

# Solution

- ⦿ **Problem:** each process has a page table that maps all pages in its address space
- ✓ **Solution:** we only need to map the portion of the address space actually being used
- ➔ Use another level of indirection : **two-level page tables**

# Two-Level Page Tables

Virtual addresses have three parts

- **a master page number** i.e the index in the master page table (a.k.a "page directory") that maps to a secondary page table
- **a secondary page number** i.e the index in the secondary page table (a.k.a "page table") that maps to the physical memory
- **an offset** that indicates where in physical page address is located