

Simplify the following regular expression:

$$(aa^*(aa^*b + b) + b + aa^*b + \varepsilon)(aa^*(aa^*b + b))^*$$

Let's call $aa^* = P$: $((P(Pb + b) + \varepsilon) + (Pb + b)) \cdot (P(Pb + b))^*$

Use distribution law: $(P(Pb + b) + \varepsilon) \cdot (P(Pb + b))^* + (Pb + b) \cdot (P(Pb + b))^*$
 $(L + \varepsilon) \cdot L^* = L^* \Rightarrow (P(Pb + b))^* + (Pb + b) \cdot (P(Pb + b))^*$

$Pb + b = aa^*b + b = (aa^* + \varepsilon) \cdot b = a^*b$ by distribution law.

$\Rightarrow (aa^*(a^*b))^* + a^*b \cdot (aa^*(a^*b))^* = (aa^*b)^* + a^*b \cdot (aa^*b)^* \quad (a^*a^* = a^*)$

By distribution law: $(a^*b + \varepsilon) \cdot (aa^*b)^*$

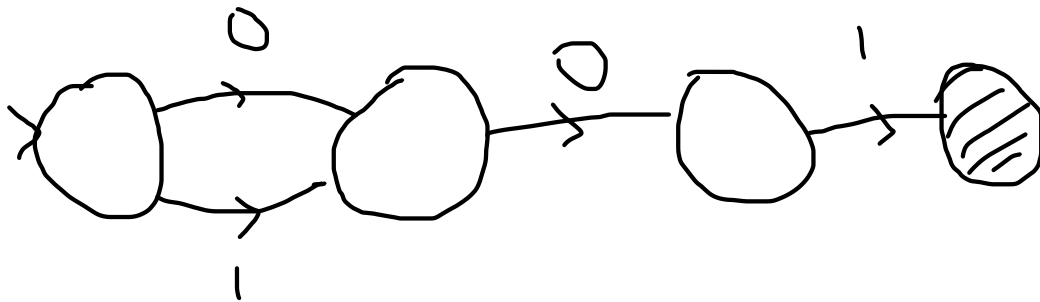
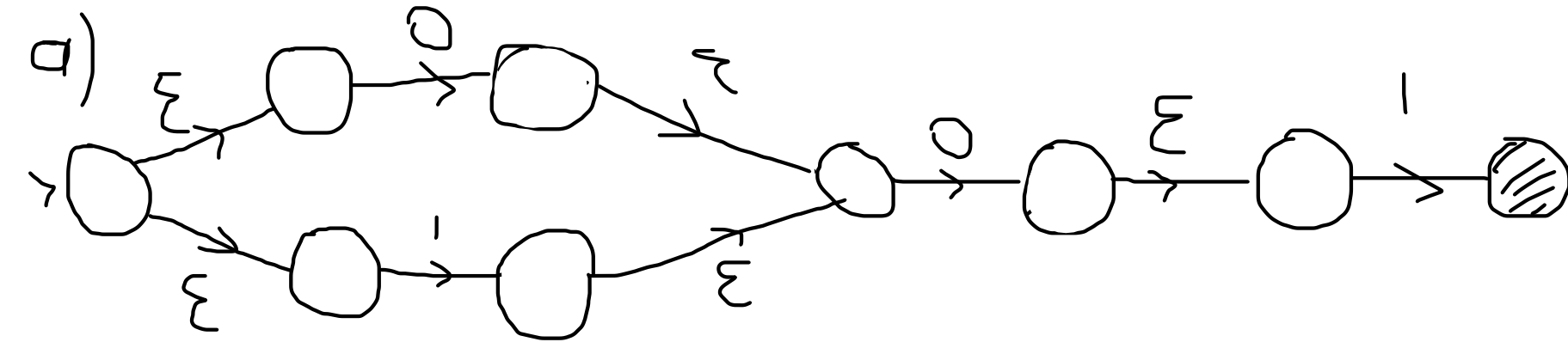
$a^* = (aa^* + \varepsilon)$: $a^*b = (aa^* + \varepsilon) \cdot b = aa^*b + b \Rightarrow (aa^*b + \varepsilon + b) \cdot (aa^*b)^*$
 $= (aa^*b + \varepsilon) \cdot (aa^*b)^* + b \cdot (aa^*b)^* = (aa^*b)^* + b \cdot (aa^*b)^* \quad (\text{by } (L + \varepsilon) \cdot L^* = L^*)$

By distribution law: $\boxed{(\varepsilon + b) \cdot (aa^*b)^*}$

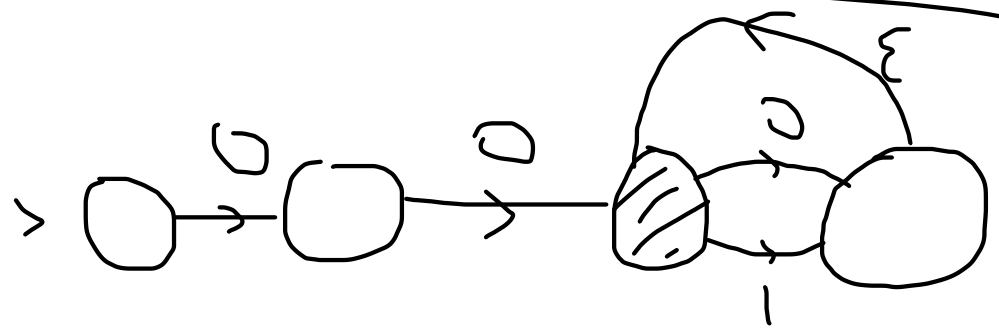
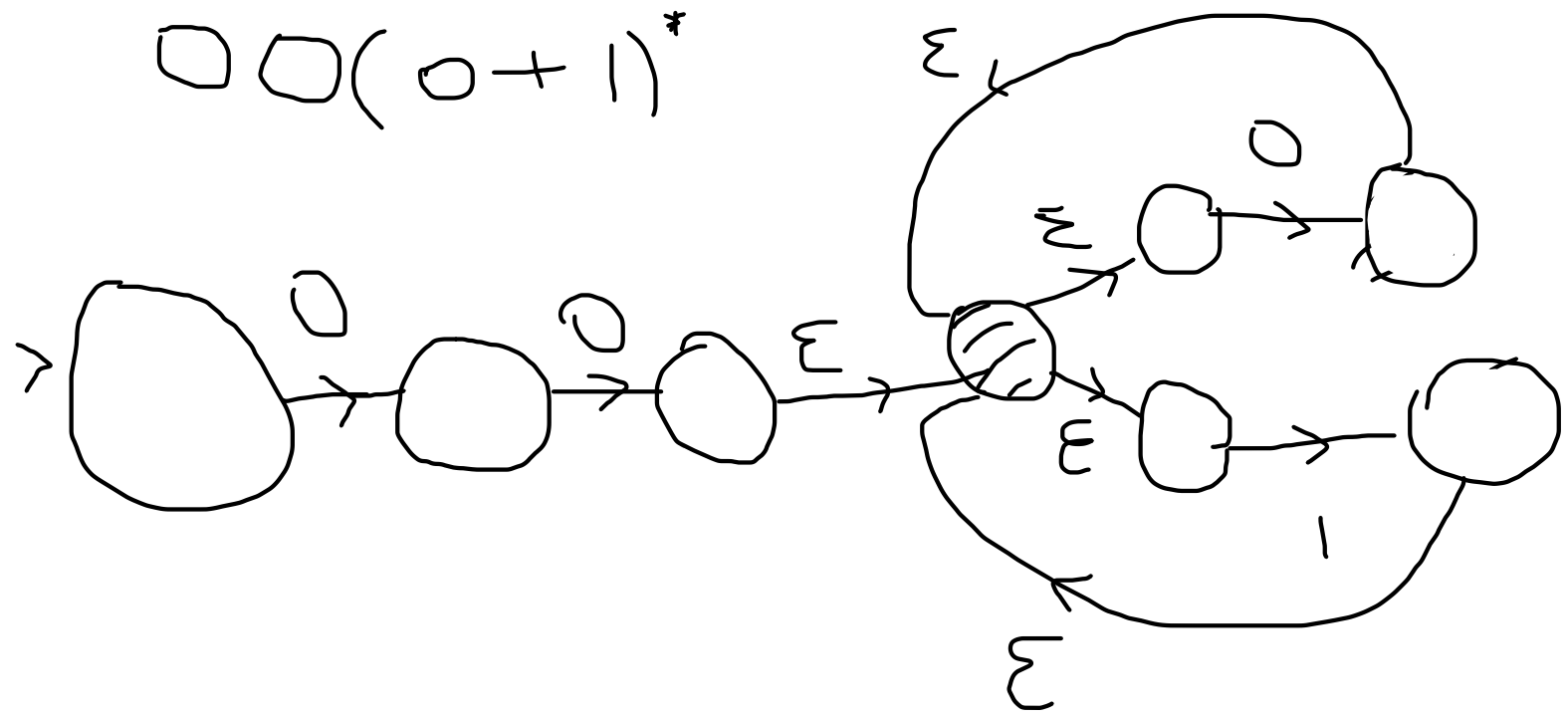
Convert the following regular expressions to ϵ -NFA's:

a) $(0 + 1)01$

b) $00(0 + 1)^*$

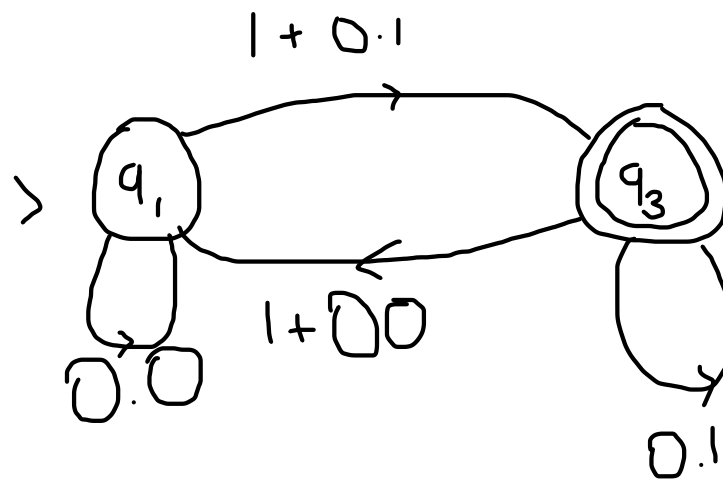
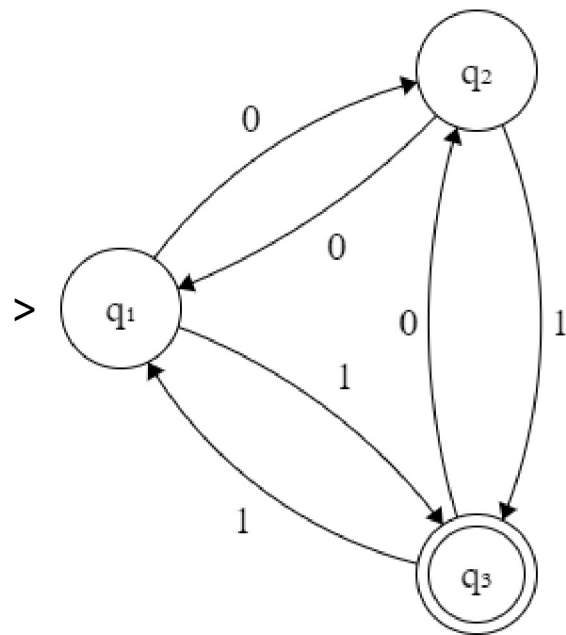


$$00(0+1)^*$$



Give a regular expression for the following DFA using:

- state elimination
- inductive formula on R_{ij}^k



$$R \in : \left((00)^* \cdot (1+01) \cdot (01)^* \cdot (1+00) \right)^* \cdot \left((00)^* \cdot (1+01) \cdot (01)^* \right)$$

$$b) R_{ij}^k = R_{ij}^{k-1} + R_{ik}^{k-1} \cdot (R_{kk}^{k-1})^* \cdot R_{kj}^{k-1}$$

$$R_{13}^3 = R_{13}^2 + R_{13}^2 \cdot (R_{33}^2)^* \cdot R_{33}^2$$

The rest is exercise.

Prove whether the language $L = \{a^m b^n c^k \mid k = m + n\}$ is regular or not.

Pumping lemma: n

$$w = a^n b^n c^{2n} \quad |w| \geq n$$

$$w = x \cdot y \cdot z$$

$$\textcircled{1} |xy| \leq n$$

$$\textcircled{2} |y| > 0$$

$$\textcircled{3} xy^i z \in L, \quad i = 0, 1, 2, \dots$$

$$xy^0 z \notin L$$

NOT REGULAR!

$$xy = a^p, \quad p \leq n$$

$$y = a^q, \quad p \geq q > 0$$

$$x = a^{p-q}$$

$$z = a^{n-p} b^n c^{2n}$$

$$xz \in L$$
$$a^{p-q} a^{n-p} b^n c^{2n} = a^{n-q} b^n c^{2n}$$
$$2n - q \neq 2n$$

Prove whether the language $L = \{w \in \{0, 1\}^* \mid w \text{ has an equal number of occurrences of "01"s and "10"s}\}$ is regular or not.

