Graphs, III

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Outline

1. Backgrounds
2. Eigenvalues
3. Information quickly spreads in expander
4. Combinatorial characterisation
5. Expander $\approx$ pseudo-random generator
6. Algorithm for constructing expanders
7. UPATH is in Logspace

Why expanders?

- Communication networks
- Pseudo random generator
- Randomness
- Derandomisation
- UPATH is Log space
- PageRank

Conventions

For simplicity, we assume that the graphs allow:

- regular
- selfloop
- parallel edges

Theory is possible for general graphs without these assumptions.

Inner product

$\langle u, v \rangle$

- $\langle xu + yv, w \rangle = x\langle u, w \rangle + y\langle v, w \rangle$
- $\langle v, u \rangle = \overline{\langle u, v \rangle}$, $\bar{z}$ is the complex conjugation of z
- For all u, $\langle u, u \rangle \geq 0$, with 0 only if $u = 0$
- $\langle u, v \rangle = 0$ means u, v are orthogonal, written $u \perp v$
- If $u^1, u^2, \cdots, u^n$ satisfy $u^i \perp u^j$ for all $i \neq j$, then they are linearly independent.

Parseval's identity: If $u^1, u^2, \cdots, u^n$ form an orthonormal basis for C^n, then for every v, if $v = \sum_i \alpha_i u^i$, then

$$\langle v, v \rangle = \sum_{i=1}^{n} |\alpha_i|^2.$$

Hilbert space: Vector spaces with inner product.

Dot product

- For $u, v \in \mathbb{F}^n$, $u \odot v = \sum_{i=1}^{n} u_i v_i$
- $S \subset \mathbb{F}^n$, $S^\perp = \{u : u \perp S\}$
- $u \perp v$, if $u \odot v = 0$, $u \perp S$, if for all $v \in S$, $u \perp v$.
- $\dim(S) + \dim(S^\perp) = n$
- $u \in \mathbb{F}^n$, $u^\perp = \{v : v \perp u\}$, and $\dim(u^\perp) = n - 1$.

Random subsum principle

For every nonzero $u \in \mathrm{GF}(2^n)$,

$$\Pr_{v \in_{\mathrm{R}} \mathrm{GF}(2^n)} [u \odot v = 0] = \frac{1}{2}.$$

Eigenvectors and eigenvalues

If A is a real, symmetric matrix, for λ and v, if $Av = \lambda v$, then

$$\lambda\langle v, v\rangle = \langle Av, v\rangle = \overline{\langle v, Av\rangle} = \overline{\langle v, \lambda v\rangle} = \bar{\lambda}\langle v, v\rangle$$

$$\Downarrow$$

$$\lambda = \bar{\lambda}$$

so λ is a real.

Norms

$$|| \, || : \mathbb{F}^n \to \mathbb{R}^{\geq 0}$$

(i) $||v|| = 0 \iff v = 0$

(ii) $||\alpha v|| = |\alpha| \cdot ||v||$

(iii) $||u + v|| \leq ||u|| + ||v||$.

L_p-norm

L_p-norm of v, $p \geq 1$,

$$||v||_p = (\sum_{i=1}^{n} |v_i|^p)^{1/p}$$

$p = 2$ – the Euclidean norm

$$||v||_2 = (\sum_{i=1}^{n} |v_i|^2)^{1/2}$$

$p = 1$,

$$||v||_1 = \sum_{i=1}^{n} |v_i|$$

$p = \infty$,

$$||v||_\infty = \max_i |v_i|.$$

Hölder inequality

For every p, q, if $\frac{1}{p} + \frac{1}{q} = 1$, then

$$\|u\|_p \cdot \|v\|_q \geq \sum_{i=1}^{n} |u_i v_i|.$$

$p = q = 2$, Cauch-Schwarz

L_1- and L_2-norms

For every vector $v \in \mathbb{R}^n$,

$$\frac{|v|_1}{\sqrt{n}} \leq ||v||_2 \leq |v|_1.$$

Adjacent matrix

- G: d-regular, n vertices,
- p: a column vector, a distribution over the vertices of G
- A_{ij}: $\frac{n_{ij}}{d}$, where n_{ij} the number of edges between i and j.
- A: the adjacent matrix. It is normalised, symmetric, stochastic
- $q = Ap$: the distribution of a random walk in G from distribution p.
- $A^l e^i$: the distribution of l-step random walk from node i
- 1: the transpose of $(\frac{1}{n}, \cdots, \frac{1}{n})$, the uniform distribution
- $1^{\perp}$: $\{v : v \perp 1\}$
- $v \perp 1 \iff \sum v_i = 0$.

$\lambda(A)$

Define

$$\lambda(A) = \lambda(G) = \max\{||Av||_2 : \ ||v||_2 = 1, v \perp 1\}.$$

Suppose that

$$\lambda_1, \lambda_2, \cdots, \lambda_n$$

are the eigenvalues of A with orthogonal eigenvectors

$$v^1, v^2, \cdots, v^n$$

respectively.
Let $|\lambda_1| \geq |\lambda_2| \geq \cdots \geq |\lambda_n|$.

$|\lambda_i| \leq 1$

For λ and v such that $Av = \lambda v$. Then $\lambda = \frac{\langle v, Av \rangle}{\langle v, v \rangle}$.
By definition,

$$\langle v, Av \rangle = \sum_{i=1}^{n} a_{ii} v_i^2 + 2 \sum_{i<j, i \sim j} a_{ij} v_i v_j$$

For $i < j$, $i \sim j$:

$$a_{ij}(v_i - v_j)^2 = a_{ij} v_i^2 - 2 a_{ij} v_i v_j + a_{ij} v_j^2$$

Summing up all such i, j's:

$$\sum_{i=1}^{n} (1 - a_{ii}) v_i^2 - 2 \sum_{i<j, i \sim j} a_{ij} v_i v_j$$

Proof - I

$$\begin{aligned}\langle v, Av\rangle &= \sum_{i=1}^{n} a_{ii}v_i^2 + \sum_{i=1}^{n}(1-a_{ii})v_i^2 - \sum_{i<j, i\sim j} a_{ij}(v_i - v_j)^2 \\ &= \sum_{i=1}^{n} v_i^2 - \sum_{i<j, i\sim j} a_{ij}(v_i - v_j)^2 \leq \sum_{i=1}^{n} v_i^2. \qquad (1)\end{aligned}$$

And:

$$\sum_{i=1}^{n} v_i^2 - \sum_{i<j, i\sim j} a_{ij}(v_i - v_j)^2 \geq -\sum_{i=1}^{n} v_i^2.$$

Therefore $-1 \leq \lambda \leq 1$.
By definition, $A1 = 1$.
So $\lambda_1 = 1$, and 1 is the eigenvector of $\lambda_1 = 1$.
By the choice of the eigenvectors, $1^{\perp} = \mathrm{Span}\{v^2, \cdots, v^n\}$.

Proof - II

Given v, with $v \perp 1$, $||v||_2 = 1$.
Let $v = \alpha_2 v^2 + \cdots + \alpha_n v^n$ with $\alpha_2^2 + \cdots + \alpha_n^2 = 1$.

$$Av = \alpha_2 A v^2 + \cdots + \alpha_n A v^n = \alpha_2 \lambda_2 v^2 + \cdots + \alpha_n \lambda_n v^n$$

$$||Av||_2^2 = \alpha_2^2 \lambda_2^2 + \cdots + \alpha_n^2 \lambda_n^2$$

Since $\lambda_2^2 \geq \cdots \geq \lambda_n^2$,

$$\max ||Av||_2^2 = \lambda_2^2.$$

Therefore

$$\lambda = \lambda(G) = |\lambda_2|.$$

Spectral gap

We call $1 - \lambda(G)$ the *spectral gap of G*.

Lemma 1

Let G be an n-vertex regular graph and p a probability distribution over G's vertices. Then,

$$||A^l p - 1||_2 \leq \lambda^l.$$

Proofs consist of the following items:
1) By definition of $\lambda = \lambda(G)$, for every $v \perp 1$,

$$||Av||_2 \leq \lambda ||v||_2.$$

Proofs - I

2) If $v \perp 1$, then so is Av.

$$\langle 1, Av \rangle = \langle A^{\mathrm{T}} 1, v \rangle = \langle 1, v \rangle = 0.$$

Note $A = A^{\mathrm{T}}$, and $A1 = 1$.
3) $A: \ 1^{\perp} \to 1^{\perp}$, and
A shrinks each $v \in 1^{\perp}$ by at least λ factor in L_2 norm.
4) By 3), A^l shrinks each $v \in 1^{\perp}$ by at least λ^l factor, giving

$$\lambda(A^l) \leq \lambda^l.$$

Proofs - II

5) Let $p = \alpha 1 + p'$, $p' \perp 1$, Since $p' \perp 1$, $\sum p'_i = 0$.
But $\sum p_i = 1$, so $\alpha = 1$.

$$A^l p = A^l(1 + p') = A^l 1 + A^l p' = 1 + A^l p'.$$

$$\begin{aligned} ||A^l p - 1||_2 &= ||A^l p'||_2 \\ &\leq ||A^l||_2 \cdot ||p'||_2 \\ &\leq \lambda^l \cdot ||p'||_2 \\ &\leq \lambda^l \cdot ||p||_2 \\ &\leq \lambda^l \cdot |p|_1 = \lambda^l. \end{aligned} \tag{2}$$

The third inequality uses $||p||_2^2 = ||1||_2^2 + ||p'||_2^2$.

Log space algorithm for connectivity in expanders

Suppose that λ is a constant significantly smaller than 1.
By the lemma above, let $l = O(\log n)$.
Then $\lambda^l \approx 0$. Therefore

$$A^l p \approx 1.$$

This means that for any two nodes i, j, the distance between i and j is within $O(\log n)$.
According to this property, we are able to design a log space algorithm to decide, for any two vertices, whether or not, they are connected.
The algorithm simply enumerates all the paths from i of length $O(\log n)$, to see if there is a path passes j. The enumeration of all the paths can be done in log space.

Randomized log space for connectivity

Lemma 2

If G is a regular connected graph with selfloop at each vertex, then

$$\lambda(G) \leq 1 - \frac{1}{4dn^2}.$$

Let $u \perp 1$, $||u||_2 = 1$.
We show that $||Au||_2 \leq 1 - \frac{1}{4dn^2}$.
Let $v = Au$. It suffices to show that $1 - ||v||_2^2 \geq \frac{1}{2dn^2}$.
Since $||u||_2 = 1$,

$$1 - ||v||_2^2 = ||u||_2^2 - ||v||_2^2.$$

Considering $\sum_{i,j} A_{ij}(u_i - v_j)^2$, we have

Proofs - I

$$
\begin{aligned}
\sum_{i,j} A_{ij}(u_i - v_j)^2 &= \sum_{i,j} A_{ij}u_i^2 - 2\sum_{i,j} A_{ij}u_i v_j + \sum_{i,j} A_{ij}v_j^2 \\
&= \sum_{i=1}^{n} u_i^2 - 2\langle Au, v\rangle + \sum_{j=1}^{n} v_j^2 \\
&= ||u||_2^2 - 2\langle Au, v\rangle + ||v||_2^2 \\
&= ||u||_2^2 - 2||v||_2^2 + ||v||_2^2 \\
&= ||u||_2^2 - ||v||_2^2 = 1 - ||v||_2^2. \qquad (3)
\end{aligned}
$$

Therefore, we only need to prove

$$\sum_{i,j} A_{ij}(u_i - v_j)^2 \geq \epsilon = \frac{1}{2dn^2}.$$

Proofs - II

By the choice of u, $\sum u_i = 0$, and $\sum u_i^2 = 1$. So there exist i, j such that $u_i u_j < 0$.
Since $||u||_2 = 1$, the average of u_i^2 is $\frac{1}{n}$, and the average of $|u_i|$ is $\frac{1}{\sqrt{n}}$.
Let i and j be such that $u_i > 0$, $u_j < 0$, and

$$u_i - u_j \geq \frac{1}{\sqrt{n}}.$$

(Such i, j are guaranteed to exist, as above)

Proofs - III

Because G is connected, there is a path P between i and j. Suppose that the path P is labelled by $1, 2, \cdots, D+1$. Then:

$$\begin{aligned}
\frac{1}{\sqrt{n}} &\leq u_1 - u_{D+1} \\
&= (u_1 - v_1) + (v_1 - u_2) + (u_2 - v_2) + \cdots + (v_D - u_{D+1}) \\
&\leq |u_1 - v_1| + |v_1 - u_2| + \cdots + |v_D - u_{D+1}| \\
&\leq \sqrt{(u_1 - v_1)^2 + (v_1 - u_2)^2 + \cdots + (v_D - u_{D+1})^2} \cdot \sqrt{2D + 1} \qquad (4)
\end{aligned}$$

Proofs - IV

Since $A_{ii}, A_{ii+1} \geq \frac{1}{d}$,

$$\begin{aligned}
&\sum_{i,j} A_{ij}(u_i - v_j)^2 \\
\geq\ & \frac{1}{d} \cdot [(u_1 - v_1)^2 + (v_1 - u_2)^2 + \cdots + (v_D - u_{D+1})^2] \\
\geq\ & \frac{1}{dn(2D+1)} \geq \frac{1}{2dn^2}. \qquad (5)
\end{aligned}$$

Random walk lemma

Lemma 3

Let G be a d-regular n-vertex graph with all vertices having a selfloop. Let s be a vertex in G. Let $l > \Omega(dn^2 \log n)$, and X_l be the distribution of the vertex of the lth step in a random walk from s. Then for every t,

$$\Pr[X_l = t] > \frac{1}{2n}.$$

Proofs

By the previous lemma,

$$||A^l p - 1||_2 \leq (1 - \frac{1}{4dn^2})^{\Omega(dn^2 \log n)} < \frac{1}{n^\alpha}$$

for some constant α.
Choose α such that for $q = A^l p$,

$$|q - 1|_1 < \frac{1}{n^2}.$$

Therefore, the probability that $X_l = t$ is at least

$$\frac{1}{n} - \frac{1}{n^2} \geq \frac{1}{2n}.$$

Run the l-step random walks for $O(n \log n)$ many times, almost surely, every vertex is visited.
This gives a randomized log space algorithm to decide the connectivity of two vertices.

(n, d, λ)-expander graph

Definition 4

It is an n-vertex, d-regular graph G, satisfying $\lambda(G) \leq \lambda$ for some $\lambda < 1$.

A family of graphs $\{G_n\}$ is an expander family, if there exist d, $\lambda < 1$ such that for every n, G_n is an (n, d, λ)-expander graph.

(n, d, ρ)-combinatorial edge expander

For every S, $|S| \le \frac{n}{2}$,

$$|E(S, \bar{S}| \ge \rho \cdot d \cdot |S|.$$

Theorem 5

For each $\epsilon > 0$, there is $d = d(\epsilon)$ and N such that for all $n > N$, there is an $(n, d, \frac{1}{2} - \epsilon)$-edge expander.

Probabilistic argument.
Random graphs are expanders with high probability.

Characterisation

Theorem 6

1) *If G is (n, d, λ)-expander, then it is $(n, d, \frac{1-\lambda}{2})$-edge expander.*
2) *If G is (n, d, ρ)-edge expander, then*

$$\lambda(G) \leq 1 - \frac{\rho^2}{2}.$$

Furthermore, if G has all self loops, it is $(n, d, 1-\epsilon)$-expander, $\epsilon = \min\{\frac{2}{d}, \frac{\rho^2}{2}\}$.

Algebraic expander implies combinatorial edge expansion

Lemma 7

Let G be an (n, d, λ)-expander. $S \subset V$, $T = \overline{S}$. Then:

$$|E(S,T)| \geq (1-\lambda)\frac{d|S| \cdot |T|}{|S|+|T|}. \tag{6}$$

Define $x \in \mathbb{R}^n$ by

$$x_i = \begin{cases} |T|, \text{ if } i \in S, \\ -|S|, \text{otherwise.} \end{cases} \tag{7}$$

Proof - I

Then:

$$||x||_2^2 = |S| \cdot |T|^2 + |T| \cdot |S|^2 = |S| \cdot |T| \cdot (|S| + |T|)$$

$$x \perp 1,$$

since $\sum x_i = 0$.
Set

$$Z = \sum_{i,j} A_{ij}(x_i - x_j)^2.$$

If i, j are all in S or T, $x_i - x_j = 0$, and if i, j are in the cut, then $(x_i - x_j)^2 = (|S| + |T|)^2$.

Proof - II

Therefore,

$$Z = \frac{2}{d} \cdot |E(S,T)| \cdot (|S| + |T|)^2.$$

On the other hand,

$$\begin{aligned} Z &= \sum_{i,j} A_{ij}(x_i - x_j)^2 \\ &= \sum_{i,j} A_{ij}x_i^2 - 2\sum_{i,j} A_{ij}x_i x_j + \sum_{i,j} A_{ij}x_j^2 \\ &= 2||x||_2^2 - 2\langle x, Ax\rangle. \end{aligned} \tag{8}$$

Proof - III

Therefore

$$\frac{1}{d} \cdot |E(S,T)|(|S|+|T|)^2 = ||x||_2^2 - \langle x, Ax \rangle.$$

Since $x \perp 1$,
(i) $\|Ax\|_2 \leq \lambda \|x\|_2$
(ii) $\langle x, Ax \rangle \leq \|x\|_2 \cdot \|Ax\|_2$.
Finally,

$$|E(S,T)| \geq (1-\lambda)\frac{d|S| \cdot |T|}{|S|+|T|}.$$

Expander mixing lemma

Lemma 8

Let $G = (V, E)$ be an (n, d, λ)-expander. Let $X, Y \subseteq V$. Then:

$$\left| |E(X,Y)| - \frac{d}{n}|X| \cdot |Y| \right| \leq \lambda d \cdot \sqrt{|X| \cdot |Y|}. \tag{9}$$

Intuition: Expander $\approx$ Pseudorandom

Proof - I

Define

$$\psi_X(x) = \begin{cases} \sqrt{d}, & \text{if } x \in X, \text{ and} \\ 0, & \text{otherwise.} \end{cases} \tag{10}$$

Then:

$$\psi_X A \psi_Y^{\mathrm{T}} = e(X, Y) = |E(X, Y)|.$$

Let $\phi_1, \phi_2, \cdots, \phi_n$ be the orthonormal eigenvectors of A. Suppose that

$$\psi_X = a_1\phi_1 + a_2\phi_2 + \cdots + a_n\phi_n$$

$$\psi_Y = b_1\phi_1 + b_2\phi_2 + \cdots + b_n\phi_n.$$

Then $\phi_1 = \sqrt{n}1$, $a_1 = \frac{\sqrt{d}}{\sqrt{n}}|X|$, and $b_1 = \frac{\sqrt{d}}{\sqrt{n}}|Y|$.

Proof - II

$$
\begin{aligned}
|e(X,Y) - a_1 b_1| &= \left| \sum_{i=2}^{n} a_i b_i \lambda_i \right| \\
&\leq \lambda(G) |\sum_{i=2}^{n} a_i b_b| \\
&\leq \lambda(G) \sqrt{\sum_{i=2}^{n} a_i^2 \cdot \sum_{i=2}^{n} b_i^2}. \qquad (11)
\end{aligned}
$$

By definition,

$$
\sum_{i=1}^{n} a_i^2 = \|\psi_X\|_2^2 = d \cdot |X|.
$$

Proof - III

Giving

$$\sqrt{\sum_{i=1}^{n} a_i^2} = \sqrt{d \cdot |X|}.$$

Therefore,

$$\begin{aligned} \left| e(X,Y) - \frac{d \cdot |X| \cdot |Y|}{n} \right| &\leq \lambda(G)\sqrt{d|X|} \cdot \sqrt{d|Y|} \\ &= \lambda d \sqrt{|X| \cdot |Y|}. \end{aligned} \tag{12}$$

Combinatorial edge expansion implies algebraic expander

Let $G = (V, E)$ be n-vertex, d-regular such that for any $S \subset V$ of size $\leq \frac{n}{2}$, $e(S, \overline{S}) \geq \rho d|S|$.
We will show that $\lambda(G) \leq 1 - \frac{\rho^2}{2}$.
Let A be the matrix of G, and λ be the second largest absolute eigenvalue of A.
Then there exists a u such that
(i) $u \perp 1$
(ii) $Au = \lambda u$.

Proof - I

Let

$$v_i = \begin{cases} u_i, & \text{if } u_i > 0, \text{ and} \\ 0, & \text{otherwise.} \end{cases}$$

$$w_i = \begin{cases} u_i, & \text{if } u_i \leq 0, \text{ and} \\ 0, & \text{otherwise.} \end{cases}$$

Then $u = v + w$. Since $u \perp 1$, $v, w \neq 0$. Suppose WLOG that the number of i's such that $v_i \neq 0$ is at most $\frac{n}{2}$.
Set

$$Z = \sum_{i,j} A_{ij} \left| v_i^2 - v_j^2 \right| .$$

We will prove

(1) $Z \geq 2\rho \|v\|_2^2$.
(2) $Z \leq \sqrt{8(1-\lambda)} \|v\|_2^2$.

The result follows.

For (1)

Suppose $v_1 \geq v_2 \geq \cdots \geq v_n$. So $v_i = 0$ for $i > \frac{n}{2}$.
For $i < j$:

$$v_i^2 - v_j^2 = \sum_{k=i}^{j-1} (v_k^2 - v_{k+1}^2).$$

By the assumption of v_i's,

$$Z = \sum_{i,j} A_{ij} \left| v_i^2 - v_j^2 \right| = 2 \sum_{i<j} A_{ij} \sum_{k=i}^{j-1} (v_k^2 - v_{k+1}^2).$$

For fixed k, for every edge $i \sim j$ with $i \leq k < j$, the term $v_k^2 - v_{k+1}^2$ appears once.

Proof - I

Therefore,

$$\begin{aligned}
Z &= 2\sum_{i<j} A_{ij} \cdot e(\{1,2,\cdots,k\},\{k+1,\cdots,n\}) \cdot (v_k^2 - v_{k+1}^2) \\
&= \frac{2}{d}\sum_{k=1}^{n/2} e(\{1,2,\cdots,k\},\{k+1,\cdots,n\})(v_k^2 - v_{k+1}^2) \\
&\geq \frac{2}{d}\sum_{k=1}^{n/2} \rho \cdot d \cdot k \cdot (v_k^2 - v_{k+1}^2) \\
&= 2\rho\sum_{k=1}^{n/2} k(v_k^2 - v_{k+1}^2) = 2\rho\|v\|_2^2. \qquad (13)
\end{aligned}$$

For (2)

$$Z \leq \sqrt{8(1-\lambda)} \cdot \|v\|_2^2$$

Recall $Au = \lambda u$, $u = v + w$, $u \perp 1$, $v \perp w$.

$$\langle Av, v\rangle + \langle Aw, v\rangle = \langle Au, v\rangle = \langle \lambda(v+w), v\rangle = \lambda \|v\|_2^2$$

Since $\langle Aw, v\rangle < 0$, $\frac{\langle Av, v\rangle}{\langle v, v\rangle} \geq \lambda$.

Therefore

$$\begin{aligned} 1 - \lambda &\geq 1 - \frac{\langle Av, v\rangle}{\|v\|_2^2} \\ &= \frac{\|v\|_2^2 - \langle Av, v\rangle}{\|v\|_2^2} \\ &= \frac{\sum_{i,j} A_{ij}(v_i - v_j)^2}{2\|v\|_2^2}. \end{aligned} \tag{14}$$

Proof - II

Note $Z = \sum_{i,j} A_{ij}(v_i - v_j)^2 = 2\|v\|_2^2 - 2\langle Av, v\rangle$.
Cauch-Schwarz: $\langle x, y\rangle \leq \|x\|_2 \cdot \|y\|_2$.
Let $x_{ij} = \sqrt{A_{ij}}(v_i - v_j)$, $y_{ij} = \sqrt{A_{ij}}(v_i - v_j)$.

$$\begin{aligned} &(\sum_{i,j} A_{ij}(v_i - v_j)^2) \cdot (\sum_{i,j} A_{ij}(v_i + v_j)^2) \\ &\geq (\sum_{i,j} A_{ij}(v_i - v_j) \cdot (v_i + v_j))^2 = Z^2. \end{aligned} \tag{15}$$

Proof - III

$$\begin{aligned} & 2\|v\|_2^2 \sum_{i,j} A_{ij}(v_i + v_j)^2 \\ = \ & 2\|v\|_2^2 \left(\sum A_{ij} v_i^2 + 2 \sum A_{ij} v_i v_j + \sum A_{ij} v_j^2\right) \\ = \ & 2\|v\|_2^2 (2\|v\|_2^2 + 2\langle Av, v\rangle). \end{aligned} \tag{16}$$

Noting
(i) $\|Av\|_2 \leq \|v\|_2$
(ii) $\langle Av, v\rangle \leq \|Av\|_2 \cdot \|v|_2 \leq \|v\|_2^2$.
Finally, $1 - \lambda \geq \frac{Z^2}{8\|v\|_2^4}$, giving

$$Z \leq \sqrt{8(1-\lambda)} \cdot \|v\|_2^2.$$

Research project

Information theoretical characterisation of expanders?

Spectral norm

For every matrix A, define the *spectral norm* of A, written $\|A\|$, as follows:

$$\begin{aligned} \|A\| &= \max\{\|Av\|_2 : \|v\|_2 = 1\} \\ &= \max\{\frac{\|Av\|_2}{\|v\|_2}\}. \end{aligned} \tag{17}$$

Proposition For any matrices A, B,
(1) $\|A+B\| \leq \|A\| + \|B\|$, and
(2) $\|AB\| \leq \|A\| \cdot \|B\|$.

Extracting randomness from expander

Theorem 9

Let A be the adjacency matrix of an (n, d, λ)-expander graph G. Let J be the $n \times n$ matrix such that $J_{ij} = \frac{1}{n}$ for all i, j. Then

$$A = (1 - \lambda)J + \lambda C$$

for some C with $\|C\| \leq 1$.

Intuition A uniformly random distribution can be extracted from an expander. If λ is small, then G is largely a random graph.

Proof - I

Solving C, we have

$$C = \frac{1}{\lambda}(A - (1-\lambda)J).$$

We prove $\|C\| \le 1$, that is, for every v, $\|Cv\|_2 \le \|v\|_2$.
Fix v.
Set
$v = u + w$, $u = \alpha 1$, $w \perp 1$.
We have
(1) $Cu = u$, easy
(2) For $w' = Aw$,

$$Cw = \frac{1}{\lambda} w'$$

Because: $w \perp 1$, so $\sum w_i = 0$, and hence $Jw = 0$.

Proof - II

(3) $Cv = C(u+w) = u + \frac{1}{\lambda}w'$
(4)

$$\begin{aligned} \|Cv\|_2^2 &= \|u\|_2^2 + \|\frac{1}{\lambda}w'\|_2^2 \\ &= \|u\|_2^2 + \frac{1}{\lambda^2}\cdot\|Aw\|_2^2 \\ &\leq \|u\|_2^2 + \frac{1}{\lambda^2}\cdot\lambda^2\cdot\|w\|_w\|_2^2 \\ &= \|v\|_2^2. \end{aligned} \tag{18}$$

Intuition of expanders

- Expander is basically a random graph
- The nice properties of expander graphs can be achieved simply by randomness
- Randomness plays an essential role for expanders:
- Information quickly spreads in expander graphs
- Viruses quickly infect the whole expander graphs
  (Here there is a dilemma to achieve both security and quick spreading of information in communication networks. Expanders may not be the best model for communication networks.

Expander walk theorem

Theorem 10

Let G be an (n, d, λ)-expander graph. Let $\mathcal{B}$ be a set of $[n]$ of size $\leq \beta n$, $0 < \beta < 1$. Let $X_1, X_2, \cdots, X_k$ be a random walk in G from X_1, where X_1 is randomly and uniformly chosen. Then:

$$\Pr[(\forall i \in [k])[X_i \in \mathcal{B}]] \leq ((1 - \lambda)\sqrt{\beta} + \lambda)^{k-1}.$$

Proof - I

For each i, $1 \leq i \leq k$, let
B_i: the event $X_i \in \mathcal{B}$.
Then:

$$\begin{aligned} & \Pr[\wedge_{i=1}^{k} B_i] \\ = \quad & \Pr[B_1] \cdot \Pr[B_2 | B_1] \cdot \cdots \cdot \Pr[B_k \,| B_1, \cdots, B_{k-1}]. \end{aligned} \tag{19}$$

Define B to be a linear transformation from $\mathbb{R}^n$ to $\mathbb{R}^n$ that keeps the values indexed in $\mathcal{B}$.
That is, for $(u_1, u_2, \cdots, u_n)$, define

$$(Bu)_i = \begin{cases} u_i, & \text{if } i \in \mathcal{B}, \text{ and} \\ 0, & \text{otherwise.} \end{cases} \tag{20}$$

Proof - II

For every probability vector p,
(i) Bp is the vector whose coordinates sum to the probability that a vertex i is chosen according to p, is in $\mathcal{B}$.
(ii) The normalised Bp is the distribution of p conditioned to the event that the vertex is in $\mathcal{B}$.

Proof - III

Let p^i be the distribution of X_i conditioned on the events $B_1, \cdots, B_i$. Then:

$$p^1 = \frac{1}{\Pr[B_1]} \cdot B1$$

$$p^2 = \frac{1}{\Pr[B_2|B_1]\Pr[B_1]} BAB1$$

$$p^i = \frac{1}{\Pr[B_i \, |B_{i-1}, \cdots B_1] \cdots \Pr[B_1]} (BA)^{i-1} B1.$$

Hence,

$$\Pr[B_1] \cdots \Pr[B_k|B_{k-1} \cdots B_1] p^k = (BA)^{k-1} B1.$$

Proof - IV

$$\Pr[\wedge_{i=1}^{k} B_i] = \Pr[B_1] \cdots \Pr[B_k | B_{k-1} \cdots B_1] = |(BA)^{k-1} B\mathbf{1}|_1.$$

Let $A = (1-\lambda)J + \lambda C$.
Then $BA = (1-\lambda)BJ + \lambda BC$.
Noting:
(i) $\|B\mathbf{1}\|_2 \leq \sqrt{\beta}\|\mathbf{1}\|_2$
(ii) $\|BJ\| \leq \sqrt{\beta}$, $\|B\| \leq 1$, $\|BC\| \leq 1$.
(iii) $\|BA\| \leq (1-\lambda)\sqrt{\beta} + \lambda$
Therefore,

$$\begin{aligned} |(BA)^{k-1} B\mathbf{1}|_1 &\leq \|(BA)^{k-1} B\mathbf{1}\|_2 \cdot \sqrt{n} \\ &\leq ((1-\lambda)\sqrt{\beta} + \lambda)^{k-1}. \end{aligned} \tag{21}$$

Research projects

- Information theoretical approach to expanders?
- There is no reason that λ_1 is the only important eigenvale What roles do the other eigenvalues play in the combinatorial properties of graph?

Rotation map

Given a d-regular graph G,

$$\widehat{G}: \ [n] \times [d] \to [n] \times [d]$$

$\widehat{G}(u, i) = (v, j)$ means:
(i) v is the i-th neighbor of u, and
(ii) u is the j-th neighbor of v.
$\widehat{G}$ is log space computed.

The matrix product

GG' corresponds to AA'

$$\lambda(GG') \leq \lambda(G) \cdot \lambda(G').$$

The tensor product

Graphs G, G'
matrices A, A'
$G \otimes G'$
$A \otimes A'$.

$$\lambda(G \otimes G') \leq \max\{\lambda(G), \lambda(G')\}.$$

Replacement product

Given:
(i) G: n vertices, degree D
(ii) G': D vertices, degree d.
Define the replacement product:

$$A \circ_R A' = \frac{1}{2}\widehat{A} + \frac{1}{2}(I_n \otimes A')$$

$\widehat{A}$ is the matrix of the rotation map of G.

Lemma 11

If $\lambda(G) \leq 1 - \epsilon$, $\lambda(H) \leq 1 - \delta$, then

$$\lambda(G \circ_R H) \leq 1 - \frac{\epsilon\delta^2}{24}.$$

The construction

1. Let H be a $(D = (2d)^{100}, d, 0.01)$-expander, d constant.
2. Let G_1 be a $((2d)^{100}, 2d, \frac{1}{2})$-expander
   G_2 be a $((2d)^{200}, 2d, \frac{1}{2})$-expander.
3. For $k > 2$,

$$G_k = (G_{\lfloor \frac{k-1}{2} \rfloor} \otimes G_{\lceil \frac{k-1}{2} \rceil})^{50} \circ_R H.$$

Theorem 12

G_k is $((2d)^{100k}, 2d, 1 - \frac{1}{50})$-expander graph.

UPATH is in RL

UPATH: Given an undirected graph G, for given s, t, decide whether or not there is a path from s to t.
Assume G is regular and has self-loop at every vertex.
By the previous theorems, for $l = n^4$, with probability $\geq \frac{2}{3}$, a random walk of length l hits t, if there is a path from s to t.
So
UPATH is in RL, randomised log space.

Connectivity of expander

For regular graphs with self-loop at each vertex, we have:
1) If G is connected and $\lambda(G) < 1$, then the diameter of G is $O(\log n)$.
2) If there is a constant $\lambda < 1$ such that for every connected component H of G, $\lambda(H) \leq \lambda$, then for every H, the diameter of H is $O(\log n)$.
For a graph with property 2), there is a deterministic log space algorithm to decide for given s, t, whether or not there is a path from s to t.

UPATH $\in$ L

Reduction: for a regular graph G,

1) Let G_0 be obtained from G by adding self-loops such that G_0 has degree d^{50} for some constant d.
2) Let H be a $(d^{50}, \frac{d}{2}, 0.01)$-expander.
3) Gor $k \geq 1$,

$$G_k = (G_{k-1} \circ_R H)^{50}.$$

Proof

Lemma 13

For every $k \geq 0$, every connected component in G_k is an $(d^{50k}n, d^{20}, 1-\epsilon)$-expander, where $\epsilon = \min\{\frac{1}{20}, \frac{1.5^k}{12n^2}\}$, there n is the number of vertices in G.

For $k = 10\log n$, ϵ is constant.
G_k is computed from G by log space, and the connectivity in G_k is decided in log space.

Conclusions and discussion

1. expander $\approx$ random graph
2. expander can be used to de-randomize
3. information quickly spreads in expander
4. explicit construction of expanders can be used in new algorithms

Open questions

- Resolving the dilemma of expander walk and security of networks
- What are the optimum communication networks?

谢谢！