

Week 8: Lecture 15 - Lecture 16

Lecture 15

#graph

Dijkstra's algo

// code implementation with priority queue

// code implementation without priority queue

Problem 1

Problem - 20C - Codeforces

Problem 2

CSES - Flight Discount

idea: use Dijkstra's algo to find the min cost route. Then half the most expensive flight on the route.

Actually this idea doesn't work.

Counter example: there are two routes from source to destination

Route 1:

Src \rightarrow A \rightarrow B \rightarrow dst

Cost (src, A) = 4

Cost (A, B) = 2

Route 2:

Src \rightarrow dst

Cost (src, dst) = 5

By the idea we would have chosen route 1 and half the route from src to A, which gives final result = $2 + 2 = 4$

However if we just half cost (src, dst), we would have gotten 2 instead.

Soln:

- Run Dijkstra from src \rightarrow dst
- Create reverse graph G'. Run Dijkstra from dst \rightarrow src
- for each edge e (u, v), half the cost (u, v). Find out cost from src \rightarrow u in G and cost from v \rightarrow dst in G'
- Take the min over all possible results

Lecture 16

#dp

Key to solve dp problems:

- Finding parameters that uniquely define a problem "*DP State*"
E.g. in merge sort parameters are: (start index, end index)
- Dividing problems into overlapping subproblems

Problem 1

CSES - Dice Combinations

idea: $\text{ways}(1) = 1$; $\text{ways}(n) = \text{ways}(n-1) + \text{ways}(n-2) + \dots + \text{ways}(n-6)$

Example:

$\text{ways}(3) = 1$ (last selection) + $\text{ways}(2)$; – actually 2 ways

2 (last selection) + $\text{ways}(1)$;

3 (last selection) + $\text{ways}(0)$;

Results = combine all three situations above = 4

DP initialization

$\text{dp}[0] = 1$

$\text{dp}[1] = 1$

Core code:

```
vector<int> dp(n+1);
dp[0] = 1;
dp[1] = 1;
int ans = 0;
for (int i = 2; i <= n; ++i) {
    for (int j = 1; j <= 6; j++) {
        if (i - j >= 0) {
            ans += dp[i-j];
        }
    }
}
```

Problem 2

CSES - Minimizing Coins

Recurrence relation:

$\text{coins}(n) = \min \{ \text{coins}(n-i) \}$ for all $i + 1$

Problem 3

CSES - Book Shop

For each book there are two decisions to make: to purchase or not to purchase

$f(n, i)$ represents max # of pages possible with budget = n considering book i

$f(n, i) = \max \{ f(n - \text{price}[i], i - 1) + \text{pages}[i], f(n, i - 1) \}$ // the second situation represents when current book is skipped.

// code to do