

Asymmetric encryption for confidentiality

Bob encrypts a message m with Alice's public key KpA \rightarrow Nobody can decrypt m, except Alice with her private key Ks_A ✓ Confidentiality without the need to exchange a secret key







Ks_A, Kp_A

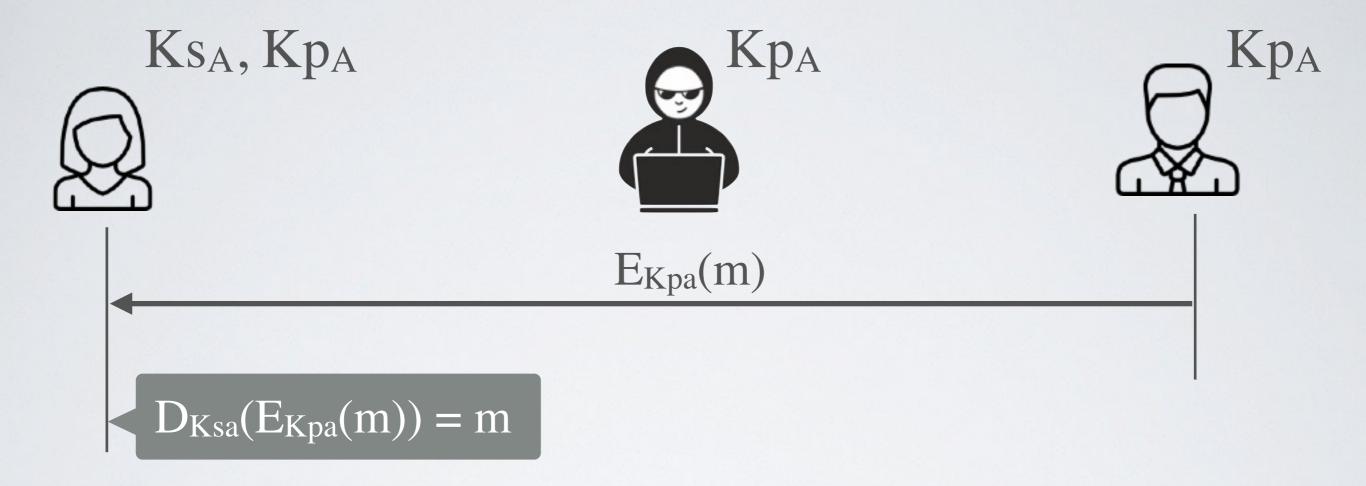






 $D_{Ksa}(E_{Kpa}(m)) = m$

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Functional Requirements

 $D_{Ks}(E_{Kp}(m)) = m$ and $D_{Kp}(E_{Ks}(m)) = m$ for every pair (Kp, Ks)

- ✓ Generating a pair (Kp, Ks) is easy to compute (polynomial)
- ✓ Encryption is easy to compute (either polynomial or linear)
- ✓ Decryption is easy to compute (either polynomial or linear)
- Finding a matching key Ks for a given Kp is hard (exponential)
- Decryption without knowing the corresponding key is hard (exponential)