

# EXTRA: SKOLEMIZATION EXERCISES. SKOLEMIZE THE FOLLOWING:

a)  $\exists x \underline{P(x, y)} \rightarrow Q(x)$        $\exists \Rightarrow x = f(y)$

$$P(f(y), y) \rightarrow Q(f(y))$$

b)  $\forall x \exists y (P(x, y) \wedge Q(y, z))$        $\exists \Rightarrow y = f(x, z)$  UNCHANGED

$$\forall x (P(\underline{x}, f(x, z)) \wedge Q(f(x, z), z)) \mapsto P(x, f(x, z)) \wedge Q(f(x, z), z)$$

c)  $\forall x \exists y P(x, y) \wedge \exists x Q(x)$   
 $\forall \Rightarrow x = a \quad f(a) = b$

$$\forall x P(x, f(x)) \wedge \exists \underline{Q(x)} \mapsto P(a, b) \wedge Q(c)$$

d)  $\exists x \forall y (P(x, y) \vee Q(y))$        $f(x) = a$  UNCHANGED

$$\forall y (P(a, y) \vee Q(y)) \mapsto P(a, y) \vee Q(y)$$

e)  $\exists x (P(x) \rightarrow \forall y Q(x, y))$        $f(x) = a$  UNCHANGED

$$P(a) \rightarrow \forall y Q(a, y) \mapsto P(a) \rightarrow Q(a, y)$$

f)  $\exists x \forall y (P(x, y) \wedge \exists z Q(y, z))$        $f(x) = a \quad f(z) = b$

$$\exists x \forall y (P(x, y)) \wedge \exists \underline{z} \forall y \underline{Q(y, z)}$$

$$\forall y (P(a, y)) \wedge \forall y Q(y, b)$$

$$P(a, y) \wedge Q(y, b)$$

PROVE THAT THE FOLLOWING FORMULAS ARE SATISFIABLE BUT NOT VALID

a)  $\forall x \forall y (\forall x y \rightarrow \exists z (\forall x z \wedge \forall y z))$

• SATISFIABILITY

$$\forall x \forall y (\forall x y \rightarrow \exists z (\forall x z \wedge \forall y z))$$

$$\begin{array}{c} / \quad \backslash \\ \forall x (\neg \forall y) \quad \forall x \exists (\exists z (\forall x z \wedge \forall y z)) \\ | \quad \checkmark \\ \neg R_{\alpha \alpha} \end{array}$$

$\Rightarrow$  SATISFIABLE

• VALIDITY

$$\neg (\forall x \forall y (\forall x y \rightarrow \exists z (\forall x z \wedge \forall y z)))$$

$$\begin{array}{c} | \\ \forall x (\forall y), \forall x (\exists z (\neg \forall x z)) \\ | \\ R_{\alpha \alpha}, \neg R_{\alpha \beta} \end{array}$$

$\checkmark$

$\Rightarrow \neg \varphi$  IS SATISFIABLE  
 $\Rightarrow$  NOT VALID

$$b) \forall x \exists y. R_{xy} \rightarrow \exists y \forall x. R_{xy}$$

- SATISFIABILITY

$$\begin{array}{c}
 \forall x \exists y R_{xy} \rightarrow \exists y \forall x R_{xy} \\
 \diagdown \qquad \qquad \qquad \diagup \\
 \neg (\forall x \exists y R_{xy}) \qquad \exists y \forall x R_{xy} \\
 | \\
 \neg (\exists y R_{ay}) \\
 | \\
 \neg R_{aa} \text{ SATISFIABLE}
 \end{array}$$

- VALIDITY

$$\neg (\forall x \exists y R_{xy} \rightarrow \exists y \forall x R_{xy}) \equiv \forall x \exists y R_{xy} \wedge \neg \exists y \forall x R_{xy}$$

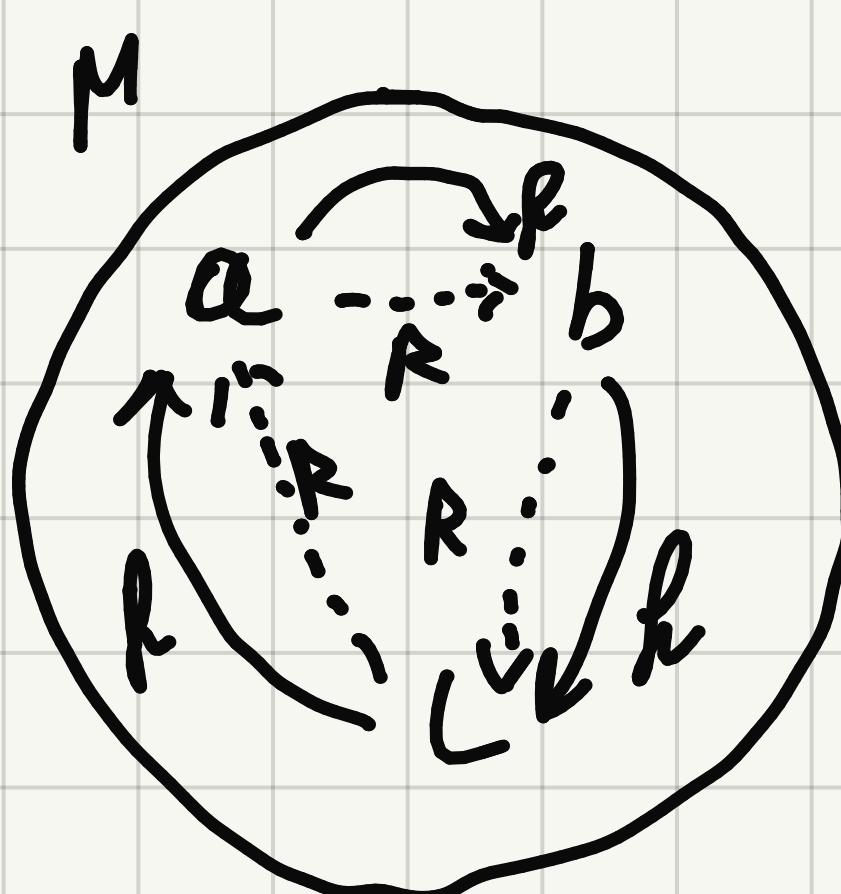
$$\equiv \forall x \exists y R_{xy} \wedge \forall x \exists y \neg R_{yx} \mapsto \forall x (R_{xy} \wedge \neg R_{yx})$$

FOR SIMPLICITY. SUPPOSE  $f(x) = g(x) \mapsto x R f(x) \wedge \neg f(x) R x$

$x = \alpha$   $f(\alpha)$ ?  $\neq \alpha$ . OTHERWISE,  $\alpha R \alpha \wedge \neg \alpha R \alpha$

$\rightarrow f(\alpha) = b$   $f(b)$ ?  $\neq \alpha$   $b R \alpha \wedge \neg \alpha R b$

$f(b) = c$   $f(c)$ ?  $= \alpha$   $b R c \wedge \neg c R \alpha \checkmark$



$$\rightarrow M = \{a, b, c\}$$

$$[[R]] = \{(a, b), (b, c), (c, a)\}$$

$$[[f]] = \{(a, b), (b, c), (c, a)\}$$

THAT MAKES  $\neg \varphi$  SATISFIABLE  $\Rightarrow \varphi$  IS NOT VALID