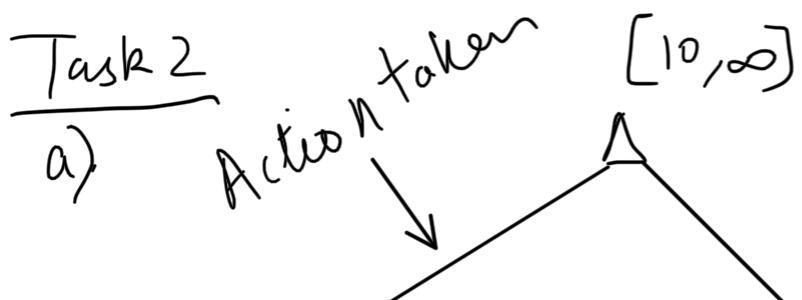
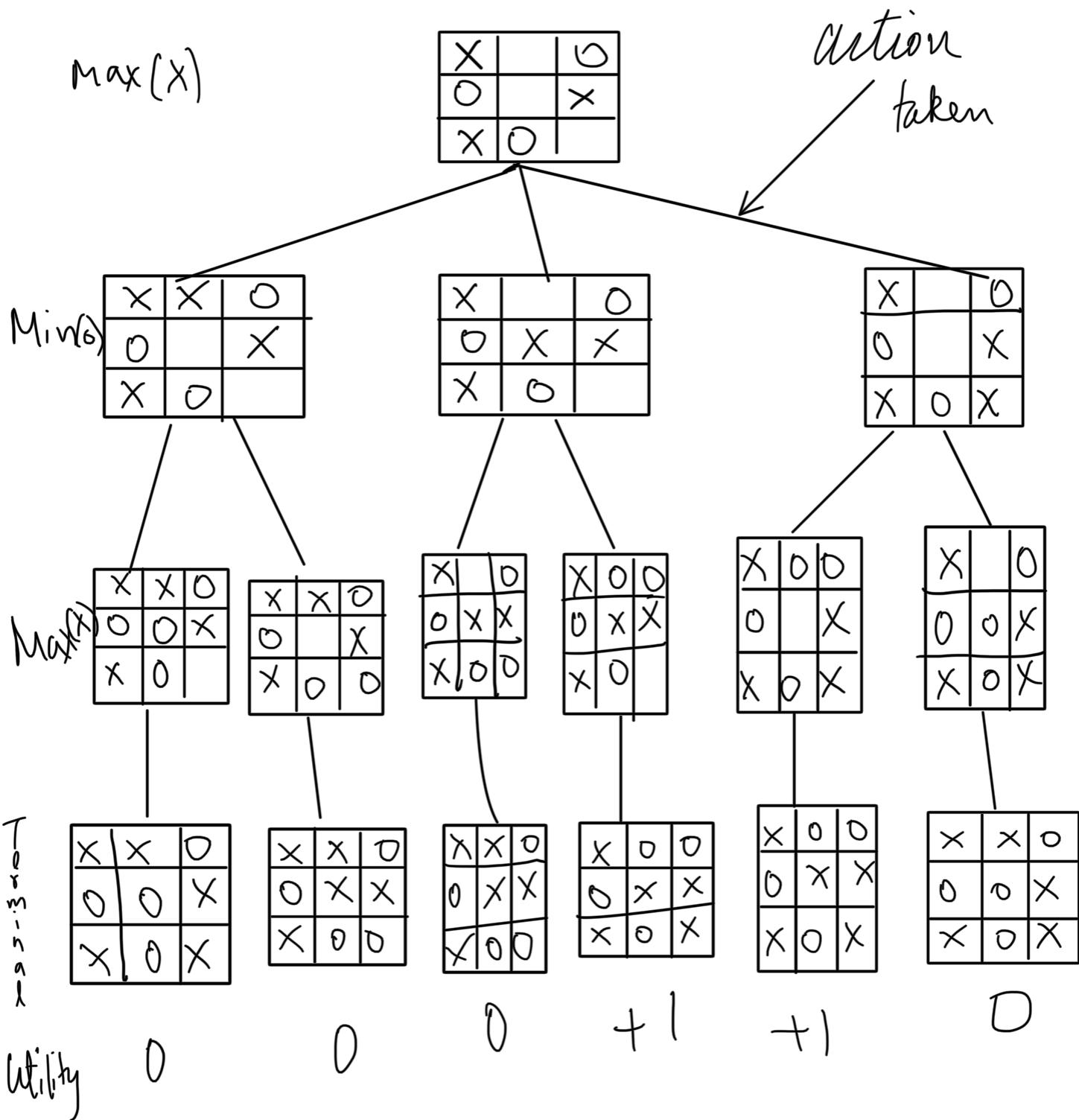
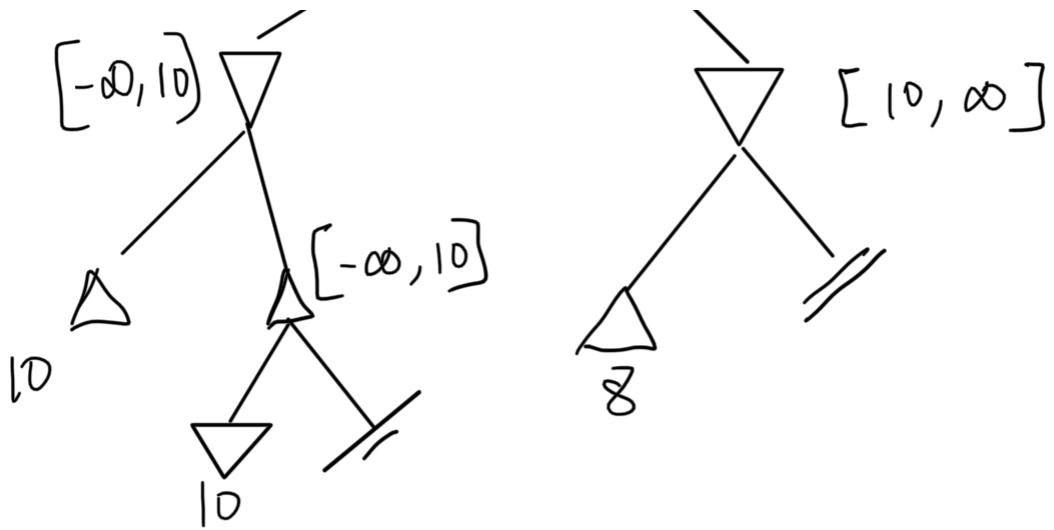


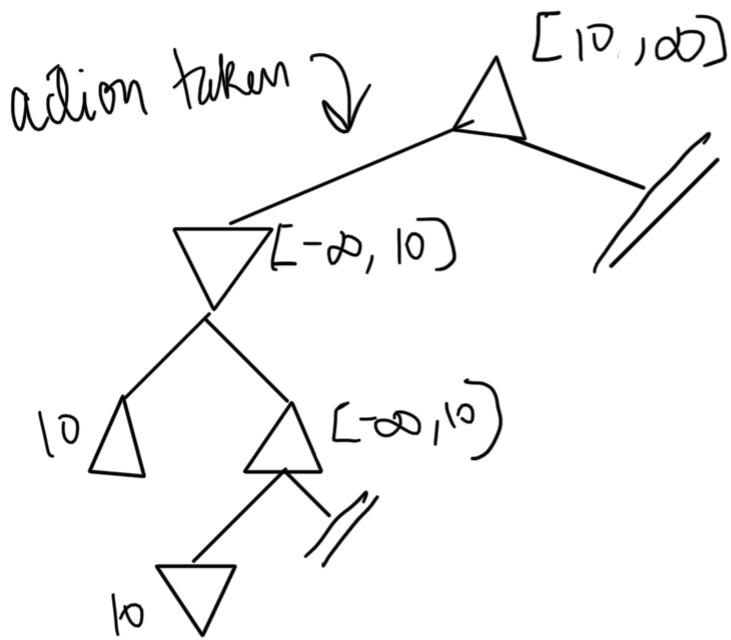
Homework 3

Task - 1





b) If we are given the max utility value is 10, Then we can prune any branch after it finds a successor with a value of 10.



Task 3

Since this is a deterministic game, we will know what move Deep Green would return. We can optimize the code by not taking all

the successors for a min node, but rather take a single successor for each node.

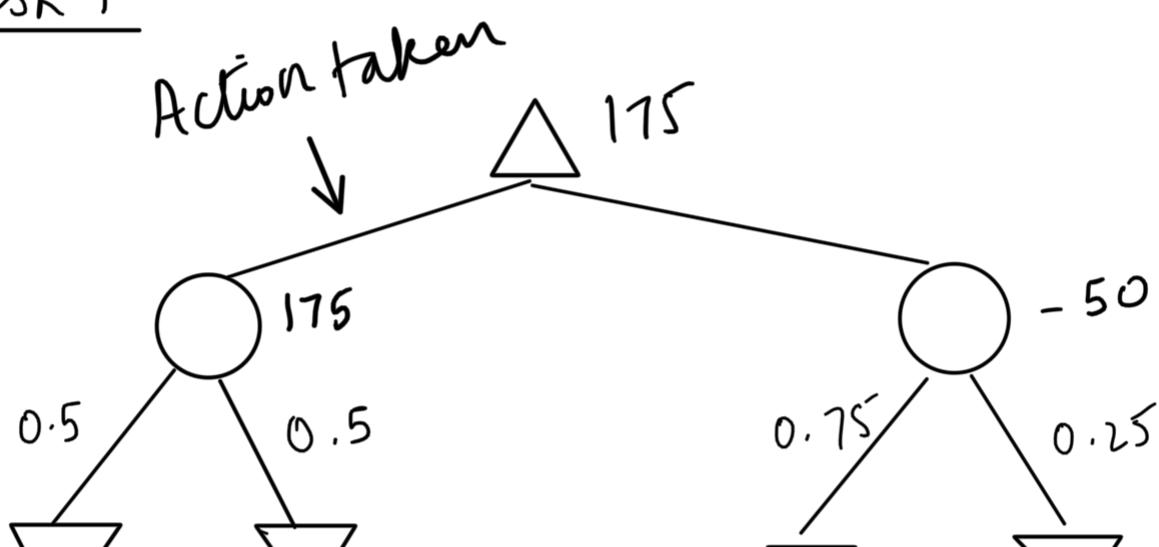
Pseudo code:

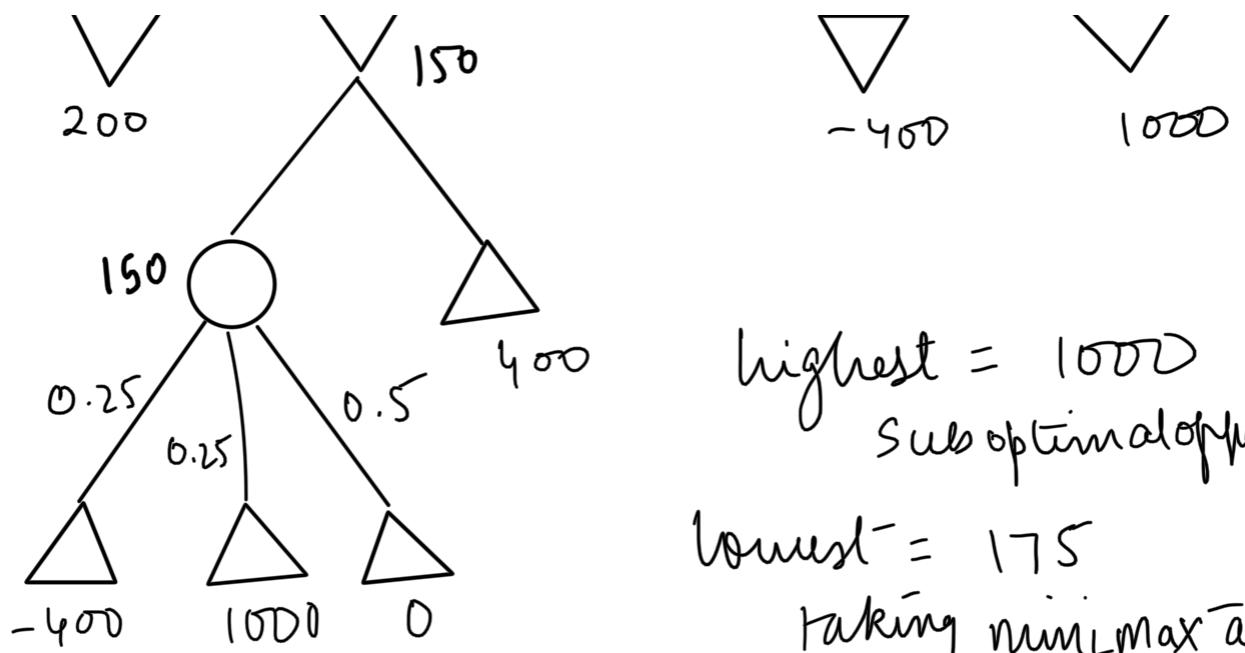
Minimax-decision and max-value function remain the same.

```
function MIN-VALUE (state) returns a utility value
    if Terminal-Test (state) then return UTILITY (state)
    return MAX-VALUE (DeepMinValue (state))
```

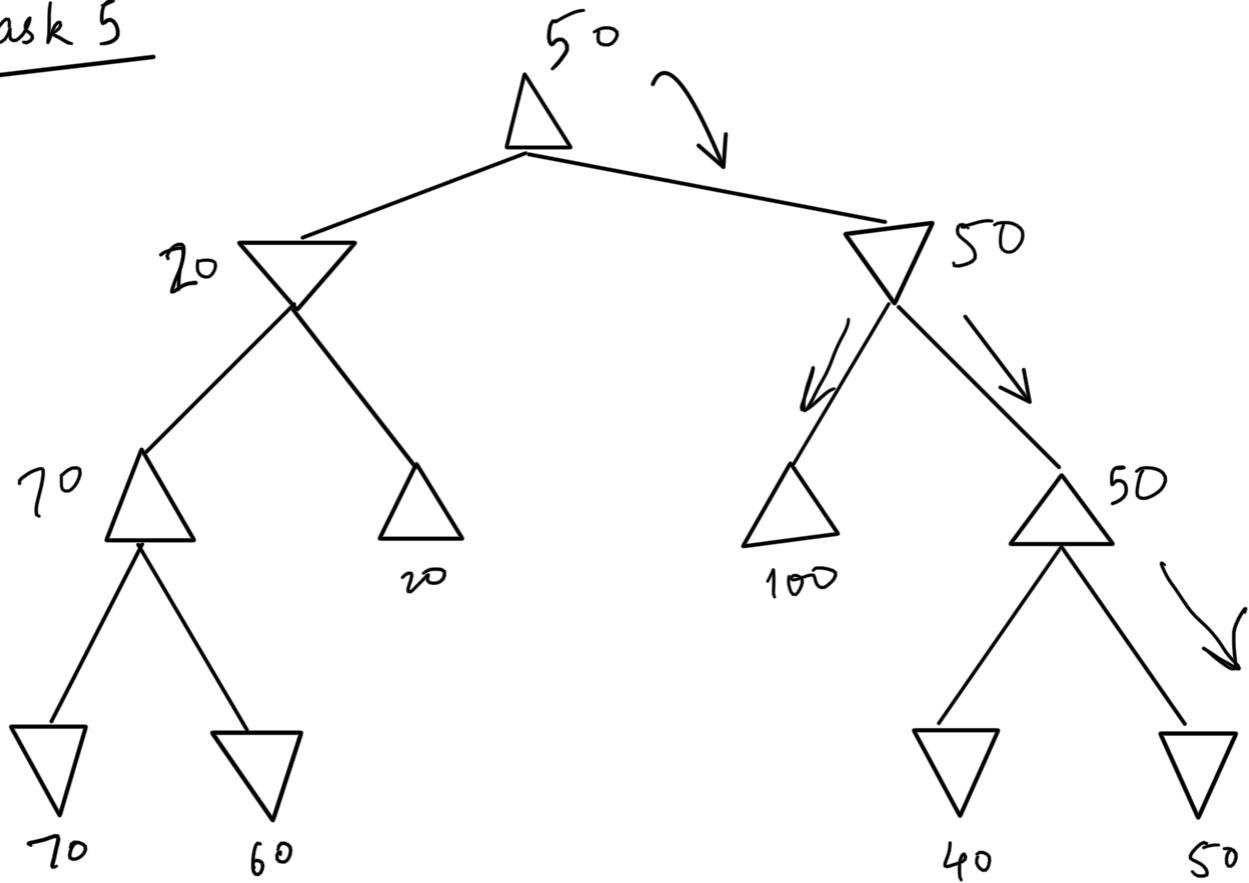
The advantage this algorithm will have is that it will explore fewer nodes as it will only return one successor for min nodes.

Task 4





Task 5



In a complete game under the assumptions stated in the question, the worst outcome will be 50 if the opponent plays optimal +.

strategy.

The best case outcome will be 100 if the opponent plays with a suboptimal strategy for the MAX player.

Task - 6

function CHECK-EQUIVALENCE (KB1, KB2)
returns true or false {

if TT-ENTAILS KB1, KB2 AND
TT-ENTAILS KB2, KB1 {

 return true

} else {

 return false

}

}

Task - 7

- a) KB entails S1 if and only if when S1 is true in all worlds where KB is true.

- b) No, it is not necessary for all rows

where KB is false to have S1 false too.
 There are more than 1 rows in the table where KB is false and S1 is not. So $\text{NOT}(\text{KB})$ does not entail $\text{NOT}(S1)$.

Task - 8

Statement when knowledge base is false.

$$\neg [(A \wedge \neg B \wedge C \wedge D) \vee (\neg A \wedge \neg B \wedge C \wedge \neg D)]$$

$$\neg (A \wedge \neg B \wedge C \wedge D) \wedge \neg (\neg A \wedge \neg B \wedge C \wedge \neg D)$$

Using De Morgan's theorem

$$(\neg A \vee \neg \neg B \vee \neg C \wedge \neg D) \wedge (\neg \neg A \vee \neg \neg B \vee \neg C \vee \neg \neg D)$$

$$(\neg A \vee B \vee \neg C \wedge \neg D) \wedge (A \vee B \vee \neg C \vee D)$$

Task - 9

$$A \Rightarrow B \quad D \Rightarrow A \quad E \Rightarrow F$$

$$B \Rightarrow C \quad E \Rightarrow D \quad F$$

$$C \Rightarrow B \quad C \Rightarrow F$$

i) Forward Chaining

$$\begin{array}{c}
 E \quad \frac{E \quad E \Rightarrow D}{D} \quad \frac{D \quad D \Rightarrow A}{A} \\
 D \\
 A \quad \frac{A \quad A \Rightarrow B}{B} \quad \frac{B \quad B \Rightarrow C}{C} \\
 B \\
 C \\
 \cancel{F} \quad \frac{C \quad C \Rightarrow F}{F}
 \end{array}$$

Therefore $KB \models F$

ii) Backward Chaining

$$\begin{array}{c}
 \boxed{E} \quad \frac{E \quad E \Rightarrow D}{D} \quad \frac{D \quad D \Rightarrow A}{A} \\
 \cancel{D} \\
 A \\
 \cancel{B} \\
 \cancel{C} \\
 F \quad \frac{A \quad A \Rightarrow B}{B} \quad \frac{B \quad B \Rightarrow C}{C} \\
 \cancel{E} \quad \frac{C \quad C \Rightarrow F}{F}
 \end{array}$$

$KB \not\models F$

iii) Resolution

$$\begin{array}{ll}
 A \Rightarrow B & \neg A \vee B \\
 B \Rightarrow C & \neg B \vee C \\
 C \Rightarrow C & \neg C \vee B \\
 D \Rightarrow A & \neg D \vee \neg A \\
 E \Rightarrow D & \neg E \vee D \\
 F \Rightarrow F & \neg F \vee F
 \end{array}$$

$\neg \vdash \neg C \Rightarrow F$ $\vdash \neg \neg C \vee F$

$(\neg A \vee B) \quad (\neg B \vee C) \quad (\neg C \vee D) \quad (\neg D \vee A)$
 $\neg E \quad (\neg E \vee D) \quad (\neg C \vee E) \quad (\neg C \vee F)$
 Add $\neg F$ $\neg C \quad \neg B \quad \neg A \quad \neg D \quad \neg E \quad \boxed{\quad}$ empty clause
 Therefore $KB \not\models F$

Task - 10

- a) A : It rains in May
 B : John gave Mary a check for \$10,000
 C : Mary mowed the lawn.

If it rains, it IMPLIES that John gave Mary a check for \$10K.

$$A \Rightarrow B$$

If John gives Mary 10K, it implies that Mary mowed the lawn.

$$B \Rightarrow C$$

b) $\neg A \wedge B \wedge C$

c) Symbols -

1. A (May)
2. B (John, Mary)
3. B (Mary, John)
4. B (Mary, Mary)
5. B (John, John)
6. C (Mary)

d) $\neg A(\text{May}) \wedge B(\text{John}, \text{Mary}) \wedge C(\text{Mary})$

e) In the above example KB is true
and α is also true.

So $KB \models \alpha$ which means constraint not violated.