

CSC 261/461 Database Systems

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Characteristics of Relations

Values in a tuple

- ▶ All values are considered atomic (indivisible).
- ▶ Each value in a tuple must be from the domain of the attribute for that column.
- ▶ A special NULL value is used to represent values that are unknown/inapplicable in certain tuples.



Constraints

Constraints

Determine which values are permissible and which are not in the database. They are of three main types:

1. Inherent or Implicit Constraints: These are based on the data model itself.
2. Schema-based or Explicit Constraints: They are expressed in the schema by using the facilities provided by the model.
3. Application based or semantic constraints: These are beyond the expressive power of the model and must be specified and enforced by the application programs.



Relational Integrity Constraints

- ▶ **Constraints** are conditions that must hold on **all** valid relation states.
- ▶ Three main types of constraints can be expressed in the relational model:
 - ▶ Key constraints
 - ▶ Entity integrity constraints
 - ▶ Referential integrity constraints



Key Constraints

- ▶ **Superkey** of R: Is a set of attributes SK of R such that:
 - ▶ No two tuples in any valid relation state $r(R)$ will have the same value for SK
- ▶ **Key** of R: A "minimal" superkey
 - ▶ a key is a superkey K such that removal of any attribute from K results in a set of attributes that is not a superkey

Question: Is the key also a superkey? Is the SK a key? [PK]



Example

Keys and Superkeys

CAR

<u>License_number</u>	Engine_serial_number	Make	Model	Year
Texas ABC-739	A69352	Ford	Mustang	02
Florida TVP-347	B43696	Oldsmobile	Cutlass	05
New York MPO-22	X83554	Oldsmobile	Delta	01
California 432-TFY	C43742	Mercedes	190-D	99
California RSK-629	Y82935	Toyota	Camry	04
Texas RSK-629	U028365	Jaguar	XJS	04



Entity Integrity

Entity Integrity

- ▶ The primary key attributes PK cannot have NULL values.
 - ▶ This is because primary key values are used to identify the individual tuples.
 - ▶ If PK has several attributes, NULL is not allowed in any of the attributes
- ▶ Other attributes of R may be constrained to disallow NULL values.



Referential Integrity

Referential Integrity

- ▶ A constraint involving two relations
- ▶ Used to specify a relationship among tuples in two relations:
 - ▶ The **referencing** relation and the **referenced** relation.
- ▶ attributes in FK have the same domain(s) as the attributes PK of R2.
- ▶ Attributes FK in R1 reference the PK attributes PK in R2.
 - ▶ A tuple $t1$ in R1 is said to reference a tuple $t2$ in R2 if $t1[FK] = t2[PK]$.



Referential Integrity

Example

EMPLOYEE

Fname	Minit	Lname	<u>Ssn</u>	Bdate	Address	Sex	Salary	Super_ssn	Dno
-------	-------	-------	------------	-------	---------	-----	--------	-----------	-----

DEPARTMENT

Dname	<u>Dnumber</u>	Mgr_ssn	Mgr_start_date
-------	----------------	---------	----------------

DEPT_LOCATIONS

<u>Dnumber</u>	<u>Dlocation</u>
----------------	------------------

PROJECT

Pname	<u>Pnumber</u>	Plocation	Dnum
-------	----------------	-----------	------

WORKS_ON

<u>Essn</u>	<u>Pno</u>	Hours
-------------	------------	-------

DEPENDENT

<u>Essn</u>	<u>Dependent_name</u>	Sex	Bdate	Relationship
-------------	-----------------------	-----	-------	--------------

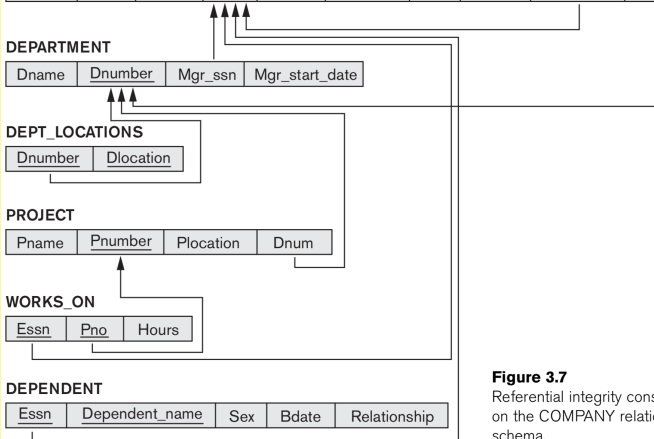


Figure 3.7

Referential integrity constraints displayed on the COMPANY relational database schema.



Structured Query Language

Declaring Keys

- ▶ There are two declarations that may be used to indicate keys
 1. **PRIMARY KEY**
 2. **UNIQUE**
- ▶ the effect of declaring a set of attributes S to be a key is:
 - ▶ two tuples in R cannot agree on all of the attributes in set S , unless one of them is **NULL**.
 - ▶ any attempt to insert or update a tuple that violates this rule is rejected.
 - ▶ if **PRIMARY KEY** is used, then attributes in S do not allow **NULL** as value.

Structured Query Language

Referential Integrity

- ▶ Referential integrity is specified via **FOREIGN KEY**.
- ▶ referential integrity can be violated
 - ▶ when tuples are inserted or deleted, or
 - ▶ when a foreign key or primary key attribute value is modified.
- ▶ The default action that SQL takes for an integrity violation is to reject the operation
 - ▶ this is known as the **RESTRICT** option.



Structured Query Language

Attribute and Tuple Constraints

In a **CREATE TABLE** statement, we can declare two kinds of constraints:

1. A constraint on a single attribute.

```
gender CHAR(1) CHECK (gender IN ('F', 'M'))
```

2. A constraint on a tuple as a whole.

```
CHECK (gender = 'F' OR name LIKE 'Mr.%')
```



Update Operations on Relations

Operations

- ▶ INSERT a tuple.
- ▶ DELETE a tuple.
- ▶ MODIFY a tuple.
- ▶ Integrity constraints *should not* be violated by the update operations.



Possible Violations

INSERT

- ▶ Domain constraint:
 - ▶ if one of the attribute values provided for the new tuple is not of the specified attribute domain
- ▶ Key constraint:
 - ▶ if the value of a key attribute in the new tuple already exists in another tuple in the relation
- ▶ Referential integrity:
 - ▶ if a foreign key value in the new tuple references a primary key value that does not exist in the referenced relation
- ▶ Entity integrity:
 - ▶ if the primary key value is NULL in the new tuple



ROCHESTER

Examples

EMPLOYEE

Fname	Minit	Lname	Ssn	Bdate	Address	Sex	Salary	Super_ssn	Dno
-------	-------	-------	-----	-------	---------	-----	--------	-----------	-----

DEPARTMENT

Dname	Dnumber	Mgr_ssn	Mgr_start_date
-------	---------	---------	----------------

EMPLOYEE

Fname	Minit	Lname	Ssn	Bdate	Address	Sex	Salary	Super_ssn	Dno
John	B	Smith	123456789	1965-01-09	731 Fondren, Houston, TX	M	30000	333445555	5
Franklin	T	Wong	333445555	1955-12-08	638 Voss, Houston, TX	M	40000	888665555	5
Alicia	J	Zelaya	999887777	1968-01-19	3321 Castle, Spring, TX	F	25000	987654321	4
Jennifer	S	Wallace	987654321	1941-06-20	291 Berry, Bellaire, TX	F	43000	888665555	4
Ramesh	K	Narayan	666884444	1962-09-15	975 Fire Oak, Humble, TX	M	38000	333445555	5
Joyce	A	English	453453453	1972-07-31	5631 Rice, Houston, TX	F	25000	333445555	5
Ahmad	V	Jabbar	987987987	1969-03-29	980 Dallas, Houston, TX	M	25000	987654321	4
James	E	Borg	888665555	1937-11-10	450 Stone, Houston, TX	M	55000	NULL	1

DEPARTMENT

Dname	Dnumber	Mgr_ssn	Mgr_start_date
Research	5	333445555	1988-05-22
Administration	4	987654321	1995-01-01
Headquarters	1	888665555	1981-06-19

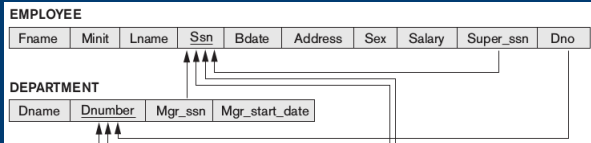
DEPT_LOCATIONS

Dnumber	Dlocation
1	Houston
4	Stafford
5	Bellaire
5	Sugarland

Insert <'Alicia', 'J', 'Zelaya', '999887777', '1960-04-05', '6357 Windy Lane, Katy, TX', F, 28000, '987654321', 4> into EMPLOYEE.

Insert <'Cecilia', 'F', 'Kolonsky', NULL, '1960-04-05', '6357 Windy Lane, Katy, TX', F, 28000, NULL, 4> into EMPLOYEE.

Examples



EMPLOYEE

Fname	Minit	Lname	Ssn	Bdate	Address	Sex	Salary	Super_ssn	Dno
John	B	Smith	123456789	1965-01-09	731 Fondren, Houston, TX	M	30000	333445555	5
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DEPT_LOCATIONS

Dnumber	Location
1	Houston
4	Stafford
5	Bellaire
5	Sugarland

Insert <'Cecilia', 'F', 'Kolonsky', '677678989', '1960-04-05', '6357 Windswept, Katy, TX', F, 28000, '987654321', 7> into EMPLOYEE.

Insert <'Cecilia', 'F', 'Kolonsky', '677678989', '1960-04-05', '6357 Windy Lane, Katy, TX', F, 28000, NULL, 4> into EMPLOYEE.

Possible Violations

DELETE

- ▶ DELETE may violate only referential integrity:
 - ▶ If the primary key value of the tuple being deleted is referenced from other tuples in the database
 - ▶ RESTRICT option: reject the deletion
 - ▶ CASCADE option: attempt to cascade the deletion by deleting tuples that reference the deleted tuple
 - ▶ SET NULL option: set the foreign keys of the referencing tuples to NULL



Possible Violations

UPDATE

- ▶ UPDATE may violate domain or NOT NULL constraint
- ▶ Any of the other constraints may also be violated:
 - ▶ Updating the primary key (PK):
 - ▶ Similar to a DELETE followed by an INSERT
 - ▶ Need to specify similar options to DELETE
 - ▶ Updating a foreign key (FK):
 - ▶ May violate referential integrity
 - ▶ Updating an ordinary attribute (neither PK nor FK):
 - ▶ Can only violate domain constraints



Questions?

