

# A Model Tripes Question

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In a propositional language  $p, q, \dots$  are the atomic formulæ;  $\neg p$  etc. are **negatomics**; a **literal** is a formula that is an atomic or a negatomic;  $A, B$  are complex formulæ. The **dual**  $\hat{A}$  of a propositional formula  $A$  is the result of replacing every literal in  $A$  by its negation. (Notice that  $\hat{A}$  is *not* usually the same as  $\neg A$ !)

A formula  $A$  **self-dual** if it is logically equivalent to its own dual: that is to say that  $A \longleftrightarrow \hat{A}$  is a tautology. (For example:  $p \text{ XOR } q$  is self-dual—‘ $p$ ’ and ‘ $q$ ’ being literals—even tho’  $A \text{ XOR } B$  is not self-dual in general.)

Show—by considering disjunctive normal forms or otherwise—that whenever  $A$  is a self-dual formula there is a formula  $B$  such that  $A$  is logically equivalent to  $B \longleftrightarrow \hat{B}$ .