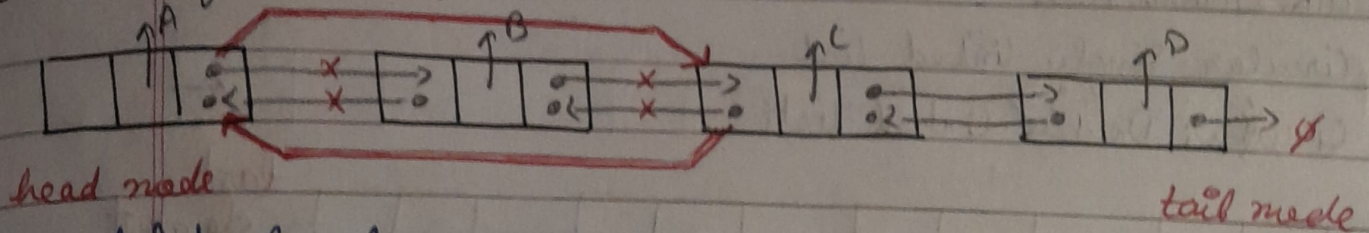


## Deletion of a node, IN B/W



- ① To delete B,  $A \cdot \text{next} \rightarrow C \cdot \text{previous}$   $A \cdot \text{next} \cdot \text{next}$
- ②  $C \cdot \text{previous} \rightarrow A \cdot \text{next}$  ~~Express~~  $C \cdot \text{prev} \cdot \text{prev}$
- ③ Decrease  $\text{size} = \text{size} - 1$ .

## Complexity of insertion

Insertion at head  $O(1)$

Insertion at tail  $O(n)$

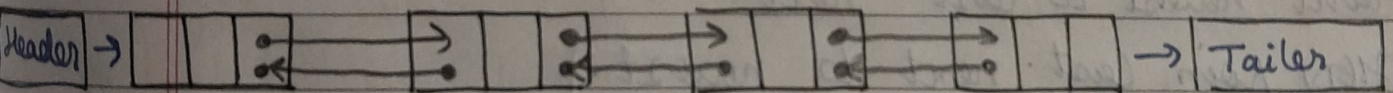
Insertion in b/w  $O(n)$

Deletion at head

Deletion at tail

Deletion in b/w

★ TO GENERALISE INSERTION,  $\text{insert\_element}(e, \text{pred}, \text{succ})$



def  $\text{insert\_in\_b/w}(\text{self}, e, \text{pred}, \text{successor})$

$\text{newest} = \text{self} \cdot \text{Node}(e, \text{pred}, \text{successor})$

$\text{pred} \cdot \text{next} = \text{newest}$

$\text{successor} \cdot \text{prev} = \text{newest}$

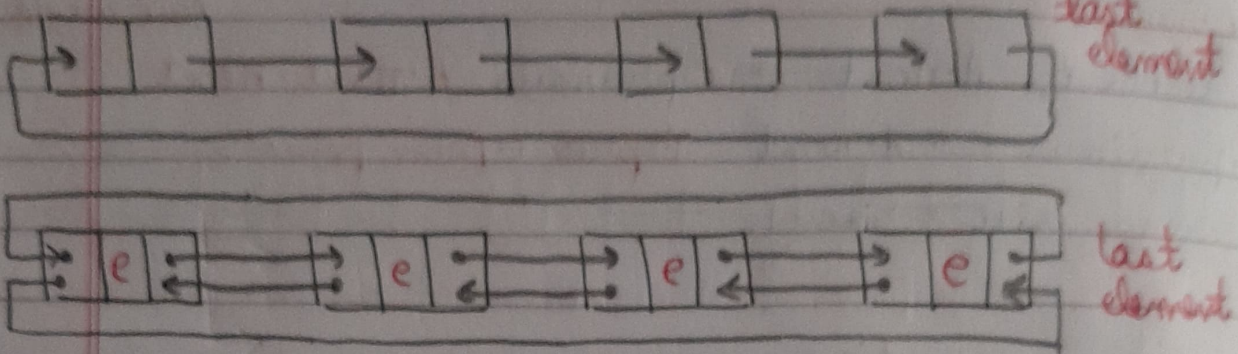
$\text{self} \cdot \text{Size} += 1$

return newest



CIRCULAR LINKED LIST

- ① Circular Singly Linked list
- ② Circular Double Linked list



Insertion at beginning -

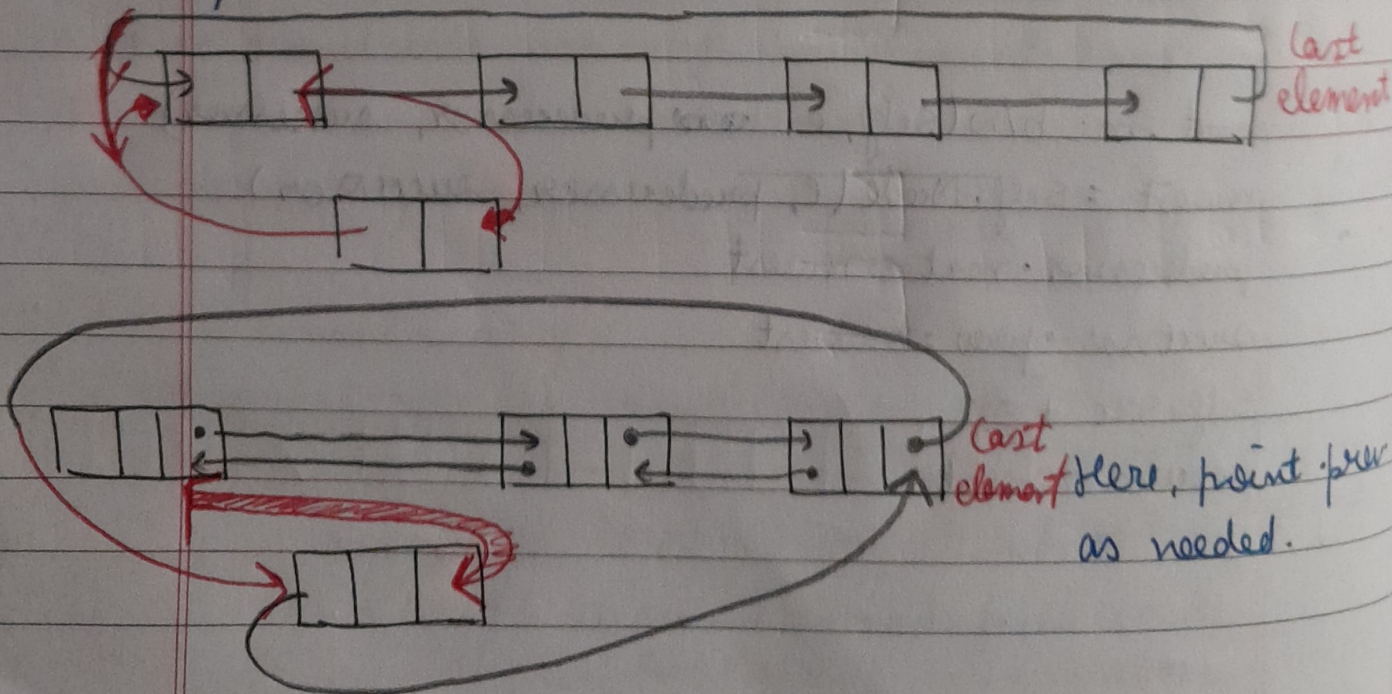
create a new node, newest.

we have explicit reference to last element.

Point last.next to newest.

Point newest to last.next.

THEN point last.next to newest.



last element here, point prev as needed.

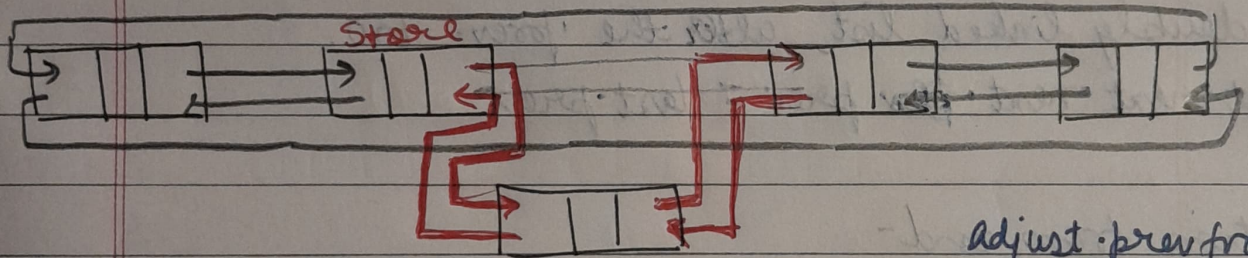
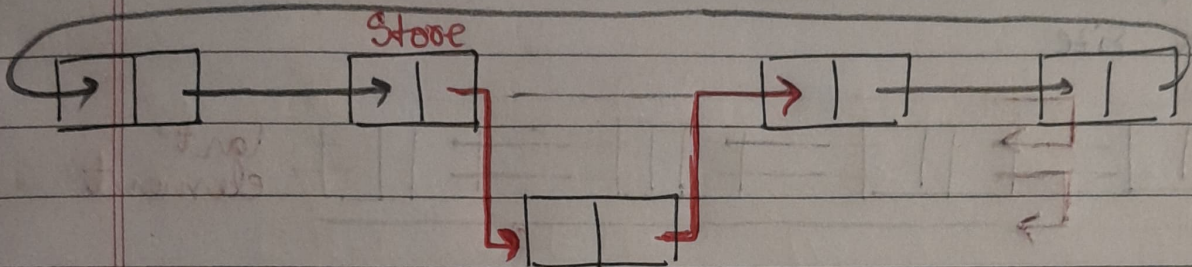


Insertion in b/w nodes -

Travel to the node where you want to enter, store it  
create a new node.

$\text{new node} \cdot \text{next} = \text{stored num} \cdot \text{next}$

$\text{stored num} \cdot \text{next} = \text{new node}$



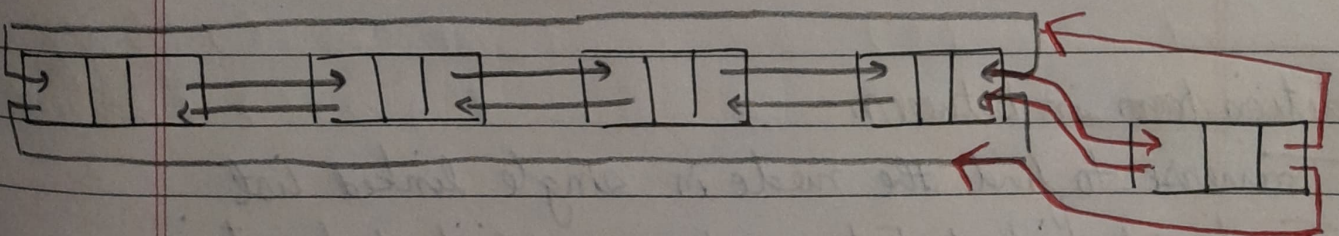
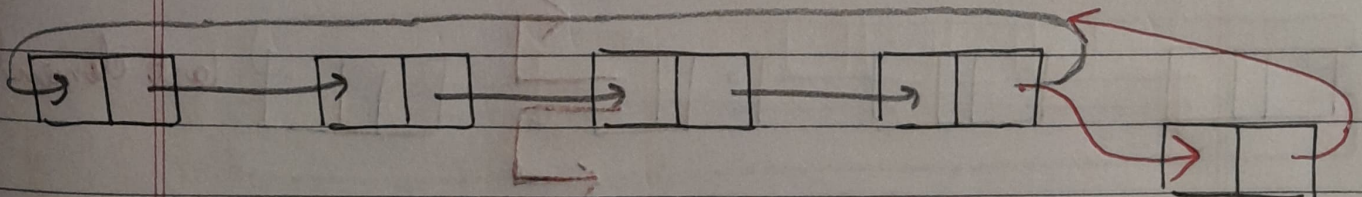
adjust prev for DLL

Insertion at the end (after last) -

Create new node.

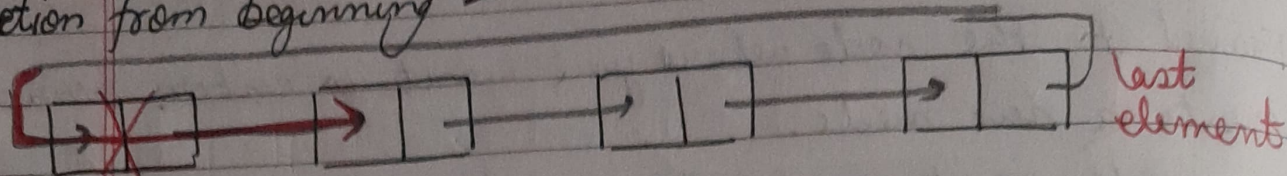
$\text{new node} \cdot \text{next} \rightarrow \text{last} \cdot \text{next}$

$\text{last} \cdot \text{next} \rightarrow \text{new node}$



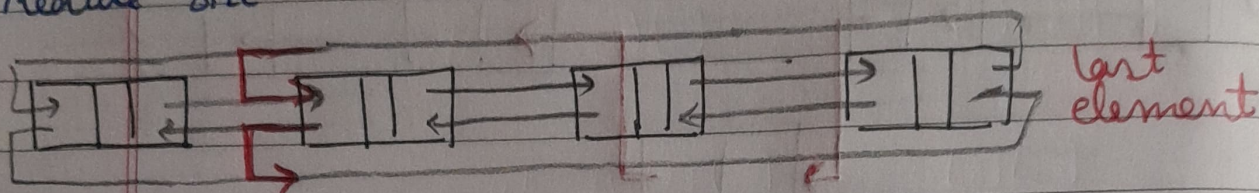


Deletion from beginning -



Simple point  $\text{last} \cdot \text{next} = \text{last} \cdot \text{next} \cdot \text{next}$ .

Reduce size



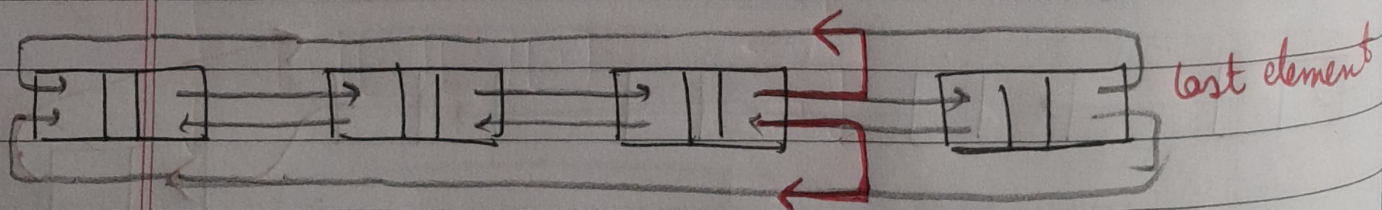
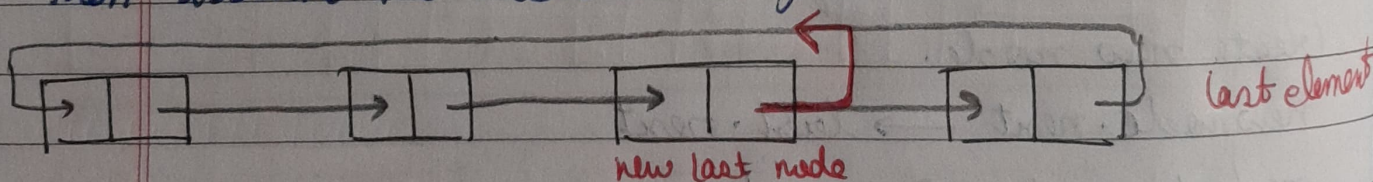
In doubly linked list, alter the 'prev' too.

$\text{last} \cdot \text{next} \cdot \text{next} \cdot \text{prev} = \text{last} \cdot \text{prev}$

Deletion from end -

Reassign the last node (traverse)

Then assume the node to be first node



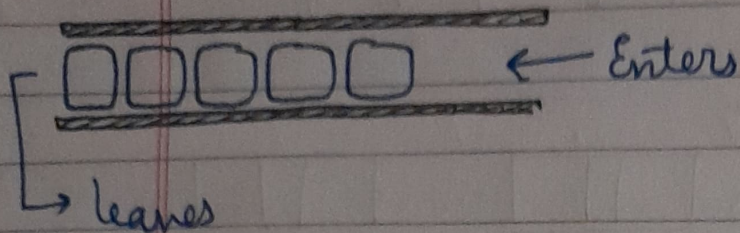
Deletion from in between -

Traverse to find the node in single linked list

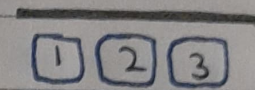
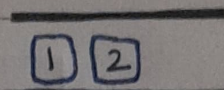
In double linked list, we have explicit declaration



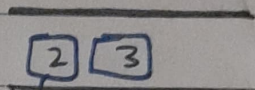
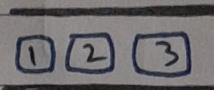
QUEUE FIFO: first in first out.



Enqueue (insert element) -



Deque (deletion of element) -



is\_full()

is\_empty()

Peek (front)

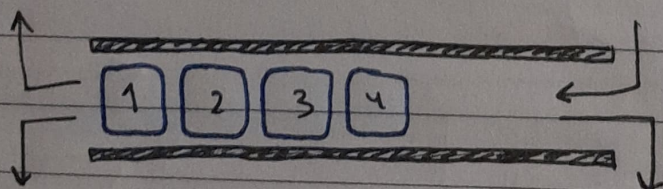
Complexity

enqueue is  $O(1)$

deque is  $O(n)$

*★ we solve this using circular and double ended queues.*

DOUBLE ENDED QUEUE



\* insert\_at\_end() and insert\_at\_start()

\* delete\_at\_end() and delete\_at\_start()

*★ Complexity*