Time Complexity, Recursion and Function Memory Allocation.

Song Liu (song.liu@bristol.ac.uk)
GA 18, Fry Building,
Microsoft Teams (search "song liu").

Previous Lecture

- In the previous lecture, we talked about:
 - Nested if-else
 - Nested Loops
 - Early Loop Break/Restart
 - break;
 - continue;

Previous Lab

- Finding the maximum number among a, b and c.
- Do pairwise comparison between a, b and c.
 - o If a>b and a>c , a is the biggest.
 - If a>b and c>=a, c is the biggest.
 - If a<=b and b<c , c is the biggest.
 - If a<=b and c<=b, b is the biggest.

Previous Lab

• Write pseudo code.

```
Given three numbers: a, b and c
declare a variable: biggest
if a > b
    if a > c
        let biggest be a
    else
        let biggest be c
else
    if b < c
        let biggest be c
    else
        let biggest be b
return biggest
```

Previous Lab

Translate it into the actual c code.

```
int max(int a, int b, int c){
    int biggest = 0;
    if(a < b){
        if(c > b){
            biggest = c;
        }else{
            biggest = b;
    }else{
        if(c > a){
            biggest = c;
        }else{
            biggest = a;
        }
```

Today's Agenda

- Time Complexity of an Algorithm
 - Look back at prime finding algorithm.
- Recursion and Memory Layout

Finding Primes

Write a program that prints out all prime numbers from 1 to 100.

```
for(int i = 1; i <= 100; i++){
   int numfactor = 0;
   for (int j = 1; j <= i; j++){
      if(i % j ==0){
        numfactor = numfactor + 1;
      }
   }
   if(numfactor == 2){
      printf("%d ", i);
   }
}</pre>
```

How many iterations do we perform using the above code?

Making Code Faster!

```
1 + 2 + 3 + 4 \dots 100 = 5050 iterations.
```

Can we make the code run faster? i.e., doing the same things but with less iterations?

Yes! We give up on checking factors when numfactor > 2.

```
for(int i = 1; i <= 100; i++){
    int numfactor = 0;
    for (int j = 1; j <= i; j++){
        if(i % j == 0){
            numfactor = numfactor + 1;
            if(numfactor > 2){
                break; // stop the current loop!
                       // go to check i+1!
    if(numfactor == 2){
        printf("%d ", i);
```

Making Code Even Faster!

The previous code runs only 1890 iterations! Can we make the code even faster?

• For each i, there cannot be any factor bigger than i/2 except itself.

```
for(int i = 1; i <= 100; i++){
    int numfactor = 1; // start counting from 1!
    for (int j = 1; j <= i/2; j++){ // loop from 1 to i/2
        if(i % j == 0){
            numfactor = numfactor + 1;
            if(numfactor > 2){
                break;
    if(numfactor == 2){
        printf("%d ", i);
```

Time Complexity

The above code runs only for 715 iterations!

- You can write different algorithms that solves the same computational problem.
- Some are faster, some are slower.
- The number of elementary computing cycles (loop iterations) that are required for an algorithm is called the time complexity of the algorithm.
 - In our prime finding example, the time complexity of alg. 1 > alg. 2 > alg. 3.
 - Time complexity is an important metric of the algorithmic efficiency.

Unsolved Mystery

Given a computational problem, what is the lowest time complexity that one can achieve?

- For many problems, we know how fast we can get.
- Printing number from 1-100 requires 100 for loop iterations. These problems are called easy problems.
- Finding the solution of a sudoku game requires many more iterations. These problems are called **hard** problems.

Can we translate all **hard** problems into **simple** problems to achieve a lower time complexity?

Nobody knows.

- You cannot define a function inside another function.
- You can call a function inside another function.
 - A function can call itself!

- How do you countdown from 10 to 1?
 - o i.e., printing out number 10 9 8 ... 1?
- You can use a for loop.

```
for(int i = 10; i >= 1; i = i - 1) {printf("%d", i);}
```

- Or, you can use the following process:
 - i. initiate the countdown from 10
 - ii. If $i \ge 1$, then
 - a. print the current number i
 - b. continue the countdown from i-1.
- Weirdly, this procedure is defined in terms of itself, i.e.,
 the countdown process refers to itself.

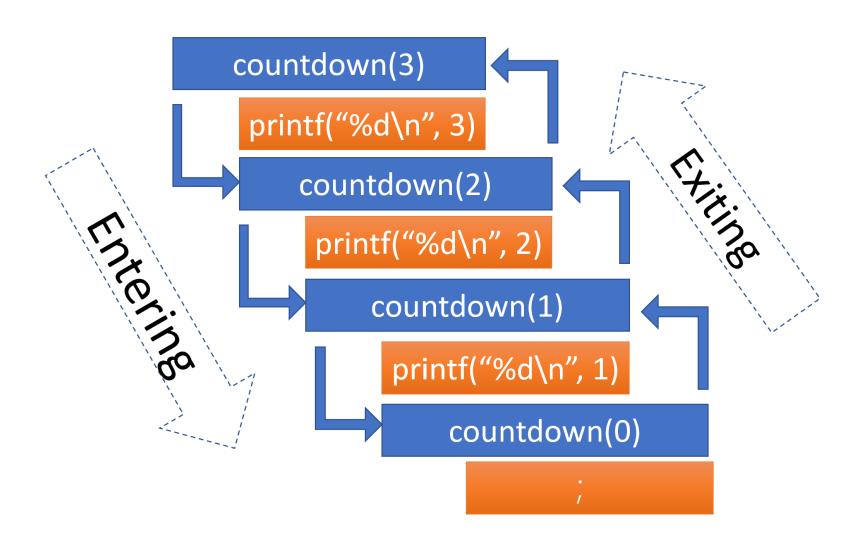
- A self-referential definition is called a recursion.
- Recursion is natural in math.
- Define \mathbb{N} , the set of all Natural Numbers.
 - \circ 0 is in \mathbb{N} .
 - \circ If i is in \mathbb{N} , i+1 is in \mathbb{N} too.
 - The smallest set that satisfies the above two criteria is the set of all natural numbers.
- Some mathematical process does not have a closed form, and can only be defined via a recursive process.

```
void countdown(int i){
   if(i >= 1 ){
      printf("%d\n", i); // print
      countdown(i - 1); //countdown from n-1
   }
}

void main(){
   countdown(3); //initiate the countdown
}
```

- Prints out 3 2 1.
- For each n > 0, it prints out n and continues to countdown with a smaller number n-1.

How does Recursion Work?



How does Recursion Work?

```
void countdown(int i){
   printf("enter countdown(%d)\n", i); //new!
   if(i >= 1 ){
      printf("%d\n", i);
      countdown(i - 1);
   }
   printf("exit countdown(%d)\n", i); //new!
}
```

```
enter countdown(3)
3
enter countdown(2)
2
enter countdown(1)
1
enter countdown(0)
exit countdown(0)
exit countdown(1)
exit countdown(1)
exit countdown(2)
exit countdown(3)
```

Memory Allocation for Functions

- When the function is being executed on the CPU, its data (such as variables declared in the function) and code are temporarily stored in the memory.
- The memory region for storing function data in the current program is called "stack".

Stack

• Stack is a data structure that the newest element is placed on top of the old elements.

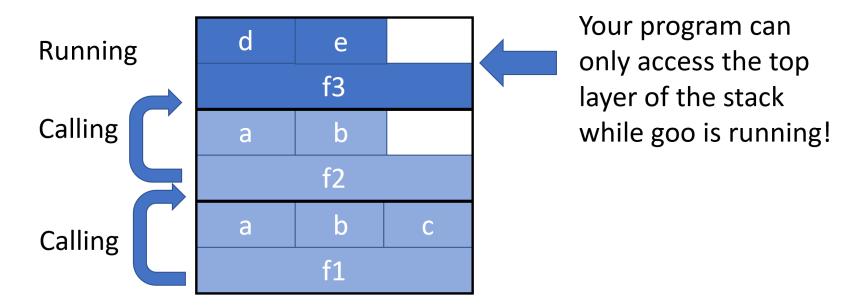


- 1. When a function is called, its data and code is added to the top of the stack.
- 2. CPU can only access data and code from the top stack.
 - It means, CPU can only access variables from the current function that is being called!
- 3. When a function finishes its execution, its data is removed from the stack and the space it occupies is freed for future calls of functions.

Consider a situation where function f1 calls f2 and f2 calls f3.

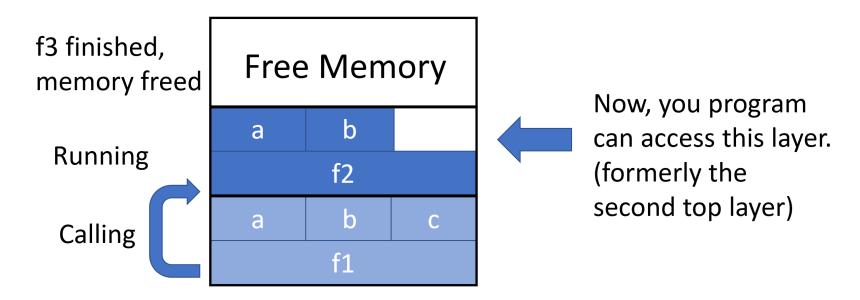
```
void f3(){
    int d = 6, e = 7;
};
void f2(){
    int a = 4, b = 5;
    f3();
void f1(){
    int a = 1, b = 2, c = 3;
    f2();
};
void main(){
    f1();
```

Below is the stack while f3 is running.



The only variables are accessible to us are d and e.

When f3 finishes running, its memory is freed.



The only variables are accessible to us are a and b in f2.

When f2 finishes running, its memory will be freed too and only variables in f1 will be accessible.

Local Variables

Variables declared inside the function are called **local variables** (This includes all input argument variables!).

- They can only be accessed by the function where it is defined.
 - They cannot be accessed by other functions.
- Why? The program can only access the top layer of the stack, which stores variables of the function that is currently running.
- In the "f1-f2-f3" example, your program cannot access a and b defined in f1 while f2 is running.

Stack is a highly efficient memory allocation/release mechanism.

- The memory allocation and release are all automatically handled by the OS.
- However, there is only a limited stack space for each program (determined by the OS). If a single function occupies a large memory space, or the call stack gets too "tall", we may run out of stack memory and an execution error will be raised by the OS.
 - This out-of-memory error is called "Stack Overflow".

Conclusion

- 1. There can be multiple algorithms that solves the same computational task.
- 2. The computing cycles required for each algorithm is called time complexity.
- 3. The function definition that refers to itself is called recursion.
- 4. Function data and code are stored in stack memory.
 - i. Local variables can only be accessed by the function where it is defined.

Labs

1. Download the lab file and unzip it to your labpack folder.

Homework 1

- 1. We are going to verify the time-complexity for different prime finding algorithms I introduced in the lecture.
- 2. Open prime1.c.
- 3. Slightly modify the code, so it prints out the number of elementary computing cycles (loop iterations) it has gone through in order to find all prime numbers <= 100.
 - For example, in the pseudo code below, the outer loop repeats 10 times and the inner loop repeats 10 times, the program will go through 10*11 elementary computing cycles before it stops.

```
for i from 1 to 10
for j from 1 to 11
print out i*j
```

Homework 2

- 1. Modify prime2.c and prime3.c and print out the number of elementary computing cycles.
- 2. Imagine a program that finds all prime numbers which are smaller than n by using the algorithm in prime1.c
 - Express the number of cycles in terms of n.

Homework 3 (21-22 Exam Question)

```
#include <stdio.h>
void g(){
    int b = 4, c = 5;
    printf("g is being called!\n");
void f(){
    int a = 1, b = 2;
    g();
    printf("g has been called!\n");
void main(){
    f();
```

- 1. Without running the program, draw the stack memory diagram (see lecture slides) when the program is printing out g is being called!
- 2. Verify your answer using debugger (see video).

Homework 4 (submit)

- O. Open prime4.c.
- 1. Write a function <code>num_factor</code> that takes input <code>i</code> and <code>j</code>, and output an integer.
 - For any input i > j > 1,
 - The function returns the number of i 's factors from 2 to j.
 - \circ e.g. i = 5, j = 4, the output is 0.
 - \circ e.g. i = 16, j = 5, the output is 2.
 - You are not allowed to use any loop. Use recursion.
- 2. Test it in main() with different input i,j.

Homework 4 (submit)

- 3. Now, make use of num_factor function you just wrote, write some additional code in main (don't modify num_factor), so it prints out all prime numbers from 1 to 100.
- 4. Rename prime4.c as you have done before, submit it to the blackboard.