LFS - Crash recovery

The check-point region (CR) must be updated atomically

- → LFS keeps two CRs and writing a CR is done in 3 steps
 - I. writes out the header with a timestamp # I
 - 2. writes the body of the CR
 - 3. writes one last block with another timestamp #2
- √ Crash can be detected if timestamp # I is after #2
- ✓ LFS will always choose the most recent and valid CR
- ✓ All logs written after a successful CR update will be lost in case of a crash

LFS - Disk Cleaning (a.k.a Garbage Collection)

LFS leaves old version of file structures on disk

- → LFS keeps information of the version of each segment and runs a disk cleaning process
 - A cleaning process removes old versions by compacting contiguous blocks in memory
 - That cleaning process runs when the disk is idle or when running out of disk space