AdvFS on-disk structures



Module 2
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Objectives



To explain AdvFS on-disk structures, you should be able to describe:

- Bitfiles
- Mcells
- Extent maps
- Tags
- Directory files and directory index files
- Fragments
- Storage bitmap bitfiles
- Quota files

Two-level implementation of AdvFS (1 of 2)

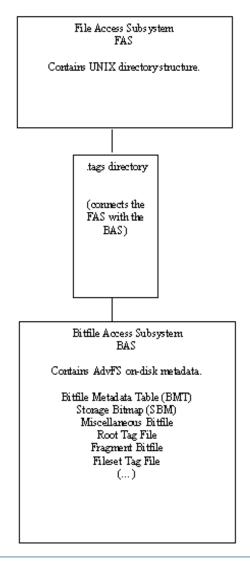


- File Access System (FAS)
 - the higher level of AdvFS
 - transforms bitfiles into normal UNIX files
 - the client layer
- Bitfile Access System (BAS)
 - the lowest level of AdvFS providing storage support

Two-level implementation of AdvFS (2 of 2)



On Disk Structures



.tags directory



- Each AdvFS file system has a .tags directory
- Allows files to be accessed by tag, sequence number
 - -/Advfs_mount_point/.tags/15374
 - -/Advfs_mount_point/.tags/0x3c0e.8001

Tag number can access any file



```
# ls -li
total 31
22894 -rwxr-xr-x 1 root system 31114 Jun 24 15:20
ob 1
#
# tail -3 ob 1
:of#169728:pf#35135:bf#8192:ff#1024:\
:og#99458:pg#149368:bg#8192:fg#1024:\
: oh#0: ph#0: bh#8192: fh#1024:
#
# tail -3 /usr/.tags/22894
:of#169728:pf#35135:bf#8192:ff#1024:\
:og#99458:pg#149368:bg#8192:fg#1024:\
:oh#0:ph#0:bh#8192:fh#1024:
#
```

BAS on-disk format: Everything is a bitfile



- Bitfiles are arrays of 8K disk pages holding user data or metadata
- A series of contiguous 8K pages in a bitfile is stored as an extent
- Each bitfile is identified by its tag
 - Tag consists of a tag number.sequence number pair
- Tag number can be used to locate the extents of a file

Note that tag2name has a new format for V5.

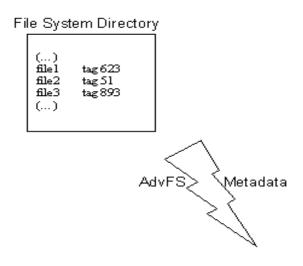
Tag number 22894 being translated by tag2name



```
# /sbin/advfs/tag2name /usr/.tags/22894
/usr/bruden/ob_1
#
# echo $PATH
/sbin:/usr/sbin:/usr/bin:/usr/ccs/bin:/usr/bin/X11:/usr
/local
#
# PATH=$PATH:/sbin/advfs
#
# tag2name usr_domain -S usr 22894
open_vol: open for volume "/dev/disk/dsk2g" failed:
Device busy
#
# tag2name -r usr_domain -S usr 22894
                                           <- Uses raw
device (-r)
bruden/ob_1
```

Using AdvFS metadata to translate FAS to BAS

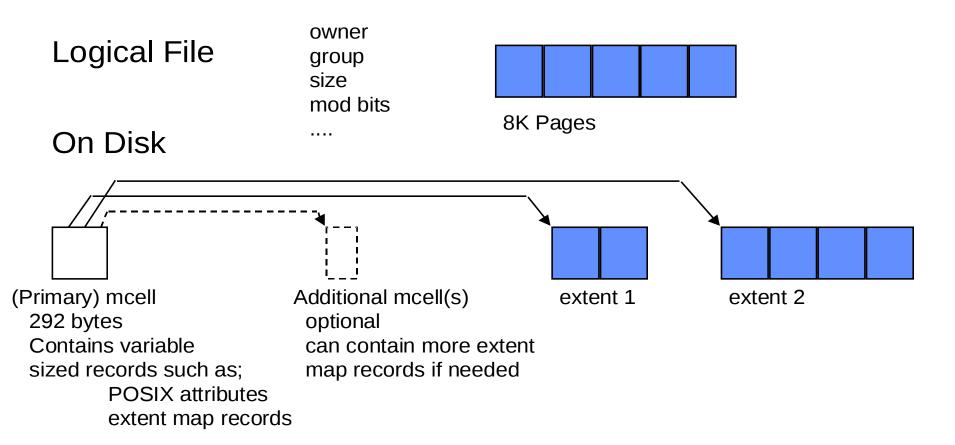




File3 on disk LBN 80334 (AdvFS sees this as a bitfile.)

BAS on-disk format





Bitfiles



- Array of 8K pages
- Stored as extents
 - groups of on-disk contiguous 8K pages
 - managed by extent maps
- Identified by a tag
 - tag.sequence, for example 4714.8001
 - tag number is similar to an inode number
 - sequence number functions as a generation number
- All sectors are free or in a bitfile
- Managed by mcell chains

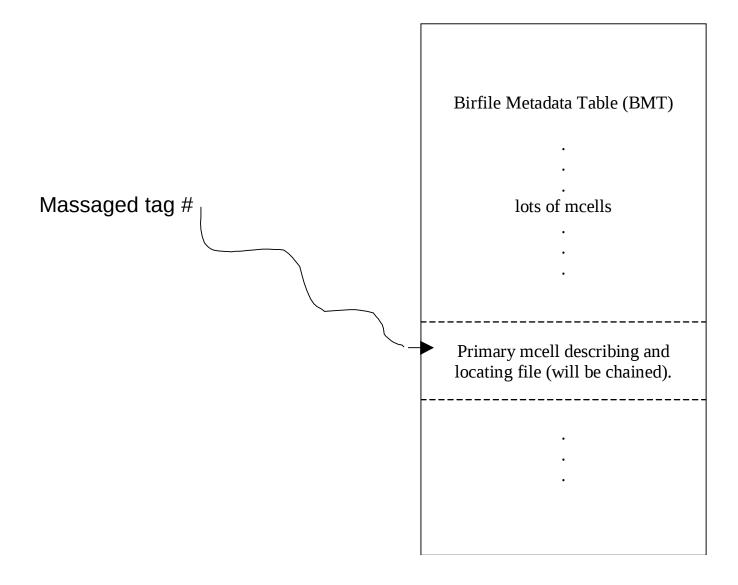
Mcells



- AdvFS locates extents by finding the file's Primary Metadata cell (Mcell) and stepping through the mcells to find the extent information
 - tag number is like an inode number, but an mcell functions like an inode
 - holds permissions, size, extent info, link count, and so forth
- Each mcell is 292 bytes
 - can fit 28 on an 8K page (plus a 16-byte header)
- Where are the mcells located?

Tag number to BMT mcell to logical blocks.





AdvFS file addresses



AdvFS files have a means of identification:

- Tag
 - very similar to the UFS inode number
 - can be discovered with the Is -i command
- Primary mcell ID
 - component of the lower BAS layer
 - start of a linked list of one or more mcells
 - well hidden from users and system administration

Bitfile-sets



- FAS fileset represents a BAS bitfile-set
- Identified by numbers
- Bitfiles are known by:
 - domain ID
 - fileset ID
 - tag, sequence number

Reusing tags



- Tag numbers can be reused
 - with file creation and deletion
 - like inode numbers
- Sequence number identifies various lives of the tag
 - tags have initial sequence number of 8001 hexadecimal or 32769 decimal
 - sequence number is incremented when tag number is reused
 - when sequence number overflows, tag is discarded
 - leftmost bit indicates 'tag in use', remaining 15 bits are used for sequencing
 - hexadecimal 8001 is binary <u>1</u>000000000000001

Domain and volume structures



- Reserved bitfile metadata table
- Bitfile metadata table
- Storage bitmap
- Misc bitfile
- Transaction log
- Root tag file

BAS on-disk metadata bitfiles



per volume

per domain

per fileset

ReservedBitfile MetadataTable

Storage Bitmap

MISC Bit File

BitfileMetadataTable

Root Tag File

On-Disk Log

ReservedBitfile MetadataTable

Storage Bitmap

MISC Bit File

BitfileMetadataTable

Tag File: filesetA

Fragment FilefilesetA

Per domain bitfiles



- On-disk log
 - contains the transaction log
 - usually 4MB in size
- Root tag file
 - lists filesets
 - one page in size

Displaying the root tag file



nvtagpg -r usr_domain

```
DOMAIN "usr_domain" VDI 1 (/dev/rdisk/dsk2g) lbn 96 root TAG page 0

currPage 0

numAllocTMaps 3 numDeadTMaps 0 nextFreePage 0 nextFreeMap 5

tMapA[1] tag 1 seqNo 1 primary mcell (vol,page,cell) 1 0 1 usr

tMapA[2] tag 2 seqNo 1 primary mcell (vol,page,cell) 1 0 13 var

tMapA[3] tag 3 seqNo 1 primary mcell (vol,page,cell) 2 2 4

ob_fset

#
```

Per volume bitfiles



Each volume is supported by these bitfiles:

- Reserved Bitfile Metadata Table (RBMT)
 - contains mcells for reserved bitfiles
 - last mcell in each RBMT page will link to the next page
 - eliminates BMT fragmentation problems

Displaying RBMT summary information (1 of 2)



```
# nvbmtpg -r -R usr_domain
DOMAIN "usr_domain" VDI 1 (/dev/rdisk/dsk2g) lbn 32
                                                      RBMT
page 0
There is 1 page in the RBMT on this volume.
There are 19 free mcells in the RBMT on this volume.
DOMAIN "usr_domain" VDI 2 (/dev/rdisk/dsk4b) lbn 32
                                                      RBMT
page 0
There is 1 page in the RBMT on this volume.
There are 19 free mcells in the RBMT on this volume.
```

Displaying RBMT summary information (2 of 2)



- Bitfile metadata table (BMT)
 - contains all the mcells for nonreserved bitfiles
 - grows as new files are created
- Storage bitmap (SBM)
 - contains 1 bit for every 8K bytes
 (1 bit per 1K in Tru64 UNIX V4.0)
- Misc bitfile
 - contains bootblocks, and so forth
 - four pages in size

Displaying BMT for volume 1 of usr_domain



Per fileset bitfiles (1 of 2)



- Tag file (not .tags)
 - translates tag # into location of primary mcell (within BMT) for the appropriate file
 - formerly called the 'Tag Directory File'
- Tags can be reused as files are deleted
 - limited to size of associated sequence number
 - 8001 is a typical sequence number showing that this tag is in use (left bit is set) and it is in use for the first time (001)
 - limits tag reuse to ~4k times before tag is dead

Per fileset bitfiles (2 of 2)



Tag file consists of 8K pages with 1022 tagmap entries

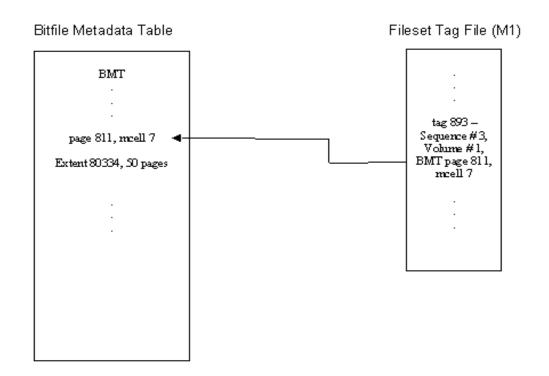
(8 bytes each) preceded by a 16-byte header

Tagmap entry contains:

- Sequence number
- Volume index
- Primary Mcell ID (BMT page # and cell # within page)

Fileset tag directory locating primary mcell





Displaying a fileset tag file



nvtagpg -r usr_domain usr -a DOMAIN "usr domain" VDI 1 (/dev/rdisk/dsk1g) lbn 112 "usr" TAG page 0 numAllocTMaps 1021 numDeadTMaps 0 nextFreePage 30 nextFreeMap 0 tMapA[0] tmFreeListHead 50 tmUninitPg 50 tMapA[1] tag 1 segNo 1 primary mcell (vol,page,cell) 1 0 3 tMapA[2] tag 2 segNo 1 primary mcell (vol,page,cell) 1 0 4 tMapA[3] tag 3 seqNo 1 primary mcell (vol,page,cell) 1 0 6 (\ldots) tMapA[1021] tag 1021 seqNo 1 primary mcell (vol,page,cell) 1 40 23 DOMAIN "usr_domain" VDI 1 (/dev/rdisk/dsk1g) lbn 128 "usr" TAG page 1 numAllocTMaps 1022 numDeadTMaps 0 nextFreePage 30 nextFreeMap 0 tMapA[0] tag 1022 segNo 1 primary mcell (vol,page,cell) 1 40 25 tMapA[1] tag 1023 primary mcell (vol, page, cell) 1 40 27 seqNo 1

Finding primary mcell thru tag directory using tag number



```
# ls -li bigfile
 2324 -rwxr-xr-x 1 root system 15429536 Nov 24 14:54
bigfile
# df .
Filesystem 512-blocks
                                 Used Available Capacity
Mounted on
var_domain#var 524288
                               273858 238416 54%
                                                           /var
# ntagpg -r var_domain var -t 2324
DOMAIN "var_domain" VDI 1 (/dev/rdisk/dsk1h) lbn 144 "var" TAG page 2
numAllocTMaps 350 numDeadTMaps 0 nextFreePage 0 nextFreeMap 337
tMapA[280] tag 2324 seqNo 49 primary mcell (vol, page, cell) \frac{1}{104}
```

Finding primary mcell thru tag directory using file name



```
# ls -li bigfile
 2324 -rwxr-xr-x 1 root system 15429536 Nov 24 14:54
bigfile
# pwd
/var/cluster/members/member0/tmp
# df .
Filesystem 512-blocks Used Available Capacity Mounted on
var_domain#var 524288 273858
                                   238416
                                         54%
                                                 /var
# ntagpg -r var_domain var nvtagpg -r var_domain var \
             cluster/members/member0/tmp/bigfile
DOMAIN "var_domain" VDI 1 (/dev/rdisk/dsk1h) lbn 144 "var" TAG page 2
numAllocTMaps 350 numDeadTMaps 0 nextFreePage 0 nextFreeMap 337
tMapA[280] tag 2324 seqNo 49 primary mcell (vol,page,cell) 1 104
```

Fragment bitfile



- Contains small files including the last parts of small files
- Size varies with number of small files
- Has tag number 1

Reserved bitfile special names (1 of 3)



All accessible under .tags directory

<u>Name</u>	<u>Function</u>	One Per
M-6, M-12,	Reserved Bitfile Metadata Table	Volume
M-7, M-13,	Storage Bitmap	Volume
M-8, M-14,	Root Tag File	Domain
M-9, M-15,	Transaction Log	Domain
M-10, M-16,	Bitfile Metadata Table	Volume
M-11, M-17,	.Miscellaneous Bitfile	Volume

Reserved bitfile special names (2 of 3)



<u>Name</u>	<u>Function</u>	One Per
M1	Fileset tag file (not .tags) for fileset #1.	Fileset
M2	Fileset tag file (not .tags) for fileset #2.	Fileset
МЗ	Fileset tag file (not .tags) for fileset #3.	Fileset
()		
Mn	Fileset tag file (not .tags) for fileset #n.	Fileset
1	Fragment Bitfile	Fileset

Reserved bitfile special names (3 of 3)



<u>Name</u>	<u>Function</u>	One Per
2	Fileset's Root Directory	Fileset
3	.tags Directory	Fileset
4	User Quota File	Fileset
5	Group Quota File	Fileset
6	User File with Tag # 6	Fileset
7	User File with Tag # 7	Fileset
()		
n	User File with tag # n	Fileset

Metadata bitfile tags



Tags for reserved bitfiles of virtual disk *i* are:

tag = - (magic-number-shown-below + (vol_index * 6))

Metadata Bitfile Tags Example

Reserved File Formu	<u>Disk</u>	<u>Disk 2</u>		
RBMT	- (0 +	(vol * 6))	-6	-12
SBM	- (1 +	(vol * 6))	-7	-13
Root tag file	- (2 +	(vol * 6))	-8	-14
Log	- (3 +	(vol * 6))	-9	-15
BMT	- (4 +	(vol * 6))	-10	-16
Misc Bitfile	- (5 +	(vol * 6))	-11	-17

 May be printed in weird tags which effectively translate to negative numbers

.tags for directory entries for metadata bitfiles



Start with an M

- For virtual disk specific files, use the negative number
- For fileset tagfiles, use the fileset ID

For example, if /usr is an AdvFS fileset:

```
File Description

/usr/.tags/1 Fragment bitfile

/usr/.tags/M-6 RBMT of disk 1

/usr/.tags/-6 RBMT of disk 1 also

/usr/.tags/M-15 Log of disk 2

/usr/.tags/M2 tag file for 2nd fileset
```

Bitfile metadata table (1 of 2)



- BMT is represented by a file found under tags directory
 - special file name is M-10 (more later)
 - contains a series of 8K pages just like any other AdvFS file
 - pages contain mcells (292 bytes each) and header information
- Each mcell contains one or more variable-length records
 - describe various file attributes, extents, permissions, frag info, etc
- BMT can grow just as any file can grow just adds another extent
 - starts with slightly more than 1M (can be tailored)
- BMT is created when the mkfdmn command is issued
 - one BMT for each volume in the domain

Bitfile metadata table (2 of 2)



- Contains all mcells for all files other than the reserved bitfiles
 - user files
 - user directories
- Mcells for the reserved bitfiles are in RBMT
- RBMT and BMT:
 - first mcell describes itself
 - as more mcells are needed, it grows using extents
 - RBMT reserves the last mcell on each page to chain to other pages of mcells

Reserved bitfiles on disk layout



Miscellaneous Bitfile (M-11) Pages 0, 1 Sectors 0-31
RBMT (M-6) Page 0 Sectors 32-47
BMT (M-10) Page 0 Sectors 48-63
Miscellaneous Bitfile (continued) Pages 2,3 Sector 64-95
Root Tag Directory (M-8) Page 1 Sectors 96-111
Storage Bitmap (M-7) (1 bit per 8k cluster) Sectors 112-?
Transaction Log (M-9) 512 Pages Sectors ?
Fileset Tag Directory File (M1) 8 Pages Sectors ?
Mount Point Directory for Fileset (2) 1 Page Sector ?
.tags Directory (3) 1 Page Sector ?
Quota.user (4) 1 Page Sector ?
Quota.Group (5) 1 Page Sector ?

Mcell records (1 of 2)



- The inodes of AdvFS
- 28 Fixed-size (292 byte) mcells are packed into 8K pages
- One or more linked mcells describe bitfiles
 - First mcell in list is primary mcell
- Each mcell contains variably sized records describing attributes of the bitfile
- Contained in the Bitfile Metadata Table (BMT) and the RBMT

Mcell records (2 of 2)



Record Types include:

- Extent Maps (of various kinds)
- Bitfile Attributes (clone, original, etc.)
- Domain Attributes
- Virtual Disk Attributes (disk ID, disk index, etc.)
- Fragment Attributes more later
- POSIX file stats (permissions, size, link count, etc.)
- Symbolic link targets

Mcell page structure



Page

page header

mcell

mcell header

record

record

<u>record</u>

mcell header

record

record

<u>record</u>

mcell header

record

record

<u>record</u>

variable sized records

per page

28 mcells

mcell

mcell

RBMT page 0



RBMT Page 0:

- Starts at sector 32 (LBN 32)
- Contains these primary mcells:
 - mcell 0 Reserved Bitfile Metadata Table (RBMT)
 - mcell 1 Storage Bitmap (SBM)
 - mcell 2 Root Tag File Optional, one per domain
 - mcell 3 Log Optional, one per domain
 - mcell 4 Bitfile Metadata Table (BMT)
 - mcell 5 Misc bitfile

BMT page 0



- Starts at Sector 48
- mcell 0 is head of the BMT page free list
- Contains mcells for nonreserved bitfiles, such as user files and directories
- All other BMT pages are found via the RBMT

BMT page format



Each 8192-byte page is composed of:

- 16-byte header followed by
- 28 292-byte mcells

BMT header consists of:

- Pointer to next free mcell on page
- Pointer to next page with free mcells
- Number of free mcells on the page
- Page number (within BMT)
- AdvFS version (now 4)

Mcell addresses



- Mcells are addressed by a 32-bit mcell ID, bfMCIdT
 - 27 bits give the mcell's BMT page number
 - 5 bits give the mcell's position within its page
- Every bitfile has a primary mcell
- Tagfiles map (tag file)
 - tag numbers can be mapped to primary mcell locations

Reserved mcell addresses



Within RBMT page 0, the slot numbers listed are used to calculate the tag numbers for the reserve bitfiles:

- 0: the RBMT itself
- 1: SBM, storage bitmap
- 2: root tag file
- 3: transaction log
- 4: BMT
- 5: misc bitfile

Mcell format



- Begins with 24-byte header
 - 32-bit ID of the next mcell in the chain and virtual disk containing next mcell
 - position of this mcell within the chain
 - tag number of the bitfile
 - tag number of the fileset
- Has 268 bytes remaining for mcell records

Mcell records



- Vary in size and type
- Begin with a 4-byte header
 - 2-byte count of record size
 - 1-byte type field (There are currently about 20 record types)
 - 1-byte version number for different versions of a type
- Have a null record whose size is 4 and type is 0
 - this is the end-of-records indicator

Utilities for viewing records



- nvbmtpg
- nvfragpg
- nvlogpg
- nvtagpg
- vsbmpg
- vfilepg
- savemeta

Extent maps for nonreserved files



For nonreserved files:

- Primary extent map record
 - within the primary mcell
 - allocated when file gets a page
 - if full, points to an extra extent map
- Extra extent map record
 - allocated as the file grows

For striped files, extent map records use a shadow extent map and may point to more than one disk.

Extent maps for reserved files



For reserved bitfiles, for example, the RBMT

- Primary extent map
 - if full, points to extra extent map
- Extra extent map
 - usually only needed for the BMT itself

All records for reserved files are in the RBMT

Encoding of extents



Extents information is captured in two fields

- page number within bitfile
- block number within virtual disks

Size of one extent is inferred from the next

- compute difference between the page numbers
- place -1 in block number to indicate a hole (-2 for clone hole)

• Try the following:

- page = 0, block = 8000
- page = 100, block = -1
- page = 200, block = 9600
- page = 300, block = -1

Tag file characteristics



Bitfiles are identified by tags

Tag files are:

- arrays of tagmap entries
- indexed by tag number

Tag File entries contain:

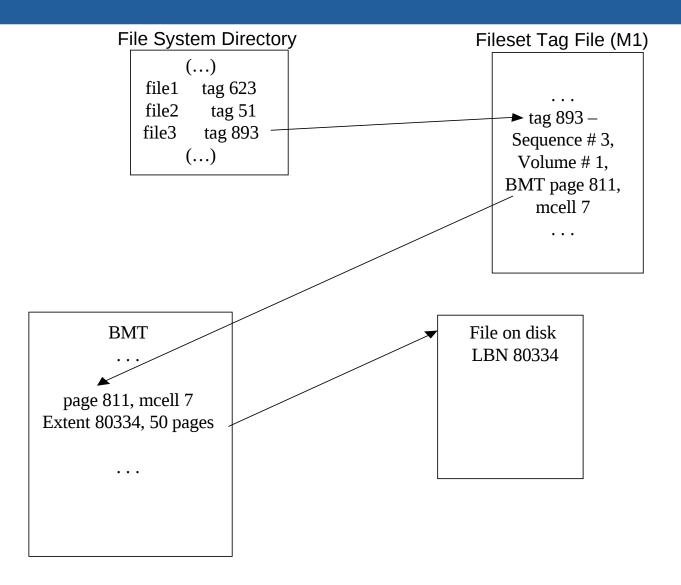
- sequence number (high bit set if in use)
- volume index
- mcell ID within volume

Justifying Tag Files

 tag files introduce an extra level of overhead in accessing file data

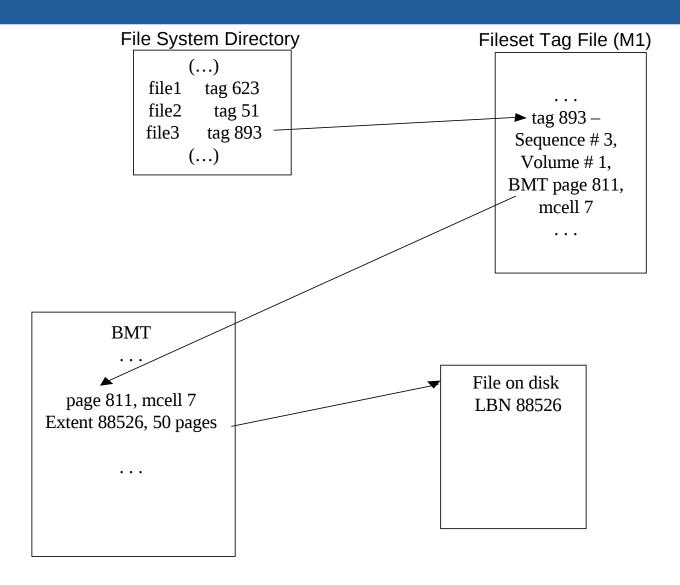
File access through tag file





Tag file allowing transparent data move





Tag file page



- Starts with a five-field, 16-byte header
 - number of this page (for sanity checking)
 - next page with free tagmap entries
 - next free tagmap within this page
 - number of of allocated tagmaps
 - number of dead tagmap entries
- Followed by 1022 tagmap entries

Tagmap entries



Head of free tagmap list

- stored in first slot of page 0
- points to:
 - first page with a free tagmap entries
 - first uninitialized page

Element of free tagmap list

- sequence number
- pointer to next free entry on this page

Allocated tagmap entry

- sequence number
- virtual disk of bitfile
- mcell ID of bitfile

Root tag file



- Entries for each fileset of the domain
- Contains mcell IDs of fileset tagfiles
- Can be used to find the list of domain filesets

Fileset tag file



- Each fileset has its own tag file
- Maps fileset's bitfiles to primary mcells
- Has these special fileset tags
 - 1 Fragment bitfile
 - 2 Root directory
 - 3 .tags
 - 4 User quota file
 - 5 Group quota file

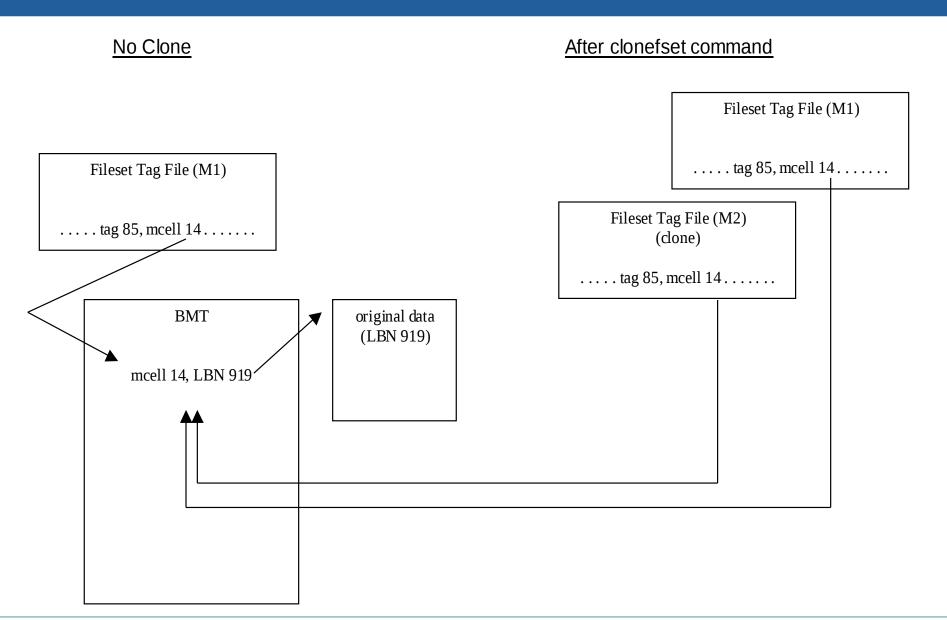
Cloning through fileset tag



- Clone fileset is a read-only virtual copy of the data as it existed at the time the clone was created
 - should be used as source for a backup operation and then removed
 - creates a copy of the Fileset Tag File (M1)
 - still points to the same BMT
- As original data changes while clone exists:
 - copy of the original data must be created
 - new mcell must be allocated in the BMT
 - appropriate entry in new copy of Fileset Tag file must be updated to point to cloned page's new mcell in BMT

Fileset tag file before and after cloning





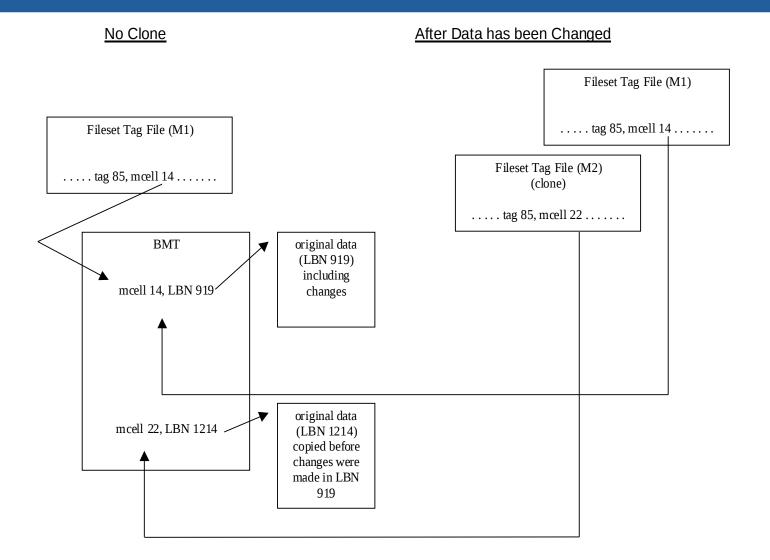
Fileset tag file before and after cloning



- No data is copied unless a change is made to original data
- A copy is only made of data that is about to be changed
- Clone is effectively a 'snapshot' of the data at a known time

Clone structures after data write





Utility for viewing tag files



The nvtagpg utility:

 Prints formatted pages of a root tag file or a fileset tag file

Displaying root tag file information



nvtagpg -r usr_domain

```
DOMAIN "usr_domain" VDI 1 (/dev/rdisk/dsk2g) lbn 96 root TAG page 0

currPage 0

numAllocTMaps 3 numDeadTMaps 0 nextFreePage 0 nextFreeMap 5

tMapA[1] tag 1 seqNo 1 primary mcell (vol,page,cell) 1 0 1

usr

tMapA[2] tag 2 seqNo 1 primary mcell (vol,page,cell) 1 0 13

var

tMapA[3] tag 3 seqNo 1 primary mcell (vol,page,cell) 2 2 4

ob_fset
```

Displaying fileset tag file



Given a particular tag number, divide by 1022:

- Quotient is the page
- Remainder is tagmap slot

```
# nvtagpg -r usr_domain usr
```

```
DOMAIN "usr_domain" VDI 1 (/dev/rdisk/dsk2g) lbn 1047552 "usr" FRAG page 0
currPage 0
numAllocTMaps 1021
                  numDeadTMaps 0 nextFreePage 23 nextFreeMap 0
tMapA[1] tag 1
                   segNo 1
                               primary mcell (vol,page,cell)
                                                            1 0 3
                               primary mcell (vol,page,cell) 1 0 4
tMapA[2] tag 2
                   seqNo 1
tMapA[3] tag 3
                   seqNo 1
                               primary mcell (vol, page, cell) 1 0 6
                               primary mcell (vol, page, cell) 1 0 7
tMapA[4] tag 4 seqNo 1
                               primary mcell (vol,page,cell)
tMapA[5]
        tag 5
                   seqNo 1
                                                            1 0 8
(\ldots)
```

Displaying individual file's tag file entry



```
# ls -li big1
22896 -rwxr-xr-x 1 root system 13729520 Jun 24 16:53 big1
#
# nvtagpg -r usr_domain -T 1 -t 22896
DOMAIN "usr_domain" VDI 1 (/dev/rdisk/dsk2g) lbn 2055136
"usr" FRAG page 22
currPage 22
numAllocTMaps 1022 numDeadTMaps 0 nextFreePage 0 nextFreeMap 0
tMapA[412] tag 22896 seqNo 1 primary mcell (vol,page,cell) 1
951 15
# bc
22896/1022
22
22896%1022
412
VD
```

UNIX directories



UNIX directories are contained in standard bitfiles

Directories are a series of directory entries

POSIX directory entry format

- inode number: 32 bits

- entry length: 16 bits

- name length: 16 bits

name: variable-length, zero-padded to fill the entry

AdvFS directory information hidden in POSIX entry

- tag numbers replace inode numbers
- the 64-bit tag. sequence ID is placed in the name "padding"

AdvFS directory index files



File name lookup through large directory files is very slow

On average, one-half of the directory entries must be searched

AdvFS uses a B-tree to index files to speed the search

- Each page of the index file has 510 entries
- Filenames are "hashed" using ELF
 - b-tree maps hash keys to directory entry locations
- Mcell chain of large directory
 - contains tag.sequence number of the index file
- Mcell chain of the index file
 - identifies the height and root of the B-tree

Displaying the location of the index file



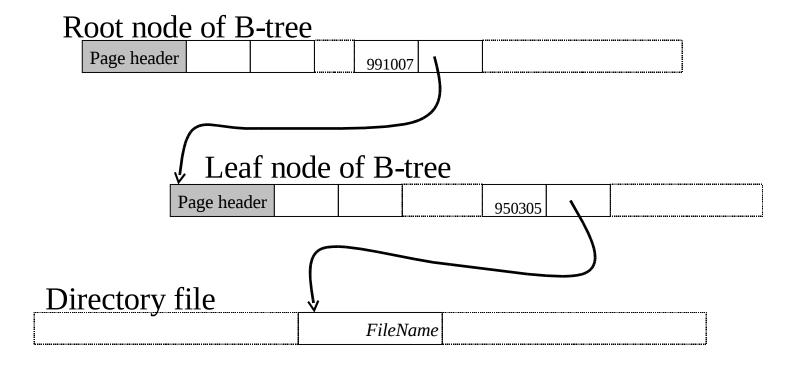
showfile -i LotsOfFiles

```
Id Vol PgSz Pages XtntType Segs SegSz I/O
Perf File
713.8001 1 16 4 simple ** ** ftx
33% index ()
# ls -ld LotsOfFiles
drwxr-xr-x 2 root system 32768 Nov 22 18:20
LotsOfFiles
```

AdvFS directory index b-tree

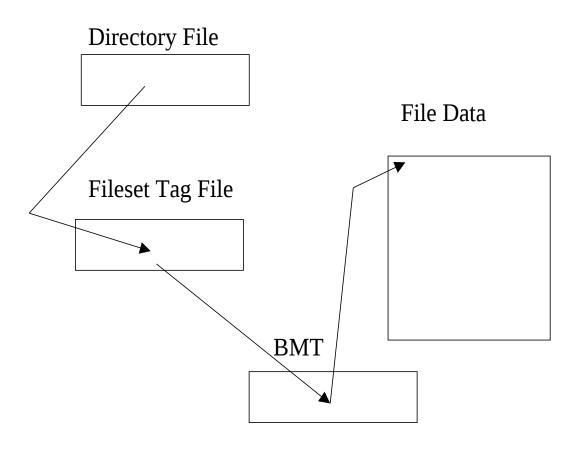


ELF hash function maps file name *FileName* to 950305



Relationship to POSIX files





AdvFS tagfiles and migration



The ability to move:

- A bitfile's data transparently this is supported by the buffer cache and I/O scheduling algorithms
- A bitfile's metadata to another volume tag files enable this

Fragment bitfile



- One per fileset
- Contains small (< 8K) ends of files
- Bitfiles are allocated in 8K (16 sector) units
- If file is less than 100K, a fragment is used if needed.
- Fragment bitfile is broken into a series of 128KB fragment groups

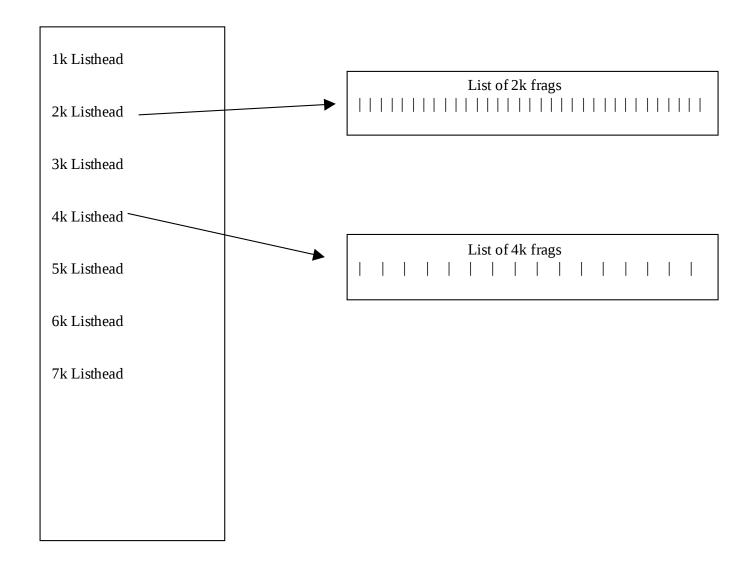
Fragment groups



- Each group consists of fragments of a particular size
 - free (0K), 1K, 2K,...7K
- Each fragment group has the form
 - one page fragment header
 - fragments sized for the group
- Fragment addresses are usually
 - in 1K pages
 - relative to start of fragment bitfile
- Fragments are not Extents!!
 - extent is a contiguous range of 8K pages
 - fragment is a chunk of disk space between 1K and 7K in size

Fragment bitfile locating fragment groups





Fragment header



- Pointer to next fragment group of this type with free space
- Page number of this page (a sanity check)
- Type of this group
- Number of free fragments
- Fileset ID
- Version
- List of free fragments in the group

Fragment group statistics display



nvfragpg -r usr_domain usr

DOMAIN "usr_domain"

reading 438 fra	ag group	headers, 400	headers read
-----------------	----------	--------------	--------------

frag type	free	1K	2K	3K	4K	5K	6K	7K	totals
groups	2	44	67	80	76	63	52	54	438
frags	-	5588	4254	3386	2413	1600	1100	979	19320
frags used	-	5469	4245	3364	2405	1581	1093	973	19130
disk space	256K	5632K	8576K	10.0M	9728K	8064K	6656K	6912K	54.8M
space used	-	5469K	8490K	9.9M	9620K	7905K	6558K	6811K	53.7M
space free	254K	119K	18K	66K	32K	95K	42K	42K	668K
overhead	2K	44K	67K	80K	76K	63K	52K	54K	438K
wasted	-	OK	67K	80K	228K	126K	52K	54K	607K
% used	-	97%	98%	98%	98%	98%	98%	84%	98%

Fragment free list display using nvfragpg (1 of 3)



nvfragpg -fr usr_domain usr

DOMAIN "usr_domain"

reading 438 frag group headers, 100 headers read reading 438 frag group headers, 200 headers read reading 438 frag group headers, 300 headers read reading 438 frag group headers, 400 headers read

frag type	free	1K	2K	3K	4K	5K	6K	7K	totals
groups	2	44	67	80	76	63	52	54	438
frags	-	5588	4254	3386	2413	1600	1100	979	19320
frags used	-	5469	4245	3364	2405	1581	1093	973	19130
disk space	256K	5632K	8576K	10.0M	9728K	8064K	6656K	6912K	54.8M
space used	-	5469K	8490K	9.9M	9620K	7905K	6558K	6811K	53.7M
space free	254K	119K	18K	66K	32K	95K	42K	42K	668K
overhead	2K	44K	67K	80K	76K	63K	52K	54K	438K
wasted	-	0K	67K	80K	228K	126K	52K	54K	607K
% used	-	97%	98%	98%	98%	98%	98%	84%	98%

Fragment free list display using nvfragpg (2 of 3)



head of f	ree lis	ts of	frag gro	oups fro	om file	set attribut	es:		
frag type	BF_FRA	G_ANY	firstFr	eeGrp	6976	lastFreeGrp	32		
frag type	BF_FRA	G_1K	firstFr	eeGrp	6960	lastFreeGrp	560		
frag type	BF_FRA	G_2K	firstFr	eeGrp	6880	lastFreeGrp	6880		
frag type	BF_FRA	G_3K	firstFr	eeGrp	6944	lastFreeGrp	848		
frag type	BF_FRA	G_4K	firstFr	eeGrp	6832	lastFreeGrp	816		
frag type	BF_FRA	G_5K	firstFr	eeGrp	6896	lastFreeGrp	864		
frag type	BF_FRA	G_6K	firstFr	eeGrp	6768	lastFreeGrp	6768		
frag type	BF_FRA	G_7K	firstFr	eeGrp	6864	lastFreeGrp	6864		
any	6976	6992							
free 1K	6960	6912	6928	560					
full 1K	0	112	128	176	224	320	400	496	944
1184									
	1616	1840	2928	3184	3280	3424	3568	3808	3856
3904									
F120	3936	3968	4192	4544	4816	4864	4912	4928	4944
5120	E244	E 4 E C	EE26	E600	E760	6240	CE 20	6600	6640
6848	5344	5456	5536	5600	5760	6240	6528	6608	6640
00-10									

Fragment free list display using nvfragpg (3 of 3)



free 2K	6880									
full 2K	64	208	256	272	288	336	480	624	784	928
	1040	1168	1264	1520	1664	2480	2736	3024	3056	3072
	3088	3104	3120	3200	3360	3504	3792	3824	3872	3888
	3952	4032	4080	4144	4160	4288	4304	4352	4400	4448
	4496	4528	4560	4592	4608	4624	4656	4688	4704	4896
	4976	5168	5360	5488	5552	5584	5616	5664	5824	5984
	6144	6320	6480	6560	6624		6656	5		
()										
free 7K	6864									
full 7K	80	384	512	544	656	704	768	832	912	1024
	1328	1440	1568	1680	1792	2528	2608	2688	2800	2944
	3168	3296	3408	3488	3536	3584	3600	3632	3664	3680
	3712	4048	4224	4432	4784	5072	5184	5200	5232	5264
	5472	5680	5792	5888	6000	6080	6128	6224	6336	6400

Fragments and files



- POSIX stats record of mcells contains
 - fragId field
 - fragId.frag is the page offset of fragment
 - fragId.type is the size of fragment
- Use showfile to determine if there is a fragment
 - does number of pages match file size?
- Use nvbmtpg to find fragment location

Tracking down a fragment (1 of 2)

fragId.type BF_FRAG_7K fragId.frag 2396



```
# ls -i /etc/disktab
1931 disktab
# nvbmtpq -r root domain root -t 1931 -c
DOMAIN "root_domain" VDI 1 (/dev/rdisk/dsk1a) lbn 222464 BMT page 82
pageId 82 megaVersion 4
freeMcellCnt 0          nextFreePg -1          nextfreeMCId page,cell 0,0
CELL 25 next mcell volume page cell 0 0 0 bfSetTag, tag 1,1931
RECORD 2 bCnt 92
                                                            BMTR FS STAT
st_mode 100755 (S_IFREG) st_uid 3 st_gid 4 st_size 30910
st_nlink 1 dir_tag 13 st_mtime Wed May 29 07:06:42 2002
```

Tracking down a fragment (2 of 2)



```
# dd if=/.tags/1 of=/tmp/dtfrag bs=1024 \
                            iseek=2396 count=7
7+0 records in
7+0 records out
# head /tmp/dtfrag
#8192:fc#1024:\
        :od#393216:pd#324698:bd#8192:fd#1024:\
        :oe#717914:pe#324698:be#8192:fe#1024:\
        :of#1042612:pf#324698:bf#8192:ff#1024:\
        :og#393216:pg#819200:bg#8192:fg#1024:\
```

Storage bitmap bitfile characteristics



- Storage bitmap file represents each 8K of on-disk storage within a volume with 1 bit (1K per bit in Tru64 UNIX V4.0)
- Basically a bunch of bits representing whether storage is in use or not
- SBM file is .tags/M-7.
- One per volume
- All storage is either free or in a bitfile
- AdvFS storage is allocated in pages
- On-disk SBM
 - little more than an array of bits (1 bit per page)

SBM format



SBM page header

- sequence number, now unused (32 bits)
- XOR field for rest of page (32 bits)

Bitmap

- 65472 bits
- enough for 8184 pages (8k each)

SBM display through .tags/M-7 (1 of 2)



```
# od -x -N 1000000 /facusers/.tags/M-7
        0000 0000 ffff 0000 ffff ffff ffff
0000000
*
        ffff ffff ffff ffff 0000 0000 0000
0000120
0000140
        0000 0000 0000 0000 0000 0000 0000
0020000
       0000 0000 0321 4ed9 0000 0000 0000 0000
0020020
        0000 0000 0000 0000 0000 0000 0000
       0000 f800 ffff f1ff ffff e7ff ff7f ffff
0033340
0033360
       ffff 9fff 3fff fc00 7fff 8020 ffff f7ff
       dfff ffff ffff 1fcf 7ffc ffff f7ff
0033400
0033420
       Offf fffe dodf fffd fff7 ffff ffff 7ffh
        fff8 ffff ffff 7fff 0000 0020 0000
0033440
0033460
       8000 0002 2040 0000 0040 0000 0000 0002
0033500
        0000 0000 0000 0000 0000 0000 0000
```

SBM display through .tags/M-7 (2 of 2)



showfile -x /facusers/.tags/M-7

SBM page allocation status display using vsbmpg



Miscellaneous bitfile



Characteristics of the misc bitfile includes:

- One per volume
- Holds pages for
 - primary and Secondary Boot Block
 - partition Table (Disk Label)
 - fake UFS Super Block with AdvFS Magic Number

Learning check





Lab 2





