

set ('n ~~test~~)LISP

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means binding a variable
name n with value ~~test~~
or n.

What is Lisp? Explain the structural organization of Lisp with a suitable example.

Lisp is a programming language developed in 1950s out of need for Artificial Intelligence programming.

where the pointer and linked list structures are its natural data-structuring methods.

Lisp was first implemented in IBM 709 and allowed 6 or 3 overlapping programming paradigms.

Structural Organization of Lisp.

1. Lisp is ^{extremely} applicative language.

Lisp is an extremely applicative language. as almost everything is a function in Lisp. It also represents programs and data in same way.

In Lisp function written as:

(f a₁ a₂ ... a_n)

where f is the function and a₁, a₂, ... a_n

operator before
operands.

{ Cambridge polish notation }

'text means passing -
"text" rather than value held by
variable text.

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are the arguments.
treating program as data

this notation is Cambridge polish notation where
operators before operands
ie (plus 5 2)

lisp would respond 5.

Also (set 'text (make to do)) is equivalent to
Algol set (text, make (to, do))

2. In lisp list is the primary data structure
Constructor.

→ In lisp list is the primary data
structure means lisp treats both program
and data in the same way. in the form of
lists.

This means lisp's focus is on allowing programmers
to manipulate list of data. which includes
comparing them, passing them to functions, putting
(set 'text to be

them together and taking them apart.

Eg.

set 'text '(to be or not)

(set 'text '(to be or not))

this means variable text is set the list
composed of values/atoms to, be, or, not.

3) programs are represented as list:

In Lisp, functional applications and lists look the same.

Eg in S-expression:

(make-table text nil)

could either be a 3 element list or it could be an application of the function make-table to which arguments named text and nil are passed. So which is it?

In Lisp its both & since Lisp program itself is a list.

Under most circumstances, an S expression is interpreted as function application which means that arguments are evaluated and function is invoked. However, if the list is quoted, then it is treated as data that is it is unevaluated.

Eg

(set 'text '(to be or not to be))

here set function is passed a "text" as data followed by a list or ~~data~~ "data s".

If, the above statement was written as

(set 'text (to be or not to be))

means the set func is equivalent to algol.

→ set (text, to (be, or, not, to, be))

47 Lisp is Interpreted.

Lisp system has an interactive interpreter.

i.e. lisp interprets a statement and prints out results.

Eg

(plus 2 3)

gives 5

(eg (plus 2 3) (difference 9 4))

the system responds with

5

which means true as $2+3 = 5$.

eg and plus are pure functions while set is a pseudo function.

Eg (set 'text '(to do or to be))

defun is a pseudofunction.

(defun f (n₁ n₂ ... n_m) B)

that defines a function with name f and formal parameters n₁, n₂ ... n_m;

5.7 An lisp program example:

```
(defun make-table (text table)
  (if (null text)
      table
      (make-table (cdr text)
                   (update-entry
                    table (car text)))))
```

```
(defun update-entry (table word)
  (cond ((null table) (list (list word)))
        ((eq word (car table))
         (cons (list word (add1 (cadr table)))
                 (cdr table)))
        (t (cons (car table)
                  (update-entry (cdr table)
                                word))))))
```

```
(defun lookup (table word)
  (cond ((null table) 0)
        ((eq word (car table)) (cadr
                                table))))
```

(t (lookup (code table) word)))

(set 'text '(to be or not to be))

(set 'Reg (make-table text nil))

←

how lisp maintained the simplicity principle?
Explain different searching techniques in lisp.

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differentiate cde and cse.

→ cde and cse are used to access parts of lists in lisp.

→ A list requires operations for building the structure and operations for taking them apart.

hence lisp has operations that build structures such as called constructors.

while those that extract their parts are called selectors.

List has one constructor as - cons - and two selectors - car and cdr.

Car and cdr are pure functions as they don't modify argument's list.

car

car is used to select first element of a list.

Eg (car '(to be or not to be))

returns atom to.

The first element can either be an atom or a list, and car returns it, whichever it is. For example, let frog be the list

((to 2) (be 2))

(car frog) (not 1))

the application

(car frog)

returns the list
(to 2)

The argument of car is always a non null list, while it may return an atom or a list depending on what its argument's first element is.

Cde

car can access only first element (either atom or list).

so cde is a selector that is used to access rest of the elements except the first.

eg: (cde '(to do not be))

returns list: (do not be)

cde also like car requires non-null argument but unlike car, cde ~~can~~ always returns list. in case of only one argument it returns a null list.

eg: (cde '())
returns ().

About Cde and Car:

Both don't modify their argument list (hence pure functions) i.e. they work on a copy of list provided to them as argument.

eg: cde doesn't delete first element of its argument list. rather, it returns a new list exactly like its

argument except its first element.

In Lisp `car` and `cdr` can be used in combination to access the components of a list.

Eg. `(set 'DS '((Don Smith) 45 3000 (August 25 1980)))`

The list `DS` contains Don Smith's name, age, salary and hire date.

to access his name (`car DS`) can be used which returns `(Don Smith)`.

while to access Don Smith's ^{hire date} ~~salary~~ the below statement can be used.

`(car (cdr (cdr (cdr DS))))`
4 3 2 1

hence,

① `cdr` ~~returns~~ ^{returns} `(45 3000 (August 25 1980))`

② `cdr` ~~returns~~ ^{returns} `(3000 (August 25 1980))`

③ `cdr` ~~returns~~ ^{returns} `((August 25 1980))`

④ `car` gives us `(August 25 1980)`.

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how lisp maintained its simplicity principle. also explain different searching techniques in lisp.

→ In lisp two selectors car and cdr are pure functions are adequate for accessing the components of any list structure. due to this lisp maintained its simplicity principle.

Along with this they provided abbreviation to prevent long and large compositions of cars and cders.

Eg:

(car (cdr (cdr (cdr DS))))

could be abbreviated to

[DS is the same list as on 4]

(caddr DS).

write

Also it is observed that these abbreviations can also be error prone.

So in lisp the solutions for this is to write special purpose functions for accessing the parts of a record.

Eg:

(define hire-date (lambda (addd e) e))

the function (hire-date DS) returns Don Smith's hire date or date of any list passed of any other personnel records with same structure.

This makes lisp program maintainable ~~why~~ while lisp maintains its simplicity.

lisp thinks of personnel records as an abstract data type that can only be accessed through the provided function.

(accessing)

→ Searching techniques in lisp are cde, ~~etc.~~ [Note here explain those]

and just like above; abbreviations of lisp are presented in form of ~~single~~ linked data structures.

Eg. (caddadde DS) where DS is a list of DS's

((Don Smith) 45 30000 (August 25 1980))

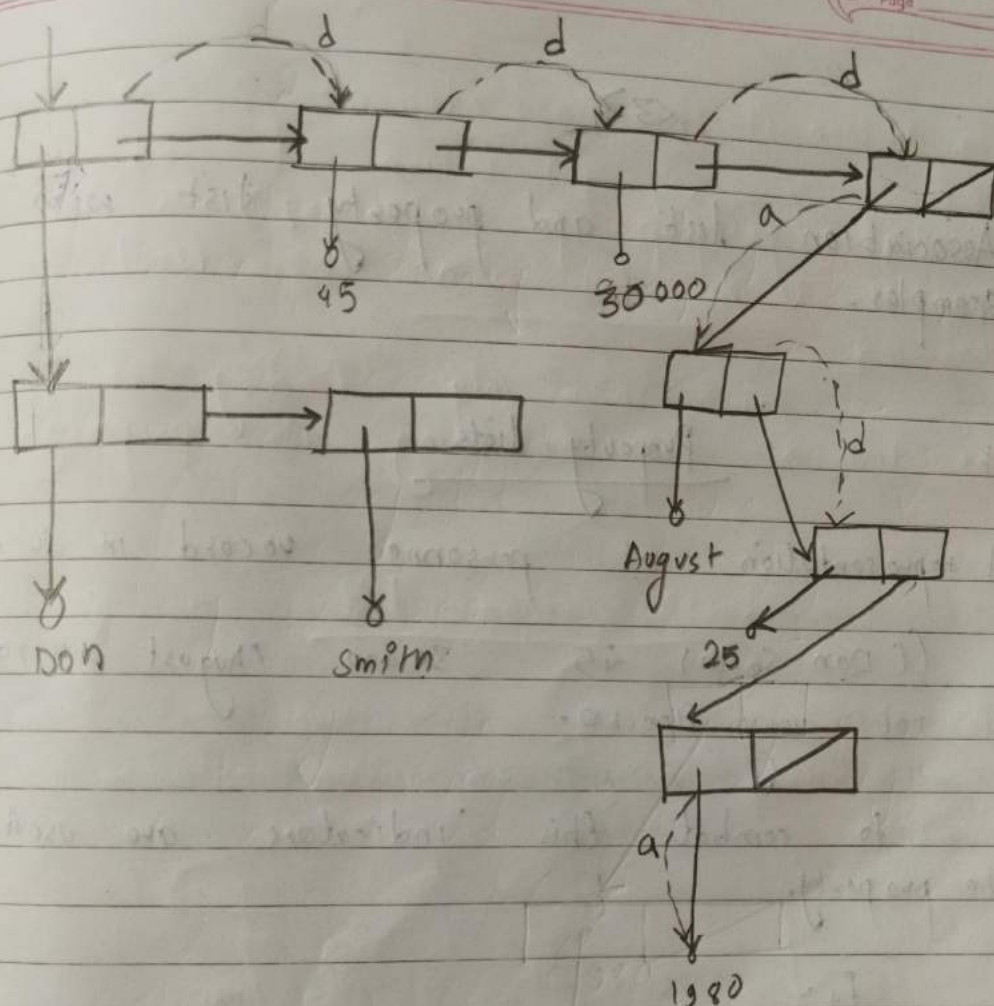


Fig. Walking down a list structure.

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Association list and property list with relevant examples.

Property lists

A representation of personnel record in form.

((Don Smith) 45 30000 (August 25 1980))
is not very flexible.

to combat this indicators are used to identify the property.

Eg:

Name (Don Smith) age 45 salary 30000 hire-date
(August 25 1980))

This method of representing information are called Association list,

general form

$(p_1 v_1 \quad p_2 v_2 \quad \dots \quad p_i v_i)$
 $p = \text{property indicators}$
 $v = \text{value of property}$

Advantages of property is flexing.

In lisp these property list can be accessed with the following code or something familiar

```
(defun getprop (p x)
  (if (eq (caddr x) p)
      (caddr x)
      (getprop p (caddr x))))
```

the above lisp code uses recursion to get values based on indicators.

where p is the ~~property~~ indicator or ~~get~~ the property we need from a property list

```
(getprop 'name DS)
→ (Don Smith)
```

```
getprop 'hire-date DS)
→ (August 25 1980).
```

Explanation of how the getprop works :

If the first element of x is the ~~property~~ indicator p then (eq (caddr x) p) returns a 't' and if statement gets validated

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to return ~~it~~ the first element; else the
(eg (car x) p) returns p and else
return (get prop p (addr x)) is called on
the list x but with the first element
removed.

hence recursively finding the indicator in
the list x and returning its property.

Association list

a data structure

A property list data structure is where one
atom or list value is associated with a property.

eg (p.v . . . p.v)

this is inconvenient at times where the
data structure has flags.

In case of flags its presence itself
is the value or information and requires not be in
form. to not be in ~~prop~~ indicator - property

eg
DS indicates retired atom in the list
the personnel being senior

DS => (name (Don Smith) age 45 salary 30000
hire-date (August 25 1980) retired)

here retired flag breaks the p-list properties and could be fixed by making the flag p-list inflexible to flags.

These inflexibility can be solved by an association list where:
it could be in form.

(name (Don Smith)) (age 45) (salary 30000) (hire-date (August 25 1980)) (retired))

The Alist is in form.

$((a_1, v_1) (a_2, v_2) \dots (a_n, v_n))$

is attribute value pairs.

Statement:

(assoc 'hire-date DS)
→ (August 25 1980)

An assoc function is used in forward association.

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how does cde and car help in searching data elements? explain with walking down diagram.

Qn 4 (about car and cde)

Searching techniques in lisp part of Qn 2 write about abbreviation and its walking down diagram equivalent.

nil is an atom
in lisp

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polish notation

→ cons in $O(n)$

how hierarchical structures processed in lisp?

→ In lisp hierarchical structures difficult to be handled by iterations are handled Recursively.

An example would be an equal function implemented by using pro function of that returns 't' or 'nil'

Example :

```
(defun equal (x y)
  (or (and (atom x) (atom y) (eq x y))
      (and (not (atom x)) (not (atom y))
            (equal (car x) (car y))
            (equal (cdr x) (cdr y)))))
```

here

(atom x)
returns t if x is an atom if - a list returns
nil

defun defines a function equal with two parameters
x y.

~~for~~

for means if any of its elements in the list
it lies is 't' it returns t ~~and~~
~~stops computing~~. else returns nil. or returns
as soon as it encounters 't' in them.

① (and (atom x) (atom y) (eq x y))

here and means if ~~the list~~ all elements besides
it self is 't' it returns a 't'.

So the statement ① means, if the arguments
passed to the equal function ~~are~~ ~~also~~ atoms,
then compare them and if their equal return
t else nil.

statement

② (and (not (atom x)) (not (atom y))
means when both x and y arguments
are not atoms return 't' else nil.

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Explain conditional Expression, the logical connectives and Mapcar and reduce Functions.

Conditional Expression

- Lisp is first language to contain a conditional expression.
- in previous languages 'Fortran' or newer 'Pascal' required the user to drop from the expression level to the statement level in order to make a choice.
- this was not efficient when representing mathematical expressions and equations.
- hence cond was introduced in Lisp to allow for conditional expression.

Eg
$$sg(x) = \begin{cases} 1 & \text{if } x > 0 \\ 0 & \text{if } x = 0 \\ -1 & \text{if } x < 0 \end{cases}$$

In Lisp is defined as:

```
(defun sg (x)
  (cond ((plusp x) 1)
        ((zerop x) 0)
        ((minusp x) -1)))
```

Format (cond p1 ... pnen)

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(null L)
if L is null
if L is null => t
else not.

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Logical connectives

Logical connectives are used to combine logical statements into more complex statements.

The basic logical connectives are AND, OR, NOT, and IMPLIES.

AND is represented by the symbol \wedge . It is true only when both statements are true.

OR is represented by the symbol \vee . It is true when at least one of the statements is true.

NOT is represented by the symbol \neg . It is true when the statement it negates is false.

IMPLIES is represented by the symbol \Rightarrow . It is true unless the first statement is true and the second is false.

Truth tables are used to determine the truth value of a logical expression for all possible combinations of the truth values of the individual statements.

Mapcar

It is a function that applies a given function to each element of a list and returns a list of results.

implementation:

```
(defun mapcar (f x)
  (if (null x)
      nil
      (cons (f (car x)) (mapcar f (cdr x)))))
```

)))

Reduction Function :

the implementation of plus-red is an example of reduction function.

~~that it is~~

plus-red is a binary function to reduce a list to a single value.

```
(defun plus-red (a)
  if (null a)
    0
    (plus (car a) (plus-red (cdr a)))
  )
```

Example

(plus-red '(1 2 3 4 5))
→ 15

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Short notes:

i) User defined function in Lisp:

User defined function in Lisp are done by using pseudo function defun which defines a function.

(defun f (n_1 n_2 ... n_m) b)

This statement defines a function f , with formal parameters (n_1 n_2 ... n_m) and a body b .

Example.

~~(defun sum (x) &~~

(defun add1-map (a)
 (if (null a)
 nil

~~(add1~~

(cons (add1 (cadr a))

(add1-map (cddr a))))))

~~The add1-map~~

(add1-map '(2 4 6 8))

3 5 7 8

→ i.e. adds 1 to all elements of list.

ii) Cde 'cde Function
← On 4.

~~23rd Storage Reclamation~~

iv) Recursive Interpreters:

A Lisp Interpreter can be written in Lisp. This is done by a universal function. A universal function is a function that can interpret any other function.

Since this universal function is written in Lisp, it can make use of the facilities of Lisp, such as cde, car, and cons.

Which in turn are used to interpret cde, car, and cons of Lisp itself.

If a recursive Interpreter is being made in any other language other than Lisp, it should have recursive procedures and ability to implement Linked List.

The Lisp universal function is called eval.

and in addition to the expression to be evaluated ~~eval~~ must have second parameter. which is a data structure representing the context in which the evaluation to take place.

eg.

$(\text{eval } 'E \ A) = A.$

Eg:

$(\text{eval } ! (\text{cons } (\text{quote } A) (\text{quote } (B C D))) \text{ nil})$

$\rightarrow (A \ B \ C \ D)$ where nil is empty context.

this is equivalent to ~~$(\text{cons } ('A) ('(B C D)))$~~
 $(\text{cons } ('A) ('(B C D)))$
 $\rightarrow (A \ B \ C \ D)$

here recursive interpreter has implementations to convert quote to ' before execution.

<10> Evaluate the lisp expressions:

$$a) \frac{1}{2i} \sqrt{x^2 - 1^2}$$

$$= \frac{(\text{times } (\text{quotient } 1 \ 2) (\text{sqrt } (\text{difference } (\text{times } (\text{pi } (\text{expt } 4 \ 2)))) (\text{expt } 1 \ 2)))}{1}$$

$$\frac{-b \pm \sqrt{(b^2 - 4ac)}}{2a}$$

$$\Rightarrow \frac{(\text{quotient } (\text{difference } (\text{minus } b) (\text{difference } (\text{expt } b \ 2) (\text{times } 4 \ a \ c))))}{2a}$$

$$= \frac{(\text{quotient } (\text{difference } (\text{minus } b) (\text{sqrt } (\text{difference } (\text{expt } b \ 2) (\text{times } 4 \ a \ c))))}{(\text{times } 2 \ a)}$$

$$c) (-1)^k k^{(1/k)}$$

Seen!
fist gring

$$\rightarrow \frac{\text{times (mins)}}{\text{times (expt (minus 1) k) (expt k (quotient 1 k))}}$$

$$d) \frac{n!}{4!(n-4)!}$$

$$\rightarrow \frac{(\text{quotient } n)}{n} (\text{fact } n) (\text{times (fact 4) (difference$$

$$= \frac{(\text{quotient } (\text{fact } n) (\text{times (fact 4) (fact (difference } n \text{ 4) })))}{1}$$

$$e) \frac{-b \pm \sqrt{b^2 - 4ac}}{2\sqrt{4a^2 - b^2}}$$

opening closing of same function.

$$\Rightarrow \frac{(\text{quotient (difference (minus b) (sqrt (difference$$

$$(\text{expt } b \text{ 2) (times } 4 \text{ a a c) (sqrt (difference (times 4 (expt a) (expt b 2)))))$$