

# Unifying Branch-and-Bound methods for $\ell_0$ -penalized problems

Théo Guyard<sup>\*</sup>, Cédric Herzet<sup>†</sup>, Clément Elvira<sup>‡</sup>, Ayşe-Nur Arslan<sup>◇</sup>

<sup>\*</sup>Inria and Insa Rennes, <sup>†</sup>Ensaï, <sup>‡</sup>CentraleSupélec, <sup>◇</sup>Inria Bordeaux

## In short

Provide necessary ingredients to implement **generic** and **efficient**  $\ell_0$ -problem solvers.

## $\ell_0$ -problem

$$p^* = \min_{\mathbf{x} \in \mathbb{R}^n} f(\mathbf{A}\mathbf{x}) + \lambda \|\mathbf{x}\|_0 + h(\mathbf{x})$$

Generic solvers

Specialized solvers

✗ Slow  
✓ Flexible w.r.t  $f/h$

✓ Fast  
✗ Only some  $f/h$

Proposed solver

✓ Fast  
✓ Flexible w.r.t  $f/h$

## Working hypotheses

- $f$  simple to work with (convex, smooth, ...)
- $h$  proper, closed, separable, cont. at  $\mathbf{x} = \mathbf{0}$
- $h$  even, convex, coercive,  $h \geq h(\mathbf{0}) = 0$

## Versatile toolbox

User-defined  $f$   
func\_value  
conj\_value  
prox\_value  
subd\_value

User-defined  $h$   
func\_value  
conj\_value  
prox\_value  
subd\_value

**E10ps**

Branch-and-Bound solver  
Tunable strategies  
Acceleration methods

## Branch-and-Bound algorithms

### Principle

“Partition the feasible space into **regions** and **prune** those that cannot contain minimizers.”

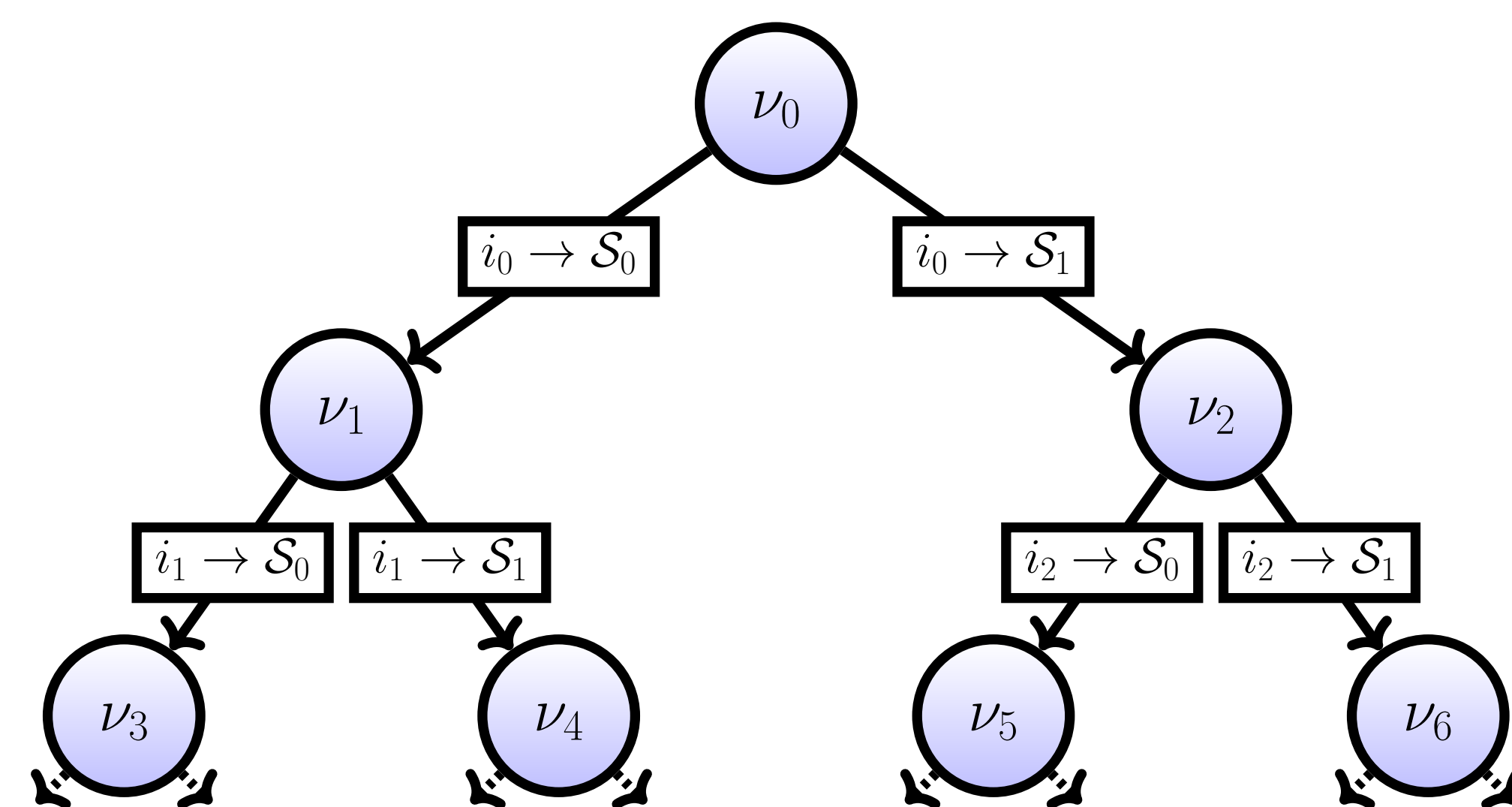
### Regions construction

- $\nu = (\mathcal{S}_0, \mathcal{S}_1)$  disjoint subsets of  $\{1, \dots, n\}$
- $\mathcal{S}_0$  entries fixed to zero
- $\mathcal{S}_1$  entries fixed to non-zero

$$\mathcal{X}^\nu = \{\mathbf{x} \in \mathbb{R}^n \mid \mathbf{x}_{\mathcal{S}_0} = \mathbf{0}, \mathbf{x}_{\mathcal{S}_1} \neq \mathbf{0}\}$$

- All regions explored: problem solved

### Feasible space exploration



### Processing nodes

- Node problem  
$$p^\nu = \min_{\mathbf{x} \in \mathcal{X}^\nu} f(\mathbf{A}\mathbf{x}) + \lambda \|\mathbf{x}\|_0 + h(\mathbf{x})$$
- Test if region  $\mathcal{X}^\nu$  contains minimizers:  $p^\nu > p^*$
- Practical **pruning test** with bounds  
$$p^\nu \geq \tilde{p}^\nu > \bar{p} \geq p^*$$
- Upper bound  $\bar{p}$  on  $p^*$   
→ Evaluate objective at any point  
→ Efficient heuristics

## Generic relaxations

### Constructing lower-bounds

#### Node problem reformulation

$$p^\nu = \min_{\mathbf{x} \in \mathbb{R}^n} f(\mathbf{A}\mathbf{x}) + g^\nu(\mathbf{x})$$

$$\text{with } g_i^\nu(x_i) = \begin{cases} \mathbf{I}(x_i = 0) & \text{if } i \in \mathcal{S}_0 \\ h_i(x_i) + \lambda & \text{if } i \in \mathcal{S}_1 \\ h_i(x_i) + \lambda \|x_i\|_0 & \text{otherwise} \end{cases}$$

#### Primal relaxation

$$\min_{\mathbf{x} \in \mathbb{R}^n} f(\mathbf{A}\mathbf{x}) + (g^\nu)^{**}(\mathbf{x})$$

#### Dual relaxation

$$\max_{\mathbf{u} \in \mathbb{R}^m} -f^*(-\mathbf{u}) + (g^\nu)^*(\mathbf{A}^T \mathbf{u})$$

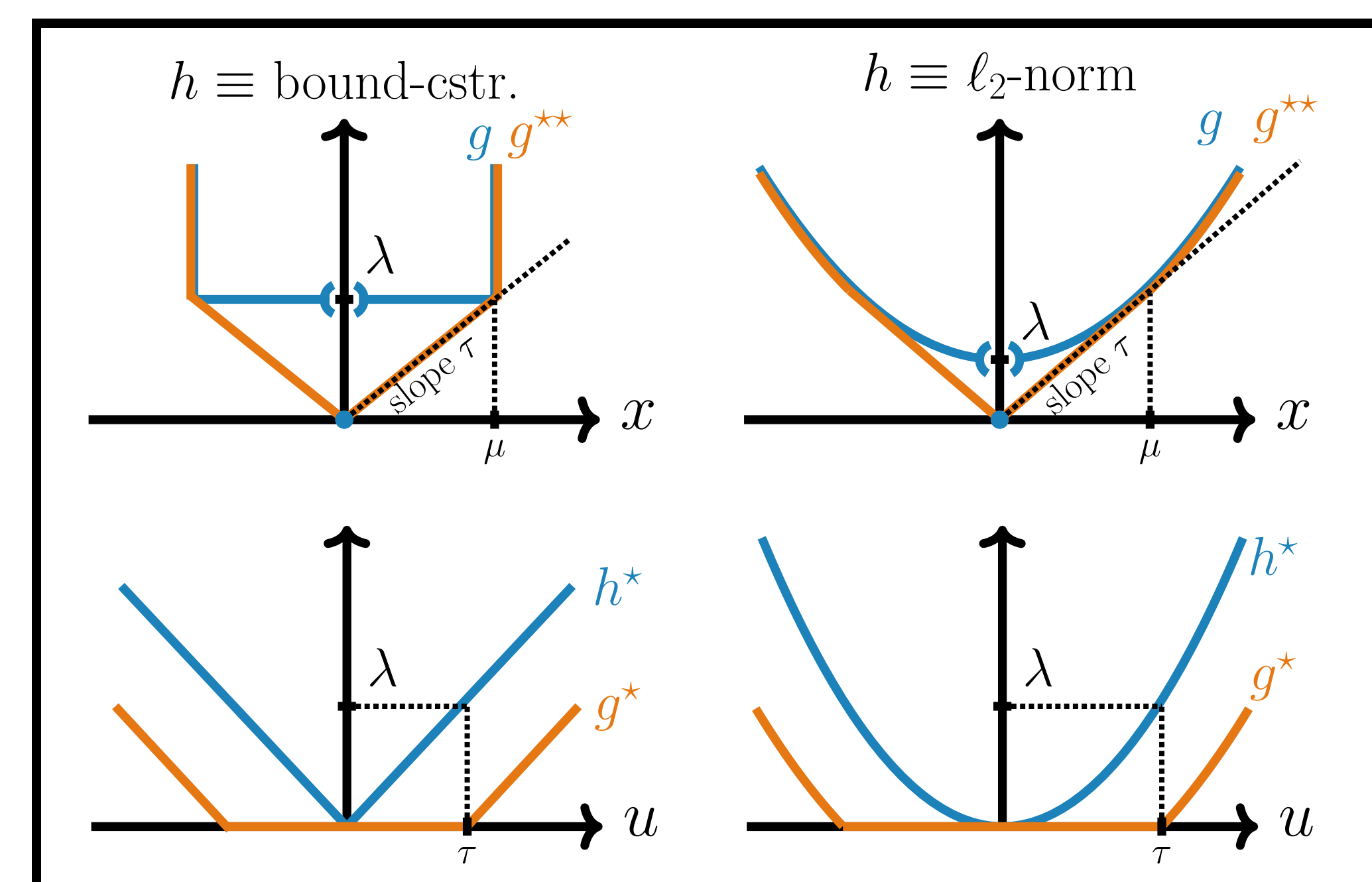
### Closed-form expressions

- 1D-study of  $g(x) = h(x) + \lambda \|x\|_0$
- Parameters  $(\tau, \mu)$  easily computable from  $h^*$
- Closed-form expressions

$$g^{**}(x) = \begin{cases} \tau|x| & \text{if } |x| \leq \mu \\ h(x) + \lambda & \text{if } |x| > \mu \end{cases}$$

$$g^*(u) = \begin{cases} 0 & \text{if } |u| \leq \tau \\ h^*(u) - \lambda & \text{if } |u| > \tau \end{cases}$$

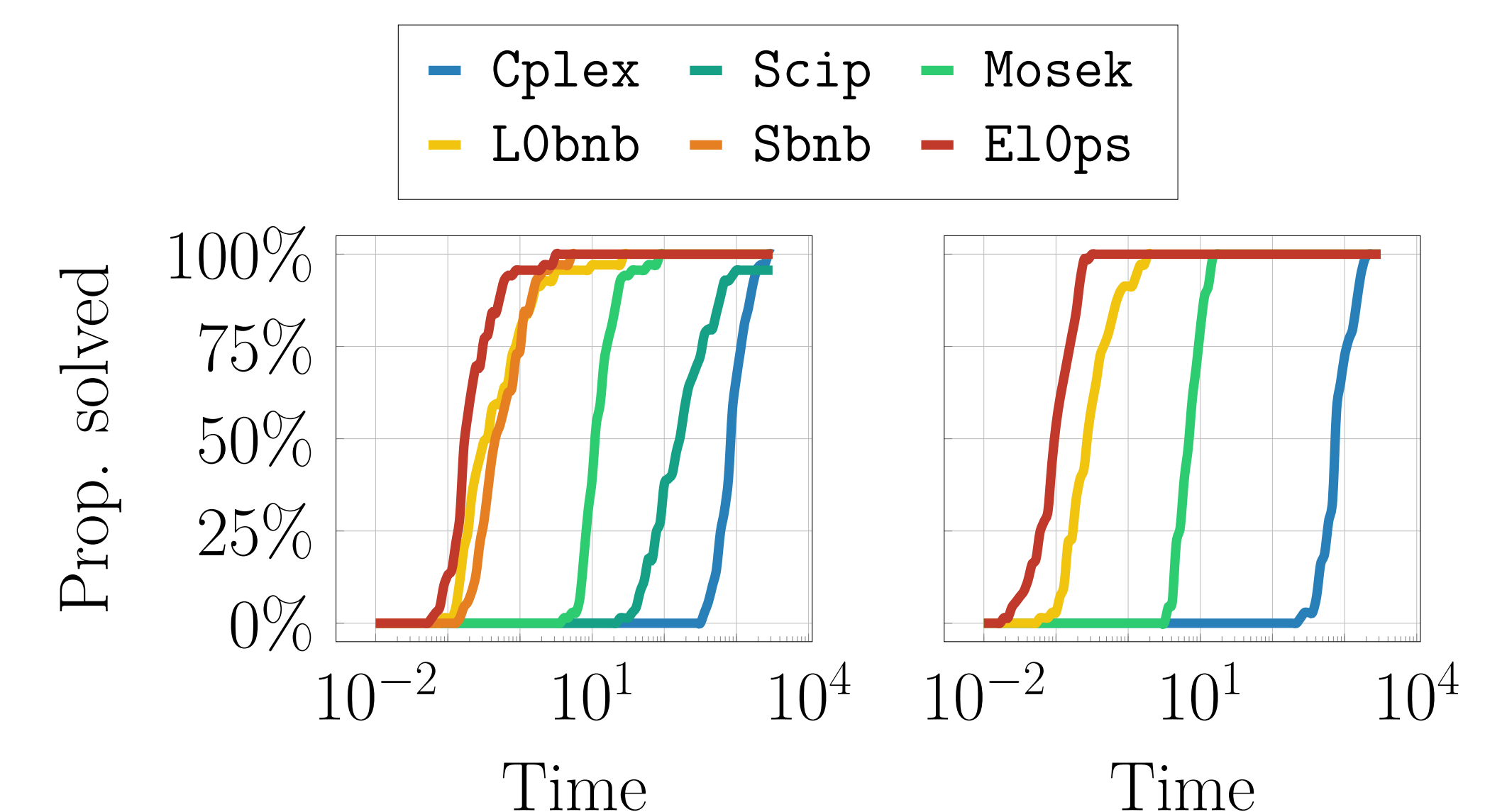
- Simple evaluation of prox, subdiff, ...



## Numerical results

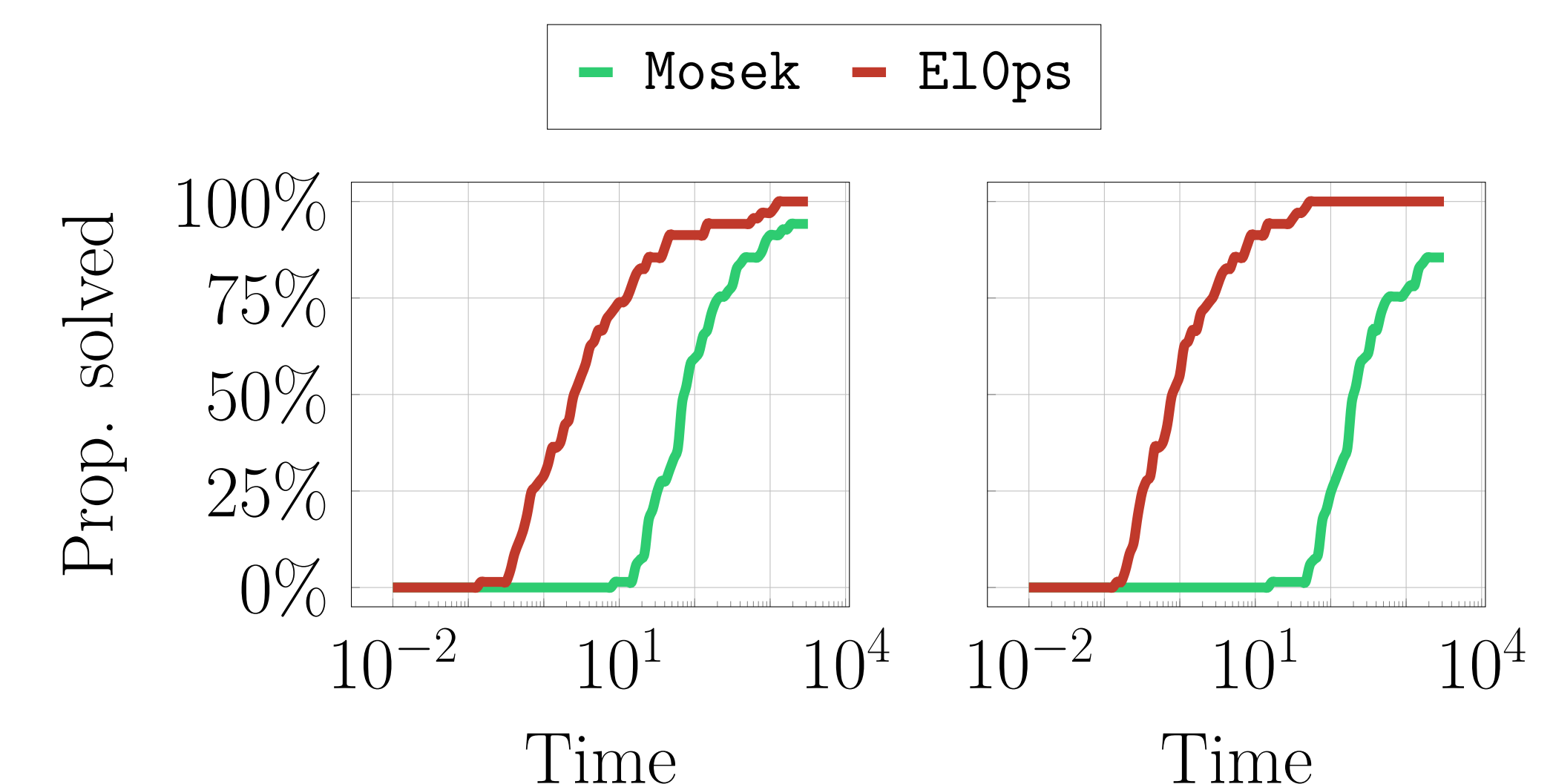
### Standard instances

- $f$ : least-squares
- $h$ : bound-ctr. or  $\ell_2$ -norm (left/right)
- $\lambda$ : tuned statistically



### New application opportunities

- $f$ : logistic
- $h$ : bound-ctr. +  $\ell_1$ - or  $\ell_2$ -norm (left/right)
- $\lambda$ : tuned statistically



✓ Fast  
✓ Flexible w.r.t  $f/h$

## Take home message

Opportunities to address new instances of  $\ell_0$ -problems efficiently.