

Chapter - I.

Goals in Problem Solving :-

Fields / Application # gnp

Goal :

- Target / Desired result / Accomplishment
What is to be achieved?

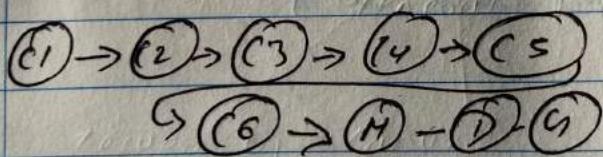
Planning

- Set / Sequence of tasks to achieve a goal.

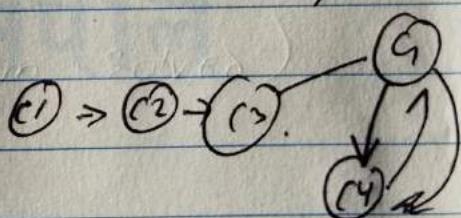
Types of Planning :- (gnp)

- 1) linear planning
- 2) non-linear planning

Linear Planning



Non-linear Planning



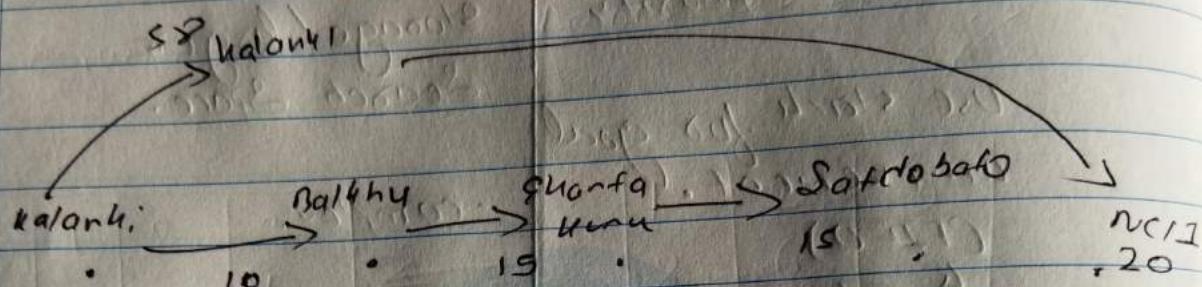
- planning algorithm that works on one goal until completely finished before moving on to next goal.
- works to achieve goal by interleaving (alternating) on up goals
- Considered Priority.

Preemptive \Rightarrow (switches) ~~if~~ when higher priority task occurs
 Non-preemptive (doesn't switch)

S.NO.	TOPICS	PAGE NO.
	- doesn't prioritize. Uses stack for Goal task / no priority Achievement. { {1,13}, {2,12}, {3,13}, {4,43,95,13} }	
-	- doesn't prioritize - Requiere more task / no priority. Storage / larger search space.	
-	- Use stack for goal achievement. - complex. (FIFO) ↳ a/c/b/c>/c/y/l/s - may produce off-path [AI not FIFO is not mal soln. stack]	
-	- complete.	
-	- Requires less storage / Reduced search space.	
-	- simple	
-	- may not produce	
-	- Incomplete.	

MEA (means-End Analysis) (Contd.)

[initial state & goal ~~final state~~ final state
 find best intermediate state choose steps towards MEA



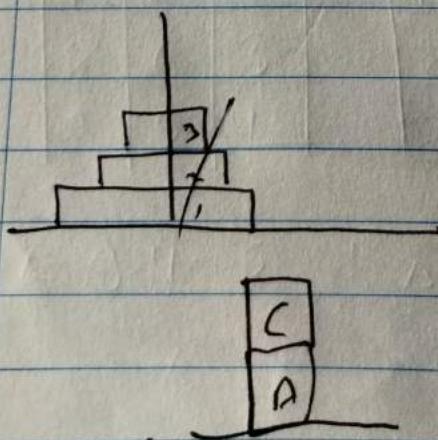
1) Current State

2) Final state

3) Goal difference (Δ)

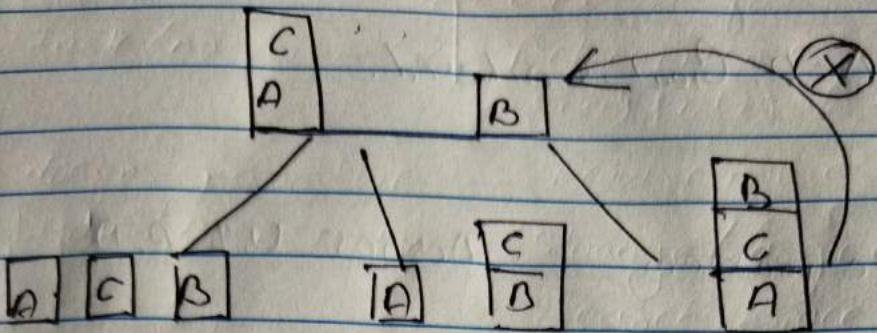
GPS (General Problem Solver)

Towers of Hanoi



A	-1
B	-1
C	-1

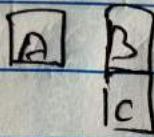
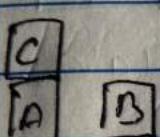
G.S



$$\Delta = 3 - 1 = 2$$

$$\Delta = 3 - 0 = 3$$

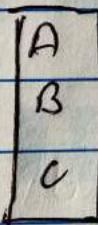
$$\Delta = 3 - 1 = 2$$



$$\Delta = 3$$

$$\Delta = 3 - 2 = 1$$

X



$$\Delta = 0$$

Algorithm that works based on 3 parameters:

1) Initial State / Current

2) Final State / Goal.

3) Goal difference (Δ)

- Reducing goal difference between I.S and G.S until goal

difference (Δ) is 0.
- When $\Delta = 0$, Goal is achieved.

8 puzzle problem:

Actions: U, L, R, D.

1	2	3
8	-	4
7	6	5

Goal state

1	2	3
8	6	-
7	5	4

Initial state

1	2	3
8	6	-
2	5	4

Down

Up

Right

1	2	-
8	6	3
7	5	4

1	2	3
8	6	4
2	5	-

1	2	3
8	-	6
2	5	4

$$\Delta = 9 - 4 = 5$$

$$\Delta = 9 - 8 = 1$$

$$\Delta = 9 - 6 = 3$$

X

S.NO.

TOPICS

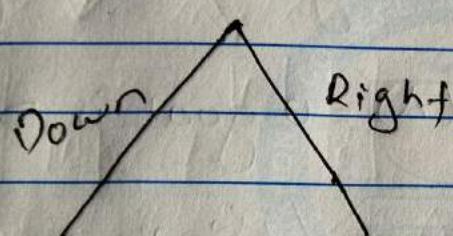
PAGE NO.

State Space Tree

-Tree which shows all possible states from current state.

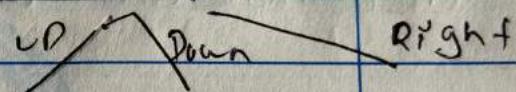
1	2	3
8	6	4
2	5	-

1	2	3
8	-	6
2	5	4



1	2	3
8	6	-
2	5	4

1	2	3
8	6	4
2	-	5

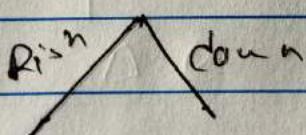


1	2	3
8	5	6
2	-	4

1	2	3
8	-	6
2	5	4

$$\Delta = 4 \quad \Delta = 5 \quad \Delta = 5$$

(Initial state
1 2 3 4 5 6
so discard)



1	2	3
8	6	4
-	2	5

$$\Delta = 0$$

$$= 3$$

X

Algorithm (v.v.g)

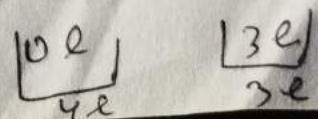
Men Algorithm:

- 1) Define initial state, goal state and calculate goal difference (Δ)
- 2) Choose action/procedure with least goal difference that will ultimately reach to goal.
- 3) When goal difference (Δ) = 0 goal is reached.

Production Rule System / Simple reflex System
[~~if~~ condition \rightarrow Action] or Action performed by production rule system

[Pattern / Precondition \rightarrow Action]

S.NO.	TOPICS	PAGE NO.
	water Jug problem.	Possible actions:
	 $\begin{matrix} & & \\ \text{4l} & & 3\text{l} \end{matrix} \rightarrow \begin{matrix} & & \\ 2\text{l} & & \end{matrix}$	1) Fill 4l Jug 2) Fill 3l Jug 3) Empty 4l Jug 4) Empty 3l Jug 5) Pour 4l to 3l 6) Pour 3l to 4l
	$\begin{matrix} & 0 & 0 \\ 0 & 3 & \\ 3 & 0 & \\ 3 & 3 & \\ 4 & 2 & \end{matrix}$	$3\text{l full} \leftarrow$ $3\text{l not full} \leftarrow$ 4 l Jug full 4 l Jug not full
#	 $\begin{matrix} & & \\ \text{4l} & & 3\text{l} \end{matrix}$	$x = \text{Quantity of water in 4l Jug}$ $y = \text{" " of water in 3l Jug}$
1)	Fill 4l Jug (F4) $x < 4 \rightarrow (x, y)$	
2)	Fill 3l Jug (F3) $y < 3 \rightarrow (x, y)$	
3)	Empty 4l Jug (E4) $x > 0 \rightarrow (0, y)$	 $\begin{matrix} 2\text{l} \\ \text{4l} \end{matrix} \quad \begin{matrix} 0\text{l} \\ \text{3l} \end{matrix}$

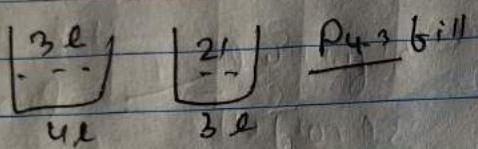


4) Empty 3e jug (E3)

$$y>0 \rightarrow (x, 0)$$

5) Pour water from 4e to 3e to fill it ($P_{4-3}(f^{\text{ill}})$)

$$x>0 \text{ and } x+y \geq 3 \rightarrow (x-(3-y), 3)$$



6) Pour water from 4e to 3e to not fill ($P_{4-3}(\text{not fill})$)

$$x>0 \text{ and } x+y < 3 \rightarrow (0, x+y) \quad \boxed{⑥}$$

7) Pour 3e to 4e to fill it $P_{3-4}(f^{\text{ill}})$

$$y>0 \text{ and } x+y \geq 4 \rightarrow (4, y-(4-x))$$

8. Pour 3e to 4e to not fill $P_{3-4}(\text{no fill})$

$$y>0 \text{ and } x+y < 4 \rightarrow (\cancel{x+y}, 0) (x+y, 0)$$

S.NO.	TOPICS	PAGE NO.
	$x = 2 - e$ $y = n$	
	$x = n$ $y = 2e$	$n = 50, 1, 2, 3, 74$
Goal State.	$x = 0$ $y = 0$	$n = 50, 1, 2, 3, 74$
	I.S	
	$x = 4$ $y = 0$	
	$x = 0$ $y = 3$	
	F_4	F_3
	F_3	F_4
$P_{4-3}(F)$	$x = 1$ $y = 3$	$P_{3-4}(NF)$
	$x = 4$ $y = 3$	
	$x = 3$ $y = 0$	
	E_3	F_4
	$x = 1$ $y = 0$	
	$x = 4$ $y = 3$	
	X	
	$(\text{Since we get } E_3 < E_4)$	
	$P_{4-3}(NF)$	
	$x = 0$ $y = 1$	
	F_4	
	$x = 4$ $y = 1$	
	$P_{4-3}(F)$	
	$x = 4$ $y = 3$	

Final.

$$\boxed{3l} + \boxed{5l} \rightarrow \boxed{4l}$$

x	y	x	y	x	y
0	0	0	0	0	0
0	5	3l	0	0	5
0	2	0	3l	3	2
3	2	3l	3l	0	2
0	2	1l	5l	2	0
2l	0	1l	0l	2	5
		0	1l.	3	4.
		3l	1l		
		0	4l.		

Steps to solve:

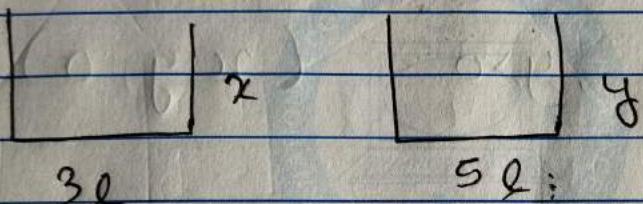
S1: State description two integers (x, y)
 $\{x = 0, 1, 2, 3\}, \{y = 0, 1, 2, 3, 4, 5\}$.

S2: Describe initial & Goal state
initial = $(0, 0)$, Goal $(0, 4)$

Possible Actions

S.NO.	TOPICS	PAGE NO.
①	Fill 3L Jug	
②	Fill 5L Jug	
③	Empty 3L Jug	
④	Empty 5L Jug	
⑤	Pours 3L to 5L	5L jug full 5L jug not full
vi	Pours 5L to 3L	3L jug full 3L jug not full.

Solving using production rule system.



① Fill 3L Jug (F3)

$$x < 3 \rightarrow (3, y)$$

② Fill 5L Jug (F5)

$$y < 5 \rightarrow (x, 5)$$

③ Empty 3L Jug (E3)

$$x > 0 \rightarrow (0, y)$$

④ Empty 5L Jug (E5)

$$y > 0 \rightarrow (x, 0)$$

x
y

(5) Pour 3L to 5L (se not full) P₃₋₅NP

$$x > 0, \quad x+y \leq 5 \quad (0, x+y)$$

(6) Pour 3L to 5L (5L full) P₃₋₅F

$$x > 0, \quad x+y \geq 5 \quad (x-(5-y), 5)$$

(7) Pour 5L to 3L (3L not full) P₅₋₃(N-F)

$$x+y \leq 3, \quad y > 0 \quad (x+y, 0)$$

(8) Pour 5L to 3L (3L full (P₅₋₃F))

$$x+y \geq 3, \quad y > 0 \quad (3, y-(3-x))$$

(b) Drawing State Space tree

$$x = 0$$

$$y = 0$$

let x denote the quantity of water in 3L Jug
let y denote the quantity of water in 5L Jug.

S.NO.	TOPICS	PAGE NO.
	$x=0$ $y=0$	10.50.
	$x=3$ $y=0$	
	$x=0$ $y=5$	
	$x=3$ $y=5$	
	$x=0$ $y=3$	
	$x=3$ $y=2$	
	$x=3$ $y=3$	
	$x=1$ $y=2$	
	$x=3$ $y=2$	
	$x=3$ $y=2$	
	$x=0$ $y=2$	
	$x=2$ $y=0$	
	$x=2$ $y=5$	
	$x=3$ $y=4$	goal

F3 FS

P3-S(NF) P5-3 (ADNF)

F3 F3

P3-S(NF) P5-3 (NF)

F3 F3

P5-3 (NF)

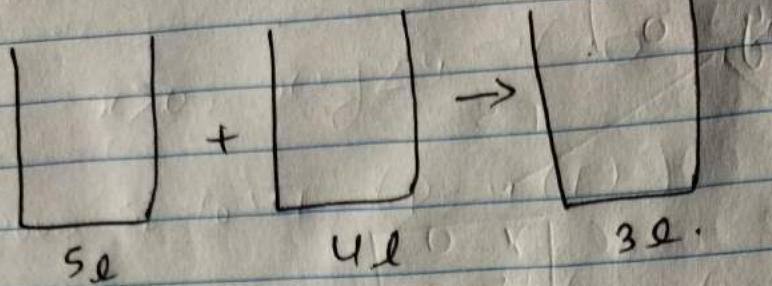
E3

P5-3 (NF)

FS

P5-3 (NF)

Solve #.



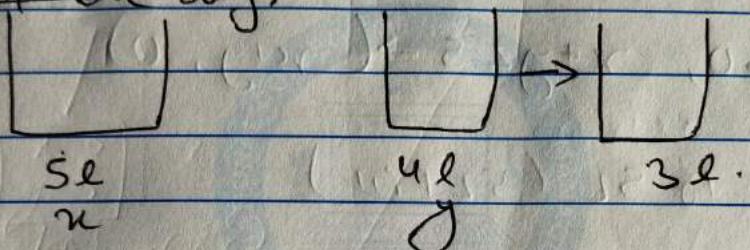
Possible actions:

0	4	0	0
5	0	4	0
1	4	0	0
1	0	3	1
0	1	2	2
5	1	0	3
2	4	1	2
0	0	4	0
2	4	0	0
4	0	0	4
4	4	0	0
5	3	0	0

Possible actions?

- 1) Fill 5l Jug
- 2) Fill 4l Jug
- 3) Empty 5l Jug
- 4) Empty 4l Jug

S.NO.	TOPICS	PAGE NO.
5)	pour 5l to 4l → 4l full → 4l not full	
6)	pour 4l to 5l → 5l full → 5l not full	
a)	now solving using production rule system.	
b)	Fill 5l jug.	



let x denote the quantity of water in 5 liters

let y denote the quantity of water in 4 liters.

1) Fill 5l jug.

$$x \leftarrow 5 \rightarrow (5, y)$$

2) Fill 4l jug.

$$y \leftarrow 4 \rightarrow (x, 4)$$

3) Empty 5l jug.

$$x \leftarrow 0 \rightarrow (0, y)$$

4) Empty 4l jug.

$$y \leftarrow 0 \rightarrow (x, 0)$$

5. Pour SL to UL (ULNF)

$$x > 0, x+y \leq 4 \rightarrow (0, x+y)$$

6. Pour SL to UL (PS-4(F))

$$x > 0, x+y \geq 4 \rightarrow (x-(4-y), 4)$$

7. Pour UL to SL (SLNF)

$$y > 0, x+y \leq 5 \rightarrow (x+y, 0)$$

8. Pour UL to SL (SLFII)

$$y > 0, x+y \leq 5 \rightarrow (5, y-(5-x))$$

Now,

Drawing state square tree.

$$x = 0$$

$$y = 0$$

let x denote the quantity of wafers in SL
let y denote the quantity of wafers in UL

farmer Fox Goose Grain problem:-

River crossing:

Left: Fa, Fo, Gu, Go

Right: Fa, Fo, Go, Gu
[Goal]

Possible actions:

- i) Fa X
- ii) Fa, Go ✓
- iii) Fa, Fo X
- iv) Fa, Go ✗ [Fox will eat Goose]

Drawing State Space tree:

L: Left side of River

R: Right side of River

A farmer has a fox, a goose and a grain. He comes to a river he must cross, and there is a boat but it can only hold him and one other thing at a time.

ge No.

He cannot leave the fox with goose, because the fox will eat the goose. ~~that too~~.

He cannot leave the goose with the grain, beca

S.NO.	TOPICS	PAGE NO.
1	$\begin{array}{ c } \hline L: fa, fo, ga, ga \\ \hline R: fa \\ \hline \end{array}$	
1	$\begin{array}{ c } \hline L: fa, fo, ga \\ \hline R: fa, go \\ \hline \end{array}$	
2	$\begin{array}{ c } \hline L: fa, fo, ga \\ \hline R: fo, ga \\ \hline \end{array}$	2
3	$\begin{array}{ c } \hline L: ga \\ \hline R: fa, fo, go \\ \hline \end{array}$	
3	$\begin{array}{ c } \hline L: fo \\ \hline R: go, fa, ga \\ \hline \end{array}$	3
4	$\begin{array}{ c } \hline L: fa, go, ga \\ \hline R: fo \\ \hline \end{array}$	
4	$\begin{array}{ c } \hline L: fo, fa, go \\ \hline R: ga \\ \hline \end{array}$	4
5	$\begin{array}{ c } \hline L: go \\ \hline R: fo, fa, go \\ \hline \end{array}$	
5	$\begin{array}{ c } \hline L: go \\ \hline R: fo, fa, go \\ \hline \end{array}$	5
6	$\begin{array}{ c } \hline L: fo, go \\ \hline R: fo, ga \\ \hline \end{array}$	6
7	$\begin{array}{ c } \hline L: \\ \hline R: fa, fo, ga, ga \\ \hline \end{array}$	7

Path cost = ?

Total no. of states = 8

Missionaries and cannibals problem :-

3 missionaries and 3 cannibals were on one side of the river.

↳ All want to cross the river.

↳ On same side of river missionaries count can't be less than cannibals.

↳ Only one boat available that can hold only two people at a time.

Possible actions:

M1

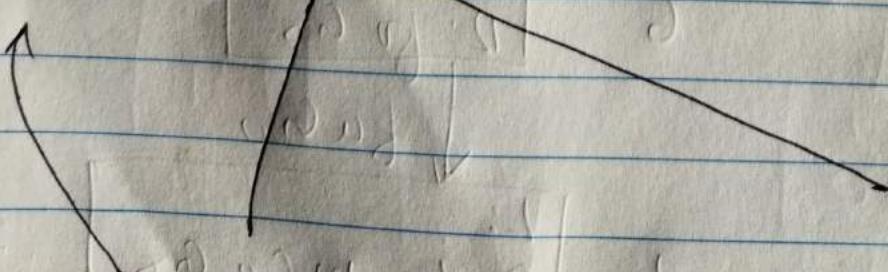
C1

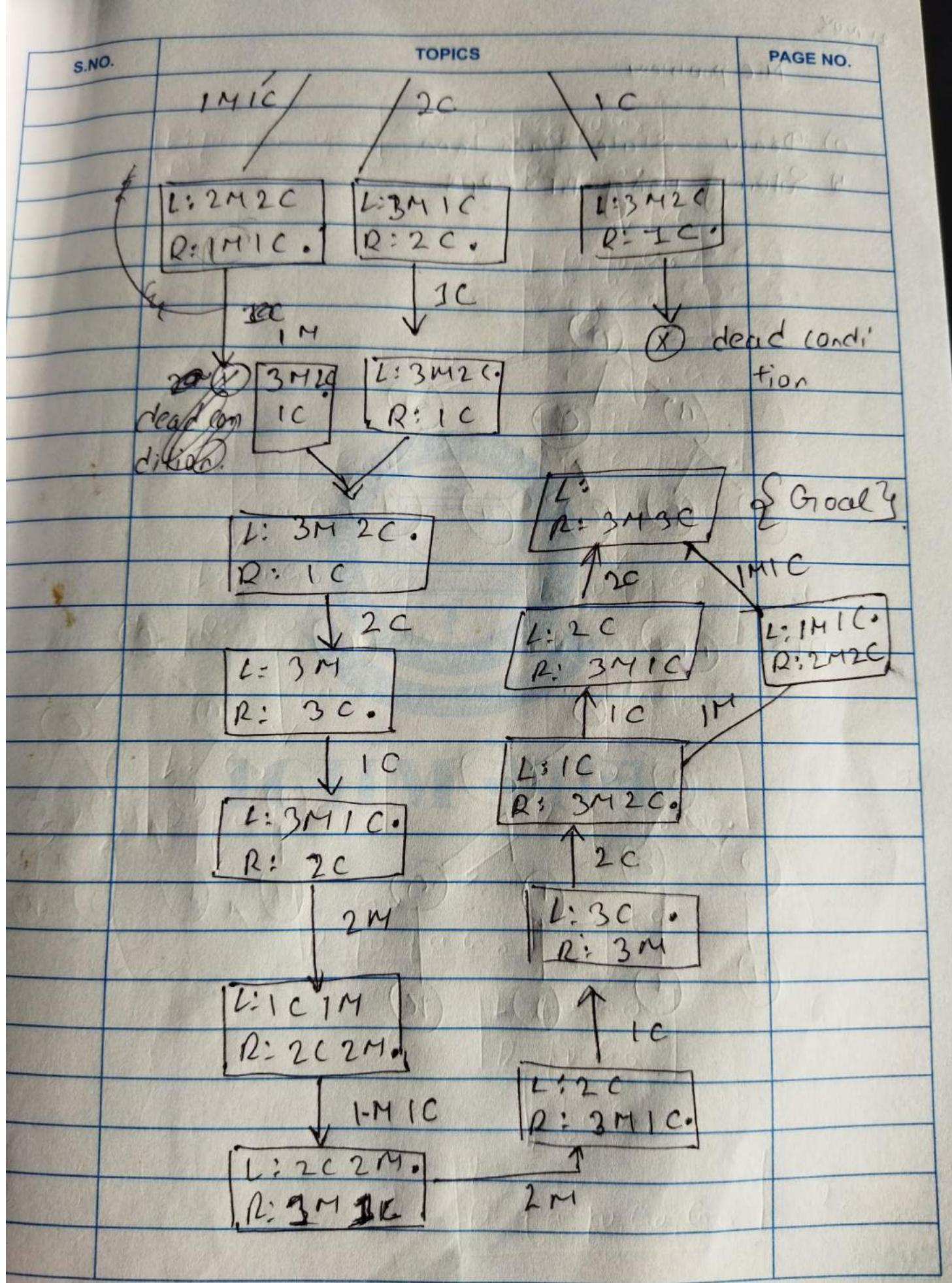
M1C1

2M

2C

L: 3M 3C .
R: []

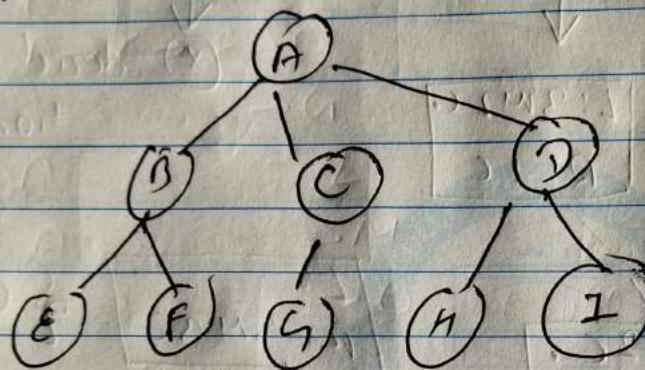




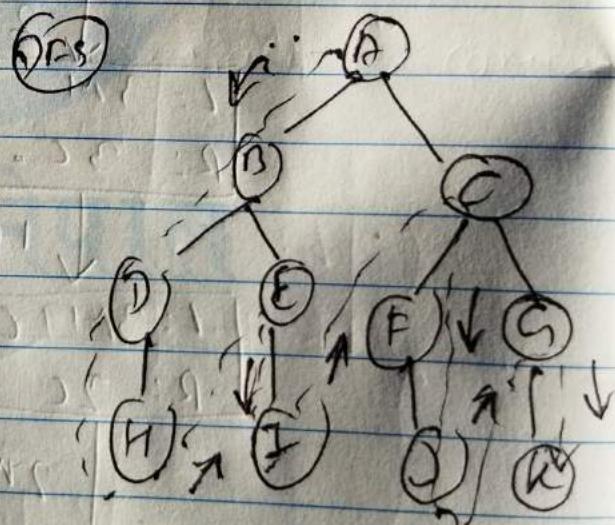
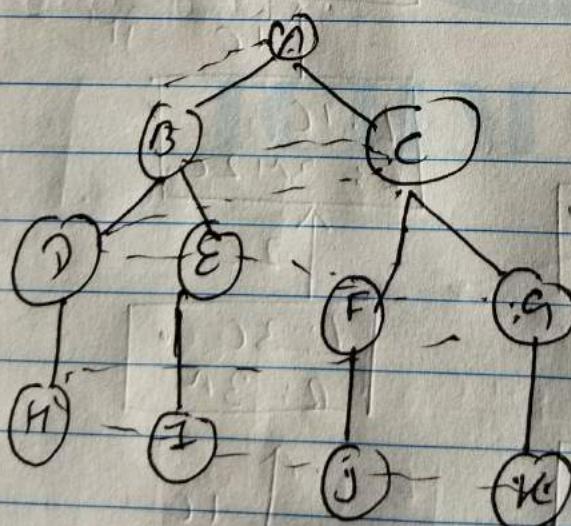
21 Nov

M-C Problem²

- a) Draw State space tree.
- b) Solving using BFS/DFS.



DFS: BFS:

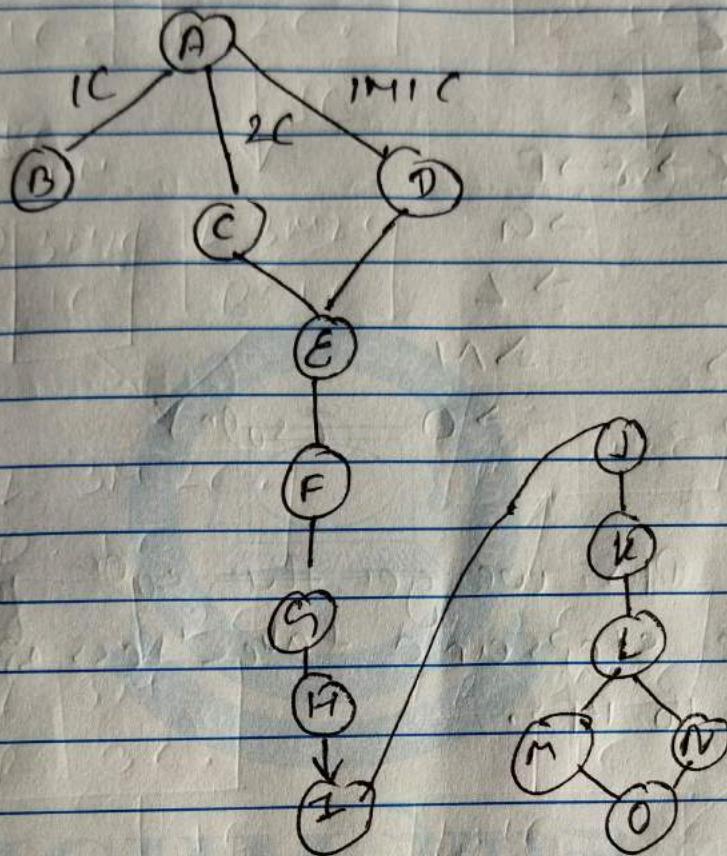


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TOPICS

PAGE NO.

BFS/DFS of Missionary / Cannibal



BFS:

A

 $A \rightarrow B$ $A \rightarrow C$ $A \rightarrow D$
 (no goal) $A \rightarrow C \rightarrow E$ $A \rightarrow D \rightarrow E$
 $A \rightarrow C \rightarrow E \rightarrow F$ $A \rightarrow D \rightarrow E \rightarrow F$
 $A \rightarrow C \dots \rightarrow G$ $A \rightarrow D \dots \rightarrow G$
 $A \rightarrow C \dots \rightarrow L$ $A \rightarrow D \dots \rightarrow L$
 $A \rightarrow C \rightarrow \dots \rightarrow N$ $A \rightarrow D \rightarrow E \rightarrow M$ $A \rightarrow C \rightarrow N$. $A \rightarrow D \rightarrow N$
 $\dots \rightarrow O$ $\dots \rightarrow O$ $\rightarrow O$ $\dots \rightarrow O$

DFS:

A

$A \rightarrow B$ (no goal)

$A \rightarrow C$

$A \rightarrow C \rightarrow E$

$A \rightarrow C \rightarrow E \rightarrow F$

--- → G

A --- → L

A --- → M

A --- → O (solⁿ)

→ Using DFS we get the sol faster and in single steps and takes less space as compared to BFS.

S.NO.	TOPICS	PAGE NO.
	Advantage of using BFS for N-Queens over A* and Hill Climbing	
-	Gives more solution (4 solutions)	

Disadvantage

- As it gives more solution it is time consuming.

Reasoning to Analogy

(Analogical reasoning / learning)

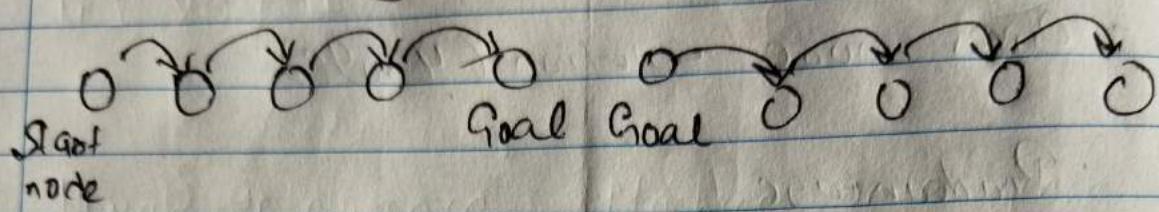
Tumor problem

Castle Problem

- Comparison of problems & applying solution from known problem to the new problem
- 4 R's :
 - Retrieval : Accessing from DB
 - Reuse : Mapping same soln.
 - Revise : Modification to the soln.
 - Rptain : Store

~~2nd Nov~~

Forward Chaining



Backwards chaining

S.NO

- Starts from initial state - starts from goal and and aims for conclusion. backward search to necessary conclusion.
- Data Driven inference - Goal driven inference technique.
Inference = to extract something
- Bottom up approach . Top- Down approach.
- Breadth first search. Depth first search

eg: Diagnosis Recommendation System in
disease diagnosis
Prediction bots.

- The number of final solution is infinite (multiple) The no. of final solution is limited.

S.NO.	TOPICS	PAGE NO.
	MycPox	
-	Expert system to identify bacterial infections and recommended antibiotics.	
-	500 production rules.	
	$\frac{A \rightarrow B}{\text{if } A, \text{ then } B}$	
-	If symptoms \rightarrow Disease If Disease \rightarrow Antibiotics	
	If strain of organism is negative and morphology is rod^+ then then Suggest <i>Enterobacteriaceae</i> .	rod^+
	<pre> graph LR UI[UI] -- "P/P" --> IE["IE
(Knowledge
Engine)"] IE --> KB["KB
Knowledge
Base"] IE --> DData["Dynamic
data
store"] KB <--> DData </pre>	

(V.G)

* HI vs MI

Human Intelligence
(Natural)Machine Intelligence
(Artificial)

- Operation ability
 - Performs declines over time
 - Multitasking with consistent performance.
- Speed:
 - Relatively slow vs comparatively faster.
- Inconsistent accuracy
- Accuracy:
 - Inconsistent accuracy comparatively.
 - Consistent accuracy.
- Storage capacity
 - Unlimited
 - Limited.
- Decision making / Reasoning
 - Might be based on emotions.
 - Unbiased and Rational.

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7 marks
FF

C-4

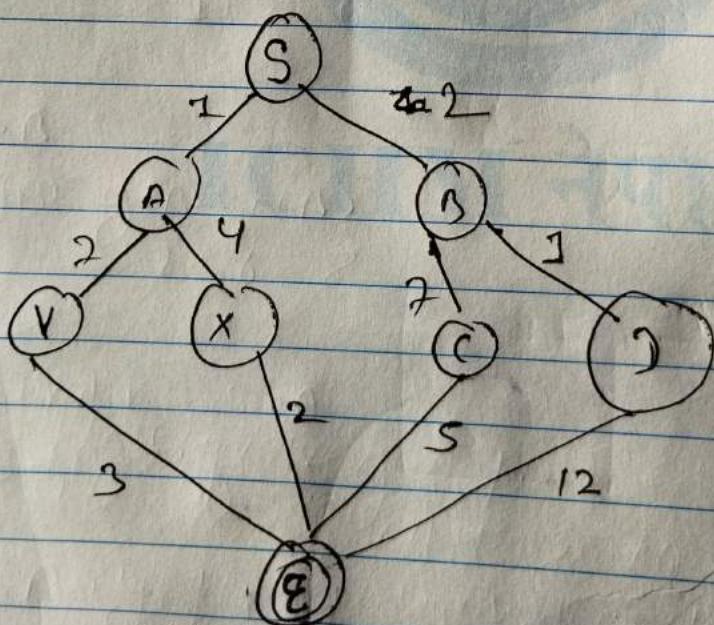
Inference and Reasoning:

Search Techniques:

- 1) Blind Search or Uninformed Search.
- 2) Heuristic Search or Informed Search. Covert A*

Uninformed = no information

Informed = ^{with} information ~~with~~ (about goal)



values for h:

A	5
B	6

TOPICS

PAGE NO.

$$c = 4$$

$$D = 15$$

$$y = 5$$

$$j = 8$$

$$\epsilon = 0$$

Given -> $\epsilon = 0$, that is
goal node

Greedy Search

where $f(n)$ = evaluation function.

$h(n)$ = heuristic value of node 'n'

= estimated path cost from 'n' to goal

Start node: S

$S \rightarrow A$

$$f(n) = h(n)$$

$$f(A) = h(A) = 5 \quad \checkmark$$

$S \rightarrow B$

$$f(B) = h(B) = 6$$

Expanding $S \rightarrow A$

$S \rightarrow A \rightarrow X: f(X) = h(X) = 5 \quad \checkmark$

$S \rightarrow A \rightarrow Y: f(Y) = h(Y) = 8$

Expanding $S \rightarrow A \rightarrow X$

$S \rightarrow A \rightarrow X \rightarrow \epsilon: f(\epsilon) = h(\epsilon) = 0$

$\therefore S \rightarrow A \rightarrow X \rightarrow \epsilon$ is the best path.

(path cost = 7)

A* Search

$$f(n) = g(n) + h(n)$$

where, $f(n)$ = evaluation function

$h(n)$ = heuristic value of node 'n'
= estimated path cost from 'n' to
goal.

$g(n)$ = path cost to reach node 'n'.

S → A

$$\begin{aligned} f(A) &= g(A) + h(A) \\ &= 1 + 5 = 6 \quad \checkmark \end{aligned}$$

S → B

$$\begin{aligned} f(B) &= g(B) + h(B) \\ &= 2 + 6 = 8 \end{aligned}$$

Expanding S → A

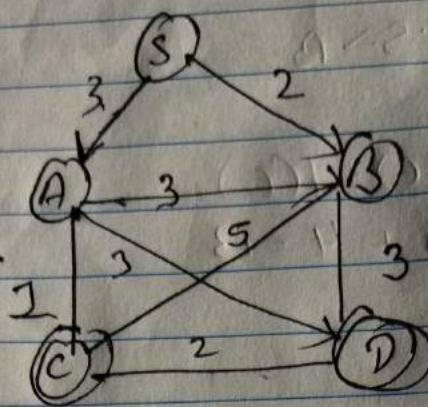
$$\begin{aligned} S \rightarrow A \rightarrow X: f(X) &= g(X) + h(X) \\ &= (1+4) + 5 = 10 \end{aligned}$$

$$\begin{aligned} S \rightarrow A \rightarrow Y: f(Y) &= g(Y) + h(Y) \\ &= (1+2) + 8 = 16 \end{aligned}$$

{ S → D has least cost so expand
S → D.

S.NO.	TOPICS	PAGE NO.
	$\text{Expanding } S \rightarrow B$ $S \rightarrow B \rightarrow C$ $f(B) = g(C) + h(C)$ $= 9 + 4 = 13$	
	$S \rightarrow B \rightarrow D$ $f(D) = g(D) + h(D)$ $= 3 + 15 = 18$	
	$S \rightarrow A \rightarrow X$ has least cost so expand $S \rightarrow A \rightarrow X$.	
	$\text{Expanding } S \rightarrow A \rightarrow X$ $S \rightarrow A \rightarrow X \rightarrow E$ $f(E) = g(E) + h(E)$ $= (1+4+2) + 0$ $= 7$	
	$\therefore S \rightarrow A \rightarrow X \rightarrow E$ is the best path.	
	Solution involving In Solution involving Out Solution involving Both	

#



$$\begin{array}{ll}
 h(S) = h(A) & h(B) = h(C) = h(D) \\
 1 & 3 \quad 3 \quad 0 \quad 0
 \end{array}$$

Values for h : 1, 3, 3, 0, 0

S	1
A	3
B	3
C	0
D	0

Greedy Search

$f(n)$ = evaluation function

$h(n)$ = heuristic value of node 'n'

= estimated path cost from 'n' to goal

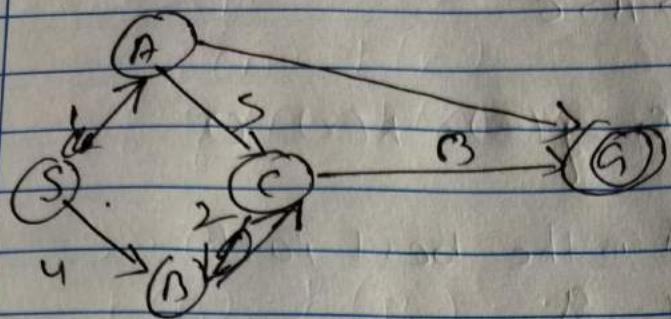
Start node S:

S.NO.	TOPICS	PAGE NO.
	$s \rightarrow A$ $f(n) = h(n)$ $f(A) = h(A) = 3$	
	$s \rightarrow B$ $f(B) = h(B) = 3$	
	Expanding $s \rightarrow A$.	
	$s \rightarrow A \rightarrow n$ $f(n) = 3$ $\boxed{s \rightarrow A \rightarrow C}$ $f(C) = 0$ $\boxed{s \rightarrow A \rightarrow D}$ $f(D) = 0$	
		least path.
	Expanding $s \rightarrow B$	
	$s \rightarrow B \rightarrow C$ $f(C) = 0$ $\boxed{s \rightarrow B \rightarrow D}$ $f(D) = 0$	
	using A'	
	$s \rightarrow A : f(A) = 6$ $s \rightarrow B : f(B) = 5$	
	Expanding $s \rightarrow B$	
	$s \rightarrow B \rightarrow C$ $f(C) = 0.7$. $s \rightarrow D \rightarrow D$ $f(D) = 5$	
	$\therefore s \rightarrow D \rightarrow D$ is the best path	

S.NO.

TOPICS

PAGE NO.



State

h

S

7

A

6

B

2

C

1

G

0

Start node : S

 $S \rightarrow A$

$$f(n) = h(n)$$

$$f(A) = h(A) = 6$$

 $S \rightarrow B$

$$f(n) = h(n)$$

$$f(B) = h(B) = 2$$

Expanding $S \Rightarrow B$ $S \Rightarrow B \Rightarrow C : f(x) = h(x) = 3 \neq$

Expanding $S \rightarrow B \rightarrow C$

$$S \rightarrow B \rightarrow C \rightarrow G = f(x) = g(x) + h(x) = 0$$

$\therefore S \rightarrow B \rightarrow C \rightarrow G$ is the best path
path cost = 9

Now,

using A* Search,

$S \rightarrow A$

$$\begin{aligned}f(A) &= g(A) + h(A) \\&= 1 + 6 \\&= 7\end{aligned}$$

$S \rightarrow B$

$$\begin{aligned}f(B) &= g(B) + h(B) \\&= 4 + 2 \\&= 6\end{aligned}$$

Expanding $S \rightarrow B$

$$\begin{aligned}S \rightarrow B \rightarrow C &= f(x) = g(x) + h(x) \\&= 6 + 1 \\&= 7\end{aligned}$$

S.NO.

TOPICS

PAGE NO.

Since $S \rightarrow A$ & $S \rightarrow B \rightarrow C$ has same cost
expand both:

expanding $S \rightarrow A$.

$$\text{f } S \rightarrow A \rightarrow C$$

$$\begin{aligned} f(C) &= g(C) + h(C) \\ &= 6 + 1 \\ &= 7 \end{aligned}$$

$$\text{ext } S \rightarrow A \rightarrow B$$

$$\begin{aligned} f(B) &= g(B) + h(B) \\ &= 3 + 2 \\ &= 5 \end{aligned}$$

$$S \rightarrow A \rightarrow G$$

$$\begin{aligned} f(G) &= g(G) + h(G) \\ &= 13 + 0 \\ &= 13 \end{aligned}$$

expanding $S \rightarrow A \rightarrow C$

$$S \rightarrow B \rightarrow C \rightarrow G$$

$$\begin{aligned} f(G) &= g(G) + h(G) \\ &= 9 + 0 \\ &= 9 \end{aligned}$$

S.NO.

$$\begin{aligned} & \text{expanding } S \rightarrow A \rightarrow D \\ & S \rightarrow A \rightarrow B \rightarrow C \\ & \cdot f(B) = g(A) + h(B) \\ & \quad \vdots \\ & \quad = 5 + 2 \\ & \quad = 7 \end{aligned}$$

Hence
Since
so

28 NOV

Now,
Since $S \rightarrow A \rightarrow C$ & $S \rightarrow A \rightarrow D$ has same final
path cost so expand both.

expanding $S \rightarrow A \rightarrow B$

$$\begin{aligned} & S \rightarrow A \rightarrow C \rightarrow G \\ & f(C) = g(A) + h(C) \\ & \quad \vdots \\ & \quad = 9 + 0 \\ & \quad = 9 \end{aligned}$$

S2AD

expanding $S \rightarrow A \rightarrow D$

$$S \rightarrow A \rightarrow B \rightarrow C \rightarrow G$$

$$f(D) = f(C) + h(D)$$

$$= \frac{8}{8}$$

 $S \rightarrow A \rightarrow B$. $S \rightarrow A \rightarrow B \rightarrow C \rightarrow G$

$$f(G) = g(A) + h(G)$$

$$= 8 + 0$$

$$= 8$$

S.NO.	TOPICS	PAGE NO.
	Hence, $S \rightarrow A \rightarrow \dots$	
	Since $S \rightarrow A \rightarrow B \rightarrow C \rightarrow G$ has least cost so, $S \rightarrow A \rightarrow B \rightarrow C \rightarrow G$ is the best path.	
28 NOV		
	Start node is a.	
	$a \rightarrow b$ $f(n) = h(n)$ $f(b) = h(b) = 14$.	
	$a \rightarrow c$ $f(n) = h(n)$ $f(c) = h(c) = 18$	
	$a \rightarrow d$ $f(n) = h(n)$ $f(d) = h(d) = 18$	

Expanding $a \rightarrow b$

$a \rightarrow b \rightarrow e : f(e) = h(e) = 20$

$a \rightarrow b \rightarrow$

Expanding, $a \rightarrow e \rightarrow a \rightarrow b \rightarrow e$

$a \rightarrow b \rightarrow e \rightarrow c : f(c) = h(c) = 32$

$a \rightarrow b \rightarrow e \rightarrow z : f(z) = h(z) = 25$

Expanding $a \rightarrow b \rightarrow e \rightarrow z$

Hence $a \rightarrow b \rightarrow e \rightarrow z$, is the final path.

Now using A*

Start node a

$a \rightarrow b$,

$$f(b) = g(b) + f^h(b)$$

$$; \quad g_f(4)$$

$$= 23 \quad \swarrow$$

S.NO.

$a \rightarrow d$

$f(d) = 2$

$a \rightarrow c$

$f(c)$

Since
expand

$a \rightarrow$

a

S.NO.	TOPICS	PAGE NO.
	$a \Rightarrow d$	
	$f(d) = g(d) + h(d)$ $= 7 + 18$ $= 25$	
	$a \Rightarrow c$	
	$f(c) = g(c) + h(c)$ $= 4 + 18$ $= 22 \checkmark$	
	<p>Since, $a \Rightarrow c$ has least path cost so we expand $a \Rightarrow c$</p>	
	$a \Rightarrow c \Rightarrow e$	
	$f(e) = g(e) + h(e)$ $= 21 + 5$ $= 26$	
	$a \Rightarrow c \Rightarrow f$	
	$f(f) = g(f) + h(f)$ $= 16 + 8$ $= 24 \checkmark$	
	<p>Since $a \Rightarrow b$ has the least path cost so we expand $a \Rightarrow b$.</p>	

$$\begin{aligned}
 & a \rightarrow b \rightarrow c \\
 f(c) &= g(c) + h(c) \\
 & = 20 + 5 \\
 & = 25
 \end{aligned}$$

~~a → b~~

Since, $a \rightarrow c \rightarrow f$ has the least path cost so,
we expand

$$a \rightarrow c \rightarrow f$$

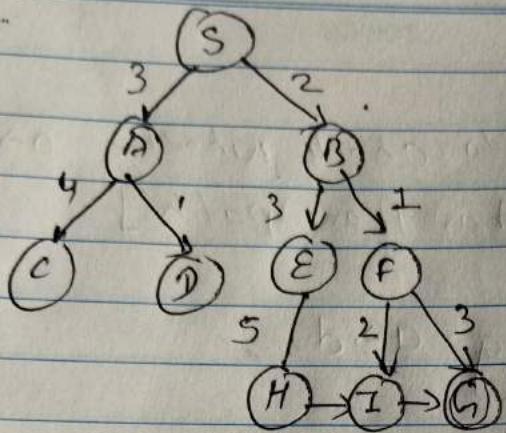
$$\begin{aligned}
 & a \rightarrow c \rightarrow f \rightarrow z \\
 f(z) &= g(z) + h(z) \\
 & = 25 + 0 \\
 & = 25
 \end{aligned}$$

Hence the best path cost is 25 and the
path is $a \rightarrow c \rightarrow f \rightarrow z$.

$$a \rightarrow c \rightarrow f \rightarrow d$$

$$\begin{aligned}
 f(d) &= g(d) + h(d) \\
 & = 30 + 18 \\
 & = 48
 \end{aligned}$$

S.NO.	TOPICS	PAGE NO.
	<p>now,</p> <p>[we have to check whether $a \rightarrow b \rightarrow c$ & $a \rightarrow c$ has least path]</p> <p>expanding $a \rightarrow d$</p> <p>$a \rightarrow d \rightarrow f$.</p> $f(f) = g(f) + h(f)$ $= 21 + 8$ $= 29$ <p>expanding $a \rightarrow b \rightarrow e$.</p> <p>$a \rightarrow b \rightarrow e \rightarrow z$</p> $f(z) = g(z) + h(z)$ $= 25 + 0$ $= 25$ <p>Hence, $a \rightarrow c \rightarrow f \rightarrow z$ and $a \rightarrow b \rightarrow e \rightarrow z$ is the best path cost.</p>	



node	$h(n)$
A	12
B	4
C	7
D	3
E	8
F	2
H	4
I	9
S	13
G	0

Using Greedy
expanding S

$S \rightarrow A$

$$f(A) = 12$$

$S \rightarrow B$

$$f(B) = 4 \Leftarrow$$

expanding $S \rightarrow B \rightarrow E$

$$f(E) = 8 \Leftarrow$$

expanding $S \rightarrow B \rightarrow F$

$$f(F) = 3 \Leftarrow$$

S.NO.	TOPICS	PAGE NO.
	expanding $S \rightarrow B \rightarrow F$	
	$S \rightarrow B \rightarrow F \rightarrow I : f(I)$ = 5	
	$S \rightarrow B \rightarrow F \rightarrow G : f(G)$ = 6	
	Hence $S \rightarrow B \rightarrow F \rightarrow G$ is the best path.	
	now, A^*	
	expanding S	
	$S \rightarrow A : f(A) = g(A) + h(A)$ = 3 + 12 = 15	

S.NO.	TOPICS	PAGE NO.
	Romania question	
	Goal node: Bucharest	
	Starting node: Arad	
	A → B shortest path?	
	<u>Greedy:</u>	
	A → Z : $f(z) = h(z) = 324$	
	A → T : $f(T) = 329$	
	A → S : $f(S) = 253$	
	expanding A → S	
	A → S → F : $f(F) = 178$	
	A → S → R : $f(R) = 193$	
	expanding A → S → F	
	A → S → F → B : $f(B) = 0$	
	∴ A → S → F → B is the shortest path.	
	(140 + 99 + 211)	
	= 450	

A*

$$A \rightarrow Z : f(Z) = 449$$

$$A \rightarrow T : f(T) = 329 + 118 = 442$$

$$A \rightarrow S : f(S) = 393$$

expanding A → S

$$A \rightarrow S \rightarrow F : f(F) = g(F) + h(F) = 412$$

$$A \rightarrow S \rightarrow R : f(R) = 415$$

expanding A → S → R

$$A \rightarrow S \rightarrow R \rightarrow C : f(C) = 526$$

$$A \rightarrow S \rightarrow R \rightarrow P : f(P) = 415$$

expanding A → S → R → P

$$A \rightarrow S \rightarrow R \rightarrow P \rightarrow C : f(C) = 615$$

$$A \rightarrow S \rightarrow R \rightarrow P \rightarrow D : f(D) = 418$$

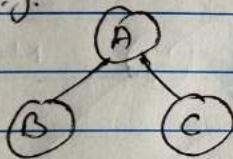
expanding A → S → F

$$A \rightarrow S \rightarrow F \rightarrow B : f(B) = 450$$

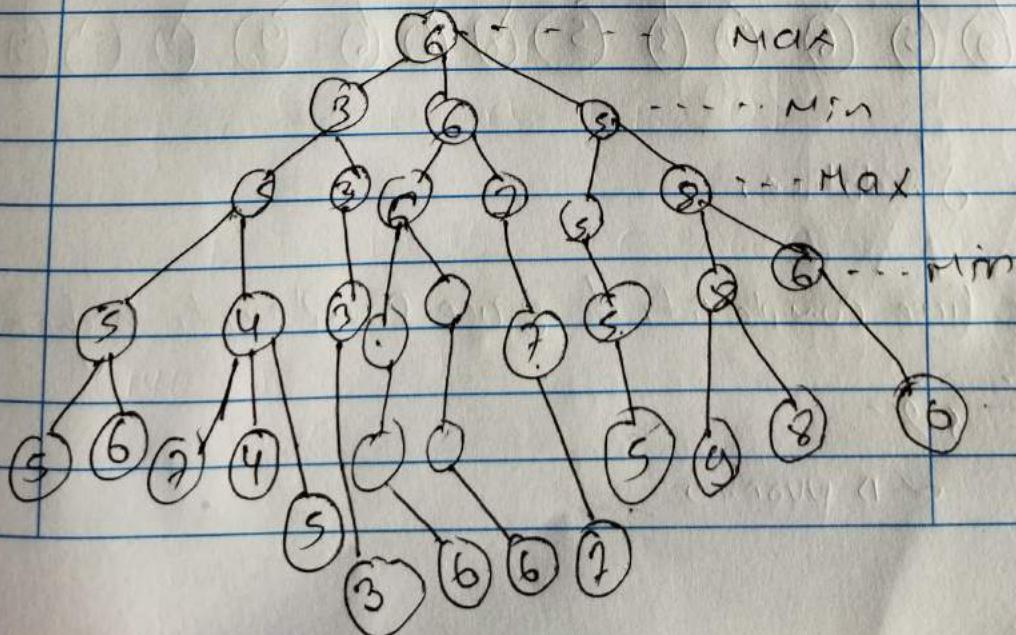
∴ B → G → R → F → D is the best path. (418)

S.NO.	TOPICS	PAGE NO.
	<p>MinMax Algorithm:</p> <p>Greedy f^*</p> <ul style="list-style-type: none"> 1) $f(n) = h(n)$ 2) Heuristic value $f(n) = g(n) + h(n)$ Also, considers Path cost. 3) not required Backtracking 4) Simple & faster Computationally slower 5) Computationally lower Time complexity \neq Space complexity high 	

e.g:



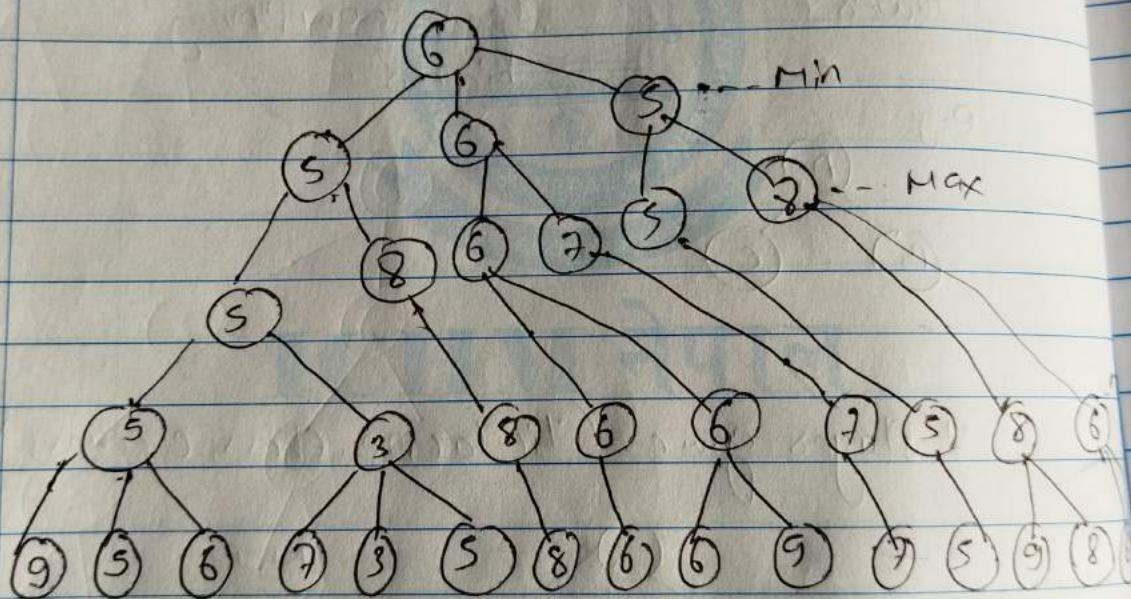
2-player games (Tic-Tac-Toe, Chess, Checkers)



MinMax Algorithm

→ Game playing algorithm used in 2 player games. (Tic Tac Toe, Chess)

→ Two players: Maximizer: Selects Max value
Minimizer: Selects Min value



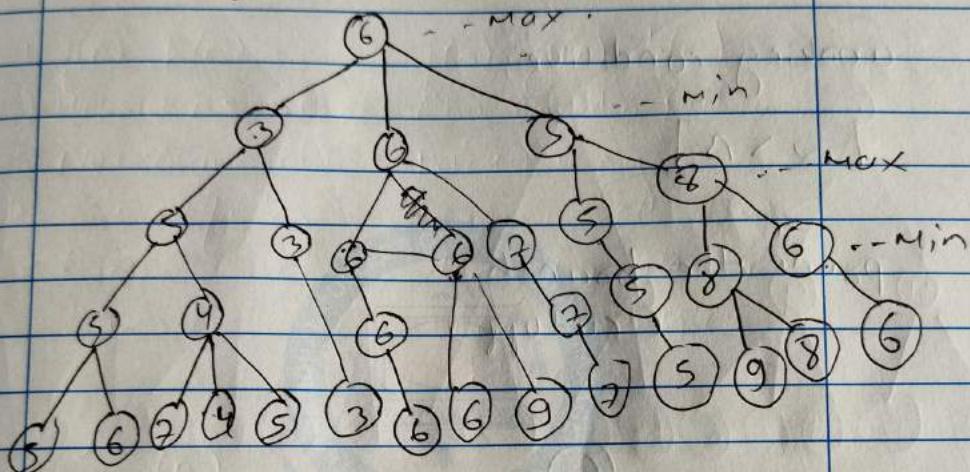
Limitations:

- Time complexity & space complexity higher

soar

S.NO.	TOPICS	PAGE NO.
#	Alpha-beta pruning overcomes these limitations.	

MinMax Algorithm



Maximizer: Selects max value
 2 players game
 Minimizer: min value

Limitations:

- Time complexity and Space complexity high.

* α - β Pruning (Alpha-Beta Pruning)

Initial value (worst case value)

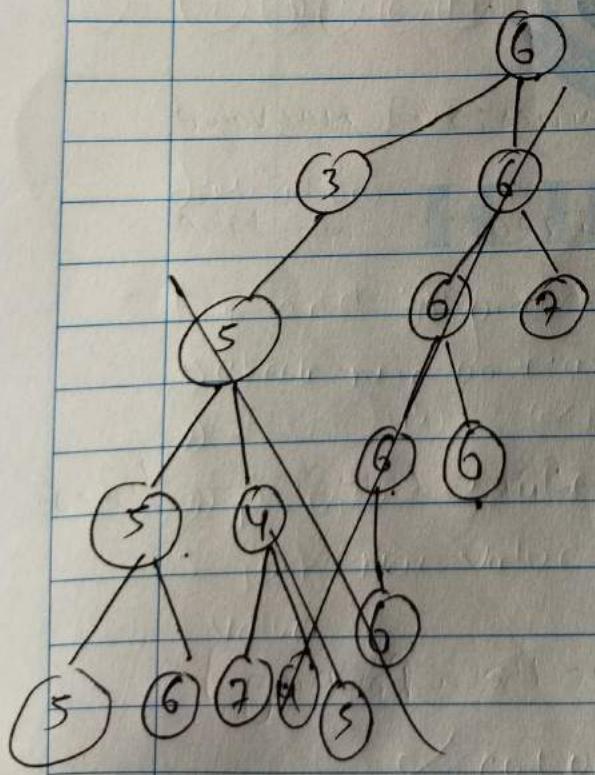
$$\begin{aligned} \alpha &= -\infty \\ \beta &= \infty \end{aligned}$$

- 1) Check min or max level
- 2) Update α - β values
- 3) Push min max value.

working condition:

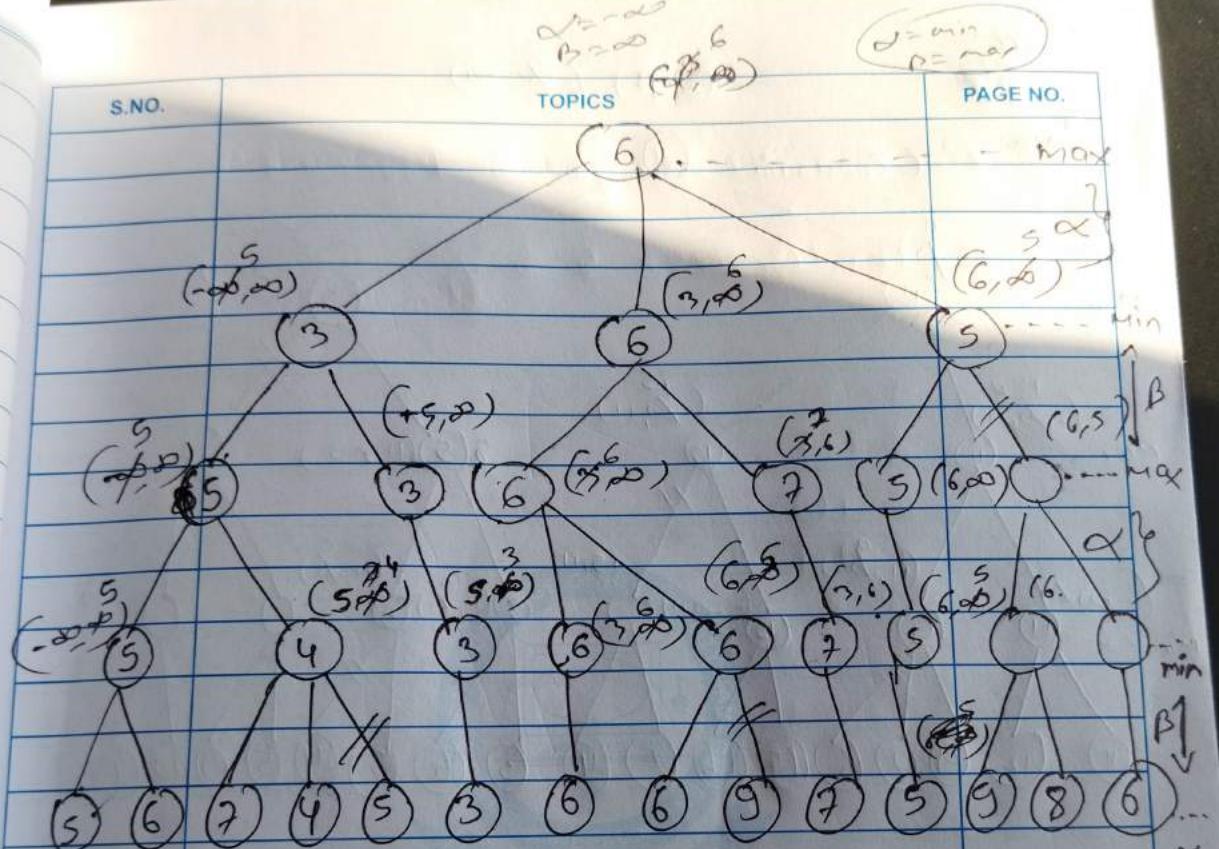
$$[\alpha \geq \beta]$$

prune the branch
(cutoff)



S.NO.



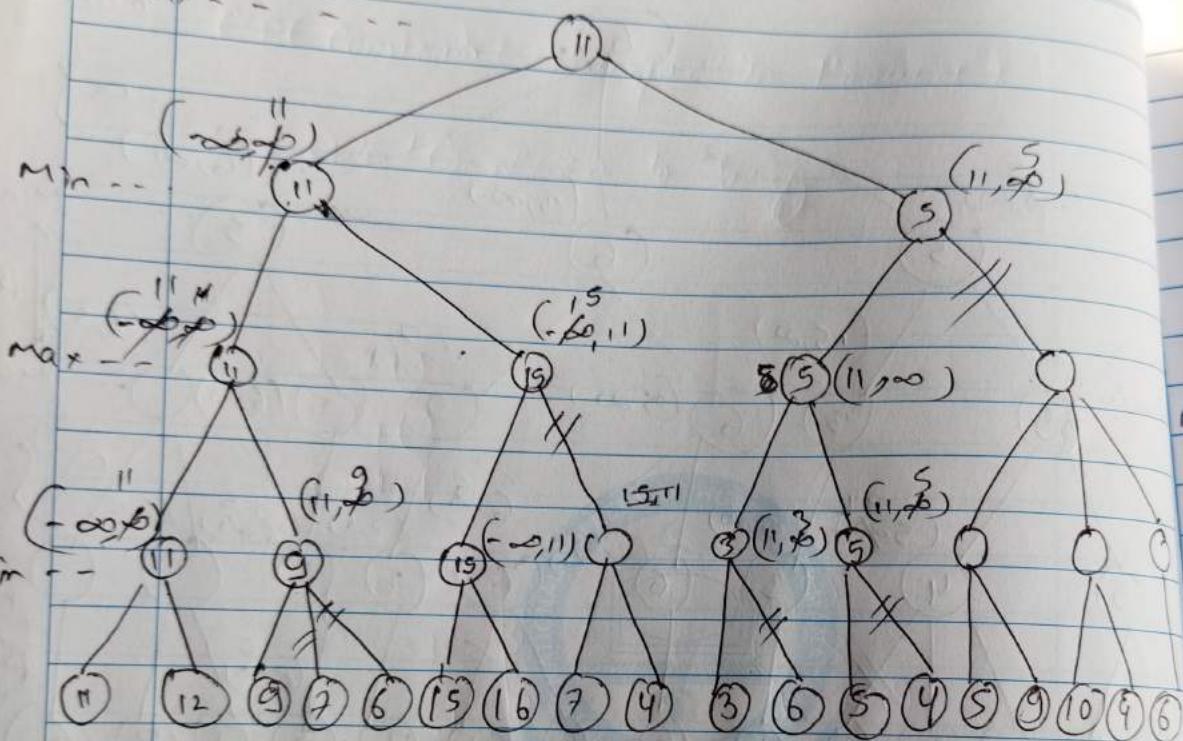


$$\alpha\text{-cutoff} = 0$$

$$\beta\text{-cutoff} = 3$$

5th dec.

$$(\alpha, \beta) = (\infty, \infty)$$



Limitations:

* Time and Algorithm

- 1) Check min or max level
- 2) α (B) β (α)
- 3) Push update α value if in α -level else update β but
3 push max value in α -level else push min value

$$\alpha\text{-cut off} = (\max) = 1$$

$$\beta\text{-cut off} = (\min) = 5$$

S.NO.	TOPICS	$(A, B) = (-\infty, \infty)$	PAGE NO.
1	Max	\emptyset	
2	Min	$(-\infty, \infty)$	
3	Max	$(-\infty, \infty)$	
4	Min	$(-\infty, \infty)$	
5		$(-\infty, \infty)$	
6		$(-\infty, \infty)$	
7		$(-\infty, \infty)$	
8		$(-\infty, \infty)$	
9		$(-\infty, \infty)$	
10		$(-\infty, \infty)$	
11		$(-\infty, \infty)$	
12		$(-\infty, \infty)$	
13		$(-\infty, \infty)$	
14		$(-\infty, \infty)$	
15		$(-\infty, \infty)$	
16		$(-\infty, \infty)$	
17		$(-\infty, \infty)$	
18		$(-\infty, \infty)$	
19		$(-\infty, \infty)$	
20		$(-\infty, \infty)$	
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94		$(-\infty, \infty)$	
95		$(-\infty, \infty)$	
96		$(-\infty, \infty)$	
97		$(-\infty, \infty)$	
98		$(-\infty, \infty)$	
99		$(-\infty, \infty)$	
100		$(-\infty, \infty)$	

S.NO.	TOPICS	PAGE NO.
	Jack [face card]	
	Uncertainty with Bayes' Rule	
	$P(A B) = \frac{P(AB)}{P(B)}$	
	$P(\text{King}/\text{Face}) = \frac{P(\text{King} \cap \text{Face})}{P(\text{Face})}$	
	$= \frac{4/52}{12/52} = \frac{1}{3}$	
	what is uncertainty? sources? How to overcome? Baye's Rule? Eg? ⇒ unable to predicted outcome / conclusion.	
	Sources of uncertainty: Incomplete data / unknown data (underfitting) 2) Vague data / Redundant data. (Overfitting) / weak implication 3) Imprecise language:	

Use of ambiguous words (hardly ever, rarely, often, sometimes)

4) Combining / multiple views of experts.

Using Bayes Rule we can overcome uncertainty

Bayes Rule / Bayesian Network & BDN (Bayesian Belief Network)

- handles uncertainty/uncertain data.
- deals with conditional probability, probability of an event 'A'

$P(A|B)$ given that 'B' has occurred prior.

$$P(A|B) = \frac{P(AB)}{P(B)} - i$$

$$P(AB) = P(A|B) \cdot P(B)$$

S.NO.	TOPICS	PAGE NO.
	$P(B/A) = \frac{P(B \cap A)}{P(A)} \quad \text{--- (2)}$	
	$P(A \cap B) = P(B/A) \cdot P(A)$	
	$P(A/B) = \frac{P(B/A) \cdot P(A)}{P(B)} \quad \text{--- (3)}$	
#	<p>While watching a game of Champion league football in a cafe, you observe someone who's clearly supporting Manchester United.</p> <ul style="list-style-type: none"> What is the probability that they were born within 25 miles of Manchester? <p>Assume that the probability that a randomly selected person in that bar is born within 25 miles of Manchester is $\frac{1}{20}$.</p> <p>The chance that a person born within 25 miles of Manchester actually supports Manchester United is $\frac{2}{10}$?</p>	

The probability that a person
not born within 25 miles of Manchester
supports Manchester United with Probability
 $\frac{1}{10}$?

Sol:

Let S & B denote,

S = Supports M.U.

B = Born within 25 miles of Manchester.

\bar{B} = Not born

by question,

$$P(B) = \frac{1}{10} \quad \therefore P(\bar{B}) = \frac{9}{10}$$

$$P(S|B) = \frac{3}{10}$$

$$P(S|\bar{B}) = \frac{1}{10}$$

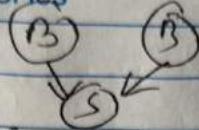
Now,

$$P(B|S) = \frac{P(S|B) \cdot P(B)}{P(S)}$$

$$\text{Now } P(S) = ?$$

(S / B)

TOPICS



PAGE NO.

$$P(S) = P(B) \cdot P(S/B) + P(B) \cdot P(S/B)$$

$$= \frac{1}{20} \cdot \frac{7}{10} + \frac{19}{20} \times \frac{1}{10}$$

$$= \frac{7}{200} + \frac{19}{200}$$

$$= \frac{26}{200}$$

(26)

Bayesian Network. (Not for exam just for understanding)

Pneumonia		(P _p)	Smoking		(S)
T	0.1	Pneumonia	T	0.2	Smoking
F	0.9		F	0.8	SMOKING
↓		↓		↓	
Temperature		Cough		caish	
Pneumonia	Y	N	Pneumonia	Y	caish
Y	0.9	0.1	T	0.95	0.05
N	0.2	0.8	F	0.8	0.2
			F	0.6	0.4
				0.08	0.95

- 1) $P(C/Sn P_n)$
- 2) $P(C)$
- 3) $P(C/S)$

\Rightarrow
 $P(C/Sn P_n)$ = Probability of Cough due to Smoking & Pneumonia

$$= 0.95$$

$P(C) = \text{prob of Cough}$

$$= P(P) \cdot P\left(\frac{C}{P}\right) + \text{(Pneumonia and smoking - Total cases)} \\ \text{2 cases} \quad \text{Total cases} \\ \text{1 case} \quad \text{Cough cases}$$

$$= P(P_n) \cdot P(S) \cdot P\left(C/P_{nns}\right) + \\ P(P_n) \cdot P(\bar{S}) \cdot P\left(C/\bar{P}_{nns}\right)$$

$$+ P(\bar{P}_n) \cdot P(S) \cdot P\left(\frac{C}{P_{nns}}\right) + P(P_n) \cdot P(\bar{S}).$$

$$P\left(\frac{C}{P_{nns}}\right)$$

$$= 0.95 \times 0.05 + 0.8 \times 0.2 + 0.6 \times 0.4 + 0.15 \times 0.95$$

=

S.NO.

$$\text{iii) } P(C/S) = P(P_n) =$$

$$\text{iv) } P(C/P_n) =$$

only smoking $P=F$ $S=V$

S.NO.	TOPICS	PAGE NO.
iii)	$P(C/S) = P(\text{prob of cough due to smoking})$ $= P(P_n) \cdot P(S) \cdot P(C/P_{ns}) + P(\bar{P}_n) \cdot P(\bar{S}) \cdot P(C/\bar{P}_{ns})$	
iv)	$P(C/P_n) = P(P_n) \cdot P(\bar{S}) \cdot P(C/\bar{P}_{ns})$ $+ P(\bar{P}_n) \cdot P(S) \cdot P(C/P_{ns})$ $= 0.95 \times 0.05$	

- 1) Unary Operator (\neg)
 2) Binary operator ($\wedge, \vee, \rightarrow, \leftrightarrow$): two operands

2) Logic

Propositional Logic
 Predicate Logic

Truth Table:

P	Q	(Not)	(Not)	(Or/Djunction)	(And/Conjunction)	"if-then"	(Implication)	(Biconditional/Double Implication)
T	T	F	F	T	T	(T)	T	(if and only if)
T	F	F	T	T	F	(F)	T	
F	T	T	F	T	F	(F)	F	*
F	F	T	T	F	F	(T)	F	

UF 50cc/100
 PNN
 (7-8-30)

$$P \leftrightarrow Q \equiv (P \rightarrow Q) \wedge (Q \rightarrow P)$$

$$P \rightarrow Q \quad | P \rightarrow Q \equiv \neg P \vee Q$$

P: "It is raining"

Hypothesis \rightarrow Conclusion

Q: "I will take ^{an} umbrella"

Condition \rightarrow Action

Antecedent \rightarrow Consequent

Logically
equivalent