RISC-V Instruction Set Summary

31:	: 25	24:20	19:15	14:12	11:7	6:0	
fun	funct7		rs1	funct3	rd	ор	R-Type
imm ₁	1:0		rs1	funct3	rd	ор	I-Type
imm ₁	1:5	rs2	rs1	funct3	imm _{4:0}	ор	S-Type
imm ₁	2,10:5	rs2	rs1	funct3	imm _{4:1,11}	op	B-Type
imm ₃	1:12				rd	ор	U-Type
imm ₂	0,10:1,11,1	9:12			rd	ор	J-Type
fs3	fmt	fs2	fs1	funct3	fd	ор	R4-Type
5 bits	2 bits	5 bits	5 bits	3 bits	5 bits	7 bits	

Figure I.1 RISC-V 32-bit instruction formats

XLEN: 32 for RV32 / 64 for RV64
i mm: signed immediate in imm_{11:0}

uimm: 5/6-bit unsigned immediate in imm_{5/4:0}
 upimm: 20 upper bits of a 32-bit immediate, in imm_{31:12}

Address: memory address: rs1 + SignExt(imm_{11:0})
 [Address]: data at memory location Address

 $\begin{array}{ll} \bullet \ \mathsf{BTA:} & \mathsf{branch} \ \mathsf{target} \ \mathsf{address:} \ \mathsf{PC} + \mathsf{SignExt}(\{\mathsf{imm}_{12:1}, \ \mathsf{1'b0}\}) \\ \bullet \ \mathsf{JTA:} & \mathsf{jump} \ \mathsf{target} \ \mathsf{address:} \ \mathsf{PC} + \mathsf{SignExt}(\{\mathsf{imm}_{20:1}, \ \mathsf{1'b0}\}) \\ \end{array}$

label: text indicating instruction address
 SignExt: value sign-extended to XLEN
 ZeroExt: value zero-extended to XLEN
 csr: control and status register

Table I.1 RV32/64I: RISC-V integer instructions

ор	funct3/ imm _{14:12}	funct7/	Туре	Instruct	ion		Description	Operation
0000011	000	imm _{11:5}	I		rd,	imm(rs1)	load byte	rd = SignExt([Address] _{7:0})
0000011	001	imm _{11:5}	ī		rd,	imm(rs1)	load half	rd = SignExt([Address] _{15:0})
0000011	010	imm _{11:5}	ı	1w	rd,	imm(rs1)	load word	rd = SignExt([Address] _{31:0})
0000011	100	imm _{11:5}	ī	1bu	rd,	imm(rs1)	load byte unsigned	rd = ZeroExt([Address] _{7:0})
0000011	101	imm _{11:5}	Ī	1hu	rd,	imm(rs1)	load half unsigned	rd = ZeroExt([Address] _{15:0})
0010011	000	imm _{11:5}	ı	addi	rd,	rs1, imm	add immediate	rd = rs1 + SignExt(imm)
0010011	001	000000*	1	slli	rd,	rs1, uimm	shift left logical immediate	rd = rs1 << uimm
0010011	010	imm _{11:5}	I	slti	rd,	rs1, imm	set less than immediate	rd = (rs1 < SignExt(imm))
0010011	011	imm _{11:5}	I	sltiu	rd,	rs1, imm	set less than imm. unsigned	rd = (rs1 < SignExt(imm))
0010011	100	imm _{11:5}	I	xori	rd,	rs1, imm	xor immediate	rd = rs1 ^ SignExt(imm)
0010011	101	000000*	I	srli	rd,	rs1, uimm	shift right logical immediate	rd = rs1 >> uimm
0010011	101	010000*	I	srai	rd,	rs1, uimm	shift right arithmetic imm.	rd = rs1 >>> uimm
0010011	110	imm _{11:5}	I	ori	rd,	rs1, imm	or immediate	rd = rs1 SignExt(imm)
0010011	111	imm _{11:5}	I	andi	rd,	rs1, imm	and immediate	rd = rs1 & SignExt(imm)
0010111	imm _{14:12}	imm _{31:25}	U	auipc	rd,	upimm	add upper immediate to PC	rd = {upimm, 12'b0} + PC
0100011	000	imm _{11:5}	S	sb	rs2,	imm(rs1)	store byte	$[Address]_{7:0} = rs2_{7:0}$
0100011	001	imm _{11:5}	S	sh	rs2,	imm(rs1)	store half	[Address] _{15:0} = $rs2_{15:0}$
0100011	010	imm _{11:5}	S	SW	rs2,	imm(rs1)	store word	[Address] _{31:0} = rs2 _{31:0}
0110011	000	0000000	R	add	rd,	rs1, rs2	add	rd = rs1 + rs2
0110011	000	0100000	R	sub	rd,	rs1, rs2	sub	rd = rs1 - rs2
0110011	001	0000000	R	sll	rd,	rs1, rs2	shift left logical	$rd = rs1 << rs2_{5/4:0} (RV64/32)$
0110011	010	0000000	R	slt	rd,	rs1, rs2	set less than	rd = (rs1 < rs2)
0110011	011	0000000	R	sltu	rd,	rs1, rs2	set less than unsigned	rd = (rs1 < rs2)
0110011	100	0000000	R	xor	rd,	rs1, rs2	xor	rd = rs1 ^ rs2
0110011	101	0000000	R	srl	rd,	rs1, rs2	shift right logical	$rd = rs1 >> rs2_{5/4:0} (RV64/32)$
0110011	101	0100000	R	sra	rd,	rs1, rs2	shift right arithmetic	$rd = rs1 >>> rs2_{5/4:0} (RV64/32)$
0110011	110	0000000	R	or	rd,	rs1, rs2	or	rd = rs1 rs2
0110011	111	0000000	R	and	rd,	rs1, rs2	and	rd = rs1 & rs2
0110111	imm _{14:12}	imm _{31:25}	U	lui	rd,	upimm	load upper immediate	rd = {upimm, 12'b0}
1100011	000	imm _{12,10:5}	В	beq	rs1,	rs2, label	branch if =	if (rs1 == rs2) PC = BTA
1100011	001	imm _{12,10:5}	В	bne	rs1,	rs2, label	branch if ≠	if (rs1 ≠ rs2) PC = BTA
1100011	100	imm _{12,10:5}	В	blt	rs1,	rs2, label	branch if <	if (rs1 < rs2) PC = BTA
1100011	101	imm _{12,10:5}	В	bge	rs1,	rs2, label	branch if ≥	if (rs1 ≥ rs2) PC = BTA
1100011	110	imm _{12,10:5}	В			rs2, label	branch if < unsigned	if (rs1 < rs2) PC = BTA
1100011	111	imm _{12,10:5}	В	bgeu	rs1,	rs2, label	branch if ≥ unsigned	if (rs1 ≥ rs2) PC = BTA
1100111	000	imm _{11:5}	I	jalr	rd,	rs1, imm	jump and link register	PC = rs1 + SignExt(imm), rd = PC + 4
1101111	imm _{14:12}	imm _{20,10:5}	J	jal	rd,	label	jump and link	PC = JTA, $rd = PC + 4$

^{* = 0} for RV32, uim m_5 for RV64.

Table I.2 RV64I: Extra integer instructions

ор	funct3	funct7/ imm _{11:5}	Туре	Instruction	Description	Operation
0000011	011	imm _{11:5}	I	ld rd, imm(rs1)	load double word	rd=[Address] _{63:0}
0000011	110	imm _{11:5}	I	lwu rd, imm(rs1)	load word unsigned	rd=ZeroExt([Address] _{31:0})
0011011	000	imm _{11:5}	I	addiw rd, rs1, imm	add immediate word	rd=SignExt((rs1+SignExt(imm)) _{31:0})
0011011	001	0000000	I	slliw rd, rs1, uimm	shift left logical immediate word	rd=SignExt((rs1 _{31:0} << uimm) _{31:0})
0011011	101	0000000	I	srliw rd, rs1, uimm	shift right logical immediate word	rd=SignExt((rs1 _{31:0} >> uimm) _{31:0})
0011011	101	0100000	I	sraiw rd, rs1, uimm	shift right arith. immediate word	rd=SignExt((rs1 _{31:0} >>> uimm) _{31:0})
0100011	011	imm _{11:5}	S	sd rs2, imm(rs1)	store double word	[Address] _{63:0} =rs2
0111011	000	0000000	R	addw rd, rs1, rs2	add word	rd=SignExt((rs1+rs2) _{31:0})
0111011	000	0100000	R	subw rd, rs1, rs2	subtract word	rd=SignExt((rs1-rs2) _{31:0})
0111011	001	0000000	R	sllw rd, rs1, rs2	shift left logical word	rd=SignExt((rs1 _{31:0} << rs2 _{4:0}) _{31:0})
0111011	101	0000000	R	srlw rd, rs1, rs2	shift right logical word	$rd = SignExt((rs1_{31:0} >> rs2_{4:0})_{31:0})$
0111011	101	0100000	R	sraw rd, rs1, rs2	shift right arithmetic word	rd=SignExt((rs1 _{31:0} >>>rs2 _{4:0}) _{31:0})

In RV64l, registers are 64 bits, but instructions are still 32 bits. The term "word" generally refers to a 32-bit value. In RV64l, immediate shift instructions use 6-bit immediates: uimm $_{50}$; but for word shifts, the most significant bit of the shift amount (uimm $_{5}$) must be 0. Instructions ending in "w" (for "word") operate on the lower half of the 64-bit registers. Sign- or zero-extension produces a 64-bit result.

Table I.3 RVF/D/Q/Zfh: RISC-V single-, double-, quad-, and half-precision floating-point instructions

						, adable, quad, and han pre	l l	
ор	funct3	funct7/ instr _{31:25}	rs2/ imm _{4:0}	Type	Flags NV DZ OF UF NX	Instruction	Description	Operation
1000011			fs2	R4	x - x x x		multiply-add	fd = fs1 * fs2 + fs3
1000111			fs2	R4	x - x x x		multiply-subtract	fd = fs1 * fs2 - fs3
1001011	_	fs3, fmt	fs2	R4	x - x x x	fnmsub.* fd,fs1,fs2,fs3	negate multiply-sub	fd = -(fs1 * fs2 - fs3)
1001111		fs3, fmt	fs2	R4	x - x x x	fnmadd.* fd,fs1,fs2,fs3	negate multiply-add	fd = -(fs1 * fs2 + fs3)
1010011	rm	00000, fmt	fs2	R	x - x - x	fadd.* fd,fs1,fs2	add	fd = fs1 + fs2
	rm	00001, fmt	fs2	R	x - x - x	fsub.* fd,fs1,fs2	subtract	fd = fs1 - fs2
1010011	rm	00010, fmt	fs2	R	x - x x x	fmul.* fd,fs1,fs2	multiply	fd = fs1 * fs2
1010011	rm		fs2	R	xxxxx	fdiv.* fd,fs1,fs2	divide	fd = fs1 / fs2
1010011	rm	01011, fmt	00000	ı	x x	fsqrt.* fd,fs1	square root	fd = sqrt(fs1)
1010011	000	00100, fmt	fs2	R		fsgnj.* fd,fs1,fs2	sign injection	fd = fs1, sign = sign(fs2)
1010011	001	00100, fmt	fs2	R		fsgnjn.* fd,fs1,fs2	negate sign injection	fd = fs1, sign = -sign(fs2)
1010011	010	00100, fmt	fs2	R		fsgnjx.* fd,fs1,fs2	xor sign injection	fd = fs1,
		·					<i>3</i> ,	sign = sign(fs2)^sign(fs1)
1010011	000	00101, fmt	fs2	R	x	fmin.* fd,fs1,fs2	min	fd = min(fs1, fs2)
1010011		,	fs2	R	x	fmax.* fd,fs1,fs2	max	fd = max(fs1, fs2)
1010011		,	fs2	R	x	feq.* rd,fs1,fs2	compare =	rd = (fs1 == fs2)
1010011	001	10100, fmt	fs2	R	x	flt.* rd,fs1,fs2	compare <	rd = (fs1 < fs2)
1010011	000	,	fs2	R	x	fle.* rd,fs1,fs2	compare ≤	rd = (fs1 ≤ fs2)
1010011	001	11100, fmt	00000	I		fclass.* rd,fs1	classify	rd = classification of fs1
0000111	010	imm _{11:5}	imm _{4:0}	I		flw fd, imm(rs1)	load float	$fd = [Address]_{31:0}$
0100111	010	imm _{11:5}	fs2	S		fsw fs2,imm(rs1)	store float	$[Address]_{31:0} = fd$
1010011	rm	,	000,lu	I	x - x - x		convert fp to int	rd = dst(fs1)
1010011	rm	,	000,lu	I	x - x		convert int to fp	fd = dst(rs1)
1010011	rm	01000, dfmt		I	x - x x x		convert fp to fp	fd = dst(fs1)
1010011	000	11100, fmt	00000	I		$fmv.x.\{h/w/d\}$ rd, fs1	move fp to int	rd = fs1
1010011	000	11110, fmt	00000	I		$fmv.{h/w/d}.x$ fd, rs1	move int to fp	fd = rs1
						RVD only		
0000111	011	imm _{11:5}	imm _{4:0}	I		fld fd, imm(rs1)	load double	fd = [Address] _{63:0}
0100111	011	imm _{11:5}	fs2	S		fsd fs2,imm(rs1)	store double	$[Address]_{63:0} = fs2$
						RVQ only		
0000111	100	imm _{11:5}	imm _{4:0}	I		flq fd, imm(rs1)	load quad	$fd = [Address]_{127:0}$
0100111	100	imm _{11:5}	fs2	S		fsq fs2,imm(rs1)	store quad	$[Address]_{127:0} = fs2$
						RVH only		
0000111	001	imm _{11:5}	imm _{4:0}	I		flh fd, imm(rs1)	load half	fd = [Address] _{15:0}
0100111	001	imm _{11:5}	fs2	S		fsh fs2,imm(rs1)	store half	$[Address]_{15:0} = fs2$

^{* = {}h/s/d/q}. fs1, fs2, fs3, fd: floating-point registers. fs1, fs2, and fd are encoded in fields rs1, rs2, and rd; only R4-type also encodes fs3. fmt: precision of computational instruction (half=10₂, single=20₂, double=21₂, quad=11₂). sfmt: source format. rm: rounding mode (0=to nearest, 1=toward zero, 2=down, 3=up, 4=to nearest (max magnitude), 7=dynamic). sign(fs1): the sign of fs1. dst or src = {h/s/d/q} for fp, {w/wu/l/lu} for 32/64 bit signed/unsigned ints. {w/wu} = 32-bit signed/unsigned int. {l/ ul} = 64-bit (long) signed/unsigned int. For fp move instructions (fmv), {h/w/d} denotes a half-/single-/double-precision fp number held in a fp register, and x is a fp number (held in an integer register), with the same-precision as the fp source/destination. Floating-point exception flags are: NV: invalid operation, DZ: divide by zero, OF: overflow, UF: underflow, and NX: inexact.

Table I.4 Register names and numbers

Table II Register names and name is							
Name	Register Number Use						
zero	x0	Constant value 0					
ra	x1	Return address					
sp	x2	Stack pointer					
gp	x3	Global pointer					
tp	x4	Thread pointer					
t0-2	x5-7	Temporary registers					
s0/fp	x8	Saved register / Frame pointer					
s1	x9	Saved register					
a0-1	x10-11	Function arguments / Return values					
a2-7	x12-17	Function arguments					
s2-11	x18-27	Saved registers					
t3-6	x28-31	Temporary registers					

15 14 13	12	11 10	9 8 7	6 5	4 3 2	1 0	
funct4		rd	/rs1	rs2		ор	CR
funct3	imm	rd	/rs1	imm		op	CI
funct3	imm		rs1'	imm	rs2'	op	CS
funct6			rd', rs1'	funct2	rs2'	op	CA
funct6			rs1'	imm	rd' / rs2'	op	CLB / CSB
funct6			rs1'	f1 imm	rd' / rs2'	op	CLH / CSH
funct6			rd', rs1'	funct5		op	CU
funct3	imm		rs1'	imm		op	CB
funct3	imm	funct	rd', rs1'	imm		ор	CB'
funct3	imm		•			ор	CJ
funct3	imm			rs2		op	CSS
funct3	imm				rd'	op	CIW
funct3	imm		rs1'	imm	rd'	op	CL
3 bits	3 bits		3 bits	2 bits	3 bits	2 bits	•

Figure I.2. RISC-V compressed (16-bit) instruction formats

Table I.5 RVM: RISC-V multiply and divide instructions

ор	funct3	funct7	Туре	Instruction	Description	Operation
0110011	000	0000001	R	mul rd, rs1, rs2	multiply	$rd = (rs1 * rs2)_{XLEN-1:0}$
0110011	001	0000001	R	mulh rd, rs1, rs2	multiply high signed signed	$rd = (rs1 * rs2)_{2*XLEN-1:XLEN}$
0110011	010	0000001	R	mulhsu rd, rs1, rs2	multiply high signed unsigned	$rd = (rs1 * rs2)_{2*XLEN-1:XLEN}$
0110011	011	0000001	R	mulhu rd, rs1, rs2	multiply high unsigned unsigned	$rd = (rs1 * rs2)_{2*XLEN-1:XLEN}$
0110011	100	0000001	R	div rd, rs1, rs2	divide (signed)	rd = rs1 / rs2
0110011	101	0000001	R	divu rd, rs1, rs2	divide unsigned	rd = rs1 / rs2
0110011	110	0000001	R	rem rd, rs1, rs2	remainder (signed)	rd = rs1 % rs2
0110011	111	0000001	R	remu rd, rs1, rs2	remainder unsigned	rd = rs1 % rs2

RV64M adds mulw, divw, remw, divuw, remuw with op = 59. These operate on only the lower 32 bits of a register.

Table I.6 RVC/Zca: RISC-V compressed (16-bit) integer instructions

				Table I.	6 RVC/ZCa	a: KISC-V C	ompressea (1	6-bit) integer instructi	ions		
ор	instr _{15:10}	funct2	Registers	Type	Compress	ed Instru	ction		32-Bit I	Equivale	ent
	C/Zca: Compressed instructions excluding floating-point loads and stores										
00	000	-	imm ≠ 0	CIW	c.addi4s	pn rd',	sp, imm	-	addi	rd',	sp, ZeroExt(imm)*4
00	010	-		CL	c.lw	rd',	imm(rs1')		1w	rd',	(ZeroExt(imm)*4)(rs1')
	110	-		CS		rs2',	imm(rs1')		SW		(ZeroExt(imm)*4)(rs1')
	000000	-	rs1 = 0, $imm = 0$	CI	c.nop				addi	хO,	x0, 0
	000	-	rd ≠ 0, imm ≠ 0	CI		rd,	imm		addi	rd,	rd, SignExt(imm)
01	010	-	rd ≠ 0	CI		rd,	imm		addi	rd,	xO, SignExt(imm)
01	011	-	rd \neq {0,2}, imm \neq 0	CI	c.lui	rd,	imm		lui	rd,	{14{imm ₅ }, imm}
01	011	-	rd = 2, imm ≠ 0	CI	c.addi16	1 1 1	imm		addi	sp,	sp, SignExt(imm)*16
01	100-00	-		CB'		rd',	imm		srli	rd',	rd', imm
01	100-01	-		CB'		rd',	imm		srai	rd',	rd', imm
01	100-10	-		CB'		rd',	imm		andi	rd',	rd', SignExt(imm)
01	100011	00		CA		rd',	rs2'		sub	rd',	rd', rs2'
01	100011	01		CA		rd',	rs2'		xor	rd',	rd', rs2'
01	100011	10		CA		rd',	rs2'		or	rd',	rd', rs2'
01	100011	11		CA		rd',	rs2'		and	rd',	rd', rs2'
01	101	-		CJ	c.j	label			jal	хO,	label
01	110	-		CB	c.beqz	rsl',	label		beq		
01	111	-		CB	c.bnez	rsl',	label		bne	rs1',	x0, label
10	000	-	rd ≠ 0	CI	c.slli	rd,	imm		slli	rd,	rd, imm
10	010	-	rd ≠ 0	CI	c.lwsp	rd,	imm		1 w	rd,	(ZeroExt(imm)*4)(sp)
10	1000	-	rs1 ≠ 0, rs2 = 0	CR	c.jr	rs1			jalr	х0,	rs1, 0
10	1000	-	rd ≠ 0, rs2 ≠ 0	CR	c.mv	rd,	rs2		add	rd,	x0, rs2
10	1001	-	rs1 = 0, rs2 = 0	CR	c.ebreak	(ebreal	k	
10	1001	-	rs1 ≠ 0, rs2 = 0	CR	c.jalr	rs1			jalr	ra,	rs1, 0
10	1001	-	rd ≠ 0, rs2 ≠ 0	CR	c.add	rd,	rs2		add	rd,	rd, rs2
10	110	_		CSS	c.swsp	rs2,	imm		SW	rs2,	(ZeroExt(imm)*4)(sp)
							Zca (RV32 only	')			
01	001	-		CJ	c.jal	label			jal	ra,	label
			C/Zca (RV64 only,	reuses	encodings	for c.jal a	bove and c.flv	, c.fsw, c.flwsp, and c.	fswsp fr		
00	011	-		CL	c.ld	rd',	imm(rs1')	(replaces c.flw)	1 d	rd',	(ZeroExt(imm)*8)(rs1')
00	111	-		CS	c.sd	rs2',	imm(rs1')	(replaces c.fsw)	sd	rs2,	(ZeroExt(imm)*8)(rs1')
01	001	-	rd ≠ 0	CI	c.addiw	rd,	imm	(replaces c.jal)	addiw		rd, (SignExt(imm))
01	100111	00		CA	c.subw	rd',	rs2'		subw	rd',	rd', rs2'
01	100111	01		CA	c.addw	rd',	rs2'		addw	rd',	rd', rs2'
10	011		rd ≠ 0	CI	c.ldsp	rd,	i mm	(replaces c.flwsp)	1d	rd,	(ZeroExt(imm)*8)(sp)
10	111	-		CSS	c.sdsp	rs2,	i mm	(replaces c.fswsp)		rs2,	(ZeroExt(imm)*8)(sp)

Table I.7 Common RISC-V pseudoinstructions

Pseudoinstruction	RISC-V Instructions	Description	Operation
nop	addi x0, x0, 0	no operation	
li rd, $imm_{11:0}$	addi rd, x0, imm _{11:0}	load 12-bit immediate	rd = SignExtend(imm _{11:0})
li rd, imm _{31:0}	lui rd, imm _{31:12} *	load 32-bit immediate	$rd = imm_{31:0}$
	addi rd, rd, imm _{11:0}		
mv rd, rs1	addi rd, rs1, 0	move (also called "register copy")	rd = rs1
not rd, rs1	xori rd, rs1, -1	one's complement	rd = ~rs1
neg rd, rs2	sub rd, x0, rs2	two's complement	rd = -rs1
seqz rd, rs1	sltiu rd, rs1, 1	set if = 0	rd = (rs1 == 0)
snez rd, rs2	sltu rd, x0, rs2	set if ≠ 0	$rd = (rs1 \neq 0)$
sltz rd, rsl	slt rd, rs1, x0	set if < 0	rd = (rs1 < 0)
sgtz rd, rs2	slt rd, x0, rs2	set if > 0	rd = (rs1 > 0)
beqz rs1, label	beq rs1, x0, label	branch if = 0	if (rs1 == 0) PC = label
bnez rs1, label	bne rs1, x0, label	branch if ≠ 0	if (rs1 ≠ 0) PC = label
blez rs2, label	bge x0, rs2, label	branch if ≤ 0	if (rs1 ≤ 0) PC = label
bgez rs1, label	bge rs1, x0, label	branch if ≥ 0	if (rs1 ≥ 0) PC = label
bltz rs1, label	blt rs1, x0, label	branch if < 0	if (rs1 < 0) PC = label
bgtz rs2, label	blt x0, rs2, label	branch if > 0	if (rs1 > 0) PC = label
ble rs1, rs2, label	bge rs2, rs1, label	branch if ≤	if (rs1 ≤ rs2) PC = label
bgt rs1, rs2, label	blt rs2, rs1, label	branch if >	if (rs1 > rs2) PC = label
bleu rs1, rs2, label	bgeu rs2, rs1, label	branch if ≤ (unsigned)	if (rs1 ≤ rs2) PC = label
bgtu rs1, rs2, label	bltu rs2, rs1, offset	branch if > (unsigned)	if (rs1 > rs2) PC = label
j label	jal x0, label	jump	PC = label
jal label	jal ra, label	jump and link	PC = label, $ra = PC + 4$
jr rs1	jalr x0, rs1, 0	jump register	PC = rs1
jalr rs1	jalr ra, rs1, 0	jump and link register	PC = rs1, $ra = PC + 4$
ret	jalr x0, ra, 0	return from function	PC = ra
call label	jal ra, label	call nearby function	PC = label, ra = PC + 4
call label	auipc ra, offset _{31:12} *	call far away function	PC = PC + offset, ra = PC + 4
	jalr ra, ra, offset _{11:0}		
la rd, symbol	auipc rd, symbol _{31:12} *	load address of global variable	rd = PC + symbol
	addi rd, rd, symbol _{11:0}		
	auipc rd, symbol $_{31:12}^*$ l{b h w} rd, symbol $_{11:0}$ (rd)	load global variable	rd = [PC + symbol]
s{b h w} rs2, symbol, rs1	auipc rs1, symbol _{31:12} *	store global variable	[PC + symbol] = rs2
cenn nd cen	$s\{b h w\}$ rs2, symbol _{11:0} (rs1)	road CCD	rd - ccr
csrr rd, csr	csrrs rd, csr, x0	read CSR	rd = csr
csrw csr, rs1	csrrw x0, csr, rs1	write CSR	csr = rs1
csrs/csrc csr, rs1	csrrs/csrrc x0, csr, rs1	set/clear bits in CSR	$ csr = csr rs1 / csr = csr \& \sim rs1$

^{*} If bit 11 of the immediate/offset/symbol is 1, the upper immediate is incremented by 1. offset/symbol are the 32-bit PC-relative addresses of a label/global variable.

Table I.8 Privileged / CSR / fence instructions

ор	funct3	imm _{11:0} /funct7	Registers	Туре	Instruction	Description	Operation
1110011	000	00000000000	rs1,rd=0	I	ecall	transfer control to OS	
1110011	000	000000000001	rs1,rd=0	I	ebreak	transfer control to debugger	
1110011	000	000100000010	rs1,rd=0	I	sret	return from supervisor exception	PC = sepc
1110011	000	001100000010	rs1,rd=0	I	mret	return from machine exception	PC = mepc
1110011	000	000100000101	rs1,rd=0	I	wfi	wait for interrupt	
0001111	000	000000110011	rs1,rd=0	I	fence	synchronize loads and stores	
0001111	001	000000000000	rs1,rd=0	I	fence.i	synchronize instruction memory	
1110011	000	0001001	rs1,rs2,rd=0	R	sfence.vma	synchronize page table	
1110011	001	csr		I	csrrw rd,csr,rs1	CSR read/write	rd=csr,csr = rs1
1110011	010	csr		I	csrrs rd,csr,rs1	CSR read/set	rd=csr,csr = rs1
1110011	011	csr		I	csrrc rd,csr,rs1	CSR read/clear	rd = csr,csr &= ~rs1
1110011	101	csr	rs1=uimm	I	csrrwi rd,csr,uimm	CSR read/write immediate	rd=csr,csr = ZeroExt(uimm)
1110011	110	csr	rs1=uimm	I	csrrsi rd,csr,uimm	CSR read/set immediate	rd=csr,csr = ZeroExt(uimm)
1110011	111	csr	rs1=uimm	I	csrrci rd,csr,uimm	CSR read/clear immediate	rd = csr,csr &= ~ZeroExt(uimm)