

EcoSearch: A Constant-Delay Best-First Search Algorithm for Program Synthesis

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Program Synthesis in the wild

	A	B	C
1	Name	First	Last
2	Ned Lanning	Ned	
3	Margo Hendrix	Margo	
4	Dianne Pugh	Dianne	
5	Earlene McCarty	Earlene	
6	Jon Voigt	Jon	
7	Mia Arnold	Mia	

Copyright Microsoft for syntactic manipulation,
based on papers by *S. Gulwani*

In practice

Logic:

$$\begin{aligned}\forall a, b \\ f(a, b) &\geq a \\ f(a, b) &\geq b \\ f(a, b) &\in \{a, b\}\end{aligned}$$

Specifications:

Examples:

$$\begin{aligned}f(1, 5) &= 5 \\ f(2, 1) &= 2 \\ f(-3, -9) &= -3\end{aligned}$$

Natural language:

‘Write a function that takes the maximum of its two arguments.’

Produced Algorithms:

```
def max(a: int, b: int) ->int:
    if a <=b:
        return b
    else:
        return a
```

Program Synthesis: Problem

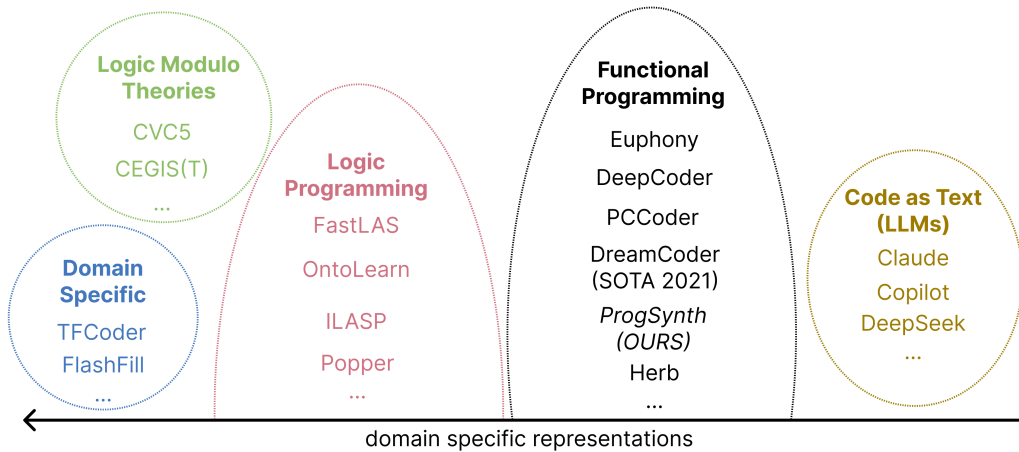
Input:

- the search space G :
a deterministic tree grammar
- a specification \mathcal{C} :
it checks if a program $p \in \mathcal{L}(G)$ matches the specification

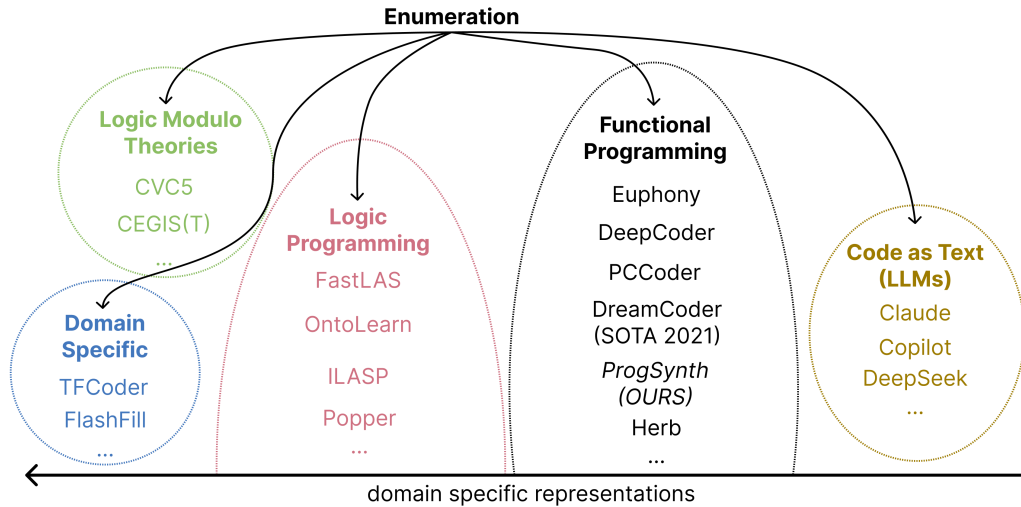
Output:

- a program in the search space that matches the specification:
a $p \in \mathcal{L}(G)$ such that $\mathcal{C}(p) = \checkmark$

Frameworks



Frameworks



Program Synthesis in the age of LLMs?

A reviewer said "isn't this research now useless?".
No, we do not solve the same problem.

LLM: takes natural language as input

Program Synthesis: takes a checker as input

Enumeration can also be used on top of LLMs

Enumeration Problem

Input:

- the search space G :
a deterministic tree grammar with a cost for each tree
- a specification \mathcal{C} :
it checks if a program $p \in \mathcal{L}(G)$ matches the specification

Goal:

Enumerate all programs in order of non-decreasing costs

Delay:

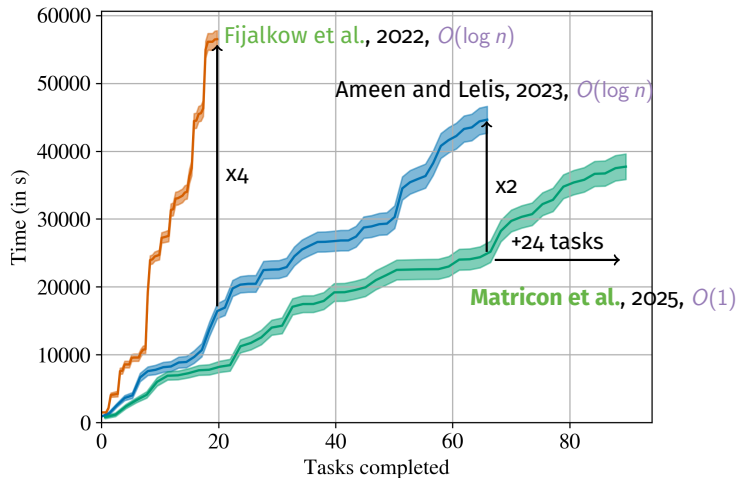
Time complexity in terms of n : number of programs enumerated

Between enumeration of the n^{th} program and the next

Overview

Major papers:

- 2017, machine learning + enumeration, *Balog et al.*, ICLR
- 2018, $O(\log n)$, *Lee et al.*, PLDI



Skeleton of an enumeration algorithm

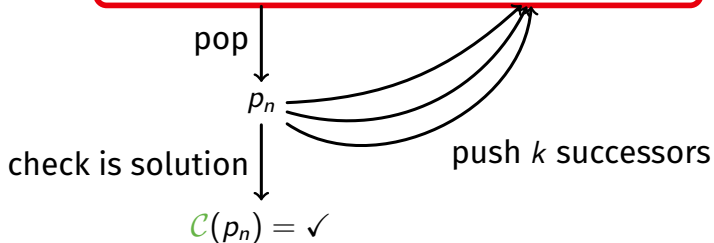
priority queue: S pairs of (program, cost)



At step n :

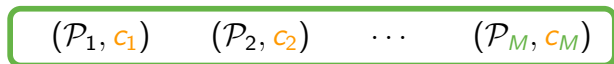
$$S = O(n)$$

$$k = O(1)$$



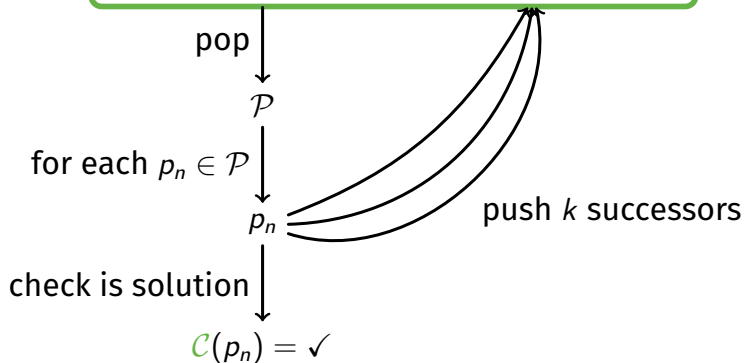
Our enumeration algorithm

bucket queue: M buckets of (programs, cost)



At step n :

$M = O(1)$
 $k = O(1)$



Our contribution

We prove bounded differences in cost:

$$\exists M, \forall n, \text{cost_next}(p_n) - \text{cost}(p_n) \leq M$$

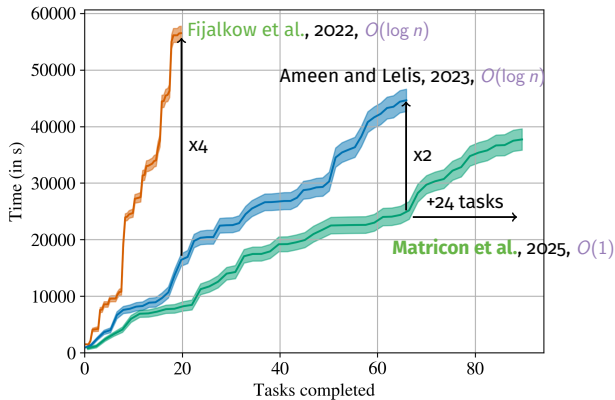
This implies: **priority queues** $O(\log n)$ \rightarrow **bucket queues** $O(1)$.

Impact

Fastest ranked enumeration for program synthesis in practice.

First algorithm with $O(1)$ delay \rightarrow closes open question.

Futures



- Combine with LLMs
- Implement in other Program Synthesis solvers
- Memoryless version?
- Parallelized version? GPU?

