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University Of Portsmouth  
BSc (Hons) Computer Science  
First Year

**Architecture and Operating Systems - Computer**

M30943

September 2022 - May 2023

20 Credits

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# S.1. INTRODUCTION TO MODULE

📅 26-09-22

🕒 16:00

🎓 Farzad

📍 RB LT1

## Division of the Module

This module is split into two parts: computer (this part) which is worth 70% and maths (the other part) which is worth 30%. The two parts are run completely independently of each other. The only time they come together is when the final overall score is calculated.

There are two separate Moodle pages (one for Computer and one for Maths)

## Computer Module assessments

For the Computer section of the module, there are two assessments. One is in January 2023, which will be a Computer Based Test (covering content taught in the first teaching block). It is worth 30% of the over module score. The second is in the May/June 2023 assessment period. It will be computer based. This assessment will be worth 40% of the overall module score.

Both assessments are closed book however a formula sheet will be provided for the January assessment. Nothing is provided for the May/June assessment.

The pass mark for the entire module is 40%, this score is generated from all the computer assessments AND all the maths assessments.

## Module structure

There will be a one hour lecture per week, where content is introduced to us. This will be delivered using worksheets for the first 10 weeks.

There will also a practical session each week where the cohort is split into smaller groups. These sessions will be a chance to practice the ideas introduced in the lectures. There will be more members of staff around at the practical sessions to help out.

### More Information on Practical Sessions

There are practical session guidelines available in the induction slides or on Moodle.

### Content in each Week

There is a teaching plan on Moodle which outlines the content covered each week as well as the weeks in which the exams will be held.

## S.2. BINARY ARITHMETIC

📅 26-09-2022

🕒 16:15

🎓 Farzad

📍 RB LT1

### Number Systems

There are a number of different number systems and different methods to convert between them.

#### Denary (Base 10)

Used most commonly, this is the one most people learn.

$10^x$	$10^3$	$10^2$	$10^1$	$10^0$
$10^x =$	1000	100	10	1
	4	2	5	1

The total of the numbers above would be calculated in the following way:

$$4251 = (1000 \times 4) + (100 \times 2) + (10 \times 5) + (1 \times 1)$$

Denary is also known as base 10, this means each column can have one of ten possible values (0, 1, 2, 3, 4, 5, 6, 7, 8, 9)

#### Binary (Base 2)

This is base 2, this means each column can have one of two possible values (0, 1). The columns are also different. Moving from right to left, the columns double each time.

$2^x$	$2^7$	$2^6$	$2^5$	$2^4$	$2^3$	$2^2$	$2^1$	$2^0$
$2^x =$	128	64	32	16	8	4	2	1
	1	0	1	1	0	0	1	1

The largest value which can be stored in 8-bits of binary is  $11111111_2$  or  $255_{10}$ .

#### Hexadecimal (Base 16)

Also known as Hex. Using this method, numbers up to 255 can be stored in two characters. This is used a lot in computing, especially in graphics and website development. Each column can have one of 16 values (1 2 3 4 5 6 7 8 9 A B C D E F). The letters are used to represent two-digit numbers as seen below.

Hex:	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
Decimal:	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

To calculate the value held in a Hex number, we calculate in a similar way to Denary and Binary as seen below.

$16^x$	$16^3$	$16^2$	$16^1$	$16^0$
$16^x =$	4096	256	16	1
	D	3	C	E

$$D3CE = (13 \times 4096) + (3 \times 256) + (12 \times 16) + (14 \times 1) = 54222$$

## Binary To Denary

## Denary To Binary

## Denary to Hex

## Binary Addition

## Basic Rules

$$\begin{array}{rcl} 0 + 0 & = & 0 \\ 0 + 1 & = & 1 \\ 1 + 1 & = & 1 \\ 1 + 1 & = & 10 \end{array}$$

## Binary Addition Example

1. Draw out the binary addition columns

$$+ \begin{array}{ccc} 1 & 1 & 0 \\ 0 & 1 & 1 \end{array}$$

$$\begin{array}{r} 1 \ 1 \ 0 \\ + \ 0 \ 1 \ 1 \\ \hline 1 \end{array}$$
$$\begin{array}{r} \phantom{+} \phantom{0} \phantom{0} \phantom{0} \\ \phantom{+} \phantom{0} \phantom{0} \phantom{0} \\ + \phantom{0} \phantom{0} \phantom{0} \phantom{0} \\ \hline \phantom{+} \phantom{0} \phantom{0} \phantom{0} \end{array}$$

- $$\begin{array}{r} 1110 \\ + 0111 \\ \hline 0011 \end{array}$$

- $$\begin{array}{r} 1110 \\ + \quad 011 \\ \hline 1001 \end{array}$$

	1	1
x	1	0

$$\begin{array}{cc} & 1 & 1 \\ x & 1 & 0 \\ \hline & 0 & 0 \end{array}$$

	1	1
x	1	0
	0	0
1	1	

		1	1
x		1	0
		0	0
+	1	1	
	1	1	0

$$\begin{array}{rrrr} & 1 & 1 & 0 \\ - & 0 & 0 & 1 \end{array}$$

$$\begin{array}{r} 101 \\ - 00 \\ \hline 101 \end{array}$$

$$\begin{array}{r} 101 \\ - 001 \\ \hline 100 \end{array}$$

$$\begin{array}{rrrr} & 1 & 0 & 1 \\ - & 0 & 0 & 1 \\ \hline & 1 & 0 & 1 \end{array}$$

This gives us the final answer of  $110 - 001 = 101$

## Binary Division

Binary division follows much the same procedure as ‘bus stop’ decimal division.

### Binary Division Example

Divide  $110 \div 10$

1. Draw out the division columns as you would for a standard decimal ‘bus stop’ division.

$$\begin{array}{r} 10 \overline{) 110} \end{array}$$

2. Start by looking for factors and find 11 is greater than 10. We then write the number of times the value goes into 11 at the top, and the value itself underneath.

$$\begin{array}{r} 1 \\ 10 \overline{) 110} \\ \underline{10} \phantom{0} \end{array}$$

3. We then subtract to see if there is a remainder. ( $11 - 10 = 01$ )  
The remainder is written up on the top line

$$\begin{array}{r} 11 \\ 10 \overline{) 110} \\ \underline{- 10} \phantom{0} \\ 01 \phantom{0} \end{array}$$

4. We then bring down the final digit in the division (0), to where we are working.

$$\begin{array}{r} 11 \\ 10 \overline{) 110} \\ \underline{- 10} \phantom{0} \\ 010 \phantom{0} \end{array}$$

5. Now, we look to see if our divisor can fit in again. It does fit again, so we subtract it. ( $010 - 10 = 0$ ) It leaves no remainder.

$$\begin{array}{r} 11 \\ 10 \overline{) 110} \\ \underline{- 10} \phantom{0} \\ 010 \phantom{0} \\ \underline{- 10} \phantom{0} \\ 0 \phantom{0} \end{array}$$

This gives us the answer of  $110 \div 10 = 11$ .



# WORKSHEET 1

📅 20-09-22

👁 Worksheet

## Basic Exercises

1. Convert the following numbers to binary

(a)  $12 = 1100$

(b)  $103 = 1100111$

(c)  $97 = 1100001$

(d)  $55 = 0110111$

(e)  $395 = 110001011$

2. Convert the following binary numbers to decimal

(a)  $1101 = 13$

(b)  $101001 = 41$

(c)  $110111 = 55$

(d)  $1000011 = 135$

(e)  $11111110 = 254$

3. Convert the following decimal numbers to hexadecimal numbers

(a) 1026

0100	0000	0010
4	0	2

$= 402$

(b) 5678

0001	0110	0010	1110
1	6	2	E

$= 162E$

(c) 9567

0010	0101	0101	1111
2	5	5	F

$= 255E$

(d) 72627

0001	0001	1011	1011	0011
1	1	B	B	3

$= 11BB3$

(e) 115497

0001	1100	0011	0010	1001
1	C	3	2	9

 $= 1C329$ 

4. Convert the following hexadecimal numbers to binary numbers

(a)  $2D = 0010\ 1101$ (b)  $F3A = 1111\ 0011\ 1010$ (c)  $1BD = 0001\ 1011\ 1101$ (d)  $ABC = 1010\ 1011\ 1100$ (e)  $D3F2 = 1101\ 0011\ 1111\ 0010$ **Core Exercises**a)  $11 + 11$ 

$$\begin{array}{r}
 \phantom{+} \phantom{1} \phantom{1} \phantom{0} \phantom{0} \\
 \phantom{+} \phantom{1} \phantom{1} \phantom{0} \phantom{0} \\
 + \phantom{1} \phantom{1} \phantom{0} \phantom{0} \phantom{0} \\
 \hline
 1 \phantom{0} 1 \phantom{0} \phantom{0} \phantom{0} \\
 \hline
 1 \phantom{0} 1
 \end{array}$$

 $= 110$ b)  $100 + 10$ 

$$\begin{array}{r}
 \phantom{+} \phantom{1} \phantom{0} \phantom{0} \phantom{0} \\
 \phantom{+} \phantom{1} \phantom{0} \phantom{0} \phantom{0} \\
 + \phantom{1} \phantom{0} \phantom{0} \phantom{0} \phantom{0} \\
 \hline
 1 \phantom{0} 1 \phantom{0} \phantom{0} \phantom{0} \\
 \hline
 1 \phantom{0} 1 \phantom{0} \phantom{0} \phantom{0}
 \end{array}$$

 $= 110$ c)  $111 + 11$ 

$$\begin{array}{r}
 \phantom{+} \phantom{1} \phantom{1} \phantom{0} \phantom{0} \phantom{0} \\
 \phantom{+} \phantom{1} \phantom{1} \phantom{0} \phantom{0} \phantom{0} \\
 + \phantom{1} \phantom{1} \phantom{0} \phantom{0} \phantom{0} \phantom{0} \\
 \hline
 1 \phantom{0} 1 \phantom{0} \phantom{0} \phantom{0} \phantom{0} \\
 \hline
 1 \phantom{0} 1 \phantom{0} \phantom{0} \phantom{0} \phantom{0}
 \end{array}$$

 $= 1010$ d)  $110 + 100$ 

$$\begin{array}{r}
 \phantom{+} \phantom{1} \phantom{1} \phantom{0} \phantom{0} \phantom{0} \\
 \phantom{+} \phantom{1} \phantom{1} \phantom{0} \phantom{0} \phantom{0} \\
 + \phantom{1} \phantom{1} \phantom{0} \phantom{0} \phantom{0} \phantom{0} \\
 \hline
 1 \phantom{0} 1 \phantom{0} \phantom{0} \phantom{0} \phantom{0} \\
 \hline
 1 \phantom{0} 1 \phantom{0} \phantom{0} \phantom{0} \phantom{0}
 \end{array}$$

 $= 1010$ 

e)

$$\begin{array}{r}
 \phantom{+} \phantom{1} \phantom{1} \phantom{1} \phantom{1} \phantom{0} \phantom{0} \\
 \phantom{+} \phantom{1} \phantom{1} \phantom{1} \phantom{1} \phantom{0} \phantom{0} \\
 + \phantom{1} \phantom{1} \phantom{1} \phantom{1} \phantom{0} \phantom{0} \phantom{0} \\
 \hline
 1 \phantom{0} 1 \phantom{0} \phantom{0} \phantom{0} \phantom{0} \phantom{0} \\
 \hline
 1 \phantom{0} 1
 \end{array}$$

 $= 11011$ f)  $11 - 01$ 

$$\begin{array}{r}
 \phantom{-} \phantom{1} \phantom{0} \phantom{0} \phantom{0} \phantom{0} \\
 \phantom{-} \phantom{1} \phantom{0} \phantom{0} \phantom{0} \phantom{0} \\
 - \phantom{1} \phantom{0} \phantom{0} \phantom{0} \phantom{0} \phantom{0} \\
 \hline
 1 \phantom{0} \phantom{0} \phantom{0} \phantom{0} \phantom{0} \phantom{0} \\
 \hline
 1 \phantom{0} \phantom{0} \phantom{0} \phantom{0} \phantom{0} \phantom{0}
 \end{array}$$

 $= 10$ g)  $11 - 10$ 

$$\begin{array}{r}
 \phantom{-} \phantom{1} \phantom{0} \phantom{0} \phantom{0} \phantom{0} \\
 \phantom{-} \phantom{1} \phantom{0} \phantom{0} \phantom{0} \phantom{0} \\
 - \phantom{1} \phantom{0} \phantom{0} \phantom{0} \phantom{0} \phantom{0} \\
 \hline
 0 \phantom{0} \phantom{0} \phantom{0} \phantom{0} \phantom{0} \phantom{0} \\
 \hline
 0 \phantom{0} \phantom{0} \phantom{0} \phantom{0} \phantom{0} \phantom{0}
 \end{array}$$

 $= 10$ h)  $111-100$ 

$$\begin{array}{r}
 \phantom{-} \phantom{1} \phantom{1} \phantom{1} \phantom{0} \phantom{0} \phantom{0} \\
 \phantom{-} \phantom{1} \phantom{1} \phantom{1} \phantom{0} \phantom{0} \phantom{0} \\
 - \phantom{1} \phantom{1} \phantom{1} \phantom{0} \phantom{0} \phantom{0} \\
 \hline
 1 \phantom{0} \phantom{0} \phantom{0} \phantom{0} \phantom{0} \phantom{0} \\
 \hline
 1 \phantom{0} \phantom{0} \phantom{0} \phantom{0} \phantom{0} \phantom{0}
 \end{array}$$

 $= 011$

i)  $101 - 011$ 

$$\begin{array}{r}
 {}^0\cancel{1} \quad {}^10 \quad 1 \\
 - \quad 0 \quad 1 \quad 1 \\
 \hline
 0 \quad 1 \quad 0 \\
 \hline
 \end{array}$$

 $= 010$ j)  $11 \times 11$ 

$$\begin{array}{r}
 \quad \quad \quad 1 \quad 1 \\
 \times \quad \quad \quad 1 \quad 1 \\
 \hline
 \quad \quad \quad 1 \quad 1 \\
 + \quad \quad 1 \quad 1 \\
 \hline
 1 \quad 0 \quad 0 \quad 1 \\
 \hline
 \end{array}$$

 $= 1001$ k)  $101 \times 111$ 

$$\begin{array}{r}
 \quad \quad \quad 1 \quad 0 \quad 1 \\
 \times \quad \quad \quad 1 \quad 1 \quad 1 \\
 \hline
 \quad \quad \quad 1 \quad 0 \quad 1 \\
 + \quad \quad 1 \quad 0 \quad 1 \\
 \hline
 1 \quad 0 \quad 0 \quad 0 \quad 1 \quad 1 \\
 \hline
 1 \quad 1 \quad 1 \\
 \hline
 \end{array}$$

 $= 100011$ l)  $1101 \times 1010$ 

$$\begin{array}{r}
 \quad \quad \quad 1 \quad 1 \quad 0 \quad 1 \\
 \times \quad \quad \quad 1 \quad 0 \quad 1 \quad 0 \\
 \hline
 \quad \quad \quad 0 \quad 0 \quad 0 \quad 0 \\
 \quad \quad 1 \quad 1 \quad 0 \quad 1 \\
 \quad 0 \quad 0 \quad 0 \quad 0 \\
 + \quad 1 \quad 1 \quad 0 \quad 1 \\
 \hline
 1 \quad 0 \quad 0 \quad 1 \quad 0 \quad 0 \quad 1 \quad 0 \\
 \hline
 1 \quad 1 \quad 1 \\
 \hline
 \end{array}$$

 $= 10000010$ m)  $110 \div 11$ 

$$\begin{array}{r}
 \quad \quad 1 \quad 0 \\
 1 \quad 1 \quad \overline{) 1 \quad 1 \quad 0} \\
 \quad \underline{1 \quad 1} \quad \\
 0 \quad 0 \quad 0
 \end{array}$$

 $= 10$ n)  $110 \div 10$ 

$$\begin{array}{r}
 \quad \quad 1 \quad 1 \\
 1 \quad 0 \quad \overline{) 1 \quad 1 \quad 0} \\
 \quad \underline{1 \quad 0} \quad \\
 0 \quad 1 \quad 0 \\
 \quad \underline{1 \quad 0} \quad \\
 0 \quad 0
 \end{array}$$

 $= 11$ o)  $1100 \div 100$ 

$$\begin{array}{r}
 \quad \quad \quad 1 \quad 1 \\
 1 \quad 0 \quad 0 \quad \overline{) 1 \quad 1 \quad 0 \quad 0} \\
 \quad \underline{1 \quad 0 \quad 0} \quad \\
 0 \quad 1 \quad 0 \quad 0 \\
 \quad \underline{1 \quad 0 \quad 0} \quad \\
 0 \quad 0 \quad 0
 \end{array}$$

 $= 11$



k)  $1100 - 1001$ 

$$\begin{array}{r}
 \begin{array}{cccc}
 & 1 & & {}^0\cancel{1} & {}^1\cancel{0} & {}^{10}\cancel{0} \\
 - & 1 & & 0 & 0 & 1 \\
 \hline
 & 0 & & 0 & 1 & 1
 \end{array}
 \end{array}$$

 $= 0011$ m)  $11 \times 11$ 

$$\begin{array}{r}
 \begin{array}{cccc}
 & & 1 & 1 \\
 \times & & 1 & 1 \\
 \hline
 & & 1 & 1 \\
 + & & 1 & 1 \\
 \hline
 1 & 0 & 0 & 1 \\
 \hline
 1
 \end{array}
 \end{array}$$

 $= 1001$ o)  $111 \times 101$ 

$$\begin{array}{r}
 \begin{array}{cccc}
 & & 1 & 1 & 1 \\
 \times & & 1 & 0 & 1 \\
 \hline
 & & 1 & 1 & 1 \\
 & 0 & 0 & 0 & \\
 + & 1 & 1 & 1 & \\
 \hline
 1 & 0 & 0 & 0 & 1 & 1 \\
 \hline
 1 & 1 & & & & 
 \end{array}
 \end{array}$$

 $= 100011$ q)  $1101 \times 1101$ 

$$\begin{array}{r}
 \begin{array}{cccc}
 & & 1 & 1 & 0 & 1 \\
 \times & & 1 & 1 & 0 & 1 \\
 \hline
 & & 1 & 1 & 0 & 1 \\
 & 0 & 0 & 0 & 0 & \\
 & 1 & 1 & 0 & 1 & \\
 + & 1 & 1 & 0 & 1 & \\
 \hline
 1 & 0 & 1 & 0 & 1 & 0 & 0 & 1 \\
 \hline
 1 & 1 & 1 & 1 & 1 & & & 
 \end{array}
 \end{array}$$

 $= 10101001$ l)  $11010 - 10111$ 

$$\begin{array}{r}
 \begin{array}{ccccc}
 & 1 & & {}^0\cancel{1} & {}^1\cancel{0} & {}^{10}\cancel{1} & {}^{10}\cancel{0} \\
 - & 1 & & 0 & 1 & 1 & 1 \\
 \hline
 & 0 & & 0 & 0 & 1 & 1
 \end{array}
 \end{array}$$

 $= 00011$ n)  $100 \times 10$ 

$$\begin{array}{r}
 \begin{array}{cccc}
 & & 1 & 0 & 0 \\
 \times & & 1 & 1 & 0 \\
 \hline
 & & 0 & 0 & 0 \\
 + & 1 & 0 & 0 & \\
 \hline
 1 & 0 & 0 & 0 & 
 \end{array}
 \end{array}$$

 $= 1000$ p)  $1000 \times 110$ 

$$\begin{array}{r}
 \begin{array}{cccc}
 & & 1 & 1 & 0 & 1 \\
 \times & & 1 & 1 & 0 \\
 \hline
 & & 0 & 0 & 0 & 0 \\
 & 1 & 0 & 0 & 1 & \\
 + & 1 & 0 & 0 & 1 & \\
 \hline
 1 & 1 & 0 & 1 & 1 & 0
 \end{array}
 \end{array}$$

 $= 110110$ r)  $1110 \times 1101$ 

$$\begin{array}{r}
 \begin{array}{cccc}
 & & 1 & 1 & 1 & 0 \\
 \times & & 1 & 1 & 0 & 1 \\
 \hline
 & & 1 & 1 & 1 & 0 \\
 & 0 & 0 & 0 & 0 & \\
 & 1 & 1 & 1 & 0 & \\
 + & 1 & 1 & 1 & 0 & \\
 \hline
 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 \\
 \hline
 1 & 1 & 1 & 1 & & & & 
 \end{array}
 \end{array}$$

 $= 10110110$

t)  $1001 \div 11$

[illegible]

$$= 11$$

$$\begin{array}{cccc} & & 1 & 1 \\ 1 & 0 & 0 & \overline{\begin{array}{cccc} 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & \\ \hline 0 & 1 & 0 & 0 \\ & 1 & 0 & 0 \\ & \hline & 0 & 0 & 0 \end{array}} \end{array}$$

$$= 11$$

## S.3. NEGATIVE NUMBERS

📅 03-10-22

🕒 16:00

🎓 Farzad

📍 RB LT1

In computers, subtraction is not possible. We must convert the calculation to be an addition. For example  $5-3$  is not possible, so it becomes  $5+(-3)$ . This means we need to be able to represent negative numbers in binary; there are three methods we can use to do this.

### Sign and Magnitude

In this method, the Most Significant Bit (MSB) is replaced to show the sign rather than a number. A 0 represents a positive number and a 1 represents a negative number. The other bits behave the same.

Converting to and from sign and magnitude binary and decimal is the same as unsigned binary.

	+/-	64	32	16	8	4	2	1
27	0	0	0	1	1	0	1	1
-27	1	0	0	1	1	0	1	1
+13	0	0	0	0	1	1	0	1
-34	1	0	1	0	0	0	1	0

### 1's Complement

To convert to 1's complement, first you need to convert to unsigned binary. You then invert the bits so that 0s become 1s and 1s become 0s.

When doing a 1s complement addition, its important that any overflow bits are carried around to the least significant bit and added on there.

#### 1's Complement subtraction example

Perform the calculation  $10-6 = 1010 - 0110$ .

First, convert the second value to 1s complement =  $1010 + 1001$ . Then draw out the addition grid and perform the addition

$$\begin{array}{r}
 \begin{array}{cccc}
 1 & 0 & 1 & 0 \\
 + & 1 & 0 & 0 & 1 \\
 \hline
 0 & 0 & 1 & 1 \\
 \hline
 1
 \end{array}
 \end{array}$$

As we have an overflowing carry, we have to add this to the least significant bit of the answer.

$$\begin{array}{r}
 \begin{array}{cccc}
 1 & 0 & 1 & 0 \\
 + & 1 & 0 & 0 & 1 \\
 \hline
 0 & 0 & 1 & 1 \\
 \hline
 + & & & & 1 \\
 \hline
 0 & 1 & 0 & 0 \\
 \hline
 1 & 1
 \end{array}
 \end{array}$$

And here we have our final answer, 4.

## 2's Complement

To convert to decimal to 2's complement binary, first convert to unsigned binary. Then work from right to left, inverting the bits so that 0 becomes 1 and 1 becomes 0. However, don't flip any bits to the right of or including the first 1. All bits to the left of should be flipped.

### 2's Complement subtraction example

Perform the calculation  $6-1 = 110-001$ .

First, convert the second value to 2's complement = 111. Then draw out the addition grids and perform the addition.

$$\begin{array}{r} \phantom{+} \phantom{1} \phantom{1} \phantom{0} \\ + \phantom{1} \phantom{1} \phantom{1} \phantom{1} \\ \hline \phantom{+} \phantom{1} \phantom{0} \phantom{1} \\ \hline 1 \phantom{1} \phantom{1} \phantom{1} \phantom{1} \end{array}$$

We have an overflow carry, we discard this. This gives us our final answer of 5.



# WORKSHEET 2/1

📅 04-10-22

👁 Worksheet

## Core Exercises

Find the negative numbers using the sign and method for the following.

- (a)  $-7 = 10000111$
- (b)  $-12 = 10001100$
- (c)  $-15 = 1001111$
- (d)  $-46 = 10101110$
- (e)  $-57 = 10111001$
- (f)  $-112 = 11110000$
- (g)  $-149 = 100010010101$
- (h)  $-179 = 100010110011$
- (i)  $-216 = 100011011000$
- (j)  $-406 = 100110010110$
- (k)  $-1001 = 11111101001$
- (l)  $-1645 = 111001101101$

Find 1's Complement for the following.

- (a)  $01 = 10$
- (b)  $10 = 01$
- (c)  $111 = 000$
- (d)  $011 = 100$
- (e)  $100 = 011$
- (f)  $101 = 010$
- (g)  $0011 = 1100$
- (h)  $1001 = 0110$
- (i)  $10111 = 01000$

Find 2's Complement for the following.

(a) 01

$$\begin{array}{r} \phantom{0}1 \phantom{0} \\ + \phantom{0}1 \\ \hline \phantom{0}1 \phantom{0} \\ \hline \end{array}$$

= 11

(b) 10

$$\begin{array}{r} \phantom{0}0 \phantom{1} \\ + \phantom{0}1 \\ \hline \phantom{0}1 \phantom{0} \\ \hline \phantom{0}1 \\ \hline \end{array}$$

= 10

(c) 111

$$\begin{array}{r}
 0 \quad 0 \quad 0 \\
 + \quad \quad \quad 1 \\
 \hline
 0 \quad 0 \quad 1 \\
 \hline
 \end{array}$$

= 001

(d) 011

$$\begin{array}{r}
 1 \quad 0 \quad 0 \\
 + \quad \quad \quad 1 \\
 \hline
 1 \quad 0 \quad 1 \\
 \hline
 \end{array}$$

= 101

(e) 100

$$\begin{array}{r}
 0 \quad 1 \quad 1 \\
 + \quad \quad \quad 1 \\
 \hline
 1 \quad 0 \quad 0 \\
 \hline
 1 \quad 1 \quad 0
 \end{array}$$

= 100

(f) 101

$$\begin{array}{r}
 0 \quad 1 \quad 0 \\
 + \quad \quad \quad 1 \\
 \hline
 0 \quad 1 \quad 1 \\
 \hline
 \end{array}$$

= 011

(g) 0011

$$\begin{array}{r}
 1 \quad 1 \quad 0 \quad 0 \\
 + \quad \quad \quad 1 \\
 \hline
 1 \quad 1 \quad 0 \quad 1 \\
 \hline
 \end{array}$$

= 1101

(h) 1001

$$\begin{array}{r}
 1 \quad 1 \quad 1 \quad 0 \\
 + \quad \quad \quad 1 \\
 \hline
 0 \quad 1 \quad 1 \quad 1 \\
 \hline
 \end{array}$$

= 0111

(i) 10111

$$\begin{array}{r}
 0 \quad 1 \quad 0 \quad 0 \quad 0 \\
 + \quad \quad \quad 1 \\
 \hline
 0 \quad 1 \quad 0 \quad 0 \quad 1 \\
 \hline
 \end{array}$$

= 01001

Carry out the following subtractions using (i) 1's complement and then (ii) 2's complement.  
*1's complement shown on the left and 2's complement shown on the right.*

(a) 11-01

$$\begin{array}{r}
 {}^1 1 \quad 1 \\
 + \quad 1 \quad 0 \\
 \hline
 0 \quad 1 \\
 + \quad \quad 1 \\
 \hline
 1 \quad 0 \\
 \hline
 1
 \end{array}$$

= 10

$$\begin{array}{r}
 {}^1 1 \quad 1 \\
 + \quad 1 \quad 1 \\
 \hline
 \cancel{1} 1 \quad 0 \\
 \hline
 \end{array}$$

= 10

(b) 11-10

$$\begin{array}{r}
 {}^1 \quad {}^1 1 \quad 1 \\
 + \quad 0 \quad 1 \\
 \hline
 \quad 0 \quad 0 \\
 + \quad \quad 1 \\
 \hline
 \quad 0 \quad 1 \\
 \hline
 \end{array}$$

$$\begin{array}{r}
 {}^1 \quad 1 \quad 1 \\
 + \quad 1 \quad 0 \\
 \hline
 \cancel{1} \quad 0 \quad 1 \\
 \hline
 \end{array}$$

$$= 01$$

$$= 01$$

(c) 101-011

$$\begin{array}{r}
 \quad 1 \quad 0 \quad 1 \\
 + \quad 1 \quad 0 \quad 0 \\
 \hline
 \quad 0 \quad 0 \quad 1 \\
 + \quad \quad \quad 1 \\
 \hline
 \quad 0 \quad 1 \quad 0 \\
 \hline
 \end{array}$$

$$\begin{array}{r}
 {}^1 \quad 1 \quad {}^1 0 \quad 1 \\
 + \quad 1 \quad 0 \quad 1 \\
 \hline
 \cancel{1} \quad 0 \quad 1 \quad 0 \\
 \hline
 \end{array}$$

$$= 010$$

$$= 010$$

(d) 101-100

$$\begin{array}{r}
 {}^1 \quad {}^1 1 \quad {}^1 0 \quad 1 \\
 + \quad 0 \quad 1 \quad 1 \\
 \hline
 \quad 0 \quad 0 \quad 0 \\
 + \quad \quad \quad 1 \\
 \hline
 \quad 0 \quad 0 \quad 1 \\
 \hline
 \end{array}$$

$$\begin{array}{r}
 {}^1 \quad 1 \quad 0 \quad 1 \\
 + \quad 1 \quad 0 \quad 0 \\
 \hline
 \cancel{1} \quad 0 \quad 0 \quad 1 \\
 \hline
 \end{array}$$

$$= 001$$

$$= 001$$

(e) 110-101

$$\begin{array}{r}
 {}^1 \quad {}^1 1 \quad 1 \quad 0 \\
 + \quad 0 \quad 1 \quad 0 \\
 \hline
 \quad 0 \quad 0 \quad 0 \\
 + \quad \quad \quad 1 \\
 \hline
 \quad 0 \quad 0 \quad 1 \\
 \hline
 \end{array}$$

$$\begin{array}{r}
 {}^1 \quad {}^1 1 \quad 1 \quad 0 \\
 + \quad 0 \quad 1 \quad 1 \\
 \hline
 \cancel{1} \quad 0 \quad 0 \quad 1 \\
 \hline
 \end{array}$$

$$= 001$$

$$= 001$$

(f) 1110 - 0011

$$\begin{array}{r}
 {}^1 \quad {}^1 1 \quad 1 \quad 1 \quad 0 \\
 + \quad 1 \quad 1 \quad 0 \quad 0 \\
 \hline
 \quad 1 \quad 0 \quad 1 \quad 0 \\
 + \\
 \hline
 \quad 1 \quad 0 \quad 1 \quad 1
 \end{array}$$

$$= 1011$$

(g) 1100 - 1001

$$\begin{array}{r}
 {}^1 \quad {}^1 1 \quad 1 \quad 0 \quad 0 \\
 + \quad 0 \quad 1 \quad 1 \quad 0 \\
 \hline
 \quad 0 \quad 0 \quad 1 \quad 0 \\
 + \\
 \hline
 \quad 0 \quad 0 \quad 1 \quad 1
 \end{array}$$

$$= 0011$$

(h) 11110 - 10011

$$\begin{array}{r}
 {}^1 \quad {}^1 1 \quad {}^1 1 \quad 1 \quad 1 \quad 0 \\
 + \quad 0 \quad 1 \quad 1 \quad 0 \quad 0 \\
 \hline
 \quad 0 \quad 1 \quad 0 \quad 1 \quad 0 \\
 + \\
 \hline
 \quad 0 \quad 1 \quad 0 \quad 1 \quad 1
 \end{array}$$

$$= 01011$$

(i) 10111 - 10110

$$\begin{array}{r}
 {}^1 \quad {}^1 1 \quad {}^1 0 \quad {}^1 1 \quad {}^1 1 \quad 1 \\
 + \quad 0 \quad 1 \quad 0 \quad 0 \quad 1 \\
 \hline
 \quad 0 \quad 0 \quad 0 \quad 0 \quad 0 \\
 + \\
 \hline
 \quad 0 \quad 0 \quad 0 \quad 0 \quad 1
 \end{array}$$

$$= 00001$$

$$\begin{array}{r}
 {}^1 \quad {}^1 1 \quad 1 \quad 1 \quad 0 \\
 + \quad 1 \quad 1 \quad 0 \quad 1 \\
 \hline
 \cancel{1} \quad 1 \quad 0 \quad 1 \quad 1
 \end{array}$$

$$= 1011$$

$$\begin{array}{r}
 {}^1 \quad {}^1 1 \quad 1 \quad 0 \quad 0 \\
 + \quad 0 \quad 1 \quad 1 \quad 1 \\
 \hline
 \cancel{1} \quad 0 \quad 0 \quad 1 \quad 1
 \end{array}$$

$$= 0011$$


$$\begin{array}{r}
 {}^1 \quad {}^1 1 \quad {}^1 1 \quad 1 \quad 1 \quad 0 \\
 + \quad 0 \quad 1 \quad 1 \quad 0 \quad 1 \\
 \hline
 \cancel{1} \quad 0 \quad 1 \quad 0 \quad 1 \quad 1
 \end{array}$$

$$= 01011$$

$$\begin{array}{r}
 {}^1 \quad {}^1 1 \quad {}^1 0 \quad 1 \quad 1 \quad 1 \\
 + \quad 0 \quad 1 \quad 0 \quad 1 \quad 0 \\
 \hline
 \cancel{1} \quad 0 \quad 0 \quad 0 \quad 0 \quad 1
 \end{array}$$

$$= 00001$$

# WORKSHEET 2/2

 05-10-22

 Worksheet

1. Convert the decimal numbers to signed and unsigned 8-bit binary numbers.

Decimal	Unsigned Binary	Signed		
		Sign-Magnitude	1's Complement	2's Complement
11	00001011	00001011	00001011	00001011
-11	—	10001011	11110100	11110101
29	00011101	00011101	00011101	00011101
-29	—	10011101	11100010	11100011
114	01110010	01110010	01110010	01110010
-114	—	11110010	10001101	10001110
111	01101111	01101111	01101111	01101111
-111	—	11101111	10010000	10010001

2. Convert the binary numbers using the required methods.

Binary Numbers	Unsigned	Signed		
		Sign-Magnitude	1's Complement	2's Complement
00001101	13	13	13	13
10001001	137	-9	-118	-119
11000100	196	-68	-59	-60
01000100	68	68	68	68
11111111	255	-127	0	-1
01111111	127	127	127	127
10100101	165	-37	-90	-91
11001001	201	-73	-54	-55

## S.4. DIGITAL GATES

📅 10-10-22

🕒 16:00

🎓 Farzad

📍 RB LT1

Digital gates are used as the building block of digital systems. They manipulate inputs to provide outputs.

Throughout this lecture, its important to remember that a signal of 1 represents on (or high voltage) and a signal of 0 represents off (or low voltage).

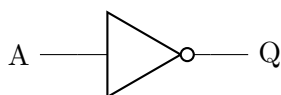
Digital gates can be represented by a circuit symbol. Their inputs and outputs can be mapped onto a Truth Table and they have symbols which can be used in expressions to describe the circuit.

### NOT Gate

This can also be called an inverter.

As the name inverter suggests, the NOT gate inverts the bits. This means a 0 inputted, will output a 1 and a 1 inputted will output a 0.

NOT gate can only have one input.



$$Q = \overline{A}$$

$$Q = A'$$

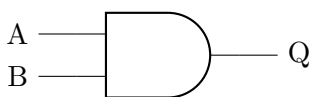
A	Q
0	1
1	0

Truth table for the NOT gate.

### AND Gate

This gate compares two or more inputs. If all its inputs are 1, then it outputs 1; otherwise it outputs 0.

The rules of binary multiplication are the same of the AND gate.



$$Q = A \bigwedge B$$

$$Q = A \cdot B$$

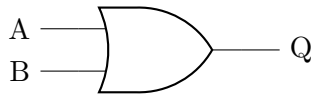
$$Q = AB$$

B	A	Q
0	0	0
0	1	0
1	0	0
1	1	1

Truth table for the AND gate.

### OR Gate

This gate compares two or more inputs. If one or more input is 1, then it outputs 1; otherwise it outputs 0.



$$Q = A \vee B$$

$$Q = A + B$$

B	A	Q
0	0	0
0	1	1
1	0	1
1	1	1

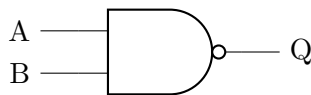
Truth table for the OR gate.

## NAND Gate

*Not AND*

This gate compares two or more inputs. It outputs 1 where at least one of the inputs are not 1 and 0 where all the inputs are 1.

Through combinations of this gate, all the other logic gates can be made, due to this, it is called a Universal Gate.



$$Q = \overline{AB} = (AB)'$$

$$Q = \overline{A \wedge B} = (A \wedge B)'$$

$$Q = \overline{A \cdot B} = (A \cdot B)'$$

B	A	Q
0	0	1
0	1	1
1	0	1
1	1	0

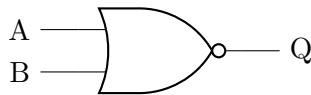
Truth table for the NAND gate.

## NOR Gate

*Not OR*

This gate compares two or more inputs. Where all the inputs are 0, it outputs 1; otherwise it outputs 0.

Through combinations of this gate, all other logic gates can be constructed, due to this it can be called a Universal Gate.



$$Q = \overline{A + B} = (A + B)'$$

$$Q = \overline{A \vee B} = (A \vee B)'$$

B	A	Q
0	0	1
0	1	0
1	0	0
1	1	0

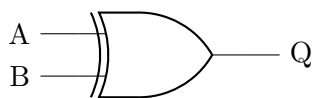
Truth table for the NOR gate.

## XOR Gate

*eXclusive OR*

This gate compares two or more inputs. For a 2-input XOR gate, where the two inputs are different, it outputs 1; otherwise it outputs 0.

This gate is used for adder circuits.



$$Q = A \oplus B$$

$$Q = AB' + A'B$$

$$Q = A \cdot \overline{B} + \overline{A} \cdot B$$

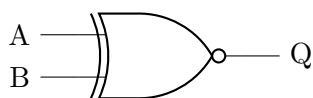
B	A	Q
0	0	0
0	1	1
1	0	1
1	1	0

Truth table for the XOR gate.

## XNOR Gate

*eXclusive Not OR*

This gate compares two or more inputs and where the inputs are the same, it outputs 1 and where the inputs are different, it outputs 0.



$$Q = \overline{A \oplus B}$$

$$Q = (AB' + A'B)'$$

B	A	Q
0	0	1
0	1	0
1	0	0
1	1	1

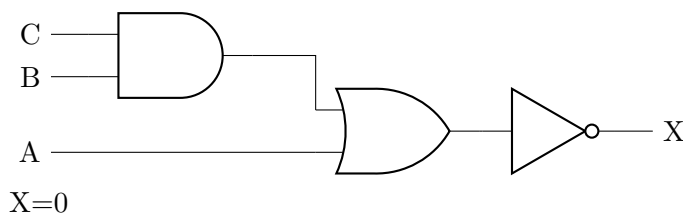
Truth table for the XNOR gate.

## Practice Questions

### Question 1

Find the value of X where the inputs have the following values.

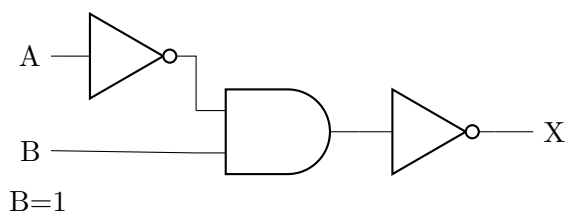
$$A = 0 \quad B = 1 \quad C = 1$$



### Question 2

Find the value of B where the logic system is as follows and has the following values.

$$X = 0 \quad A = 0$$

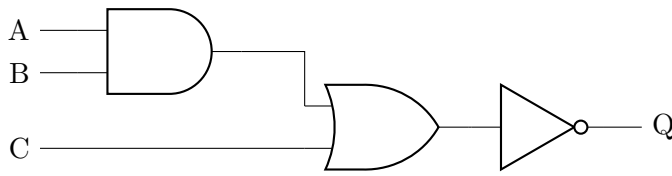




**Question 3**

Convert the following boolean logic expression to a circuit diagram.

$$\overline{A \cdot B + C}$$



# WORKSHEET 03

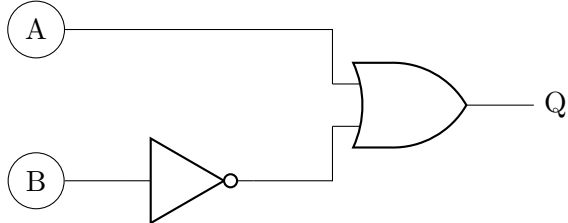
📅 15-10-22

👁 Worksheet

## Basic Exercises

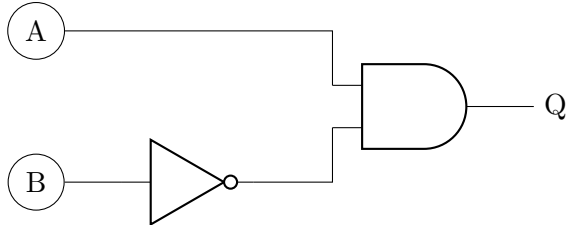
1. Translate the following boolean expressions into a digital gate circuit and construct a truth table.

a.  $A + \overline{B}$



B	A	Q
0	0	1
0	1	1
1	0	0
1	1	1

b.  $A \cdot \overline{B}$



B	A	Q
0	0	0
0	1	1
1	0	0
1	1	0