# Project Report for Seminar Course in Numerical Analysis, VT23

Junzi Zhang, Brendan O'Donoghue, Stephen Boyd: Globally Convergent Type-I Anderson Acceleration for Non-Smooth Fixed-Point Iterations

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Input : Initial value x_0 \in \mathbb{R}^n and function f : \mathbb{R}^n \to \mathbb{R}^n.

for k = 0, 1, \dots do
| Set x_{k+1} = f(x_k).
end
```

Algorithm 1: Fixed point iteration (original)

## Introduction

The following report will summarise and present the main results of the paper [1] as part of the course NUMN27, Seminar in Numerical Analysis. More precisely the report will deal with a modification of the Anderson-Accelleration-I (AA-I) algorithm as described in [1]. For this we start with a motivation of the AA-II algorithm. Then we will discuss the modifications made to the algorithm. After that we will give the main convergence result and finally test the AA-I algorithm in numerical experiments. The report, the python implementation and the corresponding presentation of the topic can be found online under [3].

# Motivation of the Anderson-Acceleration algorithm

In science it is a common problem to find a fixed point  $x = f(x) \in \mathbb{R}^n$  of a function  $f \colon \mathbb{R}^n \to \mathbb{R}^n$ . An equivalent formulation is to find a zero  $x \in \mathbb{R}^n$  of the function  $g = \operatorname{Id} - f$ . In the following we assume that f indeed has a fixed point and that f is non-expansive, i.e. we have for all  $x, y \in \mathbb{R}^n$  that

$$||f(x) - f(y)|| \le ||x - y||.$$

We however also assume that  $\nabla f$  is unknown which means we cannot use Newton iteration to solve our problem. We also assume that our problem is noisy so that we cannot take finite difference derivatives. If the cost of evaluating f is very high then line search becomes infeasable and if n is very large we are forced to work matrix free. We know that this type of problem can be solved by the fixed point iteration described in algorithm 1.

It can be shown that the fixed point iteration will in this case converge to a fixed point of f. This convergence can be very slow in practice if the Lipshitz-constant of f is close to 1. It is here were the Anderson-Acceleration (AA) methods come in.

The main idea of the AA is to use the information gained from previous function evaluations. To update  $x_{k+1}$  we form a weighted average as described in algorithm 2. For simplicity of notation we assume here and in the following that our memory is unlimited.

The General AA leaves the quation open of how to choose the  $\alpha \in \mathbb{R}^{k+1}$ . Since finding a fixed point of f is equivalent to finding a zero of  $f = \operatorname{Id} - f$  it seems sensible to require  $\alpha$  to minimise

$$\|\sum_i \alpha_i g_i\|_2$$
.

```
\begin{aligned} &\textbf{Input:} x_0 \in \mathbb{R}^n \text{ and } f \colon \mathbb{R}^n \to \mathbb{R}^n. \\ &\textbf{for } k = 0, 1, \dots \textbf{ do} \\ & \mid \text{ Set } f_k = f(x_k). \\ & \mid \text{ Choose } \alpha = \alpha^k \in \mathbb{R}^{k+1} \text{ such that } \sum_i \alpha_i = 1. \\ & \mid \text{ Set } x_{k+1} = \sum_i \alpha_i f_i. \end{aligned}
```

Algorithm 2: General AA (Anderson Acceleration)

Algorithm 3: AA-II

This then yields the AA-II algorithm given in 3.

It is shown in [1] that this update can be brought into the form of a quasi-Newton-like method as shown in algorithm 4.

Now one could expect  $H_k$  to be an approximate inverse of  $\nabla f(x_k)$ . Indeed one can show

**Proposition 1** (Approximate inverse Jacobian).  $H_k$  minimises  $||H_k - \text{Id}||_F$  under the multi-secant condition  $H_kS_k = Y_k$ .

The good Broyden method approximates the Jacobian rather than its inverse and tends to yield better results. This motivates to choose  $B_k$  to be a minimiser of  $||B_k - \text{Id}||_F$ 

```
Input: x_0 \in \mathbb{R}^n and f : \mathbb{R}^n \to \mathbb{R}^n.

Set x_1 = f(x_0).

for k = 0, 1, \dots do
\begin{cases} \text{Set } g_k = g(x_k). \\ \text{Construct } S_k = \begin{bmatrix} x_1 - x_0 & \cdots & x_k - x_{k-1} \end{bmatrix} \in \mathbb{R}^{n \times k} \text{ and } \\ Y_k = \begin{bmatrix} g_1 - g_0 & \cdots & g_k - g_{k-1} \end{bmatrix} \in \mathbb{R}^{n \times k}. \\ \text{Set } H_k = \text{Id} + (S_k - Y_k) \left(Y_k^\top Y_k\right)^{-1} Y_k^\top \in \mathbb{R}^{n \times n}. \\ \text{Set } x_{k+1} = x_k - H_k g_k. \end{cases}
```

Algorithm 4: AA-II (reformulated)

```
Input: x_0 \in \mathbb{R}^n and f: \mathbb{R}^n \to \mathbb{R}^n.

Set x_1 = f(x_0)

for k = 0, 1, \dots do

Set g_k = g(x_k).

Construct S_k from x_0, \dots, x_k and Y_k from g_0, \dots, g_k.

Set B_k = \operatorname{Id} + (Y_k - S_k) \left(S_k^\top S_k\right)^{-1} S_k^\top \in \mathbb{R}^{n \times n}.

Set H_k = H_k^{-1}.

Set H_k = H_k^{-1}.

Set H_k = H_k^{-1}.
```

Algorithm 5: AA-I

under the condition  $B_k Y_k = S_k$ . One can then show that

$$B_k = \operatorname{Id} + (Y_k - S_k) \left( S_k^{\top} S_k \right)^{-1} S_k^{\top}.$$

which yields the AA-I algorithm 5.

## **Modifications to AA-I**

The AA-I algorithm as stated in 5 has some apparent problems. For one, the approach is not matrix-free. This will be fixed by a rank-1 update formula for matrices  $B_k$  and later  $H_k$ . It may also occur in a step that the matrix  $H_k$  is not well-defined. This may occur if  $B_k$  itself is not well-defined or is singular. The well-definedness will be resolved by the Powell-type regularisation and the restarting of the iteration. The restarting of the iteration will also yield an algorithm that does not require unlimited memory. Lastly, this algorithm does not have the convergence property of the fixed point iteration algorithm. This will be resolved by the safeguarding of steps.

One can show that

**Proposition 2** (Rank-1 update for  $B_k$ ). We have

$$B_k = B_{k-1} + \frac{(y_{k-1} - B_{k-1} s_{k-1}) \hat{s}_{k-1}^{\top}}{\hat{s}_{k-1}^{\top} s_{k-1}}$$

where  $y_{k-1} = g_k - g_{k-1}$ ,  $B_0 = \text{Id}$  and

$$\hat{s}_{k-1} = s_{k-1} - \sum_{j=0}^{k-2} \frac{\hat{s}_j^{\top} s_{k-1}}{\|\hat{s}_j\|^2} \hat{s}_j$$

is the Gram-Schmidt orthogonalisation of  $s_{k-1} = x_k - x_{k-1}$ .

```
Input: x_0 \in \mathbb{R}^n and f: \mathbb{R}^n \to \mathbb{R}^n.

Set B_0 = \text{Id} and x_1 = f(x_0).

for k = 0, 1, \dots do
\begin{cases} \text{Set } g_k = g(x_k). \\ \text{Set } s_{k-1} = x_k - x_{k-1}, y_{k-1} = g_k - g_{k-1} \text{ and } \hat{s}_{k-1} = s_{k-1} - \sum_{i=0}^{k-2} \frac{\hat{s}_i^\top s_{k-1}}{\|\hat{s}_i\|^2} s_i. \end{cases}
\text{Set } B_k = B_{k-1} + \frac{(y_{k-1} - B_{k-1} s_{k-1}) \hat{s}_{k-1}^\top}{\hat{s}_{k-1}^\top s_{k-1}}.
\text{Set } H_k = B_k^{-1}.
\text{Set } x_{k+1} = x_k - H_k g_k.
```

Algorithm 6: AA-I (rank-1 update)

## Well-definedness of $H_k$ : Powell-type regularisation

To fix the singularity of  $B_k$  we use Powell-type regularisation.

```
s Input: x^0 \in \mathbb{R}^n, f: \mathbb{R}^n \to \mathbb{R}^n and \bar{\theta} \in (0,1).

Set B_0 = \text{Id} and x_1 = f(x_0).

for k = 0, 1, \dots do

Set g_k = g(x_k), s_{k-1} = x_k - x_{k-1} and y_{k-1} = g_k - g_{k-1}.

Set \hat{s}_{k-1} = s_{k-1} - \sum_{i=0}^{k-2} \frac{\hat{s}_i^{\top} s_{k-1}}{\|\hat{s}_i\|^2} s_i.

Choose \theta_{k-1} in dependence of \bar{\theta}.

Set \tilde{y}_{k-1} = \theta_{k-1} y_{k-1} - (1 - \theta_{k-1}) g_{k-1}.

Set B_k = B_{k-1} + \frac{(\tilde{y}_{k-1} - B_{k-1} s_{k-1}) \hat{s}_{k-1}^{\top}}{\hat{s}_{k-1}^{\top} s_{k-1}}.

Set H_k = B_k^{-1}.

Set x_{k+1} = x_k - H_k g_k.
```

Algorithm 7: AA-I with Powell-type regularisation

## Well-definedness of $H_k$ , memory usage: Restarting iteration

If  $\hat{s}_k = 0$  the update

$$B_k = B_{k-1} + \frac{(\tilde{y}_{k-1} - B_{k-1} s_{k-1}) \hat{s}_{k-1}^{\top}}{\hat{s}_{k-1}^{\top} s_{k-1}}$$

is ill-defined. This occurs in algorithm 7 e.g. for k > n as then  $\hat{s}_k = 0$  by linear dependence. Hence we restart the algorithm with  $x_k$  as the new starting point if

- k = m + 1 for some fixed  $m \in \mathbb{N}$  or
- $\|\hat{s}_{k-1}\| < \tau \|s_{k-1}\|$  for some fixed  $\tau \in (0,1)$ .

It can be shown that  $B_k$  is then well-defined.

Input: 
$$x^0 \in \mathbb{R}^n$$
,  $f: \mathbb{R}^n \to \mathbb{R}^n$ ,  $m \in \mathbb{N}$  and  $\bar{\theta}, \tau \in (0, 1)$   
Set  $B_0 = \operatorname{Id}, x_1 = f(x_0)$  and  $m_0 = 0$ .  
for  $k = 0, 1, \dots$  do
$$\begin{cases} \text{Set } g_k = g(x_k), m_k = m_{k-1} + 1, s_{k-1} = x_k - x_{k-1} \text{ and } y_{k-1} = g_k - g_{k-1}. \\ \text{Set } \hat{s}_{k-1} = s_{k-1} - \sum_{i=k-m_k}^{k-2} \frac{\hat{s}_i^\top s_{k-1}}{\|\hat{s}_i\|^2} s_i. \\ \text{if } m_k = m+1 \text{ or } \|\hat{s}_{k-1}\| < \tau \|s_{k-1}\| \text{ then } \\ | \text{Set } m_k = 0, \hat{s}_{k-1} = s_{k-1} \text{ and } B_{k-1} = \operatorname{Id}. \\ \text{end} \\ \text{Choose } \theta_{k-1} \text{ in dependence of } \bar{\theta}. \\ \text{Set } \tilde{y}_{k-1} = \theta_{k-1} y_{k-1} - (1 - \theta_{k-1}) g_{k-1}. \\ \text{Set } B_k = B_{k-1} + \frac{(\tilde{y}_{k-1} - B_{k-1} s_{k-1}) \hat{s}_{k-1}^\top}{\hat{s}_{k-1}^\top s_{k-1}}. \\ \text{Set } H_k = B_k^{-1}. \\ \text{Set } x_{k+1} = x_k - H_k g_k. \end{cases}$$

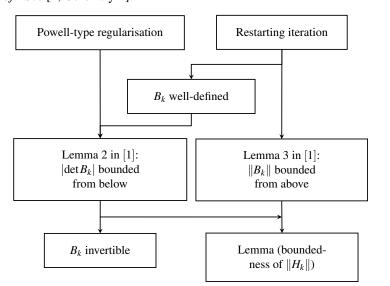
Algorithm 8: AA-I with Powell-type regularisation and Restarting

One can then show

**Lemma 3** (bound on  $||H_k||_2$ ). In algorithm 8 we have that  $H_k$  is well-defined and there exists a constant  $c_1 = c_1(m, n, \bar{\theta}, \tau) > 0$  such that

$$||H_k||_2 \leq c_1$$
.

Proof. See [1, Corollary 4].



## Computational efficiency: Rank-1 update for $H_k$

From the Sherman-Morrison formula one can obtain

**Proposition 4** (Rank-1 update for  $H_k$ ). We have

$$H_k = H_{k-1} + \frac{(s_{k-1} - H_{k-1} y_{k-1}) \hat{s}_{k-1}^{\top} H_{k-1}}{\hat{s}_{k-1}^{\top} H_{k-1} y_{k-1}}$$

Input: 
$$x^0 \in \mathbb{R}^n$$
,  $f : \mathbb{R}^n \to \mathbb{R}^n$ ,  $m \in \mathbb{N}$  and  $\bar{\theta}, \tau \in (0,1)$   
Set  $H_0 = \operatorname{Id}, x_1 = f(x_0)$  and  $m_0 = 0$ .  
for  $k = 0, 1, ...$  do

Set  $g_k = g(x_k)$ ,  $m_k = m_{k-1} + 1$ ,  $s_{k-1} = x_k - x_{k-1}$  and  $y_{k-1} = g_k - g_{k-1}$ .  
Set  $\hat{s}_{k-1} = s_{k-1} - \sum_{i=k-m_k}^{k-2} \frac{\hat{s}_i^{\top} s_{k-1}}{\|\hat{s}_i\|^2} s_i$ .  
if  $m_k = m+1$  or  $\|\hat{s}_{k-1}\| < \tau \|s_{k-1}\|$  then

Set  $m_k = 0$ ,  $\hat{s}_{k-1} = s_{k-1}$  and  $H_{k-1} = \operatorname{Id}$ .  
end

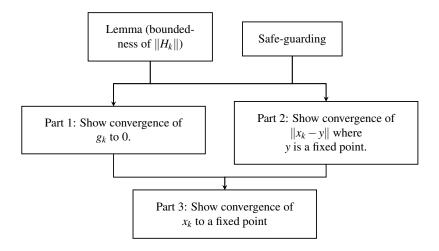
Choose  $\theta_{k-1}$  in dependence of  $\bar{\theta}$ .  
Set  $\tilde{y}_{k-1} = \theta_{k-1} y_{k-1} - (1 - \theta_{k-1}) g_{k-1}$ .  
Set  $H_k = H_{k-1} + \frac{(s_{k-1} - H_{k-1} \tilde{y}_{k-1}) \hat{s}_{k-1}^{\top} H_{k-1}}{\hat{s}_{k-1}^{\top} H_{k-1} \tilde{y}_{k-1}}$ .  
Set  $x_{k+1} = x_k - H_k g_k$ .

Algorithm 9: AA-I with Powell-type regularisation and Restarting

### **Convergence: Safeguarding steps**

To guarantee the decrease in  $||g_k||$  one can interleave the AA-I steps with Krasnosel'skii-Mann (KM) steps which are given by

$$x_{k+1} = (1 - \alpha)x_k + \alpha f_k$$



for some fixed  $\alpha \in (0,1)$ .

```
Input: x^0 \in \mathbb{R}^n, f: \mathbb{R}^n \to \mathbb{R}^n, m \in \mathbb{N}, \bar{\theta}, \tau, \alpha \in (0,1) and safe-guarding constants D, \varepsilon > 0

Set H_0 = \operatorname{Id}, x_1 = \bar{x}_1 = f(x_0), m_0 = n_{AA} = 0 and \bar{U} = \|g_0\|_2.

for k = 0, 1, \dots do
\begin{cases}
\operatorname{Set} g_k = g(x_k), m_k = m_{k-1} + 1, s_{k-1} = \bar{x}_k - x_{k-1} \text{ and } y_{k-1} = g(\bar{x}_k) - g_{k-1}.\\
\operatorname{Set} \hat{s}_{k-1} = s_{k-1} - \sum_{i=k-m_k}^{k-2} \frac{\hat{s}_i^{\top} \bar{s}_{k-1}}{\|\bar{s}_i\|^2} s_i.\\
\operatorname{if} m_k = m+1 \text{ or } \|\hat{s}_{k-1}\| < \tau \|s_{k-1}\| \text{ then}\\
| \operatorname{Set} m_k = 0, \hat{s}_{k-1} = s_{k-1} \text{ and } H_{k-1} = \operatorname{Id}.\\
\operatorname{end}
\end{cases}
Choose \theta_{k-1} in dependence of \bar{\theta}.

Set \tilde{y}_{k-1} = \theta_{k-1} y_{k-1} - (1 - \theta_{k-1}) g_{k-1}.\\
\operatorname{Set} H_k = H_{k-1} + \frac{(\bar{s}_{k-1} - H_{k-1} \bar{y}_{k-1}) \bar{s}_{k-1}^{\top} H_{k-1}}{\hat{s}_{k-1}^{\top} H_{k-1}} \text{ and } \tilde{x}_{k+1} = x_k - H_k g_k.\\
\operatorname{if} \|g_k\| \le D\bar{U}(n_{AA} + 1)^{-(1+\varepsilon)} \text{ then}\\
| \operatorname{Set} x_{k+1} = \tilde{x}_{k+1} \text{ and } n_{AA} = n_{AA} + 1.\\
\operatorname{else}\\
| \operatorname{Set} x_{k+1} = (1 - \alpha) x_k + \alpha f_k.\\
\text{end}
\end{cases}
```

Algorithm 10: AA-I with Powell-type regularisation, restarting and safeguarding

# Convergence result

**Theorem 5** (Convergence). Let  $x_k$  be generated by algorithm 10 then  $x_k$  converges to a fixed point of f.

*Proof.* The proof follows [1, Theorem 6]. In the first part we use lemma ?? on the boundededness of  $||H_k||$  and the safe-guarding step to show the convergence of  $g_k$  to 0. In part 2 we also use lemma ?? and the safe-guarding to show convergence of  $a_k = ||x_k - y||^2$  to some  $a \in \mathbb{R}$ . Here y denotes a fixed point of f. In the third part we use parts 1 and 2 to show that  $x_k$  converges to a fixed point.

Part 1. We partition  $\mathbb{N} = K_{AA} \sqcup K_{KM}$  where  $K_{AA} = \{m_0, m_1, ...\}$  denote the indices k where the algorithm chose an AA-step (a) and  $K_{KM} = \{n_0, n_1, ...\}$  where the algorithm chose a KM-step (b).

$$\begin{aligned} & \text{if } \|g_k\| \leq D\bar{U}(n_{AA}+1)^{-(1+\varepsilon)} \text{ then} \\ & | & \text{Set } x_{k+1} = \tilde{x}_{k+1} \text{ and } n_{AA} = n_{AA}+1. \\ & \text{else} \\ & | & \text{Set } x_{k+1} = (1-\alpha)x_k + \alpha f_k. \end{aligned} \tag{b}$$

**Algorithm 11:** The two cases for  $x_{k+1}$ .

Let y be a fixed point of f. We distinguish the cases

**case** (a) if  $k_i \in K_{AA}$  then

$$||x_{k_{i}+1} - y|| \le ||x_{k_{i}} - y|| + ||H_{k_{i}}g_{k_{i}}||$$

$$\le ||x_{k_{i}} - y|| + c_{1}||g_{k}||$$

$$\le ||x_{k_{i}} - y|| + c_{2}(i+1)^{-(1+\varepsilon)}$$
(1)

for some constant  $c_2 > 0$ .

**case** (b) if  $l_i \in K_{KM}$  then one can show (see [1, Theorem 6])

$$||x_{l_i+1} - y||^2 \le ||x_{l_i} - y||^2 - \alpha(1 - \alpha)||g_{l_i}||^2.$$
 (2)

Here one uses the non-expansiveness of f and that y is a fixed point.

Hence we have in any case

$$||x_k - y|| \le ||x_0 - y|| + c_2 \sum_i (i+1)^{-(1+\varepsilon)} = c_3 < \infty.$$

It then follows that

$$a_{k_{i}+1} = \|x_{k_{i}+1} - y\|^{2} \stackrel{(1),(2)}{\leq} \left( \|x_{k_{i}} - y\| + c_{2}(i+1)^{-(1+\varepsilon)} \right)^{2}$$

$$\leq \underbrace{\|x_{k_{i}} - y\|^{2}}_{=a_{k_{i}}} + \underbrace{c_{2}^{2}(i+1)^{-2(1+\varepsilon)} + 2c_{2}}_{=b_{k_{i}}} \underbrace{\|x_{k_{i}} - y\|}_{=b_{k_{i}}} (i+1)^{-(1+\varepsilon)}$$

$$= a_{k_{i}} + b_{k_{i}}$$

$$(3)$$

and thus

$$\alpha(1-\alpha)\sum_{i}||g_{l_{i}}||^{2} \stackrel{(2)}{\leq} \sum_{i}a_{l_{i}}-a_{l_{i}+1} \stackrel{(3)}{\leq} a_{0}+\sum_{k}b_{k}<\infty.$$

We therefore have  $\lim_i \|g_{l_i}\| = 0$ . It also follows from  $\|g_{k_i}\| \le D\bar{U}(i+1)^{-(1+\varepsilon)}$  that  $\lim_i \|g_{k_i}\| = 0$ . As all subsequences of  $g_k$  converge to 0 we thus have that  $g_k$  converges to 0.

Part 2. Let now  $a_{k_i}$  and  $a_{l_i}$  be subsequences such that  $k_i \leq l_i$  and

$$a_{k_i} \xrightarrow{j \to \infty} \liminf_k a_k = \underline{a} \quad \text{and} \quad a_{l_i} \xrightarrow{j \to \infty} \limsup_k a_k = \overline{a}.$$

It then follows that

$$a_{l_i} - a_{k_i} = \sum_{k=k_i}^{l_i-1} a_{k+1} - a_k \stackrel{(3)}{\leq} \sum_{k=k_i}^{\infty} b_k$$

and when taking  $l_i \rightarrow \infty$ 

$$\overline{a} - a_{k_i} \le \sum_{k=k_i}^{\infty} b_k$$

and then taking  $k_i \to \infty$  we get

$$\overline{a} - a \leq 0$$
.

Thus

$$\limsup_{k} a_k = \overline{a} \le \underline{a} = \liminf_{k} a_k$$

and so  $a_k = \|x_k - y\|^2$  converges to some  $a \in \mathbb{R}$ .

*Part 3.* Let  $k_i$  and  $l_i$  now be convergent subsequences of  $x_k$  which converge to x and  $\tilde{x}$  respectively. Since by continuity of g

$$||g(x)|| = \lim_{i} ||g(x_{k_i})|| \stackrel{\text{part } 1}{=} 0$$

we have that x is a fixed point and analogously  $\tilde{x}$  is too. Now we have

$$||x_{k_i}||^2 = ||x_{k_i} - y||^2 - ||y||^2 + 2y^{\top} x_{k_i}$$

and by part 2 we obtain when taking  $i \to \infty$ 

$$||x||^2 = a - ||y||^2 + 2y^{\top}x$$

Analogously we obtain

$$\|\tilde{x}\|^2 = a - \|y\|^2 + 2y^{\top}\tilde{x}$$

which implies

$$2y^{\top}(x-\tilde{x}) = ||x||^2 - ||\tilde{x}||^2.$$

As x and  $\tilde{x}$  are fixed points it follows for  $y \in \{x, \tilde{x}\}$  that

$$x^{\top}(x-\tilde{x}) = \tilde{x}^{\top}(x-\tilde{x})$$

and further

$$(x-\tilde{x})^{\top}(x-\tilde{x})=0.$$

We thus have  $x = \tilde{x}$ . We have shown that two convergent subsequences of  $x_k$  have the same limit and hence  $x_k$  is convergent and the limit must be a fixed point of f.

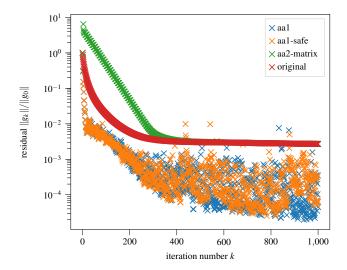


Figure 1: Residual norms for the elastic net regression problem.

## **Numerical experiments**

As part of the project some numerical experiments from [1] were replicated. Thus the obtained results are the same as in [1]. The aim of the experiments is to test the numerical performance of the alorithms. For this f is chosen so that the Lipshitz constant is close to 1. This is precisely the type of problem for which the AA algorithm was developed.

#### **Elastic net regression**

In the first experiment f originates from an elastic net regression problem and is motivated in [1, Section 5.1f]. Specifically one obtains

$$f: \mathbb{R}^n \to \mathbb{R}^n, \qquad x \mapsto S_{\alpha\mu/2} \left( x - \alpha \left( A^\top (Ax - b) + \frac{\mu}{2} x \right) \right)$$

with shrinkage operator

$$S_{\kappa}(x) = \left(\operatorname{sgn}(x_i)(|x_i| - \kappa)_+\right)_{i=1}^n$$

and  $A \in \mathbb{R}^{m \times n}$ ,  $b \in \mathbb{R}^m$  and some  $\alpha, \mu \in \mathbb{R}$ . Here we choose the parameters as in [1, Section 5.2]. In particular we choose m = 500 and n = 1000 and A, b and  $x^0$  to be randomly generated.

The results for the different methods can be seen in Figure 1. Here the method 'original' is the fixed point iteration, i.e. algorithm ??. The method 'aa1-safe' is the AA-I algorithm with Powell-type regularisation, restarting and safeguarding. The 'aa2-matrix' algorithm is an implementation of the AA-II algorithm where 'matrix' indicates that the implementation is not matrix-free. We see little difference between the 'aa1' and the

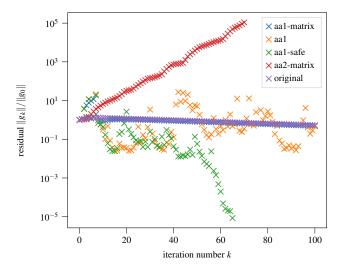


Figure 2: Residual norms for the Markov decision process problem.

'aa1-safe' methods. Both methods however outperform the 'original' and 'aa2-matrix' methods for this problem.

## Markov decision process

In a second experiment f originates from a random Markov decision process which is motivated in [1, Section 5.1f]. Here our aim is to find a fixed point of the Bellman operator

$$f \colon \mathbb{R}^n \to \mathbb{R}^n, \qquad x \mapsto \left( \max_a \left( R_{sa} + \gamma \sum_{s'} P_{sas'} x_{s'} \right) \right)_{s=1}^n$$

with some  $R \in \mathbb{R}^{S \times A}$ ,  $P \in \mathbb{R}^{S \times A \times A}$  and  $\gamma \in \mathbb{R}$ . Here the parameter  $\gamma$  determines the Lipshitz-constant of f. Again, we choose the parameters as in [1, Section 5.2]. In particular we have n = 1000, A = 200, S = 300 and A and R to be randomly generated.

The performance for the various methods can be seen in figure 2. In contrast to the 'aa1-matrix' method the the 'aa2-matrix' method does not converge here. In this problem the fixed point iteration ('original') converges very slowly. The 'aa1-safe' method outperforms all others. This confirms numerically that the 'aa1-safe' algorithm can deal with the problems it was specifically designed for and for which the fixed point iteration fails.

In figure 3 we see how the performance of the 'aa1-safe' method depends on the memory parameter m. In particular one sees that the algorithm performs best for this problem with a parameter of  $m \approx 10$ . We also see that increasing the parameter m does not necessarily improve performance of the method as in this plot the choice m = 50 performs worst.

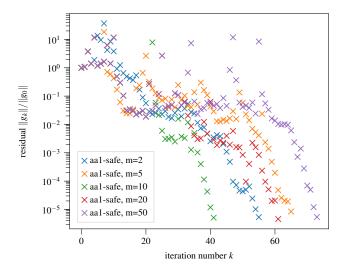


Figure 3: Residual norms for the Markov decision process problem.

## **Summary**

The AA-I algorithm is specifically taylored to find the fixed point of a function f which is expensive to evaluate, noisy, has an unknown gradient and where the dimension n is large. The main idea of the AA-I and AA-II algorithms is to generalise the fixed point iteration by setting  $x_{k+1} = \sum_i \alpha_i f_i$  for some clever choice of  $\alpha = \alpha^k \in \mathbb{R}^{k+1}$ . The AA-I algorithm one obtains requires some modifications. More specifically, one applies Powell-type regularisation and a restarting of the iteration for well-definedness. One builds in a mechanism for safeguarding of the steps for convergence and one uses a rank-1 update formula to make the implementation matrix-free. One can then show the convergence of the algorithm under the assumption that f is non-expansive and that there exists an algorithm. The numerical experiments then show that the AA-I algorithm with the modifications outperforms the fixed point iteration for the problems tested.

## **Bibliography**

#### Main source

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