

Exercise 2.18

$f(n) = \{0 \text{ if } n = 0; 2 + f(n-1) \text{ otherwise}\}$

base case:

$f(0)$

= <definition of $f(0)$ >

0

=> <0 = $2 * n$ [$n := 0$]>

$2 * n //$

inductive case, for $f(n + 1) = 2 * (n + 1)$ assuming $f(n) = 2 * n$ as the inductive hypothesis

$f(n + 1)$

= <definition of $f(n)$ [$n := n + 1$]>

$2 + f(n)$

= <inductive hypothesis>

$2 + 2 * n$

= < $2 = 2 * 1$ >

$2 * 1 + 2 * n$

= <combine like terms>

$2 * (n + 1) //$

Exercise 3.14(a)

The `closest-power` process generates a recursive process because the function calls itself on line 5 and uses itself to do work.

Exercise 3.15

(a)

$(f\ 1) \Rightarrow (g\ 0) \Rightarrow 1$

$(f\ 2) \Rightarrow (g\ 1) \Rightarrow (f\ 0) \Rightarrow 0$

$(f\ 3) \Rightarrow (g\ 2) \Rightarrow (f\ 1) \Rightarrow (g\ 0) \Rightarrow 1$

(b)

Even arguments cause `f` to return 0, and odd arguments consider `f` to return 1.

(c)

The process is iterative because the final function call does the work, rather than having the first call doing the return.