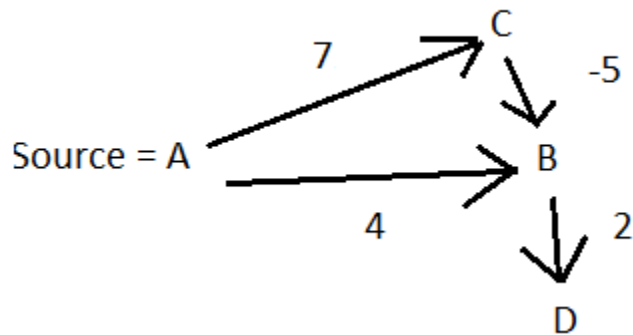


HW8 Theory

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Problem 6.6



First A is popped out of the queue and edges (A, C) and (A, B) are relaxed. B is popped off and relax the edge (B, D) giving a distance of 6. D is now at the top, but has no edges. C is the last vertex left. The edge (C, B) is relaxed giving a total distance to B as 2. Now the queue is empty, but the shortest path to B from A is not correct.

Problem 7.1

```
Insert( val, v, adjList)
```

```
    assert  $v \in \text{adjList}$ 
```

```
    if  $\text{val} \notin \text{adjList}$ 
```

```
        add(adjList, val)
```

```
        ListInsert(adjList[ v ], val)
```

```
    else
```

```
        ListInsert(adjList[ v ], val)
```

```
Delete(val, vertex, adjList)
```

```
    assert val and vertex ∈ adjList
```

```
    ∀ e ∈ adjList [ val ]
```

```
        if e == vertex
```

```
            remove(adjList[ vertex ], e)
```

```
        break
```

```
    ∀ e ∈ adjList [ vertex ]
```

```
        if e == val
```

```
            remove(adjList[ val ], e)
```

```
        break
```

Problem 7.2

Restructuring the adjacency list as a matrix would allow for removal of the first edge in constant time. The down side is that this derivative requires more space than a linked list implementation.

	1	2	3	4	5	6	7
1	0	1	1	1	0	0	0
2	1	0	1	1	0	0	0
3	1	1	0	0	1	1	1
4	1	1	0	0	1	0	0
5	0	0	1	1	0	1	0
6	0	0	1	0	1	0	1
7	0	0	1	0	0	1	0

```
DeleteFirstEdge(vertex, adjMatrix)
```

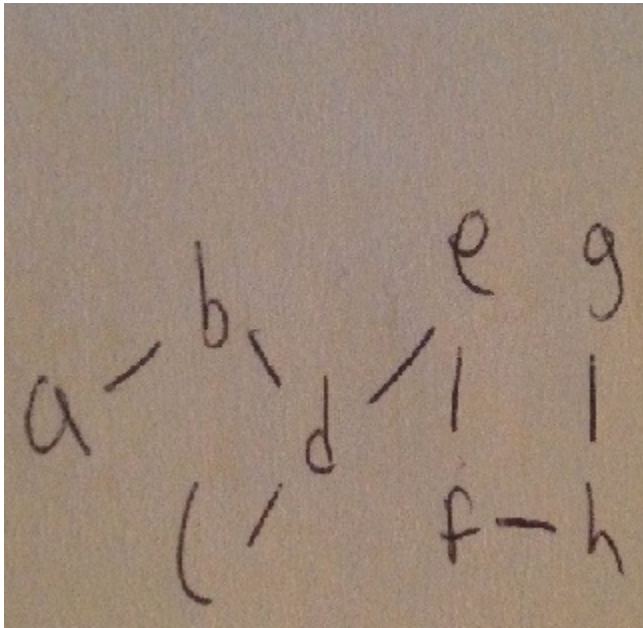
```
    assert vertex ∈ adjList
```

```
    adjMatrix[ vertex[ 0 ] ] [ vertex ] = Null
```

```
    adjMatrix[ vertex ] [ 0 ] = Null
```

Problem 7.3

Part a



Part b

