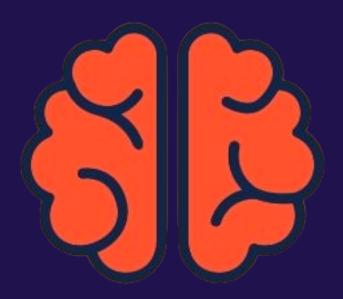
# Κρυπτογραφία- Απαλλακτική Εργασία



# **CRYPTOHACK**

Κοντοπούλου Δέσποινα

Πανεπιστήμιο Πειραιώς- Τμήμα Πληροφορικής 15/01/2024 Σύνδεσμος profile στο CryptoHack:
https://cryptohack.org/user/thespianexe
Σύνδεσμος repository στο Github:
https://github.com/Thespiann/CryptoHack
Doc

# Σύνολο πόντων εώς 15/1/2024:



- RSA Starter 6 (25 pts)
- Diffie-Hellman Starter 5 (40 pts)
- Modes of Operation Starter (15 pts)
  960pts μετρήσιμοι στην εργασία

# CHALLENGES INTRODUCTION

# Finding Flags

Εισαγωγή flag σε form.

crypto{y0ur\_f1rst\_f14g}

#### Great Snakes

Εκτέλεση συννημένου.

crypto{z3n Of pythOn}

#### Network Attacks

For this challenge, connect to socket.cryptohack.org on port 11112. Send a JSON object with the key buy and value flag. Επίλυση

```
from pwn import *
import json

HOST = "socket.cryptohack.org"
PORT = 11112

r = remote(HOST, PORT)
```

```
def json_recv():
    line = r.readline()
    return json.loads(line.decode())

def json_send(hsh):
    request = json.dumps(hsh).encode()
    r.sendline(request)

print(r.readline())
print(r.readline())
print(r.readline())
print(r.readline())

request = {
    "buy": "flag"
}
json_send(request)
response = json_recv()
print(response)
```

#### Εξήγηση επίλυσης

- 1. Χρήση βιβλιοθήκης pwn για σύνδεση στον διακομιστή.
- 2. Αποστολή αντικειμένου json με json\_send()
- 3. Ανάγνωση απαντησης με json\_recv()
- 4. Εκτύπωση απαντήσεων

```
crypto{sh0pp1ng f0r f14g5}
```

# **GENERAL**

#### **ENCODING**

#### **ASCII**

Using the below integer array, convert the numbers to their corresponding ASCII characters to obtain a flag.

#### Επίλυση

```
flag = "".join([chr(i) for i in [99, 114, 121, 112, 116,
111, 123, 65, 83, 67, 73, 73, 95, 112, 114, 49, 110, 116,
52, 98, 108, 51, 125]])
print(flag)
```

#### Εξήγηση επίλυσης

To flag είναι η ένωση (join) των αντίστοιχων ASCII χαρακτήρων των ακεραίων(chr(int)).

crypto{ASCII\_pr1nt4bl3}

#### HEX

Included below is a flag encoded as a hex string. Decode this back into bytes to get the flag. Επίλυση

```
hex='63727970746f7b596f755f77696c6c5f62655f776f726b696e67
5f776974685f6865785f737472696e67735f615f6c6f747d'
print(bytes.fromhex(hex))
```

#### BASE 64

Take the below hex string, *decode* it into bytes and then *encode* it into Base64.

Επίλυση

```
import base64
hex='72bca9b68fc16ac7beeb8f849dca1d8a783e8acf9679bf9269f7
bf'
bytes=bytes.fromhex(hex)
print(base64.b64encode(bytes).decode())
```

crypto/Base+64+Encoding+is+Web+Safe/

#### BYTES AND BIG INTEGERS

Convert the following integer back into a message: Επίλυση

```
from Crypto.Util.number import bytes_to_long,
long_to_bytes
int_value =
115151950638623188999316854888137473957755162872896826364
99965282714637259206269 #given int
byte_data =long_to_bytes(int_value) #convert int to bytes
ascii_bytes=bytes(byte_data) #convert bytes to ascii
print(ascii_bytes.decode('utf-8')) #print ascii
```

crypto{3nc0d1n6 411 7h3 w4y d0wn}

#### ENCODING CHALLENGE

Now you've got the hang of the various encodings you'll be encountering, let's have a look at automating it. Can you pass all 100 levels to get the flag? Connect at socket.cryptohack.org 13377

#### Επίλυση

```
import base64
import codecs
import json
from pwn import *
import Crypto.Util.number
HOST = "socket.cryptohack.org"
PORT = 13377
r= remote(HOST, PORT)
#functions for reading lines and handling json objects
def readline():
   return r.readuntil(b"\n")
def json recv():
   line = readline()
   return json.loads(line.decode())
def json_send(hsh):
   request = json.dumps(hsh).encode()
   r.write(request)
for i in range(100):
   received = json recv()#receive data
   encoding = received['type']
   encoded = received['encoded']
#decode using the specified encoding
   if encoding == "base64":
       decoded = base64.b64decode(encoded).decode()
   elif encoding == "hex":
       decoded = bytes.fromhex(encoded).decode()
   elif encoding == "rot13":
       decoded = codecs.decode(encoded, "rot13")
   elif encoding == "bigint":
       decoded =
Crypto.Util.number.long_to_bytes(int(encoded,
16)).decode()
   elif encoding == "utf-8":
       decoded = "".join(chr(o) for o in encoded)
#prints
   print(f"{i + 1}) {encoding}:")
   print(f"{encoded} ---> {decoded}\n")
#send decoded back to get the flag
```

#### Εξήγηση Επίλυσης

Το πρόγραμμα συνδέεται στον διακομιστή και την θύρα που δίνονται και έπειτα λαμβάνει 100 κρυπτογραφημένα json objects οπου το κάθε ένα εχει τα πεδία type και encoded. Αποκωδικοποιεί τα μηνύματα, στέλνει πίσω το decoded και έπειτα διαβάζει και εκτυπώνει το flag.

crypto{3nc0d3\_d3c0d3\_3nc0d3}

#### **XOR**

#### XOR STARTER

Given the string label, XOR each character with the integer 13. Convert these integers back to a string and submit the flag as crypto {new\_string}. Επίλυση

```
xor= [chr(ord(char)^13) for char in "label"]
new=''.join(xor)
print("crypto{"+new+"}")
```

crypto{aloha}

#### XOR PROPERTIES

Below is a series of outputs where three random keys have been XOR'd together and with the flag. Use the above properties to undo the encryption in the final line to obtain the flag.

#### Επίλυση

```
k1 =
bytes.fromhex('a6c8b6733c9b22de7bc0253266a3867df55acd
e8635e19c73313')
k2k1 =
bytes.fromhex('37dcb292030faa90d07eec17e3b1c6d8daf94c
35d4c9191a5e1e')
k2k3 =
bytes.fromhex('c1545756687e7573db23aa1c3452a098b71a7f
bf0fddddde5fc1')
flagk1k2k3=
bytes.fromhex('04ee9855208a2cd59091d04767ae47963170d1
660df7f56f5faf')
def xor(bytes1,bytes2):
   return bytes([b1 ^ b2 for b1, b2 in
zip(bytes1,bytes2)])
k2 = xor(k2k1, k1)
k3=xor(k2k3,k2)
flag=xor(xor(flagk1k2k3,k1),xor(k2,k3))
print(flag)
```

#### FAVOURITE BYTE

I've hidden some data using XOR with a single byte, but that byte is a secret. Don't forget to decode from hex first.

#### Επίλυση

```
from pwn import xor
hex='73626960647f6b206821204f21254f7d694f7624662065622127
234f726927756d'
bytes=bytes.fromhex(hex)
for i in range(255):
    flag= xor(bytes,i)
    if flag.startswith(b'crypto'):
        print(xor(bytes,i))
        print("secret key is: " + str(i))
        break
```

## Εξήγηση επίλυσης

Το πρόγραμμα δοκιμάζει τις δυνατές τιμές ενός byte i ως κλειδί στην χοι με το δοθέν hex. Εαν το αποτέλεσμα ξεκινά άπο crypto, εκτυπώνει το αποκρυπτογραφημένο κείμενο και το κλειδι.

```
crypto{0x10 15 my f4v0ur173 by7e}
```

# YOU EITHER KNOW, XOR YOU DON'T

I've encrypted the flag with my secret key, you'll never be able to guess it.

#### Επίλυση

```
from pwn import xor
hex='0e0b213f26041e480b26217f27342e175d0e070a3c5b103e
2526217f27342e175d0e077e263451150104'
bytes=bytes.fromhex(hex)
crypto='crypto{'.encode()
#print(xor(bytes,crypto)) :
#b'myXORke+y_Q\x0bHOMe$~seG8bGURN\x04DFWg)a|\x1dTM!an
#\x7f'
key='myXORkey'.encode()
print(xor(bytes,key))
```

#### Εξήγηση επίλυσης

Κάνοντας xor με το format των flag, προκύπτει το

```
b'myXORke+y_Q\\x0bHOMe$\sim seG8bGURN\\x04DFWg)a|\\x1dTM!an\\x7f'
```

αρα μαντέυουμε πως το κλειδί ειναι myxorkey.

#### LEMUR XOR

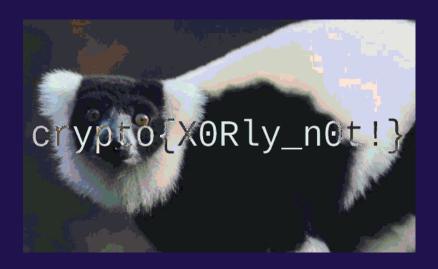
I've hidden two cool images by XOR with the same secret key so you can't see them!

#### Επίλυση

```
from PIL import Image
import numpy as np
def image_to_array(image_path):#for getting each pictures pixels
   img = Image.open(image_path)
   return np.array(img)
def array to image(array, output path):#for turning the pixels into images
   img = Image.fromarray(array)
   img.save(output_path)
flag_path = "flag.png"
lemur path = "lemur.png"
flag_array = image_to_array(flag_path)
lemur_array = image_to_array(lemur_path)
new_array = np.bitwise_xor(flag_array, lemur_array) #xor-ing the pixels of
the images
output_path = "new_flag.png" #save in new file
array_to_image(new_array, output_path)
```

#### Εξήγηση Επίλυσης

Το πρόγραμμα μετατρέπει τις φωτογραφίες σε arrays με pixel με την συνάρτηση image\_to\_array(). Έπειτα εφαρμόζεται xor στα arrays και μετατρέπεται το array που προέκυψε σε .jpg με την array\_to\_image(). new\_flag.jpg:



#### **MATHEMATICS**

#### GREATEST COMMON DIVISOR

Now calculate gcd(a,b) for a = 66528, b = 52920

#### Επίλυση

```
def gcd(a,b):#function for getting the greatest common
divisor
   if a==0:
      return b #if a is 0 then b is the gcd
   else:
      return gcd(b%a,a) #recursion, we call gcd() with the
updated values
print(gcd(66528,52920))
```

1512

#### EXTENDED GCD

Using the two primes p = 26513, q = 32321, find the integers u, v such that p \* u + q \* v = gcd(p,q). Enter whichever of u and v is the lower number as the flag.

#### Επίλυση

```
p = 26513
q = 32321
#p * u + q * v = gcd(p,q)
def egcd(a,b):
    if a==0:
        return b,0,1
    else:
        gcd, u, v= egcd(b%a,a)
        return gcd, v-(b//a)*u, u
gcd,u,v=egcd(p,q)
print(min(u,v))
```

#### Εξήγηση Επίλυσης:

Υλοποιειται ο επαυξημένος αλγόριθμος του Ευκλείδη για τον υπολογισμό του gcd. Συγκεκριμένα:

- αν α=0 -> μεγιστος κοινος διαιρέτης b
- av a!= 0-> egcd(b mod a, a) -return gcd, u, v -(b//a)\*u, u

εκτυπώνω το ελάχιστο απο τα u,v

#### MODULAR ARITHMETIC 1

Calculate the following integers:

 $11 \equiv x \mod 6$  $8146798528947 \equiv y \mod 17$ 

Επίλυση

```
print(min(11%6,8146798528947%17))
```

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#### MODULAR ARITHMETIC 2

Now take the prime p = 65537. Calculate 273246787654<sup>65536</sup> mod 65537.

#### Επίλυση

```
p=65537
a=273246787654
#fermats little theorem
#if p=prime, a^p mod p= a
#if a%p!=0 (a not divisble by p), then a^(p-1)-1 =
x*p, so a^(p-1)=1 mod p
if a%p !=0:
    print(1)
```

#### MODULAR INVERTING

What is the inverse element:  $3 * d \equiv 1 \mod 13$ ?  $E\pi(\lambda u \sigma \eta)$ 

```
#For all elements g in the field, there exists a unique integer d such that g * d \equiv 1 mod p d=0 while (3*d)%13 !=1: d=d+1 print(d)
```

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#### DATA FORMATS

#### PRIVACY-ENHANCED MAIL

Extract the private key *d* as a decimal integer from this PEM-formatted RSA key.

#### Επίλυση

```
from Crypto.PublicKey import RSA
with open('privacy_enhanced_mail.pem', 'r') as key_file:
    key = RSA.importKey(key_file.read())
    private_key_d_decimal = key.d
    print(private_key_d_decimal)
```

 $15682700288056331364787171045819973654991149949197959929860861\\ 22818002170731685192445620554366556581089267419005983133023143\\ 69709144747745627149456205191443897851589089941819513488460174\\ 32506464163564960993784254153395406799101314760033445065193429\\ 59251234995202098293221852446234100210206343548931881331646451\\ 16217369439384407104706949123362376802197462045951289591618005\\ 95216366237538296447335375818871952520026993102148328897083547\\ 18428649324119150595360166885894112979096690923694112785137020\\ 24211358970910867635698847600991122910720569706363804173490195\\ 79768748054760104838790424708988260443926906673795975104689$ 

#### CERTAINLY NOT

Find the modulus of the certificate, giving your answer as a decimal.

#### Επίλυση

```
from Crypto.PublicKey import RSA
from cryptography import x509
from cryptography.hazmat.primitives import serialization
from cryptography.hazmat.backends import default_backend

with open('2048b-rsa-example-cert.der', 'rb') as der_file:
    cert = x509.load_der_x509_certificate(der_file.read(),
    default_backend())

with open('2048b-rsa-example-cert.pem', 'wb') as pem_file:

pem_file.write(cert.public_bytes(encoding=serialization.Encoding.PEM))
with open('2048b-rsa-example-cert.pem', 'r') as key_file:
    key = RSA.import_key(key_file.read())
    print(type(key))
    print(key.n)
```

#### Εξήγηση Επίλυσης

Το πρόγραμμα χρησιμοποιει το cryptography module για την φόρτωση του der πιστοποιητικού cert. Επειτα το μετατρέπει σε μορφή PEM με την χρήση της

```
cert.public_bytes(encoding=serialization.Encoding.PEM)
```

και χρησιμοποιεί την RSA για να φορτώσει το κλειδι σε δεκαδικό.

228253736920195308043062128646095127753741718239937085165098976315
475136346358563756240037370680345490476779993109418374543788293513
983023826296582640787754568386262075077254940306005168728523061912
554929264959655363792718753104573191079360207300504762352786715282
658175714339195611756650961711897584061364539879662552369637826660
669626546784649500759230603273586913566329086064982317559635673823
390109852226232055869234664058092174266703334100144299051469416522
933662129037336300830163988108873560199774094673747422662762671375
470215768742048095060459149644910633938004991674164719490219954477
22415959979785959569497

#### SSH KEYS

Extract the modulus *n* as a decimal integer from Bruce's SSH public key.

#### Επίλυση

```
from Crypto.PublicKey import RSA
pub = open("bruce_rsa.pub").read()
key = RSA.importKey(pub)
print(key.n)
```

#### TRANSPARENCY

Find the subdomain of cryptohack.org which uses these parameters in its TLS certificate, and visit that subdomain to obtain the flag.

#### Επίλυση

```
import requests, json, sys
target = 'cryptohack.org'
req =
requests.get("https://crt.sh/?q=%.{d}&output=json".fo
rmat(d=target))
json_data = json.loads(req.text)
for (index, certificate) in enumerate(json_data):
    url = str(certificate['name_value'])
```

#### Επεξήγηση επίλυσης

Με την βοήθεια του ιστοτόπου crt.sh γίνεται αναζήτηση για πιστοποιητικά που να περιέχουν 'cryptohack.org' με ένα http request.

Περιηγούμαστε στην λίστα απο τα πιστοποιητικά που επιστράφηκε με το json και για κάθε ένα:

- 1. Ελέγχουμε αν το url λήγει σε 'cryptohack.org'
- 2. Αν ναι, κάνουμε ξανα http request στον ιστοτοπο.
- 3. Αν η απάντηση ξεκινά απο 'crypto' εκτυπώνω το flag και γίνεται break.

```
crypto{thx redpwn for inspiration}
```

# **MATHEMATICS**

#### **MODULAR MATH**

## QUADRATIC RESIDUES

Find the quadratic residue and then calculate its square root. Of the two possible roots, submit the smaller one as the flag.

#### Επίλυση

```
p = 29
ints = [14, 6, 11]
#for i in range(29):
# print(pow(i,2,29))
# if pow(i,2,29)==18:
# print("i="+str(i))

for i in ints:
  for a in range(29):
    if pow(a,2,29)==i:
        print("a="+str(a))
```

#### LEGENDRE SYMBOL

Given the following 1024 bit prime and 10 integers, find the quadratic residue and then calculate its square root; the square root is your flag. Of the two possible roots, submit the larger one as your answer. Επίλυση

```
p = 1015...
ints = [25081841...]
```

```
def legendre(a, p):
    legendre=pow(a,(p-1)//2,p)
    if(legendre==p-1):
        legendre=-1
    return legendre

for i in ints:
    l= legendre(i,p)
    if l==1:
        sqrt= pow(i,(p+1)//4,p)
        print(sqrt)
```

#### Εξήγηση επίλυσης.

Αρχικά υλοποιείται συνάρτηση για τον υπολογισμό του Legendre, ωστε να βρεθεί ποιος απο τους αριθμούς είναι quadratic residue. Ελέγχονται οι ακέραιοι, και στην περίπτωση που βρεθει το quadratic residue υπολογίζουμε την τετραγωνική του ρίζα, υψώνοντας τον ακέραιο εις την (p+1)//4 με mod p.

#### MODULAR SQUARE ROOT

Find the square root of a modulo the 2048-bit prime **p**. Give the smaller of the two roots as your answer. Επίλυση (σε περιβάλλον sageMath)

```
from sage.rings.finite_rings.integer_mod import
square_root_mod_prime
a=8479994...
p=305318...
square_root_mod_prime(mod(a, p))
```

 $236233930768304863832777329858048929893213750552050038833827105205373474\\786235177964731417681795335907187156004112528991924714607490715161276264\\086819962118655952206833803260099131188222401602122267224313936218046123\\264673246584884042545825793088785658337960096776173859678287785131848935\\567982281315512304570528511209944814642675511016000251559241885043210364\\181581107154845628426350780558944507365756538185052136796967569976075531\\078462357707644003774768176030243492493211364006173877760119462224419275\\802418085391624442725406544196255728257284916277274079898964794864520734\\9737457445440405057156897508368531939120$ 

# SYMMETRIC CIPHERS

#### **HOW AES WORKS**

#### KEYED PERMUTATIONS

What is the mathematical term for a one-to-one correspondence?

crypto{bijection}

#### RESISTING BRUTEFORCE

What is the name for the best single-key attack against AES?

crypto{biclique}

#### STRUCTURE OF AES

Included is a bytes2matrix function for converting our initial plaintext block into a state matrix. Write a matrix2bytes function to turn that matrix back into bytes, and submit the resulting plaintext as the flag.

Επίλυση

```
def bytes2matrix(text):
   """ Converts a 16-byte array into a 4x4 matrix.
   return [list(text[i:i + 4]) for i in range(0, len(text),
4)7
def matrix2bytes(matrix):
   """ Converts a 4x4 matrix into a 16-byte array.
   bytes array = bytearray()
   for i in matrix:
       for j in i:
           bytes array.append(j)
   return bytes_array
matrix = \Gamma
   [99, 114, 121, 112],
   \lceil 116, 111, 123, 105 \rceil
   [110, 109, 97, 116],
   [114, 105, 120, 125],
print(matrix2bytes(matrix))
```

crypto{inmatrix}

#### ROUND KEYS

Complete the add\_round\_key function, then use the matrix2bytes function to get your next flag. Επίλυση

```
from StructureOfAES import matrix2bytes
state = [
   [206, 243, 61, 34],
   [171, 11, 93, 31],
   [16, 200, 91, 108],
   [150, 3, 194, 51],
round key = [
   [173, 129, 68, 82],
   [223, 100, 38, 109],
   [32, 189, 53, 8],
   [253, 48, 187, 78],
def add round key(s, k):
   added_round_key = [
       [scolumn ^ kcolumn for scolumn, kcolumn in zip(scol, kcol)]
       for scol, kcol in zip(s, k)
   return added round key
print("The added round key is: \n ", add round key(state,
round_key))
print("The byte array is: \n ", matrix2bytes(add_round_key(state,
round key)))
```

crypto{r0undk3y}

#### CONFUSION THROUGH SUBSTITUTION

Implement sub\_bytes, send the state matrix through the inverse S-box and then convert it to bytes to get the flag.

#### Επίλυση

```
import base64
from StructureOfAES import matrix2bytes
from RoundKeys import add_round_key
s box = (
   0x63,...)
inv_s_box = (
  0x52,...
state = [
   [251, 64, 182, 81],
   [146, 168, 33, 80],
   [199, 159, 195, 24],
   [64, 80, 182, 255],
def sub_bytes(s, sbox=s_box):
   sub_bytes =""
   for i in range(4):
       for j in range(4):
           sub_bytes+=chr(sbox[s[i][j]])
   return sub bytes
print(sub_bytes(state, sbox=inv_s_box))
```

#### Εξήγηση επίλυσης

Ορίζεται η κενή αρχικα συμβολοσειρα sub bytes. Επειτα υλοποιειται η διαδικασία της αντικαταστασης καθε byte του πίνακα s με το sbox και ο χαρακτήρας που προκύπτει προστίθεται στην συμβολοσειρά.

crypto{l1n34rly}

#### DIFFUSION THROUGH PERMUTATION

We've provided code to perform MixColumns and the forward ShiftRows operation. After implementing inv\_shift\_rows, take the state, run inv\_mix\_columns on it, then inv\_shift\_rows, convert to bytes and you will have your flag. Επίλυση

```
def shift_rows(s):
   s[0][1], s[1][1], s[2][1], s[3][1] = s[1][1], s[2][1], s[3][1],
s[0][1]
   s[0][2], s[1][2], s[2][2], s[3][2] = s[2][2], s[3][2], s[0][2],
s[1][2]
   s[0][3], s[1][3], s[2][3], s[3][3] = s[3][3], s[0][3], s[1][3],
s[2][3]
def inv shift rows(s):
   s[1][1], s[2][1], s[3][1], s[0][1] = s[0][1], s[1][1], s[2][1],
s[3][1]
   s[2][2], s[3][2], s[0][2], s[1][2] = s[0][2], s[1][2], s[2][2],
s[3][2]
   s[3][3], s[0][3], s[1][3], s[2][3] = s[0][3], s[1][3], s[2][3],
s[3][3]
# learned from http://cs.ucsb.edu/~koc/cs178/projects/JT/aes.c
xtime = lambda a: (((a << 1) ^ 0x1B) & 0xFF) if (a & 0x80) else (a + 1) ^ 0x1B
<< 1)
def mix_single_column(a):
   # see Sec 4.1.2 in The Design of Rijndael
  t = a[0] ^ a[1] ^ a[2] ^ a[3]
   u = a[0]
   a[0] ^= t ^ xtime(a[0] ^ a[1])
   a[1] ^= t ^ xtime(a[1] ^ a[2])
   a[2] ^= t ^ xtime(a[2] ^ a[3])
   a[3] ^= t ^ xtime(a[3] ^ u)
def mix columns(s):
   for i in range(4):
       mix single_column(s[i])
def inv mix columns(s):
   # see Sec 4.1.3 in The Design of Rijndael
```

```
for i in range(4):
       u = xtime(xtime(s[i][0] ^ s[i][2]))
       v = xtime(xtime(s[i][1] ^ s[i][3]))
       s[i][0] ^= u
       s[i][1] ^= v
       s[i][2] ^= u
       s[i][3] ^= v
   mix columns(s)
state = [
   [108, 106, 71, 86],
   [96, 62, 38, 72],
   [42, 184, 92, 209],
   [94, 79, 8, 54],
inv mix columns(state)
inv shift rows(state)
for i in range(4):
   for j in range(4):
       print(chr(state[i][j]), end="")
```

#### Εξήγηση Επίλυσης

Δημιουργείται η συνάρτηση inv\_shift\_rows(s) με βαση την δοθεν shift\_rows() ωστε να αντιστρέφονται οι αλλαγές που προκάλεσε η δεύτερη στο s. Έπειτα ο τελικός πίνακας εκτυπώνεται σε χαρακτήρες ASCII για την ευρεση του flag.

crypto{d1ffUs3R}

#### BRINGING IT ALL TOGETHER

Copy in all the building blocks you've coded so far, and complete the decrypt function that

implements the steps shown in the diagram. The decrypted plaintext is the flag.
Επίλυση

```
#Only showing decrypt func
def decrypt(key, ciphertext):
   round keys = expand key(
       key) # Remember to start from the last round key and
work backwards through them when decrypting
   # Convert ciphertext to state matrix
   text = bytes2matrix(ciphertext)
   # Initial add round key step
   add_round_key(text, round_keys[10])
   for i in range(N ROUNDS - 1, 0, -1):
       inv shift rows(text)
       inv sub bytes(text)
       add round key(text, round keys[i])
       inv mix columns(text)
   # Run final round (skips the InvMixColumns step)
   inv shift rows(text)
   inv sub bytes(text)
   add round key(text, round keys[0])
   # Convert state matrix to plaintext
   plaintext = matrix2bytes(text)
   return plaintext
print(decrypt(key, ciphertext))
```

crypto{MYAES128}

\*Υπάρχει η επίλυση για το *Modes of Operation Starter* στο github, όμως δεν προσμετρούνται οι πόντοι.

# **RSA**

#### **STARTER**

#### RSA STARTER 1

Find the solution to  $101^{17} \mod 22663$   $E\pi i \lambda u \sigma \eta$ 

```
print(pow(101,17,22663))
```

19906

#### RSA STARTER 2

"Encrypt" the number 12 using the exponent e = 65537 and the primes p = 17 and q = 23. What number do you get as the ciphertext? Ení $\lambda u \sigma \eta$ 

```
p=17
q=23
e=65537
N=p*q #calculation of N
print(pow(12,e,N))#modular exponantion with my
parameters
```

301

#### RSA STARTER 3

Given N = p\*q and two primes:

p = 857504083339712752489993810777

q = 1029224947942998075080348647219

What is the totient of N?

#### Επίλυση

```
p=857504083339712752489993810777 q=1029224947942998075080348647219 N=p*q """Euler's totient function is a multiplicative function, meaning that if two numbers and are relatively prime, then, \varphi(p*q)=\varphi(p)*\varphi(q). If p is a prime, then f(p)=p-1, because all the numbers up to p are not divisible by p""" totient=(p-1)*(q-1) print(totient)
```

882564595536224140639625987657529300394 956519977044270821168

#### RSA STARTER 4

Given the two primes:

p = 857504083339712752489993810777

q = 1029224947942998075080348647219

and the exponent:

e = 65537

What is the private key d? Επίλυση

```
p = 857504083339712752489993810777

q = 1029224947942998075080348647219

e = 65537

totient = (p-1)*(q-1)# (f(N) = f(p*q) = totient. if

gcd(e, f(N)) = 1, d = e^{-1} mod totient. In the context of

RSA, it is typical to choose an encryption exponent e

that is coprime with the totient, so d is:

print(pow(e, -1, totient))
```

#### 121832886702415731577073962957377780195 510499965398469843281

#### RSA STARTER 5

Use the private key that you found for these parameters in the previous challenge to decrypt this ciphertext.

### Επίλυση

```
n=
882564595536224140639625987659416029426239230804614613279163
#from previous challenge
fn=8825645955362241406396259876575293003949565199770442708211
68
e = 65537
key= pow(e,-1,fn)
c =
77578995801157823671636298847186723593814843845525223303932
#decrypted=c^key mod n
decrypted=pow(c,key,n)
print(decrypted)
```

13371337

\*Υπάρχει η επίλυση για το *RSA Starter 6* στο github, όμως δεν προσμετρούνται οι πόντοι.

#### PRIMES PART 1

#### **FACTORING**

Factorise the 150-bit number

510143758735509025530880200653196460532653

147 into its two constituent primes. Give the smaller one as your answer.

#### Επίλυση

```
#from sympy import factorint wont run for me
#i would do
#print(min(factorint(5101437587355090255308802006531964605326
5#3147))), instead
#i put the number in this calculator
#https://www.alpertron.com.ar/ECM.HTM
a = 19704762736204164635843
b = 25889363174021185185929
print(min(a,b))
```

19704762736204164635843

#### MONOPRIME

Why is everyone so obsessed with multiplying two primes for RSA. Why not just use one? Επίλυση

```
from Crypto.Util.number import long_to_bytes
n = 1717313...
e = 65537
ct = 1613675...
d=pow(e,-1,n-1)# the private key is calculated, the totient of prim N is
N-1
ciphertext=pow(ct,d,n)
print(long_to_bytes(ciphertext))
```

crypto{0n3\_pr1m3\_41n7\_pr1m3\_101}

#### **MANYPRIME**

Using one prime factor was definitely a bad idea so I'll try using over 30 instead.

#### Επίλυση

```
from Crypto.Util.number import inverse, long_to_bytes
import gmpy2
n = 58064...
e = 65537
ct = 3207214...
factors=[9282...]
#using http://factordb.com/index.php?query=58064239...
f=1
for x in factors:
    f *= (x-1)
#d=pow(e,-1,f-1)
d=gmpy2.invert(e,f)
print(long_to_bytes(pow(ct,d,n)))
```

crypto{700 m4ny 5m4ll f4c70r5}

#### PUBLIC EXPONENT

#### SALTY

Smallest exponent should be fastest, right? Επίλυση

```
from Crypto.Util.number import long_to_bytes
n = 110581...
e = 1
ct = 449812...
print(long_to_bytes(ct))# since d = 1 mod (p-1)(q-1) if
e=1
```

crypto{saltstack fell for this!}

# **DIFFIE-HELLMAN**

#### **STARTER**

#### DIFFIE-HELLMAN STARTER 1

Given the prime p = 991, and the element g = 209, find the inverse element d such that  $g * d \mod 991 = 1$ .

#### Επίλυση

```
p= 991
g=209
# g*d mod p =1
for x in range(p):
    res= (g*x)%p
    if res==1:
        print(x)
        break
```

569

#### DIFFIE-HELLMAN STARTER 2

For the finite field with p = 28151 find the smallest element g which is a primitive element of  $F_p$ .  $E\pi i \lambda u \sigma \eta$ 

```
def is_primitive_root(g, p):#Check if g is a
primitive root modulo p
  for i in range(2, p - 1):
     if pow(g, i, p) == 1:
        return False
  return True

def min_prim(p):# Find the smallest primitive element
```

```
in the finite field Fp
  for g in range(2, p):

    if is_primitive_root(g, p):
        return g

p = 28151
primitive_element = min_prim(p)
print(primitive_element)
```

#### DIFFIE-HELLMAN STARTER 3

Calculate the value of g^a mod p. Επίλυση

```
g=2
p=24103...
a=97210...
print(pow(g,a,p))
```

 $\frac{180685769784072652332258672182091135848942012812924807867393365353}{393068167618175384941171571417360435232355655878375925266106118632}\\027421488310488605016436812919171970740229157733048549951352236828\\939535952390140613802502252241242923897159127216051914467238953239\\367383226507005731948539979310118268217746536439627742471754343401\\766634380727697086447583039177640395755067836236831977656602511849\\206219694145126563805440017724857227134254861610396741199043735792$ 

#### DIFFIE-HELLMAN STARTER 4

You and Alice are now able to calculate a shared secret, which would be infeasible to calculate knowing only {g,p,A,B}. What is your shared secret? Επίλυση

```
g=2
p=2410...
```

7

```
A= 7024...
b= 120...
B=518...
#person A number a, person B num b, a,b private
#g,p public
#k=g^(ab) mod p= messfromA^b mod p
print(pow(A,b,p))
```

 $11741307404138206565338327460348419858773020863163883801659844 \\ 36672307692443711310285014138545204369495478725102882673427892 \\ 10453912095239378896105199290164969406317985359831147382034121 \\ 58799653431363514364105228507174084458020430031646583480065774 \\ 08558693502220285700893404674592567626297571222027902631157072 \\ 14333004311841846709423796559119844080397072660453780714670376 \\ 35716068614483546075026546647003904537944931767946789173526340 \\ 29713320615865940720837909466$ 

\*Υπάρχει η επίλυση για το *Diffie-Hellman Starter 5* στο github, όμως δεν προσμετρούνται οι πόντοι.

# CRYPTO ON THE WEB

#### **JSON WEB TOKENS**

#### TOKEN APPRECIATION

To prove them wrong, decode the JWT above to find the flag. There are online tools to do this quickly, but working with Python's <a href="PvJWT">PvJWT</a> library will prepare you best for future challenges.

Επίλυση

```
import jwt
encoded="eyJ0eX.."
print(jwt.decode(encoded,options={"verify_signature":
False}))#i can decode without a signature
```

crypto{jwt contents can be easily viewed}