

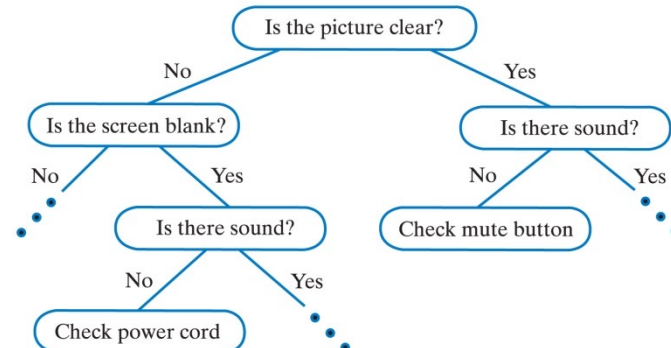
Binary Trees

Trees

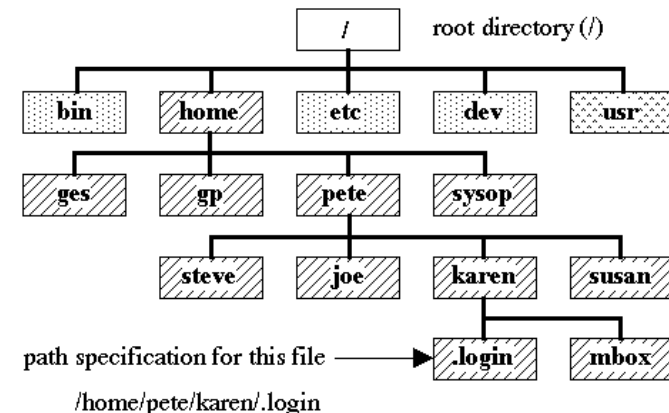
Binary Trees

Binary Tree Nodes

Recursive tree algorithms



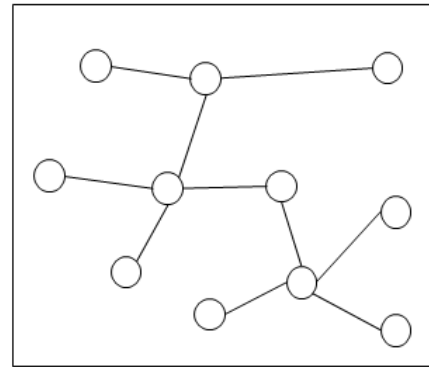
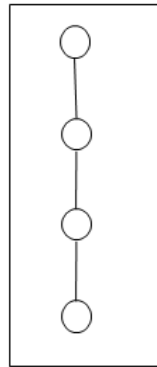
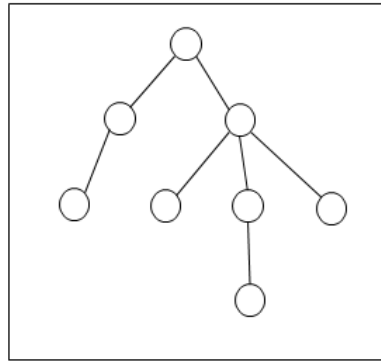
UNIX File System Hierarchy (sample)



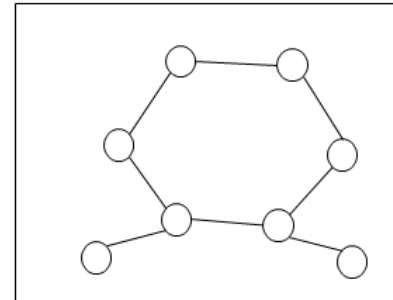
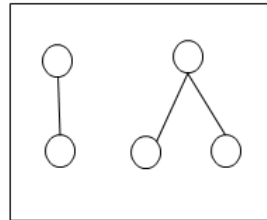
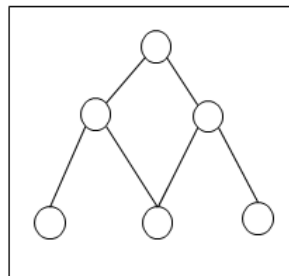
Trees

A **tree** is a connected undirected simple graph with no cycles

trees:



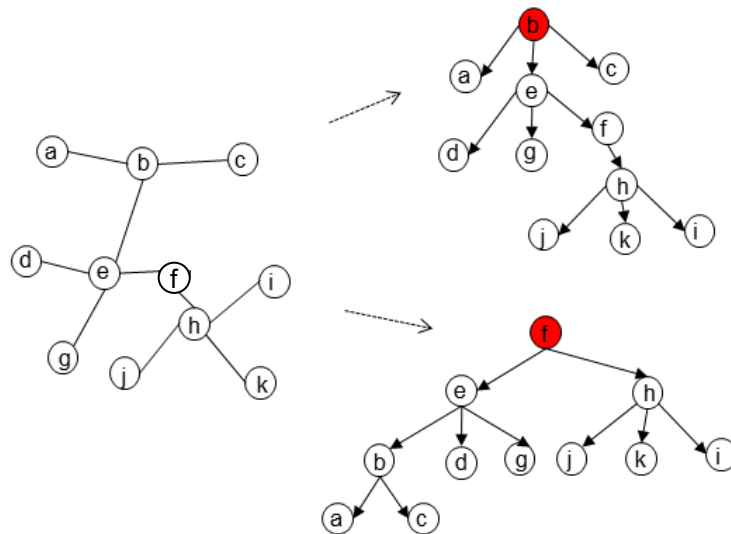
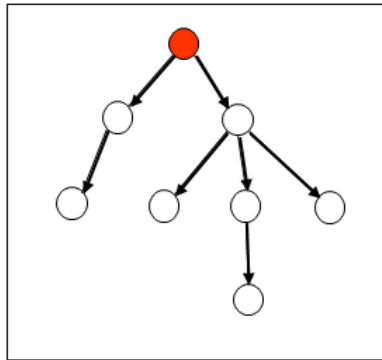
not
trees:



Rooted Trees

Usually, when we think of trees, we assume there is a root.

A **rooted tree** is a tree in which one of the vertices is designated the **root**, and all edges are then directed away from that root.



Describing vertices in rooted trees

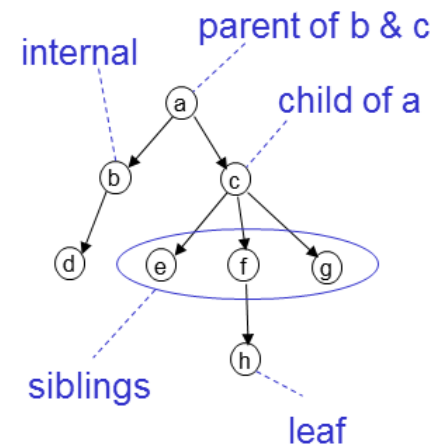
If there is a directed edge from x to y , then x is the **parent** of y , and y is a **child** of x .

If two vertices y and w have the same parent, then they are **siblings** of each other.

A vertex with no children is a **leaf**. A vertex with children is an **internal** vertex.

The **ancestors** of a vertex v are all the vertices in the path from v to the root (except for v itself).

The **descendants** of a vertex v are all the vertices that have v as an ancestor.

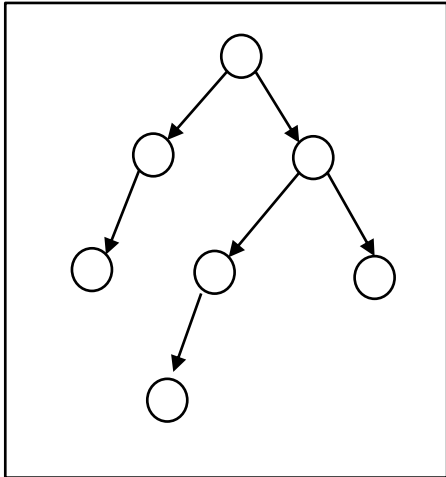


Binary Trees

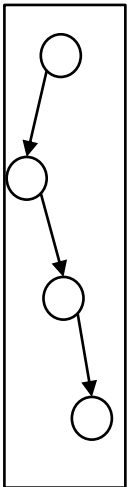
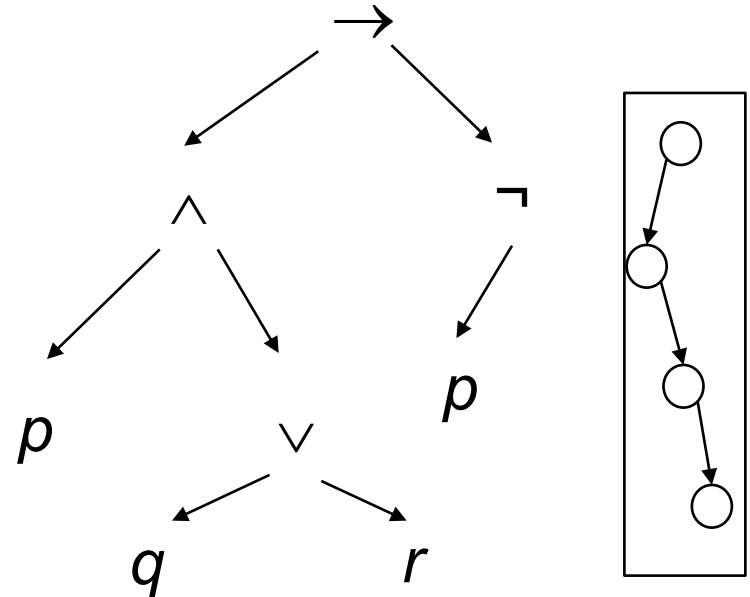
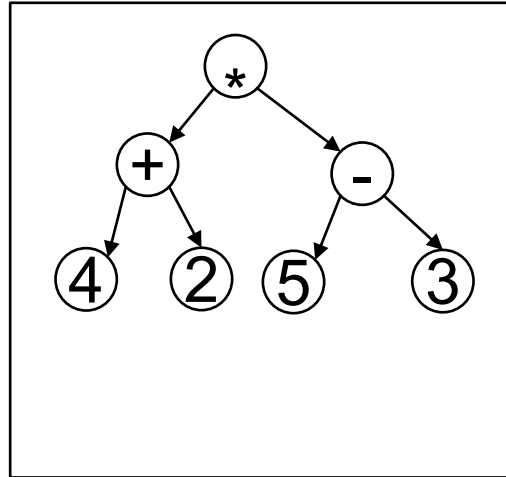
A *binary tree* is a rooted tree in which:

- every node has at most 2 children
- the children of a node are identified as left child and right child

The depth of a tree is the length of the longest path from the root node to a leaf node



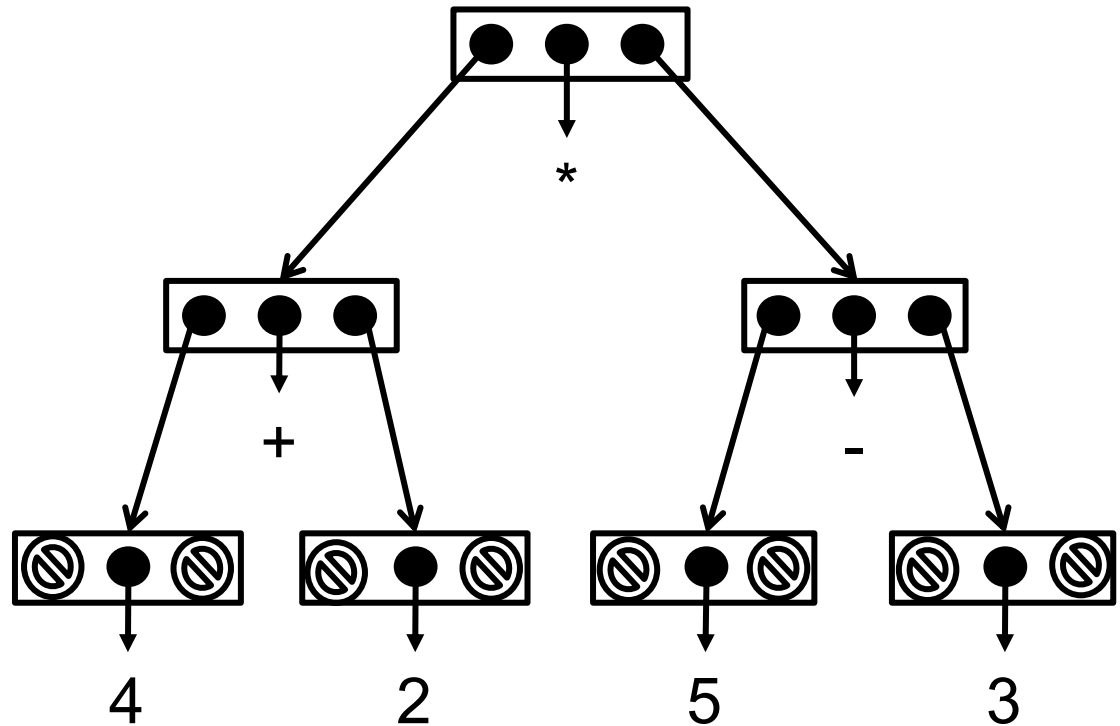
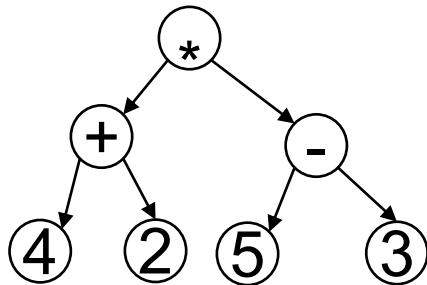
depth == 3



BinaryTreeNode

BinaryTreeNode

element
leftchild
rightchild

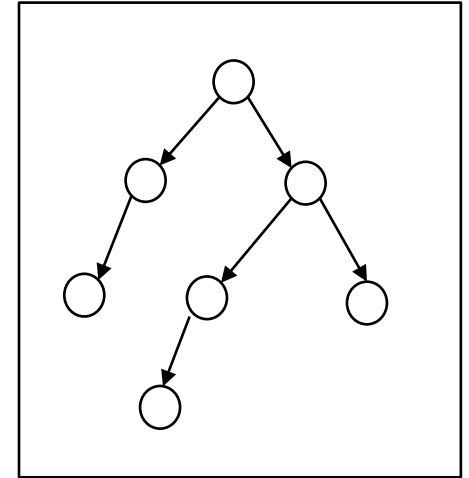


Computing the height of a node


The height of a node is the length of its longest (directed) path to a leaf.

Recursive definition:

height(node) = 0 if node is a leaf

$$\text{height}(\text{node}) = 1 + \max(\text{height}(\text{left}), \text{height}(\text{right}))$$


```
def height(node):  
    if node == None:  
        return 0  
    elif node.leftchild == None and node.rightchild == None:  
        return 0  
    else:  
        return 1 + max(height(node.leftchild),  
                        height(node.rightchild))
```

 recursion

recursive

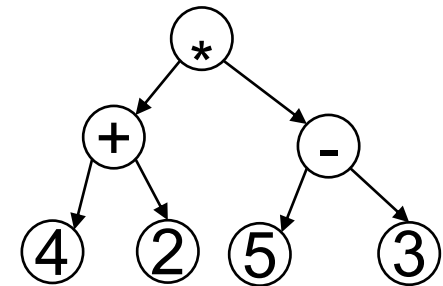
Preorder traversal

to visit a node:

read the element first, then visit the children in left-to-right order

'preorder' because we do the
parent's element before the children

```
def preorder_print(node):  
    if node:  
        print(node.element)  
        preorder_print(node.leftchild)  
        preorder_print(node.rightchild)
```



*
+
4
2
-
5
3

← recursive

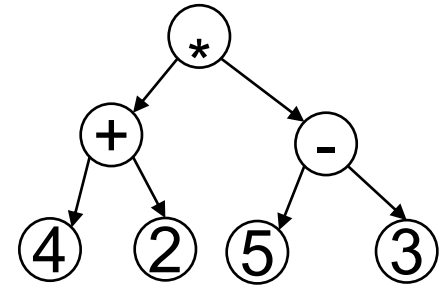
```
def preorder_str(node):  
    if node:  
        outstr = str(node.element)  
        outstr = outstr + preorder_str(node.leftchild)  
        outstr = outstr + preorder_str(node.rightchild)  
        return outstr  
    else:  
        return ''
```

* + 4 2 - 5 3

Inorder traversal

to visit a node:

visit the leftchild, then the parent's element,
then the right child



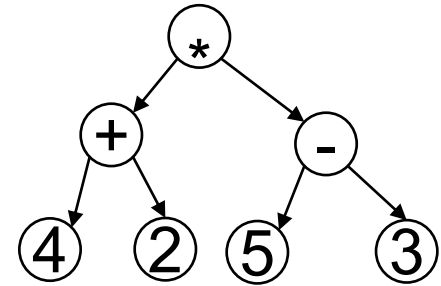
$((4 + 2) * (5 - 3))$

```
def inorder_str(node):  
    if node:  
        if node.leftchild or node.rightchild:  
            outstr = '(' + inorder_str(node.leftchild)  
            outstr += node.element  
            outstr += inorder_str(node.rightchild) + ')'  
            return outstr  
        else:  
            return str(node.element)  
    else:  
        return ''
```

Post-order traversal

Exercise: what would 'post-order' traversal mean?

Write Python code to produce a post-order traversal print of the tree on the right.



Evaluating Expression Trees

evaluate node:

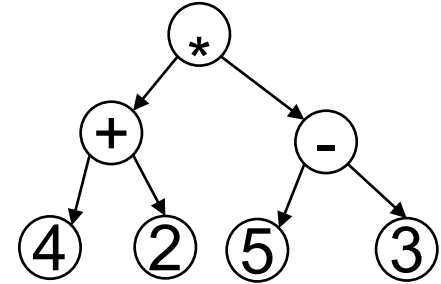
if node element is a number, return it
else

evaluate left child

evaluate right child

determine the operator in the node

apply leftvalue operator rightvalue



Exercise: implement this in Python

Exercise

How would you implement the preorder traversal *without* using explicit recursion?

Next Lecture

Binary Search Trees