 'e', 'f', 'g', 'h', 'i'}
```







# Strongly Connected Components (SCC) Weakly Connected Components (WCC)

Um componente é fortemente conectado, quando em um subgrafo todos os nós se acessam.



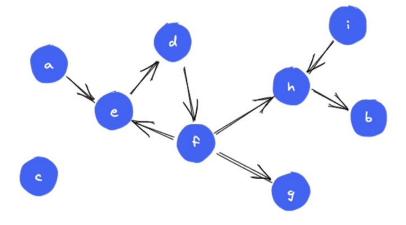
```
nx.is_strongly_connected(G)
False

nx.is_weakly_connected(G)
False

list(nx.weakly_connected_components(G))
[{'a', 'b', 'd', 'e', 'f', 'g', 'h', 'i'}, {'c'}]

list(nx.strongly_connected_components(G))
[{'b'}, {'h'}, {'g'}, {'d', 'e', 'f'}, {'a'}, {'i'}, {'c'}]

nx.number_strongly_connected_components(G)
7
```





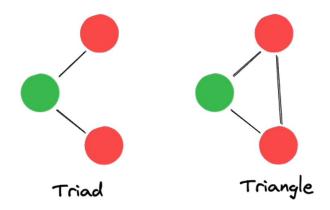


#### Análise de Redes - Colab



http://colab.research.google.com/github/terrematte/network\_analysis/blob/main/notebooks/02\_hubs\_path\_and\_structures.ipynb

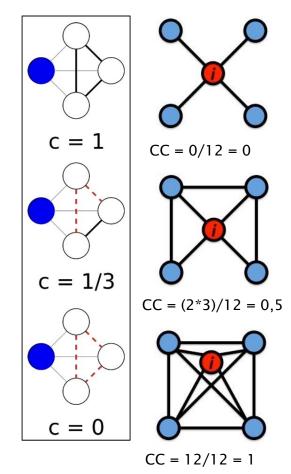




## Coeficiente de Clusterização

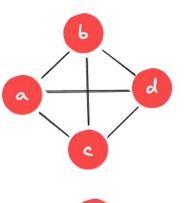
- O coeficiente de clusterização é a fração de todos possíveis triângulos que contém aquele nó.
- Mede o quanto os nós se relacionam.
- CC n\u00e3o expressa uma propriedade do v\u00e9rtice e sim dos seus vizinhos!

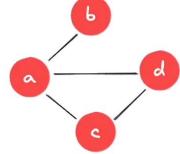
$$C(i) = rac{ au(i)}{ au_{max}(i)} = rac{ au(i)}{\left(egin{array}{c} k_i \ 2 \end{array}
ight)} = rac{2 au(i)}{k_i(k_i-1)} \hspace{1cm} C = rac{\sum\limits_{i;k_i>1} C(i)}{N_{k>1}}$$



## Coeficiente de Clusterização

```
nx.triangles(G)
{'a': 3, 'b': 3, 'c': 3, 'd': 3}
nx.clustering(G,"a")
1.0
nx.clustering(G)
{'a': 1.0, 'b': 1.0, 'c': 1.0, 'd': 1.0}
nx.average_clustering(G)
G.remove_edge("b","c")
G.remove_edge("b","d")
nx.clustering(G)
nx.average_clustering(G)
0.5833333333333333
```

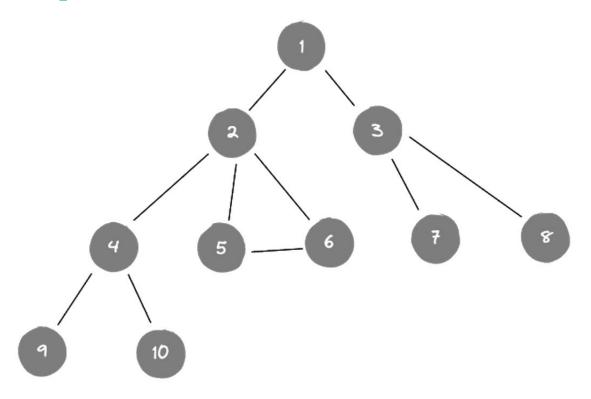








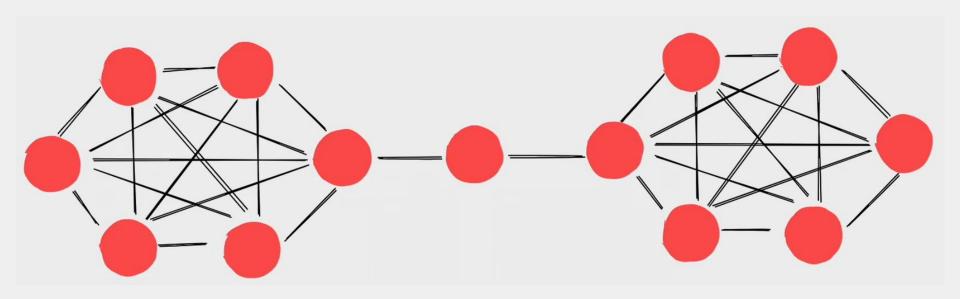
# Breadth First Search (BFS) versus Depth First Search (DFS)





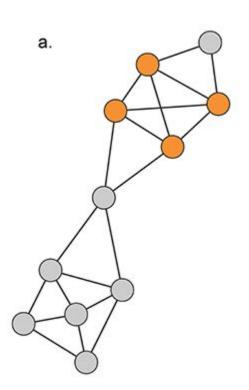
Key	Breadth First Search (BFS)	Depth First Search (DFS)
Definition	BFS stands for Breadth First Search.	DFS stands for Depth First Search.
Data structure	Uses a Queue to find shortest path. FIFO (First In First Out)	Uses a Stack to find the shortest path. LIFO (Last In First Out)
Source	BFS is better when target is closer to Source.	DFS is better when target is far from source.
Suitability for decision tree	As BFS considers all neighbor so it is not suitable for decision tree used in puzzle games.	DFS is more suitable for decision tree. As with one decision, we need to traverse further to augment the decision. If we reach the conclusion, we won.
Speed	BFS is slower than DFS.	DFS is faster than BFS.
Time Complexity	Time Complexity of BFS = O(V+E) where V is vertices and E is edges.	Time Complexity of DFS is also O(V+E) where V is vertices and E is edges.
Memory	BFS requires more memory space.	DFS requires less memory space.

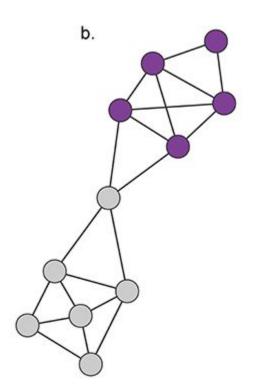
## Como medir a importância de um nó?

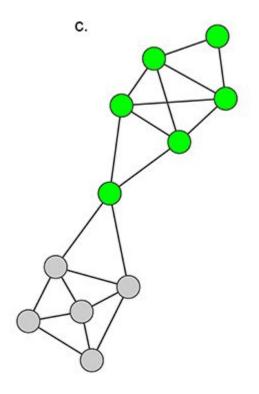




## Como medir a importância de um nó?







### Degree Centrality

Number of connections

#### Closeness centrality

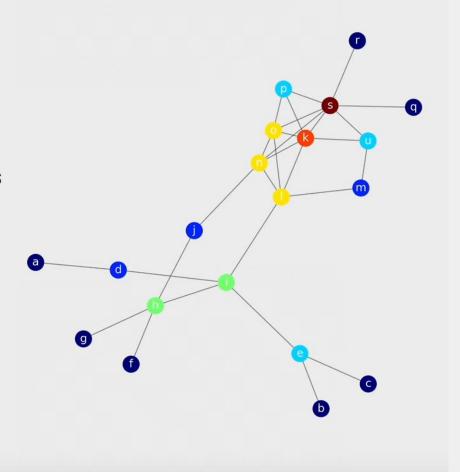
Average distance to all other vertices

#### Betweenness Centrality

Position on the shortest path

#### Eigenvector Centrality

Authority score based on the score of the neighbors





### Degree Centrality

Number of connections

### Closeness centrality

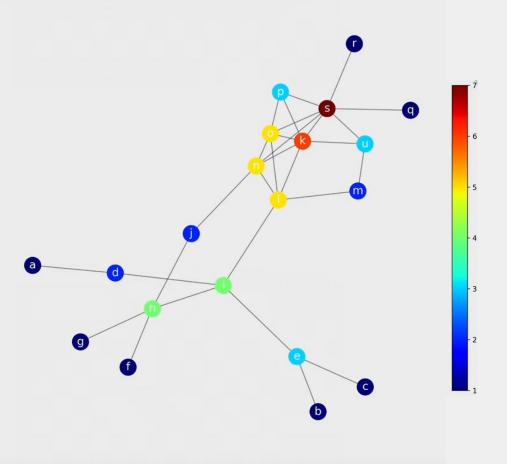
Average distance to all other vertices

### Betweenness Centrality

Position on the shortest path

#### Eigenvector Centrality

Authority score based on the score of the neighbors

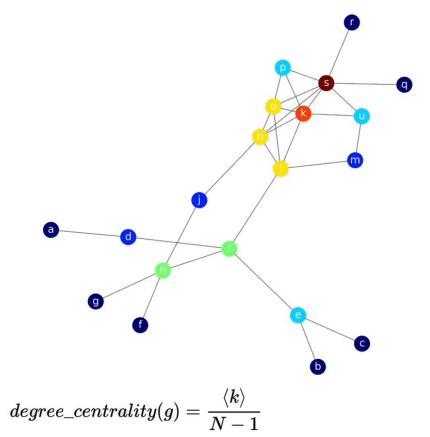




#### Degree Centrality

#### Number of connections

```
nx.degree centrality(g)
{'a': 0.05263157894736842,
 'b': 0.05263157894736842,
 'c': 0.05263157894736842,
 'd': 0.10526315789473684,
 'e': 0.15789473684210525,
 'f': 0.05263157894736842,
 'g': 0.05263157894736842,
 'h': 0.21052631578947367,
 'i': 0.21052631578947367,
 'j': 0.10526315789473684,
 'k': 0.3157894736842105,
 '1': 0.2631578947368421,
 'm': 0.10526315789473684,
 'n': 0.2631578947368421,
 'o': 0.2631578947368421,
 'p': 0.15789473684210525,
 'q': 0.05263157894736842,
 'r': 0.05263157894736842,
 's': 0.3684210526315789,
 'u': 0.15789473684210525}
```



Degree of the node / Number of all other nodes.



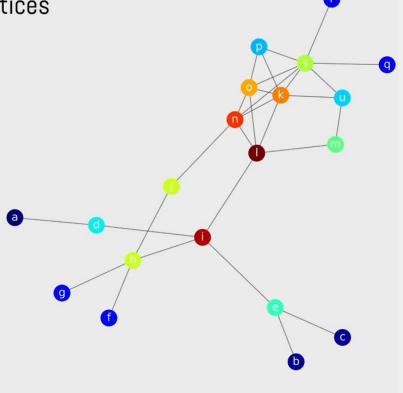
## Closeness Centrality

Average distance to all other vertices

Another way to measure the centrality of a node is by determining how "close" it is to the other nodes. This can be done by summing the distances from the node to all others.

$$closeness\_centrality(g,i) = rac{N-1}{\displaystyle\sum_{j 
eq i} l_{ij}}$$

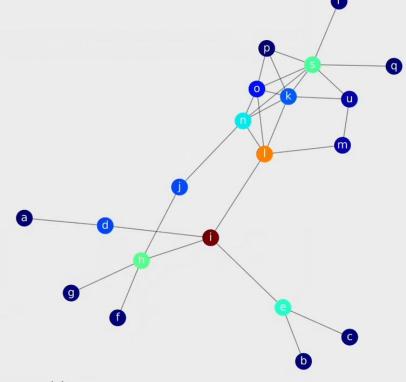
nx.closeness\_centrality(g)



Betweenness Centrality

Position on the shortest path

Many phenomena taking place in networks are based on diffusion processes. Examples include the transmission of information across a social network, the traffic of goods through a port, and the spread of epidemics in the network of physical contacts between the individuals of a population. This has suggested a third notion of centrality, called betweenness: a node is the more central, the more often it is involved in these processes.



 $betweenness\_centrality(g,i) = \sum_{h 
eq j 
eq i} rac{\sigma_{hj}(i)}{\sigma_{hj}}$ 

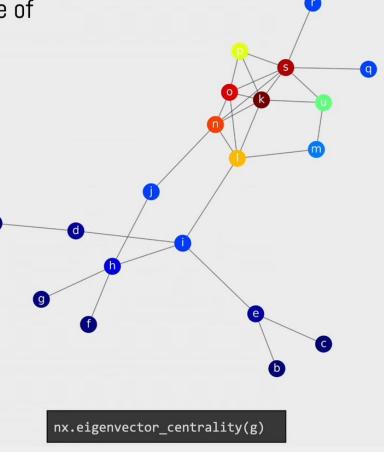
nx.betweenness\_centrality(g)

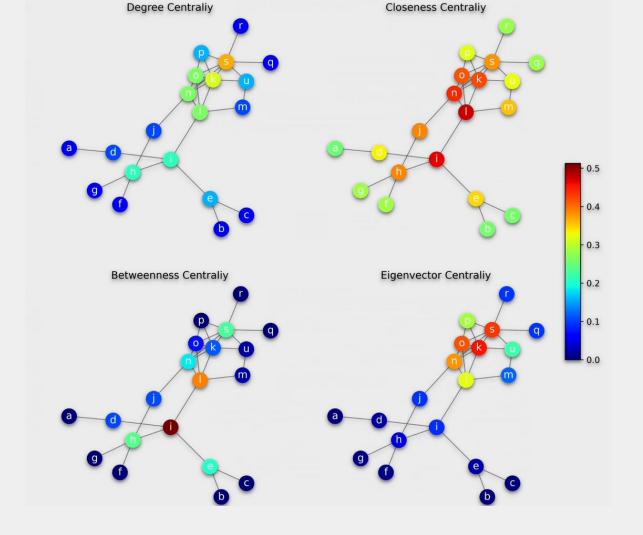


## Eigenvector Centrality

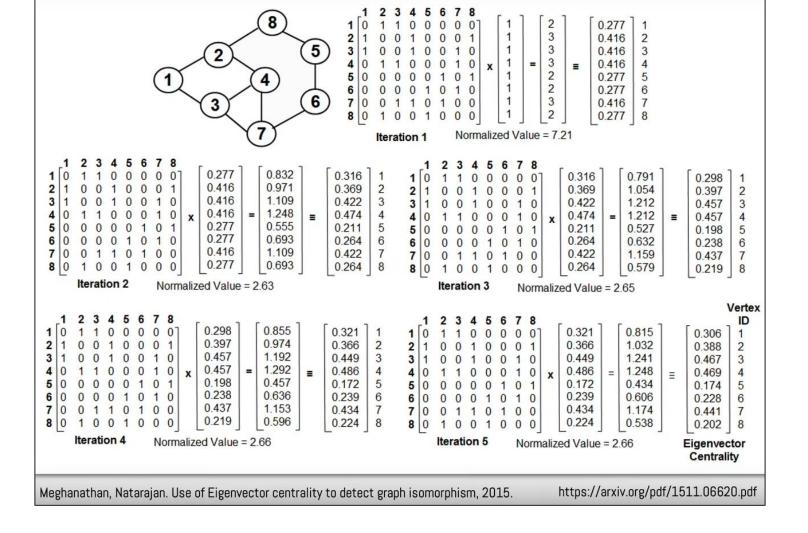
Authority score based on the score of the neighbors

In many circumstances a node's importance in a network is increased by having connections to other nodes that are themselves important. For instance, you might have only one friend in the world, but if that friend is the CR7 then you yourself may be an important person. Thus centrality is not only about how many people you know but also who you know.











#### Referências

- Network Science by Albert-László Barabási <u>http://networksciencebook.com/</u>
- <u>Network Analysis Course Prof. Ivanovitch (DCA/UFRN)</u>
   <u>YouTube Playlist</u>



## **Exercícios**

https://rosalind.info/problems/topics/graph-algorithms/

https://rosalind.info/problems/topics/graphs/

