Solutions to MATH3411 Problems 76–83

76. We first factor n = 8051:

$$x_0 \equiv 1$$
 $\gcd(x_{2i} - x_i, n)$
 $x_1 \equiv 2$ $\gcd(x_{2i} - x_i, n)$
 $x_2 \equiv 5$ $x_2 - x_1 \equiv 3$ 1
 $x_3 \equiv 26$ $x_4 \equiv 677$ $x_4 - x_2 \equiv 672$ 1
 $x_5 \equiv 7474$ $x_6 \equiv 2839$ $x_6 - x_3 \equiv 2813$ 97

We have found a factor d = 97. It is a prime and so is $\frac{n}{d} = 83$, so $n = 91643 = 83 \times 97$. Now let us factor n = 201001:

$$x_0 \equiv 1$$
 $\gcd(x_{2i} - x_i, n)$
 $x_1 \equiv 2$
 $x_2 \equiv 5$ $x_2 - x_1 \equiv 3$ 1
 $x_3 \equiv 26$ $x_4 \equiv 677$ $x_4 - x_2 \equiv 672$ 1
 $x_5 \equiv 56328$ $x_6 \equiv 42800$ $x_6 - x_3 \equiv 42774$ 1
 $x_7 \equiv 117888$ $x_8 \equiv 170404$ $x_8 - x_4 \equiv 169727$ 19

We have found a prime factor d = 19 of n = 201001.

We now factor the quotient $\frac{n}{d} = 10579$:

We have found another prime factor d = 71, and the quotient $\frac{201001}{19 \times 71} = 149$ is also prime, so $n = 201001 = 19 \times 71 \times 149$.

77. Let us now first use Fermat's Method: For n = 92131, $\lceil \sqrt{n} \rceil =$, so

t	2t + 1	$s^2 = t^2 - n$	$s \in \mathbb{Z}$?
304	609	285	×
305	611	894	×
306	613	1505	×
307	615	2118	×
308	617	2733	×
309	619	3350	×
310	621	3969	\checkmark

so t = 310 and $s = \sqrt{3969} = 63$, and a = s + t = 373 and b = t - s = 247; hence, $92131 = n = ab = 373 \times 247$. The number 373 is prime but 247 is not; we need to factorise the latter: $\lceil \sqrt{247} \rceil = 16$, so

$$\begin{array}{c|cccc} t & 2t+1 & s^2=t^2-n & s\in\mathbb{Z}?\\ \hline 16 & 33 & 9 & \checkmark \end{array}$$

so t = 16 and $s = \sqrt{9} = 3$, and a = s + t = 19 and b = t - s = 13; these are prime factors, so $92131 = 13 \times 19 \times 373$.

Let us now use the Pollard ρ -Method:

$$x_0 \equiv 2$$
 $\gcd(x_{2i} - x_i, n)$
 $x_1 \equiv 5$ $\gcd(x_{2i} - x_i, n)$
 $x_2 \equiv 26$ $x_2 - x_1 \equiv 21$ 1
 $x_3 \equiv 677$ $x_4 \equiv 89806$ $x_4 - x_2 \equiv 89780$ 1
 $x_5 \equiv 62028$ $x_6 \equiv 82225$ $x_6 - x_3 \equiv 81548$ 19

We have found a factor d = 19 and factor the quotient $\frac{n}{d} = 4849$:

$$x_0 \equiv 2$$
 $gcd(x_{2i} - x_i, n)$
 $x_1 \equiv 5$
 $x_2 \equiv 26$ $x_2 - x_1 \equiv 21$ 1
 $x_3 \equiv 677$
 $x_4 \equiv 2524$ $x_4 - x_2 \equiv 2498$ 1
 $x_5 \equiv 3840$
 $x_6 \equiv 4641$ $x_6 - x_3 \equiv 3964$ 1
 $x_7 \equiv 4473$
 $x_8 \equiv 756$ $x_8 - x_4 \equiv 3081$ 13

We have found another prime factor d = 13, and $\frac{92131}{19 \times 13} = 373$ is also prime, so $n = 92131 = 13 \times 19 \times 373$. Finally, we use Shor's Algorithm on n = 92131: $3^{558} \equiv 1 \pmod{n}$, so $n \mid (3^{558} - 1) = (3^{279} - 1)(3^{279} + 1)$. Now, it is possible to calculate $\gcd(n, 3^{279} - 1) = 13$ and $\gcd(n, 3^{279} + 1) = 7087$ to find the factors 13 and $7087 = 19 \times 373$.

Again, we see that $92131 = 13 \times 19 \times 373$.

78.

so t = 159 and $s = \sqrt{784} = 28$, and a = s + t = 187 and b = t - s = 131. Therefore, $24497 = N = ab = 131 \times 187$.

- b) Write $N = 131 = 2^s t + 1$ with s = 1 and t = 65. Since $a^{2^r} t = a^t \equiv -1 \pmod{131}$ for r = 0, n = 131 is a strong pseudo-prime base 2 and therefore quite likely to be a prime.
- c) Note that $2^{40} \equiv 1 \pmod{187}$, so $187|(2^{40}-1)=(2^{20}-1)(2^{20}+1)$. Calculating $\gcd(187,2^{20}-1)=11$ and $\gcd(n,2^{20}+1)=17$ to find the factors 11 and 17; these are prime and are thus the factors of 187.

79. A very simple method could just be to test the p smallest primes $2, 3, 5, \ldots, p$, to see whether $\frac{n}{2^p - 1}$ is an (odd) integer, and if so, whether it is of the form $2^q - 1$ for some prime q. For n = 16646017, we find that p = 17:

$$\frac{n}{2^{17} - 1} = 127 = 2^7 - 1$$

Here, q = 7 is prime, and so $n = (2^{17} - 1)(2^7 - 1) = 131071 \times 127$.

80.

- **a.** 0,1,0,1,1
- **b.** As soon as any consecutive 3 bits x_i, x_{i+1}, x_{i+1} reappear later in the list, the output will repeat itself. The numbers x_0, \ldots, x_9 are as follows:

$$x_0 = 1, 1, 0, 0, 1, 0, 1, (x_7 =) 1, 1, 0$$

We see that the first sequence 1,1,0 reappears 7 steps later; 7 is therefore the period of this LSFR.

81.

a. The output numbers x_0, \ldots, x_{19} are as follows:

i	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19
x_i	1	8	3	4	11	6	7	14	9	10	17	12	13	2	15	16	5	0	1	8

As soon as any number reappears later in the list, the output will repeat itself. We see that $x_0 = 1$ reappears as $x_{18} = 1$; the period of this LSFR is therefore 18.

b. The output numbers x_0, \ldots, x_{19} are as follows:

i		0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19
x_i	i	1	0	0	0	1	1	1	1	0	1	0	1	1	0	0	1	0	0	0	1

As soon as any consecutive 3 bits x_i, \ldots, x_{i+4} reappear later in the list, the output will repeat itself. We see that the first sequence 1,0,0,0 reappears 15 steps later; 15 is therefore the period of this LSFR.

82.

a. First, write the message as it has been received:

FSOTU

OHFOI

UIJNP

RPUTM

TELHE

HQYEN

You can use a number of easy heuristic arguments (for instance, "fourth" goes with "of July", as do "q" and "u") to permute the columns above in the right way but will get the following correct answer:

SUTFO

HIOOF

IPNUJ

PMTRU

EEHTL

QNEHY

This gives us the message "SHIP EQUIPMENT ON THE FOURTH OF JULY".

b.
$$\sigma = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 4 & 1 & 5 & 3 & 2 \end{pmatrix}$$

c. First, write the message as it has been received:

SLAFK

EOROE

LIERO

LLSEV

ASBTE

LHEAR

Now permute the columns:

LKFSA

OEOER

IORLE

LVELS

SETAB

HRALE

The enciphered message is then "LKFSA OEOER IORLE LVELS SETAB HRALE".

83. Looking at the first block HSTII, we might guess that it was enciphered from "THIS I", using one of the permutations

$$\sigma_1 = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 3 & 1 & 4 & 2 & 5 \end{pmatrix} \quad \text{and} \quad \sigma_2 = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 3 & 1 & 5 & 2 & 4 \end{pmatrix}$$

Applying the inverse σ_1^{-1} of the first permutation to the 1st three blocks, we get "THIS IS A TON HERAR". This does not seem to make sense, so lets try the inverse $\sigma_2^{-1} = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 2 & 4 & 1 & 5 & 3 \end{pmatrix}$ of the second permutation: We now get a much more lucid message: THIS IS ANOTHER ARTICLE ON SECRET CODES PQ