Spectral Collocation Methods for Differential-Algebraic Equations with Arbitrary Index

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Abstract In this paper, a symmetric Jacobi–Gauss collocation scheme is explored for both linear and nonlinear differential-algebraic equations (DAEs) of arbitrary index. After standard index reduction techniques, a type of Jacobi–Gauss collocation scheme with N knots is applied to differential part whereas another type of Jacobi–Gauss collocation scheme with N+1 knots is applied to algebraic part of the equation. Convergence analysis for linear DAEs is performed based upon Lebesgue constant of Lagrange interpolation and orthogonal approximation. In particular, the scheme for nonlinear DAEs can be applied to Hamiltonian systems. Numerical results are performed to demonstrate the effectiveness of the proposed method.

 $\textbf{Keywords} \quad \text{Spectral collocation method} \cdot \text{Differential-algebraic equation} \cdot \text{Hamiltonian} \\ \text{systems}$

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1 Introduction

In this paper, we consider the numerical solution of linear differential-algebraic equations (DAEs) with form

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$$\begin{cases} E(t)\dot{x}(t) = A(t)x(t) + f(t) & \text{fall all } t \in \mathbb{I}, \\ Bx(t) + Dx(\bar{t}) = r, \end{cases}$$
 (1)

where $\mathbb{I} = [\underline{t}, \overline{t}], E, A : \mathbb{I} \to R^{n \times n}$ and $f : \mathbb{I} \to R^n$ are sufficiently smooth and $B, D \in R^{d \times n}, d \leq n$ is the number of differential equations. We also consider nonlinear DAEs of the form

$$\begin{cases}
F(t, x, \dot{x}) = 0, \\
f(x(t), x(\bar{t})) = 0,
\end{cases}$$
(2)

where F and f are smooth functions. We assume that these equations possesses a unique solution x^* .

Collocation methods for linear DAEs has been extensively studied, see Radau collocation [15], projected piecewise polynomial collocation [1], perturbed collocation [3] and symmetric collocation [9]. Collocation methods for general nonlinear DAE is explored in [10]. In [9–11], by index reduction technique, one can separate algebraic equations from differential ones. Furthermore, an h-version Legendre–Gauss scheme for differential part with N knots and h-version Legendre–Gauss–Lobatto scheme for algebraic part with N+1 knots are applied to obtain a symmetric method which guarantees consistent approximations at the mesh points. In this work, however, we apply a p-version Jacobi–Gauss scheme (for differential part) and Jacobi–Gauss–Lobatto scheme (for algebraic part) to solve linear or nonlinear DAEs. Based upon the scheme, a convergence result is established for linear DAEs under standard assumptions and is confirmed by our numerical examples. Furthermore, the scheme for nonlinear DAEs can be applied for numerical solution of Hamiltonian systems. We compare phase plot of some Hamiltonian systems as DAEs with those in [7], in which a spectral collocation method is designed for Hamiltonian systems as ODEs.

This paper is organized as follows: In Sect. 2, some preliminary knowledge on the theory of DAEs and Jacobi polynomials are given. In Sect. 3, we formulate our algorithm for linear DAEs and provide a convergence analysis. Two numerical experiments are also included in this section. Sect. 4 is devoted to the algorithm and numerical experiments for nonlinear DAEs. We then give some conclusions in Sect. 5.

Throughout the paper, C stands for a generic constant that is independent of the number of collocation points N but may depend on the length of time interval $\bar{t} - \underline{t}$ and the dimension n.

2 Preliminaries

Application of index reduction techniques on the linear DAEs (1) and nonlinear DAEs (2), one obtain

$$\hat{E}(t)\dot{x}(t) = \hat{A}(t)x(t) + \hat{f}(t), \quad t \in [\underline{t}, \overline{t}]$$
(3)

with

$$\hat{E} = [\hat{E}_1, 0]^T, \ \hat{A} = [\hat{A}_1, \hat{A}_2]^T, \ \hat{f} = [\hat{f}_1, \hat{f}_2]^T.$$
 (4)

and

$$\begin{cases} \hat{F}_1(t, x, \dot{x}) = 0, \\ \hat{F}_2(t, x) = 0, \quad t \in [t, \bar{t}]. \end{cases}$$
 (5)



Since solutions of (1) and (3), (2) and (5) are equivalent respectively, we assume that both the linear and nonlinear DAE that we work on is of the index reduced form and without causing any confusion, we keep the same notation as those in (1) and (2). We denote a and d the size of algebraic and differential part of (3) and (5). Interested readers are referred to [9,10] for more details. For the linear DAEs, we have

Lemma 2.1 ([9]) For $E, A \in C^k(\mathbb{I}, R^{n \times n})$, there exist point-wise nonsingular $P \in C^{k-1}(\mathbb{I}, R^{n \times n})$, $Q \in C^k(\mathbb{I}, R^{n \times n})$ such that

$$PEQ = \begin{bmatrix} I_d & 0 \\ 0 & 0 \end{bmatrix}, PAQ - PE\dot{Q} = \begin{bmatrix} 0 & 0 \\ 0 & I_a \end{bmatrix}.$$
 (6)

In particular, P has the special structure

$$P = \begin{bmatrix} P_{11} & P_{12} \\ 0 & P_{22} \end{bmatrix}, \quad with \quad P_{11}(t) \in \mathbb{R}^{d \times d}, P_{12} \in \mathbb{R}^{d \times a}, P_{22}(t) \in \mathbb{R}^{a \times a}. \tag{7}$$

In addition, if $f \in C^{k-1}(\mathbb{I}, \mathbb{R}^n)$, then $x \in C^{k-1}(\mathbb{I}, \mathbb{R}^n)$ for every solution x of (1).

Next, we introduce Jacobi polynomials $P_n^{(\alpha,\beta)}(x)$, which plays a center role throughout the paper. It is well-known that $P_n^{(\alpha,\beta)}(x)$ is a class of orthogonal polynomials associate with weight $\omega(x) = (1-x)^{\alpha}(1+x)^{\beta}$, $\alpha, \beta > -1$. Under the normalization $P_k^{(\alpha,\beta)}(1) = {k+\alpha \choose k}$, one has the expression,

$$P_k^{(\alpha,\beta)}(x) = \frac{1}{2^k} \sum_{l=0}^k {k+\alpha \choose k-l} (k+\beta l) (x-1)^l (x+1)^{k-l}.$$

Jacobi polynomials satisfy the three-term recursive relations:

$$\begin{split} P_0^{(\alpha,\beta)}(x) &= 1, P_1^{(\alpha,\beta)}(x) = \frac{1}{2} [(\alpha - \beta) + (\alpha + \beta + 2)x], \\ a_{1,k} P_{k+1}^{(\alpha,\beta)}(x) &= a_{2,k} P_k^{(\alpha,\beta)}(x) - a_{3,k} P_{k-1}^{(\alpha,\beta)}(x), \end{split}$$

where

$$a_{1,k} = 2(k+1)(k+\alpha+\beta+1)(2k+\alpha+\beta),$$

$$a_{2,k} = (2k+\alpha+\beta+1)(\alpha^2-\beta^2) + x\Gamma(2k+\alpha+\beta+3)/\Gamma(2k+\alpha+\beta),$$

$$a_{3,k} = 2(k+\alpha)(k+\beta)(2k+\alpha+\beta+2).$$

Especially, Legendre polynomial is exactly a class of Jacobi polynomial with $\alpha = \beta = 0$ and Chebyshev polynomials of the first kind is another class with $\alpha = \beta = -1/2$ up to a constant. It is clear that $P_n^{(\alpha,\beta)}(x)$ forms a complete orthogonal system associate with the norm

$$\|v\|_{\omega^{\alpha,\beta}} = \left(\int_{-1}^{1} |v(x)|^2 \omega^{\alpha,\beta}(x) dx\right)^{1/2}.$$



Denote $\omega(x) = \omega^{-1/2, -1/2}(x)$ the Chebyshev weight function and introduce a semi-norm for the simplicity of analysis

$$|v|_{H^{m,N}_{\omega}(-1,1)} = \left(\sum_{k=\min(m,N+1)}^{m} |\partial_x^k v|_{L^2_{\omega}(-1,1)}^2\right)^{1/2}.$$

Let I_N^c denote the interpolation operator on Chebyshev points, then we have [4]

$$||u - I_N^c u||_{\infty} \le C N^{1/2 - m} |u|_{H_{\alpha}^{m,N}(-1,1)}.$$
 (8)

Let $h(t) \in C^k(\mathbb{I}, R^n)$ and write $u(x) = h\left[\frac{\overline{t}-t}{2}(1+x) + \frac{\overline{t}+t}{2}\right]$, $x \in [-1, 1]$. Let $P_N^{(\alpha, \beta)}$ be the Jacobi polynomial with degree N and $I_N u \in P_N[-1, 1]$ interpolate u at N+1 Jacobi–Gauss points, Jacobi–Gauss–Radau points or Jacobi–Gauss–Lobatto points, then we have the following lemma on the Legesgue constant, see [12].

Lemma 2.2 Let $\{l_j(x)\}_{j=0}^N$ be the Lagrange interpolation polynomials at Jacobi–Gauss points, then

$$||I_N^{\alpha,\beta}||_{\infty} := \max_{x \in [-1,1]} \sum_{j=0}^{N} |I_j(x)|$$

$$= \begin{cases} \mathcal{O}(\log N), & -1 < \alpha, \beta \le -\frac{1}{2}, \\ \mathcal{O}(N^{\gamma+1/2}), & \gamma = \max(\alpha, \beta), \end{cases} \text{ otherwise.}$$

$$(9)$$

The remainder of the interpolation is

$$u(x) - I_N u(x) = u[x_0, x_1, \dots, x_N, x] v(x),$$

where $v(x) = (x - x_0)(x - x_1) \dots (x - x_N)$ and $u[x_0, x_1, \dots, x_N, x]$ is the divided difference of u. Hence,

$$u(x) - I_N u(x) = \frac{u[x_0, x_1, \dots, x_N, x]}{L_{N+1}} P_{N+1}^{(\alpha, \beta)}(x),$$

where $L_{N+1} = \frac{\Gamma(2N+\alpha+\beta+3)}{2^{N+1}\Gamma(\alpha+\beta+N+2)(N+1)!} \approx \frac{2^{\alpha+\beta}}{\sqrt{\pi N}} 2^{N+1}$ is the leading coefficient of $P_{N+1}^{(\alpha,\beta)}(x)$. Moreover, if u is sufficiently smooth, then the divided difference

$$u[x_0, x_1, \dots, x_N, x] = \frac{u^{(N+1)}(\xi_x)}{(N+1)!}, \quad \xi_x \in (-1, 1),$$

which implies

$$||u(x) - I_N u(x)||_{L^{\infty}} \le C \frac{\sqrt{N+1}}{2^{N+1}(N+1)!} ||u^{(N+1)}(\xi_x)||_{L^{\infty}}.$$
 (10)

For the sake of analysis, we define the Lagrange interpolation at (N + 1) interpolation points, i.e.,

$$I_N u(x_i) = \sum_{i=0}^{N} u(x_i) l_i(x),$$
 (11)

where $l_i(x)$ is the Lagrange interpolation basis function at these points.

It is obvious that for all $t \in [t, \bar{t}]$ and $x \in [-1, 1]$

$$h(t) = u(x), h_N(t) = u_N(x),$$



Therefore,

$$(h - h_N)(t) = (u - u_N)(x) := e(x).$$
(12)

3 Linear Differential-Algebraic Equations

This section is devoted to the symmetric Jacobi collocation methods for linear DAEs and its convergence analysis.

Firstly, we transform the problem (3) onto an equivalent problem on interval [-1, 1],

$$\frac{2}{b-a}E(\xi)\dot{x}(\xi) = A_1(\xi)x(\xi) + f_1(\xi)$$

$$0 = A_2(\xi)x(\xi) + f_2(\xi), \ \xi \in [-1, 1].$$
(13)

Here, for simplicity of notation, the hats are omitted and subindex of E is suppressed.

3.1 Formulation of Algorithm

On the interval [-1, 1], let ρ_j 's be the zeros of Jacobi polynomial $P_N^{(\alpha, \beta)}(x)$, and σ_i 's be the associated Jacobi–Gauss–Lobatto points.

$$-1 \le \rho_0 < \dots < \rho_{N-1} \le 1, \ -1 = \sigma_0 < \dots < \sigma_N = 1.$$
 (14)

Then, we have the interlacing property [6,13],

$$-1 = \sigma_0 < \rho_0 < \sigma_1 < \dots < \rho_{N-1} < \sigma_N = 1. \tag{15}$$

In order to fix the idea, we use T_k , the Chebyshev polynomial of the first kind in algorithms. Nevertheless, our method is applicable to all Jacobi polynomials. Approximate the true solution of (13) by

$$x_N^m(\xi) = \sum_{k=0}^N C_k^m T_k(\xi), \ m = 1, \dots, n.$$
 (16)

The Jacobi coefficients $C_k^m \in R^{(N+1)\times 1}$ is determined by the following set of conditions:

$$r = Bx_N(-1) + Dx_N(1),$$

$$\frac{2}{b-a}E(\rho_i)\dot{x}_N(\rho_i) = A_1(\rho_i)x_N(\rho_i) + f_1(\rho_i), \quad i = 0, \dots, N-1,$$

$$0 = A_2(\sigma_i)x_N(\sigma_i) + f_2(\sigma_i), \quad j = 0, \dots, N.$$
(17)

Note that the above system has exactly (N+1)n unknowns and d+Nd+(N+1)a = (N+1)n equations.

By the fact that $T'_k(\xi) = kU_{k-1}(\xi)$, where $U_k(\xi)$ is the Chebyshev polynomial of the second kind and denoting for all i = 1, ..., n, j = 1, ..., a and k = 1, ..., d

$$W_{1} = \begin{bmatrix} 1, -1, \dots, (-1)^{N} \end{bmatrix}, \quad W_{2} = \begin{bmatrix} 1, 1, \dots, 1^{N} \end{bmatrix},$$

$$f_{1,k} = \begin{bmatrix} f_{1,k}(\rho_{0}), \dots, f_{1,k}(\rho_{N-1}) \end{bmatrix}, \quad f_{2,j} = \begin{bmatrix} f_{2,j}(\sigma_{0}), \dots, f_{2,j}(\sigma_{N}) \end{bmatrix},$$

$$f_{1} = \begin{bmatrix} f_{1,1}, \dots, f_{1,d} \end{bmatrix}^{T}, \quad f_{2} = \begin{bmatrix} f_{2,1}, \dots, f_{2,a} \end{bmatrix}^{T}.$$



$$U = \begin{bmatrix} 0 & 1 & \cdots & NU_{N-1}(\rho_0) \\ 0 & 1 & \cdots & NU_{N-1}(\rho_1) \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 1 & \cdots & NU_{N-1}(\rho_{N-1}) \end{bmatrix},$$

$$T = \begin{bmatrix} 1 & \rho_0 & \cdots & T_N(\rho_0) \\ 1 & \rho_1 & \cdots & T_N(\rho_1) \\ \vdots & \vdots & \ddots & \vdots \\ 1 & \rho_{N-1} & \cdots & T_N(\rho_{N-1}) \end{bmatrix}, \quad \tilde{T} = \begin{bmatrix} 1 & \sigma_0 & \cdots & T_N(\sigma_0) \\ 1 & \sigma_1 & \cdots & T_N(\sigma_1) \\ \vdots & \vdots & \ddots & \vdots \\ 1 & \sigma_N & \cdots & T_N(\sigma_N) \end{bmatrix}, \quad (18)$$

we have a system AC = F, where

$$C = \begin{bmatrix} C_{1}^{1}, \dots, C_{N}^{1}, \dots, C_{1}^{n}, \dots, C_{N}^{n} \end{bmatrix}^{T},$$

$$A = \begin{bmatrix} \frac{B \otimes W_{1} + D \otimes W_{2}}{\frac{2}{b-a}E_{11}U - A_{1,11}T & \cdots & \frac{2}{b-a}E_{1n}U - A_{1,1n}T} \\ \vdots & \ddots & \vdots \\ \frac{\frac{2}{b-a}E_{d1}U - A_{1,d1}T & \cdots & \frac{2}{b-a}E_{dn}U - A_{1,dn}T}{-A_{2,11}\tilde{T} & \cdots & -A_{2,1n}\tilde{T}} \end{bmatrix}, F = \begin{bmatrix} \frac{r}{f_{1}} \\ f_{1} \\ \vdots \\ f_{2} \end{bmatrix}, (19)$$

$$-A_{2,a1}\tilde{T} & \cdots & -A_{2,an}\tilde{T} \end{bmatrix}$$

and \bigotimes denotes the Kronecker product. Once C is determined, we plug it back to (16) to obtain x_N .

3.2 Convergence Results

We try to derive an error estimate in the L^{∞} -norm of the collocation scheme (17).

Since N-point Jacobi–Gauss–Lobatto points with parameters (α, β) are zeros of $P_{N-2}^{(\alpha+1,\beta+1)}(t)$ (± 1 excluded) [14], throughout the paper we only need to consider the case when both ρ_i 's and σ_j 's are zeros of certain type of Jacobi polynomials. Moreover, it is well-known that both ρ_j 's and σ_j 's are dense in [-1,1] as $N\to\infty$. Note that the distance between any two consecutive zeros of any type of Jacobi polynomials is bound by $\frac{\sqrt{1-t^2}}{N}$, $t\in[-1,1]$ up to a constant independent of N and t [8]. We derive that $\rho_i-\sigma_i=\mathbb{O}(N^{-1})$, $i=1,\ldots,N-1$.

To obtain the solvability of our algorithm, let us rewrite our algorithm in its equivalent form. For each component of x^m , its approximation

$$x_N^m = \sum_{k=0}^N C_k^m T_k(\xi) = \sum_{k=0}^N x_{k,m}^N l_k(\xi),$$
 (20)

where $l_k(\xi)$ is the Lagrange interpolation polynomial on point set σ_j 's. Denote $v_{ki} = l'_i(\rho_k)$, $u_{ki} = l_i(\rho_k)$, $E_i = E(\rho_i)$, $A_{1i} = A_1(\rho_i)$, $A_{2j} = A_2(\sigma_j)$, and

$$x_N = \left[x_{0,1}^N, x_{0,2}^N, \dots, x_{0,n}^N, \dots, x_{N,1}^N, \dots, x_{N,n}^N \right]^T$$



for k = 1, ..., N, i = 1, ..., N. Plugging (20) into (13) yields a system $A'x_N = F'$, where

 $F' = [f_1(\rho_0) \ f_2(\sigma_0) \dots f_1(\rho_{N-1}) \ f_2(\sigma_{N-1}) \ r \ f_2(1)]^T$. Here, a missing condition $-A_2(1)x_N(1) = f_2(1)$ is used. Hence, if A' is regular, our collocation algorithm is stable. To reorder rows and columns, we define

Remark 1 As is well-known, Lagrange interpolation is not stable for large *N*. Moreover, it requires a large computational cost to find the base functions. Hence, the equivalent algorithm is exhibited only for a theoretical purpose.

Lemma 3.1 For N sufficiently large,

$$l_{j}(\rho_{i}) = \begin{cases} \mathbb{O}(N^{-1}), & i \neq j \text{ and } i \neq j-1, \\ \mathbb{O}(1), & \text{otherwise.} \end{cases}$$

Proof Let δ be a fixed number such that $\delta < 1 - |\rho_i|$. By the interlacing property, we can choose δ so small that $B(\sigma_i, \delta) \cap \rho = {\rho_{i-1}, \rho_i}$. From [15] (pp. 336), if $\sigma_i - \rho_i > \delta$,

$$l_{j}(\rho_{i}) = \mathbb{O}(N^{-1/2})|P_{N+1}^{(\alpha,\beta)}(\sigma_{j})|^{-1},$$

$$= \mathbb{O}(N^{-1/2})(j^{\alpha+3/2}N^{-\alpha-2} + j^{\beta+3/2}N^{-\beta-2}),$$

$$\leq \mathbb{O}(N^{-1}),$$
(23)

otherwise,

$$l_j(\rho_i) = \mathbb{O}\left(N^{-1/2}\right) \frac{P_{N+1}^{(\alpha,\beta)}(\sigma_j) - P_{N+1}^{(\alpha,\beta)}(\rho_i)}{\sigma_j - \rho_i} = \mathbb{O}(1).$$
 (24)

As an illustration, we plot an $l_{12}(x)$ for N=25 and $(\alpha, \beta)=(1/2, 2/3)$, see Fig. 1. It is clear that $l_{12}(\rho_i)$ has the largest modulus at ρ_{11} and ρ_{12} .

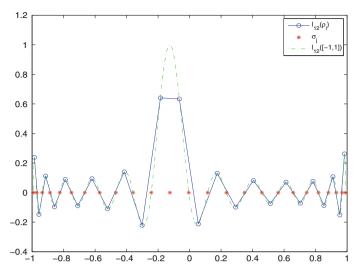


Fig. 1 A plot of $l_{12}(x)$ for N = 25 and $(\alpha, \beta) = (1/2, 2/3)$

Next, we prove the invertibility of A' by multiplying it from the left and from the right, respectively, by

$$T_{P} = \operatorname{diag}\left(\operatorname{diag}\left(\begin{array}{cc} P_{11}(\rho_{i}) & P_{12}(\sigma_{i}) \\ 0 & P_{22}(\sigma_{i}) \end{array}\right)_{i=0,\dots,N-1}, \begin{bmatrix} I & 0 \\ 0 & P_{22}(1) \end{bmatrix}\right),$$

$$T_{Q} = \operatorname{diag}\left(Q(\sigma_{0}), \dots, Q(\sigma_{N-1}), Q(1)\right). \tag{25}$$

Lemma 3.2 Assume A_2 , $Q \in C^2$. Then A' is invertible for sufficiently large N.

Proof Since $l_i(x) = \frac{P_{N+1}(x)}{P'_{N+1}(\sigma_i)(x-\sigma_i)}$, after a simple calculation, we have

$$v_{ij} = l_i'(\rho_j) = -\frac{l_i(\rho_j)}{|\rho_j - \sigma_i|}.$$
 (26)

Clearly, from Lemma (3.1),

$$v_{ij} = \begin{cases} \mathbb{O}(N), \ j = i, \text{ or } i - 1; \\ \mathbb{O}(1), \ j \neq i, \text{ or } i - 1 \text{ and } |\sigma_j - \rho_i| = \mathbb{O}(N^{-1}); \\ o(1), \text{ otherwise.} \end{cases}$$
 (27)

Hence, for $j \neq i$ or i - 1,

$$I_{1} = [P_{11}(\rho_{i}) \quad P_{12}(\sigma_{i})] \begin{bmatrix} \frac{2v_{ij}}{b-a} E_{i} - u_{ij} A_{1i} \\ 0 \end{bmatrix} Q(\sigma_{j})$$

$$= \frac{2v_{ij}}{b-a} (P_{11}E)(\rho_{i}) Q(\sigma_{j}) - u_{ij} (P_{11}A_{1})(\rho_{i}) Q(\sigma_{j})$$

$$= \frac{2v_{ij}}{b-a} (P_{11}EQ)(\rho_{i}) - u_{ij} (P_{11}A_{1})(\rho_{i}) Q(\sigma_{j}) + e_{1},$$
(28)



where $e_1 = \frac{2v_{ij}}{b-a}(P_{11}E)(\rho_i)(Q(\sigma_j) - Q(\rho_i))$. Apply the smoothness assumption of P, E and Q,

$$e_{1} = \frac{2l'_{i}(\rho_{j})}{b-a}(P_{11}E)(\rho_{i})\dot{Q}(\xi)(\rho_{j}-\sigma_{i}),$$

$$= -l_{i}(\rho_{j})(P_{11}E)(\rho_{i})\dot{Q}(\xi),$$

$$= \mathbb{O}(1)l_{i}(\rho_{j}). \tag{29}$$

However, in this case, $l_i(\rho_j)/l_i'(\rho_j) \approx 0$ by the L'Hopital's rule. Therefore, (29) and the definition of u_{ij} yield

$$I_1 = \frac{2v_{ij}}{b - a}[I\ 0] + o(v_{ij}). \tag{30}$$

Similarly, if j = i - 1, $I_1 = \frac{2v_{ij}}{b-a}[I \ 0] + \mathbb{O}(1)$. If j = i,

$$I_{2} = [P_{11}(\rho_{i}) \quad P_{12}(\sigma_{i})] \begin{bmatrix} \frac{2v_{ij}}{b-a} E_{i} - u_{ij} A_{1i} \\ -A_{2i} \end{bmatrix} Q(\sigma_{i})$$

$$= \frac{2v_{ii}}{b-a} (P_{11}E)(\rho_{i}) Q(\sigma_{i}) - u_{ii} (P_{11}A_{1})(\rho_{i}) Q(\sigma_{i}) - (P_{12}A_{2}Q)(\sigma_{i})$$

$$= \frac{2v_{ii}}{b-a} (P_{11}EQ)(\rho_{i}) - u_{ii} (P_{11}A_{1})(\rho_{i}) Q(\sigma_{i}) - (P_{12}A_{2}Q)(\sigma_{i}) + e_{2},$$

$$= \frac{2v_{ii}}{b-a} [I \ 0] + \mathbb{O}(1). \tag{31}$$

At last,

$$\begin{bmatrix} I & 0 \\ 0 & P_{22}(1) \end{bmatrix} \begin{bmatrix} B \\ 0 \end{bmatrix} Q(1) = \begin{bmatrix} B_{11} & B_{12} \\ 0 & 0 \end{bmatrix},$$
$$\begin{bmatrix} I & 0 \\ 0 & P_{22}(1) \end{bmatrix} \begin{bmatrix} D \\ -A_2(1) \end{bmatrix} Q(1) = \begin{bmatrix} D_{11} & D_{12} \\ 0 & -I \end{bmatrix}.$$

Hence, we have

$$M = T_{P}A'T_{Q} = \begin{bmatrix} \frac{2v_{00}}{b-a}[I_{d\times d}\ 0] + \mathbb{O}(1) & \cdots & \frac{2v_{0,N-1}}{b-a}[I_{d\times d}\ 0] + \mathbb{O}(\frac{1}{N}) & \frac{2v_{0N}}{b-a}[I_{d\times d}\ 0] + \mathbb{O}(\frac{1}{N}) \\ -I_{a\times a} & \cdots & 0 & 0 \\ & \ddots & & \ddots & \\ \frac{2v_{N-1,0}}{b-a}[I_{d\times d}\ 0] + \mathbb{O}(\frac{1}{N}) & \cdots & \frac{2v_{N-1,N-1}}{b-a}[I_{d\times d}\ 0] + \mathbb{O}(1) & \frac{2v_{N-1,N}}{b-a}[I_{d\times d}\ 0] + \mathbb{O}(1) \\ 0 & \cdots & -I_{a\times a} & 0 \\ \hline B_{11} & B_{12} & D_{11} & D_{12} \\ 0 & 0 & & D_{1n} & D_{12} \\ 0 & 0 & & 0 \end{bmatrix}$$

$$(32)$$

Rewrite the above matrix as

$$M = \left[\frac{\tilde{A} \, \big| \, \tilde{B}}{\tilde{C} \, \big| \, \tilde{D}} \right].$$

Clearly, if D_{11}^{-1} exists,

$$\tilde{D}^{-1} = \begin{bmatrix} D_{11}^{-1} & D_{11}^{-1} D_{12} \\ 0 & -I \end{bmatrix}.$$

Note that

$$\left[\begin{array}{cc} \tilde{A} & \tilde{B} \\ \tilde{C} & \tilde{D} \end{array} \right] \left[\begin{array}{cc} I & 0 \\ -\tilde{D}^{-1}\tilde{C} & I \end{array} \right] = \left[\begin{array}{cc} \tilde{A} - \tilde{B}\tilde{D}^{-1}\tilde{C} & \tilde{B} \\ 0 & \tilde{D} \end{array} \right].$$

Therefore, the invertibility of $\tilde{A} - \tilde{B}\tilde{D}^{-1}\tilde{C}$ implies the invertibility of A'. By a simple matrix manipulation,

$$\tilde{A} - \tilde{B}\tilde{D}^{-1}\tilde{C} = \begin{bmatrix} \frac{2v_{00}}{b-a}[I_{d\times d}\ 0] - \frac{2v_{0N}}{b-a}D_{11}^{-1}B_{11} + \mathbb{O}(1) & \cdots & \frac{2v_{0,N-1}}{b-a}[I_{d\times d}\ 0] + \mathbb{O}(\frac{1}{N}) \\ -I_{a\times a} & \cdots & 0 \\ & \ddots & & \ddots \\ \frac{2v_{N-1,0}}{b-a}[I_{d\times d}\ 0] - \frac{2v_{N-1,N}}{b-a}D_{11}^{-1}B_{11} + \mathbb{O}(1) & \cdots & \frac{2v_{N-1,N-1}}{b-a}[I_{d\times d}\ 0] + \mathbb{O}(1) \\ 0 & \cdots & -I_{a\times a} \end{bmatrix}$$

It is clear that the matrix is diagonal dominant, hence invertible.

This result indicates that our algorithm is well-posed. In the following, we find the error estimates for the algorithm.

Theorem 3.3 Let $x(\xi)$ be the true solution of transformed DAE (13) and $x_N(\xi)$ be the approximation given by (17) based upon our collocation method. If the given data $E(\xi)$, $A(\xi) \in C^{m+1}([-1, 1], R^{n \times n})$, and $f \in C^m([-1, 1], R^n)$, then for sufficiently large N,

$$\|x - x_N\|_{\infty} \le C \begin{cases} (\log N) N^{3/2 - m}, -1 < \alpha, \beta < -1/2, \\ N^{\gamma + 2 - m}, \quad \gamma = \max(\alpha, \beta), \text{ otherwise.} \end{cases}$$
 (33)

Proof Let I_N be the interpolation operator on $\sigma_0, \ldots, \sigma_N$ and \tilde{I}_{N-1} be the one on $\rho_0, \ldots, \rho_{N-1}$. Clearly, the true solution satisfies

$$r = Bx(-1) + Dx(1),$$

$$\frac{2}{b-a}E(\rho_i)\dot{x}(\rho_i) = A_1(\rho_i)x(\rho_i) + f_1(\rho_i), \quad i = 0, \dots, N-1,$$

$$0 = A_2(\sigma_i)x(\sigma_i) + f_2(\sigma_i), \quad j = 0, \dots, N.$$
(34)

Multiplying both sides of the first and the third equation of (34) by $l_j(\xi)$, the second equation by $\tilde{l}_i(\xi)$, respectively, and summing over i, j yield

$$r = BI_N x(-1) + DI_N x(1),$$

$$\frac{2}{b-a} \tilde{I}_{N-1}(E\dot{x}) = \tilde{I}_{N-1}(A_1 x) + \tilde{I}_{N-1} f_1,$$

$$0 = I_N(A_2 x) + I_N f_2.$$
(35)

We follow the same fashion for (17) to obtain a system for x_N , then subtract the system obtained from (35),

$$0 = BI_N(x(-1) - x_N(-1)) + DI_N(x(1) - x_N(1)),$$

$$0 = \frac{2}{b - a} \tilde{I}_{N-1}(E_1 \dot{x} - E \dot{x}_N) - \tilde{I}_{N-1}(A_1 x - A_1 x_N),$$

$$0 = I_N(A_2 x - A_2 x_N).$$
(36)



Recall the definition of I_{N}^{c} . Since $I_{N}A_{2}$ and $I_{N}x$ are polynomials, We have the following estimate from (5.5.28) of [4],

$$||I_{N}(A_{2}x) - (I_{N}A_{2})(I_{N}x)||_{\infty} = ||I_{N}^{c}(A_{2}x) - (I_{N}^{c}A_{2})(I_{N}^{c}x)||_{\infty}$$

$$= ||I_{N}^{c}((I_{N}^{c}A_{2})(I_{N}^{c}x)) - (I_{N}^{c}A_{2})(I_{N}^{c}x)||_{\infty}$$

$$\leq CN^{1/2 - 2m}|I_{N}^{c}A_{2} \cdot I_{N}^{c}x|_{H_{\infty}^{m,N}(-1,1)}.$$
(37)

Similar estimation can be obtained for $||I_N(A_2x_N) - (I_NA_2)(I_Nx_N)||_{\infty}$, $||\tilde{I}_{N-1}(E\dot{x}) - (\tilde{I}_{N-1}E)(\tilde{I}_{N-1}\dot{x})||_{\infty}$, $||\tilde{I}_{N-1}(E\dot{x}_N) - (\tilde{I}_{N-1}E)(\tilde{I}_{N-1}x_N)||_{\infty}$, $||\tilde{I}_{N-1}(A_1x) - (\tilde{I}_{N-1}A_1)(\tilde{I}_{N-1}x_N)||_{\infty}$. However, all these terms are high-order terms. Combine all these estimates, (36) can be written as

$$\begin{cases}
0 = BI_N(x(-1) - x_N(-1)) + DI_N(x(1) - x_N(1)), \\
\mathcal{O}(N^{1/2 - 2m}) = \frac{2}{b - a} (I_{N-1} E_1) [\tilde{I}_{N-1} \dot{x} - \tilde{I}_{N-1} \dot{x}_N] - (\tilde{I}_{N-1} A_1) [\tilde{I}_{N-1} x - \tilde{I}_{N-1} x_N], \\
\mathcal{O}(N^{1/2 - 2m}) = (I_N A_2) (I_N x - I_N x_N).
\end{cases}$$
(38)

Also, by simple calculation,

$$\|\tilde{I}_{N-1}\dot{x} - \tilde{I}_{N-1}\dot{x}_N - (\hat{I}_Nx) + (\hat{I}_Nx_N)\|_{\infty} = \|\tilde{I}_{N-1}\dot{x} - (\hat{I}_Nx)\|_{\infty}$$

$$< CN^{3/2-m}.$$
(39)

We obtain

$$\begin{cases} 0 = BI_N(x(-1) - x_N(-1)) + DI_N(\dot{x}(1) - x_N(1)), \\ \mathcal{O}(N^{3/2-m} || I_N ||_{\infty}) = \frac{2}{b-a} (I_N E_1) [(\dot{I_N} \dot{x}) - \dot{x}_N] - (I_N A_1) [I_N x - x_N], \\ \mathcal{O}(N^{1/2-2m}) = (I_N A_2) (I_N x - x_N). \end{cases}$$

Writing the right hand side of the last two groups of equations as w and applying the transformation in Lemma 2.1 yields that

$$\begin{bmatrix} I_d & 0 \\ 0 & 0 \end{bmatrix} \begin{bmatrix} \dot{y}_1 \\ \dot{y}_2 \end{bmatrix} = \begin{bmatrix} 0 & 0 \\ 0 & I_a \end{bmatrix} \begin{bmatrix} y_1 \\ y_2 \end{bmatrix} + \begin{bmatrix} g_1 \\ g_2 \end{bmatrix}$$
(40)

$$[B_1 \ B_2] \begin{bmatrix} y_1(-1) \\ y_2(-1) \end{bmatrix} + [D_1 \ D_2] \begin{bmatrix} y_1(1) \\ y_2(1) \end{bmatrix} = 0, \tag{41}$$

where

$$I_N x - x_N = Q \begin{bmatrix} y_1 \\ y_2 \end{bmatrix}, \quad Pw = \begin{bmatrix} g_1 \\ g_2 \end{bmatrix}, \quad BQ(-1) = [B_1 \ B_2], \quad DQ(1) = [D_1, \ D_2].$$

Matrix functions P and Q are given in Lemma 2.1. The differences $P(E - I_N E)Q = \mathcal{O}(N^{-m-1})$ and $P(A - I_N A)Q = \mathcal{O}(N^{-m-1})$ are absorbed into functions g. Clearly,

$$y_1(\xi) = y_1(-1) + \int_{-1}^{\xi} g_1(s)ds, \ y_2(\xi) = -g_2(\xi).$$
 (42)

The equation is equivalent to [11]

$$\begin{cases} \dot{y}_1 = g_1(\xi), \\ y_1(-1) = -(B_1 + D_1)^{-1} \int_{-1}^1 g_1(s) ds. \\ y_2(\xi) = -g_2(\xi). \end{cases}$$
(43)

Hence, it is enough for us to consider the corresponding initial value condition for our problem. Obviously, $||y|| \le C||g||$, which implies

$$||I_N x - x_N||_{\infty} \le C||I_N||_{\infty} N^{3/2 - m}. \tag{44}$$

Therefore,

$$||x - x_N||_{\infty} \le ||x - I_N x||_{\infty} + ||I_N x - x_N||_{\infty}$$

$$\le C \begin{cases} (\log N) N^{3/2 - m}, & -1 < \alpha, \beta < -1/2, \\ N^{\gamma + 2 - m}, & \gamma = \max(\alpha, \beta), & \text{otherwise.} \end{cases}$$
(45)

3.3 Numerical Experiments

We present two examples here to illustrate the effectiveness of our algorithm. Both examples are from [9]. Unconfined to the type of points in our theory, we use three different types of collocation points, which is presented in Table 1.

Example 1 Consider the DAE

$$\begin{bmatrix} 0 & 0 & 0 \\ 1 & -t & 0 \\ -1 & t & 1 \end{bmatrix} \dot{x} = \begin{bmatrix} -1 & t & 0 \\ 0 & 0 & 0 \\ 0 & t^2 & 1 \end{bmatrix} x + \begin{bmatrix} e^{t/2} \\ 0 \\ 0 \end{bmatrix}, \quad t \in [-5, 0]$$
$$[1 \ 7 \ 0] \ x(-5) + [0 \ 4 \ 1] \ x(0) = 6.$$

This is an index-two problem with solution

$$x(t) = e^{t/2} \left(1 - \frac{t}{2}, -\frac{1}{2}, t^2 + 4t + 8 \right).$$

Since the solution is smooth, from Table 2 and Fig. 2, we actually observe a super geometric rate of convergence $(\mathcal{O}(\frac{\gamma}{N})^{\sigma N})$, where γ , σ are some constants [16].

Table 1 Three sets of collocation points

	ρ	σ
Set 1	Gauss-Legendre points	Gauss-Legendre-Lobatto points
Set 2	Gauss-Legendre-Lobatto points	Gauss-Legendre-Radau points
Set 3	Gauss-Legendre points	Gauss-Legendre points

Table 2 Errors for Example 1

Set 1	Set 2	Set 3
5.7025e-2	8.2836e-2	5.9604e-2
9.7657e-6	2.2007e-5	9.9000e-6
1.8526e-10	5.3087e-10	1.8701e-10
6.9944e-15	7.9936e-15	1.3323e-14
	5.7025e-2 9.7657e-6 1.8526e-10	5.7025e-2 8.2836e-2 9.7657e-6 2.2007e-5 1.8526e-10 5.3087e-10



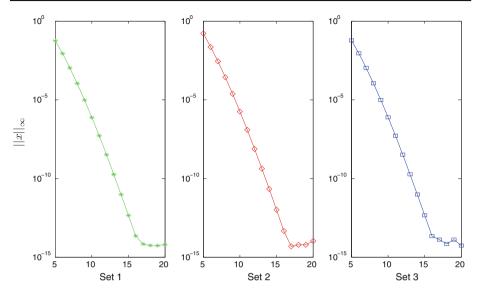


Fig. 2 Numerical errors for Example 1

Example 2 Consider the problem of form (1) with

$$E(t) = \begin{bmatrix} 1 & t & 0 \\ t & 1 & -t \\ p(t) - 2 & -t(p(t) - 2) & 0 \end{bmatrix},$$

$$A(t) = \begin{bmatrix} \kappa - \frac{1}{2-t} & \frac{2}{2-t} - \kappa t & (2-t)\kappa \\ \frac{\kappa - 1}{2-t} - t - 1 & -t\frac{\kappa - 1}{2-t} - 1 & t + \kappa - \frac{\kappa p(t)}{2+t} \\ \kappa t(t^2 - 3) - \frac{p(t) - 2}{2-t} & 2\frac{p(t) - 2}{2-t} - 4\kappa & \kappa(p(t)(2-t) - t^3 + 6t - 4) \end{bmatrix},$$

$$f(t) = \begin{bmatrix} \frac{3-t}{2-t} \\ 2 + \frac{(\kappa + 2)p(t) + \dot{p}(t)}{t^2 - 4} - 2\frac{tp(t)}{(t^2 - 4)^2} \\ (p(t) - 2)\frac{3-t}{2-t} - \kappa(t^2 + t - 2) \end{bmatrix} e^t,$$

$$(46)$$

where $t \in [0, 1]$, parameter $\kappa \in R$ and a smooth function p is a C^1 smooth function. The boundary condition is $x_1(0) = 1$. We choose $\kappa = 20$ and

$$p(t) = -\left(1 + \operatorname{erf}\left(\frac{t - 1/3}{\sqrt{2\varepsilon}}\right)\right), \quad \varepsilon = 10^{-2}$$

There exists an inner layer at t = 1/3 in p and also in the solution

$$x(t) = \frac{e^t}{t^2 + 1} \begin{bmatrix} 1 + t - \frac{t^2}{2 - t} + \frac{tp(t)}{t^2 - 4} \\ 1 - t - \frac{t}{2 - t} + \frac{p(t)}{t^2 - 4} \\ - \frac{t^2 + 1}{2 - t} \end{bmatrix}.$$

If ε is too small, the solution is not smooth. Hence, it is not advantageous to use the *p*-version method. For the choice of $\varepsilon = 10^{-2}$, Table 3 records the numerical errors and Fig. 3 shows the rate of convergence as desired.



Table 3	Errors for	Example 2
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Number of points	Set 1	Set 2	Set 3
10	1.0926e-2	1.0902e-2	1.9384e-2
20	1.0531e-5	4.5635e-5	9.6400e-5
30	6.3180e-9	1.1820e-8	6.5315e-8
40	7.0655e-13	5.3695e-12	7.9923e-12

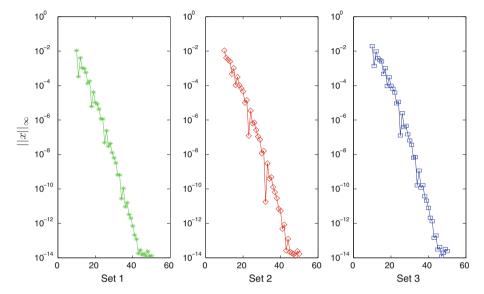


Fig. 3 Numerical errors for Example 2

4 Nonlinear Differential-Algebraic Equations

Similar as linear DAEs, we first transform (5) into its canonical form,

$$\begin{cases}
F_1(\xi, x, \frac{2}{b-a}\dot{x}) = 0, \\
F_2(\xi, x) = 0, \quad \xi \in [-1, 1], \\
f(x(-1), x(1)) = 0.
\end{cases} (47)$$

4.1 Formulation of Algorithm

The collocation points that we use are exactly the same as those in (14), where ρ is for

differential part and σ is for algebraic part. Denote $x_N^m(\xi) = \sum_{k=0}^N C_k^m T_k(\xi), \ m = 1, \dots, n$ the approximation of x in (47). We have the following set of conditions nonlinear equations:

$$f(x_N(-1), x_N(1)) = 0,$$

$$F_1(\rho_i, x_N(\rho_i), \frac{2}{b-a} \dot{x}_N(\rho_i)) = 0, \quad i = 0, \dots, N-1,$$

$$F_2(\sigma_i, x_N(\sigma_i)) = 0, \quad j = 0, \dots, N.$$
(48)



Now we convert it into a nonlinear equation of $C = (C^1, ..., C^n)^T$.

$$\begin{cases}
f\left([W_{1}C^{1}, \dots, W_{1}C^{n}]^{T}, [W_{2}C^{1}, \dots, W_{2}C^{n}]^{T}\right) = 0, \\
F_{1}\left(\rho, [TC^{1}, \dots, TC^{n}]^{T}, \frac{2}{b-a}[UC^{1}, \dots, UC^{n}]^{T}\right) = 0, \\
F_{2}\left(\sigma, [\tilde{T}C^{1}, \dots, \tilde{T}C^{n}]^{T}\right) = 0,
\end{cases} (49)$$

where T, U and \tilde{T} are defined in (18). We use Newton-iteration to solve the above system.

4.2 Numerical Experiments for Nonlinear DAEs

We present four examples in this section to illustrate the application of spectral collocation method for nonlinear DAEs. One important field of application is to solve Hamiltonian systems [5].

Given a Hamiltonian system:

$$\begin{cases}
\frac{dp(t)}{dt} = -\frac{\partial H}{\partial q}, \ p(0) = p_0, \\
\frac{dq(t)}{dt} = \frac{\partial H}{\partial p}, \quad q(0) = q_0,
\end{cases}$$
(50)

where $p(t), q(t) \in \mathbb{R}^n$. Setting $(x_1, \dots, x_{2n}) = (p, q)^T, x_{2n+1} = H$, we obtain a nonlinear DAE on the interval [0, T]

$$\begin{cases} \frac{dx_{i}(t)}{dt} = -\frac{\partial H}{\partial q_{i}}, & i = 1, \dots, n \ (Free) \\ \frac{dx_{j}(t)}{dt} = \frac{\partial H}{\partial p}, & j = n+1, \dots 2n \ (Free) \\ \frac{dx_{2n+1}(t)}{dt} = 0, & (51) \\ H(x_{1}, \dots, x_{n}, x_{n+1}, \dots, x_{2n}) - x_{2n+1} = 0, & \\ x_{i}(0) = x_{0}, & i = 1, \dots, 2n, \\ x_{2n+1}(0) = H_{0}. & \end{cases}$$

It should be pointed out that the system above is overdetermined and in practice, one need to abandon one equation among all free equations. Throughout this subsection, we choose ρ as Legendre–Gauss points and σ as Legendre–Gauss–Lobatto points.

Remark 2 On the interval [nT, (n+1)T], taking the transform $t = \frac{2nT+T}{2} + \frac{T}{2}\xi$, $\xi \in [-1, 1]$, we can obtain the same system as (51) except that $x_i(-1) = \hat{x}_0, x_{2n+1}(-1) = \hat{H}_0$, where \hat{x}_0 and \hat{H}_0 are the value of $x_i(1)$ and $x_{2n+1}(1)$ obtained from the preceding interval, respectively. To better approximate energy, we may specify initial condition as $x_i(0) = x_0$, $i = 1, \ldots, 2n$ and take a poststablization process [2] at the end of each step:

$$\tilde{x}_i(-1) = \hat{x}_0,$$

$$x_i(-1) = \tilde{x}_i - F(\tilde{x})h(\tilde{x}),$$

where $h = H_x$ and $F = H^T (HH^T)^{-1}$.

Example 3 We first consider an example from [10].

$$\begin{cases} \dot{x}_1 = (\epsilon + x_2 - p_2(t))x_4 + \dot{p}_1(t), \\ \dot{x}_2 = \dot{p}_2(t), \\ \dot{x}_3 = x_4, \\ 0 = (x_1 - p_1(t))(x_4 - e^t). \end{cases}$$



Number of points	5	10	20	30
$\max_{1 \le i \le 4} \ x(\sigma_i) - x_i\ $	1.9566e0	2.6436e-3	1.1258e-10	4.2188e-15

Table 4 Errors for Example 3

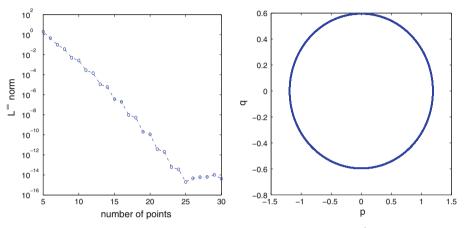


Fig. 4 (Left) numerical errors for example 3; (Right) a phase plot of (p, q) on $[0, 10^4]$ —for Example 4

Choosing the boundary condition

$$x_1(0) = P_1(0) + \epsilon$$
, $x_3(0) = 1$, $x_2(1) = p_2(1)$,

the exact solution of the problem is

$$x^*(t) = \left(\epsilon e^t + p_1(t), p_2(t), e^t, e^t\right).$$

We choose $p_1(t) = \sin(4\pi t)$, $p_2(t) = \sin(t)$, and $\epsilon = 1/2$ as parameters. Using our algorithm, we obtain Table 4 and Fig. 4(Left). Clearly, the convergence rate is geometric.

Example 4 Consider a simple linear Hamiltonian system [7]:

$$\begin{cases} p'(t) = -4q(t), \\ q'(t) = p(t). \end{cases}$$

with initial condition p(0) = 0, q(0) = 0.6 and Hamiltonian $H(p,q) = 1/2p^2 + 2q^2 = 0.72$. Although the Hamiltonian system itself is linear, if we consider the Hamiltonian as a variable and apply our algorithm, we obtain a nonlinear DAE. In [7], the system is treated as ODEs and numerically solved by spectral collocation method. On $[0, 10^5]$, their algorithm has an energy error at the terminal point 2.2949×10^{-10} whereas our algorithm has an error 3.2968×10^{-11} . Furthermore, applying the poststablization approach mentioned above, we can reduce the error to machine epsilon. Hence, our algorithm is much more accurate from this point of view. Figure 4(Right) shows that our algorithm keeps the symplectic structure of the system also. However, our algorithm takes longer CPU time to obtain the result, which is 1, 142 s while theirs takes only 621 s for 20 collocation points.



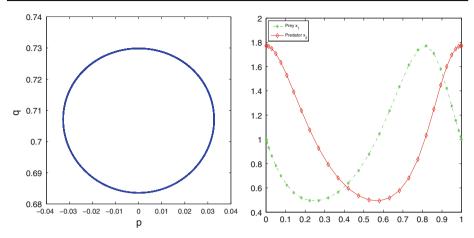


Fig. 5 (Left) a phase plot of (p, q) on [0, 5000]—for Example 6; (Right) prey/predator interaction in one cycle—for Example 5

Example 5 Consider the Lotka–Volterra system for a predator/prey interaction, which consists of two equations [10].

$$\begin{cases} \dot{x}_1 = x_1(1 - x_2), \\ \dot{x}_2 = -cx_2(1 - x_1), \end{cases}$$
 (52)

where c > 0. In the system, the quantity,

$$H = c(x_1 - \log x_1) + (x_2 - \log x_2)$$

remains constant along time, which forces solutions x_1 and x_2 to be periodic. But the period is unknown. As in [10], we introduce new variables $x_3 = H$ and $x_4 = T$ and transform the equation into the interval [0, 1].

$$\begin{cases} \dot{x}_1 = x_1(1 - x_2)x_4, \\ \dot{x}_3 = 0, \\ \dot{x}_4 = 0, \\ c(x_1 - \log x_1) + (x_2 - \log x_2) - x_3 = 0, \ t \in [0, 1] \end{cases}$$
(53)

with initial conditions $x_1(0) = 1$, $x_1(0) = x_1(1)$, and $x_3(0) = 2.2$. For the choice of c = 1, we obtain T converges to 6.4943297198 as the number of collocation points increases. Figure 5(Right) represents the computed interaction between prey and predator in one cycle for 30 collocation points.

Example 6 Consider a system with Hamiltonian $H(p,q) = p^2 - q^2 + q^4$ [7]. The corresponding system of nonlinear ODEs is

$$\begin{cases} \dot{p} = 2q - 4q^3, \\ \dot{q} = 2p \end{cases}$$

with initial condition p(0) = 0, q(0) = 0.73. Hence, the Hamiltonian is -0.24891759. Our DAE algorithm takes only 17s to obtain a result on [0, 5000], with an energy error 4.72×10^{-13} . It can be further reduced to machine epsilon by applying the poststablization process. A phase plot can be found in Fig. 5(Left).



5 Conclusion

In this work, a Jacobi spectral collocation method is applied to solve both linear and nonlinear DAE with arbitrary index. We apply a type of Jacobi–Gauss collocation scheme with N knots to differential part of DAEs whereas another type of Jacobi–Gauss collocation scheme with N+1 knots to algebraic part of the equation. Convergence analysis for linear DAEs shows that the spectral collocation method works efficiently if the solution itself is smooth enough. In particular, the scheme for nonlinear DAEs can be applied to solve Hamiltonian systems and it keeps the symplectic structure of the system. However, it may take more CPU seconds to obtain a result compared with the method in [7] since we enhance the dimension of the problems.

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