LéPiX

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https://github.com/ThePhD/lepix

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Introduction

LéPiX is a small, general-purpose programming language whose goal is to make working with parallel computation and multidimensional arrays simple. Featuring a preprocessor, namespacing, sliced multidimensional array syntax, parallel blocks based on an invocation count and ID, and bottom-up type derivation (automatic type deduction), the goal of the language is to produce an environment that is heavily statically checked and ensures a degree of correctness the user of the language can rely on.

However, this implementation of LéPiX specifically departed from some of the original design goals due to time constraints and team issues. Therefore, while parallelism and arrays were on the table, neither made it into the implementation in the given time frame I had to put the rest of the language (entire semantic analyzer plus all of the codegen) (about 2 weeks, plus postmortem time). The implementation here instead set out to demonstrate how 4 techniques can be achieved with the language:

- 1. Namespaces namespacing as defined in the original language, but lacking using statements inside definition blocks
- 2. Overloading having multiple functions assigned to the same name, separated internally by a name-mangling scheme based on arity and arguments.
- 3. Bottom-up Type Derivation deduction of return types for functions based on the expression of returns (or none therein).

4. Call targets as expressions - DICE¹ and other languages – including our own, at first – implemented function calls as an identifier plus necessary function all syntax. This implementation of LéPiX treats it as an expression, which presents unique challenges for the previous 2 goals.

1.1 Language Proposal

Below is our initial language proposal, in its entirety. It was very ambitious, and Professor Edwards told us to scale our goals back considerably. From it, we threw out GPU code generation in exchange for Parallelism and code generation for LLVM IR alone, and also wanted to focus on syntax for multidimensional arrays.

 $^{^{1}} http://www.cs.columbia.edu/\ sedwards/classes/2015/4115\text{-}fall/reports/Dice.pdf$

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An Overview

1.1 Introduction

Heterogeneous computing and graphics processing is an area of intense research. Many existing solutions – such as C++ AMP [5] and OpenCL [3] – leverage the power of an existing language and add preprocessors and software libraries to connect a user to allow code to be run on the GPU. Due to its massively parallelizable nature, code executed on the GPU can be orders of magnitude faster, but comes at the cost of having to master a specific programming library and often learn new framework-specific or platform-specific language subsets in order to compute on the GPU.

1.2 Enter LéPiX

We envision LéPiXto be a graphics processing language based loosely on a subset of the C language. Using an imperative style with strong static typing, we plan to support primitives that enable quick and concise programs for image creation and manipulation. The most novel feature we have planned for the language is the ability to compile to both the CPU as well as the GPU. The reason is to enable high performance and ease the pain most notably found with trying to write applications which leverage the power of the GPU at the cost of a steep learning curve for explicit non-CPU device computation APIs, e.g. OpenGL Compute Shaders, DirectCompute, OpenCL, CUDA, and others. The final goal is to enable the writing of computer vision and computer graphics algorithms in LéPiX with relative ease compared to other languages.

The LéPiX Language

2.1 Language

The language itself is meant to follow loosely from imperative C, but it subject to change as we refine our desired set of primitives and base operations. What follows is a loose definition of the primitive types, operations we would like to implement to get a baseline for the language, and how would would like to put those together syntactically and grammatically. All of these definitions will be mostly informal.

The goal of LéPiX is to provide a strongly and statically typed language by which to perform image manipulation easily. The hope is that powerful algorithms can be expressed in the language by providing a useful set of basic types, including the concept of a pixel and a natively slice-able matrix type that will serve as the basis for an image.

- Primitive Data Type: the core types defined by the language itself and given a set of supportable operations
- Built-ins: some of the built-in functions to help users
- Function Definitions: how to define a function in the language and use it
- Operators: which operators put into the language and operate on the primitive types

2.1.1 User Defined Types

Support for user-defined types is planned, but will be a stretch goal (5.1). We want to support having user-defined types that can be used everywhere,

and for it to be able to overload the conceivable set of all operations.

2.1.2 Syntax

This is a basic guide to the syntax of the language. We are striving to develop a C-Like imperative language. It will follow many of the C conventions, but with some differences that we think will better fit the domain we are striving to work within. As we do not have a formal grammer just yet, we will present potential programs that we wish to allow to generate appropriate code. You can find these example programs in 3. Below are some quick points about the LéPiX language.

Namespacing As a stretch goal, LéPiX will attempt to support namespacing, to avoid record collision problems as present in OCaml (without the use of modules) and to formalize the good practice of of prepending the short name of the module / library to all functions in C code. Other languages have explicit support. LéPiX will attempt to encourage code sharing and reuse by including the use of namespacing.

Keywords The following words will be reserved for use with the language: namespace, struct, class, typename, typedef, for, while, break, if, elseif, else, void, unit, int, float, uint, pixel, image, vec, vector, mat, matrix. In addition, all identifiers containing __ (two underscores) are reserved for the use of the compiler and the standard library.

The standard library reserves the usage of the namespace lib. The language will place intrinsics and built-ins within the lpx namespace.

Function Definitions Typical function definitions will follow a usual C-style syntax. An example in pseudo-lexer code:

code/func-def.lex

Ideally, we would like the order of function definition not to matter, so long as it appears *somewhere* in the whole source code listing. Early versions of the LéPiX compiler might require definition before use, for sanity purposes.

As a stretch goal, there are plans for lambda functions (anonymous function values) to be generated by a much more terse syntax.

Control Flow Control flow will follow a C-style syntax as well. An example in pseudo-lexer code:

code/flow-control.lex

Operations The LéPiX language will support most of the basic mathematical operators. Other operators will be provided VIA functions. These include:

- Mathematical plus (+), minus (-), multiply (*), divide (/), modulus (%), power of (**)
- Logical and (&&), or (||), less than (<), greater than (>), equal to (==), not equal to (!=) negate (!)

Some of these will use both the symbol and the name, such as "not" for negation. Support for bitwise operations, such as left shift and right shift as well as bitwise and / bitwise or, will come as a stretch goal, depending on whether we can handle these basics. Ternary conditionals may also prove useful, but will not be immediately supported.

Comments Single-line comments will begin with //. Multi-line and *nestable* comments will begin with /**/ will be supported as well.

2.1.3 Built-ins

Some useful built-in functions that will be provided with the language:

• Trigonometric functions: lib.sin, lib.cos, lib.tan, lib.asin, lib.acos, lib.atan, lib.atan2

- Power Functions: lib.sqrt, lib.cbrt, lib.pow
- Exponential Functions: lib.expe, lib.exp2, lib.loge, lib.log2, lib.log10

All degree arguments will come in radians. Other functions that operate on built-in types will be provided as library functions, to allow for replacement if necessary.

2.1.4 Primitive Types

Primitives are integral types and multi-dimensional array types. Vectors, Matrices, and Pixels are all subsets of N-dimensional array types. String will be presented as a built-in type, but may be implemented either as a built-in or just an always-included library type. The purpose of string will be specifically to handle reading in and writing out from the file system, as that is the only way to handle such a case. See 2.1 for details.

Table 2.1: Primitive Types in LéPiX

Type	variants	Purpose
void	nothing	built in single-value / empty type
int#	# can be $8/16/32/64$ (defaults to 32 without name)	signed integral type
uint#	# can be $8/16/32/64$ (defaults to 32 without name)	unsigned integral type
float#	# can be $16/32/64$ (defaults to 32 without name)	floating type
<type name="">[< optional #>]</type>	array type; can be multidimensional by adding [<optional #="">]</optional>	basic primitive that will help us represent an image in its decoded form; can be sliced to remove 1 dimension
vec# <type name=""></type>	aliases for 1-dimensional fixed-size arrays	base type for pixels
mat# <type name=""></type>	aliases to 2-dimensional fixed-size arrays	can be sliced into vectors
string	utf8-encoded string; potentially more later	primarily, to address the file system; not really interested in complex text handling; basically treated as an array of specially-typed uint8
pixel	rgba (red-green-blue-alpha); hsv (hue saturation value)	in the future, this will be something that will need to be customizable to support varying image types of different bit depths

Examples

The following are some examples of programs that we would like to be written in LéPiX. This does not reflect final syntax and is mostly based on equivalent or near-equivalent C-style code:

```
1 function image red(int width, int height) {
       // Create an image of a specific width / height
       image ret = image(width, height);
       for (int i = 0; i < width; i++) {
             for (int j = 0; j < height; j++) {
                  // r, g, b creation
                  ret[i][j] = pixel(1,0,0);
             }
       }
9
10
       return ret;
11
12 }
13
14 function int main () {
       image img = red(img);
       lib.save(img);
17
       return 0;
18 }
```

Listing 3.1: red.lpx

The red example shows up some simple conditionals in a for loop to write all-red to an image. It is not the most exciting code, but it has a lot of moving parts and will help us test several parts of the library, from how to call functions to basic iteration techniques.

```
pixel[] toprow = img[ri];
4
             // 0-based indexing
             pixel[] bottomrow = img[img.height - ri - 1];
             for (int ci = 0; ci < img.width; ++ci) {
                  // generic swap call in library
                  lib .swap(toprow[ci], bottomrow[ci]);
9
             }
10
       }
11
12
13
  function int main () {
14
       image img = lib.read("meow.png");
15
        flip (img);
16
        lib . save(img);
17
18
       // implicit return 0:
19
       // we want this to be able to work with
20
21
        // command-line environments as well
22 }
```

Listing 3.2: flip.lpx

This above code is a bit more complicated. It shows that we can save a slice of a matrix's (image's) row, operate on it, and even call the library function lib.swap on its pixel elements. It also demonstrates a string literal, and passing it to the lib.read function to pull out a regular image from a PNG, and then saving that same image. It also shows off an implicit return 0 (we expect our programs to be run in the context of a shell environment, and to play nice with the existing C tools in that manner).

Codegen

4.1 LLVM IR

The current goal for generating the code for this language is to use LLVM and serialize to LLVM IR. This will allow our language to work on a multiple of platforms that LLVM supports, provided we can successfully connect our AST / DAG with the our code generator.

4.2 SPIR-V

For the GPU, we still want to compile to LLVM IR. But, with the caveat that we make it work to push out to SPIR-V code using the Khronos LLVM i-i-i-SPIR-V Bidirectional Translator[1]¹. The good news is that everything from our source code processing steps to our DAG / AST generation can be done in OCaml. However, SPIR-V is new and OCaml bindings for this relatively new project are not something that is quite established: it is conceivable that our code generator will be written in C++ as opposed to any other language, simply because of the library power behind what is already present is written in C and C++, and interfacing that with OCaml might be exceptionally difficult. Granted, we could also write an OCaml Code Generator for SPIR-V from scratch, but this does not seem like a prudent use of our time.

It is also very important to note that the LLVM IR $_{i-\-\-\-\-\-}$ SPIR-V Translator produces SPIR-V, which by itself cannot be run on anything. We would need a C or C++ compiled Vulkan Driver to take that bit of our program and run it in SPIR-V land. Furthermore, many operations – such as data reading –

¹Available here: https://github.com/KhronosGroup/SPIRV-LLVM

are not instructions we can exactly slot into SPIR-V code, and would require a bootstrapper of some sort nontheless.

Stretch Goals

5.1 User Defined Types

User-defined structures are a stretch goal of this project. The core idea is that if we can manage to create pixel, vec#*, mat#*, etc. types using the language, we would be able to simply make this kind of functionality available to users. Currently, we plan to hard code these types in at the moment, however.

5.2 Namespacing

Similar to user defined types, namespacing allows an element of organization to be brought to written code. Currently, we are going to hard code built ins to the lib and lpx namespaces.

5.3 Named Parameters

This is an entirely fluff goal to make it easier to call certain functions. The idea is that arguments not yet initialized by the ordered list of arguments to a function call can be specified out-of-order – as long as others inbetween are defaulted – by passing a name=(expression) pairing, separated by commas like regular function arguments.

5.4 Lambda Functions

As mentioned in 2.1.2, we would like to support lambdas as a way of definition functions. Currently, we do not know what the most succinct and terse syntax for our language would be.

5.5 Standard Library Implementation

It would be nice to fill out a standard library implementation, to vet the LéPiX compiler. Candidates would include some basic functions in the lib namespace for manipulation of the image type.

5.6 Movies: Encoding

If we can have built in types for image and the like, then LéPiX could theoretically handle movies by presenting to the user frames of data in sequential order. Doing this is orders of magnitude difficult.

5.7 Windowing: Realtime Visuals

Part of the magic of the graphics card is its ability to perform specific kinds of computation very quickly. It would be very beneficial to have some sort of way to display those visuals without having to serialize them to disk (e.g., a display function or a window of some sort which can be backed by a write-only image).

5.8 Error Noises

The compiler should make a snobby "Ouhh Hooo!" noise in french when the user puts in ill-formed code.

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Tutorial

2.1 Invoking the Compiler

Obtaining the LéPiX compiler – lepixc – in order to build LéPiX programs with it is simple, once the dependencies are set up. It requires a working OCaml compiler of version >= 4.0.3. Building requires ocamlbuild, ocamlfind, menhir and OPAM for an easy time, but if you are brave and willing to figure out the nightmare it takes to get these dependencies working on Windows than it can work on Windows machines as well.

From the root of the repository, run make –C source to create the compiler, or cd into the *source* directory and invoke make from there.

When the compiler is made, it will be within the ./source folder of the repository. An invocation without any arguments or filenames will explain to the user how to use it, like so:

```
1 source/lepixc
з Help:
       lepix [options] filename [filenames...]
            filename | filenames can have one option -i or --input
       options:
            -h
                 --help
                                   print the help message
                 --preprocess
                                   Preprocess and display source
            -i
                 --input
                                   Take input from standard in (
      default: stdin)
                --output <value> Set the output file (default:
10
```

```
stdout)
                  --tokens
                                     Print the stream of tokens
11
                                     Print the parsed Program
12
             -a
                  --ast
                                     Print the Semantic Program
                  --semantic
             -1
                  --llvm
                                     Print the generated LLVM code
14
                  --compile
                                     Compile the desired input and
             -c
      output the final LLVM
                  --verbose
                                     Be as explicit as possible with
             -\mathbf{v}
16
      all steps
```

history/lepixc.txt

Users can stack one-word options together using syntax like source/lepixc -ls, which will pretty-print the Semantic Analysis tree and also output the LLVM IR code.

2.2 Writing some Code

Function definitions are fairly C-Like, with the exception that all typeannotations appear on the right-hand-side. Users can define variables and functions by using the var and fun keywords, respectively.

To receive access to the standard library, put import lib in the program as well. It will end up looking somewhat like this:

```
import lib

fun main () : int {
    lib.print_n("hello world");
    return 0;
}
```

One function must always be present in your code, and that is the main function. It must return an int. If the user does not return an integer value from main, then a return 0 will be automatically done for you. Also of importance is that functions can have their return type figured out from their return statements. For example, you can remove the : int above and the code will still compile and run:

```
import lib

fun main () {
    lib.print_n("hello world");
    return 0;
```

6 }

This goes for more than just the main function, but any function you define!

Language Reference Manual

On the next page is the language reference manual for the original LéPiX language. The final implementation did not meet all of the requirements due to having to work alone for the last two weeks and insufficient teammate contribution during the project, but thankfully we still capture the majority of the language's constructs in every thing but the final code generation and semantic analysis stages, leaving room to improve the implementation in the future.

LéPiX Language Specification

Ceci n'est pas un Photoshop

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Part I Introduction

Tutorial

1.1 Hello, World!

This is an example of a Hello World program in LéPiX. It creates an array from an initializer, and then proceeds to save it to the directory of the running program under the name "hello.bmp":

Listing 1.1: hello world

1.2 Variables and Declarations

1.2.1 Variables

Variables are made with the var declaration. You can declare and assign variables by giving them a name and then referencing that name in other places.

```
1 fun main () : int {
2     var a : int = 24 * 2 + 1;
3    // a == 49
4    var b : int = a % 8;
5    // b == 1
6    var c : int[[5, 2]] = [
7    0, 2, 4, 6, 8, 10;
8    1, 3, 5, 7, 9, 11;
9    ];
10    var value : int = a + b + c[0, 4];
11    // value == 58
12    return value;
13 }
```

Listing 1.2: variable declaration and manipulation

1.2.2 Mutability

Variables can also be declared immutable or unchanging by declaring them with let. That is, let is the same as a var const, and var is the same as let mutable.

```
1 fun main () : int {
2     let a : float = 31.5;
3     var const b : float = 0.5;
4     var c : int = 0;
5     c = lib.trunc(a + b);
6     // compiler error: 'var const' variable is immutable
7     b = 2.5;
8     // compiler error: 'let' variable is immutable
9     a = 1.1;
10     return c;
11 }
```

Listing 1.3: mutability

1.3 Control Flow

Control flow is important for programs to exhibit more complex behaviors. LéPiX has for and while constructs for looping, as well as if, else if, else statements. They can be used as in the following sample:

```
1 fun main () : int {
2     for (var x : int = 0 to 10) {
3         var x : int = lib.random_int(0, 40);
4         if (x < 20) {
5             lib.print("It's less than 20!");
6         }
7         else {
8             lib.print("It's equal to or greater than 20");
9         }
10     }
11 }</pre>
```

"intro/tutorial/code/flow.hak"

1.4 Functions

1.4.1 Defining and Declaring Functions

Functions can be called with a simple syntax. The goal is to make it easy to pass arguments and specify types on those arguments, as well as the return type. All functions are defined by starting with the fun keyword, followed by an identifier including the name, before an optional list of parameters.

Listing 1.4: functions

1.4.2 Parameters and Arguments

All arguments given to a function for a function call are passed by value, unless the reference symbol & is written just before the argument, as shown in the below example. This allows a person to manipulate a value that was passed in directly, rather than receiving a copy of it the argument.

```
1 fun fibonacci_to (n : int, &storage : int[]) : int {
        int index = 0;
        var result : int = 0;
3
        var n_2 : int = 0;
4
        var n_1 : int = 1;
5
        while (n > 0) {
             result = n_1 + n_2;
             storage[index] = result;
             n_2 = n_1;
             n\_1 \, = \, result \, ;
10
             --n;
11
             ++index;
13
        return result;
14
15 }
16
17 fun main () : int {
        var storage : int [3] = [];
18
        var x : int = fibonacci_to(3, storage);
19
        return x;
20
21 }
```

Listing 1.5: arguments

Part II Reference Manual

Chapter 2

Expressions, Operations and Types

2.1 Variable Names and Identifiers

2.1.1 Identifiers

- 1. All names for all identifiers in a LéPiX program must be composed of a single start alpha codepoint followed by either zero or more of a digit or an alpha codepoint. Any identifier that does not follow this scheme and does not form a valid keyword, literal or definition is considered ill-formed.
- 2. All identifiers that containing two underscores __ in any part of the name are reserved for usage by the compiler implementation details and may not be used by programs. If an identifier has two underscores the program is considered ill-formed.
- 3. All identifiers prefixed by 'lib.' (i.e., belong in the lib namespace) are reserved by the standard to the standard library and nothing may be defined in that namespace by the program, aside from implementations of the standard library.

2.2 Literals

2.2.1 Kinds of Literals

There are many kinds of literals. They are:

literal:

boolean-literal integer-literal floating-literal string-literal

2.2.2 Boolean Literals

1. A boolean literal are the keywords true or false.

2.2.3 Integer Literals

- 1. An integer literal is a valid sequence of digits with some optional alpha characters that change the interpretation of the supplied literal.
- 2. A decimal integer literal uses digits '0' through '9' to define a base-10 number.
- 3. A hexidecimal integer literal uses digits '0' through '9', 'A' through 'F' (case insensitive) to define a base-16 number. It must be prefixed by 0x or 0x.
- 4. An octal integer literal uses digits '0' and '7' to define a base-8 number. It must be prefixed by 0c or 0C.
- 5. A binary integer literal uses digits 0 and 1 to define a base-2 number. It must be prefixed by 0b (case sensitive).
- 6. An n-digit integer literal uses the characters below to define a base-n number. It must be prefixed by 0n or 0N. It must be suffixed by #n, where n is the desired base. The character set defined for these bases

goes up to 63 characters, giving a maximum arbitrary base of 63. The characters which are:

$$0 - 9$$
, $A - Z$, $a - z$,

- 7. Arbitrary bases for n-digit must be base-10 numbers.
- 8. Groups of digits may be separated by a ' and do not change the integer literal at all.

2.2.4 Floating Literals

- 1. A floating literal has two primary forms, utilizing digits as defined in 2.2.3.
- 2. The first form must have a dot '.' preceded by an integer literal and/or suffixed by an integer literal. It must have one or the other, and may not omit both the prefixing or suffixing integer literal.
- 3. The second form follows 2, but includes the exponent symbol e and another integer literal describing that exponent. Both the exponent and integer literal must be present in this form, but if the exponent is included then the dot is not necessary and may be prefixed with only an integer literal or just an integer literal and a dot.

2.2.5 String Literals

- 1. A string literal is started with a single ''' or double '"' quotation mark and does not end until the next matching single ''' or double '"' quotation mark character, with respect to what the string was started with. This includes any and all spacing characters, including newline characters.
- 2. Newline characters in a multi-line string will be included in the string as an ASCII Line Feed \n character.
- 3. A string literal must remove the leading space on each line that are equivalent to all other lines in the text, and any empty leading space at the start of the string.

4. A string literal may retain the any leading space and common indentation by prefixing the opening single or double quotation mark with an 'R'.

2.3 Variable Declarations

2.3.1 let and var declarations

variable-initialization:

```
let | var ( mutable | const )optional < identifier> : <type>;
```

- 1. A variable can be declared using the let and var keywords, an identifier as defined in 2.1.1 and optionally followed by a colon ':' and type name. This is called a variable declaration.
- 2. A variable declared with let is determined to be immutable. Immutable variables cannot have their values re-assigned after declaration and initialization.
- 3. A variable declared with var is immutable. Mutable variables can have their values re-assigned after declaration and initialization.
- 4. let mutable is equivalent to var const.
- 5. It is valid to initialize or assign to a mutable variable from an immutable variable.
- 6. A declaration can appear at any scope in the program.

2.4 Initialization

2.4.1 Variable Initialization

variable-declaration:

```
 \begin{array}{lll} let & | \ var \ ( \ mutable \ | \ const \ ) \\ optional & < identifier > : < type > = ( \ expression \\ ); \\ \end{array} )
```

- 1. Initialization is the assignment of an expression on the right side to a variable declaration.
- 2. If the expression cannot directly initialize or be coerced to initialize the type on the left, then the program is ill-formed.

2.4.2 Assignment

```
assignment\hbox{-} expression:
```

```
expression = expression
```

2.5 Access

2.5.1 Member Access

member-access-expression:

```
( expression ) . < identifier>
```

- 1. Member access is performed with the dot '.' operator.
- 2. If the expression does not evaluate to a type that can be accessed with the dot operator, the program is ill-formed.
- 3. If the identifier is not available per lookup rules in 2.5.2 on the evaluated type, the program is ill-formed.

2.5.2 Member Lookup

1. When a member is accessed through the dot operator as in 2.5.1, a name must be found that matches the supplied identifier. If there is none,

2.6 Parenthesis

1. Parentheses define expression groupings and supersede precedence rules in 2.1.

2.7 Arithmetic Expressions

2.7.1 Binary Arithmetic Operations

```
addition-expression:
expression + expression
subtraction-expression:
expression - expression
division-expression:
expression / expression
multiplication-expression:
expression * expression
modulus-expression:
expression % expression
```

- 1. Symbolic expression to perform the commonly understood mathematical operations on two operands.
- 2. All operations are left-associative.

2.7.2 Unary Arithmetic Operations

```
unary-minus-expression:
-expression
```

- 1. Unary minus is typically interpreted as negation of the single operand.
- 2. All operations are left-associative.

2.8 Incremental Expressions

2.8.1 Incremental operations

```
post-increment-expression:
    ( expression )++
pre-increment-expression:
    ++( expression )

post-decrement-expression:
    ( expression )--
pre-decrement-expression:
    --( expression )

1. Symbolic expression that should semantically evaluate to ( expression )
    = ( expression ) + 1.
2. ( expression ) is only evaluated once.
```

2.9 Logical Expressions

2.9.1 Binary Compound Boolean Operators

and-expression:

```
expression and expression
expression && expression
or-expression:
expression or expression
expression || expression
```

- 1. Symbolic expressions to check for logical conjunction and disjunction.
- 2. For the and—expression, short-circuiting logic is applied if the expression on the left evaluates to false. The right hand expression will not be evaluated.
- 3. For the or—expression, short-circuiting logic is applied if the expression on the left evaluates to true. The right hand expression will not be evaluated.
- 4. All operations are left associative.

2.9.2 Binary Relational Operators

```
equal-to-expression:
    expression == expression
not-equal-to-expression:
    expression != expression
less-than-expression:
    expression < expression
greater-than-expression:
    expression > expression
less-than-equal-to-expression:
    expression <= expression</pre>
```

```
expression >= expression
```

- 1. Symbolic expression to perform relational operations meant to do comparisons.
- 2. All operations are left-associative.

2.9.3 Unary Logical Operators

```
inversion-expression:
    !expression
complement-expression:
    ~expression
```

1. Symbolic expression to perform unary logic operations, such as logical complement and logical inversion.

2.10 Bitwise Operations

2.10.1 Binary Boolean Operators

```
\begin{tabular}{ll} bitwise-and-expression: \\ expression & expression \\ bitwise-or-expression: \\ expression & | expression \\ bitwise-xor-expression: \\ expression & | expression \\ \end{tabular}
```

- 1. Symbolic expressions to perform logical / bitwise and, or, and exclusiveor operations.
- 2. All operations are left associative.

2.11 Operator and Expression Precedence

Precedence is defined as follows:

2.12 Expression and Operand Conversions

2.12.1 Boolean Conversions

- 1. Expressions that are expected to evaluate to booleans for the purposes of Flow Control as defined in 5.3 and for common relational and logical operations as in 2.9 will have their rules checked against the following:
 - (a) If the evaluated value is already a boolean, use the value directly.
 - (b) If the evaluated value is of an integral type, then any such type which compares equivalent to the integral literal 0 will be false: otherwise, it is true.
 - (c) If the evaluated value is of a floating point type, then any such type which compares equivalent to the floating point literal 0.0 will be false: otherwise, it is true.
 - (d) Otherwise, if there is no defined conversion, then the program is ill-formed.

2.12.2 Mathematical Conversions

int to float	float variable has the same value as integer.		
float to int	integer has largest integral value less than the float.		
bool to int	integer has value 1 if true, otherwise it will be 0.		
int to bool	bool is true if int is not equal to 0 and false otherwise.		

- 1. Implicit type conversions are carried out only for compatible types. The implicit casting occurs during assignment or when a value is passed as a function argument.
- 2. The four types of conversions that are supported are summarized in the table below.

Precedence	Operator	Variants	Associativity
1	++ () []	Postfix	Left to Right
2	++ + - ! ~	Prefix, Unary Operations	Right to Left
3	* / %		
4	+		
5	<<>>>	Binary Operations	Left to Right
6	< <= >> >=		
7	==!=		
8	&		
9	^		
10			
11			
12	= += -= *= /= %= <<= >>= &= ^= =	Assignments	Right-To-Left

3. If there is no conversion operator defined for those two types exactly, then the program is ill-formed.

Chapter 3

Functions

3.1 Functions and Function Declarations

Functions are independent code that perform a particular task and can be reused across programs. They can appear in any order and in one or many source files, but cannot be split among source files.

Function declarations tell the compiler how a function should be called, while function definitions define what the function does.

3.1.1 Function Definitions

```
fun <identifier > ([<parameter_declarations >]) : <
    return_type > {
        <function_body >
            [return <expression >;]
}
```

- 1. All function definitions in LéPiX are of the above form where they begin with the keyword fun, followed by the identifier, a list of optional parameter declarations enclosed in parentheses, optionally the return type, and the function body with an optional return statement.
- 2. return types can be variable types or void.

3. Functions that return void can either omit the return statement or leave it in or return the value unit:

```
fun zero ( &arr:int[] ) : void {
    for (var i : int = 0 to arr.length) {
        arr[i] = 0;
    }
}

fun zero ( &arr:int[] ) : void {
    for (var i : int = 0 to arr.length) {
        arr[i] = 0;
    }
    return;
}
```

4. Functions that return any other variable type must include a return statement and the expression in the return statement must evaluate to the same type as the return type or be convertible to the return type:

- 5. In the function add, arg1 and arg2 are passed by value. In the function zero, arr is passed by reference.
- 6. Function input parameters can be passed by value, for all variable types, or by reference, only for arrays and array derived variable types. See 3.1.3 for more about passing by value and reference.

3.1.2 Function Declarations

1. All function declarations in LéPiX are of the form

```
fun <identifier > ([<parameter_declarations >]) : <
    return_type >;
```

2. The function declaration for the add function from 3.1.1 would be

```
fun add ( arg1:float , arg2:float ) : float;
```

- 3. Function declarations are identical to function definitions except for the absence or presence of the code body.
- 4. Function declarations are optional, but useful to include when functions are used across multiple translation units to ensure that functions are called appropriately.

3.1.3 Function Scope and Parameters

- 1. Variables are declared as usual within the body of a function. The variables declared within the body of a function exist only in the scope of the function and are discarded when they go out of scope.
- 2. External variables are passed into functions as parameters. All variable types except arrays and array derived variable types are passed by value. Arrays and array derived variable types can be passed by both value and reference.
- 3. Passing value copies the object, meaning changes are made to the copy within the function and not the original. Passing by reference gives a pointer to the original object to the function, meaning changes are to the original within the function.
- 4. To pass by value to a function, use the variable name: add (x, y);
- 5. To pass by reference to a function, use the symbol & and the variable name, as in zero (&arr);.

Chapter 4

Data Types

4.1 Data Types

The types of the language are divided into two categories: primitive types and data types derived from those primitive types. The primitive types are the boolean type, the integral type int, and the floating-point type float. The derived types are struct, Array, and image and pixel, which are both special instances of arrays.

4.1.1 Primitive Data Types

1. int

By default, the int data type is a 32-bit signed two's complement integer, which has a minimum value of -2^{31} and a maximum value of 2^{32} .

2. float

The float data type is a single precision 32-bit IEEE 754 floating point.

3. boolean

The boolean data type has possible values true and false.

4.1.2 Derived Data Types

Besides the primitive data types, the derived types include arrays, structs, images, and pixels.

1. array

An array is a container object that holds a fixed number of values of a single type. Multi-dimensional arrays are also supported. They need to have arrays of the same length at each level.

2. pixel

A pixel data type is a wrapper for an array that will contain the representation for each pixel of an image. It will contain the rgb values, each as a separate int, and the gray value of a pixel.

3. image

The image data type is just an alias for a 2-dimensional array. The 2-d array will define the size of an image and contains a pixel as each of its data elements.

4. struct

A structure is a collection of one or more variables, possibly of different types, grouped together under a single name for convenient handling. Structures help to organize data because they permit a group of related variables to be treated as a unit instead of as separate entities.

Chapter 5

Program Structure and Control Flow

5.1 Statements

1. Any expression followed by a semicolon becomes a statement. For example, the expressions x = 2, lib.save(...), return x become statements:

```
x = 2;
lib.save(...);
return x;
```

2. The semicolon is used in this way as a statement terminator.

5.2 Blocks and Scope

Braces { and } are used to group statements in to blocks. Braces that surround the contents of a function are an example of grouping statements like this. Statements in the body of a for, while, if or switch statement are also surrounded in braces, and therefore also contained in a block. Variables declared within a block exist only in that block. A semicolon is not required after the right brace.

5.2.1 Blocks

1. At any point in a program, braces can be used to create a block. For example,

```
var x: int = 2;
var y: int = 4;
var result: int;
{
    var z: int = 6
    result = x + y + z;
}
```

- 2. In this trivial example, the statements on lines 5 and 6 live within their own block.
- 3. Blocks do have access to named definitions of their surrounding scope, however variables define within a block exist only within that block.

5.2.2 Scope

- 1. Scopes are defined as the collection of identifiers and available within the current lexicographic block¹.
- 2. Every program is implicitly surrounded by braces, which define the **global block**.

5.2.3 Variable Scope

- 1. Variables are in scope only within their own block².
- 2. In the example in Section 5.2.1, z is declared within the braces.
- 3. Variables declared within blocks last only within lifetime of that block.
- 4. If we attempted to access z outside of this block, this would cause an error to occur.
- 5. If a variable with a particular identifier has been declared and the identifier is re-used within a nested block.

¹This is usually between two curly braces {}

²E.g., between the brackets {}

- 6. The original definition of the identifier is **shadowed** and the new one is used until the end of the block.
- 7. Variables are constructed, that is, stored in memory when they are first encountered in their scope, and destructed at the scope's end in the reverse order they were encountered in.

5.2.4 Function Scope

- 1. Function definitions define a new block, which each have their own scope.
- 2. Function definitions have access to any variables within their surrounding scope, however anything defined in the function definition's block is not accessible in the surrounding blocks.
- 3. Variables defined in a parameter list belong to the definition-scope of the function.

5.2.5 Control Flow Scope

- 1. Control flow also introduces a new block with its own scope.
- 2. Variables initialized in any control flow statement, that is within the parenthesis before the block, belong to the control flow block and are not accessible in the surrounding block.
- 3. In the statement for $(\text{var } x = 0 \text{ to } 5) \{ ... \}$, x only exists within that for loop and destructed after the loop ends.

5.3 Namespaces

- 1. Namespaces are essentially blocks that allow identifiers to be prefix with an arbitrary nesting of names. They are declared with the namespace keyword, followed by several identifiers delimited by a dot '.' symbol.
- 2. Accessing variables and functions inside of a namespace must have the name of the namespace prefixed before the name of the desired identifier.

- 3. The namespaces lib and compiler is reserved for use by standard library implementations and the compiler.
- 4. Namespaces are the only bracket-delimited lexical scope that do not dictate the lifetime of the variables associated with them. These variables are part of the **global scope**.

5.4 if

```
if (expression; expression; ...)
    statements
else
    alternative - statements
```

- 1. if statements are used to make decisions in control flow.
- 2. Variations on this syntax are permitted, e.g. The else block of the if statement is optional.
- 3. If the expression is evaluated and returns true, then the first portion of the if statement is executed. Otherwise, if there is an else the portion after it is executed, and if there is none then the function continues at the next statement.
- 4. Parenthesis are optional after the if block if there is a single statement. If there are multiple statements, parenthesis are needed.
- 5. Variables can be initialized inside the expression portion of the if statement as long as the final expression in a semi-colon delimited list evaluates to a Boolean value.

```
if (var x = 20; var y = 50; x < y) { statements }
```

- 6. The scope for variables x and y is within that particular if statement.
- 7. If statements can also be nested so that multiple conditions can be tested:

```
\begin{array}{ll} \text{if } (x < 0) \\ y = -1 \\ \text{else if } (x > 0) \\ y = 1 \\ \text{else} \\ y = 0 \end{array}
```

5.5 switch

```
switch (variable) {
    case (constant expression):
        statements
    end;
    default: statements
}
```

- 1. switch statements can be used as an alternative to a nested if statements.
- 2. The variable is compared against the constant expression for each case, and if it is equal to this expression then the statements in that case are executed.
- 3. If the variable does not match any of the cases then the default case is executed.
- 4. The statements in each case must be followed by an end statement.

As with if statements, if a variable is declared within the switch like switch (var x =other_variable; x) { ... }, the scope for variable x is within that particular switch block.

5.6 while

```
while (expression; expression; ...; condition) {
    statements
}
```

- 1. while loops are used to repeat a block of code until some condition is met.
- 2. Every time a loop condition evaluates to true, the while loop's block and statements are executed.
- 3. When the condition evaluates to false, the while loop's execution is stopped.
- 4. Expressions before condition are evaluated only once. For example: while (var x = 20; x < 30) {...} is a valid the while loop, and only the final x < 30 is evaluated on each loop execution.
- 5. Loops are dangerous because they can potentially run forever. Make sure your conditions are done properly, or use Flow Control keywords and primitives discussed in 5.8:

```
while (var x = 1; x <= 10) {
    arr[x] = 1;
    x = x + 1;
}</pre>
```

5.7 for

```
for (variable = lower_bound to upper_bound by size) {
    statements
}
```

- 1. For loops are another way to repeat a group of statements multiple times.
- 2. The by keyword and argument size are optional and used to specify how much the variable should change by each iteration of the loop: for $(x = 1 \text{ to } 10 \text{ by } 2) \{ \dots \}$ will increment x by two each iteration rather than the default value of 1.
- 3. Variables can be declared in the loop declaration, as in for (var x = 1 to 10) $\{ \dots \}$.

4. For loops can also be used to decrement by swapping the positions of the lower_bound and upper_bound arguments, and using a negative value for the size (if using the by keyword) The while loop in Section 5.6 could be expressed as a for loop as follows.

```
for (var x = 1 to 10) {
    arr[x] = 1;
```

5. C-style for loops are also supported:

```
for (var x = 1; x \le 10; x++) { arr [x] = 1;
```

5.8 break and continue

Break and continue statements are used to exit a loop immediately, before the specified condition has been reached.

5.8.1 break

1. Break statements exit the block of a loop immediately.

```
while (...) {
    statements_above
    break;
    statements_below
}
```

2. In the example above, statements_above would be executed only once. The statements below would never be executed.³

5.8.2 break N

1. Break statements can be used to exit nested loops by jumping out of multiple scopes by adding an integral constant after the break keyword.

 $^{^3}$ Break statements are usually included inside of an if statement within the loop to immediately exit on a particular condition.

2. The example below will allow the user to break out of both for loops with only one break statement.⁴

```
for (...) {
    for (...) {
        statements_above
        if (condition)
            break 2;
        statements_below
    }
}
```

5.8.3 continue

1. Continue statements jump to the end of the loop body and begin the next iteration.

```
for (...) {
    statements_above
    if (expressions...) {
        continue;
    }
    statements_below
}
```

- 2. When a continue statement is executed, statements below the continue keyword are not executed, and the loop post-action and condition are immediately re-evaluated.
- 3. In the example above, statements_above would always be executed. The statements_below would be executed on iterations where the if condition was false, since when the if condition were true execution would jump back to the for loop's top.

 $^{^4}$ This could be considered a structured version of goto for loops and should be used with the programmer's utmost discretion.

Chapter 6

Parallel Execution

Since a large number of elementary operations in the realm of image processing are embarrassingly parallel matrix operations, the LéPiX language supports a simple parallelization scheme.

6.1 Parallel Execution Model

- 1. Parallel Execution is when two computations defined by the language are run at the exact same time by the abstract virtual machine, capable of accessing the same memory space.
- 2. The primary parallel primitive is a parallel-marked block.
- 3. Use of parallel primitives does not guarantee parallel execution: computation specified to run in parallel may run sequentially.¹

6.2 Syntax

1. The syntax for parallel code is code simply marked with the keyword parallel.

¹This could be due to hardware limitations, operating system limitations, and other factors of the machine.

- 2. In the situation where there are variables that must be shared by all the threads, a comma separated list of variable identifiers can be specified in parentheses using the keyword shared as in parallel { <block> }.
- 3. In the case of nested for loops, only the outermost loop carrying the parallel keyword is parallel.

6.3 Threads

- 1. The code inside of a parallel block can be dispatched to multiple executing threads.
- 2. Each thread that is spawned in this way will have its own scope, which is created when the thread is spawned and destroyed when the thread is killed.
- 3. Each thread has its own copy of each variable that is declares within the scope of the loop statements.
- 4. All variables are shared by default, except the ones declared in the parallel scope.

Part III Grammar Specification

Chapter 7

Grammar

7.1 Lexical Definitions and Conventions

A program consists of one or more translation units, which are translated in two phases, namely the preprocessing step and the lexing step. The preprocessing step entails carrying out directives which begin with # in a C-like style. The lexing step reduces the program to a sequence of tokens.

7.1.1 Tokens

- 1. Tokens belong to _categories. These are whitespace, keywords, operators, integer literals, floating point literals, string literals, identifiers, and brackets.
- 2. Whitespace tokens are used to separate other tokens and are ignored in any case where they do not occur between other non-whitespace tokens.

7.1.2 Comments

- 1. Comments come in two flavors: single-line and multi-line.
- 2. Single line comment begin with // and continue until the next newline character is found. Multi-line comments begin with /* and end with

- */. They are nested.
- 3. Comments are treated as whitespace tokens, but for various purposes may still appear between other whitespace tokens in a program's token stream.

7.1.3 Identifiers

1. Identifiers are composed of letters, numbers and the underscore character (_) but must begin with a letter. Identifiers beginning with underscores and numbers will be reserved for use within the implementation of the language.

7.1.4 Keywords

1. A set of identifiers has been reserved for use as keywords and cannot be used in other cases. The list of keywords is in the table below.

int	float	void	bool
unit	char	$\operatorname{codepoint}$	string
vector	matrix	vector	vec
var	let	if	else
for	while	by	to
return	true	false	mutable
const	fun	struct	maybe
protected	public	private	shared
as	of	parallel	atomic

7.1.5 Literals

1. Literals are of three types: integer literals, floating literals, and string literals, as detailed in 2.2.1. All of them use the following definitions for their digits:

```
\langle decimal-digit \rangle ::=  one of 0 1 2 3 4 5 6 7 8 9
```

```
\langle hexidecimal-digit \rangle ::=  one of
           0\ 1\ 2\ 3\ 4\ 5\ 6\ 7\ 8\ 9
           ABCDEF
           abcdef
    \langle binary-digit \rangle ::=  one of
           0.1
    \langle octal\text{-}digit \rangle ::= \text{ one of }
           0\ 1\ 2\ 3\ 4\ 5\ 6\ 7
    \langle n\text{-}digit \rangle ::= \text{ one of }
           0\; 1\; 2\; 3\; 4\; 5\; 6\; 7\; 8\; 9
           ABCDEFGHIJKLMNOPQRSTUVWXYZ
           abcdefghijklmnopqrstuvwxyz
    \langle decimal-digit-sequence \rangle ::= \langle \rangle
           \langle decimal-digit \rangle \langle decimal-digit-sequence \rangle
    \langle binary-digit-sequence \rangle ::= \langle \rangle
           \langle binary-digit \rangle \langle binary-digit-sequence \rangle
    \langle octal\text{-}digit\text{-}sequence \rangle ::= \langle \rangle
           \langle octal-digit \rangle \langle octal-digit-sequence \rangle
    \langle hexidecimal-digit-sequence \rangle ::= \langle \rangle
           \langle hexidecimal-digit \rangle \langle hexidecimal-digit-sequence \rangle
    \langle n\text{-}digit\text{-}sequence \rangle ::= \langle \rangle
           \langle n\text{-}digit \rangle \langle n\text{-}digit\text{-}sequence \rangle
2. Integer literals consist of sequences of digits are always interpreted as
    decimal numbers. They can be represented by the following lexical
    compositions:
    \langle integer-literal \rangle ::= \langle decimal-literal \rangle
           \langle binary-literal \rangle
            \langle octal\text{-}literal \rangle
            \langle hexidecimal-literal \rangle
           \langle n\text{-}digit\text{-}literal \rangle
    \langle decimal-literal \rangle ::= \langle decimal-digit-sequence \rangle
```

```
 \langle binary\text{-}literal \rangle ::= 0 b \ \langle binary\text{-}digit\text{-}sequence} \rangle   | 0B \ \langle binary\text{-}digit\text{-}sequence} \rangle   \langle octal\text{-}literal \rangle ::= 0 c \ \langle octal\text{-}digit\text{-}sequence} \rangle   | 0C \ \langle octal\text{-}digit\text{-}sequence} \rangle   \langle hexidecimal\text{-}literal \rangle ::= 0 x \ \langle hexidecimal\text{-}digit\text{-}sequence} \rangle   | 0X \ \langle hexidecimal\text{-}digit\text{-}sequence} \rangle   \langle n\text{-}digit\text{-}literal} \rangle ::= 0 n \ \langle n\text{-}digit\text{-}sequence} \rangle   | 0N \ \langle n\text{-}digit\text{-}sequence} \rangle
```

3. Floating point literals can be specified using digits and a decimal points or in scientific notation. The following regular expression represents the set of acceptable floating-point constants.

```
\langle e\text{-}part \rangle ::= e \langle + |-\text{``}-\rangle \langle integral\text{-}literal \rangle
\langle floating\text{-}literal \rangle ::= \langle integral\text{-}literal \rangle_{opt} \cdot \langle integral\text{-}literal \rangle_{opt} \langle e\text{-}part \rangle_{opt}
| \langle integral\text{-}literal \rangle_{opt} \langle e\text{-}part \rangle
```

4. String literals are sections of quote-delimited items. They are defined as follows:

```
\langle single\text{-}quote \rangle ::= '
\langle double\text{-}quote \rangle ::= ''
\langle raw\text{-}specifier \rangle ::= R_{opt}
\langle character \rangle ::= \langle escape\text{-}character \rangle \langle source\text{-}character \rangle
\langle character\text{-}sequence \rangle ::= \langle \rangle
| \langle character \rangle \langle character\text{-}sequence \rangle
\langle string\text{-}literal \rangle ::= \langle raw\text{-}specifier \rangle \langle double\text{-}quote \rangle \langle character\text{-}sequence \rangle
| \langle raw\text{-}specifier \rangle \langle single\text{-}quote \rangle \langle character\text{-}sequence \rangle \langle single\text{-}quote \rangle
```

7.2 Expressions

The following sections formalize the types of expressions that can be used in a LéPiX program and also specify completely, the precedence of operators and left or right associativity.

7.2.1 Primary Expression

```
\langle primary\_expression \rangle ::= \langle identifier \rangle
| \langle integer\_constant \rangle
| \langle float\_constant \rangle
| (expression)
```

1. A primary expression are composed of either a constant, an identifier, or an expression in enclosing parentheses.

7.2.2 Postfix Expressions

```
 \langle postfix\_expression \rangle ::= \langle primary\_expression \rangle \\ | \langle postfix\_expression \rangle \text{ (argument\_list )} \\ | \langle postfix\_expression \rangle \text{ [expression ]} \\ | \langle postfix\_expression \rangle \text{ . identifier } \langle argument\_list \rangle ::= \langle \rangle \\ | \langle argument\_list \rangle \text{ , } \langle postfix\_expression \rangle
```

1. A postfix expression consist of primary expression followed by postfix operators. The operators in postfix expressions are left-associative.

Indexing

- 1. Array indexing consists of a postfix expression, followed by an expression enclosed in square brackets. The expression in the brackets must evaluate to an integer which will represent the index to be accessed.
- 2. The value returned by indexing is the value in the array at the specified index.

Function Calls

- 1. A function call is a postfix expression (representing the name of a defined function) followed by a (possibly empty) list of arguments enclosed in parentheses.
- 2. The argument list is represented as a comma separated list of postfix expressions.

Structure access

1. The name of a structure followed by a dot and an identifier name is a postfix expression. The expression's value is the named member's of the structure that is being accessed.

7.2.3 Unary Expression

```
 \begin{array}{l} \langle unary\_operator \rangle ::= \tilde{\phantom{a}} \\ | \  \, | \  \, | \\ | \  \, - \\ | \  \, * \\ \\ \langle unary\_expression \rangle ::= \langle unary\_operator \rangle \ \langle postfix\_expression \rangle \\ \end{array}
```

- 1. A unary expression consists of ¡postfix_expression; preceded by a unary operator (- , , !, * ,&).
- 2. Unary expressions are left-associative.
- 3. The unary operation is carried out after the postfix expression has been evaluated.

The function of each unary operator has been summarized in the table below:

-	Unary minus	
~	Bitwise negation operator	
^	Logical negation operator	
*	Indirection operator	

7.2.4 Casting

The LéPiX language supports the casting of an integer to a floating point value and vice versa. It also supports casting of an integer value to a boolean value and vice versa. Integer to float casting creates a floating point constant with the same value as the integer. Casting a floating point value to an integer rounds down to the nearest integral value. Casting a boolean value

to an integer gives 1 if the value is true and 0 if it is false. Casting an integer to a boolean value yields false if the value is 0 and true otherwise.

```
\langle cast\_expression \rangle ::= \langle unary\_expression \rangle
| \langle unary\_expression \rangle \text{ as } \langle type\_name \rangle
```

7.2.5 Multiplicative Expressions

The multiplication (*), division (/) and modulo (%) operators are left associative.

```
\langle multiplicative\_expression \rangle ::= \langle cast\_expression \rangle
| \langle multiplicative\_expression \rangle * \langle cast\_expression \rangle
| \langle multiplicative\_expression \rangle / \langle cast\_expression \rangle
| \langle multiplicative\_expression \rangle % \langle cast\_expression \rangle
```

7.2.6 Additive Expressions

The addition (+) and subtraction (-) operators are left associative.

```
\langle additive\_expression \rangle ::= \langle multiplicative\_expression \rangle

| \langle additive\_expression \rangle + \langle cast\_expression \rangle

| \langle additive\_expression \rangle - \langle cast\_expression \rangle
```

7.2.7 Relational Expressions

The relational operators less than (i), greater than (i), less than or equal to (i=) and greater than or equal to (i=) are left associative.

7.2.8 Equality Expression

```
\langle equality\_expression \rangle ::= \langle relational\_expression \rangle
| \langle equality\_expression \rangle != \langle relational\_expression \rangle
| \langle equality\_expression \rangle == \langle relational\_expression \rangle
```

7.2.9 Logical AND Expression

```
\langle logical\_and\_expression \rangle ::= \langle equality\_expression \rangle
| \langle logical\_and\_expression \rangle \&\& \langle equality\_expression \rangle
```

1. The logical and operator (&&) is left associative and returns true if both its operands are not equal to false.

7.2.10 Logical OR Expression

```
\langle logical\_or\_expression \rangle ::= \langle logical\_and\_expression \rangle 
| \langle logical\_or\_expression \rangle || \langle logical\_and\_expression \rangle
```

1. The logical OR operator (——) is left associative and returns true if either of its operands are not equal to false.

7.2.11 Assignment Expressions

```
\langle assignment\_expression \rangle ::= \langle logical\_or\_expression \rangle \\ | \langle unary\_expression \rangle = \langle assignment\_expression \rangle
```

1. The assignment operator (=) is left associative.

7.2.12 Assignment Lists

```
\langle assignment\_list \rangle ::= \langle assignment\_expression \rangle
| \langle assignment\ list \rangle ::= \langle assignment\ list \rangle, \langle assignment\ expression \rangle
```

1. Assignment lists consist of multiple assignment statements separated by commas.

7.2.13 Declarations

- 1. Declarations of a variable specify a type for each identifier and a value to be assigned to the identifier.
- 2. Declarations do not always allocate memory to be associated with the identifier.

7.2.14 Function Declaration

```
\langle function\_declaration \rangle ::= \text{ fun } \langle identifier \rangle \ ( \langle params\_list \rangle \ ) : \langle type\_name \rangle
\langle params\_list \rangle ::= \langle \ \rangle
| \ \langle identifier \rangle : \langle type\_name \rangle
| \ \langle params\_list \rangle \ , \langle identifier \rangle : \langle type\_name \rangle
```

- 1. Function declarations consist of the keyword fun followed by an identifier and a list of parameters enclosed in parentheses.
- 2. The list of arguments is followed by a colon and a type name which represents the return type for the function.
- 3. The arguments list is specified as a comma-separated list of identifier, type pairs.

7.3 Statements

Statements are executed sequentially in all cases except when explicit constructs for parallelization are used. Statements do not return values.

```
 \langle statement \rangle ::= \langle expression\_statement \rangle \\ | \langle branch\_statement \rangle \\ | \langle compound\_statement \rangle \\ | \langle iteration\_statement \rangle \\ | \langle return\_statement \rangle
```

7.3.1 Expression Statements

```
\langle expression\_statement \rangle ::= \langle \rangle 
| \langle expression \rangle ;
```

- 1. Expression statements are either empty or consist of an expression.
- 2. These effects of one statement are always completed before the next is executed.
- 3. This guarantee is not valid in cases where explicit parallelization is used.
- 4. Empty expression statements are used for loops and if statements where not action is to be taken.

7.3.2 Statement Block

```
 \langle block \rangle ::= \{ \langle compound\_statement \rangle \} 
 | \{ \langle block \rangle \langle compound\_statement \rangle \} 
 \langle compound\_statement \rangle ::= \langle declaration \rangle 
 | \langle statement \rangle 
 | \langle compound\_statement \rangle ; \langle declaration \rangle 
 | \langle compound\_statement \rangle ; \langle statement \rangle
```

- 1. A statement block is a collection of statements declarations and statements.
- 2. If the declarations redefine any variables that were already defined outside the block, the new definition of the variable is considered for the execution of the statements in the block.
- 3. Outside the block, the old definition of the variable is restored.

Branch Statements

```
\langle branch\_statement \rangle ::= if (\langle expression \rangle) \langle statement \rangle fi
| if (\langle expression \rangle) \langle statement \rangle else \langle statement \rangle fi
```

1. Branch statement are used to select one of several statement blocks based on the value of an expression.

7.3.3 Loop Statements

```
\langle loop\_statement \rangle ::= \text{ while } (\langle expression \rangle) \langle statement \rangle \\ | \text{ for } (\langle identifier \rangle = \langle expression \rangle \text{ to } \langle expression \rangle) \langle statement \rangle \\ | \text{ for } (\langle assignment\_expression \rangle = \langle expression \rangle \text{ to } \langle expression \rangle \text{ by } \langle expression \rangle \\ | \text{ for } (\langle expression \rangle; \langle expression \rangle; \text{ expression}; \text{ }) \langle statement \rangle \\
```

1. Loop statements specify the constructs used for iteration and repetition.

7.3.4 Jump Statements

```
\langle jump\_statement \rangle ::= break \langle integer-literal \rangle_{opt} | continue
```

1. Jump statements are used to break out of a loop or to skip the current iteration of a loop.

7.3.5 Return Statements

```
\langle return\_statement \rangle ::= return
\mid return \langle expression \rangle
```

1. Return statements are used to denote the end of function logic and the also to specify the value to be returned by a call to the function in question.

7.4 Function Definitions

```
\langle function \ definition \rangle ::= \langle function \ declaration \rangle \langle block \rangle
```

1. Function definitions consist of a function declaration followed by a statement block.

7.5 Preprocessor

```
⟨preprocessor_directive⟩ ::= #define ⟨identifier⟩ ⟨expression⟩
| #ifdef ⟨identifier⟩
| #ifndef ⟨identifier⟩
| #endif
| #import ⟨identifier⟩
| #import "⟨file_name⟩"
| #import string ⟨file_name⟩
```

- 1. Before the source for a LePix program is compiled, the program is consumed by a preprocessor, which expands macro definitions and links libraries and other user-defines to the current file, as specified by appropriate preprocessor directives.
- 2. define macros create an alias for a value or expression.
- 3. ifdef and ifndef macros are used to check if a particular alias has already been assigned. Import directives are used to link files/libraries with the current program.

7.6 Grammar Listing

```
\langle primary \ expression \rangle ::= \langle identifier \rangle
       \langle integer\text{-}constant \rangle
       \langle float\text{-}constant \rangle
       (expression)
\langle postfix \ expression \rangle ::= \langle primary \ expression \rangle
       \(\langle postfix_expresion \rangle \) (argument_list )
       \langle postfix \ expression \rangle \ [expression]
       \langle postfix\_expression \rangle. identifier
\langle argument \ list \rangle ::= \langle \rangle
     \langle argument\_list \rangle, \langle postfix\_expression \rangle
\langle unary \ operator \rangle := \tilde{\ }
\langle unary \ expression \rangle ::= \langle unary \ operator \rangle \langle postfix \ expression \rangle
\langle cast\_expression \rangle ::= \langle unary\_expression \rangle
     \langle unary\ expression \rangle as \langle type\ name \rangle
\langle multiplicative \ expression \rangle ::= \langle cast \ expression \rangle
       \langle multiplicative\_expression \rangle * \langle cast\_expression \rangle
       \langle multiplicative\_expression \rangle / \langle cast\_expression \rangle
       \langle multiplicative\_expression \rangle \% \langle cast\_expression \rangle
```

```
\langle additive \ expression \rangle ::= \langle multiplicative \ expression \rangle
       \langle additive\_expression \rangle + \langle cast\_expression \rangle
       \langle additive\_expression \rangle - \langle cast\_expression \rangle
\langle relational\_expression \rangle ::= \langle additive\_expression \rangle
       \langle relational \ expression \rangle < \langle additive \ expression \rangle
       \langle relational \ expression \rangle \le \langle additive \ expression \rangle
       \langle relational \ expression \rangle > \langle additive \ expression \rangle
       \langle relational \ expression \rangle >= \langle additive \ expression \rangle
\langle equality | expression \rangle ::= \langle relational | expression \rangle
       \langle equality\_expression \rangle != \langle relational\_expression \rangle
       \langle equality | expression \rangle == \langle relational | expression \rangle
\langle logical \ and \ expression \rangle ::= \langle equality \ expression \rangle
       \langle logical \ and \ expression \rangle \&\& \langle equality \ expression \rangle
\langle logical \ or \ expression \rangle ::= \langle logical \ and \ expression \rangle
      \langle logical \ or \ expression \rangle \mid \langle logical \ and \ expression \rangle
\langle assignment \ expression \rangle ::= \langle logical \ or \ expression \rangle
     \langle unary \ expression \rangle = \langle assignment \ expression \rangle
\langle assignment \ list \rangle ::= \langle assignment \ expression \rangle
       \langle assignment \ list \rangle ::= \langle assignment \ list \rangle, \langle assignment \ expression \rangle
\langle declaration \rangle ::= \operatorname{var} \langle storage\_class \rangle \langle identifier \rangle : \langle type\_name \rangle = \langle postfix\_expression \rangle
\langle type \ name \rangle ::= bool
      int
       float
       \langle type\_name \rangle \langle array \rangle
\langle array \rangle ::= [\langle int\_list \rangle]
  | [\langle array \rangle]
\langle int\_list \rangle ::= \langle integer \rangle
  |\langle int\_list \rangle, \langle integer \rangle
\langle function\_declaration \rangle ::= \text{fun } \langle identifier \rangle \ (\langle params\_list \rangle) : \langle type\_name \rangle
```

```
\langle params\_list \rangle ::= \langle \rangle
        \langle identifier \rangle : \langle type\_name \rangle
        \langle params\_list \rangle, \langle identifier \rangle: \langle type\_name \rangle
\langle statement \rangle ::= \langle expression \ statement \rangle
        \langle branch \ statement \rangle
        \langle compound statement \rangle
        \langle iteration \ statement \rangle
        \langle return \ statement \rangle
\langle expression \ statement \rangle ::= \langle \rangle
   |\langle expression \rangle ;
\langle block \rangle ::= \{ \langle compound\_statement \rangle \}
      \{ \langle block \rangle \langle compound \ statement \rangle \}
\langle compound \ statement \rangle ::= \langle declaration \rangle
        \langle statement \rangle
        \langle compound \ statement \rangle \ ; \ \langle declaration \rangle
        \langle compound \ statement \rangle \ ; \langle statement \rangle
\langle parallel \ block \rangle ::= parallel (\langle parallel \ control \ variables \rangle) \langle block \rangle
\langle branch\_statement \rangle ::= if (\langle expression \rangle) \langle statement \rangle fi
      if (\langle expression \rangle) \langle statement \rangle else \langle statement \rangle fi
\langle branch \ statement \rangle ::= if (\langle expression \rangle) \langle statement \rangle fi
       if \langle expression \rangle \langle statement \rangle else \langle statement \rangle fi
\langle loop \ statement \rangle ::= while (\langle expression \rangle) \langle statement \rangle
       for \langle identifier \rangle = \langle expression \rangle to \langle expression \rangle \rangle \langle statement \rangle
       for \langle assignment\_expression \rangle = \langle expression \rangle to \langle expression \rangle by \langle expression \rangle
        ) \langle statement \rangle
       for \langle expression \rangle; \langle expression \rangle; expression; \rangle \langle statement \rangle
\langle identifier | list \rangle ::= \langle identifier \rangle
       \langle identifier\_list \rangle, \langle identifier \rangle
\langle jump \; statement \rangle ::= break \langle integer-literal \rangle_{opt}
      continue
\langle return\_statement \rangle ::= return
     return \langle expression \rangle
```

Chapter 4

Plan

Our plan was developed slowly and mostly solidified around the making of our Language Reference Manual and a bit afterwards. We met once a week, sometimes a second time if our Advisor had the time for it, and occasionally held extra meetings to help get things done.

4.1 Process

Most of our planning was done in-person via weekly meetings. We also used Github Issues to track things and also bikeshed some of our progress and implementation. We closed issues as they passed and had issues tied to Milestones in the project:

4.2 Timeline

Our timeline was given by the milestones we had for the project. We opened them early, meaning that one of the milestones (GPU Codegen) was scrapped when our team decided that we would not pursue such an avenue.

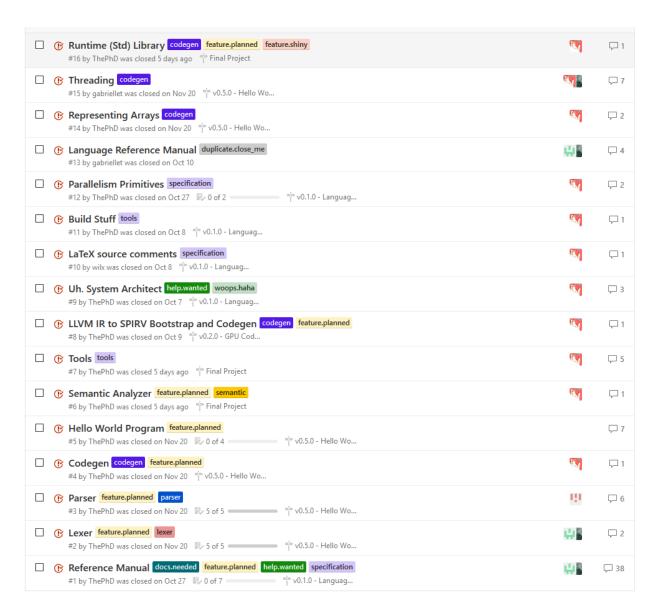


Figure 4.1: Closed issues throughout the project https://github.com/ThePhD/lepix/issues?q=is%3Aissue+is%3Aclosed.

4.3 Tools

Everyone was free to develop in whatever IDE or editor they wished, just so long as they could invoke the makefile. As I was originally the System Architect, I put together a list of all the tools someone would need to invoke the build process in Figure 4.3. The command line dependencies here helped me figure out what was needed when we started to do testing 6.

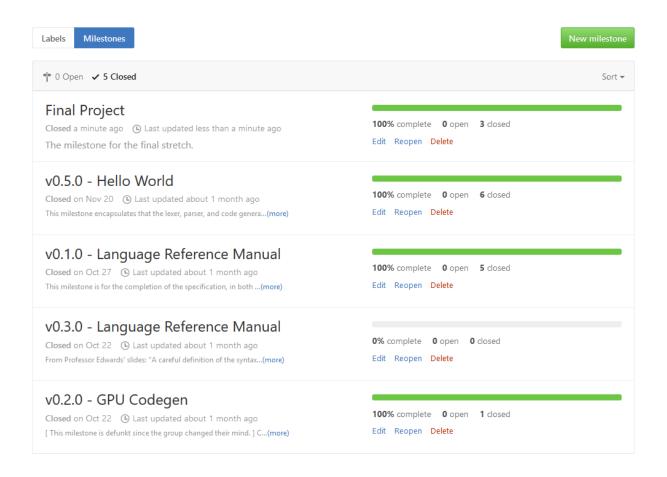


Figure 4.2: Milestones for the project https://github.com/ThePhD/lepix/milestones?state=closed.

4.4 Project Log

Asides from issues being closed and comments being made, the best project log that shows how I did is the git commit log for all the branches, included below. It was generated from git using the command git —no-pager log —graph —abbrev-commit —decorate —date=relative —all.

```
1 * commit 5d983f7 (HEAD -> master, origin/master, origin/HEAD)
2 | Author: ThePhD <phdofthehouse@gmail.com>
3 | Date: 27 minutes ago
4 |
5 | Final clean implementation of bottom-type type derivation for function returns, good literals, and overloading
6 |
7 * commit e323adc
8 | Author: ThePhD <phdofthehouse@gmail.com>
9 | Date: 22 hours ago
```

```
10
        overloading tests among other things
11
12
13 * commit 714d07f
    Author: ThePhD <phdofthehouse@gmail.com>
14
    Date:
            3 days ago
15
16
        last pdfs and reports and readme update
17
18
    commit aa558db
19 *
    Author: ThePhD <phdofthehouse@gmail.com>
20
            3 days ago
21
22
        remove temporaries and debug print statements
23
24
    commit a7e8fa2 (origin/feature/semantic, feature/semantic)
    Author: ThePhD <phdofthehouse@gmail.com>
    Date:
            3 days ago
28
        Buh.
29
30
* commit 816bad8
    Author: The PhD < phd of the house @gmail.com>
            4 days ago
34
        All I have left ....
35
36
37 * commit ecbdf74
  | Author: ThePhD <phdofthehouse@gmail.com>
            4 days ago
    Date:
40
        We're so close to the end. Don't give up. Trust in
41
      yourself, and fight for what was right...
42
* commit fb3b94d
    Author: ThePhD <phdofthehouse@gmail.com>
44
            8 days ago
    Date:
46
        It finally builds... the basics of the Semantic AST,
47
      finally more or less in place...!
48
49 * commit c00e2ab
  | \ Author: \ The PhD < phd of the house @gmail.com >
    Date:
             2 weeks ago
52
        Beef up the semantic AST and fully complete the pretty
53
      printer for it.
54
55 * commit 817f6b8
```

```
Author: ThePhD <phdofthehouse@gmail.com>
    Date:
             3 weeks ago
57
58
         Ensure lowercase acceptance as well.
59
60
61 * commit 79 fc 476
    Author: ThePhD <phdofthehouse@gmail.com>
62
    Date:
             3 weeks ago
63
64
         Fixing up the parser for integer literals and other more
65
      useful things.
66
67 * commit 2bfb992
    Author: ThePhD <phdofthehouse@gmail.com>
    Date:
             3 weeks ago
70
71
         scanner and parser are up to snuff
72
73 * commit 77847d7
    Author: ThePhD <phdofthehouse@gmail.com>
    Date:
             3 weeks ago
75
76
         handle extra cases in the parser for increment, decrement,
77
       and assignment-ops
78
79 * commit 92 ee 4 d 6
    Author: ThePhD <phdofthehouse@gmail.com>
             3 weeks ago
81
82
         Beat up the parser lots.
83
84
       commit 4eb1aeb
85 *
      Merge: 5efaebe 989e1d2
86 \
      Author: ThePhD <phdofthehouse@gmail.com>
      Date:
               3 weeks ago
89
           Merge branch 'master' into feature/semantic
90
91
         commit 989e1d2
92
         Merge: 9ccd7e0 bebe5ad
93
         Author: Jackie Lin < jackielin13@gmail.com>
         Date:
                 3 weeks ago
95
96
             Merge branch 'master' of https://github.com/ThePhD/
      lepix
98
  commit 9ccd7e0
         Author: Jackie Lin < jackielin 13@gmail.com>
101 | Date:
                 3 weeks ago
```

```
102
             tests
103
         commit 5efaebe
         Author: ThePhD <phdofthehouse@gmail.com>
         Date:
                 3 weeks ago
108
             [ci skip] heavily modify the parser to handle type
109
       qualifications, improve the AST, and begin to consider
       scoping rules
110
         commit 6c4629a
111 *
112
         Author: ThePhD <phdofthehouse@gmail.com>
         Date:
                 3 weeks ago
113 //
114
             [ci skip] commit so I can jump back to helping on
115
       master
116
       commit bebe5ad (origin/testing/travis, testing/travis)
       Author: ThePhD <phdofthehouse@gmail.com>
       Date:
               3 weeks ago
119
120
           REALLY fuck you, python3
121
122
       commit d5cb6ef
123
       Author: ThePhD <phdofthehouse@gmail.com>
124
               3 weeks ago
       Date:
125
126
127
           Fuck you too, python3
       commit fc38fe8
       Author: ThePhD <phdofthehouse@gmail.com>
130
       Date:
               3 weeks ago
131
132
           specifically invoke python3 because environments are
133
       stupid?
134
       commit a10dfaa
135
136
       Author: ThePhD <phdofthehouse@gmail.com>
       Date:
               3 weeks ago
137
138
           51 builds later, it should work...
139
       commit 38f438b
       Author: ThePhD <phdofthehouse@gmail.com>
       Date:
               3 weeks ago
143
144
           ensure only python 3 is available on the system
145
146
```

```
commit 7eb9368
       Author: ThePhD <phdofthehouse@gmail.com>
                3 weeks ago
            Only need one of either ---rm or -d
151
152
       commit d2f844f
153 *
       Author: \ The PhD < phd of the house @gmail.com >
154
       Date:
                3 weeks ago
156
            Proper travis ci with safety net for made directories
157
       for tests
158
       commit aebffd0
159
       Author: \ The PhD < phd of the house @gmail.com >
160
                3 weeks ago
       Date:
            make sure the shell is configured with eval opam...
164
       commit cc6dca8
165
       Author: ThePhD <phdofthehouse@gmail.com>
166
                3 weeks ago
       Date:
167
168
            This is getting a tad tiresome, but it's my fault for
       not having a good handle on travis-ci
170
       commit cbbbe02
171
       Author: ThePhD <phdofthehouse@gmail.com>
172
                3 weeks ago
       Date:
            Proper escaping?
175
176
       commit df0a2f1
       Author: ThePhD <phdofthehouse@gmail.com>
178
                3 weeks ago
       Date:
179
180
            escaped operators
181
182
       commit 48 edeca
183
       Author: ThePhD <phdofthehouse@gmail.com>
184
       Date:
                3 weeks ago
185
186
            Party with the cd comms
       commit 14788e0
189
       Author: ThePhD <phdofthehouse@gmail.com>
190
       Date:
                3 weeks ago
191
192
            "Docker never dies!" (Sleep infinity)
193
```

```
194
       commit\ 35\,af11a
       Author: ThePhD <phdofthehouse@gmail.com>
                3 weeks ago
       Date:
            Attemping without heredoc and just docker exec...
199
200
       commit d9d03ee
201
       Author: ThePhD <phdofthehouse@gmail.com>
202
       Date:
                3 weeks ago
203
204
           No -c on bash when using heredoc
205
206
       commit 2f6003c
207
       Author: \ The PhD < phd of the house @gmail.com >
208
       Date:
                3 weeks ago
            Attempting to make docker behave better?
212
       commit 9d0cf03
       Author: \ The PhD < phd of the house @gmail.com >
       Date:
                3 weeks ago
216
            explicit printing
218
       commit c6e8299
219
       Author: ThePhD <phdofthehouse@gmail.com>
220
                3 weeks ago
221
            Fix the tests because I'm bad at writing python code,
       weee
224
       commit f27c410
       Author: The PhD < phd of the house @gmail.com>
226
       Date:
                3 weeks ago
227
228
            blot out the Semantic Stuff until its time
229
230
231
         commit 71c87ad
         Merge: b032732 b55ae1c
232
         Author: The PhD < phd of the house @gmail.com>
         Date:
                  3 weeks ago
234
              Merge branch 'feature/semantic' into testing/travis
237
              # Conflicts:
              #
                   .travis.yml
239
240
         commit b55ae1c
241
```

```
Author: ThePhD <phdofthehouse@gmail.com>
         Date:
                  3 weeks ago
243
244
              default llvm fails because Opam is a heaping pile of
       shit
246
         commit 7a13678
247
         Author: The PhD < phd of the house @gmail.com>
         Date:
                  3 weeks ago
249
250
              Attempt to properly propogate bash errors and fix
251
       travis files
252
         commit 2ae921c
253
         Author: The PhD < phd of the house @gmail.com>
254
         Date:
                  3 weeks ago
              update test harness
258
         commit 81fb441
259
         Author: ThePhD <phdofthehouse@gmail.com>
260
                  3 weeks ago
         Date:
261
262
              semantic analyizer start
263
264
         commit a9011e3 (origin/feature/preprocessor, feature/
265
       preprocessor)
         Author: The PhD < phd of the house @gmail.com>
266
         Date:
                  3 weeks ago
              fix gitignore
269
270
         commit 11b86f8
271
         Author: ThePhD <phdofthehouse@gmail.com>
272
         Date:
                  3 weeks ago
273
274
              Full preprocessor implementation
275
276
         commit 63a06f6
277
          Author: ThePhD <phdofthehouse@gmail.com>
278
          Date:
                  3 weeks ago
280
              update gitignore
281
         commit b032732
         Author: ThePhD <phdofthehouse@gmail.com>
                  4 weeks ago
         Date:
286
              Go fuck yourself, travis, and your rules against tabs
287
```

```
288
          commit\ 75\,b8679
          Author: ThePhD <phdofthehouse@gmail.com>
          Date:
                  4 weeks ago
              Let's try this again...
293
294
          commit\ 1\,e827\,e2
295
          Author: \ The PhD < phd of the house@gmail.com >
296
          Date:
                  4 weeks ago
297
298
              messing around with docker
299
300
       commit 01adc38
301
       Author: \ The PhD < phd of the house @gmail.com >
302
                4 weeks ago
       Date:
            One more missing package from depext
305
306
       commit 8fa39e9
307
       Author: ThePhD <phdofthehouse@gmail.com>
308
                4 weeks ago
       Date:
309
310
311
            package names were wrong
312
       commit 337c74b
313
       Author: ThePhD <phdofthehouse@gmail.com>
314
       Date:
                4 weeks ago
315
316
            source script to propogate errors better
            autoshit and mcrap tools need to be there
318
319
       commit 871197c
320 *
       Author: The PhD < phd of the house @gmail.com>
321
       Date:
                4 weeks ago
322
323
            travis gooo
324
325
326
       commit 868dc82
       Author: ThePhD <phdofthehouse@gmail.com>
327
       Date:
                4 weeks ago
328
329
            assume yes for ALL cases...
330
331
       commit 81d341a
332
       Author: ThePhD <phdofthehouse@gmail.com>
333
       Date:
                4 weeks ago
334
335
            say yes, all the time
336
```

```
337
       commit b4dba58
        Author: ThePhD <phdofthehouse@gmail.com>
       Date:
                4 weeks ago
341
            update ignore
342
343
       commit 5176344
344
        Author: ThePhD <phdofthehouse@gmail.com>
345
       Date:
                 4 weeks ago
346
            update properly
348
349
       commit e9fb2af
350
       Author: \ The PhD < phd of the house @gmail.com >
351
       Date:
                4 weeks ago
            Keeep trying with docker ....
354
355
       commit\ 7\,b\,4\,1\,eef
356
       Author: \ The PhD < phd of the house @gmail.com >
357
                4 weeks ago
       Date:
358
359
            Poking at things 'till it works...
361
       commit f4c7c3b
362
        Author: ThePhD <phdofthehouse@gmail.com>
363
                4 weeks ago
364
365
            Properly bind the mount volume with the -v command, then
        swap into it
367
       commit d6b090f
368 *
        Author: ThePhD <phdofthehouse@gmail.com>
369
       Date:
                4 weeks ago
370
371
            poke at env to understand what's going on
372
373
374
       commit 667d0d1
        Author: ThePhD <phdofthehouse@gmail.com>
375
       Date:
                 4 weeks ago
376
377
            travis_run file and friends
       commit 2e8aab2
380
       Author: ThePhD <phdofthehouse@gmail.com>
381
       Date:
                4 weeks ago
382
383
            try it from a file now...
384
```

```
385
       commit\ 83\,afd3\,b
386 *
        Author: ThePhD <phdofthehouse@gmail.com>
        Date:
                4 weeks ago
389
            super duper docker
390
391
       commit 78f655d
392
        Author: ThePhD <phdofthehouse@gmail.com>
393
                4 weeks ago
        Date:
394
395
            Goddamn tabs
396
397
       commit cca9cd2
398
       Author: The PhD < phd of the house @gmail.com>
399
                4 weeks ago
       Date:
            update travis work
402
403
       commit \ c0e85a3
404
        Author: ThePhD <phdofthehouse@gmail.com>
405
                4 weeks ago
       Date:
406
407
            Use a different language to attempt to stay out of the
       python shell
409
       commit\ 9e00c5f
410 *
       Author: The PhD < phd of the house @gmail.com>
411
                4 weeks ago
       Date:
            update tests file and travis CI yaml file
414
415
       commit 8b71e6d
        Author: ThePhD <phdofthehouse@gmail.com>
417
                4 weeks ago
       Date:
418
419
            Add .travis file to start CI
420
421
422
       commit 0f25039 (origin/feature/codegen, feature/codegen)
        Author: ThePhD <phdofthehouse@gmail.com>
423
       Date:
                4 weeks ago
424
425
            Testing fixture
       commit e200757
       Author: ThePhD <phdofthehouse@gmail.com>
       Date:
                4 weeks ago
430
431
```

```
Example code for linking an external library. Various
       small changes to the driver of lepix and the codegen in
       preparation for the Semantic Analyzer and the AST.
           We still need something to preprocess source code ...
       another regular parser, perhaps?
434
       commit\ f6a208d
435 *
       Author: The PhD < phd of the house @gmail.com>
436
       Date:
                4 weeks ago
437
438
           Re-raise any bad errors we don't know how to catch
439
440
441
       commit 6b73260
       Author: ThePhD <phdofthehouse@gmail.com>
442
       Date:
                4 weeks ago
443
444
           CARAT DIAGNOSTICS YEEEAAH
       commit\ 07c1a08
       Author: ThePhD <phdofthehouse@gmail.com>
       Date:
               5 weeks ago
449
450
           better polyfill code
451
452
       commit cf4f1a5
453
       Author: ThePhD <phdofthehouse@gmail.com>
454
       Date:
               5 weeks ago
455
456
            Better options parser, again, mostly for the sake of
       writing clearer, better code
458
       commit \ c50176d
459 *
       Author: ThePhD <phdofthehouse@gmail.com>
460
       Date:
                5 weeks ago
461
462
           properly guard additions to sub
463
464
       commit 9c6eca8
465
       Author: ThePhD <phdofthehouse@gmail.com>
466
       Date:
                5 weeks ago
467
468
           More functional string_split
469
       commit 251c6ea
       Author: ThePhD <phdofthehouse@gmail.com>
       Date:
                5 weeks ago
473
474
           "cleaner" polyfill ...?
475
476
```

```
477 * | commit 058cd89
       Author: ThePhD <phdofthehouse@gmail.com>
               5 weeks ago
           goofing off with trying to write better functional code
481
482
     commit f5cc25f
483 *
     Author: The PhD < phd of the house @gmail.com>
484
             5 weeks ago
     Date:
485
486
         Full on driver and options implementation
487
         Polyfill layer to replace any missing batteries / core
488
         IO layer for opening and writing to a file
489
490
491 * commit e97a9c4
     Author: ThePhD <phdofthehouse@gmail.com>
     Date:
             5 weeks ago
494
         More comments, restructuring, and lexer-error handling.
495
496
     commit \ 5933\,da5
497
     Author: ThePhD <phdofthehouse@gmail.com>
498
             5 weeks ago
500
         We now have a driver that handles the code
501
         There is now an option to print out the token stream
502
         The lepix top level performs a basic amount of error
503
       handling now
         The parser and lexer now do a very thorough job of
504
       tracking line information; may want to propogate into the AST
       somehow
505
506 * commit b60947a
     Author: ThePhD <phdofthehouse@gmail.com>
     Date:
             5 weeks ago
508
509
         It might be beneficial to mess with how the lexing and
510
       parsing are run through, so we can generate the proper line
       numbers and token lists.
         We should also look into reading from and writing to files
511
       , even if we don't have the Batteries library and other bits
       set up for this.
         One day ...
512
513
514 * commit 5fb47e1
   Author: ThePhD <phdofthehouse@gmail.com>
             5 weeks ago
     Date:
516
517
```

```
counted arrays, namespace declarations and proper parallel
518
        binding declarations
519
     commit ce06c8f
     Author: ThePhD <phdofthehouse@gmail.com>
     Date:
              5 weeks ago
522
523
          It works.
524
          Now I'm going to redo the whole goddamn AST and Parser so
       we can really get going ...
526
   * commit 82b848f
527
     Author: ThePhD <phdofthehouse@gmail.com>
528
     Date:
              5 weeks ago
529
530
531
          Segmentation fault.
           S \ E \ G \ M \ E \ N \ T \ A \ T \ I \ O \ N \ F \ A \ U \ L \ T \ L \ A \ D \ I \ E \ S \ . 
533
        commit c35d50d
534
       Merge: 677d5bb c8e22b9
535
       Author: The PhD < phd of the house @gmail.com>
536
                5 weeks ago
       Date:
537
538
            Merge branch 'feature/codegen'
540
       commit c8e22b9
541
        Author: ThePhD <phdofthehouse@gmail.com>
542
                5 weeks ago
543
544
            skeleton of semantic analyzer
545
       commit 677d5bb
547 *
        Author: fennilin <jackielin13@gmail.com>
548 /
        Date:
                5 weeks ago
549
550
            Create 11-17-16
551
552
553 * commit f736ab5
     Author: ThePhD <phdofthehouse@gmail.com>
554
555
     Date:
              5 weeks ago
556
          properly append qualified id to list
557
558
        commit 748f5c0
       Merge: a347044 e120364
560
       Author: ThePhD <phdofthehouse@gmail.com>
561
       Date:
                5 weeks ago
562
563
            Merge branch 'feature/codegen'
564
```

```
565
           # Conflicts:
566
           # .gitignore
       commit e120364
569
       Author: ThePhD <phdofthehouse@gmail.com>
570
       Date:
               5 weeks ago
571
572
            Allow for interwoven function and data declarations
           Properly concatenate qualified IDs
574
           Start on code generation (nothing actually appears)
           Make sure top-level does not trigger semantic analyzer (
576
       its empty right now)
577
       commit 018952e
       Author: ThePhD <phdofthehouse@gmail.com>
       Date:
               5 weeks ago
581
           Array type now takes a number plus a type, rather than
582
       having a separate type for each one
           scratch source example that can be modified and
583
       committed to any current contention for a person working on
       the compiler
       commit e8b60ed
585
       Author: ThePhD <phdofthehouse@gmail.com>
586
       Date:
               5 weeks ago
587
588
589
           More ignore files and an empty main test.
       commit f5dc624
       Author: ThePhD <phdofthehouse@gmail.com>
592
                5 weeks ago
       Date:
593
594
           Remove built files (please don't commit these again...)
595
596
       commit 6f0db16
597
       Author: ThePhD <phdofthehouse@gmail.com>
598
               5 weeks ago
599
600
           Clean up these commits...
601
602
         commit 9b1384b
603
         Merge: ae4b9f6 907c6b0
         Author: ThePhD <phdofthehouse@gmail.com>
605
606
         Date:
                  5 weeks ago
607
             Merge remote-tracking branch 'origin/master' into
608
       feature/codegen
```

```
609
             # Conflicts:
             #
                   source/parser.mly
                   source/scanner.mll
613
         commit ae4b9f6
614
         Author: ThePhD <phdofthehouse@gmail.com>
615
         Date:
                  5 weeks ago
616
617
             Add qualified ID handling (we will improve it later to
618
        handle arbitrarily long strings)
619
620
         commit\ a347044
         Author: Akshaan Kakar <akshaan.crackers@gmail.com>
621
         Date:
                 5 weeks ago
622 | / |
623
              Edited .gitignore to omit built files in the source
624
       directory
625
       commit 907c6b0
626
       Author: Akshaan Kakar <akshaan.crackers@gmail.com>
627
                5 weeks ago
628
629
           Added atomic statement blocks
631
       commit 5956c1b
632
       Author: Akshaan Kakar <akshaan.crackers@gmail.com>
633
                5 weeks ago
634
635
           Simple parallel blocks (without atomic sections) and
       array literals now work in Parser+AST
637
638 * commit 7cc276a
     Author: Akshaan Kakar <akshaan.crackers@gmail.com>
639
     Date:
             5 weeks ago
640
641
         AST complete with mildly-pretty printing
642
643
644
     commit d1dcd73
     Author: Akshaan Kakar <akshaan.crackers@gmail.com>
645
     Date:
             5 weeks ago
646
647
         JK function declarations work too LOL
648
650 * commit f810dd3
   | Author: Akshaan Kakar <akshaan.crackers@gmail.com>
     Date:
             5 weeks ago
652
653
```

```
654
         Almost done with AST and pretty printing for all language
       constructs. Only function and variable decls to go
655
     commit 71 a6fcb
     Author: Akshaan Kakar <akshaan.crackers@gmail.com>
             5 weeks ago
658
659
         Added simple top level. Edited ast, parser and lexer but
660
      some errors remain. Pretty printing needs to be set up.
661
     commit d29fa01
662
     Author: Fatima <fatimakoli14@gmail.com>
663
             6 weeks ago
664
665
         Parallelblock and jump statements added
666
667
       commit 772 f3b9
       Merge: e039350 2a939a4
       Author: Fatima <fatimakoli14@gmail.com>
670
       Date:
                6 weeks ago
671
672
           Merge branch 'master' of https://github.com/ThePhD/lepix
673
674
       commit 2a939a4
       Author: Akshaan Kakar <akshaan.crackers@gmail.com>
676
                6 weeks ago
677
678
           Deleted intermediate files and yacc output
679
680
       commit 80e8614
       Author: Akshaan Kakar <akshaan.crackers@gmail.com>
                6 weeks ago
683
684
           Parser simplified. Multiple expression grammar rules
685
       collapsed into single rule
686
         commit 7fd0c99
687
         Merge: bb3f3bb 1dcf3f5
688
         Author: ThePhD <phdofthehouse@gmail.com>
689
         Date:
                  6 weeks ago
690
691
             Merge remote-tracking branch 'origin/master'
692
693
         commit bb3f3bb
         Author: ThePhD <phdofthehouse@gmail.com>
695
         Date:
                  6 weeks ago
696
697
             Makin' a bootstrapper....
698
699
```

```
commit\ ec7f5dc
         Author: ThePhD <phdofthehouse@gmail.com>
         Date:
                  6 weeks ago
             example LLVM code for the Linux architecture.
704
705
         commit e039350
706 *
         Author: Fatima <fatimakoli14@gmail.com>
707
         Date:
                  6 weeks ago
708
709
              Completed-needs to be tested
710
711
       commit 1dcf3f5
712 *
       Author: Fatima <fatimakoli14@gmail.com>
713
       Date:
                6 weeks ago
           Pretty printer started and array nodes added
         commit\ aad 0d 34
         Merge: f093c1a 4fe3fa3
       | Author: Fatima <fatimakoli14@gmail.com>
                  6 weeks ago
         Date:
721
             Merge branch 'master' of https://github.com/ThePhD/
       lepix
724
         commit 4fe3fa3
725
         Author: Akshaan Kakar <akshaan.crackers@gmail.com>
726
                  6 weeks ago
727
         Date:
             Added empty files for semantic checker and codegen
729
730
       commit 4660520
731
       Author: \ The PhD < phd of the house@gmail.com >
732
       Date:
                6 weeks ago
733
734
            floating point hello world and other examples as well
735
736
     * commit e87ca95
737
738
       Author: ThePhD <phdofthehouse@gmail.com>
       Date:
                6 weeks ago
739
740
           example code in C for many of the hello worlds and basic
741
        examples
742
    * commit 32bb043
       Author: ThePhD <phdofthehouse@gmail.com>
       Date:
                6 weeks ago
745
746
```

```
That's some thick LLVM IR...
747
748
    * commit 9b966c6
749
       Author: ThePhD <phdofthehouse@gmail.com>
                6 weeks ago
751
752
            Perfect parallel_2d example
753
754
       commit 6211aa9
755
       Author: \ The PhD < phd of the house @gmail.com >
       Date:
                6 weeks ago
757
758
            It works, uguu.
759
760
     * commit a5dfc40
761
       Author: \ The PhD < phd of the house@gmail.com >
                6 weeks ago
       Date:
764
            Additional hello world and the beginnings of a fleshed
765
       out parallel looping structure
766
       commit \ f093c1a
767
       Author: Fatima <fatimakoli14@gmail.com>
                7 weeks ago
       Date:
770
            edited ast
771
772
773 * commit ad02f05
     Author: Gabrielle A Taylor <gat2118@columbia.edu>
              7 weeks ago
776
          Simple C threading program, sums 2d array vertically
777
778
779 * commit cd6f557
     Author: Gabrielle A Taylor <gat2118@columbia.edu>
780
     Date:
              7 weeks ago
781
782
          Simple C threading program that sums 2d array
783
784
     commit f9e1a85
785 *
     Author: Gabrielle A Taylor <gat2118@columbia.edu>
786
     Date:
              7 weeks ago
787
788
          Simple C threading program
789
790
     commit 2e71b8e
     Author: Akshaan Kakar <akshaan.crackers@gmail.com>
     Date:
              7 weeks ago
793
794
```

```
Added top level file (lepix.ml) and deleted intermediate
795
       files from lex and yacc
796
       commit 6d409f9
797 *
       Merge: 542ee64 dbf4809
798
       Author: Akshaan Kakar <akshaan.crackers@gmail.com>
799
                7 weeks ago
800
801
           Merge \ branch \ 'master' \ of \ https://github.com/ThePhD/lepix
802
803
       commit dbf4809
804
       Author: fennilin <jackielin13@gmail.com>
805
                7 weeks ago
806
807
           Create 11-03-16
808
809
       commit \ 542\,ee64
       Author: Akshaan Kakar <akshaan.crackers@gmail.com>
                7 weeks ago
812
813
           Added missing tokens to parser
814
815
     commit b47e043
     Author: Akshaan Kakar <akshaan.crackers@gmail.com>
     Date:
              7 weeks ago
818
819
         All rules added to parser. No S/R or R/R conflicts. Need
820
       to defined entry point for compiler (i.e. 'main')
821
       commit 980\,e2a3
       Merge: 0ecdb5f f6283ac
       Author: Akshaan Kakar <akshaan.crackers@gmail.com>
824
                8 weeks ago
       Date:
825
826
           Merge branch 'master' of https://github.com/ThePhD/lepix
827
828
           Edited parser.mly
829
830
       commit f6283ac
831
       Author: ThePhD <phdofthehouse@gmail.com>
832
       Date:
                8 weeks ago
833
834
           preprocessor is still eluding me with a parse error.
       Need to get more info about this.
836
       commit 847c13c
       Author: ThePhD <phdofthehouse@gmail.com>
                8 weeks ago
       Date:
839
840
```

```
smaller array size, return value at end of main function
841
842
    * commit 4592035
843
       Author: ThePhD <phdofthehouse@gmail.com>
                8 weeks ago
846
            fix my dumb math
847
848
       commit\ 0\,ecd\,b\,5\,f
849
       Author: Akshaan Kakar <akshaan.crackers@gmail.com>
850
                8 weeks ago
851
852
            Implemented parser for all expression types. No shift
853
       reduce errors
854
     commit b38b568
     Author: ThePhD <phdofthehouse@gmail.com>
     Date:
              8 weeks ago
858
          ignore intermediate files
859
860
     commit \ 458\,c07\,e
861
     Author: ThePhD <phdofthehouse@gmail.com>
              8 weeks ago
864
          parallel example and preprocessor code
865
866
       commit a1129c4
867
       Merge: b46b06f 2b3d9d7
       Author: ThePhD <phdofthehouse@gmail.com>
       Date:
                8 weeks ago
871
            Merge branch 'master' into feature/preprocessor
872
873
           # Conflicts:
874
           # .gitignore
875
876
       commit 2b3d9d7
877
       Author: ThePhD <phdofthehouse@gmail.com>
878
       Date:
                8 weeks ago
879
880
            update toplevel display file
881
       commit\ aba7a94
       Author: ThePhD <phdofthehouse@gmail.com>
884
       Date:
                8 weeks ago
885
886
            updated specification source files
887
888
```

```
* commit 35483aa
       Author: ThePhD <phdofthehouse@gmail.com>
                8 weeks ago
           Toplevel PDFs we can link to.
893
894
       commit 42168dd
895
       Author: ThePhD <phdofthehouse@gmail.com>
896
                8 weeks ago
       Date:
897
898
            final specificaiton before submission
899
900
       commit e29b45e
901
       Author: ThePhD <phdofthehouse@gmail.com>
902
       Date:
                8 weeks ago
903
904
            update specification commit and ignore files
       commit 8b811a3
907
       Author: ThePhD <phdofthehouse@gmail.com>
908
       Date:
                8 weeks ago
909
910
           update specification
911
       commit cb1e9f1
       Author: Akshaan Kakar <akshaan.crackers@gmail.com>
914
       Date:
                9 weeks ago
915
916
           Added appropriate rules for integer and float literals
       in expr grammar in parser.mly
918
         commit f1c2181
         Merge: a53a16f 8bee3d7
         Author: Akshaan Kakar <akshaan.crackers@gmail.com>
         Date:
                  9 weeks ago
922
923
              Merge\ branch\ 'master'\ of\ https://github.com/ThePhD/
       lepix
925
926
              Merging
927
         commit a53a16f
         Author: Akshaan Kakar <akshaan.crackers@gmail.com>
         Date:
                  9 weeks ago
931
              Added type, loops, conditionals and array access
932
       grammars to parser
933
         commit b46b06f
934 *
```

```
Author: ThePhD <phdofthehouse@gmail.com>
         Date:
                  8 weeks ago
936
              update specification
939
         commit \ 5740\,bf0
940 *
         Author: ThePhD <phdofthehouse@gmail.com>
941
         Date:
                  9 weeks ago
942
943
              Not quite there yet. Need to ask about it.
944
         commit 8e666f1
946
         Author: ThePhD <phdofthehouse@gmail.com>
947
         Date:
                  9 weeks ago
948
949
              the skeleton of the preprocessor for all of this stuff
950
         commit 3cb13fa
         Author: ThePhD <phdofthehouse@gmail.com>
953
         Date:
                  9 weeks ago
954
955
              begin preparing the bootstrap.py
956
957
       commit 8bee3d7 (origin/specification, specification)
       Author: ThePhD <phdofthehouse@gmail.com>
959
                9 weeks ago
960
961
            specification updates
962
963
       commit \ c3d3fb9
       Author: ThePhD <phdofthehouse@gmail.com>
                9 weeks ago
966
967
            re-add specification to align git submodules without
968
       breaking anything
969
       commit 25a2e7d
970
       Author: ThePhD <phdofthehouse@gmail.com>
971
                9 weeks ago
972
973
            remove specification source since it was bugged
974
975
       commit 96\,a20\,b7
       Author: ThePhD <phdofthehouse@gmail.com>
                9 weeks ago
978
979
            update git modules
980
981
       commit daec995
982 *
```

```
Merge: 3d945dd d7811ca
983
        Author: Akshaan Kakar <akshaan.crackers@gmail.com>
                9 weeks ago
            Merge branch 'master' of https://github.com/ThePhD/lepix
987
988
        commit d7811ca
989
        Author: ThePhD <phdofthehouse@gmail.com>
990
        Date:
                 9 weeks ago
991
992
            remove old specification files
993
994
        commit 604101b
995
        Author: ThePhD <phdofthehouse@gmail.com>
996
        Date:
                9 weeks ago
997
998
            make correct overleaf bridge in right place
1000
          commit 5418373
1001
          Merge: d539d38 bcea4f1
1002
          Author: The PhD < phd of the house @gmail.com>
1003
                  9 weeks ago
          Date:
1004
1005
              Merge branch 'master' of github.com: ThePhD/lepix
1007
              Fix deletion of everything
1008
1009
          commit d539d38
1010
          Author: ThePhD <phdofthehouse@gmail.com>
          Date:
                  9 weeks ago
              make overleaf bridge
1014
1015
          commit 3d945dd
1016 *
          Author: Akshaan Kakar <akshaan.crackers@gmail.com>
                   9 weeks ago
1018
          Date:
1019
              Added rules for single line comments as well as for
1020
        nesting multi-line comments
        commit bcea4f1
1022 *
        Author: Akshaan Kakar <akshaan.crackers@gmail.com>
1023
        Date:
                9 weeks ago
            Added regex for floating pointer literals to scanner.mll
1026
1027
        commit 36e9ae5
1028
        Author: Akshaan Kakar <akshaan.crackers@gmail.com>
1029
        Date:
                9 weeks ago
1030
```

```
Added augmented version of the MicroC scanner
1032
1033
     commit f9d771a
      Author: Fatima <fatimakoli14@gmail.com>
              2 months ago
1036
          Basic tokens added to Parser
1038
     commit\ 8174\,ed7
      Author: ThePhD <phdofthehouse@gmail.com>
1041
      Date:
              3 months ago
1042
1043
          Remove SPIRV-LLVM setup.
1044
1045
     commit bed07b3
     Author: ThePhD <phdofthehouse@gmail.com>
     Date:
              3 months ago
1049
          As it stands... we will not be doing SPIRV stuff. Since
1050
       the focus will JUST be on multicore, which can be done fine
       on the CPU itself.
1051
     commit\ 805\,e0\,c5
      Author: ThePhD <phdofthehouse@gmail.com>
1053
              3 months ago
1054
          This commit allows for grammar basics.
1056
1057
          Need to figure out how to wrap threads in LLVM IR code.
1060 * commit de5101d
     Author: ThePhD <phdofthehouse@gmail.com>
1061
              3 months ago
      Date:
1062
1063
          SPIRV-LLVM node
1064
     commit 83d693e
1066 *
      Author: ThePhD <phdofthehouse@gmail.com>
1067
      Date:
              3 months ago
1068
1069
          Submodule LLVM <-> SPIRV
1070
1071
     commit 2a032f1
     Author: ThePhD <phdofthehouse@gmail.com>
1073
1074
     Date:
              3 months ago
1075
          remove old lepix file name
1076
1077
```

```
1078 * commit 31a94b4
      Author: ThePhD <phdofthehouse@gmail.com>
              3 months ago
          Spoopy language specification
1082
1083
1084 * commit 547718b
      Author: \ The PhD < phd of the house@gmail.com>
1085
      Date:
              3 months ago
1086
1087
          Skeleton files, to get ready to work.
1088
1089
1090 * commit 2497bbd
      Author: \ The PhD < phd of the house@gmail.com >
1091
      Date:
              3 months ago
1092
1093
1094 | Purged.
```

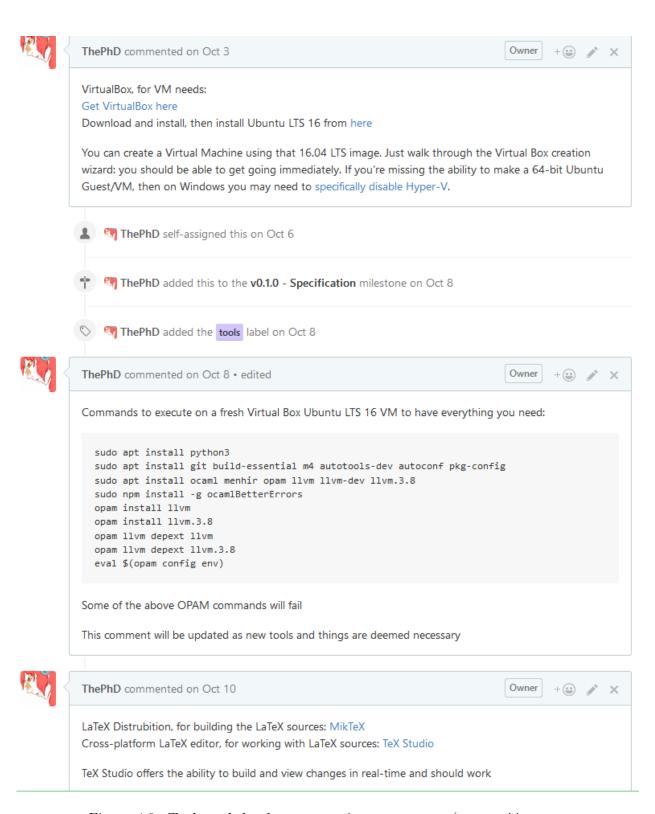


Figure 4.3: Tools and development environment setup (https://github.com/ThePhD/lepix/issues/7).

Design

5.1 Interface

The overall interface works by simply inferring more and more information from the previous step, in a manner like so:

```
Input (String) ⇒ Preprocessor [Separate Lexer, Parser] (String) ⇒ Lexer (Token Stream) ⇒ Parser (Abstract Syntax Tree) ⇒ Semantic Analyzer (Program Attributes, Semantic Syntax Tree) ⇒ Code Generation (LLVM IR Module)
```

Each step feeds a slimmed-down step to the next parser. The diagram for the workflow can be seen in Figure 5.1.

5.1.1 Top Level Work-flow

The way it works is simple at the highest level. Each stage produces one piece of work and hands it off to the next. To support error-reporting, a context argument is also provided to certain stages, geared to hold tracking information for that stage.

Each component flows from the next, with Preprocessing being an optional step that took in an input file and produced a source string. Because of the

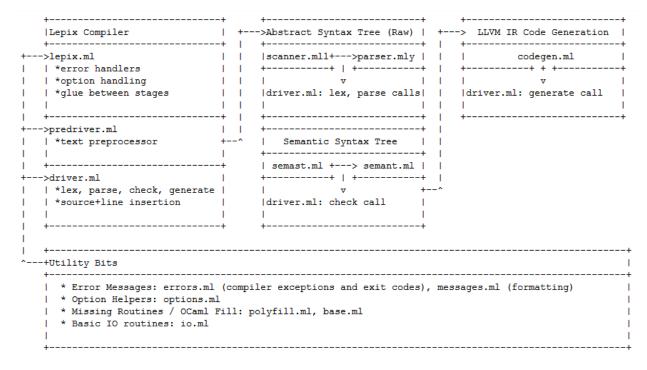


Figure 5.1: Compiler and implementation organization.

way I handled input to the lexer and parser, defining an input channel for either all the text or using an input stream such as stdin was simple.

5.1.2 Namespaces

Namespaces became trivial to implement. Since namespaces themselves are not allowed to be a type – just a meta-construct for organization – they needed no physical entry in the final Semantic Syntax Tree. All names of definitions were simply folded to be fully qualified. It would behoove me in further implementation to keep track of the original unqualified name, to ensure that I can easily access both versions of the information easily.

The only other thing that would have come in handy to implement is using namespace syntax. As of right now, accessing things in namespaces requires fully qualified names. This could get very annoying very quickly, and therefore some form of a using statement would be good.

5.1.3 Bottom-up Type Derivation

Bottom-up Type Derivation was, essentially, just analyzing the return values of a function. The implementation does not have checkers for expressions that initialize a variable, so the mechanism is not in place for automatic variable deduction, even though the Semantic Syntax Tree will properly tag a variable definition without a type specifier to be of automatically-deduced type.

Thankfully, return values for functions were done. Here, I simply re-ran the semantic analyzer after going through once and getting the return types of all functions and return expressions, and then ran it again to re-gather symbols into the global environment and program attributes. The first Semantic Syntax Tree was thrown out entirely in this process. It is a cheap programmatic way to get the derivations to resolve for all the top-level symbols gathered earlier. A much more effective and composed method would be nice for later, but I had to implement things as fast as possible.

5.1.4 Overloading

Overloading was conceptually difficult to figure out. Many languages have overloading, but a lot of the code examples I had seen were done at runtime. That is, argument arity and strict type matching were done, overloads ranked by that, and then any conversion paths were taken into consideration before determining if a single function call was the best function call. We implemented similar here, without taking into account conversions (because that would require ranking all overloads fully, which is a veritable nightmare to do properly and the source of hundreds – perhaps thousands – of bugs in C++, C#, and other language compilers).

It was important to note that overloading is not part of the original AST, just the Semantic AST as a type for Qualified IDs (only identifiers can be overloaded) and for codegen purposes become mangled names to prevent name lookup confusion (the source of a bug that plagued even past the original deadline and only fixed recently).

5.1.5 Error Handling

True error handling with notices and carat diagnostics were only implemented for the first 3 stages of the compiler: preprocessing, lexing and parsing. Every thing else only has basic exception handlers and no context object to propagate source information or provide carat diagnostics. Thankfully, the test programs were small enough that it was easy to know what was producing errors. The downside is that this means the compiler is not very friendly to users beyond the initial parsing stages, and errors can be even more cryptic than OCaml's.

My primary motivation for good error handling came from OCaml's lacking error messages. Dozens upon dozens of "syntax error" messages that did not even seem to point to the right line, where let statements would chain well with inner expressions and only error at the end of the program, even though the error that threw off the parser in the first place was much further up in the program. Using and definitions helped in that regard, but there was still a lot of lost implementation time.

Unfortunately, our error handling again does not do a good job for the semantic errors, which – once you get used to OCaml's error messages – are actually quite good. This would take a lot more time to do appropriately, so it is unfortunate that I did not get to do more of it. I really liked implementing carat diagnostics and good error messages with line and character information, and I think it helped me fix the parser and lexer much faster and iterate over it better.

5.2 Division of Labor

I wrote essentially the entire implementation, with little kept from older commits. At one point, Fatima Koly and Akshaan Kakar's for the parser and lexer remained.

Testing and Continuous Integration

6.1 Test Code

Some of the more interesting test cases include one to include a preprocessing directive (a temporary replacement for a decent module system), bottom-up type derivation for return values from functions, and a demonstration of overloading. The test cases are very involved and often nest elements to reveal bugs or other inconsistencies in the code generator (for example, properly implementing Llvm.build_load only in conditions where the type being asked for is a form of pointer). Most of these tests also had a failure case on the other side of it as well, especially in the case of overloading and bad literals. There are still bugs with expressions not quite being checked when assigned back to the original for proper convertibility, but I managed to cover a small but good area of code for working on this by myself.

```
1 import lib
3 #import "imported.lepix"
5 fun main () : int {
        var f : int = x.d();
        lib . print_n(f);
        return 0;
9 }
                        Listing 6.1: preprocess.lepix
1 fun two () {
        return 2;
2
3 }
4
5 fun main () {
        return two();
7 }
                           Listing 6.2: auto.lepix
1 import lib
3 namespace n.s {
        var global : int = 8;
7 }
9 namespace n {
        namespace s {
             var stuff : float = 3.5;
12
13 }
14
15 fun s () : int {
        return 2;
16
17 }
19 fun s (x : int) : int {
20
        return x + 2;
21 }
22
23 fun main () : int {
        var local : int = n.s.global;
        var svalue : int = s();
        lib . print_n(local);
27
        lib . print_n ( svalue ) ;
        lib.print_n(s(2));
28
        lib .print_n(n.s.global);
29
        lib . print_n(n.s.stuff);
```

```
31     return svalue;
32 }
```

Listing 6.3: overloads.lepix

6.2 Test Automation

6.2.1 Test Suite

Our test suite is a Python 3 Unit Test¹ suite, using the subprocess module to write code that called the lepix compiler, lepixc. In the case of return code 0 (success), it would then call LLVM's IR interpeter lii with the -c flag to run the program.

6.2.2 Online Automation

This was a decent bit of automation, but to further enhance our capability to know what was broken and what was fixed, I implemented Travis Continuous-Integration (travis-ci) ² support through a travis.yml file in the top level our repository. Travis-ci is free for any publicly available, open-source github repository (the code is MIT Licensed).

6.2.3 Online Automation Tools

Docker came in handy when travis-ci had not updated their own pool of images for a very long time. We configured travis-ci to run all our commands in a small docker container using the latest ubuntu, ensuring that we had the proper OPAM, OCaml, and other development tools we needed. This was extremely helpful, and if anyone has problems in the future docker is a good way to get around old and un-updated environments. It took quite a few commits to get it working (see the testing/travis-ci branch and the plenty of frustrated commits trying to work with docker, bash and everything else to behave properly), but when it worked it was quite helpful for catching

¹unittest is built into the python standard library, ensuring less installation steps to get going: https://docs.python.org/3/library/unittest.html.

²Our builds are here: https://travis-ci.org/ThePhD/lepix/builds.

any bad changes and keeping a log of things that went wrong so it could be looked at later to fix problems.

```
1 dist: trusty
2 sudo: required
4 language: cpp
6 services:
7 - docker
9 before_install:
10 - docker pull ubuntu:latest
11 - docker run -v${PWD}:/ci_repo -d --name lepix_ci ubuntu:latest
      sleep infinity
12 - docker exec lepix_ci bash -e -v -c "apt-get update"
_{13} - docker exec lepix_ci bash -e -v -c "apt install -y git python3
       build-essential m4 autotools-dev autoconf pkg-config ocaml
      menhir opam llvm llvm-dev llvm.3.8"
14 - docker exec lepix_ci bash -e -v -c "opam init -y"
15 - docker exec lepix_ci bash -e -v -c "opam install -y core
      depext llvm.3.8"
17 script:
18 - docker exec lepix_ci bash -e -v -c "source ci_repo/ci/travis.
      sh"
20 after_script:
21 - docker stop lepix_ci
22 - docker rm lepix_ci
24 notifications:
25 email:
26 on_success: change
27 on failure: change
```

Listing 6.4: .travis.yml

6.3 Division of Labor

I wrote a large number of examples and also wrote tests, implemented travis-ci integration, and wrote the python bootstrapper and test suite code.

Post-Mortem and Lessons Learned

This is going to be the most in-depth section because it is here where I can explain primarily why I think the group did not meet its target and why I felt like splitting off would be more worth it than staying with the team. While I individually put in a lot of effort and achieved some very good technical goals, the divide with my group near the end was still a problem and resulted in a lot of codegen for constructs successfully put into the parser and AST to not be implemented.

7.1 Talk to your Teammates, Early

When I experienced problems with my teammates not hitting deadlines, I at first was confused. I did not know why they were not delivering the portions of code they said they would deliver on the deadlines they imposed on themselves, and at certain points when they did deliver I had to constantly revise what they had done. Here are some examples of how I did not optimally handle bad situations:

For one, the Parser and Lexer for this LéPiX implementation look nothing like the one committed and declared to our advisor as "complete". It did not parse our language and there were obvious holes in its syntax: for loops

variable initialization did not work, initializer lists for control flow did not work, parallel block initializers were not considered, the parallel for syntax we changed for a parallel block were not changed, namespaces were not recognized and qualified identifiers did not exist.

Rather than tell my teammates what was wrong and what needed to be fix and divide the work, I instead implemented all of the things mentioned above, committed them, and then moved on. I felt that if my teammates would not run the code against the example LéPiX code we had to see if it works properly, that they were not doing the bare minimum to even know if what they wrote was correct or good. I had to learn everything, put it all together under pressure, and then fix it in time for the next Milestone.

7.2 Manage Expectations, Know What You Want

One of the next major issues is that team members had differing expectations about the quality of work. In particular, I was expecting a very thorough, consistent applied effort from my team and not things done a few weeks after the Professor, TA, and others had urged us needed to be done long before we had begun to look at it.

On the good side, the Language Reference Manual was done on-time with participation from everyone. It was the one part of the project where – even if we were working up to the deadline – everyone participated, took a section, made their work clear and actually did their work during the times they said they would.

Unfortunately, this flopped for actual implementation. One of our group members held onto the Semantic AST for nearly five weeks of time, refusing to commit code when asked and spinning down the time of myself and other group mates eager to get started on Code Generation. The Lexer and Parser were not up to parsing our language. Many disconnects appeared in how the implementation was done, which was entirely strange because we had specifically said we would wait for the Language Reference Manual to be done to begin working so everyone would have a very clear goal and standard.

Talking to your teammates about what exactly is expected, even with a document like the Language Reference Manual, would be helpful in the future. You and your teammates should be able to look at previous projects,

and see

- 1. To achieve X feature it took Y lines of code.
- 2. Is that feasible if you give yourself Z amount of time to write Y lines with W people?
- 3. What quality of implementation do you want? Proof of concept? Fully vetted with compiler errors?

As an example, I wanted full source code information and carat diagnostics throughout the program. I only managed to add that to the first half of the project, and in my lack of help and time for the second half did not implement it for Semantic AST and Codegen errors.

Other groups would consider this silly and not bother with it at all. Your team should agree on just how much effort and polish your implementation deserves, and have a frank discussion about whether people will do that work.

If people impose deadlines on themselves and do not mean them, talk to them immediately about it rather than just implementing it yourself in frustration. Only when they do not respond to your inquiries do you turn to outside sources and begin to re-evaluate what can and cannot be done with your time.

7.3 Start Confrontations

When people in my group slipped deadlines, I vented my frustrations elsewhere while implementing the code just in time for deadlines or pulling together LaTeX documents and editing them furiously. I confronted my team only once very early in September and CCed the professor and a TA with an e-mail, where I demanded they never put me in a situation similar to the one where I wrote the entire LéPiX proposal by myself and then have them – only an hour or so before the deadline – tell me grammatical edits that I needed to fix.

After that, I did not expect to have to send them anymore particularly strongly-worded e-mails. They had agreed not to do something like that

again and indeed everyone participated in the Language Reference Manual. We had communication over GroupMe about why the AST and Semantic AST were not being done on time, but I had not made it clear that their lack of implementation was unacceptable: I only patched it over in the days before the deadline after I had grown tired of waiting and needed to have implementation work done to do my part.

You must have confrontations. You must but heads. Do this early, and do it often when a group member does not hand in their work. Growing frustrated in silence while implementing things you would have expected your teammates to do will only wear you out and ultimately lead you to a place where you will want to discard anything your team does, good or bad, and not take their suggestions in because you feel like they will just let you down.

Appendix

8.1 Source Code Listing

```
1 .PHONY: default
2 default: all;
4 # Clean intermediate files
6 ocamlbuild -use-menhir -build-dir obj -clean
7 rm -rf lepixc lepix
8 rm -rf scanner.ml parser.ml parser.mli
9 rm -rf prescanner.ml preparser.ml preparser.mli
10 rm -rf *.cmx *.cmi *.cmo *.cmx *.o
11 rm -rf parser.automaton preparser.automaton
12 rm -rf parser.output preparser.output parser.conflicts preparser
      . conflicts
14 # Build top level lepix executable
16 ocambuild -use-ocambind -use-menhir -tag thread -pkgs core,
      llvm, llvm.analysis -build-dir obj lepix.native
17 cp -f obj/lepix.native lepixc
19 install:
20 cp lepixc /usr/local/bin/lepixc
22 uninstall:
23 rm -f /usr/local/bin/lepixc
25 .PHONY: all
```

```
26 all: lepix
                      Listing 8.1: source/Makefile
   (* LePiX Language Compiler Implementation
  Copyright (c) 2016- ThePhD
   Permission is hereby granted, free of charge, to any person
       obtaining a copy of this
   software and associated documentation files (the "Software"
      ), to deal in the Software
   without restriction, including without limitation the
      rights to use, copy, modify,
  merge, publish, distribute, sublicense, and/or sell copies
      of the Software, and to
   permit persons to whom the Software is furnished to do so,
      subject to the following
   conditions:
9
10
  The above copyright notice and this permission notice shall
       be included in all copies
   or substantial portions of the Software.
  THE SOFTWARE IS PROVIDED "AS IS", WITHOUT WARRANTY OF ANY
14
      KIND, EXPRESS OR IMPLIED,
  INCLUDING BUT NOT LIMITED TO THE WARRANTIES OF
      MERCHANTABILITY, FITNESS FOR A
  PARTICULAR PURPOSE AND NONINFRINGEMENT. IN NO EVENT SHALL
      THE AUTHORS OR COPYRIGHT
  HOLDERS BE LIABLE FOR ANY CLAIM, DAMAGES OR OTHER LIABILITY
      , WHETHER IN AN ACTION
  OF CONTRACT, TORT OR OTHERWISE, ARISING FROM, OUT OF OR IN
      CONNECTION WITH THE
  SOFTWARE OR THE USE OR OTHER DEALINGS IN THE SOFTWARE. *)
19
20
   (* In Javascript, there's a concept called 'Polyfill'. It's
       the concept that
   stuff that's missing can be filled over by libraries
      implemented by regular people
  because the committee that oversees Javascript can't just
      decide to
  make certain implementations and other things standard.
25
  This is that thing, for OCrapml. *)
26
   (* Algorithm *)
  let foldi f value start index len =
```

```
let end\_index = start\_index + len - 1 in
30
     if start index >= end index then value else
31
     let accumulated = ref value
32
     for i = start_index to end_index do
34
35
       accumulated := ( f !accumulated i )
     done;
36
     !accumulated
37
38
   let foldi to f value start index end index =
39
     foldi f value start_index (end_index - start_index)
40
41
   (* Integer *)
   let rec powin = function
43
     | 0 -> 1
45
       1 \rightarrow n
       x \rightarrow n * (powin x - 1)
   let int of bool b = if b then 1 else 0
49
   let int_of_string_base b s =
50
     let len = (String.length s) in
     let acc num i = let c = s.[i] in
       let v = if c >= '0' \mid \mid c <= '9' then
            int of char c - int of char '0'
54
          else
            if c >= 'A' \mid \mid c <= 'Z'  then
56
              int\_of\_char\ c - int\_of\_char\ 'A' + 10
              if c >= 'a' \mid \mid c <= 'z' then
                int\_of\_char\ c - int\_of\_char\ '0' + 10
              else 0
       and place = len - 1 - i
62
       num + (v * (powi b place))
64
65
     foldi acc 0 0 len
66
67
   (* Num *)
68
69
   exception BadBase of string
   exception DigitGreaterThanBase of string
71
72
   let num_of_string_base_part b s =
73
     if b > 36 | b < 1 then raise (BadBase "num_of_string_base
```

```
: base cannot be greater than 36 or less than 1") else
      let n0 = Num.num of int 0 in
      let len = (String.length s) in
76
      if len < 1 then n0 else
77
      let (mid, starter) = try ( (String.index s '.'), 1)
       with \_ \rightarrow (len - 1, 0)
     and nb = Num.num_of_int b
79
80
      let acc (n, skipval) i = let c = s.[i] in
81
        if c = '.' then (n, skipval - 1) else
82
        let v = if c >= '0' \&\& c <= '9' then
            ( int_of_char c ) - ( int_of_char '0' )
84
          else
            if c >= 'A' \&\& c <= 'Z' then
              (int\_of\_char c) - (int\_of\_char 'A') + 10
            else
88
              if c >= 'a' \&\& c <= 'z' then
                int\_of\_char c - (int\_of\_char 'a') + 10
90
              else 0
        and place = mid - i - skipval
92
93
        if v > b then raise (DigitGreaterThanBase ("
94
       num_of_string_base: digit '" ^ (String.make 1 c) ^ "' ("
        ^ (string_of_int v) ^ ") is higher than what base '" ^
       (string_of_int b) ^" can handle")) else
        let nv = Num.num of int v
95
       and nplace = Num.num_of_int_place
96
97
        (Num.add_num n ( Num.mult_num nv ( Num.power_num nb
98
       nplace ) ), skipval)
99
      let (n, \_) = foldi acc (n0, starter) 0 len in
100
   let num_of_string_base b s =
      num_of_string_base_part b s
104
   let num of string s =
106
      let slen = String.length s in
107
108
        let eidx = String.index s 'e' in
109
        if eidx < 1 then raise (Not found);
110
        let eidxp1 = (eidx + 1) in
111
        let nval = num_of_string_base_part 10 (String.sub s 0
112
       eidx)
```

```
and eval = if eidxp1 < slen then
       num of string base part 10 (String.sub s eidxp1 (slen -
        eidxp1)) else ( Num.num_of_int 0 )
114
        Num.mult_num nval (Num.power_num (Num.num_of_int 10)
115
       eval)
      with
116
        Not found -> num of string base part 10 s
117
118
119
    (* Char *)
    let is\_whitespace = function
120
        -> true
121
        '\t' -> true
122
        '\n' -> true
123
        '\r' -> true
125
        _ -> false
126
    (* String *)
127
    let string\_to\_list s =
128
      let l = ref [] in
129
130
      let acc c =
        1 := c :: !1; ()
132
      in
      String.iter acc s;
133
      List.rev !1
134
135
    let iteri f start_index len =
136
      let end\_index = start\_index + len - 1 in
137
      if start_index < end_index then
138
        for i = start\_index to end\_index do
          ( f i )
140
141
        done
142
    type split_option =
143
        RemoveDelimeter
144
        KeepDelimeter
145
146
    let string_split_with v s opt =
147
      let e = String.length s
148
      and vlen = String.length v
149
      in
150
      if vlen >= e then [s] else
151
      let forward search start =
        let acc found idx =
153
          found && (s.[start + idx] = v.[idx])
```

```
in
        foldi acc true 1 (vlen -1)
156
157
      let add sub len slist start =
158
        if len < 1 then ( start, slist ) else
159
160
        let fresh = (String.sub s start len)
        and last = start + len + vlen in
161
        begin match opt with
162
            RemoveDelimeter -> ( last, fresh :: slist )
163
            KeepDelimeter ->
164
            let slist = v :: slist in
165
            ( last, fresh :: slist )
166
        end
167
      in
168
      let acc (last, slist) start =
        if (start < last) then (last, slist) else
        if (s.[start] = v.[0]) then
171
          if (forward_search start) then
            let len = start - last in
            (add_sub len slist last)
174
175
          else
            (last, slist)
176
        else
          if (start = (e - 1)) then
178
            let len = e - last in
179
            (add sub len slist last)
180
181
          else
            (last, slist)
182
      in
183
      let (\_, slist) = (foldi acc (0, []) 0 e) in
      (* Return complete split list *)
185
      List.rev slist
186
187
   let string_starts_with str pre =
      let prelen = (String.length pre) in
189
      prelen <= (String.length str ) && pre = (String.sub str 0
190
        prelen)
191
   let string\_split v s =
192
      string_split_with v s RemoveDelimeter
193
                           ../source/polyfill.ml
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19
20
   (* Base types and routines. *)
21
22
   type token_source = {
23
        token source name: string;
24
25
     token number : int;
        token_line_number : int;
26
        token_line_start : int;
27
     token_column_range : int * int;
28
     token_character_range : int * int;
29
   }
30
31
   type target =
32
     | Pipe
33
     | File of string
34
35
   let target to string = function
36
       Pipe -> "pipe"
37
     | File(s) -> "file: " ^ s
```

```
39
   let target_to_pipe_string i b = match i with
       Pipe -> if b then "stdin" else "stdout"
41
      File(s) -> "file: " ^ s
42
43
44
   type action =
       Help
45
        Preprocess
46
       Tokens
47
48
        Ast
       Semantic
49
       Llvm
50
       Compile
51
   let action_to_int = function
        Help \rightarrow -1
54
        Preprocess -> 0
55
        Tokens \rightarrow 1
56
        Ast ~-\!\!>~ 10
       Semantic -> 100
58
       Llvm \rightarrow 1000
59
       Compile -> 10000
60
61
   let entry_point_name = "main"
62
63
   (* Core options *)
64
65
   let default_integral_bit_width = 32
66
   let default_floating_bit_width = 64
67
   (* Error message helpers *)
69
70
71
   let line_of_source src token_info =
     let ( absb , abse ) = token_info.token_character_range
     and linestart = token_info.token_line_start
73
74
     in
      let (lineend, _ ) =
75
        let f (endindex, should_skip) idx =
76
          let c = src.[idx] in
77
          let skip\_this = c = '\n' in
78
          if should_skip || skip_this then
79
            (endindex, true)
          else
81
            (endindex + 1, false)
82
       in
```

```
Polyfill.foldi f ( linestart , false ) linestart ( (
84
       String.length src ) - linestart )
85
      let srcline = String.sub src linestart (max 0 (lineend -
86
       linestart - 1) in
      let srclinelen = String.length srcline in
87
      let ( srcindent , _ ) =
        let f(s, should skip) idx =
89
          let c = srcline.[idx] in
90
          let nws = not ( Polyfill.is_whitespace c ) in
91
          if should_skip || nws then
            (s, false)
93
          else
94
            (s ^ (String.make 1 c), true)
95
        in
        Polyfill.foldi f ("", false ) 0 srclinelen
97
98
      let indentlen = String.length srcindent
99
     and tokenlen = lineend - absb
100
      in
      ( srcline, srcindent, (max ( srclinelen - indentlen -
102
       tokenlen ) 0 )
104
   let brace tabulate str tabs =
105
      let len = ( String.length str ) in
106
      let lines = Polyfill.string_split_with "\n" str Polyfill.
107
       KeepDelimeter in
      let lineslen = ( List.length lines ) in
108
      let buf = Buffer.create (len + (lineslen * 4)) in
109
      let acc (buf, t) line =
110
        let tmod = 0 - ( Polyfill.int_of_bool ( String.contains
111
        line '}' ) ) in
        let t = t + tmod in
        Buffer.add\_string\ buf\ (String.make\ t\ '\t');\ Buffer.
       add_string buf line;
       let t = t + ( Polyfill.int_of_bool ( String.contains
114
       line '{' ) ) in
        (buf, t)
115
116
      let (buf, _) = List.fold_left acc ( buf, tabs ) lines in
117
      Buffer.contents buf
118
                            ../source/base.ml
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19
20
   (* Options / argument parser *)
21
22
   type option =
      Dash of string
24
       DoubleDash of string
25
      Argument of int * string
26
27
   type options_context = {
28
        mutable options help: string -> string;
29
30
31
   let read_options ocontext sys_argv =
     let argc = Array.length sys argv - 1 in
     (* Skip first argument one (argv 0 is the path
34
     of the exec on pretty much all systems) *)
35
     let argv = ( Array.sub sys_argv 1 argc )
```

```
and action = ref Base. Help
37
     and verbose = ref false
38
     and input = ref Base. Pipe
39
     and output = ref Base. Pipe
40
     and specified = ref []
41
     and seen_stdin = ref false
42
     (* Our various options *)
44
     let update action a =
45
       specified := a :: !specified;
46
       if ( Base.action_to_int !action ) < ( Base.
       action_to_int a ) then
         action := a;
48
     in
49
     let options = [
       (1, "h", "help", "print the help message",
51
         fun _ _ -> ( update_action(Base.Help) )
       (1, "p", "preprocess", "Preprocess and display source"
         fun _ _ -> ( update_action(Base.Preprocess) )
56
       ( 1, "i", "input", "Take input from standard in (
57
       default: stdin)",
         fun -> ( input := Base.Pipe; seen stdin := true )
58
59
       ( 2, "o", "output", "Set the output file (default:
60
       stdout)",
         fun _o \rightarrow (output := Base.File(o))
61
       ( 1, "t", "tokens", "Print the stream of tokens",
63
         fun _ _ -> ( update_action(Base.Tokens) )
65
       (1, "a", "ast", "Print the parsed Program",
         fun _ _ -> ( update_action(Base.Ast) )
67
       ( 1, "s", "semantic", "Print the Semantic Program",
         fun \ \_ \ \_ \ -> \ ( \ update\_action ( Base. Semantic ) \ )
70
71
       ( 1, "l", "llvm", "Print the generated LLVM code",
72
         fun _ _ -> ( update_action(Base.Llvm) )
73
       (1, "c", "compile", "Compile the desired input and
75
       output the final LLVM",
         fun _ _ -> ( update_action(Base.Compile) )
```

```
77
        ('1, "v", "verbose", "Be as explicit as possible with
       all steps",
          fun _ - -> (verbose := true)
79
        );
80
81
      and position_option arg_index positional_index arg =
82
        if Sys. file exists arg then
83
          input := Base. File ( arg )
84
85
          raise (Errors. OptionFileNotFound (arg))
86
      in
87
      let help tabulation =
        let value_text = "<value>" in
89
        let value_text_len = String.length value_text in
        let longest_option =
91
          let acc len o = match o with
            | (sz, _, long, _, _) ->
              let newlen = (String.length long)
                + if sz = 2 then 1 + value_text_len else 0
95
96
              if newlen > len then newlen else len
97
98
          let l = 1 + ( List.fold_left acc 1 options ) in
99
          if I < value text len then value text len else I
100
102
        let concat_options t =
103
          let builder s o = match o with
            | ( sz, short, long, desc, _ ) ->
104
              let long_len = String.length long in
              let spacing_size = longest_option - long_len - (
106
       if sz = 2 then 1 + value_text_len else 0 ) in
              let spacing_string = (String.make spacing_size '
       ') in
              s
^ "\n" ^ t ^ "-" ^ short
108
109
               \hat{t} = 0  \int t dt = 0  \int t dt = 0  \int t dt = 0 
       value_text)
              ^ spacing_string
111
              \hat{} desc
113
          (List.fold left builder "" options)
114
        in
115
        let (_, input_short, input_long, _, _) = List.nth
116
       options 2 in
```

```
let msg = "Help:"
117
          ^ "\n" ^ tabulation ^ "lepix [options] filename [
118
       filenames ...]"
          ^ "\n" ^ tabulation ^ "\t" ^ "filename | filenames
119
       can have one option -" ^ input_short ^ " or --" ^
       input_long
          `"\n" `tabulation `"options:"
120
          ^ ( concat_options (tabulation ^ "\t") )
121
       in
122
       msg
123
124
     in
     ocontext.options_help <- help;
125
      (* Exit early if possible *)
126
      if argc < 1 then
127
        (!input, !output, !action, !specified, !verbose)
      else
129
130
      let to_option idx arg =
       let arglen = String.length arg in
       match arg with
          when Polyfill.string_starts_with arg "--" ->
134
       DoubleDash ((String.sub arg 2 (arglen - 2)))
          | _ when Polyfill.string_starts_with arg "-" -> Dash
135
       ((String.sub arg 1 (arglen - 1)))
          -> Argument (idx, arg)
136
137
138
      (* Convert all arguments to the Option type first *)
      let options_argv = Array.mapi to_option argv in
139
140
      (* Function for each argument *)
      let f (index, positional_index, skip_next) option_arg =
142
143
        if skip_next then (1 + index, positional_index, false)
       let execute_on_match_sub_option ( opt_failure ,
       should_block ) opt_string pred = match opt_failure with
          (* There is some failure, so just propagate it
145
       through *)
          | Some(x) -> ( opt_failure, should_block )
146
          (* There is no failure, so now work with the list *)
147
          | None -> begin match List. filter pred options with
148
            (* We use filter instead of find because find is
149
       dumb and throws an
            exception instead of just returning an optional
150
            whoever designed the OCaml standard library is an
151
```

```
absolute
            bell end. *)
            \mid ( 1, _, _, _, f ) :: tail \rightarrow (* Only needs 1
       argument *)
               (f opt_string "");
154
               ( opt_failure, should_block )
             | (2, _-, _-, _-, _f) :: tail \rightarrow (* Needs 2 arguments) |
156
       , look ahead by 1 *
               if (index + 1) >= argc then
157
                 raise (Errors. Missing Option (opt_string));
158
               let nextarg = ( options_argv.(1 + index) ) in
               let _ = match nextarg with
                   Argument(idx, s) -> (f opt_string s)
                   _ -> raise (Errors.BadOption(opt_string))
162
163
               in
               ( opt_failure, true )
164
              _ -> (* Unhandled case: return new failure string
165
        *)
               (Some opt string, should block)
166
            \quad \text{end} \quad
167
168
        and on_failure dashes opt arglist arg =
          let msg = dashes \hat{ } opt
169
              if (List.length arglist) > 1 then "(in "^
170
       dashes ^ arg ^ " )" else ""
          in
171
          raise (Errors . BadOption (msg))
172
173
        let ( should_skip_next , was_positional ) = match
174
       option_arg with
          | Dash(arg) ->
            (* if it has a dash only *)
176
            (* each letter can be its own thing *)
177
            let perletter (opt_failure, should_break) c =
178
               let opt_string = ( String.make 1 c ) in
               let short_pred (_, short, _, _, _) =
180
                 short = opt\_string
181
182
               execute_on_match_sub_option (opt_failure,
183
       should_break) opt_string short_pred
184
            (* look at every character. If there's 1 match
185
       among them, go crazy *)
            let arglist = (Polyfill.string_to_list arg) in
186
            let (opt_failure, causes_skip) = ( List.fold_left
187
       perletter ( None, false ) arglist ) in
```

```
begin match opt_failure with
188
               | None ->
189
                 ( causes_skip, 0 )
190
               | Some(opt) -> let _- = (on_failure "--" opt
191
       arglist arg) in
                 ( causes_skip, 0 )
192
            end
193
            DoubleDash (arg) ->
194
            (* if it has a double dash... *)
195
            (* each comma-delimeted word can be its own option
196
       *)
            let perword (opt_failure, problems) opt_string =
197
              let long_pred (_, _, long, _, _) =
                long = opt_string
199
              in
200
              execute_on_match_sub_option (opt_failure,
201
       problems) opt_string long_pred
            in
202
            (* look at word character. If there's 1 match among
203
        them, go crazy *)
            let arglist = Polyfill.string_split "," arg in
204
            let (opt_failure, causes_skip) = ( List.fold_left
205
       perword ( None, false ) arglist ) in
            begin match opt_failure with
206
               \mid None \rightarrow
207
                 ( causes_skip, 0 )
208
               | Some(opt) -> let _- = (on_failure "---" opt
209
       arglist arg) in
                 ( causes_skip, 0 )
210
            end
211
          (* otherwise, it's just a positional argument *)
212
          | Argument(idx, arg) ->
            (position_option index positional_index arg);
214
             ( skip_next, 1 )
216
        (1 + index, positional_index + was_positional,
217
       should_skip_next)
218
      (* Iterate over the arguments *)
219
      let _ = Array.fold_left f (0, 0, false) options_argv in
220
      (* Return tuple of input, output, action *)
221
      (!input, !output, !action, !specified, !verbose)
                           ../source/options.ml
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19
20
   (* Message formatters and helpers. *)
21
22
23
   let preprocessing error prontext =
24
     let (t, info) = pcontext. Predriver. token in
     let ( source_line , source_indentation ,
26
      columns_after_indent ) =
       ( Base.line of source pcontext.Predriver.source code
      info )
     in
28
     let column_range = info.Base.token_column_range in
29
     let (column_text, is_columns_wide) = Representation.
      token range to string column range in
     let msg = "Preprocessing Error in " ^ pcontext. Predriver.
31
      source_name ^ ":"
     ^ "\n" ^ "\t" ^ "Unrecognizable parse pattern at token \#"
```

```
\hat{\ } ( <code>string_of_int pcontext.Predriver.token_count )</code>
       \hat{\ } ": [id " \hat{\ } string_of_int info.Base.token_number
33
       ^ Representation.preparser_token_to_string t ^ "]"
     ^ "\n" ^ "\t" ^ "Line: " ^ string_of_int info.Base.
34
      token_line_number
     ^ "\n" ^ "\t" ^ ( if is_columns_wide then "Columns: "
35
      else "Column: " ) ^ column_text
     ^ "\n"
36
     ^ "\n" ^ source_line
37
     ^ "\n" ^ source_indentation ^ ( String.make
38
      columns_after_indent ' ') ^ "^~~"
39
     msg
40
41
   let preprocessing_lexer_error pcontext core_msg c s e =
     let ( t, info ) = pcontext.Predriver.token in
43
     let ( source_line , source_indentation ,
44
      columns_after_indent ) =
       ( Base.line of source pcontext.Predriver.source code
      info )
46
     let column_range = info.Base.token_column_range in
47
     let (column_text, is_columns_wide) = Representation.
48
      token_range_to_string column_range in
     let msg = "Preprocessing Lexing Error in " ^ pcontext.
49
      Predriver.source_name ^ ":"
     ^ "\n" ^ "\t" ^ core_msg ^ " at token#" ^ ( string_of_int
50
       pcontext.Predriver.token_count )
        ": [id " ^ string_of_int info.Base.token_number ^ ":"
       ^ Representation.preparser_token_to_string t ^ "]"
     ^ "\n" ^ "\t" ^ "Line: " ^ string_of_int info.Base.
      token_line_number
     ^ "\n" ^ "\t" ^ ( if is_columns_wide then "Columns: "
      else "Column: " ) ^ column_text
     ^ "\n"
54
     ^ "\n" ^ source_line
     columns_after_indent ' ') ^ ; ^ ~ ~ ?
     in
57
58
59
   let lexer error context core msg c s e =
60
     let abspos = s. Lexing.pos cnum in
61
     let endabspos = e. Lexing. pos cnum in
62
     let relpos = 1 + abspos - s.Lexing.pos_bol in
```

```
let endrelpos = 1 + endabspos - e.Lexing.pos_bol in
64
               let (column text, is columns wide) = Representation.
65
                  token_range_to_string ( relpos, endrelpos ) in
               let msg = "Lexing Error in " ^ context.Driver.source_name
              ^ "\n" ^ "\t" ^ core_msg ^ " at character: " ^ \ensuremath{\mathbf{c}}
67
               ^ "\n" ^ "\t" ^ "Line: " ^ string_of_int s.Lexing.
68
                 pos lnum
               ^{\text{``}} \ ^{\text{``}} 
69
                  else "Column: " ) ^ column text
70
              msg
71
72
        let parser_error context core_msg =
73
               let ( t, info ) = context.Driver.token in
               let ( source_line , source_indentation ,
                  columns_after_indent ) =
                    ( Base.line_of_source context.Driver.source_code info )
76
77
               let column_range = info.Base.token_column_range in
78
               let (column text, is columns wide) = Representation.
                  token_range_to_string column_range in
               let msg = "Parsing Error in " ^ context.Driver.
80
                 source_name ^ ":"
               ^{\circ} "\n" ^{\circ} "\t" ^{\circ} core msg ^{\circ} " at token #" ^{\circ} (
81
                  string of int context. Driver. token count )
                    ^ ": [id " ^ string_of_int info.Base.token_number ^ ":"
82
                     ^ Representation.parser_token_to_string t ^ "]"
               ^ "\n" ^ "\t" ^ "Line: " ^ string_of_int info.Base.
83
                 token_line_number
               ^ "\n" ^ "\t" ^ ( if is_columns_wide then "Columns: "
84
                  else "Column: " ) ^ column_text
85
               ^ "\n" ^ source_line
              ^ "\n" ^ source_indentation ^ ( String.make
87
                  columns_after_indent ' ') ^ ", ^ ~ ~ ~ ,
              in
88
89
                                                                       ../source/message.ml
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19
20
   (* Top-level of the LePiX compiler: scan & parse the input,
   check the resulting AST, generate LLVM IR, and dump the
      module *)
23
24
     let input = ref Base. Pipe in
25
     let output = ref Base. Pipe in
27
     let action = ref Base. Compile in
     let verbose = ref false in
     let specified = ref [] in
29
     let context = {
       Driver.source name = "";
31
       Driver.source_code = "";
32
       Driver.original_source_code = "";
33
       Driver.token count = 0;
34
       Driver.token = ( Parser.EOF,
         { Base.token_source_name = ""; Base.token_number = 0;
         Base.token line number = 0; Base.token line start =
37
         Base.token_column_range = (0, 0); Base.
```

```
token\_character\_range = (0, 0) 
       );
39
     }
40
     and pcontext = \{
41
          Predriver.source_name = "";
42
       Predriver.source_code = "";
43
       Predriver.original_source_code = "";
44
       Predriver.token count = 0;
       Predriver.token = ( Preparser.EOF,
46
         { Base.token source name = ""; Base.token number = 0;
47
         Base.token_line_number = 0; Base.token_line_start =
       0;
         Base.token_column_range = (0, 0); Base.
49
       token\_character\_range = (0, 0) 
50
51
     in
     let ocontext = {
       Options.options help = fun (s) \rightarrow ("");
     } in
     (* Call options Parser for Driver *)
56
     let _ =
57
     try
58
       let ( i, o, a, s, v ) = ( Options.read_options ocontext
59
        Sys.argv ) in
         input := i;
60
61
         output := o;
         action := a;
62
         specified := s;
63
         verbose := v
     with
65
       err -> let _ = match err with
          | Errors.BadOption(s) ->
67
           let msg = "Options Error:"
            ^ "\n" ^ "\t" ^ "Unrecognized option: " ^s
            ^{"} "\" ^{"} ( ocontext. Options.options_help "\t" ) in
            prerr_endline msg
           Errors. NoOption ->
72
           let msg = "Options Error:"
73
            ^ "\n" ^ "\t" ^ "No inputs or options specified"
74
            ^ "\n" ^ ( ocontext.Options.options_help "\t" ) in
            prerr endline msg
76
          | Errors. MissingOption(o) ->
            let msg = "Options Error:"
            ^ "\n" ^ "\t" ^ "Flag " ^ o ^ " needs an additional
```

```
argument after it that is not dashed"
            80
            prerr_endline msg
81
          | Errors.OptionFileNotFound(f) ->
82
            let msg = "Options Error:"
^ "\n" ^ "\t" ^ "File " ^ f ^ " was not found"
84
            ^{"}n" ^{"} ( ocontext. Options.options_help ^{"}t" ) in
            prerr endline msg
           err ->
87
            let msg = "Unknown Error during Option parsing:"
88
            ^{"}n"^{"}t"^{"} Contact the compiler vendor for
89
       more details and possibly include source code, or try
       simplifying the program"
90
            prerr_endline msg;
            raise (err)
92
        in
93
        (* Exit if arguments are wrong *)
94
        ignore ( exit Errors.option error exit code ) )
96
     in
      (* Perform actual lexing and parsing using the Driver
97
       here *)
     try
98
       let allactions = !specified in
99
       let source name = ( Base.target to pipe string !input
100
       true ) in
       let output_to_target ( s ) = match !output with
101
            Base.Pipe -> ( print_endline s )
102
            Base.File(f) -> ( Io.write_file_text s f )
103
       in
104
        let print_predicate b =
105
          fun v \rightarrow (v = b)
106
        let print_help () =
          let msg = ( ocontext.Options.options_help "\t" ) in
          print_endline msg
       in
111
        if !action = Base. Help then begin
112
          print_help ()
       end else
114
        (* Since we do the actions in these functions multiple
115
       We refactor them out here to make our lives easier
116
       while we tweak
        stuff *)
117
```

```
118
        let get_source () =
          let pre source text = match !input with
119
              Base.Pipe -> Io.read_text stdin
120
              Base.File(f) -> ( Io.read_file_text f )
          in
          let source_text = Predriver.pre_process pcontext !
123
       input pre_source_text in
          context. Driver. source name <- source name;
124
          context.Driver.original source code <-
125
       pre source text;
          context.Driver.source_code <- source_text;</pre>
126
          source_text
127
        in
128
        let dump_tokens f tokenstream =
129
          if (List.exists (print_predicate Base.Tokens)
130
       allactions ) then f( Representation.
       parser_token_list_to_string_tokenstream )
        and dump_ast f program =
          if (List.exists (print predicate Base.Ast)
       allactions ) then f( Representation.string_of_program
       program )
        and dump_semantic f semanticprogram =
          if ( List.exists (print_predicate Base.Semantic)
       allactions ) then f( Representation.string_of_s_program
       semanticprogram )
        and dump module f m =
135
          f(Llvm.string_of_llmodule m)
136
        in
137
            _{-} = match ! action with
138
            Base. Help -> print_help ()
            Base. Preprocess ->
140
            let source_text = get_source () in
            output_to_target( source_text )
142
            Base. Tokens \rightarrow
            let source_text = get_source () in
144
            let lexbuf = Lexing.from_string source_text in
145
            let tokenstream = Driver.lex source_name lexbuf in
146
            ( dump tokens output to target tokenstream )
147
            Base . Ast \rightarrow
148
            let source_text = get_source () in
149
            let lexbuf = Lexing.from_string source_text in
150
            let tokenstream = Driver.lex source name lexbuf in
151
            ( dump tokens print endline tokenstream );
152
            let program = Driver.parse context tokenstream in
153
            ( dump_ast output_to_target program)
```

```
Base. Semantic ->
            let source text = get source () in
156
            let lexbuf = Lexing.from_string source_text in
            let tokenstream = Driver.lex source name lexbuf in
158
            ( dump_tokens print_endline tokenstream );
159
            let program = Driver.parse context tokenstream in
160
            ( dump_ast print_endline program );
            let semanticprogram = Driver.analyze program in
            ( dump semantic output to target semanticprogram )
163
            Base . Llvm \rightarrow
164
            let source_text = get_source () in
165
            let lexbuf = Lexing.from_string source_text in
            let tokenstream = Driver.lex source_name lexbuf in
            ( dump tokens print endline tokenstream );
168
            let program = Driver.parse context tokenstream in
            ( dump_ast print_endline program );
            let semanticprogram = Driver.analyze program in
            ( dump_semantic print_endline semanticprogram );
            let m = Codegen.generate semanticprogram in
            if !verbose then ( dump_module print_endline m );
174
175
            ( dump module output to target m )
            Base. Compile ->
            let source_text = get_source () in
            let lexbuf = Lexing.from_string source_text in
178
            let tokenstream = Driver.lex source name lexbuf in
179
            ( dump tokens print endline tokenstream );
180
            let program = Driver.parse context tokenstream in
181
            ( dump_ast print_endline program );
182
            let semanticprogram = Driver.analyze program in
183
            ( dump_semantic print_endline semanticprogram );
            let m = Codegen.generate semanticprogram in
185
            Llvm_analysis.assert_valid_module m;
            ( dump module output to target m )
187
        in
        ()
189
      with
190
          err -> let _ = match err with
191
          (* Preprocessor-Specific Errors *)
192
          (* Preprocessing Parser Errors *)
193
           Preparser. Error ->
194
            let msg = Message.preprocessing_error pcontext in
            prerr endline msg
196
          | Errors.PreUnknownCharacter(c, (s, e)) ->
197
            let msg = Message.preprocessing_lexer_error
198
       pcontext "Unrecognized character in program" c s e in
```

```
prerr_endline msg
199
200
          (* General Compiler Errors *)
201
          (* Lexer Errors *)
202
          | Errors. UnknownCharacter(c, (s, e)) ->
203
            let msg = Message.lexer_error context "Unrecognized
204
        character in program" c s e in
            prerr endline msg
205
206
           Errors.BadNumericLiteral(c, (s, e)) ->
207
            let msg = Message.lexer_error context "Bad
208
       character in numeric literal" c s e in
            prerr_endline msg
209
210
          (* Parser Errors *)
            Parser. Error
212
            Parsing.Parse_error ->
213
            let msg = Message.parser_error context "
214
       Unrecognizable parse pattern" in
            prerr_endline msg
216
          | Errors.MissingEoF ->
            let msg = "Parsing Error in" ^ context. Driver.
217
                     ^ ":"
       source name
            ^{-13.13} ^{\circ} "\n" ^{\circ} "\t" ^{\circ} "Missing EoF at end of token stream
218
         (bad lexer input?)"
219
            prerr_endline msg
220
221
          (* Semantic Analyzer and Codegen Errors *)
222
          (* Semantic Errors *)
          (* TODO: positional information should be tracked
224
       through the AST and SemAST,
          all the way to codegen, as well... *)
           | Errors.BadFunctionCall(s) ->
            let msg = "Bad Function Call error: " ^ s
227
            in
            prerr_endline msg
            Errors. FunctionAlreadyExists(s) ->
230
            let msg = "Function Already Exists error: " ^ s
231
232
            prerr_endline msg
233
            Errors. Variable Already Exists (s) ->
234
            let msg = "Variable Already Exists error: " ^ s
235
236
            prerr_endline msg
237
```

```
Errors. TypeMismatch(s) ->
238
            let msg = "Mismatched types error: " ^ s
239
240
            prerr_endline msg
241
           Errors.IdentifierNotFound(s) ->
242
            let msg = "Identifier Not Found error: " ^ s
243
            prerr endline msg
245
            Errors. Invalid Function Signature (s, n) ->
246
            let msg = "Invalid signature: " ^ s ^ " in " ^ n
247
            in
            prerr_endline msg
249
          | Errors.InvalidMainSignature(s) ->
            let msg = "Invalid signature: " ^ s
251
            prerr_endline msg
253
            Errors. InvalidBinaryOperation(s)
            Errors.InvalidUnaryOperation(s) ->
255
            let msg = "Invalid operation: " ^ s
256
            in
257
258
            prerr_endline msg
259
          (* Direct Codegen Errors *)
260
            Errors. UnknownVariable(s) ->
261
            let msg = "Codegen (LLVM IR) error: " ^ s
262
263
            prerr_endline msg
264
          | Errors. UnknownFunction(s) ->
265
            let msg = "Codegen (LLVM IR) error: " ^ s
266
            in
            prerr_endline msg
268
            Errors. VariableLookupFailure (name, _) ->
            let msg = "Codegen (LLVM IR) error: could not
270
       properly find variable with the name " ^ name
            in
            prerr_endline msg
272
            Errors.FunctionLookupFailure(name, mangledname) ->
            let msg = "Codegen (LLVM IR) error: looking for the
274
        function with the name " ^ name ^ " (mangled name: " ^
       mangledname ^ ")"
            in
275
            prerr endline msg
276
          | Errors.BadPrintfArgument ->
277
            let msg = "Codegen (LLVM IR) error: lib.print and
278
       related functions only take either a string, an integer,
```

```
or a floating point argument"
279
            prerr_endline msg
280
281
          (* Common Errors *)
282
          (* Missing File/Bad File Name, Bad System Calls *)
283
            Sys_error(s) ->
            let msg = "Sys\_error: \n\t" \hat{s}
285
286
            prerr_endline msg
287
          (* Unsupported features *)
289
            Errors. Unsupported(s) ->
290
            let msg = "Unsupported (ran out of implementation
291
       time): \n \ s
            in
292
            prerr_endline msg
293
294
          (* Unknown Errors *)
295
           err ->
296
            let msg = "Unknown Error during Compilation:"
297
            ^ "\n" ^ "\t" ^ "Contact the compiler vendor for
298
       more details and possibly include source code, or try
       simplifying the program"
299
            prerr endline msg;
300
            raise (err)
301
302
        in
        ignore( exit Errors.compiler_error_exit_code )
303
                            ../source/lepix.ml
   (* LePiX Language Compiler Implementation
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  OF CONTRACT, TORT OR OTHERWISE, ARISING FROM, OUT OF OR IN
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20
   (* A listing of exceptions and the methods that power them
   to make the parser more expressive *)
   (* Driver and Related class of errors *)
24
   let option error exit code = 1
25
26
   (* Option Errors *)
27
   exception NoOption
   exception BadOption of string
   exception MissingOption of string
   exception OptionFileNotFound of string
31
32
   (* Compiler class of Errors *)
33
  let compiler_error_exit_code = 2
   (* Lexer Errors *)
  exception PreUnknownCharacter of string * ( Lexing.position
       * Lexing. position )
   exception UnknownCharacter of string * ( Lexing.position *
      Lexing. position )
   exception BadNumericLiteral of string * ( Lexing.position *
       Lexing.position )
39
   (* Parser Errors *)
   exception MissingEoF
41
   exception BadToken
42
  (* Semantic and Codegen Errors *)
  exception Unsupported of string
   exception FunctionAlreadyExists of string
```

```
exception VariableAlreadyExists of string
   exception IdentifierNotFound of string
   exception TypeMismatch of string
   exception BadFunctionCall of string
   exception InvalidMainSignature of string
   exception InvalidFunctionSignature of string * string
   exception InvalidBinaryOperation of string
   exception InvalidUnaryOperation of string
   (* Codegen Errors *)
56
   exception UnknownVariable of string
   exception UnknownFunction of string
   exception BadPrintfArgument
  exception FunctionLookupFailure of string * string
   exception VariableLookupFailure of string * string
                          ../source/errors.ml
   (* LePiX Language Compiler Implementation
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```

```
SOFTWARE OR THE USE OR OTHER DEALINGS IN THE SOFTWARE. *)
20
   (* Drives the typical lexing and parsing algorithm
21
   while adding pertinent source, line and character
      information. *)
23
   type context = \{
24
        mutable source name: string;
25
        mutable source_code : string;
26
27
     mutable original source code : string;
     mutable token_count : int;
28
        mutable token : Parser.token * Base.token_source;
29
   }
30
31
   let lex sourcename lexbuf =
     let rec acc lexbuf tokennumber tokens =
33
       let next_token = Scanner.token lexbuf
       and startp = Lexing.lexeme_start_p lexbuf
35
       and endp = Lexing.lexeme end p lexbuf
       in
       let line = startp. Lexing.pos lnum
       and relpos = (1 + startp.Lexing.pos\_cnum - startp.
39
       Lexing.pos bol)
       and endrelpos = (1 + endp.Lexing.pos_cnum - endp.Lexing
40
       .pos bol)
       and abspos = startp.Lexing.pos cnum
41
42
       and endabspos = endp. Lexing.pos_cnum
       in
43
       let create_token token =
44
         let t = ( token, { Base.token_source_name =
      sourcename; Base.token number = tokennumber;
           Base.token_line_number = line; Base.
       token line_start = startp.Lexing.pos_bol;
           Base.token_column_range = (relpos, endrelpos); Base
       .token_character_range = (abspos, endabspos) }
         ) in
         t
49
       in
       match next_token with
       | Parser.EOF as token -> ( create_token token ) ::
       tokens
       token -> ( create token token ) :: ( acc lexbuf ( 1 +
       tokennumber ) tokens )
54
     acc lexbuf 0 []
```

```
let parse context token list =
     (* Keep a reference to the original token list
     And use that to dereference rather than whatever crap we
      get from
60
     the channel *)
     let tokenlist = ref(token_list) in
61
     let tokenizer _ = match !tokenlist with
     (* Break each token down into pieces, info and all*)
     | (token, info) :: rest \rightarrow
64
       context.source_name <- info.Base.token_source_name;</pre>
       context.token_count <- 1 + context.token_count;</pre>
66
       context.token <- ( token, info );
       (* Shift the list down by one by referencing
       the beginning of the rest of the list *)
       tokenlist := rest;
70
       (* return token we care about *)
       token
72
     (* The parser stops calling the tokenizer when
     it hits EOF: if it reaches the empty list, WE SCREWED UP
      * )
     [] -> raise (Errors.MissingEoF)
75
76
     (* Pass in an empty channel built off a cheap string
77
     and then ignore the fuck out of it in our 'tokenizer'
     internal function *)
79
     let program = Parser.program tokenizer (Lexing.
      from_string "") in
     program
81
   let analyze program =
83
     (* TODO: other important checks and semantic analysis
     that will create a proper checked program type*)
     let sem = Semant.check program in
86
                          ../source/driver.ml
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  SOFTWARE OR THE USE OR OTHER DEALINGS IN THE SOFTWARE. *)
19
20
   (* Routines for preprocessing source code. *)
21
22
   type pre context = {
23
24
        mutable source_name : string;
        mutable source_code : string;
25
        mutable original_source_code : string;
26
     mutable token_count : int;
        mutable token : Preparser.token * Base.token_source;
28
29
   }
30
   let pre_lex sourcename lexbuf =
     let rec acc lexbuf tokens tokennumber =
32
       let next_token = Prescanner.token lexbuf
       and startp = Lexing.lexeme start p lexbuf
34
       and endp = Lexing.lexeme end p lexbuf
       in
36
       let line = startp.Lexing.pos_lnum
37
       and relpos = (1 + startp.Lexing.pos_cnum - startp.
      Lexing.pos bol)
       and endrelpos = (1 + endp. Lexing.pos cnum - endp. Lexing
39
      . pos bol)
       and abspos = startp.Lexing.pos_cnum
40
```

```
and endabspos = endp. Lexing.pos_cnum
41
42
       let create_token token =
43
         let t = ( token, { Base.token_source_name =
44
      sourcename; Base.token_number = tokennumber;
           Base.token line number = line; Base.
45
      token line start = startp.Lexing.pos bol;
           Base.token column range = (relpos, endrelpos); Base
46
      .token character range = (abspos, endabspos) }
         ) in
47
         t
48
       in
49
       let rec matcher = function
           [] -> raise (Errors.MissingEoF)
         | Preparser.EOF :: [] -> ( create_token Preparser.EOF
       ) :: tokens
         token :: [] -> ( create_token token ) :: ( acc
      lexbuf tokens (1 + tokennumber)
         | token :: rest -> ( create token token ) :: (
      matcher rest )
       in matcher next token
     in acc lexbuf [] 0
56
57
   let pre_parse context token_list =
58
     (* Keep a reference to the original token list
59
     And use that to dereference rather than whatever crap we
60
      get from
     the channel *)
61
     let tokenlist = ref(token_list) in
62
     let tokenizer _ = match !tokenlist with
     (* Break each token down into pieces, info and all*)
64
     (token, info) :: rest ->
       context.source_name <- info.Base.token_source_name;</pre>
66
       context.token_count <- 1 + context.token_count;</pre>
       context.token <- ( token, info );</pre>
68
       (* Shift the list down by one by referencing
       the beginning of the rest of the list *)
       tokenlist := rest;
       (* return token we care about *)
72
73
     (* The parser stops calling the tokenizer when
74
     it hits EOF: if it reaches the empty list, WE SCREWED UP
      * )
     [] -> raise (Errors.MissingEoF)
76
     in
77
```

```
(* Pass in an empty channel built off a cheap string
      and then ignore the fuck out of it in our 'tokenizer'
      internal function *)
80
      let past = Preparser.source tokenizer (Lexing.from string
81
        "") in
      past
82
83
   let rec pre process context source source text =
84
      let source name = Base.target to string source in
85
      context.source name <- source name;</pre>
86
      context.source_code <- source_text;</pre>
      let reldir = match source with
88
          Base.Pipe -> ( Sys.getcwd () )
          Base. File (f) -> Filename. dirname f
90
      in
      let generate v p = match p with
92
          Preast.Text(s) \rightarrow v \hat{s}
        | Preast.ImportString(f) \rightarrow v ^{\circ} "\"" ^{\circ} Io.
94
       read_file_text (Filename.concat reldir f) ^ "\""
        | Preast.ImportSource(f) \rightarrow let realf = (Filename.)
95
       concat reldir f) in
          let ftext = Io.read_file_text realf in
96
          let processedtext = ( pre_process context ( Base.File
97
       (f) ) ftext ) in
          v ^ processedtext
98
99
      let tokenstream = pre_lex source_name ( Lexing.
100
       from_string source_text ) in
      (*TODO: debug shit tokens at a later date*)
101
      (*print_endline ( Representation.
102
       preparser_token_list_to_string tokenstream );*)
      let past = pre_parse context tokenstream in
103
104
      List.fold left generate "" past
                           ../source/predriver.ml
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  SOFTWARE OR THE USE OR OTHER DEALINGS IN THE SOFTWARE. *)
19
20
   (* Source types for preprocessing LePiX source code. *)
21
22
   type pre_blob =
23
       Text of string
24
       ImportString of string
      ImportSource of string
26
27
   type pre_source = pre_blob list
                          ../source/preast.ml
   (* LePiX Language Compiler Implementation
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19
   (* Types and routines for the abstract syntax tree and
21
   representation of a LePiX program. *)
   type id = string
24
25
   type qualified id = id list
26
27
   type builtin_type =
28
       Auto
29
       Void
30
       Bool
31
       Int of int
       Float of int
       String
34
     Memory
35
36
   type constness = bool
   type referenceness = bool
38
   type type_qualifier = constness * referenceness
40
41
   type type_name =
42
       BuiltinType of builtin_type * type_qualifier
       Array of type_name * int * type_qualifier
44
      SizedArray of type_name * int * int list *
45
      type_qualifier
     | Function of type name * type name list * type qualifier
46
47
   let no_qualifiers = (false, false)
48
```

```
let void_t = BuiltinType(Void, no_qualifiers)
   let string t = BuiltinType (String, no qualifiers)
   let int32_t = BuiltinType(Int(Base.
       default_integral_bit_width), no_qualifiers)
   let float64_t = BuiltinType(Float(64), no_qualifiers)
54
   type binding = id * type_name
55
56
   let add const (id, t) = match t with
57
     | BuiltinType(bt, tq) \rightarrow let( _, refness) = tq in
58
        ({\tt id}\;,\;\; {\tt BuiltinType}\,(\,{\tt bt}\;,\;\;(\,{\tt true}\;,\;\;{\tt refness}\,)\,)\,)
59
     | Array(tn, d, tq) \rightarrow let (_, refness) = tq in
60
        (id, Array(tn, d, (true, refness)))
     | SizedArray(tn, d, il, tq) -> let (_, refness) = tq in
62
        (id, SizedArray(tn, d, il, (true, refness)))
     | Function(tn, pl, tq) \rightarrow let (_, refness) = tq in
64
        (id, Function(tn, pl, (true, refness)))
65
66
   type binary op = Add | Sub | Mult | Div | Modulo
67
     | AddAssign | SubAssign | MultAssign | DivAssign |
68
       ModuloAssign
       Equal | Neq | Less | Leq | Greater | Geq
69
       And | Or
70
71
   type prefix op =
     | Neg | Not | PreIncrement | PreDecrement
73
74
   type postfix_op =
75
     PostIncrement | PostDecrement
76
   type literal =
78
       BoolLit of bool
        IntLit of int64 * int
80
        FloatLit of float * int
       StringLit of string
82
83
   type expression =
84
        Literal of literal
85
        ObjectInitializer of expression list
86
        ArrayInitializer of expression list
87
        QualifiedId of qualified_id
       Member of expression * qualified id
       Call of expression * expression list
90
       Index of expression * expression list
91
       BinaryOp of expression * binary_op * expression
```

```
PrefixUnaryOp of prefix_op * expression
93
        Assignment of expression * expression
94
       Noop
95
96
   type parallel_expression =
97
        Invocations of expression
98
       ThreadCount of expression
99
100
   type variable definition =
      | VarBinding of binding * expression
102
   type general_statement =
104
       ExpressionStatement of expression
        VariableStatement of variable_definition
106
   type control_initializer = general_statement list *
108
       general_statement
   type statement =
110
        General\_statement
        Return of expression
112
       Break of int
113
        Continue
114
       ParallelBlock of parallel_expression list * statement
115
       list
       AtomicBlock of statement list
116
117
       IfBlock of control_initializer * statement list
      | IfElseBlock of control_initializer * statement list *
118
       statement list
       WhileBlock of control_initializer * statement list
119
       ForBlock of general_statement list * expression *
       expression list * statement list
      | ForByToBlock of expression * expression * expression *
121
       statement list
   type function_definition =
123
      qualified_id (* Name *)
124
      * binding list (* Parameters *)
      * type_name (* Return Type *)
126
      * statement list (* Body *)
127
128
   type basic definition =
129
        Variable Definition of variable definition
130
       FunctionDefinition of function_definition
131
132
```

```
type import_definition =
     | LibraryImport of qualified id
134
135
   type definition =
136
       Import of import_definition
137
       Basic of basic definition
138
       Namespace of qualified_id * definition list
139
140
   type program =
141
     | Program of definition list
142
143
   (* Useful destructuring and common operations *)
144
   let binding_type = function
     | (\underline{}, qt) -> qt
146
   let binding_name = function
148
     | (n, _) -> n
                            ../source/ast.ml
   (* LePiX Language Compiler Implementation
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```

```
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19
20
   (* Semantic checking for the Lepix compiler that will
21
      produce a new
   SemanticProgram type with things like locals group into a
       single type
   and type promotions / conversions organized for operators.
24
   module StringMap = Map. Make(String)
25
26
   type s_prefix_op = Ast.prefix_op
27
   type s_binary_op = Ast.binary_op
   type s_qualified_id = Ast.qualified_id
   type s_id = Ast.id
   type s_type_qualifier = Ast.type_qualifier
   type s builtin type = Ast. builtin type
33
34
   type s type name =
       SBuiltinType of s_builtin_type * s_type_qualifier
35
       SArray of s_type_name * int * s_type_qualifier
36
       SSizedArray of s_type_name * int * int list *
37
      s type qualifier
     | SFunction of s type name * s type name list *
       s_type_qualifier
     | SOverloads of s_type_name list
39
     SAlias of s_qualified_id * s_qualified_id
40
   let no_qualifiers = Ast.no_qualifiers
42
43
   let void t = SBuiltinType (Ast. Void, Ast. no qualifiers)
44
   let auto_t = SBuiltinType(Ast.Auto, Ast.no_qualifiers)
   let string_t = SBuiltinType(Ast.String, Ast.no_qualifiers)
   let bool_t = SBuiltinType(Ast.Bool, Ast.no_qualifiers)
   let \ int 32\_t \ = \ SBuiltin Type \left(\,Ast\,.\, Int \,(32)\,\,,\,\, Ast\,.\, no\,\_qualifiers\,\right)
   let int64 t = SBuiltinType(Ast.Int(64), Ast.no qualifiers)
   let float64_t = SBuiltinType(Ast.Float(64), Ast.
       no qualifiers)
51
   type s binding = s id * s type name
53
   type s literal =
54
     | SBoolLit of bool
```

```
SIntLit of int64 * int
56
       SFloatLit of float * int
57
       SStringLit of string
58
   type s_expression =
60
61
       SObjectInitializer of s_expression list * s_type_name
       SArrayInitializer of s_expression list * s_type_name
       SLiteral of s_literal
63
       SQualifiedId of s_qualified_id * s_type_name
64
65
       SMember of s expression * s qualified id * s type name
       SCall of s_expression * s_expression list * s_type_name
      SIndex of s_expression * s_expression list *
67
      s_type_name
     | SBinaryOp of s_expression * s_binary_op * s_expression
68
      * s_type_name
     | SPrefixUnaryOp of s_prefix_op * s_expression *
69
      s_type_name
     | SAssignment of s_expression * s_expression *
      s type name
     | SNoop
71
72
   type s_locals =
73
     | SLocals of s_binding list
74
75
   type s parameters =
76
     | SParameters of s binding list
77
78
   type s_variable_definition =
79
     SVarBinding of s_binding * s_expression
80
   type s_general_statement =
82
       SGeneralBlock of s_locals * s_general_statement list
       SExpressionStatement of s_expression
84
       SVariableStatement of s_variable_definition
86
   type s_capture =
87
     | SParallelCapture of s_binding list
88
89
   type s_control_initializer =
90
     SControlInitializer of s_general_statement *
91
      s_{expression}
   type s parallel expression =
93
       SInvocations of s_expression
94
       SThreadCount of s_expression
```

```
96
   type s statement =
       SBlock of s_locals * s_statement list
98
        SGeneral of s_general_statement
99
       SReturn of s_expression
100
        SBreak of int
101
       SContinue
102
103
      | SIfBlock of s control initializer (* Init statements
104
       for an if block *)
        * s_statement (* If code *)
105
106
      | SIfElseBlock of s_control_initializer (* Init
107
       statements for an if-else block *)
        * s_statement (* If code *)
108
        * s_statement (* Else code *)
109
      | SWhileBlock of s_control_initializer (* Init statements
111
        plus ending conditional for a while loop *)
        * s_statement (* code inside the while block, locals
       and statements *)
      | SForBlock of s_control_initializer (* Init statements
114
       plus ending conditional for a for loop *)
        * s expression list (* Post-loop expressions (increment
115
       /decrement) *)
        * s_statement (* Code inside *)
116
117
      | SParallelBlock of s_parallel_expression list (*
118
       Invocation parameters passed to kickoff function *)
       * s_capture (* Capture list: references to outside
119
       variables *)
        * s_statement (* Locals and their statements *)
      | SAtomicBlock of s_statement (* code in the atomic block
        *)
   type s_function_definition = {
125
     func_name : s_qualified_id;
126
      func_parameters : s_parameters;
127
         func_return_type : s_type_name;
128
     func body: s statement list;
129
130
   }
131
```

```
type s_basic_definition =
      | SVariableDefinition of s variable definition
133
      | SFunctionDefinition of s_function_definition
135
    type s_builtin_library =
136
137
      | Lib
138
    let builtin library names = [
139
      ("lib", Lib)
140
141
142
    type s_loop =
143
        SFor
      SWhile
145
    type s_module =
147
        SCode of string
148
        SDynamic of string
149
        SBuiltin of s builtin library
150
    type s definition =
152
      | SBasic of s_basic_definition
153
154
    type s_attributes = {
155
         attr parallelism : bool;
156
      attr arrays : int;
157
158
      attr_strings : bool;
159
160
    type s_environment = {
      env_usings : string list;
162
      env_symbols : s_type_name StringMap.t;
163
      env_definitions : s_type_name StringMap.t;
164
      env_imports : s_module list;
      env_loops : s_loop list;
166
167
168
    type s program =
169
      | SProgram of s_attributes * s_environment * s_definition
170
171
    (* Helping functions *)
    let rec coerce_type_name_of_s_expression injected =
      | SObjectInitializer(a, _) -> SObjectInitializer(a,
174
```

```
injected)
        SArrayInitializer(a, _) -> SArrayInitializer(a,
175
          injected)
          SQualifiedId(a, _) \rightarrow SQualifiedId(a, injected)
176
          SMember(a\,,\ b\,,\ \_) \ -\!\!\!> \ SMember(a\,,\ b\,,\ injected\,)
177
          SCall(a, b, \_) \rightarrow SCall(a, b, injected)
178
          SIndex\left(\begin{smallmatrix} a \\ \end{smallmatrix}, \begin{smallmatrix} b \\ \end{smallmatrix}, \begin{smallmatrix} - \end{smallmatrix}\right) \; -\!\!\!\!> \; SIndex\left(\begin{smallmatrix} a \\ \end{smallmatrix}, \begin{smallmatrix} b \\ \end{smallmatrix}, \; injected \right)
179
          SBinaryOp\left(\,a\,,\ b\,,\ c\,,\ \_\right) \ -\!\!\!> \ SBinaryOp\left(\,a\,,\ b\,,\ c\,,\ injected\,\right)
180
          SPrefixUnaryOp(a, b, _) -> SPrefixUnaryOp(a, b,
181
          injected)
        | SAssignment(a, b, _) -> SAssignment(a, b, injected)
182
        | e -> e
183
184
     let unqualify = function
185
          SBuiltinType(bt, _) -> SBuiltinType(bt, no_qualifiers)
          SArray(tn, d, \_) \rightarrow SArray(tn, d, no\_qualifiers)
187
          SSizedArray(tn, d, il, _) -> SSizedArray(tn, d, il,
188
          no_qualifiers)
        | SFunction(tn, pl, _) -> SFunction(tn, pl, no_qualifiers
189
190
        | t -> t
191
     let string_of_qualified_id qid =
192
       (String.concat"." qid)
193
194
     let parameter bindings = function
195
196
       | SParameters (bl ) -> bl
197
     let type_name_of_s_literal = function
198
          SBoolLit(_) -> bool_t
199
          SIntLit(_, b) -> SBuiltinType( Ast.Int(b),
200
          no_qualifiers )
        | SFloatLit(_,b ) -> SBuiltinType( Ast.Float(b),
201
          no_qualifiers )
        | SStringLit(_) -> string_t
202
203
     let rec type_name_of_s_expression = function
204
          SObjectInitializer(\_, t) \rightarrow t
205
          SArrayInitializer(\_, t) \rightarrow t
206
          SLiteral(lit) -> type_name_of_s_literal lit
207
          SQualifiedId(_, t) \rightarrow t
208
          SMember(\underline{\ }, \underline{\ }, t) \rightarrow t
209
          SCall(\_, \_, t) \rightarrow t
210
          SIndex(\_, \_, t) \rightarrow t
211
          SBinaryOp(\_, \_, \_, t) \rightarrow t
```

```
SPrefixUnaryOp(\_, \_, t) \rightarrow t
213
        SAssignment(\_, \_, t) \rightarrow t
214
       SNoop -> void t
215
216
   let return_type_name = function
217
        SFunction(rt,_, _) -> rt
218
      | t -> t
219
220
   let args type name = function
221
        SFunction (_, args, _) -> args
222
      | t -> []
223
224
   let mangled_name_of_type_qualifier = function
      (_, referencess) -> if referencess then "p!" else "!"
226
   let type_name_of_s_function_definition fdef =
228
      let bl = parameter_bindings fdef.func_parameters in
229
      let argst = List.map ( fun (_, t) -> t ) bl in
230
      let rt = fdef.func return type in
231
      SFunction (rt, argst, no_qualifiers)
233
   let mangled_name_of_builtin_type = function
234
        Ast. Void -> "v"
235
        Ast.Auto -> "a"
236
        Ast. Bool -> "b"
237
        Ast. Int(n) -> "i" ^ string of int n
238
        Ast. Float (n) -> "f" ^ string of int n
239
        Ast. String -> "s"
240
        Ast. Memory -> "m"
241
   let rec mangled_name_of_s_type_name = function
243
      | SBuiltinType (bt, tq ) ->
       mangled_name_of_type_qualifier tq ^
       mangled_name_of_builtin_type bt
      | SArray( tn, dims, tq ) ->
245
       mangled_name_of_type_qualifier_tq ^ "a" ^ string_of_int
       dims ^ ";" ^ mangled_name_of_s_type_name tn
      | SSizedArray( tn, dims, sizes, tq ) ->
246
       mangled_name_of_type_qualifier tq ^ "a" ^ string_of_int
            ^ ";" ^ mangled_name_of_s_type_name tn
      | SFunction( rt, pl, tq ) ->
247
       mangled_name_of_type_qualifier tq ^ "r;" ^ ( String.
       concat ";" ( List.map mangled name of s type name pl ) )
        ^ ";" ^ mangled_name_of_s_type_name rt
      _ -> "UNSUPPORTED"
248
```

```
let mangle name args qid tnl =
     string_of_qualified_id qid ^
251
     if (List.length tnl) > 0 then
252
       "_" ^ ( String.concat "_" ( List.map
253
       mangled_name_of_s_type_name_tnl ) )
     else
254
255
   let mangle name qid = function
256
       SFunction(rt, pl, tq) -> mangle name args qid pl
257
       _ -> string_of_qualified_id qid
258
                          ../source/semast.ml
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19
20
   (* Contains routines for string-ifying various parts of the
21
   infrastructure of the compiler, to make it easy to
```

249

```
understand what the fuck we're doing. *)
23
   (* Lexer types: dumping and pretty printing tokens *)
24
25
   module StringMap = Map. Make(String)
26
27
       preparser token to string = function
28
       Preparser.HASH -> "HASH"
29
       Preparser .IMPORT -> "IMPORT"
30
       Preparser .STRING -> "STRING"
31
       Preparser.TEXT(s) -> "TEXT(" ^ s ^ ")"
32
       Preparser.STRINGLITERAL(s) -> "STRINGLITERAL(" ^ s ^ ")
33
     | Preparser .EOF -> "EOF"
34
36
   let parser_token_to_string = function
37
       Parser.LPAREN -> "LPAREN"
38
       Parser.RPAREN -> "RPAREN"
       Parser.LBRACE -> "LBRACE"
40
41
       Parser .RBRACE -> "RBRACE"
       Parser .LSQUARE -> "LSQUARE"
       Parser.RSQUARE -> "RSQUARE"
43
       Parser.SEMI -> "SEMI"
44
       Parser.COMMA -> "COMMA"
45
       Parser.PLUSPLUS -> "PLUSPLUS"
46
47
       Parser . MINUSMINUS -> "MINUSMINUS"
       Parser.PLUS -> "PLUS"
48
       Parser.MINUS -> "MINUS"
49
       Parser.TIMES -> "TIMES"
       Parser.DIVIDE -> "DIVIDE"
       Parser .MODULO -> "MODULO"
       Parser . PLUSASSIGN -> "PLUSASSIGN"
       Parser.MINUSASSIGN \rightarrow "MINUSASSIGN"
       Parser.TIMESASSIGN -> "TIMESASSIGN"
       Parser . DIVIDEASSIGN -> "DIVIDEASSIGN"
       Parser.MODULOASSIGN -> "MODULOASSIGN"
       Parser.ASSIGN -> "ASSIGN"
       Parser.EQ -> "EQ"
59
       Parser.NEQ -> "NEQ"
60
       Parser.LT -> "LT"
61
       Parser.LEQ -> "LEQ"
       Parser.GT -> "GT"
63
       Parser .GEQ -> "GEQ"
64
       Parser .AND -> "AND"
```

```
Parser.OR -> "OR"
66
        Parser.NOT -> "NOT"
        Parser.DOT -> "DOT"
68
        Parser.AMP -> "AMPERSAND"
69
        Parser .COLON -> "COLON"
70
71
        Parser.PARALLEL -> "PARALLEL"
        Parser .INVOCATIONS -> "INVOCATIONS"
        Parser.THREADCOUNT -> "THREADCOUNT"
73
        Parser.ATOMIC -> "ATOMIC"
74
75
        Parser.VAR -> "VAR"
        Parser.LET -> "LET"
76
        Parser.CONST -> "CONST"
        Parser.FUN -> "FUN"
        Parser.NAMESPACE -> "NAMESPACE"
79
        Parser.IF -> "IF"
        Parser.ELSE -> "ELSE"
81
        Parser.FOR -> "FOR"
        Parser.TO -> "TO"
83
        Parser.BY -> "BY"
        Parser.WHILE -> "WHILE"
85
        Parser .RETURN -> "RETURN"
        Parser.INT(b) -> "INT" ^ string_of_int b
        Parser.FLOAT(b) -> "FLOAT" ^ string_of_int b
88
        Parser.BOOL -> "BOOL"
89
        Parser.STRING -> "STRING"
90
        Parser.VOID -> "VOID"
91
        Parser.AUTO -> "AUTO"
92
        Parser.MEMORY -> "MEMORY"
93
        Parser.TRUE -> "TRUE"
94
        Parser.FALSE -> "FALSE"
        Parser.BREAK -> "BREAK"
96
        Parser.CONTINUE -> "CONTINUE"
        Parser .IMPORT -> "IMPORT"
98
        Parser.STRINGLITERAL(s) \rightarrow "STRINGLITERAL(" ^ s ^ ")"
        Parser.INTLITERAL(i) -> "INTLITERAL(" ^ Num.
100
       string_of_num i ^ ")"
      | \  \, Parser.FLOATLITERAL(\,f\,) \  \, -\!\!\!> \  \, "FLOATLITERAL(\,"\  \, \hat{\  \, } \  \, Num.
       string_of_num f ^ ")"
        Parser.ID(s) -> "ID(" ^ s ^ ")"
        Parser.EOF -> "EOF"
103
104
    let token_range_to_string (x, y) =
105
      let range_is_wide = (y - x > 1) in
106
      if range_is_wide then
107
        ( string_of_int x ^ "-" ^ string_of_int y,
108
```

```
range_is_wide )
      else
109
        ( string_of_int x, range_is_wide )
111
    let token_source_to_string t =
112
113
      let (s, _) = token_range_to_string t.Base.
       token_column_range in
      string of int t. Base. token line number
114
      ^ ":"
115
116
117
    let preparser_token_list_to_string token_list =
118
      let rec helper = function
119
      | (token, pos) :: tail \rightarrow
120
         [" ^ ( preparser_token_to_string token ) ^ ":"
          token_source_to_string pos ^ "] "
         helper tail
123
      [] -> "\n" in helper token_list
124
125
    let parser_token_list_to_string token_list =
126
127
      let rec helper = function
      | (token, pos) :: tail ->
128
        "[" ^ ( parser_token_to_string token ) ^ ":"
129
          token_source_to_string pos ^ "] "
130
        ^ helper tail
131
      | [] -> "\n" in helper token_list
132
133
134
    (* Program types:
135
    dumping and Pretty-Printing *)
136
137
138
    let string\_of\_id i = i
139
    let string_of_qualified_id qid = ( String.concat "." ( List
140
        .map string_of_id qid ) )
141
    let string_of_binary_op = function
142
        Ast.Add \rightarrow "+"
143
        Ast.Sub -> "-"
144
        Ast.Mult -> "*"
145
        Ast. Div -> "/"
146
        Ast.\,Modulo \,\to\, "\%"
147
        Ast. AddAssign -> "+="
148
        Ast. SubAssign -> "-="
149
        Ast. MultAssign -> "*="
```

```
Ast. DivAssign -> "/="
151
        Ast. Modulo Assign -> "%="
        Ast. Equal -> "=="
        Ast. Neq -> "!="
154
        Ast. Less \rightarrow "<"
        Ast.Leq -> "<="
156
        Ast. Greater -> ">"
        Ast.Geq -> ">="
158
        Ast . And -> "&&"
159
        Ast. Or -> " | | "
160
161
    let string_of_unary_op = function
162
        Ast. Neg -> "-"
        Ast. Not -> "!"
164
        Ast. PreDecrement -> "--"
        Ast. PreIncrement -> "++"
166
167
    let rec string_of_expression = function
168
        Ast.Literal(Ast.IntLit(1, _)) -> Int64.to_string l
169
        Ast. Literal (Ast. BoolLit (true)) -> "true"
        Ast. Literal (Ast. BoolLit (false)) -> "false"
171
        Ast.\,Literal\,(\,Ast.\,StringLit\,(\,s\,)\,) \  \, -\!\!\!> \,\,"\,\setminus"\,"\,\,\,\hat{}\,\,\,s\,\,\,\hat{}\,\,\,"\,\setminus"\,"\,\,
        Ast. Literal (Ast. FloatLit (f, _)) -> string_of_float f
173
        Ast. QualifiedId (qid) -> string_of_qualified_id qid
174
        Ast. BinaryOp (e1, o, e2) \rightarrow
175
        string of expression e1 ^ " " ^ string of binary op o ^
176
                string_of_expression e2
      | Ast.PrefixUnaryOp(o, e) -\!\!\!> string\_of\_unary\_op o ^
177
        string_of_expression e
      | Ast.Index(e, 1) -> string_of_expression e ^ "[" ^ (
        String.concat ", " (List.map string_of_expression 1)) ^
      | Ast.Member(e, qid) -> string_of_expression e ^ "." ^
179
        string_of_qualified_id qid
      | Ast.Assignment(e1, e2) -> string_of_expression e1 ^ " =
180
         " ^ string_of_expression e2
      | Ast. Call(e, el) ->
181
        string_of_expression e ^ "(" ^ String.concat ", " (List
182
        .map string_of_expression el) ^ ")"
        Ast. Noop -> " { noop }"
183
      | Ast. ArrayInitializer(el) -> "[ " ^ String.concat ", " (
184
        List.map string_of_expression el) ^ " ]"
      | Ast. ObjectInitializer(el) -> "{ " String.concat ", "
185
        (List.map string_of_expression el) ^ " }"
186
```

```
let string_of_parallel_expression = function
      Ast. ThreadCount(e) -> "thread count = " ^
       string_of_expression e
      | Ast. Invocations (e) -> "invocations = " ^
189
       string_of_expression e
190
   let rec string_of_expression_list el =
191
      String.concat ", " ( List.map string_of_expression el )
192
193
   let rec string of builtin type = function
194
       Ast. Auto -> "auto"
195
       Ast. Bool -> "bool"
196
       Ast. Int(b) -> "int" ^ string_of_int b
197
       Ast. Float(b) -> "float" ^ string_of_int b
198
       Ast. String -> "string"
        Ast. Memory -> "memory"
200
       Ast. Void -> "void"
202
   let string_of_type_qualifier = function
203
     \mid (c, r) -\!\!> ( if c then "const" else "" ) ^ ( if r then "
204
       &" else "")
205
   let rec string_of_type_name tn =
206
      let tqual tq =
207
        let s = string_of_type_qualifier tq in
208
        if s = "" then "" else s " "
209
210
      in match to with
      | Ast.BuiltinType(t, tq) -> tqual tq ^
211
       string_of_builtin_type t
      | Ast.Array(t, d, tq) -> tqual tq ^ string_of_type_name t
212
        (String.make d'['') (String.make d']'')
213
      Ast. SizedArray(t, d, il, tq) -> tqual tq ^
       string_of_type_name t ^ ( String.make d '[' ) ^ ( String
       .concat ", " ( List.map string_of_int il ) ) ^ ( String.
       make d '] ' )
      | Ast.Function(r, args, tq) -> tqual tq ^ "(" ^ ( String
214
       .concat ", " ( List.map ( fun v -> string_of_type_name v
        ) args ) ) ^ ")" ^ string_of_type_name r
215
   let string_of_binding = function
216
     | (n, t) -> n ^ " : " ^ string_of_type_name t
217
218
   let string_of_variable_definition = function
219
     | Ast. VarBinding (b, Ast. Noop) -> "var" ^
       string_of_binding b
```

```
| Ast. VarBinding (b, e) -> "var" ^ string_of_binding b ^
221
        = " ^ string_of_expression e
222
   let string_of_general_statement = function
223
     | Ast. ExpressionStatement(e) -> string_of_expression e
224
      | Ast. VariableStatement(v) ->
225
       string_of_variable_definition v
226
   let \ string\_of\_condition\_initializer = function
227
     (il, cond) -> (String.concat"; "(List.map
228
       string_of_general_statement il ) )
       ( if ( List length il ) > 0 then ";" else "" )
229
        ( string_of_general_statement cond )
230
231
   let rec string_of_statement s =
      let string_of_statement_list sl = String.concat "" (List.
233
       map string_of_statement sl) in
     match s with
        | Ast. General(b) -> string of general statement b ^ ";\
235
       n";
       | Ast. Return (expr) -> "return" ^ string of expression
236
       \exp r \hat{"}; \ n";
       | Ast. IfBlock (ilcond, s) -> "if (" ^
237
       string_of_condition_initializer ilcond ^ ")"
          "{\n" 	 string_of_statement_list s 	 "}\n"
238
        | Ast.IfElseBlock(ilcond, s, s2) -> "if ("
239
       string_of_condition_initializer ilcond ^ ")"
          "{\n" \cdot "}
240
          ^{\circ} "else {\n" ^{\circ} string_of_statement_list s2 ^{\circ} "}\n"
241
        | Ast.ForBlock(gsl, cond, incrl, sl) -> "for ("
       String.concat ", " (List.map string_of_general_statement
        gsl) ) ^ "; "
           string_of_expression cond ^ "; "
243
           string_of_expression_list incrl ^ ") {\n" ^
244
       string_of_statement_list_sl ^ "}\n"
       | Ast.ForByToBlock(e1, e2, e3, s1) -> "for (" ^
245
       string_of_expression el ^ " to " ^ string_of_expression
        e2 ^ " by " ^ string_of_expression e3 ^ ") {\n"
       string\_of\_statement\_list sl ^ "}\n"
       | Ast. WhileBlock (ilcond, s) -> "while ("
246
       string_of_condition_initializer ilcond ^ ") {\n"
          \hat{string}_{of\_statement\_list} \hat{s}^{n} 
247
        Ast. Break(n) \rightarrow ( if n == 1 then "break" else "break"
248
          string_of_int n ) ^ ";\n"
        Ast. Continue -> "continue;\n"
249
```

```
| Ast. ParallelBlock (pel, sl) -> "parallel (" ^ (String.
250
        concat ", " (List.map string_of_parallel_expression pel)
       ) ^ " ) {"
          ^ "\n" ^ string_of_statement_list_sl ^ "}\n"
251
        | Ast.AtomicBlock(sl) \rightarrow "atomic {\n"}
252
        string\_of\_statement\_list sl ^ "}\n'
253
    let string of statement list sl =
254
      String.concat "" (List.map string of statement sl)
255
256
    {\tt let} \ string\_of\_function\_definition = function
257
      | ( name, parameters, return_type, body) ->
"fun " ^ string_of_qualified id name
258
259
          ^ "(" ^ (String.concat ", " (List.map
260
        string_of_binding parameters)) ^ ") : "
          string_of_type_name return_type
261
          ^ string_of_statement_list body
262
          ^ "}\n"
263
264
    let rec string_of_basic_definition = function
265
266
      | Ast. Function Definition (fdef) ->
       string_of_function_definition fdef
      | Ast. Variable Definition (vdef) ->
267
       string_of_variable_definition vdef ^ ";\n"
268
    let string of import definition = function
269
270
      | Ast. LibraryImport (qid) -> "import" ^
       string_of_qualified_id qid ^ "\n"
271
    let rec string_of_definition = function
272
        Ast.Import(idef) -> string_of_import_definition idef
273
        Ast. Basic (bdef) -> string_of_basic_definition bdef
        Ast. Namespace (qid, defs) -> "namespace" ^
275
        string_of_qualified_id qid ^ " {\n"
        ^ (String.concat "" (List.map string_of_definition defs
       let string of program = function
278
      | Ast.Program(p) \rightarrow let s = (String.concat "" (List.map)
279
       string_of_definition p) ) in
      Base.brace_tabulate s 0
280
281
    (* Semantic Program types:
282
    dumping and pretty printing *)
283
```

```
let rec string_of_s_type_name tn =
285
      let tqual tq =
        let s = string_of_type_qualifier tq in
287
        if s = "" then "" else s ^{\circ} ""
288
      in match to with
289
      | Semast.SBuiltinType(t, tq) -> tqual tq ^
290
       string_of_builtin_type t
      | Semast. SArray(t, d, tq) \rightarrow tqual tq \hat{}
291
       string_of_s_type_name t ^ ( String.make d '[' ) ^ (
       String.make d '|' )
      | Semast.SSizedArray(t, d, il, tq) -> tqual tq ^
292
       string_of_s_type_name t ^ ( String.make d '[' ) ^ (
       String.concat ", " ( List.map string_of_int il ) )
       String.make d '|' )
      | Semast. SFunction (r, args, tq) -> tqual tq ^ "(" ^ (
293
       String.concat ", " ( List.map ( fun v ->
       string_of_s_type_name v ) args ) ) ^ ")" ^
       string_of_s_type_name r
      | Semast. SOverloads (fl) -> "overloads ["
294
          (String.concat", "(List.map string_of_s_type_name
295
        fl ) )
        ~ "j"
296
      | Semast.SAlias(target, source) -> "using" \hat{\ }
297
       string_of_qualified_id target ^ " -> "
       string of qualified id source
298
    299
300
301
    let string_of_s_locals = function
      | Semast.SLocals(bl) -> if (List.length bl < 1) then ""
303
        else (String.concat";\n" (List.map
       string_of_s_binding bl ) ) ^ ";"
    let string_of_s_literal = function
305
        Semast.SBoolLit(b) -> string_of_bool b
306
        Semast.\,SIntLit\,(i\;,\;\_) \;-\!\!>\; Int64.\,to\_string\;\;i\\Semast.\,SFloatLit\,(f\;,\;\_) \;-\!\!>\; string\_of\_float\;\;f
307
308
        Semast. SStringLit(s) -> "\"" ^ s ^ "\""
309
310
    let rec string_of_s_expression = function
311
      | Semast.SObjectInitializer(el, tn) ->
312
        string_of_s_type_name tn ^ "{
313
        ^ String.concat ", " ( List.map string_of_s_expression
314
       el)
```

```
^ "}"
315
      | Semast.SArrayInitializer(el, tn) ->
316
        string_of_s_type_name tn ^ "[ "
317
        ^ String.concat ", " ( List.map string_of_s_expression
318
        el)
319
        Semast. SLiteral(1) -> string_of_s_literal 1
320
        Semast. SQualifiedId (qid, tn) ->
321
       string_of_s_type_name tn ^ " ]] "
       string of qualified id qid
      | Semast.SMember(e, qid, tn) -> "[[ " ^
string_of_s_type_name tn ^ " ]] " ^
string_of_s_expression e ^ "." ^ string_of_qualified_id
322
      | Semast.SCall(e, el, tn) -> "[[" ^ string_of_s_type_name
323
        tn ^ "]] " ^ string_of_s_expression e ^ "( " ^ ( String.
        concat ", " ( List.map string_of_s_expression el ) ) ^
        ) "
      | Semast.SIndex(e, el, tn) \rightarrow "[[ "
324
       string_of_s_type_name tn ^ " ]] " ^ string_of_s_expression e ^ "[ " ^ ( String.concat ", " (
      List.map string_of_s_expression el ) ) ^ " ]" | Semast.SBinaryOp(1, op, r, tn) -> "[[ " ^
325
       string_of_s_type_name tn ^ " ]] " ^ string_of_s_expression l ^ " " ^ string_of_binary_op op
        ^ " " ^ string_of_s_expression r
      326
       op ^ string_of_s_expression r
      327
       string_of_s_expression l ^ " = " ^
       string_of_s_expression r
      | Semast.SNoop -> "(noop)"
328
329
    let string_of_s_capture = function
330
      | Semast.SParallelCapture(bl) -> let capturecount = List.
331
       length bl in
        if capturecount = 0 then "[[no captures]]\n" else
332
        "[[captures]] { " ^ ( String.concat ", " ( List.map
333
        string_of_s_binding bl ) )
        ^ " }\n"
334
335
    let string_of_s_variable_definition = function
336
      | Semast.SVarBinding(b, e) -> "var" ^
```

```
string_of_s_binding b ^ " = " ^ string_of_s_expression e
338
    let rec string_of_s_general_statement = function
339
      | Semast.SGeneralBlock(locals, gsl) -> "{\n" ^
340
       string_of_s_locals locals ^ "\n" ^ ( String.concat ";\n"
        (List.map string_of_s_general_statement gsl) ) ^ "\n}"
       Semast. SExpressionStatement (sexpr) ->
341
       string of s expression sexpr
      | Semast.SVariableStatement(v) ->
342
       string of s variable definition v
343
    let string_of_s_general_statement_list gsl =
344
      String.concat ";\n" (List.map
       string_of_s_general_statement gsl)
346
    let string_of_s_parallel_expression = function
347
        Semast. SInvocations (e) -> string_of_s_expression e
348
        Semast.SThreadCount(e) -> string_of_s_expression e
349
    let rec string_of_s_statement s =
351
352
      let initializer_begin = function
          Semast.SGeneralBlock(locals, gsl) ->
353
          let precount = List.length gsl in
354
          if precount > 1 then
355
356
            ^string_of_s_locals locals
                      string_of_s_general_statement_list_gsl
358
            ^{n}
359
          else
360
          _ -> ""
362
      in
      let initializer_end = function
364
          Semast.SGeneralBlock(locals, gsl) ->
          let precount = List.length gsl in
366
          if precount > 1 then
367
            " }\n"
368
          else
369
370
        | _ -> ""
371
      in match s with
372
      | Semast.SBlock(locals, sl) -> "{\n" ^ string of s locals
373
        locals ^ "\n" ^ ( String.concat "\n" (List.map
       string\_of\_s\_statement sl) ) ^ "\n}\n"
      | Semast.SGeneral(g) -> ( string_of_s_general_statement
374
```

```
g ) ^ ";"
      | Semast.SReturn(sexpr) -> "return" ^
375
       string_of_s_expression sexpr ^ ";"
      | Semast.SBreak(n) \rightarrow if n < 2 then "break;" else "break
376
       " ^ string_of_int n ^ ";"
      | Semast.SContinue -> "continue;"
377
        Semast.SAtomicBlock(s) -> "atomic {"
          string of s statement s
379
        ^ "}\n"
380
      | Semast.SParallelBlock(pel, captures, s) ->
381
        "parallel(" ^ (String.concat ", " (List.map
       string_of_s_parallel_expression_pel)) ^ " ) {"
        ^{\circ} "\n" ^{\circ} string_of_s_capture captures
        ^ "\n" ^ string_of_s_statement s
384
        ^ "}\n"
      | Semast.SIfBlock(Semast.SControlInitializer(inits, cond)
386
        , s) \rightarrow
        initializer_begin inits
387
        ^ "if (" ^ string of s expression cond ^ ") {"
388
        ^ "\n" ^ string_of_s_statement s
389
        `"}\n"
390
        ^ initializer_end inits
391
       Semast. SIfElseBlock (Semast. SControlInitializer (inits,
392
       cond), is, es) \rightarrow
        initializer begin inits
393
        ^ "if (" ^ string of s expression cond ^ ") {"
        ^ "\n" ^ string_of_s_statement is
395
        `"}\n"
        ^ "else {"
397
        ^ "\n" ^ string_of_s_statement es
        ^ "}\n"
399
        ^ initializer_end inits
       Semast. SWhileBlock (Semast. SControlInitializer (inits,
401
       cond), s) \rightarrow
        initializer_begin inits
402
          "while (" ^ string_of_s_expression cond ^ ") {"
403
        ^ "\n" ^ string_of_s_statement s
404
        ^ "}\n"
405
          initializer_end inits
406
       Semast.SForBlock(Semast.SControlInitializer(inits, cond
407
       ), increxprl, s) \rightarrow
        let incrl = String.concat ", " ( List.map
408
       string_of_s_expression increxprl ) in
        initializer_begin inits
409
        ^ "for (;" ^ string_of_s_expression cond ^ "; " ^ incrl
410
```

```
^ "\n" ^
                 string_of_s_statement s
411
        `"}\n"
412
        ^ initializer_end inits
413
414
    let string_of_s_statement_list sl =
415
      ( String.concat "\n" ( List.map\ string\_of\_s\_statement\ sl
416
       ) ) ^ "\n"
417
    let string of s block = function
418
       (locals, sl) -> "{\n" ^ string_of_s_locals locals
419
         "\n" \hat{} string_of_s_statement_list_sl
420
        n \setminus n \setminus n
421
422
    let string_of_s_parameters = function
      | Semast.SParameters(parameters) -> String.concat ", " (
424
        List.map string_of_s_binding parameters)
425
    let string\_of\_s\_function\_definition f =
426
      "fun " ^ string_of_qualified_id f.Semast.func_name
427
        "(" \hat{} string_of_s_parameters f.Semast.func_parameters \hat{}
428
        ") : "
        string_of_s_type_name f.Semast.func_return_type
429
      ^ string_of_s_statement_list f.Semast.func_body
430
      ^ "}\n"
431
432
    let string_of_s_basic_definition = function
433
      | Semast. SVariable Definition (v) ->
434
       string_of_s_variable_definition v
      | Semast. SFunctionDefinition(f) ->
435
       string_of_s_function_definition f
436
    let string_of_s_builtin_library = function
      | Semast.Lib -> "lib"
438
439
    let string_of_s_module = function
440
        Semast.SCode(s) \rightarrow "import \ [[code]] \ " \ `s
441
        Semast.SDynamic(s) -> "import [[dynamic]]
442
        Semast.SBuiltin(bltin) -> "import [[builtin]] "
443
       string_of_s_builtin_library bltin
    let rec string_of_s_definition = function
445
      | Semast.SBasic(b) -> string_of_s_basic_definition b ^ "\
446
       n"
```

```
447
    let string of s program = function
448
        Semast.SProgram(attr, env, sdl) ->
449
      let symbolacc k tn l =
450
        let entry = ( string_of_s_type_name tn ) ^ " | " ^ k in
451
452
        entry :: 1
      and importance m =
453
        string of s module m
454
455
      let implist = List.map importace env. Semast.env imports
456
      and symbollist = StringMap.fold symbolacc env. Semast.
457
       env_symbols []
      and defsymbollist = StringMap.fold symbolacc env. Semast.
458
       env definitions []
459
      in
      let i = "imports: \n\t" ^ ( String.concat "\n\t" implist )
460
      and s = "symbols: \n\t" ^ ( String.concat "\n\t"
        symbollist )
      and d = "code symbols: \n\t" ^ ( String.concat "\n\t")
462
        defsymbollist )
      and a = "strings: " ^ string_of_bool attr.Semast.
463
        attr_strings
         ^ "\narrays: " ^ string_of_int attr.Semast.attr_arrays
464
        ^ "\nparallelism: " ^ string_of_bool attr.Semast.
465
        attr parallelism
466
      let p = String.concat "" (List.map string of s definition
467
         sdl) in
        Base.brace_tabulate ( a ^ "\n\n" ^ i ^ "\n\n" ^ s ^ "\n
468
       n ^{\circ} ^{\circ}
                          ../source/representation.ml
    (* LePiX - LePiX Language Compiler Implementation
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      CONNECTION WITH THE
  SOFTWARE OR THE USE OR OTHER DEALINGS IN THE SOFTWARE. *)
19
20
   (* Ocamllex Scanner for LePiX Preprocessor *)
21
   let whitespace = [' ' ' ' t' ' r']
   let newline = [' \ n']
24
   rule token = parse
26
   newline as c { Lexing.new_line lexbuf; let b = Buffer.
      create 1024 in Buffer.add char b c; sub token b lexbuf }
                  { let t = Preparser.HASH in t :: hash token
      lexbuf
   | _ as c
                  { let b = Buffer.create 1024 in Buffer.
      add_char b c; sub_token b lexbuf }
                  { [Preparser.EOF] }
   eof
30
31
  and sub_token text_buf = parse
   newline as c { Lexing.new_line lexbuf; Buffer.add_char
      text_buf c; sub_token text_buf lexbuf }
                   { let s = Buffer.contents text_buf in let t
       = Preparser.TEXT(s) in let l = Preparser.HASH :: (
      hash_token lexbuf ) in t :: 1 }
                   { Buffer.add_char text_buf c; sub_token
      text_buf lexbuf }
                   { let s = Buffer.contents text buf in let t
36
       = Preparser.TEXT(s) in [t; Preparser.EOF] }
  and hash token = parse
               { Lexing.new_line lexbuf; token lexbuf }
   | whitespace { hash_token lexbuf }
```

```
\{ \text{ let b} = ( \text{ Buffer.create } 128 ) \text{ in } \}
      string literal b lexbuf }
     "import"
                { Preparser.IMPORT :: hash token lexbuf }
42
                { Preparser.STRING :: hash_token lexbuf }
     "string"
     eof
                 { [Preparser.EOF] }
44
                { raise (Errors. UnknownCharacter(String.make 1
45
       c, (Lexing.lexeme_start_p lexbuf, Lexing.lexeme_end_p
      lexbuf ) )) }
46
47
  and string literal string buf = parse
   | newline as c { Lexing.new_line lexbuf; Buffer.add char
      string_buf c; string_literal string_buf lexbuf }
                  { let sl = Preparser.STRINGLITERAL( Buffer.
49
      contents string_buf ) in [sl] }
   "\\"" as s { Buffer.add_string_buf s;
      string_literal string_buf lexbuf }
                 { Buffer.add_char string_buf c;
      string_literal string_buf lexbuf }
                        ../source/prescanner.mll
  %{
  (* LePiX - LePiX Language Compiler Implementation
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20
21
   (* Parser for the LePiX preprocessor: compatible with both
      ocamlyacc and
  menhir, as we have developed against both for testing
      purposes. *)
24
  %}
25
26
  %token HASH
27
  %token IMPORT STRING
  %token <string> TEXT
  %token <string> STRINGLITERAL
  %token EOF
31
  %start source
  %type<Preast.pre source> source
35
36
  blob:
37
   | HASH IMPORT STRINGLITERAL { Preast.ImportSource($3) }
   | HASH IMPORT STRING STRINGLITERAL { Preast.ImportString($4
   | TEXT { Preast. Text($1) }
40
41
   blob_list: { [] }
   | blob_list blob { $2 :: $1 }
43
44
   source:
45
   | blob_list EOF { List.rev $1 }
                        ../source/preparser.mly
   (* LePiX – LePiX Language Compiler Implementation
2
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19
20
   (* Ocamllex Scanner for LePiX Preprocessor *)
21
22
23
     open Parser
24
25
26
   let whitespace = [' ' ' ' t' ' r']
   let newline = [' \ n']
   let binary_digit = '0' | '1'
   let \ hex\_digit = ['0'-'9'] \ | \ ['A'-'F'] \ | \ ['a'-'f']
   let octal_digit = ['0'-'7']
   let \operatorname{decimal\_digit} = ['0' - '9']
   let uppercase_letter = ['A'-'Z']
   let lowercase_letter = ['a'-'z']
   let n_digit = decimal_digit | uppercase_letter |
      lowercase letter
36
   rule token = parse
     whitespace { token lexbuf }
              { Lexing.new_line lexbuf; token lexbuf }
     newline
     "/*"
                multi comment 0 lexbuf }
                single comment lexbuf }
     1(1
                LPAREN }
   | ')'
                RPAREN }
```

```
{ LBRACE }
44
      1}'
                { RBRACE }
              { LSQUARE }
46
      ٠١٠
               RSQUARE }
47
                { SEMI }
48
      1:1
49
              { COLON }
                  COMMA }
50
      1+1
                  PLUS }
51
                  MINUS }
      ' * '
53
                  TIMES }
      1/1
                  DIVIDE }
54
     "+="
55
                  PLUSASSIGN }
                  MINUSASSIGN }
56
                  TIMESASSIGN }
57
     "/="
                  DIVIDEASSIGN }
58
      "%="
                  MODULOASSIGN }
59
     "++"
60
                  PLUSPLUS }
                  MINUSMINUS }
61
     "%"
                  MODULO }
62
      '='
                  ASSIGN }
63
      1&1
                  AMP
64
     "___"
                  EQ }
65
     "!="
                  NEQ }
66
      ' < '
                  LT }
67
     "<="
                  LEQ }
68
     ">"
                  GT }
69
     ">="
70
                  GEQ }
      "&&"
                  AND }
71
      1.1
                  DOT }
72
     1&1
                  AMP }
     " | "
                  OR 
74
     "!"
                  NOT }
75
     " i f "
76
                  IF }
     " else"
                  ELSE }
77
     "for"
                  FOR }
78
     "while"
                  WHILE }
79
     "by"
                  BY }
80
     " to " \,
                  TO }
81
     "return"
                  RETURN }
82
                \{ AUTO \}
83
     "int" ((decimal_digit+)? as s) { let bits = if s = ""
       then \ Base.default\_integral\_bit\_width \ else \ (
       int_of_string s ) in INT(bits) }
   "float" ((decimal_digit+)? as s) { let bits = if s = ""
       then Base.default_floating_bit_width else (
```

```
int_of_string s ) in FLOAT(bits) }
      "bool"
                  BOOL }
      "string"
                  STRING }
87
      "void"
88
                  VOID }
      "memory"
                  MEMORY }
89
      "true"
                  TRUE }
90
      "false"
                  FALSE }
91
      "var"
                  VAR }
92
                  LET }
      "let"
93
      "const"
                \{ CONST \}
94
                 { FUN }
      "fun"
95
      "parallel" { PARALLEL }
96
      "break" { BREAK }
97
      "continue" { CONTINUE }
98
      "invocations" { INVOCATIONS }
      "thread_count" { THREADCOUNT }
100
      "atomic" { ATOMIC }
      "namespace" { NAMESPACE }
      "import" { IMPORT }
                { string_literal ( Buffer.create 128 ) lexbuf }
105
      decimal_digit+ as lxm { INTLITERAL(Num.num_of_string lxm)
      "0c"
              "OC" { octal_int_literal lexbuf }
106
      "0x"
              " 0X"
                     hex_int_literal lexbuf }
107
              "0B" { binary_int_literal lexbuf }
108
      ("0n" | "0N") (decimal digit + as b) ("n" | "N") {
109
       n_int_literal (int_of_string b) lexbuf }
      '.' [0'-9']+(e'(+')'-1)?[0'-9']+? as s {
110
       FLOATLITERAL(Polyfill.num_of_string s) }
     [0'-9']+(0'-9']*(e'(+')-9')?[0'-9']+?
          ('e''('+'|'-')? ['0'-'9']+)?) as s { try FLOATLITERAL(
        Polyfill.num_of_string_base 10 s) with _ -> raise (
        Errors.BadNumericLiteral(s, (Lexing.lexeme_start_p)
       lexbuf , Lexing.lexeme_end_p lexbuf ) )) }
     \left[ \left[ \left[ a' - \left[ z' \right] \right] \right] \left[ \left[ a' - \left[ z' \right] \right] \right] \left[ \left[ a' - \left[ z' \right] \right] \right] + \text{as s } \left[ \text{ID} \right] 
        (s)
      eof { EOF }
      _ as c { raise (Errors.UnknownCharacter(String.make 1 c,
114
        ( Lexing.lexeme_start_p lexbuf, Lexing.lexeme_end_p
       lexbuf ) )) }
115
116
    and octal int literal = parse
     octal_digit+ as s { try INTLITERAL( Polyfill.
        num_of_string_base 8 s ) with _ -> raise (Errors.
```

```
BadNumericLiteral(s, (Lexing.lexeme_start_p lexbuf,
       Lexing.lexeme end p lexbuf ) )) }
   _ as c { raise (Errors.BadNumericLiteral(String.make 1 c,
        ( Lexing.lexeme_start_p lexbuf, Lexing.lexeme_end_p
       lexbuf ) )) }
120
   and hex int literal = parse
122
   hex digit+ as s { try INTLITERAL( Polyfill.
       num_of_string_base 16 s ) with _ -> raise (Errors.
       BadNumericLiteral(s, (Lexing.lexeme_start_p lexbuf,
       Lexing.lexeme_end_p lexbuf ) )) }
   _ as c { raise (Errors.BadNumericLiteral(String.make 1 c,
        ( Lexing.lexeme_start_p lexbuf, Lexing.lexeme_end_p
       lexbuf ) )) }
126
   and binary_int_literal = parse
127
   | binary_digit+ as s { try INTLITERAL( Polyfill.
       num_of_string_base 2 s ) with _ -> raise (Errors.
       BadNumericLiteral(s, (Lexing.lexeme_start_p lexbuf,
       Lexing.lexeme_end_p lexbuf ) )) }
   _ as c { raise (Errors.BadNumericLiteral(String.make 1 c,
       (Lexing.lexeme_start_p lexbuf, Lexing.lexeme_end_p
       lexbuf ) )) }
130
131
   and n_int_literal base = parse
132
                     { INTLITERAL( Polyfill.
   n_digit+ as s
      num_of_string_base base s ) }
     _ as c { raise (Errors.BadNumericLiteral(String.make 1 c,
134
       ( Lexing.lexeme_start_p lexbuf, Lexing.lexeme_end_p
       lexbuf ) )) }
136
   and string_literal string_buffer = parse
137
     newline as c { Lexing.new_line lexbuf; Buffer.add_char
       string_buffer c; string_literal string_buffer lexbuf }
     '"' { let v = STRINGLITERAL( Buffer.contents
139
       string buffer ) in v }
   | "\\\"" as s { Buffer.add_string_buffer s;
       string_literal string_buffer lexbuf }
   _ as c { Buffer.add_char string_buffer c; string_literal
141
       string_buffer lexbuf }
142
```

```
143
   and multi comment level = parse
   newline { Lexing.new_line lexbuf; multi_comment level
       lexbuf
   "*/" { if level = 0 then token lexbuf else multi_comment
146
       (level-1) lexbuf \}
     "/*" { multi_comment (level+1) lexbuf }
          { multi comment level lexbuf }
148
149
150
   and single_comment = parse
   newline { Lexing.new_line lexbuf; token lexbuf }
152
          { single_comment lexbuf }
                         ../source/scanner.mll
   %{
   (* LePiX - LePiX Language Compiler Implementation
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```

```
21
   (* Parser for the LePiX language: compatible with both
      ocamlyacc and
  menhir, as we have developed against both for testing
23
      purposes. *)
24
  %}
25
26
  %token SEMI LPAREN RPAREN LBRACE RBRACE COMMA
  %token LSQUARE RSQUARE COLON
  %token DOT
  %token PARALLEL INVOCATIONS ATOMIC THREADCOUNT
  %token PLUSPLUS MINUSMINUS
  %token PLUS MINUS TIMES DIVIDE ASSIGN
  %token MODULO
  %token PLUSASSIGN MINUSASSIGN TIMESASSIGN DIVIDEASSIGN
35 %token MODULOASSIGN
  %token NOT AND OR EQ NEQ LT LEQ GT GEQ
  %token TRUE FALSE
  %token VAR LET
39 %token FUN TO BY
40 %token RETURN CONTINUE BREAK IF ELSE FOR WHILE
  %token AMP CONST
42 %token <int> INT
43 %token <int> FLOAT
  %token BOOL VOID STRING MEMORY
  %token AUTO
46 %token NAMESPACE
47 %token IMPORT
  %token <string> ID
  \%token <string> STRINGLITERAL
50 %token < Num.num> INTLITERAL
51
  %token <Num.num> FLOATLITERAL
  %token EOF
  %right ASSIGN
55 %right PLUSASSIGN MINUSASSIGN
  %right TIMESASSIGN DIVIDEASSIGN MODULOASSIGN
57 %left OR
58 %left AND
  %left EQ NEQ
  %left LT GT LEQ GEQ
61 %left PLUS MINUS
62 %left TIMES DIVIDE MODULO
63 %right NOT NEG
```

```
%left LSQUARE
   %left LPAREN
   %left MINUSMINUS
   %left PLUSPLUS
67
68
   %start program
69
   %type<Ast.program> program
70
71
72
73
    qualified_id_builder:
    | ID { [$1] }
74
    | qualified_id_builder DOT ID { $3 :: $1 }
75
76
    qualified id:
77
    | qualified_id_builder { List.rev $1 }
78
79
    builtin_type:
      AUTO { Ast.Auto }
81
      VOID { Ast. Void }
82
      BOOL { Ast. Bool }
83
84
      INT \{ Ast. Int(\$1) \}
      FLOAT \ \{ \ Ast. \, \textcolor{red}{\textbf{Float}} \, (\,\$1\,) \ \ \}
85
      STRING { Ast. String
86
     MEMORY { Ast. Memory }
87
88
    array_spec:
89
     LSQUARE RSQUARE { 1 }
90
    | LSQUARE array_spec RSQUARE { 1 + $2 }
91
92
    int_literal_list:
     INTLITERAL { [ ( Num.int_of_num $1 ) ] }
94
     INTLITERAL int_literal_list { ( Num.int_of_num $1 ) :: $2
        }
96
    sized_array_spec:
97
    LSQUARE int_literal_list RSQUARE { ( 1, $2 ) }
     LSQUARE sized_array_spec RSQUARE { let ( d, el ) = $2 in
99
        (1 + d, el)
100
    type_category: { (false, false) }
101
    | AMP { (false, true) }
102
      CONST AMP { (true, true) }
103
    | CONST { (true, false) }
104
105
   sub_type_name:
```

```
| type_category builtin_type
                                               { Ast. Builtin Type (
       $2, $1) }
   type_category builtin_type array_spec { Ast.Array(Ast.
108
       BuiltinType ($2, Ast. no qualifiers), $3, $1) }
   type_category builtin_type sized_array_spec { let (d, el
109
       ) = \$3 in
      if d \Leftrightarrow ( List.length el ) then
110
        raise (Parsing. Parse error)
111
112
        Ast. SizedArray (Ast. BuiltinType ($2, Ast. no qualifiers),
113
       d, el, $1)
114
115
   sub_type_name_list_builder: { [] }
116
     sub_type_name { [$1] }
117
     sub_type_name_list_builder COMMA sub_type_name { $3 :: $1
118
119
   sub type name list:
120
   sub_type_name_list_builder { List.rev $1 }
121
   type_name:
123
     sub_type_name { $1 }
    type_category LPAREN sub_type_name_list RPAREN
125
       sub type name { Ast.Function(\$5, \$3, \$1) }
126
   expression_comma_list:
127
     expression
                                                  [$1] }
128
     expression COMMA expression_comma_list { $1 :: $3 }
129
130
   op expression:
131
     expression TIMES expression { Ast. BinaryOp($1, Ast. Mult,
   expression DIVIDE expression { Ast.BinaryOp($1, Ast.Div,
       $3)
   expression PLUS expression { Ast.BinaryOp($1, Ast.Add, $3
       ) }
   expression MINUS expression { Ast. BinaryOp($1, Ast. Sub,
135
       $3) }
   expression MODULO expression { Ast. BinaryOp($1, Ast.
136
       Modulo, $3)
   expression TIMESASSIGN expression { Ast. BinaryOp($1, Ast.
137
       MultAssign, $3) }
   expression DIVIDEASSIGN expression { Ast.BinaryOp($1, Ast
       . DivAssign, \$3) }
```

```
expression PLUSASSIGN expression { Ast.BinaryOp($1, Ast.
       AddAssign, $3) }
   expression MINUSASSIGN expression { Ast.BinaryOp($1, Ast.
140
       SubAssign, $3)
     expression MODULOASSIGN expression { Ast.BinaryOp($1, Ast
141
       . ModuloAssign, $3) }
     expression LT expression { Ast. BinaryOp($1, Ast. Less, $3)
   expression GT expression { Ast. BinaryOp($1, Ast. Greater,
143
   expression LEQ expression { Ast.BinaryOp($1, Ast.Leq, $3)
144
   expression GEQ expression { Ast.BinaryOp($1, Ast.Geq, $3)
145
     expression NEQ expression { Ast.BinaryOp($1, Ast.Neq, $3)
     expression EQ expression { Ast.BinaryOp($1, Ast.Equal, $3
     expression AND expression { Ast.BinaryOp($1, Ast.And, $3)
148
      expression OR expression { Ast. BinaryOp($1, Ast.Or, $3) }
149
      expression ASSIGN expression { Ast.Assignment($1, $3) }
     MINUS expression %prec NEG { Ast.PrefixUnaryOp(Ast.Neg,
       $2) }
     NOT expression { Ast. PrefixUnaryOp(Ast. Not, $2) }
152
     PLUSPLUS expression { Ast. PrefixUnaryOp(Ast. PreIncrement,
        $2) }
    | MINUSMINUS expression { Ast. PrefixUnaryOp(Ast.
154
       PreDecrement, $2) }
   value expression:
156
157
     INTLITERAL { let v = \text{match } \$1 \text{ with }
      Num. Int(i) -> Ast. IntLit((Int64.of_int i), Base.
158
       default_integral_bit_width)
      Num. Big_int(bi) ->
        begin try
160
          Ast.IntLit(Int64.of_int(Big_int.int_of_big_int bi)
161
         32)
        with
162
            -> Ast.IntLit ( Big_int.int64_of_big_int bi, 64 )
163
        \quad \text{end} \quad
164
      | n -> Ast. FloatLit ( Num. float of num n, Base.
165
       default floating bit width )
166
      Ast. Literal (v)
167
```

```
168
     FLOATLITERAL { Ast. Literal (Ast. FloatLit (Num. float of num
         $1 ), Base.default_floating_bit_width )) }
170
     STRINGLITERAL { Ast. Literal (Ast. StringLit ($1)) }
     TRUE { Ast. Literal (Ast. BoolLit (true)) }
171
172
     FALSE { Ast. Literal (Ast. BoolLit (false)) }
     LSQUARE expression_comma_list RSQUARE { Ast.
173
       ArrayInitializer($2)}
     LBRACE\ expression\_comma\_list\ RBRACE\ \{\ Ast.
174
       ObjectInitializer ($2 ) }
175
    postfix_expression:
176
     expression LSQUARE expression_comma_list RSQUARE { Ast.
       Index ($1, $3)
     expression LPAREN expression_comma_list RPAREN { Ast. Call
       (\$1, \$3)
                 }
179
    expression:
180
      qualified_id { Ast.QualifiedId($1) }
181
      value_expression { $1 }
182
      value_expression DOT qualified_id { Ast.Member($1, $3) }
183
      op_expression { $1 }
184
      postfix_expression { $1 }
185
      postfix_expression DOT qualified_id { Ast.Member($1, $3)
186
     LPAREN expression RPAREN { $2 }
187
188
     LPAREN expression RPAREN DOT qualified id { Ast. Member ($2
       , $5) }
189
    type_spec:
190
    | COLON type_name { $2 }
191
192
   maybe_type_spec: { Ast.BuiltinType(Ast.Auto, Ast.
193
       no_qualifiers) }
    | COLON type_name { $2 }
194
195
    binding:
196
    | ID type_spec { ($1, $2) }
197
198
    binding_list: { [] }
199
     binding { [$1] }
200
    | binding COMMA binding list { $1 :: $3 }
201
202
   var_binding:
203
    | ID type_spec { ($1, $2) }
```

```
ID { ($1, Ast.BuiltinType(Ast.Auto, Ast.no_qualifiers)) }
205
206
    variable definition:
207
    VAR var_binding ASSIGN expression { Ast. VarBinding($2, $4
208
    LET var_binding ASSIGN expression { Ast. VarBinding(Ast.
209
       add_const($2), $4) }
     VAR var_binding { Ast. VarBinding($2, Ast. Noop) }
210
    LET var binding { Ast. VarBinding (Ast. add const($2), Ast.
211
       Noop) }
212
    statement_list_builder: { [] }
213
    | statement_list_builder statement { $2 :: $1 }
215
    statement_list :
216
    statement_list_builder { List.rev $1 }
217
218
    parallel_binding:
219
     INVOCATIONS ASSIGN expression { Ast. Invocations ($3) }
220
     THREADCOUNT ASSIGN expression { Ast. ThreadCount($3) }
221
222
    parallel_binding_list_builder: { [] }
223
      parallel_binding { [$1] }
224
      parallel_binding_list_builder COMMA parallel_binding { $3
225
        :: $1 }
226
    parallel_binding_list:
227
    | parallel_binding_list_builder { List.rev $1 }
228
229
   sub_general_statement:
     expression { Ast. ExpressionStatement($1) }
231
232
     variable_definition { Ast. VariableStatement($1) }
    general_statement:
234
    | sub_general_statement SEMI { Ast.General($1) }
235
236
    control_initializer_builder:
237
     sub_general_statement { ( [$1], 1 ) }
238
     control_initializer_builder SEMI sub_general_statement {
239
       let (1, c) = \$1 in (\$3 :: 1, 1 + c)
240
    control initializer:
241
    | control initializer builder \{ let ( il , c ) = \$1 in if c
       < 2 then ([], List.hd il) else (List.rev (List.tl il),
        List.hd il) }
```

```
243
   sub general statement list builder: { [] }
    | sub_general_statement { [$1] }
245
    | sub_general_statement_list_builder COMMA
246
       sub_general_statement { $3 :: $1 }
247
   sub_general_statement_list:
    sub general statement list builder { List.rev $1 }
249
250
   statement:
251
     general_statement { $1 }
    IF LPAREN control_initializer RPAREN LBRACE
253
       statement_list RBRACE { Ast.IfBlock($3,$6) }
   | IF LPAREN control_initializer RPAREN LBRACE
       statement_list RBRACE ELSE LBRACE statement_list RBRACE
       { Ast. If Else Block ($3,$6,$10)
   | WHILE LPAREN control_initializer RPAREN LBRACE
       statement_list RBRACE { Ast.WhileBlock($3, $6) }
   FOR LPAREN sub general statement list SEMI expression
       SEMI_expression_comma_list_RPAREN_LBRACE_statement_list
       RBRACE { Ast.ForBlock($3, $5, $7, $10) }
    FOR LPAREN expression TO expression BY expression RPAREN
       LBRACE statement_list RBRACE { Ast.ForByToBlock($3, $5,
       $7, $10) }
     RETURN expression SEMI { Ast.Return($2) }
258
     RETURN SEMI { Ast. Return (Ast. Noop) }
259
260
     BREAK SEMI { Ast. Break(1) }
     BREAK INTLITERAL SEMI { Ast.Break( Num.int_of_num $2 ) }
261
     CONTINUE SEMI { Ast. Continue }
262
     PARALLEL LPAREN parallel_binding_list RPAREN LBRACE
       statement_list RBRACE { Ast.ParallelBlock($3, $6) }
     PARALLEL LBRACE statement_list RBRACE { Ast.
       ParallelBlock ([Ast.ThreadCount(Ast.Literal(Ast.IntLit(
       Int64.of_int(-1), Base.default_integral_bit_width)));
       Ast. Invocations (Ast. Literal (Ast. IntLit (Int64. of _int(-1),
        Base.default_integral_bit_width)))], $3) }
    ATOMIC LBRACE statement list RBRACE { Ast.AtomicBlock($3)
265
266
   function definition:
267
    FUN ID LPAREN binding_list RPAREN maybe_type_spec LBRACE
       statement list RBRACE { ([$2], $4, $6, $8) }
269
   import definition:
   | IMPORT qualified_id { $2 }
```

```
272
   definition list : { [] }
     definition_list import_definition { Ast.Import(Ast.
       LibraryImport(\$2)) :: \$1 
   | definition_list function_definition { Ast.Basic(Ast.
275
       Function Definition (\$2) :: \$1
   | definition_list_variable_definition_SEMI { Ast.Basic(Ast.
276
       Variable Definition ($2) :: $1 }
   definition list NAMESPACE qualified id LBRACE
       definition list RBRACE { Ast.Namespace($3, List.rev $5)
       :: $1 }
278
   program:
279
   definition list EOF { Ast.Program(List.rev $1) }
280
                          ../source/parser.mly
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```

```
20
   (* Drives the typical lexing and parsing algorithm
   while adding pertinent source, line and character
      information. *)
23
   type context = {
24
        mutable source_name : string;
25
        mutable source code: string;
26
     mutable original source code: string;
27
     mutable token_count : int;
28
        mutable token : Parser.token * Base.token_source;
29
   }
30
31
   let lex sourcename lexbuf =
32
     let rec acc lexbuf tokennumber tokens =
       let next_token = Scanner.token lexbuf
34
       and startp = Lexing.lexeme_start_p lexbuf
       and endp = Lexing.lexeme_end_p lexbuf
36
       let line = startp.Lexing.pos_lnum
38
39
       and relpos = (1 + startp.Lexing.pos\_cnum - startp.
       Lexing.pos_bol)
       and endrelpos = (1 + endp. Lexing.pos_cnum - endp. Lexing
40
       .pos_bol)
       and abspos = startp. Lexing.pos cnum
41
       and endabspos = endp. Lexing.pos cnum
43
       in
       let create_token token =
44
         let t = ( token, { Base.token_source_name =
45
      sourcename; Base.token_number = tokennumber;
           Base.token_line_number = line; Base.
46
       token_line_start = startp.Lexing.pos_bol;
           Base.token_column_range = (relpos, endrelpos); Base
47
       .token_character_range = (abspos, endabspos) }
         ) in
48
         t
49
       in
50
       match next token with
       | Parser.EOF as token -> ( create_token token ) ::
       | token -> ( create_token token ) :: ( acc lexbuf ( 1 +
       tokennumber ) tokens )
     in
54
     acc lexbuf 0 []
55
```

```
let parse context token_list =
     (* Keep a reference to the original token list
     And use that to dereference rather than whatever crap we
59
      get from
     the channel *)
60
     let tokenlist = ref(token list) in
61
     let tokenizer = match !tokenlist with
     (* Break each token down into pieces, info and all*)
     (token, info) :: rest ->
64
       context.source name <- info.Base.token source name;
65
       context.token_count <- 1 + context.token_count;</pre>
       context.token <- ( token, info );</pre>
67
       (* Shift the list down by one by referencing
       the beginning of the rest of the list *)
69
       tokenlist := rest;
       (* return token we care about *)
71
       token
72
     (* The parser stops calling the tokenizer when
73
     it hits EOF: if it reaches the empty list, WE SCREWED UP
      *)
     [] -> raise (Errors.MissingEoF)
     in
76
     (* Pass in an empty channel built off a cheap string
     and then ignore the fuck out of it in our 'tokenizer'
     internal function *)
     let program = Parser.program tokenizer (Lexing.
80
      from_string "") in
     program
81
82
   let analyze program =
     (* TODO: other important checks and semantic analysis
84
      here
     that will create a proper checked program type*)
85
     let sem = Semant.check program in
                          ../source/driver.ml
   (* LePiX Language Compiler Implementation
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19
20
   (* Semantic checking for the Lepix compiler that will
21
      produce a new
   SemanticProgram type with things like locals group into a
      single type
   and type promotions / conversions organized for operators.
23
      *)
24
   module StringMap = Map. Make(String)
25
26
   let extract_binding = function
     | Ast. VarBinding (b, _) \rightarrow b
28
   let extract_s_binding = function
30
     | Semast.SVarBinding(b, _) -> b
31
   let extract s binding name = function
33
    | (n, _) -> n
34
35
   let extract_s_binding_type = function
     (_, tn) \rightarrow tn
37
38
   let create_s_attributes () = {
39
     Semast.attr_parallelism = false;
```

```
Semast.attr\_arrays = 0;
41
     Semast.attr strings = false;
42
   }
43
44
   let create_s_environment () = {
45
     Semast.env_usings = [];
46
     Semast.env_symbols = StringMap.empty;
47
     Semast.env definitions = StringMap.empty;
     Semast.env\_imports = [];
49
50
     Semast.env\_loops = [];
   }
51
   let enter_block envl locals parameters =
     let acc_symbols m l =
54
       let ( n, tn ) = (extract_s_binding_name 1,
       extract_s_binding_type 1) in
       StringMap.add n tn m
56
     in
57
     let symbols = List.fold left acc symbols StringMap.empty
       parameters in
     let symbols = List.fold_left acc_symbols symbols locals
      in
     let env = {
60
       Semast.env\_usings = [];
61
       Semast.env symbols = symbols;
62
       Semast.env_definitions = StringMap.empty;
63
64
       Semast.env\_imports = [];
       Semast.env\_loops = [];
65
     } in
66
     env :: envl
67
68
   let enter_parameter_block envl parameters =
70
     let acc_symbols m l =
       let ( n, tn ) = (extract_s_binding_name l,
       extract_s_binding_type 1) in
       StringMap.add n tn m
72
     in
     let m k v1 v2 = match (v1, v2) with
74
         Some(x), Some(\underline{\ }) \rightarrow Some(x)
75
          Some(x) as s, None \rightarrow s
76
         None, (Some(y) as s) \rightarrow s
          _ -> None
78
     in
79
     let symbols = List.fold_left acc_symbols StringMap.empty
       parameters in
```

```
let env = List.hd envl in
81
      let env = \{ env with \}
82
        Semast.env_symbols = (StringMap.merge m symbols env.
83
       Semast.env_symbols);
      } in
84
85
      env :: envl
86
    let lookup id name mapl =
87
      let rec find = function
88
          [] -> None
89
        | h :: tl -> try Some ( StringMap.find name h )
90
          with | _ -> find tl
91
      in
92
      find mapl
93
    let env_lookup_id name envl =
95
      let mapl = ( List.map ( fun env -> env.Semast.env_symbols
        envl ) in
      lookup id name mapl
97
98
    let accumulate_string_type_bindings syms (n, qt) =
99
      let rec list_cmp v1 v2 = match v1, v2 with
100
          hl::tll, hr::tlr \rightarrow hl = hr && list\_cmp tll tlr
101
          [], [] -> true
102
          _, [] -> false
103
          [], _ -> false
104
105
      in
106
        let check_f l =
107
          let qt_argst = Semast.args_type_name qt in
108
          let pred t =
109
            let argst = Semast.args_type_name t in
            ( List.length argst ) = ( List.length qt_argst )
111
            && ( list_cmp argst qt_argst )
          in
          begin try
114
            let _ = List.find pred l in
            raise (Errors. Function Already Exists ("an id with name
116
        " \hat{} n \hat{} " and type " \hat{} ( Representation.
       string_of_s_type_name qt ) ^ " is already present"))
          with
117
            Not found -> Semast.SOverloads( qt :: 1)
118
          end
119
120
        let v = StringMap.find n syms in
121
```

```
let vt = match v with
            Semast. SOverloads (tl) -> check f tl
123
            Semast. SFunction (\_,\_,\_) as t \rightarrow check_f [t]
124
            _ -> raise (Not_found)
126
        in
        StringMap.add n vt syms
127
      with Not_found ->
128
        StringMap.add n qt syms
129
130
   let import_builtin_module symbols = function
131
      | Semast.Lib -> begin
        let c\_bindings = [
133
          ("lib.print", Semast.SFunction(Semast.void_t, [
134
       Semast.int32_t], Semast.no_qualifiers));
          ("lib.print", Semast.SFunction(Semast.void_t,
135
       Semast.string_t], Semast.no_qualifiers));
          ("lib.print", Semast.SFunction(Semast.void_t,
136
       Semast.float64_t], Semast.no_qualifiers));
          ("lib.print_n", Semast.SFunction(Semast.void_t, [
137
       Semast.int32_t], Semast.no_qualifiers));
          ("lib.print_n", Semast.SFunction(Semast.void_t, [
138
       Semast.string_t], Semast.no_qualifiers));
          ("lib.print\_n"\ ,\ Semast.SFunction(\ Semast.void\_t\ ,\ [
139
       Semast.float64_t], Semast.no_qualifiers));
140
       in
141
142
        print_endline "Here!";
        let symbols = List.fold_left
143
       accumulate_string_type_bindings symbols c_bindings in
        (symbols)
145
   let rec type_name_of_ast_type_name = function
147
       Ast. BuiltinType(bt, tq) -> Semast. SBuiltinType(bt, tq
148
      Ast. Array(tn, d, tq) -> Semast. SArray(
149
       type_name_of_ast_type_name_tn ), d, tq )
      | Ast. SizedArray(tn, d, dims, tq) -> Semast. SSizedArray(
150
       ( type_name_of_ast_type_name tn ), d, dims, tq )
      | Ast.Function(tn, pl, tq) -> Semast.SFunction(
151
       type_name_of_ast_type_name tn ), ( List.map
       type name of ast type name pl), tq)
152
   let type_name_of_ast_literal attrs envl astlit =
153
      let t = match astlit with
```

```
| Ast. BoolLit(_) -> Ast. BuiltinType( Ast. Bool, Ast.
       no qualifiers )
        | Ast.IntLit(_, b) -> Ast.BuiltinType( Ast.Int(b), Ast.
156
       no qualifiers )
        | Ast.FloatLit(_, b) -> Ast.BuiltinType( Ast.Float(b),
157
       Ast.no_qualifiers )
        | Ast. StringLit(_) -> Ast. BuiltinType( Ast. String, Ast.
158
       no qualifiers )
159
     type name of ast type name t
160
161
   let check_binary_op_common_type lt bop rt = match (lt, rt)
162
       with
       (Semast. SBuiltinType (Ast. Int (n), ltq), Semast.
163
       SBuiltinType(Ast.Int(m), rtq)) -> Semast.SBuiltinType(
       Ast. Int (max n m), Semast. no qualifiers)
       (Semast. SBuiltinType (Ast. Float (n), ltq), Semast.
164
       SBuiltinType (Ast.Int(m), rtq)) -> Semast.SBuiltinType (
       Ast. Float (max n m), Semast. no qualifiers)
      (Semast. SBuiltinType (Ast. Int (n), ltq), Semast.
165
       SBuiltinType(Ast.Float(m), rtq)) -> Semast.SBuiltinType(
       Ast. Float (max n m), Semast.no_qualifiers)
       (Semast. SBuiltinType (Ast. Float (n), ltq), Semast.
166
       SBuiltinType (Ast. Float (m), rtq)) -> Semast. SBuiltinType (
       Ast. Float (max n m), Semast. no qualifiers)
      (Semast. SBuiltinType (Ast. Bool, ltq), Semast.
167
       SBuiltinType (Ast. Bool, rtq)) -> Semast. SBuiltinType (Ast.
       Bool, Semast.no_qualifiers)
      _ -> raise (Errors. InvalidBinary Operation ("cannot
168
       perform a binary operation on two non-numeric types"))
170
   let check_unary_op uop sr = match sr with
      | Semast. SBuiltinType (Ast. Int (n), ltq) -> Semast.
171
       SBuiltinType (Ast. Int (n), Semast. no qualifiers)
      | Semast. SBuiltin Type (Ast. Float (n), ltq) \rightarrow Semast.
       SBuiltinType (Ast. Float (n), Semast. no qualifiers)
      _ -> raise (Errors. InvalidUnary Operation ("cannot perform
        a unary operation on this type"))
174
   let overload resolution args = function
175
      | Semast.SFunction(__, tnl, _) -> let argslen = ( List.
176
       length args ) in
        let overloadlen = ( List.length tnl ) in
177
        argslen = overloadlen &&
178
        if (argslen > 0) then (List.exists (fun a -> List.
```

```
mem a tnl ) args )
        else true
180
      _ -> raise (Errors. TypeMismatch ("cannot resolve an
181
       overload that includes a non-function in its type
       listing"))
182
   let check_function_overloads args overloadlist =
183
184
        let ft = List.find ( overload resolution args )
185
       overloadlist
        in
186
        ( ft , Semast.return_type_name ft )
187
      with _{-} \rightarrow
188
        let argslist = "( " ^ ( String.concat "," ( List.map
189
       Representation.string_of_s_type_name args ) ) ^ " )" in
        raise (Errors. BadFunctionCall ("could not resolve the
190
       specific overload for this set of " ^ ( Representation.
       string_of_s_type_name ( Semast.SOverloads(overloadlist)
       ) ) ^ " using " ^ argslist ))
      end
191
192
   let rec type_name_of_ast_expression attrs envl astexpr =
193
      let t = match astexpr with
194
        Ast. Literal (lit) -> type_name_of_ast_literal attrs
195
       envl lit
        Ast. QualifiedId (qid) -> let qualname = Semast.
196
       string_of_qualified_id qid in
          begin match ( env\_lookup\_id qualname envl ) with
197
            None -> raise (Errors.IdentifierNotFound ("
198
       identifier '" ^ qualname ^ "' not found"))
            | Some(stn) \rightarrow stn
199
200
          end
        | Ast. Call(e, args) -> let ft =
201
       type_name_of_ast_expression attrs envl e in
          (* TODO: check arguments to make sure it matches *)
202
          begin match ft with
203
            | Semast.SFunction(rt, pl, tq) ->
204
205
              Semast. SOverloads (fl) ->
206
              let sargs = List.map (
207
       type_name_of_ast_expression attrs envl ) args in
              let (ft, r) = check function overloads sargs fl
208
       in
209
            _ -> raise (Errors. TypeMismatch ("expected a
210
```

```
function type, but received something else."))
          end
211
          Ast. Noop -> Semast. void t
212
          Ast. Binary Op (1, bop, r) \rightarrow
213
          let sl = type_name_of_ast_expression attrs envl l
214
          and sr = type_name_of_ast_expression attrs envl r
215
216
          check binary op common type sl bop sr
217
          Ast. PrefixUnaryOp (uop, r) ->
218
          let sr = type name of ast expression attrs envl r in
219
          check_unary_op uop sr
220
         Ast. Assignment (lhs, rhs) -> let lhst =
221
       type_name_of_ast_expression attrs envl lhs in
          lhst
222
         Ast. Member (_, _) -> raise (Errors. Unsupported ("member
       access is not supported"))
        _ -> raise (Errors. Unsupported ("expression conversion
       currently unsupported"))
225
      (* TODO: some type checks to make sure weird things like
226
       void& aren't put in place ... *)
227
228
   let process_ast_import prefix symbols defs imports =
229
       function
      | Ast.LibraryImport(qid) ->
230
        let qualname = Semast.string_of_qualified_id qid in
231
        let (v, impsymbols) = match List.filter (fun (n, _) ->
232
        n = qualname ) Semast.builtin_library_names with
            (_, bltin) :: [] -> let b = Semast.SBuiltin(bltin)
233
       in
234
            let ( bltinsymbols ) = import_builtin_module
       symbols bltin in
            (b, bltinsymbols)
            _ -> ( Semast.SDynamic(qualname), symbols )
236
        in
237
        ( prefix, impsymbols, defs, v :: imports )
238
239
240
   let generate_global_env = function
241
     Ast. Program (ast_definitions) ->
242
      let rec acc ( prefix , symbols , defs , imports ) def =
243
        match def with
244
          Ast.Import(imp) -> process_ast_import_prefix
245
       symbols defs imports imp
```

```
| Ast.Basic(Ast.FunctionDefinition((qid, args, rt, _)
246
       )) ->
            let argst = List.map Ast.binding_type args in
247
            let qualname = prefix ^ Semast.
248
       string_of_qualified_id qid in
            let qt = Ast.Function(rt, argst, Ast.no_qualifiers)
249
        in
            let sqt = (type name of ast type name qt) in
250
            print endline "here";
251
            let nsymbols = accumulate string type bindings
252
       symbols (qualname, sqt)
            and ndefs = accumulate_string_type_bindings defs (
253
       qualname, sqt)
            in
254
            ( prefix, nsymbols, ndefs, imports )
            Ast. Basic (Ast. Variable Definition (v)) ->
256
            let (name, qt) = extract_binding v in
257
            let qualname = prefix ^ name in
258
            if StringMap.mem prefix symbols then raise (Errors.
259
       VariableAlreadyExists(qualname)) else
260
            let nsymbols = (StringMap.add qualname (
       type_name_of_ast_type_name_qt ) symbols )
            and ndefs = (StringMap.add qualname (
261
       type_name_of_ast_type_name_qt_) defs_)
            in
262
            ( prefix, nsymbols, ndefs, imports )
263
            Ast. Namespace (n, dl) ->
264
            let qualname = prefix ^ Semast.
265
       string_of_qualified_id n in
            let (_, innersymbols, innerdefs, innerimports ) =
266
       List.fold_left acc ( qualname ^ ".", symbols, defs,
       imports ) dl in
            ( prefix, innersymbols, innerdefs, innerimports )
267
      _{
m in}
268
      let (_, symbols, defs, imports) = List.fold_left acc ("",
269
        StringMap.empty, StringMap.empty, []) ast_definitions
       in
      let attrs = create s attributes () in
270
      let env = {
271
        Semast.env\_usings = [];
272
        Semast.env_symbols = symbols;
273
        Semast.env\_definitions = defs;
274
        Semast.env imports = imports;
275
        Semast.env loops = [];
276
```

```
in
278
      ( attrs, env )
279
280
   let check qualified identifier attrs envl sl t =
281
      (attrs, envl, Semast.SQualifiedId(sl, t))
282
283
   let check_function_call attrs envl target args =
284
      let (t, rt) = match Semast.type name of s expression
285
       target with
          Semast. SFunction (tn, tnl, tq) as f \rightarrow f, tn
286
          Semast. SOverloads (fl) ->
          let args_t = (List.map Semast.
288
       type_name_of_s_expression args ) in
          check_function_overloads args_t fl
289
          _ -> raise (Errors. BadFunctionCall ("cannot invoke an
       expression which does not result in a function type of
       some sort"))
      in
291
      (attrs, envl, Semast.SCall((Semast.
292
       coerce_type_name_of_s_expression t target ), args, rt ))
293
   let generate_s_binding prefix attrs envl = function
294
       (name, tn) -> (attrs, envl, (Semast.
295
       string_of_qualified_id ( prefix @ [name] ),
       type name of ast type name tn))
296
   let gather_ast_locals attrs envl sl pl =
297
      let acc locals = function
298
         Ast. General (Ast. Variable Statement (v)) ->
299
          let (_, _, sb) = generate_s_binding [] attrs envl (
       extract_binding v ) in
          sb :: locals
        \mid _ \rightarrow locals
302
      in
303
      let l = List.rev( List.fold_left acc [] sl ) in
304
      if (List.length 1) > 0 then begin
        let envl = ( enter_block envl l pl ) in
306
        (true, attrs, envl, 1)
307
      end else
308
        let envl = ( enter_parameter_block envl pl ) in
309
        (false, attrs, envl, 1)
310
311
   let generate s literal attrs envl = function
312
       Ast. BoolLit(b) -> (attrs, envl, Semast. SBoolLit(b))
313
       Ast.IntLit(i, b) -> (attrs, envl, Semast.SIntLit(i, b))
```

```
| Ast.FloatLit(f, b) -> (attrs, envl, Semast.SFloatLit(f,
315
      | Ast. StringLit(s) -> (attrs, envl, Semast. SStringLit(s))
316
317
   let rec generate_s_expression attrs envl astexpr =
318
      let acc_s_expression (attrs, envl, sel) e =
319
        let (attrs, envl, se) = generate_s_expression attrs
320
       envl e in
        (attrs, envl, se :: sel)
321
322
      let (attrs, envl, se) = match astexpr with
323
        | Ast. Literal(lit) ->
324
          let ( attrs, envl, slit ) = generate_s_literal attrs
       envl lit in
          ( attrs, envl, Semast. SLiteral ( slit ) )
326
        | Ast. QualifiedId(sl) ->
327
          let t = type_name_of_ast_expression attrs envl
328
       astexpr in
          check qualified identifier attrs envl sl t
329
         Ast. Call(e, el) ->
330
331
          let ( attrs, envl, target ) = ( generate_s_expression
        attrs envl e ) in
          let ( attrs, envl, args ) = ( List.fold_left
       acc_s_expression (attrs, envl, []) el ) in
          let args = List.rev args in
333
          check function call attrs envl target args
334
335
         Ast. BinaryOp (lhs, bop, rhs) ->
          let ( attrs, envl, slhs ) = ( generate_s_expression
336
       attrs envl lhs ) in
          let ( attrs, envl, srhs ) = ( generate_s_expression
       attrs envl rhs ) in
338
          let rhst = ( Semast.type_name_of_s_expression srhs )
       in
          let lhst = ( Semast.type_name_of_s_expression slhs )
       in
          ( attrs, envl, Semast.SBinaryOp( slhs, bop, srhs, (
340
       check_binary_op_common_type lhst bop rhst ) ) )
        | Ast.PrefixUnaryOp(uop, rhs) ->
341
          let ( attrs, envl, srhs ) = ( generate_s_expression
342
       attrs envl rhs ) in
          let rhst = ( Semast.type_name_of_s_expression srhs )
343
          ( attrs, envl, Semast.SPrefixUnaryOp( uop, srhs, (
344
       check unary op uop rhst ) )
        | Ast. Assignment (lhs, rhs) ->
345
```

```
let ( attrs, envl, slhs ) = ( generate_s_expression
346
       attrs envl lhs ) in
          let ( attrs, envl, srhs ) = ( generate_s_expression
347
       attrs envl rhs ) in
          let lhst = ( Semast.type_name_of_s_expression slhs )
348
            attrs, envl, Semast.SAssignment(slhs, srhs, lhst)
349
          Ast. Noop -> ( attrs, envl, Semast. SNoop )
350
         _ -> raise (Errors. Unsupported ("expression generation
351
       for this type is current unsupported"))
352
      let t = Semast.type_name_of_s_expression se in
353
      let attrs = match t with
354
          Semast . SArray(\_,d,\_)
          Semast. SSizedArray(\_,d,\_,\_) \rightarrow \{
356
            Semast.attr_strings = attrs.Semast.attr_strings;
            Semast.attr_arrays = max d attrs.Semast.attr_arrays
358
            Semast.attr_parallelism = attrs.Semast.
359
       attr_parallelism;
360
        | Semast.SBuiltinType( Ast.String, _ ) -> {
361
            Semast.attr_strings = true;
362
            Semast.attr arrays = attrs.Semast.attr arrays;
363
            Semast.attr parallelism = attrs.Semast.
364
       attr_parallelism;
365
            \rightarrow attrs
366
      ( attrs, envl, se )
368
   let generate_s_variable_definition prefix attrs envl =
370
       function
      | Ast. VarBinding (b, e) ->
371
        let (attrs, envl, sb) = generate_s_binding prefix attrs
372
        envl b in
        let (attrs, envl, se) = generate s expression attrs
373
       envl e in
        (attrs, envl, Semast.SVarBinding(sb, se))
374
375
   let generate s general statement attrs envl = function
376
      | Ast. ExpressionStatement(e) ->
377
        let (attrs, envl, se) = generate_s_expression attrs
378
       envl e in
```

```
( attrs, envl, Semast.SExpressionStatement( se ) )
379
       Ast. VariableStatement (v) ->
380
        let (attrs, envl, sv ) = generate_s_variable_definition
381
        attrs envl v in
        ( attrs, envl, Semast.SVariableStatement( sv ) )
382
383
   let generate_s_statement attrs envl = function
384
       Ast. General (g) ->
385
        let (attrs, envl, sgs) = generate s general statement
386
        attrs envl g in
        ( attrs, envl, Semast.SGeneral( sgs ) )
       Ast.Return(e) ->
388
        let ( attrs, envl, se ) = generate_s_expression attrs
       envl e in
        ( attrs, envl, Semast.SReturn( se ) )
       _ -> raise (Errors. Unsupported ("statement type not
391
       supported"))
392
   let check returns name ssl rt =
393
      (* Todo: recursively inspect all inner blocks for return
394
       types as well *)
      let acc rl = function
395
        | Semast.SReturn(se) -> ( Semast.
396
       type_name_of_s_expression se ) :: rl
         _ -> rl
397
      in
398
      let returns = List.fold_left acc [] ssl in
399
      let returnlength = List.length returns in
400
      if name = "main" then begin
401
        let sret0 = Semast. SReturn (Semast. SLiteral (Semast.
       SIntLit (Int64.zero, 32))) in
        let mainpred = function
            Semast. SBuiltinType (Ast. Auto, _) ->
404
            Semast. SBuiltin Type (Ast. Int (32), (\_, r)) \rightarrow
406
            if r then raise (Errors. InvalidMainSignature ("Cannot
407
        return a reference from main"));
            ()
408
              ->
409
            raise (Errors. Invalid Main Signature ("You can only
410
       return an int from main"))
411
        let ssl = if returnlength < 1 then begin ssl @ [sret0]
412
       end else ssl
        in
413
```

```
let ssl = match rt with
414
            Semast. SBuiltinType (Ast. Int (32), (\_, r)) \rightarrow
415
             if r then raise (Errors. InvalidMainSignature ("Cannot
416
        return a reference to and integer"));
             let _ = List.iter mainpred returns in
417
418
          _ -> let _ = List.iter mainpred returns in
419
420
        in
421
        (ssl, Semast.int32_t)
422
      end
423
      else begin
424
        let generalpred rt r = match rt with
            Semast.SBuiltinType(Ast.Auto, _) -> r
426
            \_ \rightarrow let r = Semast.unqualify r in
            let urt = Semast.unqualify rt in
428
            if r \lor urt then raise (Errors.
429
       InvalidFunctionSignature ("return types do not match
       across all returns", name))
             else rt
430
431
        let rt = List.fold_left generalpred rt returns in
        let (ssl, rt) = match rt with
433
          | Semast.SBuiltinType(Ast.Auto, _) ->
434
             if returnlength < 1 then
435
               (ssl @ [Semast.SReturn(Semast.SNoop)], Semast.
       void_t )
             else
437
               (ssl, rt)
438
          | Semast.SBuiltinType(Ast.Void, _) ->
             if returnlength < 1 then
440
                 ssl @ [Semast.SReturn(Semast.SNoop)], Semast.
       void_t )
             else
              (ssl, rt)
443
            _ ->
444
            if returnlength < 1 then
445
               raise (Errors. InvalidFunctionSignature ("function
446
       was expected to return a value: returned no value", name
       ))
             else
               (ssl, rt)
448
        in
449
        (ssl, rt)
450
      end
```

```
452
   let generate s function definition prefix attrs envl
453
       astfdef =
      let acc_ast_statements (attrs, envl, ssl) s =
454
        let ( attrs, envl, ss ) = ( generate_s_statement attrs
455
       envl s ) in
        ( attrs, envl, ss :: ssl )
456
457
      let acc_ast_parameters (attrs, envl, pl) p =
458
        let (attrs, envl, sp) = generate s binding [] attrs
459
       envl p in
        (attrs, envl, sp :: pl)
460
      in
461
      let (qid, astparameters, astrt, body) = astfdef in
462
      let fqid = prefix @ qid in
      let fqn = Semast.string_of_qualified_id fqid in
464
      let (attrs, envl, parameters) = List.fold_left
       acc_ast_parameters (attrs, envl, []) astparameters in
      let rt = type_name_of_ast_type_name astrt in
466
      let (has_locals, attrs, envl, bl) = gather_ast_locals
467
       attrs envl body parameters in
      let (attrs, envl, ssl) = List.fold_left
468
       acc_ast_statements (attrs, envl, []) body in
      let ssl = List.rev ssl in
469
      let (ssl, rt) = check returns fqn ssl rt in
470
      let sfuncdef = if has locals then
471
472
        Semast.func_name = fqid;
473
        Semast.func_parameters = Semast.SParameters(parameters)
474
       Semast.func return type = rt;
475
        Semast.func_body = [Semast.SBlock(Semast.SLocals(bl),
476
       ssl);
      }
      else
478
479
        Semast.func name = fqid;
480
        Semast.func parameters = Semast.SParameters(parameters)
481
        Semast.func_return_type = rt;
482
        Semast.func\_body = ssl;
483
484
     in
485
      ( attrs, envl, Semast.SFunctionDefinition( sfuncdef ) )
486
```

```
let generate s basic definition prefix attrs envl =
489
       function
      | Ast. Function Definition (fdef) ->
490
        let (attrs, envl, sfdef) =
491
       generate_s_function_definition prefix attrs envl fdef in
        (attrs, envl, Semast.SBasic(sfdef))
       Ast. Variable Definition (vdef) ->
493
        let (attrs, envl, svdef) =
494
       generate s variable definition prefix attrs envl vdef in
        (attrs, envl, Semast.SBasic(Semast.SVariableDefinition(
495
       svdef)))
496
497
   let define_libraries attrs env =
498
      let fi = Semast.SFunction(Semast.void_t, [Semast.int32_t
499
       , Semast.no_qualifiers) in
      let ff = Semast.SFunction(Semast.void_t, [Semast.
500
       float64 t], Semast.no qualifiers) in
      let fs = Semast.SFunction(Semast.void_t, [Semast.string_t
501
       , Semast.no qualifiers) in
      let fo = Semast. SOverloads ([fi; ff; fs]) in
502
      let lib_printn_defint = {
        Semast.func_name = ["lib"; "print_n"];
504
        Semast.func_parameters = Semast.SParameters([("i",
505
       Semast.int32 	ext{ t)};
        Semast.func_return_type = Semast.void_t;
506
        Semast.func\_body = [
507
          Semast. SGeneral (Semast. SExpressionStatement (
508
            Semast. SCall (Semast. SQualifiedId (["lib"; "print"],
509
       fo), [Semast.SQualifiedId(["i"], Semast.int32_t)],
       Semast.void_t)
          ));
          Semast. SGeneral (Semast. SExpressionStatement (
            Semast.SCall(Semast.SQualifiedId(["lib"; "print"],
       fo), [Semast.SLiteral(Semast.SStringLit("\n"))], Semast.
       void_t)
          ));
          Semast. SReturn (Semast. SNoop);
514
        ];
      }
516
     and lib printn deffloat = \{
517
        Semast.func name = ["lib"; "print n"];
518
        Semast.func_parameters = Semast.SParameters([("i",
519
       Semast. float 64_t);
```

488

```
Semast.func_return_type = Semast.void_t;
        Semast.func body = [
521
          Semast. SGeneral (Semast. SExpressionStatement (
             Semast.\,SCall\,(\,Semast.\,S\,Qualified\,Id\,(\,[\,"\,li\,b\,"\,;\,\,"\,print\,"\,]\,\,,
        fo), [Semast.SQualifiedId(["i"], Semast.float64_t)],
       Semast. void t)
          ));
524
          Semast. SGeneral (Semast. SExpression Statement (
525
             Semast. SCall (Semast. SQualifiedId (["lib"; "print"],
526
        fo), [Semast.SLiteral(Semast.SStringLit("\n"))], Semast.
       void_t)
          ));
          Semast. SReturn (Semast. SNoop);
        ];
530
      and lib_printn_defstr = {
        Semast.func_name = ["lib"; "print_n"];
        Semast.func_parameters = Semast.SParameters([("i",
        Semast.string t);
        Semast.func_return_type = Semast.void_t;
535
        Semast.func body = [
          Semast. SGeneral (Semast. SExpressionStatement (
536
             Semast. SCall (Semast. SQualifiedId (["lib"; "print"],
537
        fo), [Semast.SQualifiedId(["i"], Semast.string_t)],
       Semast. void t)
          ));
538
          Semast. SGeneral (Semast. SExpression Statement (
539
             Semast. SCall (Semast. SQualifiedId (["lib"; "print"],
540
        fo), [Semast.SLiteral(Semast.SStringLit("\n"))], Semast.
       void_t)
          ));
541
542
          Semast. SReturn (Semast. SNoop);
        ];
543
      }
      in
545
      let libdefs = [
546
        Semast. SBasic (Semast. SFunctionDefinition (
       lib printn defstr));
        Semast. SBasic (Semast. SFunctionDefinition (
548
        lib printn defint));
        Semast. SBasic (Semast. SFunction Definition (
549
       lib printn deffloat));
      ] in
550
      let acc defs = function
551
          Semast. SBasic (Semast. SFunction Definition (fdef)) ->
```

```
let n = (Semast.string_of_qualified_id fdef.Semast.
       func name )
          and qt = (Semast.type_name_of_s_function_definition
554
       fdef )
          in
          if (StringMap.mem n defs ) then defs else
556
       accumulate string type bindings defs (n, qt)
         Semast. SBasic (Semast. SVariable Definition (Semast.
557
       SVarBinding((n, tn), _)) \rightarrow
          (StringMap.add n tn defs)
558
559
      in
      let ndefs = List.fold_left acc env.Semast.env_definitions
560
        libdefs in
      ( { env with Semast.env definitions = ndefs; }, libdefs)
561
562
   let direct_code_inject attrs globalenv imp sdl =
563
      let s = match imp with
564
        | Ast.LibraryImport(qid) -> Semast.
565
       string of qualified id qid
     in
566
567
      match s with
        "lib" -> let (globalenv, library_defs) =
568
       define libraries attrs globalenv in
          (globalenv, library_defs @ sdl)
569
        -> (globalenv, sdl)
572
   let generate semantic attrs globalenv = function
    | Ast. Program (dl) ->
573
      let envl = [globalenv] in
574
      let rec acc_ast_definitions (prefix, attrs, envl, sdl) =
       function
         Ast.Import(imp) ->
          let globalenv = List.hd ( List.rev envl ) in
577
          let (globalenv, sdl) = direct_code_inject attrs
       globalenv imp sdl in
          let tail = List.tl (List.rev envl) in
579
          (prefix, attrs, (globalenv:: tail), sdl)
580
         Ast. Namespace(n, dl) -> let qualname = prefix @ n in
581
          let (_, attrs, envl, nssdl) = List.fold_left
582
       acc ast definitions (qualname, attrs, envl, sdl) dl in
          ( prefix, attrs, envl, nssdl)
583
         Ast. Basic(b) ->
584
          let (attrs, envl, sb) = (generate s basic definition
585
        prefix attrs envl b ) in
          ( prefix, attrs, envl, sb :: sdl)
586
```

```
in
587
      let (, attrs, envl, sdefs) = List.fold left
588
       acc_ast_definitions ([], attrs, envl, []) dl in
      let globalenv = List.hd ( List.rev envl ) in
589
      Semast.SProgram( attrs, globalenv, List.rev sdefs)
590
591
    let modify_symbols = function
592
     Semast. SProgram (attrs, env, sdls) ->
593
      let rec acc ( symbols, defs ) def =
594
        match def with
595
           | Semast.SBasic(Semast.SFunctionDefinition(f)) ->
596
            let qualname = Semast.string_of_qualified_id f.
       Semast.func_name in
            let sqt = (Semast.
598
       type_name_of_s_function_definition f ) in
            let nsymbols = accumulate_string_type_bindings
599
       symbols (qualname, sqt)
            and ndefs = accumulate_string_type_bindings defs (
600
       qualname, sqt)
            in
601
602
            ( nsymbols, ndefs )
            Semast. SBasic (Semast. SVariable Definition (v)) ->
603
             ( symbols, defs )
604
605
      in
      let (symbols, defs) = List.fold left acc (StringMap.empty
606
         StringMap.empty) sdls in
607
      let m k l r = match (l, r) with
          Some(1), Some(\underline{\ }) \rightarrow Some(1)
608
          None, (Some(r) as s) \rightarrow s
609
          (Some(1) as s), None -> s
            \rightarrow None
611
612
      in
613
      let env = \{ env with \}
        Semast.env_symbols = StringMap.merge m symbols env.
       Semast.env_symbols;
        Semast.env_definitions = defs;
615
      }
616
      in
617
      Semast. SProgram (attrs, env, sdls)
618
619
    let check astprogram =
620
      (* Pass 1: Gather globals inside of all the namespaces
621
      so they can be referenced even before they're defined (
622
       just so long as
      they're in the same lateral global scope, not necessarily
623
```

```
in vertical order) *)
624
      let (attrs, env) = generate global env astprogram in
625
      (* Pass 2: Generate the actual Semantic Tree based on
626
       what
      is inside the AST program ... *)
627
      print endline "Pass 2";
628
      let sprog = generate_semantic attrs env astprogram in
      (* Pass 3: Update any symbols that were resolved during
630
     bottom-up type derivation ... *)
631
      print endline "Pass 3";
632
      let Semast.SProgram(attrs, env, _) = modify_symbols sprog
633
      (* Pass 4: Finalize everything with new information *)
634
     print endline "Pass 4";
635
      let sprog = generate_semantic attrs env astprogram in
637
     sprog
                          ../source/semant.ml
   (* LePiX Language Compiler Implementation
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```

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```
SOFTWARE OR THE USE OR OTHER DEALINGS IN THE SOFTWARE. *)
20
   (* Code generation: translate takes a semantically checked
21
      AST and produces
22
  LLVM IR:
   http://llvm.org/docs/tutorial/index.html
   http://llvm.moe/ocaml/ *)
24
25
   (* Linked code after the c bindings from the makefile
26
   compiled the ll for the c bindings *)
27
28
   module StringMap = Map. Make(String)
29
30
   type li_universe = {
31
     lu_attrs : Semast.s_attributes;
32
     lu_env : Semast.s_environment;
33
     lu_module : Llvm.llmodule;
34
     lu context : Llvm.llcontext;
     lu_builder : Llvm.llbuilder;
36
     lu variables : Llvm.llvalue StringMap.t;
37
     lu_functions : Llvm.llvalue StringMap.t;
38
     lu_named_values : Llvm.llvalue StringMap.t;
39
     lu_named_params : Llvm.llvalue StringMap.t;
40
     lu handlers : ( li universe -> ( Semast.s expression list
41
       ) -> ( li universe * Llvm.llvalue list ) ) StringMap.t;
42
43
   let create_li_universe = function | Semast.SProgram(attrs ,
44
      env, _) ->
     let context = Llvm.global context() in
45
     let builder = Llvm.builder context in
47
     let m = Llvm.create module context "lepix" in
       lu_attrs = attrs;
49
       lu_{env} = env;
       lu \mod ule = m;
       lu context = context;
       lu_builder = builder;
       lu variables = StringMap.empty;
54
       lu_functions = StringMap.empty;
       lu named values = StringMap.empty;
       lu named params = StringMap.empty;
57
       lu_handlers = StringMap.empty;
58
```

```
60
   let rec llvm type of s type name lu st =
     let f32 t
                 = Llvm.float_type
                                      lu.lu context
62
     and f64_t
               = Llvm.double type
63
                                      lu.lu context
     (* for 'char' type to printf —
                                      even if they resolve to
64
      same type, we differentiate*)
     and char_t = Llvm.i8_type
                                      lu.lu_context
65
     and i16 t
                 = Llvm.i16 type
                                      lu.lu context
66
                 = Llvm.i32 type
     and i32 t
                                      lu.lu context
67
     and i64 t
                 = Llvm.i64_type
                                          lu.lu context
68
     (* LLVM treats booleans as 1-bit integers, not distinct
      types with their own true / false *)
     and bool_t = Llvm.i1_type
                                      lu.lu_context
70
     and void t = Llvm.void type
                                      lu.lu context
71
     let p_char_t = Llvm.pointer_type
                                            char t
73
74
     match st with
75
     (* TODO: handle reference-ness (e.g., make it behave like
       a pointer here) *)
       Semast. SBuiltinType (Ast. Bool, tq ) -> bool t
     | Semast.SBuiltinType(Ast.Int(n), tq) -> begin match n
78
      with
         64 \rightarrow i64_t
79
         32 -> i32 t
80
         16 -> i16 t
         _ -> Llvm.integer_type lu.lu_context n
82
     end
83
     | Semast.SBuiltinType(Ast.Float(n), tq ) -> begin match
84
      n with
         64 -> f64 t
85
         32 -> f32 t
        16 -> (* LLVM actually has support for this, but
      shitty OCaml bindings *)
         (* TODO: Proper Error *)
88
         raise (Failure "Cannot have a Half Float because
      OCaml binding for LLVM is garbage" )
       -> (* TODO: Proper Error *)
90
         raise (Failure "Unallowed Float Width")
91
92
       Semast.SBuiltinType(Ast.String, tq) -> p_char_t
93
       Semast.SBuiltinType(Ast.Void, tq) -> void t
     | Semast.SArray(t, d, tq) -> Llvm.array type (
95
      llvm_type_of_s_type_name_lu_t) d
     | Semast.SSizedArray(t, d, szs, tq) -> Llvm.array_type (
```

```
llvm_type_of_s_type_name_lu_t) d
       Semast. SFunction (rt, argst, tq) ->
97
        let lrt = llvm_type_of_s_type_name lu rt
98
        and largst = Array.map (llvm_type_of_s_type_name lu)
       ( Array.of_list argst )
        in
100
        Llvm.function_type lrt largst
101
       -> (* TODO: Proper Error *)
102
        raise (Errors. Unsupported ("This type is not convertible
103
       to an LLVM type"))
104
    let should_reference_pointer = function
105
        Semast. SBuiltinType (Ast. String, _) -> true
        Semast.SArray(\_, \_, \_) \rightarrow true
107
        Semast.SSizedArray(_, _, _, _) -> true
108
        Semast.SFunction(\_, \_, \_) \rightarrow true
109
        _ -> false
    let find argument handler lu target =
      let hn = Llvm.value_name target in
      try Some (StringMap.find hn lu.lu_handlers)
114
      with _ -> None
116
    let llvm_lookup_function lu name t =
117
      let mname = Semast.mangle name [name] t in
118
      match Llvm.lookup_function mname lu.lu_module with
119
120
          Some(v) \rightarrow v
          None -> raise ( Errors. FunctionLookupFailure ( name,
121
       mname ) )
    let llvm_lookup_variable lu name t =
123
124
      match Llvm.lookup_global name lu.lu_module with
          Some(v) \rightarrow v
          None -> raise ( Errors. VariableLookupFailure ( name,
       name ) )
127
    let dump_s_qualified_id lu qid t =
128
      let fqn = Semast.string_of_qualified_id qid in
129
      let lookup n =
130
        try
          let v = StringMap.find n lu.lu_named_values in
          Some (v)
        with | Not found ->
134
135
          try
            let v = StringMap.find n lu.lu_named_params in
```

```
Some(v)
        with | Not found -> try
138
             let v = StringMap.find n lu.lu_variables in
139
             Some (v)
140
        with | Not_found -> None
141
142
      let lookup_func n =
143
        try
144
           let v = StringMap.find n lu.lu named values in
145
146
        with | Not_found ->
           try
148
             let v = StringMap.find n lu.lu_named_params in
149
             Some (v)
150
        with | Not_found -> try
151
             let v = StringMap.find n lu.lu_functions in
152
             Some (v)
153
        with | Not_found -> None
155
      let overload_lookup qid ft =
156
157
        let mangled = Semast.mangle name qid ft in
        begin match lookup_func mangled with
158
             Some(v) as s \rightarrow s
159
             None -> begin match lookup_func fqn with
160
               Some(v) as s \rightarrow s
161
               None -> None
162
163
             end
        end
164
      in
165
      let overload_acc op ft =
        match op with
167
168
             Some(v) as s \rightarrow s
             None -> overload_lookup qid ft
169
170
      in
      let idval = match t with
          Semast. SFunction (rt, tnl, tq) as ft ->
           begin match overload_lookup qid ft with
               Some(v) \rightarrow v
174
               None -> raise (Errors. UnknownFunction(fqn))
175
176
          Semast. SOverloads (fl)->
           begin match List.fold left overload acc None fl with
178
               Some(v) \rightarrow v
179
               None -> raise (Errors. UnknownFunction (fqn))
180
           end
```

```
_ -> match lookup fqn with
182
            | Some(v) \rightarrow v
183
             | None -> raise (Errors.UnknownVariable fqn)
184
185
      (lu, idval)
186
187
    let dump_s_literal lu lit =
188
      let f64 t = Llvm.double type
                                          lu.lu context
189
      and bool t = Llvm.i1 type
                                         lu.lu context
190
191
      let v = match lit with
192
        | Semast.SBoolLit(value) -> Llvm.const_int bool_t (if
193
       value then 1 else 0) (* bool_t is still an integer, must
        convert *)
          Semast. SIntLit(value, b) -> Llvm.const_of_int64 (Llvm
194
       .integer_type lu.lu_context b) value true (* bool for
       signedness *)
        | Semast. SStringLit (value) ->
195
          let str = Llvm.build global stringptr value "str lit"
196
        lu.lu_builder in
197
        | Semast.SFloatLit(value, b) -> Llvm.const_float f64_t
198
       value
      in
199
      (lu, v)
200
201
202
    let dump_temporary_value lu ev v =
      let v = match ( Llvm.classify_type ( Llvm.type_of v ) )
203
       with
          Llvm. TypeKind. Pointer ->
          if ( should_reference_pointer ev ) then
205
206
            \mathbf{V}
          else
207
            Llvm.build_load v "tmp" lu.lu_builder
          _ -> v
209
      in
210
      v
211
212
    let dump_s_expression_temporary_gen f lu e =
213
      let (lu, v) = (f lu e) in
214
      let v = dump_temporary_value lu ( Semast.
215
       type name of s expression e ) v in
      (lu, v)
216
217
    let dump_arguments_gen f lu el =
```

```
let acc_expr (lu, vl) e =
219
        let (lu, v) = dump s expression temporary gen f lu e in
220
        (lu, v :: vl)
221
222
      let (lu, args) = List.fold_left acc_expr (lu, []) el in
223
      (lu, args)
224
225
    let rec dump s expression lu e =
226
      match e with
227
        Semast. SLiteral (lit) -> dump_s_literal_lu_lit
228
        Semast.SQualifiedId(qid, t) ->
229
        let (lu, v) = dump_s_qualified_id lu qid t in
230
        (lu, v)
231
        Semast. SCall(e, el, t) ->
232
        let (lu, target) = dump_s_expression lu e in
        let oparghandler = find_argument_handler lu target in
234
        let (lu, args) = match oparghandler with
          \mid None \rightarrow
236
            let ( lu, args ) = ( dump\_arguments\_gen (
237
       dump_s_expression ) lu el ) in
238
            (lu, List.rev args)
          | Some(h) \rightarrow
239
            let (lu, args) = (h lu el) in
240
             (lu, args)
241
        in
242
        let arr args = Array.of list args in
243
244
        let v = match t with
          | Semast.SBuiltinType(Ast.Void, _) -> Llvm.build_call
245
        target arr_args "" lu.lu_builder
          _ -> Llvm.build_call target arr_args "tmp.call" lu.
246
       lu_builder
247
        in
        (lu, v)
248
        Semast. SBinaryOp(l, bop, r, t) \rightarrow
        let (lu, lv) = dump_s_expression lu l in
        let (lu, rv) = dump_s_expression lu r in
        let opf = match bop with
252
            Ast.Add -> Llvm.build add
253
            Ast.Sub -> Llvm.build_sub
254
            Ast. Mult -> Llvm. build mul
255
            Ast.Div -> Llvm.build_sdiv
256
            _ -> raise (Errors. Unsupported ("This binary
257
       operation type is not supported for code generation"))
258
        let \ v = opf \ lv \ rv \ "tmp.bop" \ lu.lu\_builder \ in
259
```

```
(lu, v)
260
      -> raise (Errors. Unsupported ("This expression is not
261
       supported for code generation"))
262
   let dump_s_expression_temporary lu e =
263
      let (lu, v) = (dump_s_expression lu e) in
264
      let v = dump_temporary_value lu ( Semast.
265
       type name of s expression e ) v in
      (lu, v)
266
267
   let dump_s_locals lu locals =
268
      let acc lu (n, tn) =
269
        let lty = llvm_type_of_s_type_name lu tn in
270
        let v = Llvm.build_alloca lty n lu.lu_builder in
       { lu with lu_named_values = StringMap.add (n) v lu.
       lu_named_values }
      let Semast. SLocals (bl) = locals in
274
      let lu = List.fold left acc lu bl in
275
     lu
277
   let dump_s_parameters lu llfunc parameters =
278
      let paramarr = Llvm.params llfunc in
279
      let paraml = Array.to_list paramarr in
280
      let Semast. SParameters (bl) = parameters in
281
      let nameparam i p =
        let (n, _) = (List.nth bl i) in
283
        ( Llvm.set_value_name n p )
284
285
      let _ = Array.iteri nameparam paramarr in
      let acc lu (n, tn) =
287
       let v = List.find ( fun p -> ( ( Llvm.value_name p ) =
       n ) ) paraml in
       { lu with lu_named_params = StringMap.add n v lu.
       lu_named_params }
290
      let lu = List.fold_left acc lu bl in
291
292
293
   let dump store lu lhs lhst rhs rhst =
294
      let _ = Llvm.build_store rhs lhs lu.lu_builder in
295
      lhs
296
297
   let dump assignment lu lhse rhse lhst
298
      let rhst = Semast.type_name_of_s_expression rhse in
```

```
let (lu, rhs) = dump_s_expression_temporary lu rhse in
300
      let (lu, lhs) = dump s expression lu lhse in
301
      let v = dump_store lu lhs lhst rhs rhst in
302
303
      (lu, v)
304
    let dump_s_variable_definition lu = function
305
        Semast.SVarBinding((n, tn), rhse) \rightarrow let lhse = Semast.
306
       SQualifiedId([n], tn) in
        let lhst = tn in
307
        let rhst = Semast.type_name_of_s_expression rhse in
308
        let (lu, rhs) = dump_s_expression_temporary lu rhse in
        let (lu, lhs) = dump_s_expression lu lhse in
        let v = dump_store lu lhs lhst rhs rhst in
        ( lu, v )
312
    let rec dump_s_general_statement lu gs =
314
      let acc lu bgs =
315
        dump_s_general_statement lu bgs
316
317
      match gs with
318
319
        Semast.SGeneralBlock(locals, gsl) ->
        let lu = dump_s_locals lu locals in
        let lu = List.fold_left acc lu gsl in
321
322
        Semast. SExpressionStatement (e) ->
323
        let (lu, _) = dump_s_expression lu e in
324
325
326
      | Semast.SVariableStatement(vdef) ->
        let (lu, _) = dump_s_variable_definition lu vdef in
327
        lu
330
    let rec dump_s_statement lu s =
      let acc lu s =
331
        dump_s_statement lu s
333
      match s with
334
        Semast.SBlock( locals, sl ) ->
335
        let lu = dump s locals lu locals in
336
        let lu = List.fold_left acc lu sl in
337
338
      | Semast.SGeneral(gs) ->
339
        dump s general statement lu gs
340
       Semast. SReturn (e) ->
341
        let lu = match e with
342
            Semast.SNoop ->
```

```
let _ = Llvm.build_ret_void lu.lu_builder in
344
345
            e -> let (lu, v) = dump_s_expression_temporary lu e
346
            let _ = Llvm.build_ret v lu.lu_builder in
347
348
        in
349
        lu
350
       -> raise (Errors. Unsupported ("This statement type is
351
       unsupported"))
352
   let dump_s_variable_definition_global lu = function
353
       Semast. SVarBinding ((n, tn), e) ->
        let (lu, rhs) = dump_s_expression lu e in
355
        let v = Llvm.define_global n rhs lu.lu_module in
        let lu = \{ lu with \}
357
          lu_variables = StringMap.add n v lu.lu_variables
        } in
359
        let lty = llvm_type_of_s_type_name lu v in
361
362
        let v = Llvm.declare_global lty k lu.lu_module in
        { lu with lu_variables = StringMap.add k v lu.
363
       lu_variables }
       let v = llvm_lookup_variable lu n tn in
364
        let = Llvm.set initializer v rhs in
365
        let lu = { lu with lu variables = StringMap.add n v lu.
366
       lu_variables } in
        *)
367
        ( lu, v)
368
   let dump_s_function_definition lu f =
371
      let acc lu s =
        dump_s_statement lu s
372
      in
      (* Generate the function with its signature *)
374
      (* Which means we just look it up in the llvm module *)
375
      let ft = Semast.type_name_of_s_function_definition f in
376
      let n = Semast.string_of_qualified_id f.Semast.func_name
377
       in
      let llfunc = llvm_lookup_function lu n ft in
378
      (* generate the body *)
379
      let entryblock = Llvm.append block lu.lu context "entry"
380
       llfunc in
         Llvm.position_at_end entryblock lu.lu_builder;
381
      let lu = dump_s_parameters lu llfunc f.Semast.
382
```

```
func_parameters in
      let lu = List.fold left acc lu f.Semast.func body in
383
      let lu = \{ lu with \}
384
        lu_named_params = StringMap.empty;
385
      } in
386
      lu
387
   let dump s basic definition lu = function
389
       Semast. SVariable Definition (v) -> let (lu, ) =
390
       dump s variable definition global lu v in
391
      | Semast. SFunctionDefinition(f) ->
392
       dump_s_function_definition_lu_f
393
   let dump_s_definition lu = function
394
        Semast.SBasic(b) -> dump_s_basic_definition lu b
395
396
   let dump_array_prelude lu =
397
      (* Unfortunately, unsupported ... *)
398
      lu
399
400
   let dump_parallelism_prelude lu =
401
      (* Unfortunately, unsupported ... *)
402
403
404
   let dump global string lu n v =
405
      let rhs = Llvm.const_stringz lu.lu_context v in
406
      let v = Llvm.define_global n rhs lu.lu_module in
407
      (lu, v)
408
   let dump_builtin_lib lu =
410
411
      let char_t = Llvm.i8_type
                                         lu.lu_context
      and i32 t
                  = Llvm.i32\_type
                                         lu.lu_context
412
      (* LLVM treats booleans as 1-bit integers, not distinct
       types with their own true / false *)
414
      let p_char_t = Llvm.pointer_type
                                              char t
415
      and llzero = Llvm.const_int i32_t 0
416
      in
417
      let f_acc lu (n, lv) =
418
        { lu with lu_functions = StringMap.add n lv lu.
       lu functions }
      in
420
      let print_lib lu =
421
        let printf_t = Llvm.var_arg_function_type i32_t []
```

```
p_char_t || in
       let printf func = Llvm.declare function "printf"
423
       printf_t lu.lu_module in
       let (_, int_format_str) = dump_global_string lu "__ifmt
424
       " "%d"
       and (_, str_format_str) = dump_global_string lu "__sfmt
425
       and (_, float_format_str) = dump_global_string lu "
426
        __ffmt" "%f"
       in
427
        let handler_name = "printf" in
        let handler lu el =
429
          let ( lu , exprl ) = ( dump_arguments_gen (
430
       dump_s_expression ) lu el ) in
          if (List.length el) < 1 then (lu, exprl) else
431
          let hdt = Semast.type_name_of_s_expression ( List.hd
432
       el ) in
          let insertion = match hdt with
433
            | Semast.SBuiltinType(Ast.String, ) ->
434
       str_format_str
            | Semast.SBuiltinType(Ast.Float(n), _) ->
435
       float_format_str
            | Semast.SBuiltinType(Ast.Int(n), _) ->
436
       int_format_str
            -> raise (Errors . BadPrintfArgument)
437
438
          let fptr = Llvm.build_gep insertion [ llzero; llzero
439
        ] "tmp.fmt" lu.lu_builder in
          (lu, fptr :: exprl)
440
        in
        let libprintfuncs = [
442
          (( Semast.mangle_name ["lib"; "print"] ( Semast.
       SFunction (Semast.void_t, [Semast.string_t], Semast.
       no qualifiers))), printf_func);
          (( Semast.mangle_name ["lib"; "print"] ( Semast.
444
       SFunction(Semast.void_t, [Semast.float64_t], Semast.
       no_qualifiers)) ), printf_func);
          (( Semast.mangle_name ["lib"; "print"] ( Semast.
445
       SFunction(Semast.void_t, [Semast.int32_t], Semast.
       no_qualifiers)) ), printf_func);
       in
446
       let lu = List.fold left f acc lu libprintfuncs in
447
       let lu = \{ lu with \}
448
          lu_handlers = ( StringMap.add handler_name ( handler
449
       ) lu.lu_handlers )
```

```
} in
450
        lu
451
      in
452
      let math_lib lu =
453
        lu
454
455
      in
      let lu = print_lib lu in
456
      let lu = math lib lu in
457
      lu
458
459
    let dump_builtin_module lu = function
460
        Semast.Lib -> dump_builtin_lib lu
461
462
        dump_module_import lu = function
463
        Semast.SBuiltin(lib) -> dump_builtin_module lu lib
        Semast.SCode(\_) \rightarrow lu
465
        Semast.SDynamic(_) -> lu
466
467
    let dump declarations lu =
468
      let rec declare k lu t = match t with
469
470
          Semast. SOverloads (fl) ->
           ( List.fold_left (declare k) lu fl )
          Semast. SFunction (rt, args, tq) as ft -> let lty =
472
       llvm_type_of_s_type_name_lu_t_in
          let mk = Semast.mangle name [k] ft in
473
          let v = Llvm.declare function mk lty lu.lu module in
          { lu with lu functions = StringMap.add mk v lu.
475
       lu_functions }
         _ -> lu
476
      in
      let acc_def k t lu =
478
        declare k lu t
479
480
      let toplevel = lu.lu_env.Semast.env_definitions in
      let lu = StringMap.fold acc_def toplevel lu in
482
      lu
483
    let dump_prelude lu sprog =
485
      let lu = dump_array_prelude lu in
486
      let lu = dump_parallelism_prelude lu in
487
      let lu = List.fold_left dump_module_import lu lu.lu_env.
       Semast.env imports in
      let lu = dump declarations lu in
489
      lu
490
491
```

```
let generate sprog =
492
      let acc_def lu d =
493
        let lu = dump_s_definition lu d in
494
        lu
495
      in
496
      let lu = create_li_universe sprog in
497
      let lu = dump_prelude lu sprog in
498
      let lu = match sprog with
499
        | Semast.SProgram(_, _, defs) -> ( List.fold_left
500
       acc_def lu defs )
      in
501
      {\tt lu.lu\_module}
502
                            ../source/codegen.ml
```