CSE 330: Operating Systems

Class: 15 **Date: 10/6**

Fall 2016

Note Title

Semaphores

→ critical sections (locking)

→ synchronization

Monitor → critical section

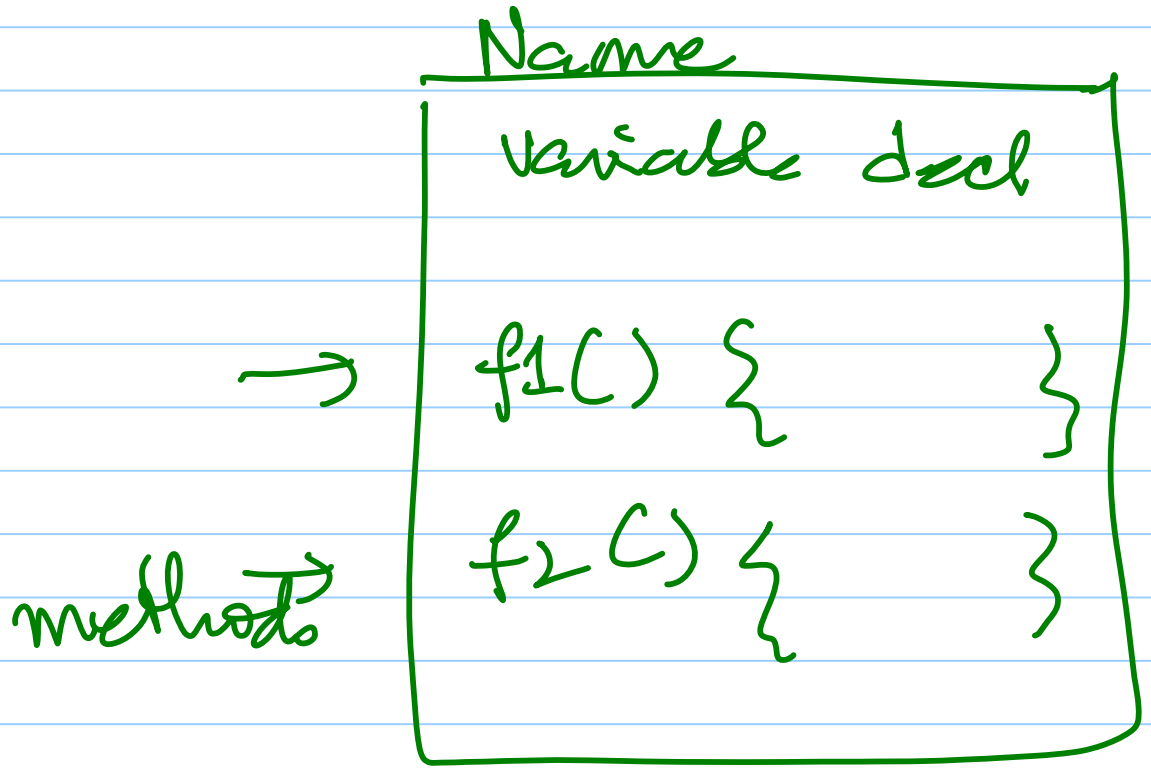
→ structured (object oriented)

Synchonization

→ conditions

↳ variables in a monitor

Monitor → class



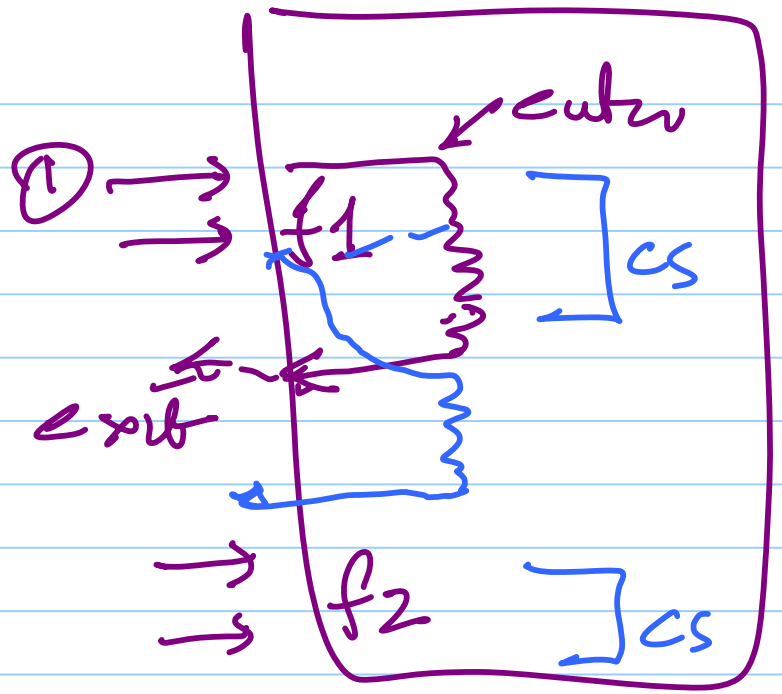
threads can call monitor methods

Name.f1()
Name.f2()

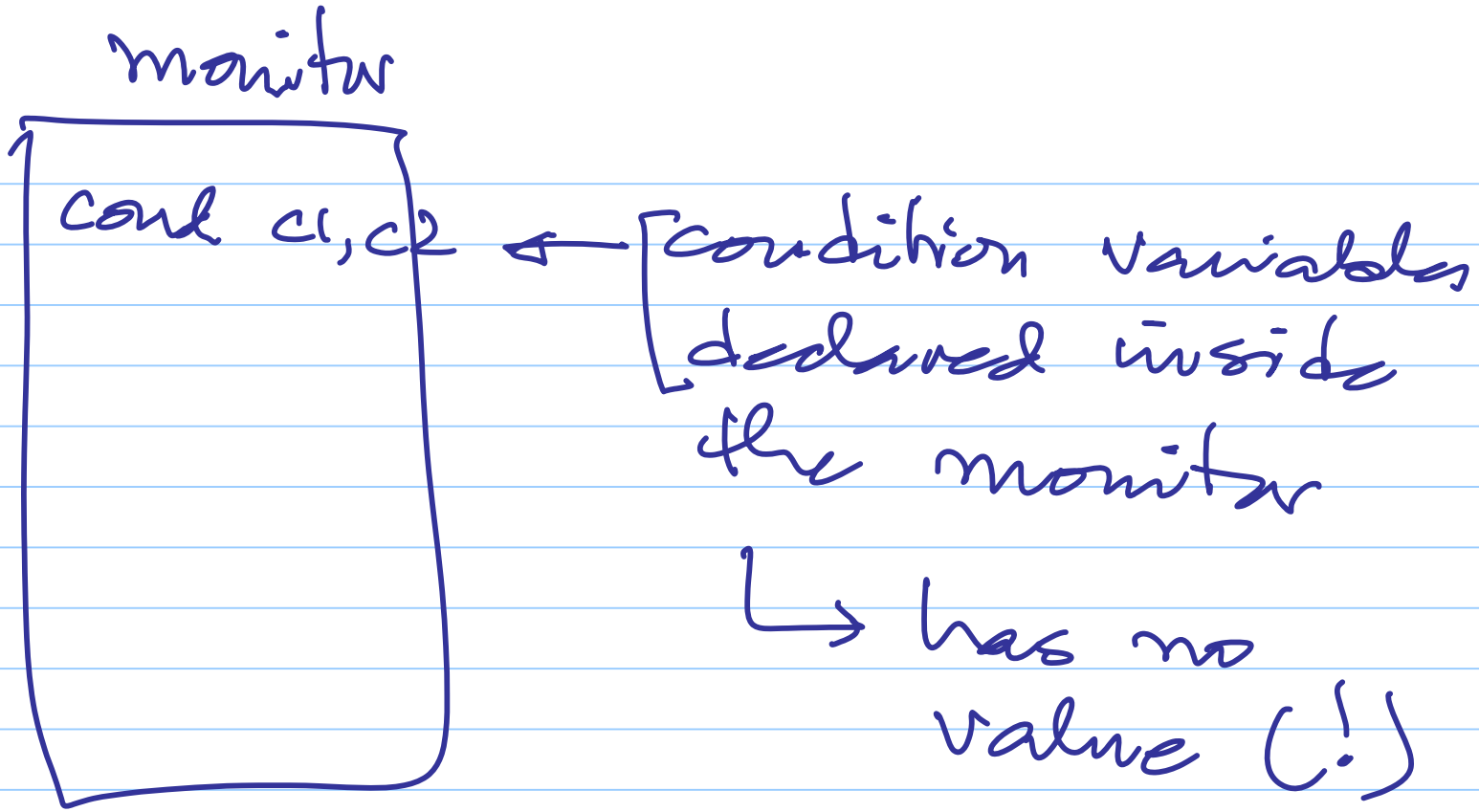
Monitor methods

→ can be called by multiple threads — concurrently

→ BUT only 1 thread is allowed to execute inside one monitor @ any point in time

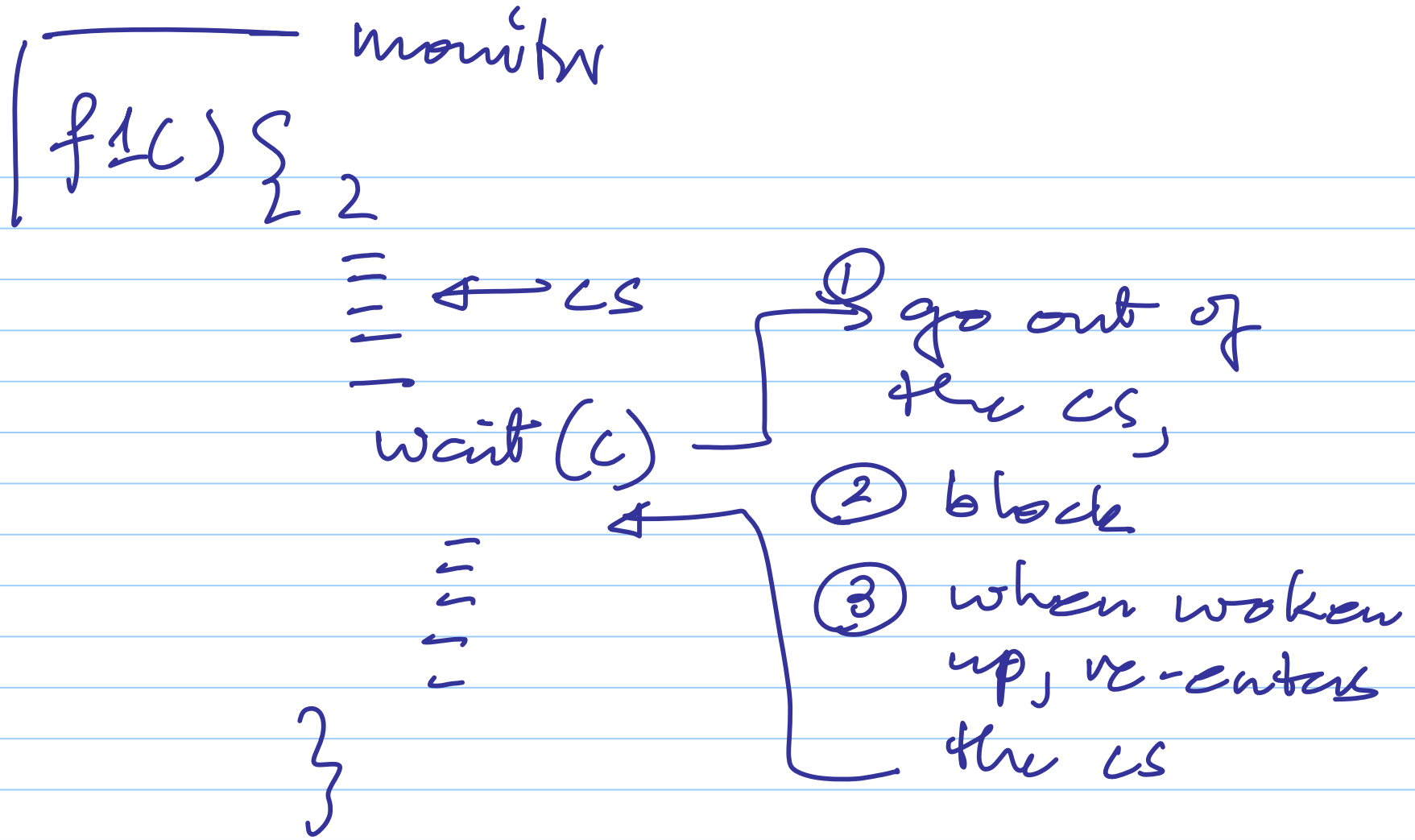


Big critical section


monitor
cond c1, c2
condition variables
declared inside
the monitor
has no
value (!)

2 functions are defined on a condition

1. wait(c) → block (unconditional block ... "forever")
2. signal(c)


monitor
f1(c) {
cs
wait(c)
}
1 go out of the cs,
2 block
3 when woken up, re-enters the cs

```
f1() {
    ...
    signal(c)
    ...
}
```

→ wake up a thread blocked on that condition – if any

↳ if no thread is blocked

→ NO_OP

if T_1 signals T_2 → T_2 is to wakeup.

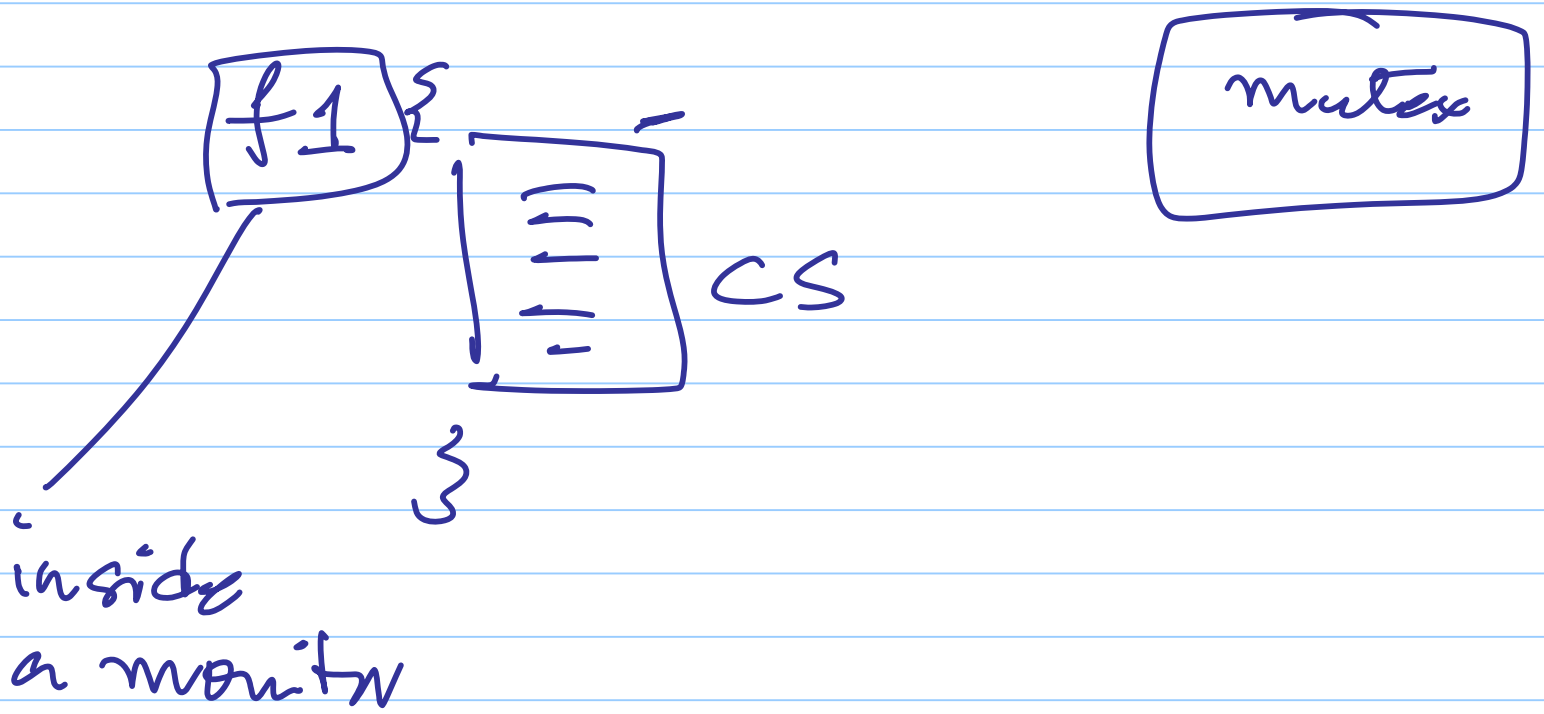
option 1 — T_1 continues to execute and T_2 has to wait till T_1 exits the monitor or T_1 blocks.

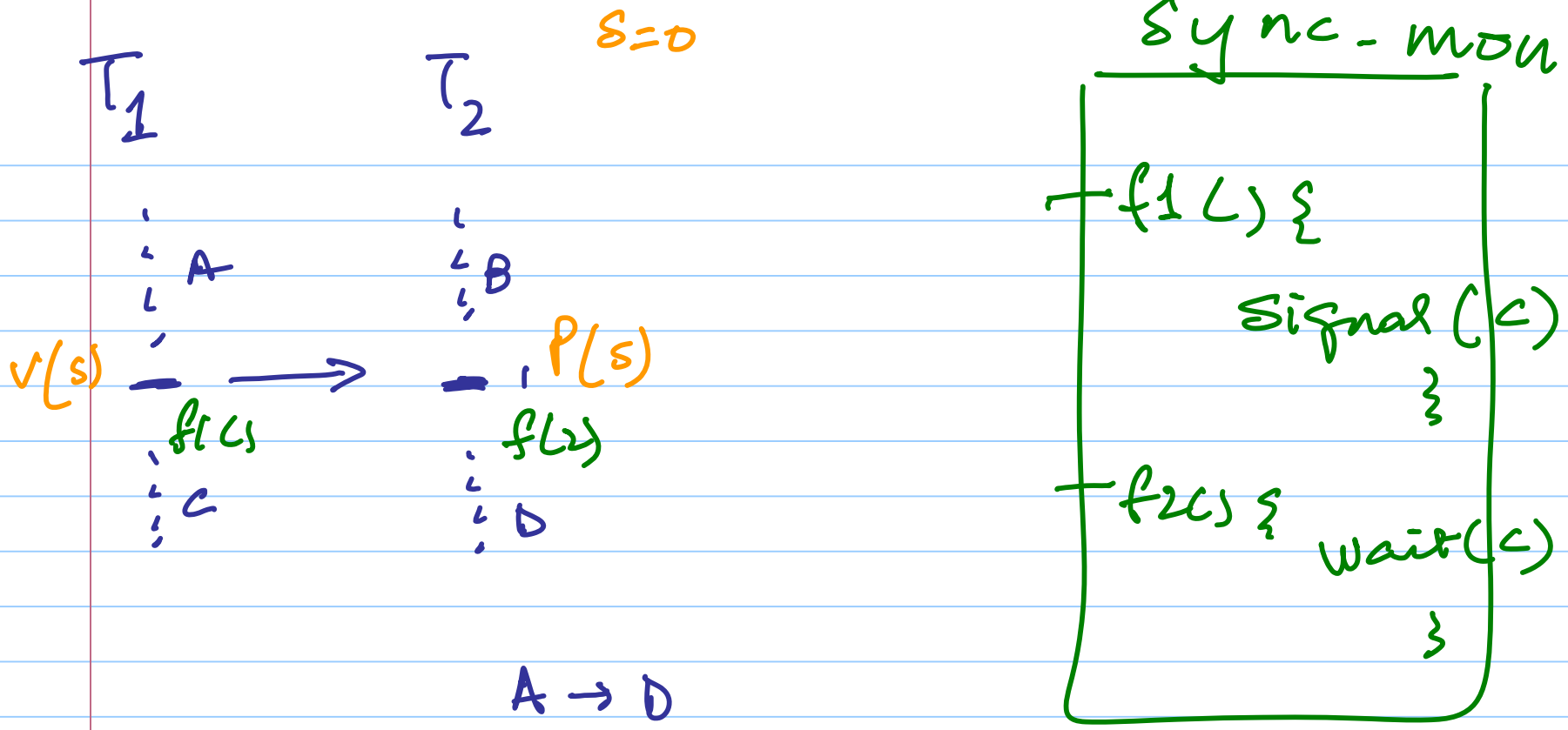
pthreads →

option 2 → (T_2 has priority over other threads)

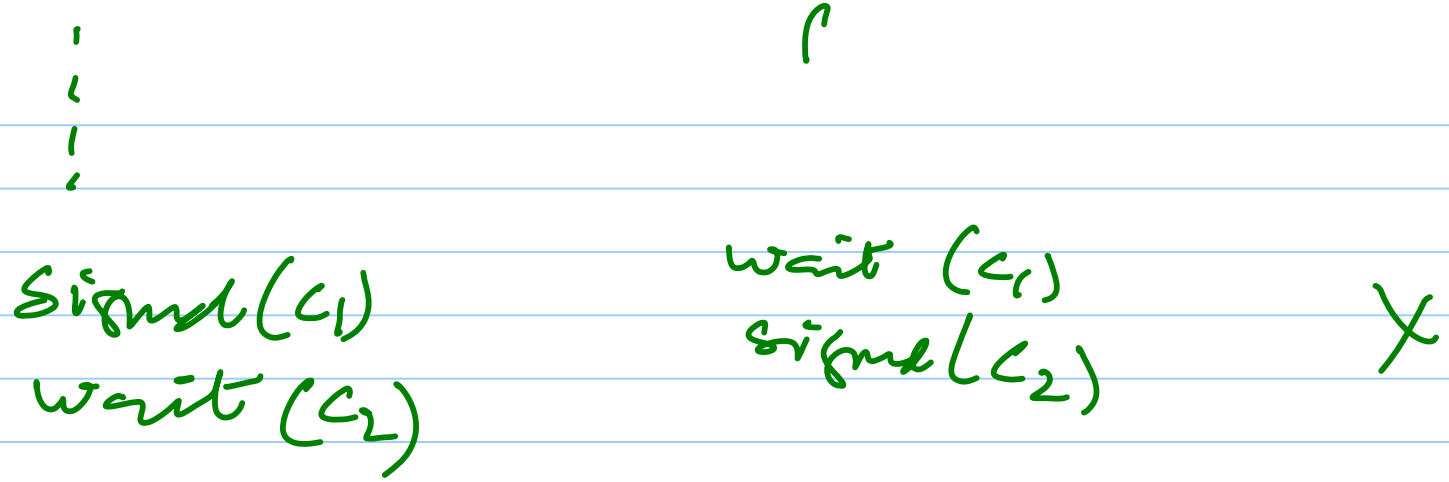
T_1 blocks. T_2 executes, T_1 executes when T_2 leaves / blocks.

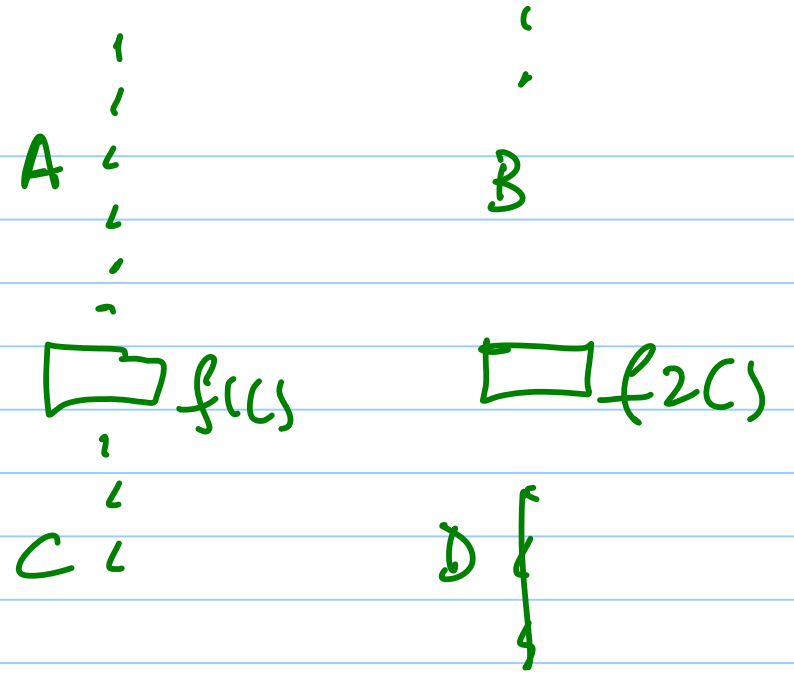
programming with monitors




S=0
T1
T2
A
B
V(S)
P(S)
f1()
f2()
C
D
A → D

Sync-mon
f1() {
signal (c)
}
f2() {
wait(c)
}

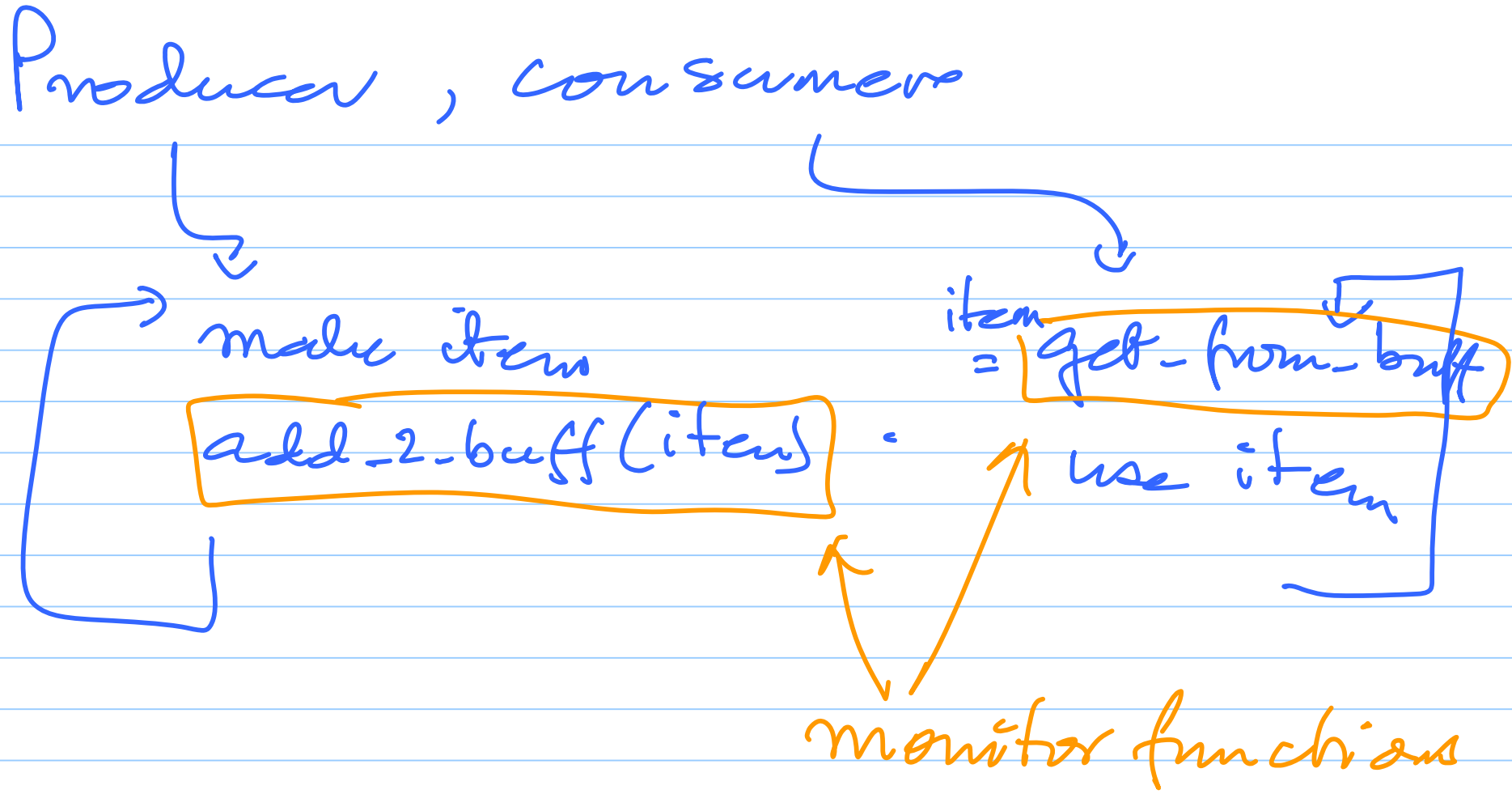

signal(c1)
wait(c2)
wait(c1)
signal(c2)
X




```
int done = 0 ;
cond c ;
```

```
f1 { done = 1 ;
     signal (c) ; }
```

```
f2 { if (done != 1)
       wait (c) }
```


Producer , consumer
make item
add_2_buff(item)
item
= get_from_buff
use item
monitor functions

MoN

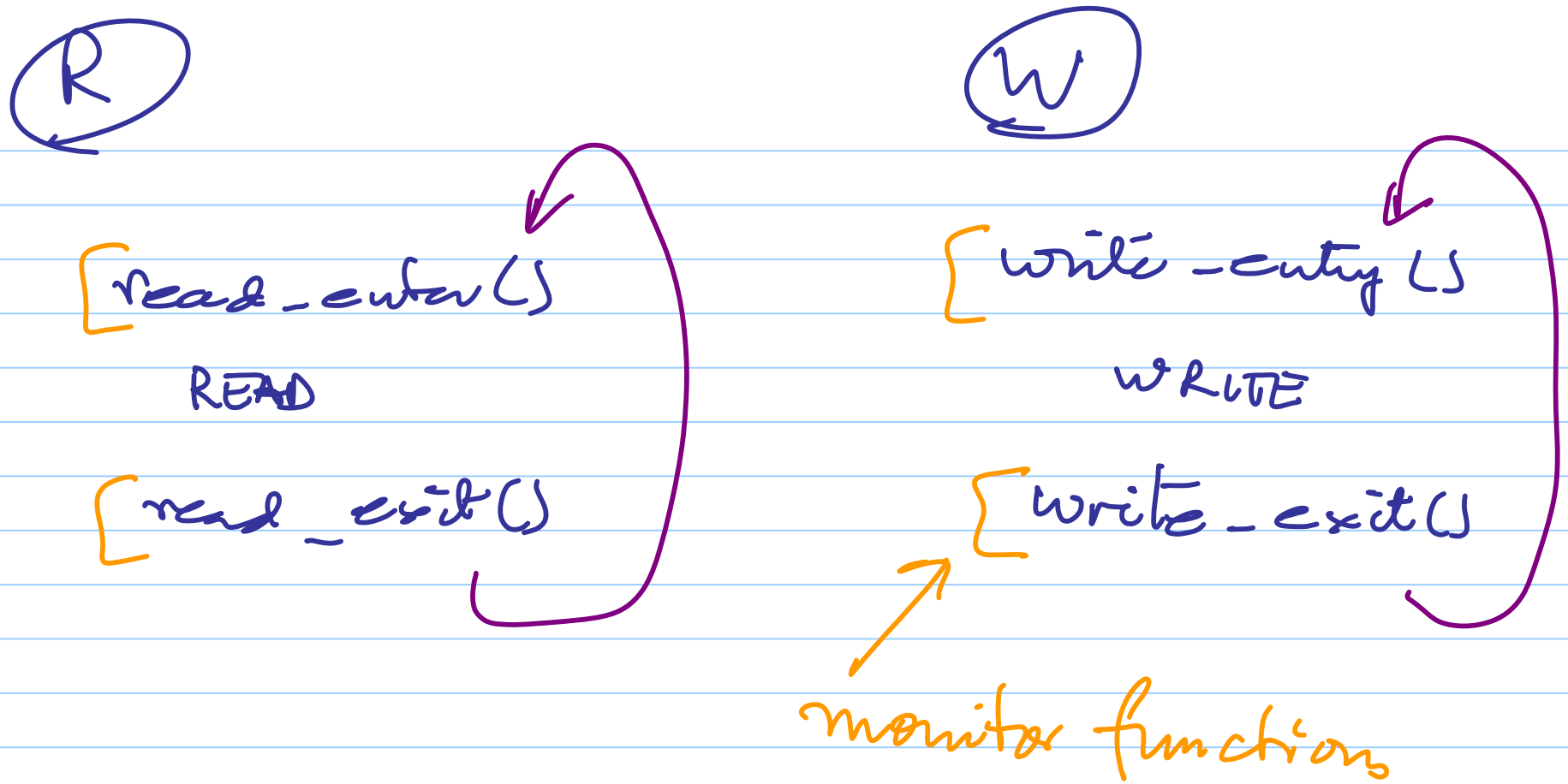
buffer[N], in, out = 0 count = 0

[item] [int] [int]

conditions → prod, cons

```
add 2 buff()
{ if (count==N) wait (prod)
  count++
  buff[in] = item
  in = (in+1) % N
  signal (cons);
}
```

```
get from buff()
{
  if (count==0)
    wait (cons)
  count--
  item = buff[out]
  out = (out++) % N
  signal (prod)
}
```


R
read_enter()
READ
read_exit()
W
write_entry()
WRITE
write_exit()
monitor functions

```
READ_ENTER
{ if((wc>0) OR (wwc>0))
        {rwc++ wait(R) rwc--}
  rc++; signal(R);
}

READ_EXIT
    rc--
  { if(rc==0) { signal(W) }
    }
```

```
write_enter()
  { if (rc > 0) or (wc > 0)
        { wwc++; wait(w); wwc--}
    wc++
  }
write exit  wc--
  { if (rwc > 0) signal (R);
      else signal (w); }
```

pickup

```
while (!(chop[LEFT] && chop[RIGHT]))
    wait(self[i]);
chp[LEFT] = chp[Right] = false
```

putdown

```
chp[L] = chp[R] = true
signal(L); signal(R)
```