Extracted models and the Independence of Extensionality

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First, some slang. If T is a name for a system of axiomatic set theory (with extensionality of course), then TU is the name for the result of weakening extensionality to the assertion that nonempty sets with the same elements are identical. 'U' is for 'Urelemente'—German for 'atoms'.

We start with a model $\langle V, \in \rangle$ of ZF. The traditional method is to define a new membership relation by taking everything that wasn't a singleton to be empty, and then set y IN z iff $z = \{x\}$ for some x such that $y \in x$: it turns out that the structure $\langle V, IN \rangle$ is a model of ZFU. However there is nothing special about the singleton function here. Any injection from the universe into itself will do. So let's explore this. We start with a model $\langle V, \in \rangle$ of ZF, and an injection $f: V \to V$ which is not a surjection (such as ι).

We then say $x \in_f y$ is false unless y is a value of f and $x \in f^{-1}(y)$. (So that everything that is not an (as it might be) singleton has become an empty set (an *urelement*) in the sense of \in_f).

This gives us a new structure: its domain is the same universe as before, but the membership relation is the new \in_f that we have just defined.

Now we must prove that the structure $\langle V, \in_f \rangle$ is a model of ZF with extensionality weakened to the assertion that *nonempty* sets with the same elements are identical.

What is true in $\langle V, \in_f \rangle$? Try pairing, for example: what is the pair of x and y in the sense of \in_f ? A moment's reflection shows that it must be $f\{x,y\}$: if you are a member of $f\{x,y\}$ in the sense of \in_f then you are a member of $f^{-1} \cdot f\{x,y\}$, so you are obviously x or y. Think about this until you are happy about it. Then try power set and sumset. (The power set of x in the new sense must be f of the set of those things that are subsets-of-x-in-the-new-sense ...). Only later should you start worrying about proving a theorem about what statements are preserved.

For use later

We define an embedding $e: V \to V$ by recursion on \in by

$$e(x) =: f(e "x).$$

It's easy to show that $x \in y$ iff $e(x) \in_f e(x)$ so e is rather nice. You are probably comfortable with the idea of an **end-extension** in connection with, say, linear orders. "All the new stuff is put on the end". There is a corresponding notion of an end-embedding: the thing being embedded-into is an end-extension of the range of the embedding. There is a corresponding notion of end-extension in models of set theory. You have a model of set theory. Add some new sets to it. As long as none of the original sets acquire new members when you do this you say we have an end-extension. "No new members of old sets!". This is an important notion.

The injection we have just defined—e—is an end-embedding. It's also what we call a \mathcal{P} -embedding, namely an embedding that not only adds no new members of old sets but doesn't even add any new subsets of old sets. A \mathcal{P} -embedding preserves $\Delta_0^{\mathcal{P}}$ formulæ, where the class of $\Delta_0^{\mathcal{P}}$ formulæ is the smallest class containing atomics and closed under boolean operations and restricted quantification AND $(\forall x \subseteq Y)(\dots$ and $(\exists x \subseteq Y)(\dots$ A \mathcal{P} -embedding not only adds no new members of old sets (so it's an end-extension) it also adds no new subsets.

There are \mathcal{P} -embeddings all over the place: the embedding from a well-founded model of ZF into any Rieger-Bernays permutation model of it is always a \mathcal{P} -embedding.

I claim e is a \mathcal{P} -embedding. It is a simple matter to check that the range of e is the whole of M. It is easy too to check that anything which \mathcal{M} believes to be a subset of a thing in the range of e is also in the range of e, and this makes e into a \mathcal{P} -embedding.

For T a theory in the language of set theory let T^* be T with an extra unary function letter: *, and two new axioms, (i) $(\forall x)(x^*$ is an urelement); (ii) $(\forall x, y)(x^* = y^* \longleftrightarrow x = y)$. We interpret T^* in T by means of a map σ defined on formulæ in the extended language as follows:

- σ of \in is \in_f .
- σ of $y = x^*$ is to be y = g(x), where g is any old injective function whose range is disjoint from the range of f.

and we extend σ to all other formulæ by recursion. We then check that σ of an axiom of ZF(C)U is a theorem of ZF(C)*.