

3rd TECHNIQUE (NOT ALWAYS WORKING): HORN RESOLUTION

"IT DOESN'T COME FROM HORN, HORN WAS A MATHEMATICIAN" - LUCA MAVA, 29/11/2023

$$CCF = \{ \{ R(x, f(x)) \}, \{ \neg R_{xy}, R_{yx} \}, \{ \neg R_{xy}, \neg R_{yz}, R_{xz} \}, \{ \neg R_{aa} \} \}$$

A CLAUSE IS OF HORN TYPE IF IT HAS AT MOST ONE POSITIVE LITERAL

IF WE KNOW THE SET HAS A CLOSED EXPANSION, WE CAN CHOOSE ALGORITHMICALLY

ONE OF THE NEGATIVE LITERAL. WE'LL ALWAYS USE THE MOST LEFT LITERAL

$$\{ \neg R_{aa} \}, \{ \neg R_{xy}, \neg R_{yz}, R_{xz} \}$$

$$\{ \neg R_{ax}, \neg R_{ya} \}$$

$$\{ R(x, f(x)) \}$$

$$\{ \neg R(fa, a) \}$$

$$\{ \neg R_{xy}, R_{yx} \}$$

$$\{ \neg R(a, fa) \}$$

$$\{ R(x, f(x)) \}$$

$$\emptyset$$

USE RESOLUTION TO PROVE THE FOLLOWING FACTS:

a) EVERY REFLEXIVE, EUCLIDEAN BINARY RELATION IS AN EQUIVALENCE RELATION

• REFLEXIVE $\forall x R_x$

• EUCLIDEAN $\forall x, y, z (R_{xy} \wedge R_{xz} \rightarrow R_{yz})$

• EQUIVALENCE RELATION:

- REFLEXIVE ✓ (IN PREMISES)



- SYMMETRIC $\forall x, y (R_{xy} \rightarrow R_{yx})$

- TRANSITIVE $\forall x, y, z (R_{xy} \wedge R_{yz} \rightarrow R_{xz})$

$\forall x (R_{xx}),$

$\{\forall x, y, z (R_{xy} \wedge R_{xz} \rightarrow R_{yz}) \vdash \forall x, y (R_{xy} \rightarrow R_{yx}), \forall x, y, z (R_{xy} \wedge R_{yz} \rightarrow R_{xz})\}$

WE FIRST VERIFY THE FIRST FORMULA IN THE CONCLUSIONS, THAT IS

$\forall x (R_{xx}), \forall x, y, z (R_{xy} \wedge R_{xz} \rightarrow R_{yz}) \vdash \forall x, y (R_{xy} \rightarrow R_{yx})$

$\{\forall x (R_{xx}), \forall x, y, z (R_{xy} \wedge R_{xz} \rightarrow R_{yz}), \neg \forall x, y (R_{xy} \rightarrow R_{yx})\}$

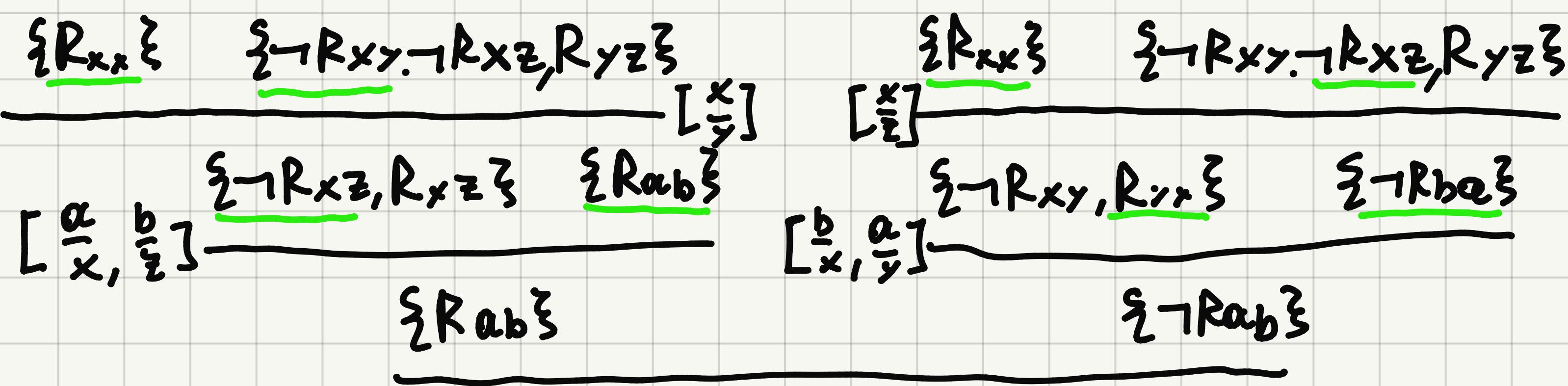
UNIFICATION

• $\{\forall x (R_{xx})\} \vdash \{R_{xx}\}$

• $\{\forall x, y, z (R_{xy} \wedge R_{xz} \rightarrow R_{yz})\} \vdash \{\neg R_{xy}, \neg R_{xz}, R_{yz}\}$

• $\{\neg \forall x, y (R_{xy} \rightarrow R_{yx})\} \vdash \{\exists x, y. \neg (R_{xy} \rightarrow R_{yx})\} \rightarrow \{R_{ab}\}, \{\neg R_{ba}\}$

$L(F) = \{\{R_{xx}\}, \{\neg R_{xy}, \neg R_{xz}, R_{yz}\}, \{R_{ab}\}, \{\neg R_{ba}\}\}$



Now we can consider it with the set of premises and prove

$\{ \forall x R_{xx}, \forall y z (R_{xy} \wedge R_{xz} \rightarrow R_{yz}), \forall xy (R_{xy} \rightarrow R_{yx}) \wedge \forall xyz (R_{xy} \wedge R_{yz} \rightarrow R_{xz}) \}$

HERBRAND

STEP	FORMULA	REASON
1	$\{\forall x R_{xx}\}$	ASSUMPTION
2	$\{\forall y z (R_{xy} \wedge R_{xz} \rightarrow R_{yz})\}$	ASSUMPTION
3	$\{\forall y (R_{xy} \rightarrow R_{yx})\}$	ASSUMPTION
4	$\{\neg \forall y z (R_{xy} \wedge R_{yz} \rightarrow R_{xz})\}$	ASSUMPTION
5	$\{\neg R_{aa}\}$	I, Y-EXPANSION
6	$\{\neg \forall y z (R_{ay} \wedge R_{az} \rightarrow R_{yz})\}$	ZY-EXPANSION
7	$\{\neg \forall z (R_{ab} \wedge R_{az} \rightarrow R_{bz})\}$	6, Y-EXPANSION
8	$\{R_{ab} \wedge R_{az} \rightarrow R_{bz}\}$	7, Y-EXPANSION
9	$\{\neg (R_{ab} \wedge R_{az}), R_{bz}\}$	8, P-EXPANSION
10	$\{\neg R_{ab}, \neg R_{az}, R_{bz}\}$	9, P-EXPANSION