# Question 1 – How would you compare RR and CS

Repeatable Read is DB2’s highest level of isolation, and the closest equivalent to a serializable schedule (also avoids phantom problem, with an insert not being caught by the lock). Cursor Stability is a sort of Read committed isolation level. RR will be surer of correct results, but it will lock a lot more than CS and then force the DB to work less in parallel, making it slower than CS.

The currently commited semantics feature of CS makes it able to read from the currently commited result, instead of the read having to wait for an update lock to be release. So instead of waiting it can fetch the last commited, as if it didn’t know anything was updating the particular row. This is okay for most situations, and a lot faster in most scenarios. This has a slight performance penalty because of additional logging.

Almost always CS will be fine, unless you are working in a business where you have to be absolutely sure than everything is done atomically and serializable.

# Question 2

# Question 3

# Question 4

# Question 5