HW9 Theory

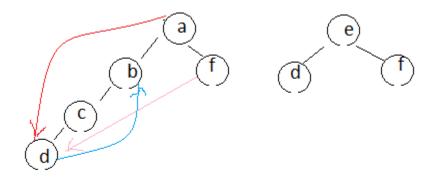
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Problem 6.2

If each T_i is a vertex, and the direction of the graph points to each T_i that needs completion before T_j DFS can be leveraged. First DFS on each vertex should be computed, with each set generated off of each vertex used as a reference for every other vertex referencing that set. This will generate at minimum one topologically ordered set, or at maximum T_n sets. Each set should be walked backwards, with their times t_i summed.

Problem 6.8



Forward Edge Cross Edge Back Edge

Problem 6.10

 $\label{eq:countDepth} \begin{tabular}{l} CountDepth(\ G,\ r\) \\ \\ label\ r\ as\ discovered \\ \\ count\ =\ 1 \\ \\ for\ all\ edges\ from\ v\ to\ w\ in\ G.adjacentEdges(\ r\) \\ \\ \end{tabular}$

```
if w not discovered
                  count += CountDepth(G, w)
      return count
IsRooted (G)
      for all v in G. vertices
            result = CountDepth( G, v )
            if result == length( G.verticies )
                  return true
      return false
Problem 6.14
under the assumption that the paths are unweighted
Time complexity is v^2 + v * e
Paths(G, v, paths = [])
      result = []
      for e in G[v].edges
            result.append( Paths(G, e, paths + [v]))
      return result
LongestPath(G)
      result = []
      for v in G.vertices
            result.append( Paths( G, v ) )
      return max( result )
Problem 6.15
(f,d,b,c)
```

Problem 6.20

too late to write this algo, but it should utilize DFS

Problem 7.3

Part c

