

T-603-THYD Fall 2017, PROJECT – part 2 (updated grammar)

Marked as orange are grammar symbols that were intentionally not left-factored as they are “harmless”, that is, a recursive-descent LL(1) parser should not have any problems parsing them. Leaving the grammar like that will make it easier to build an efficient parser, for example, in some cases, allow you to implement right-recursion (tail-recursion) as iteration.

```
program           ::= class id { variable_declarations method_declarations }

variable_declarations ::= type variable_list ; variable_declarations | ε

type              ::= int | real

variable_list     ::= variable | variable , variable_list

variable          ::= id

method_declarations ::= method_declaration method_declarations | method_declaration

method_declaration ::= static method_return_type id ( parameters )
                   { variable_declarations statement_list }

method_return_type ::= type | void

parameters        ::= parameter_list | ε

parameter_list    ::= type variable | type variable , parameter_list

statement_list    ::= statement statement_list | ε

statement         ::= if ( expr ) statement_block optional_else
                   | for ( variable = expr ; expr ; variable op_incr_decr ) statement_block
                   | return optional_expr ;
                   | break ;
                   | continue ;
                   | statement_block
                   | id_start_stm

id_start_stm      ::= id ( expr_list ) ;
                   | id op_incr_decr ;
                   | id = expr ;

optional_expr     ::= expr | ε

statement_block   ::= { statement_list }

optional_else     ::= else statement_block | ε

expr_list         ::= expr more_expr | ε
```

<i>more_expr</i>	::= , <i>expr</i> <i>more_expr</i> ϵ
<i>expr</i>	::= <i>expr_and</i> <i>expr'</i>
<i>expr'</i>	::= <i>expr_and</i> <i>expr'</i> ϵ
<i>expr_and</i>	::= <i>expr_eq</i> <i>expr_and'</i>
<i>expr_and'</i>	::= && <i>expr_eq</i> <i>expr_and'</i> ϵ
<i>expr_eq</i>	::= <i>expr_rel</i> <i>expr_eq'</i>
<i>expr_eq'</i>	::= <i>op_eq</i> <i>expr_rel</i> <i>expr_eq'</i> ϵ
<i>expr_rel</i>	::= <i>expr_add</i> <i>expr_rel'</i>
<i>expr_rel'</i>	::= <i>op_rel</i> <i>expr_add</i> <i>expr_rel'</i> ϵ
<i>expr_add</i>	::= <i>expr_mult</i> <i>expr_add'</i>
<i>expr_add'</i>	::= <i>op_add</i> <i>expr_mult</i> <i>expr_add'</i> ϵ
<i>expr_mult</i>	::= <i>expr_unary</i> <i>expr_mult'</i>
<i>expr_mult'</i>	::= <i>op_mult</i> <i>expr_unary</i> <i>expr_mult'</i> ϵ
<i>expr_unary</i>	::= <i>op_unary</i> <i>expr_unary</i> <i>factor</i>
<i>factor</i>	::= num (<i>expr</i>) id id (<i>expr_list</i>)
<i>op_eq</i>	::= == !=
<i>op_rel</i>	::= < <= > >=
<i>op_add</i>	::= + -
<i>op_mul</i>	::= * / %
<i>op_unary</i>	::= + - !
<i>op_incr_decr</i>	::= -- ++

*Note, some of the **id** terminal symbols above are referring to a variable, and should result in a VariableExprNode node to be generated. However, for ease of parsing, it is better in those situations to create the VariableExprNode explicitly without calling the variable method (this is because, to know that you were expecting a variable, you will already have had to match the **id** token).*