

# Binary Search Trees

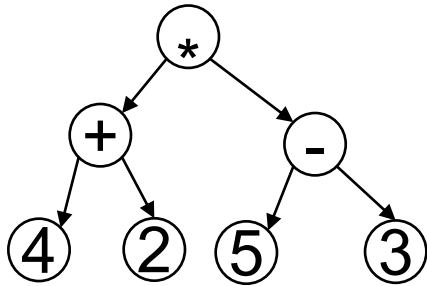
## Recap

Definition of binary search tree

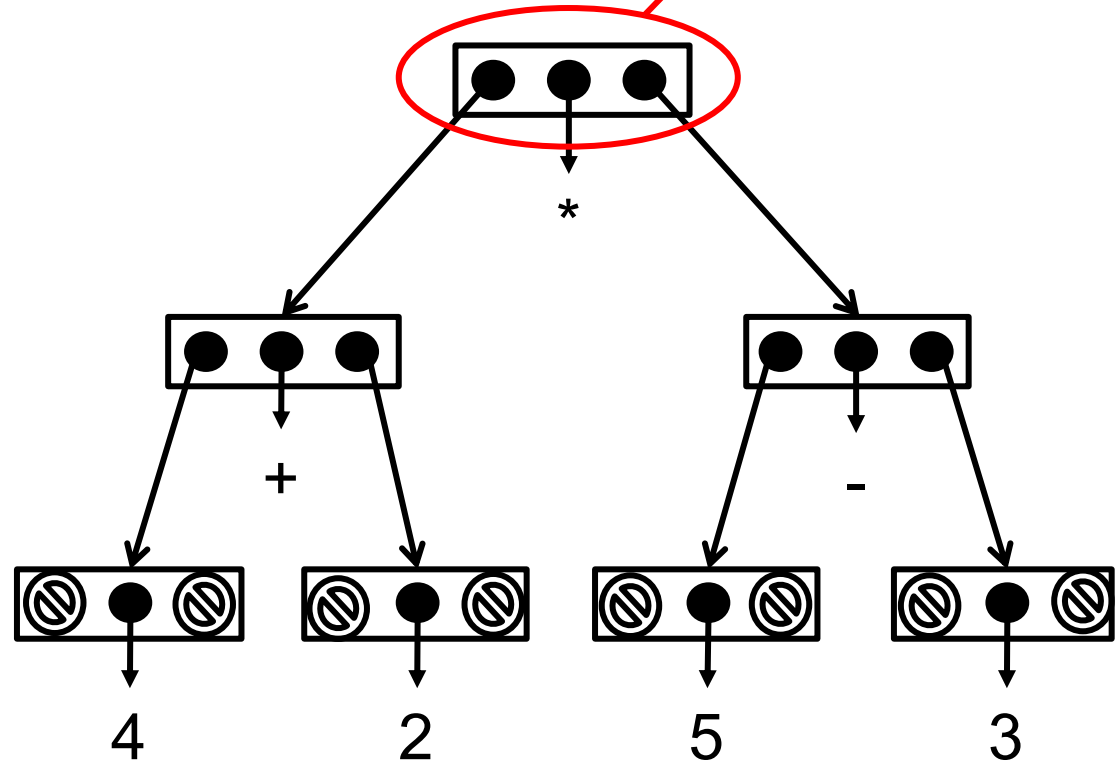
Searching a binary search tree

Adding to a binary search tree

# Reminder: BinaryTreeNode



Note: these are not lists;  
they are object instances  
of the BinaryTreeNode  
class



BinaryTreeNode

element  
leftchild  
rightchild

# Let's look at searching again ...

|                    | Array-based  | Linked list   |
|--------------------|--|---|
| Arbitrary access   | $O(1)$   | $O(n)$  |
| Add at end         | $O(n)$ worst case<br>$O(1)$ on average             | $O(1)$  |
| Add in middle      | $O(n)$   | $O(1)$ if given reference<br>$O(n)$ if given position |
| Delete from end    | <b>Bad!</b> $O(n)$ worst case<br>$O(1)$ on average | $O(1)$ if given reference<br>$O(n)$ if given position |
| Delete from middle | $O(n)$   | $O(1)$ if given reference<br>$O(n)$ if given position |
| sorting            | (see semester 2)                                   | (see semester 2)                                      |
| searching          | Unordered: $O(n)$<br>Ordered: $O(\log n)$          | Unordered: $O(n)$<br>Ordered: $O(n)$                  |

Good! (circled in green)

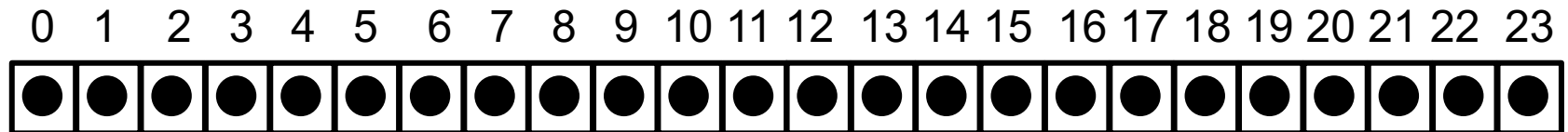
Good! (circled in green)

Bad! (circled in red)

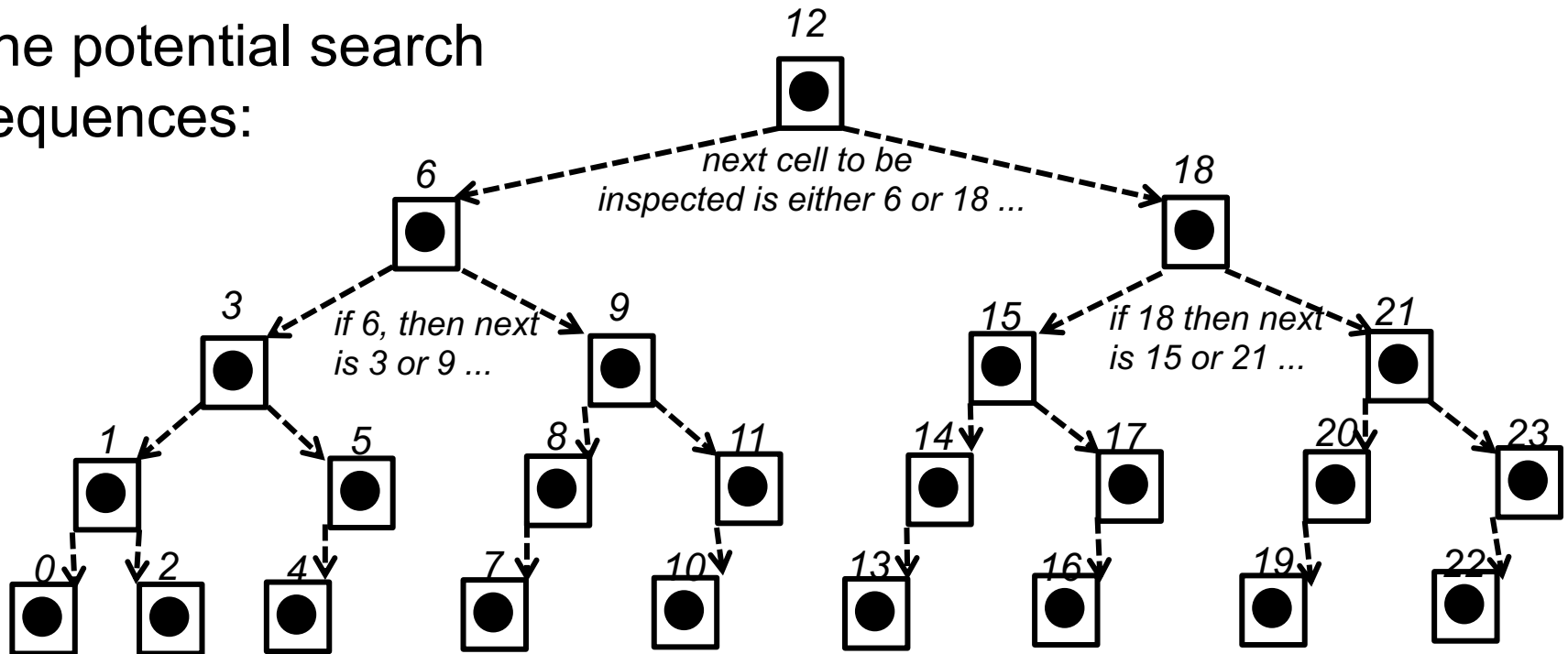
Can we build a linked structure to get the search complexity of array-based lists, but with faster add/delete?

# Sequence of cell lookups in binary search

A sorted list:



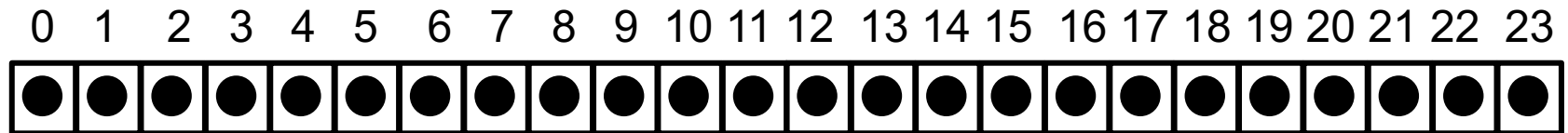
The potential search sequences:



So can we create a linked tree to replicate these possible search sequences (and then try to have efficient add/delete)?

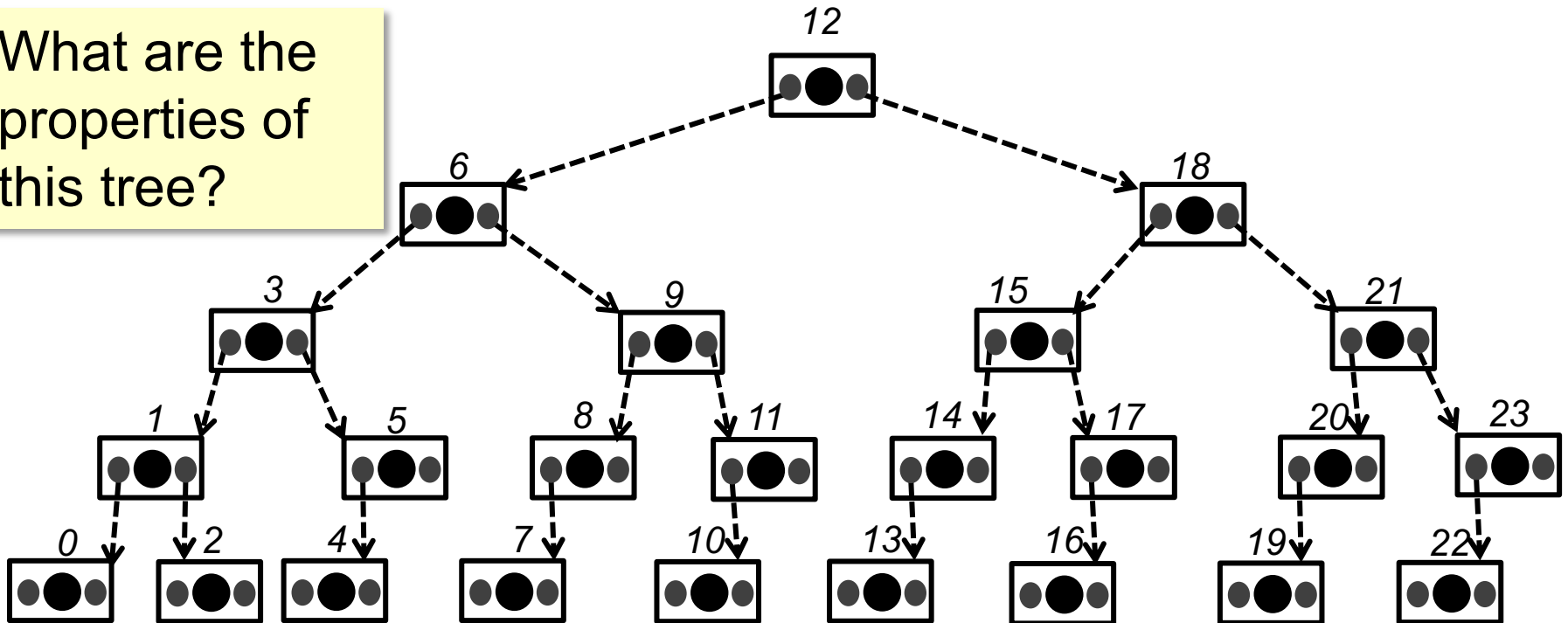
# Linked structures for efficient searching?

A sorted list:



And its (ideal) corresponding binary search tree:

What are the properties of this tree?

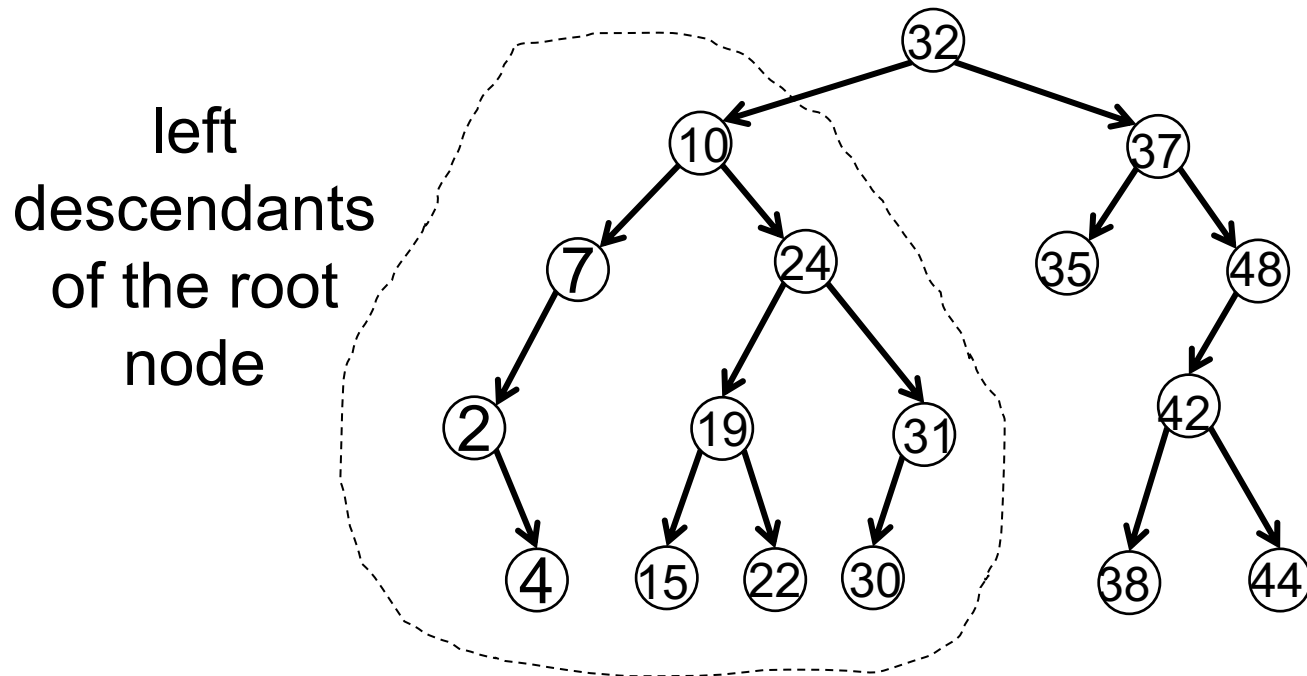


# Binary Search Tree

A *binary search tree* is a binary tree representing an ordered sequence of elements, where:

all left descendants of a node have values less than the node's value

all right descendants of a node have values greater than the node's value



**Note!** The arrows represent the tree structure, not the order of the elements in the sequence

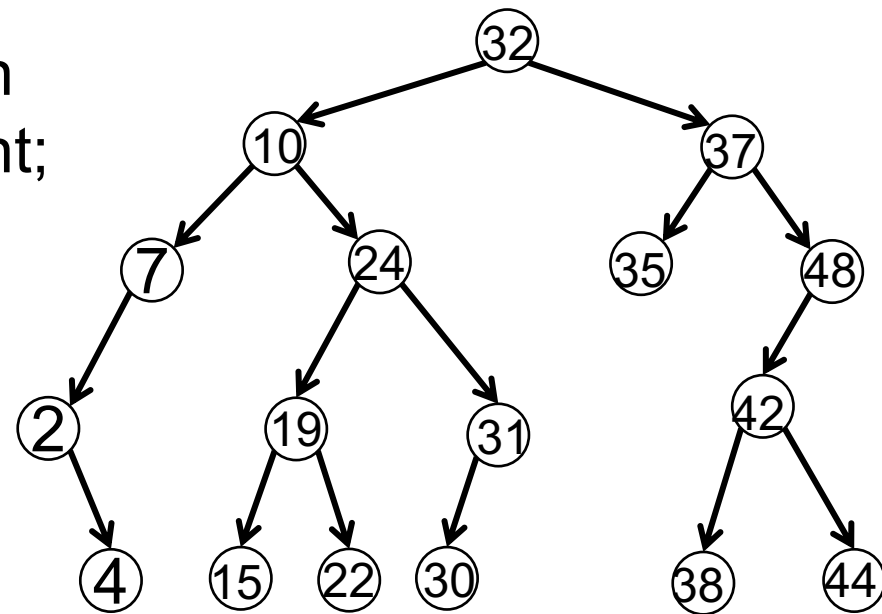
**Exercise:** how would you print the ordered sequence represented by this tree?

# Searching a binary search tree

Input: a reference to the root node of the tree; a target to search for  
Goal: if the target we are searching for is in the tree, return the node that has it as its element; else return None

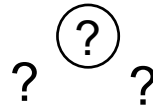
- all left descendants of a node have values less than the node's value
- all right descendants of a node have values greater than the node's value

From now on, we only have a BST; we do not have any list



# Searching a binary search tree

Is 27 in this tree





# Searching a binary search tree

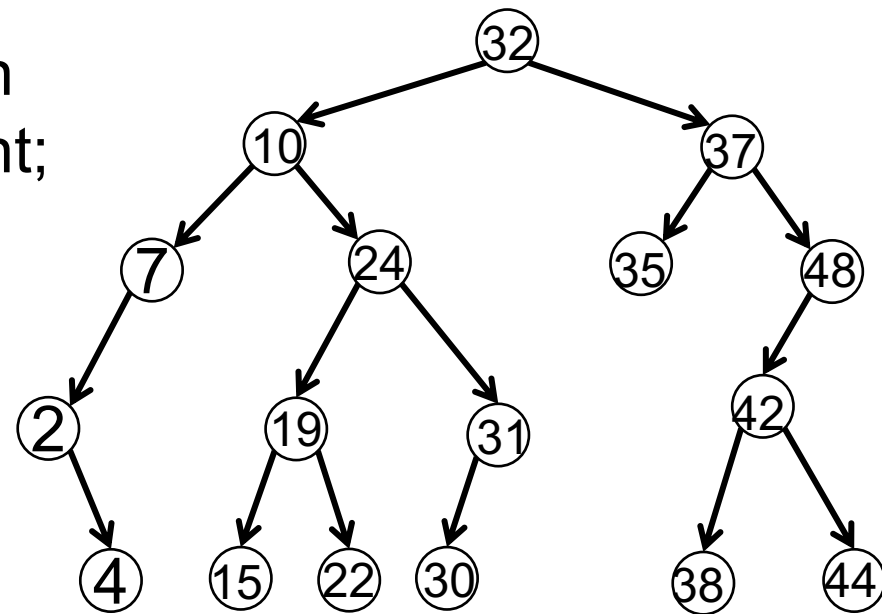
Input: a reference to the root node of the tree; a target to search for

Goal: if the target we are searching for is in the tree, return the node that has it as its element; else return None

- all left descendants of a node have values less than the node's value
- all right descendants of a node have values greater than the node's value

```
search(node, item):  
    if node == None  
        return None  
    if node's element > item  
        return search(leftchild, item)  
    else if node's element < item  
        return search(rightchild, item)  
    else return node
```

From now on, we only have a BST; we do not have any list

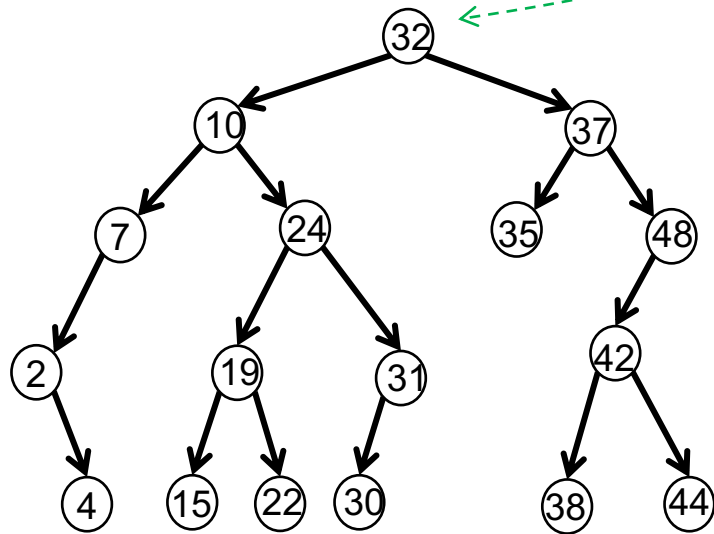


If  $h$  is the height of the root, then this is  $O(h)$

# Adding a node to a BST

*We know the value we want to add:  $x$*

*We know the location of this node:*



Requirement: maintain the order property

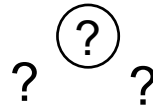
Aim: minimise the work

When asked to add, if we allow only one copy of each element, then

- 1) first we need to check that the element is not already there
- 2) then we need to add it

# Searching a binary search tree

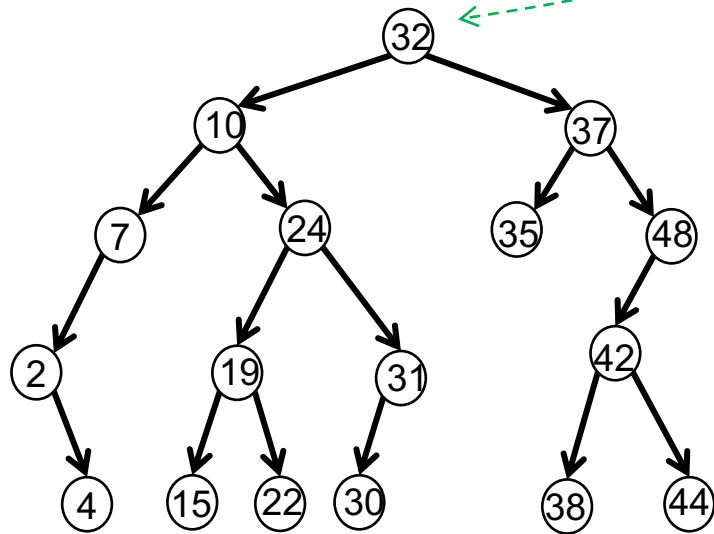
Add 27 to this tree



# Adding a node to a BST (ii)

*We know the value we want to add:  $x$*

*We know the location of this node:*



If a node is not in the tree, then the search ends when we reach a null value.

Solution: add the element there, and don't change anything else in the tree

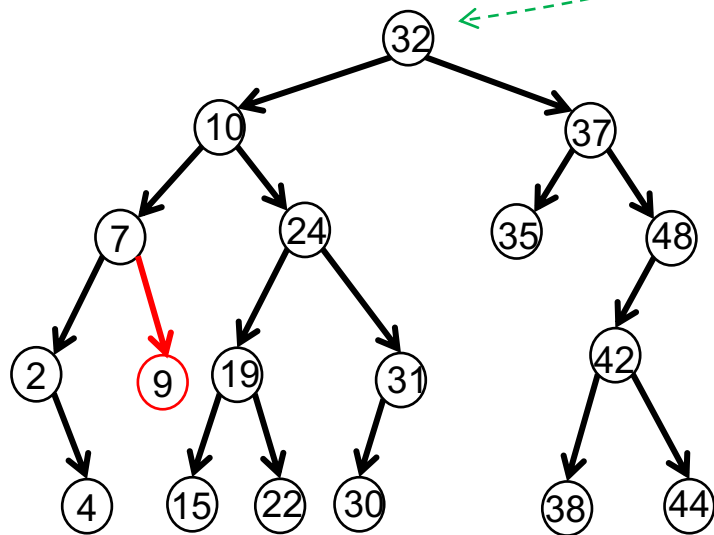
Class exercise: add 9

Class exercise: add 43

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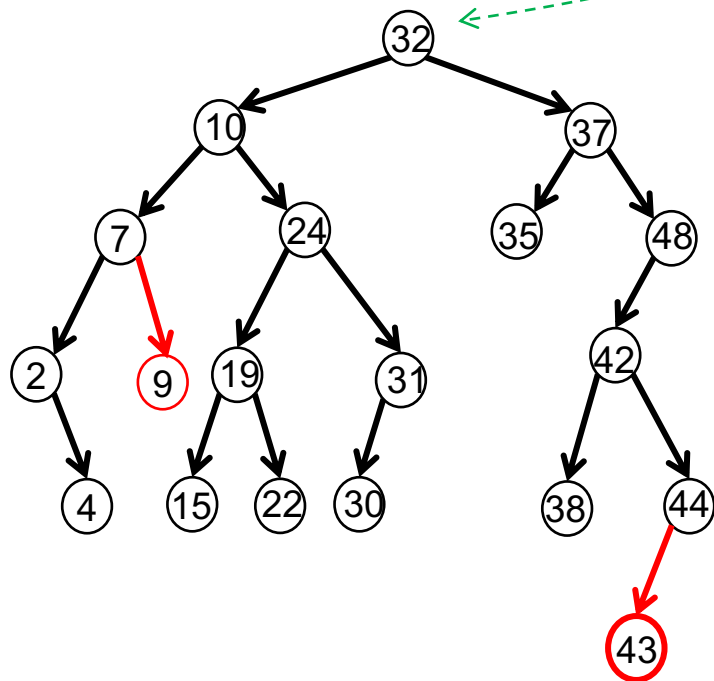
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# Adding a node to a BST (ii)

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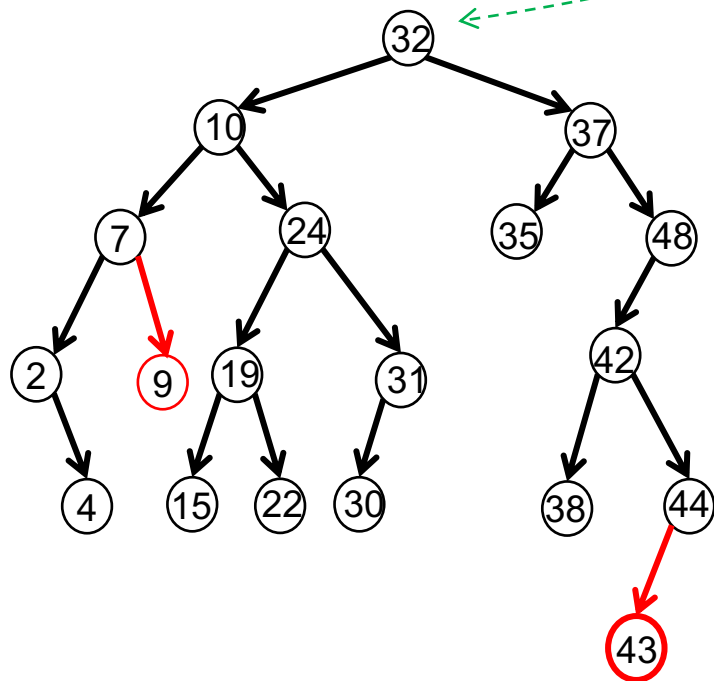
Class exercise: add 9

Class exercise: add 43

# Adding a node to a BST (iii)

We know the value we want to add:  $x$

We know the location of this node:

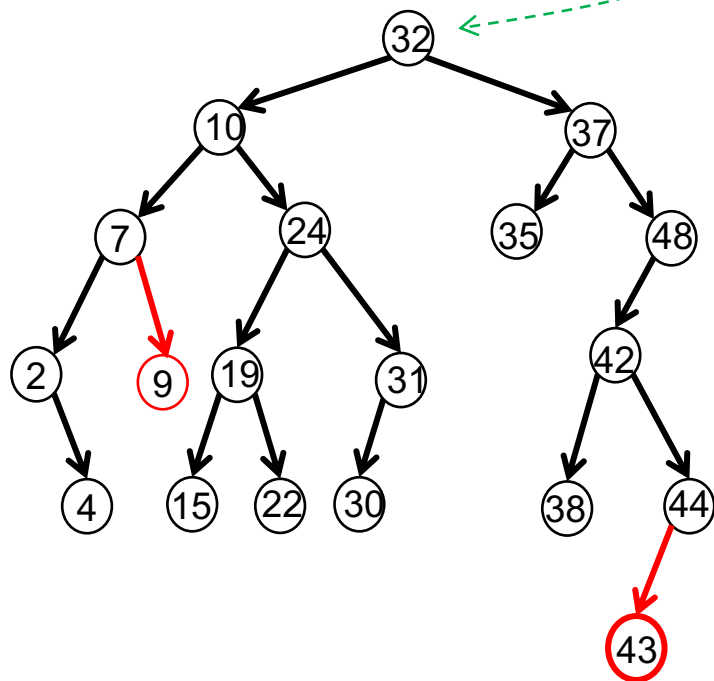


```
add(node, x):  
    if x < current element  
        if no left child  
            add x as new left child  
        else  
            add(node.left, x)  
    else if x > current element  
        if no right child  
            add x as new right child  
        else  
            add(node.right, x)  
    else  
        #do nothing - already there
```

# Adding a node to a BST (iv)

*We know the value we want to add:  $x$*

*We know the location of this node:*



For a given BST, for each possible addition, if we are restricted to simply adding the new element in an empty place, then there is only one possible location in the tree.

Class exercise: what are the empty places in the tree on the left, and what values can they take?



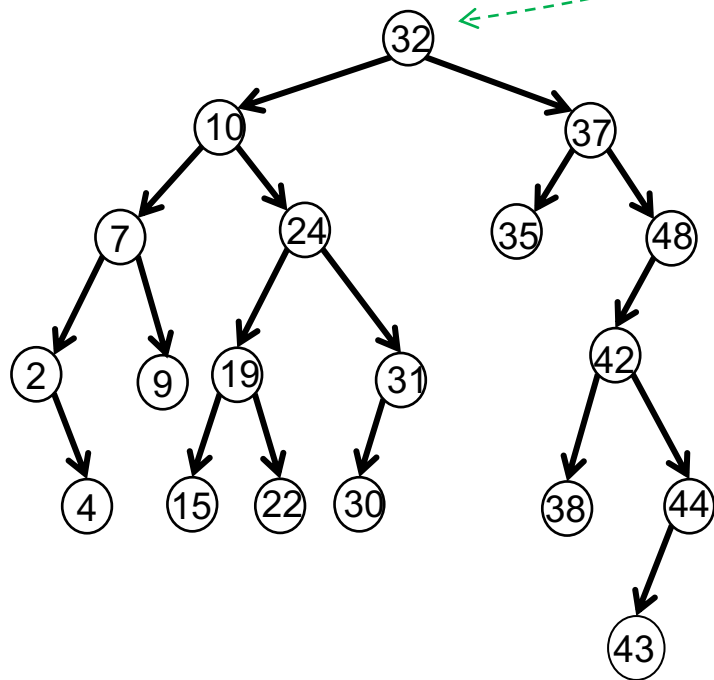
# Adding a node to a BST: complexity

To find the node or its natural position is  $O(\text{height of tree})$   
Adding the node at that position is a constant number of operations.

# Removing a node from a BST

*We know the value we want to remove:  $x$*

*We know the location of this node:*



Requirement: maintain the BST property

Aim: minimise the work

Start by finding the node in the tree (and if it is not there, stop)

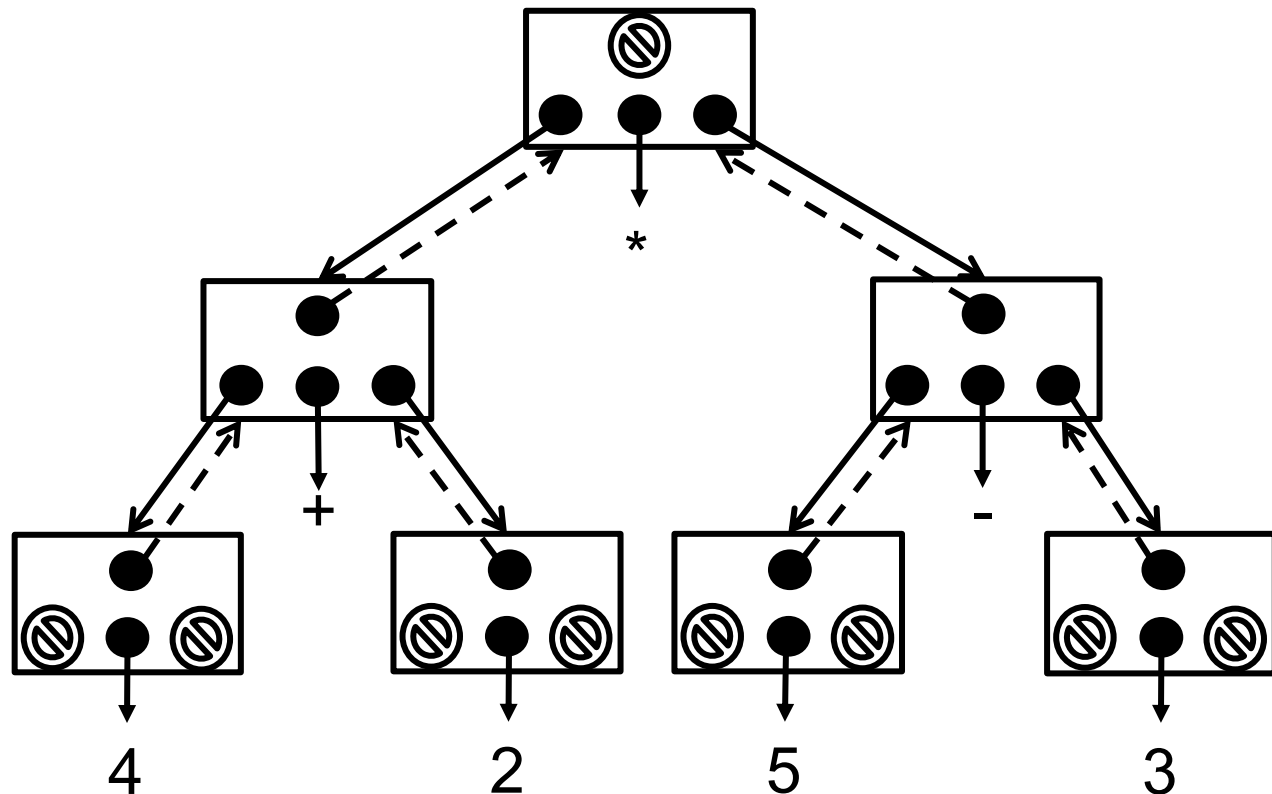
Handle the removal by breaking it down into different cases

Change the representation so that each node also points to its parent.

# BinaryTreeNode

BinaryTreeNode2

element  
leftchild  
rightchild  
parent

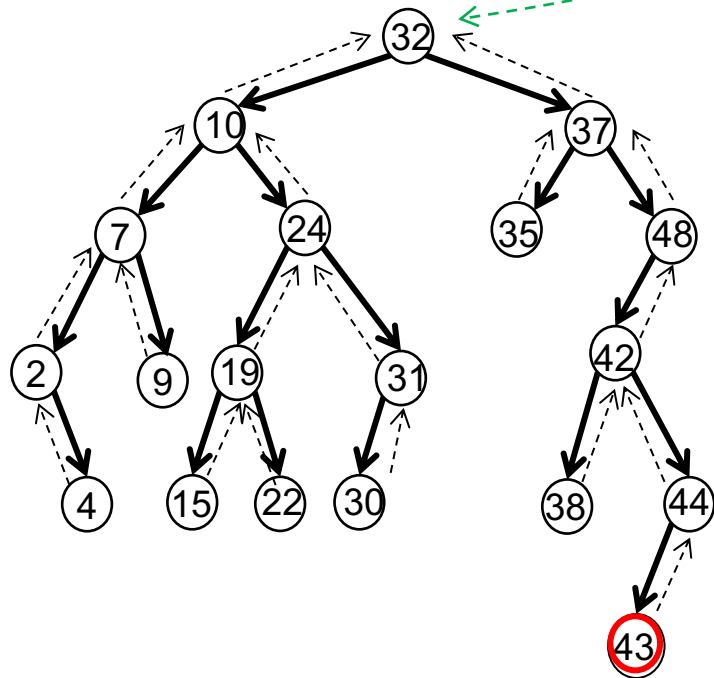


This is now a *doubly-linked* tree ...

# Removing a **leaf** node from a BST

We know the value we want to remove:  $x$

We know the location of this node:



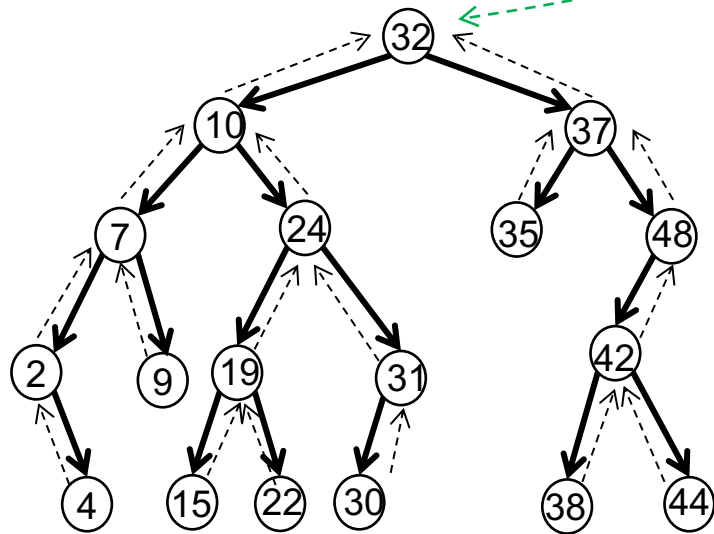
Example: remove 43

We have found the node  
we want to remove:  
and it is a leaf

# Removing a **leaf** node from a BST

*We know the value we want to remove:  $x$*

*We know the location of this node:*



Example: remove 43

Easy:

update the parent's appropriate child  
reference

set the node's parent to None

remember the element

set the node's element to None

return the element

# Removing a **root & leaf** node from a BST

*We know the value we want to remove:  $x$*

*We know the location of this node:*

③2



But be careful if the leaf node is also the root node ...

The root node has no parent.

Easy:

~~update the parent's child reference~~

set the node's parent to None

remember the element

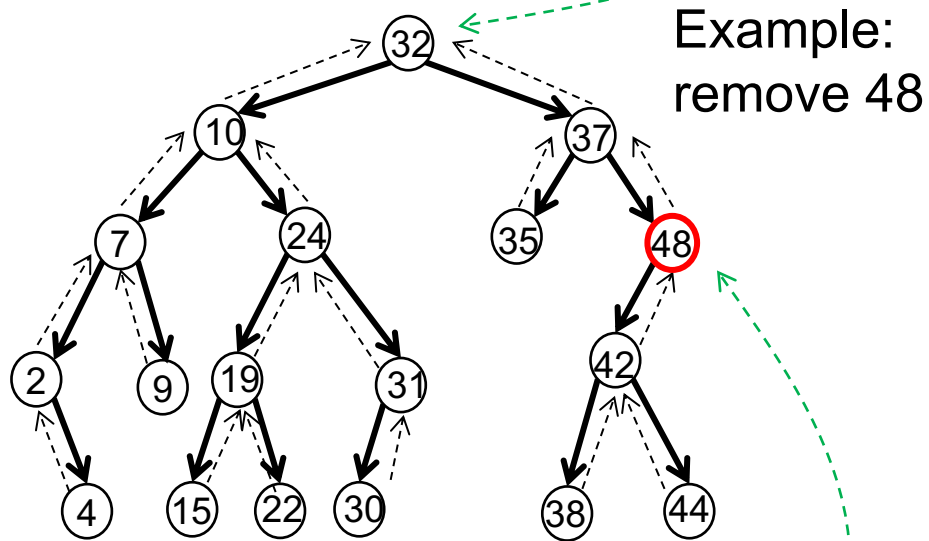
set the node's element to None

return the element

# Removing a **semi-leaf** node from a BST

We know the value we want to remove:  $x$

We know the location of this node:



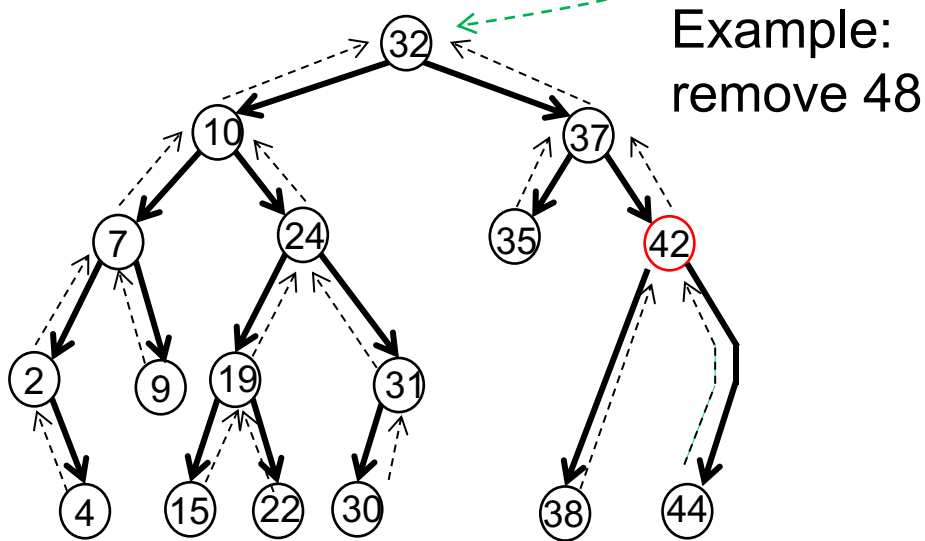
A *semileaf* is a node with only one child

We have found the node we want to remove: and it is a semileaf

# Removing a **semi-leaf** node from a BST

We know the value we want to remove:  $x$

We know the location of this node:



A *semi-leaf* is a node with only one child

Remember the semi-leaf's item  
Copy the semi-leaf's child's item into the semi-leaf.  
Remember the semi-leaf child node  
Rearrange the references to bypass the child node  
Wipe out the child  
Return the original element

Order properties are maintained

- semi-leaf and all descendants were larger than semi-leaf's parent, so all descendants are still larger
- no other relationship has changed



# Next Lecture

Full version of deleting from binary search trees