



MAKE
SCHOOL

HASH TABLES

The Ultimate Data Structure

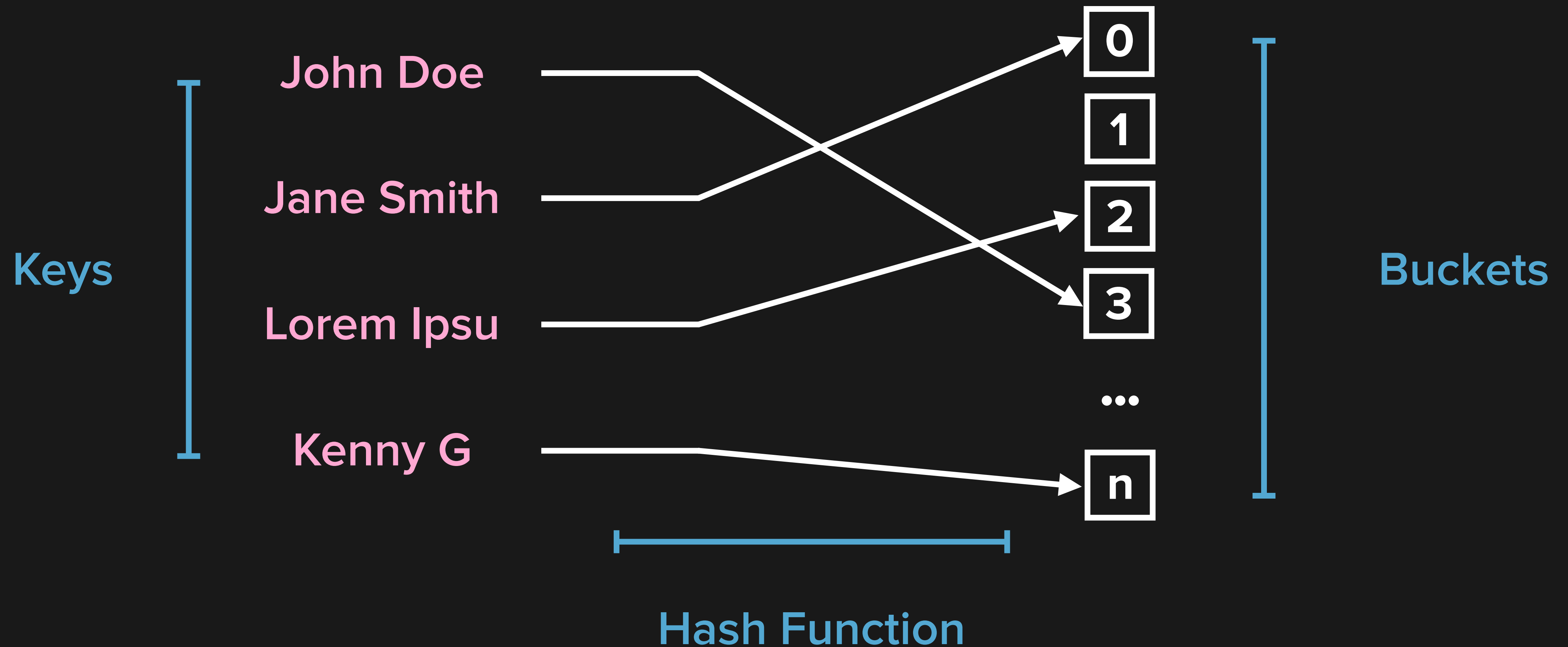
HASH TABLES

Maps keys → objects

`dict()` creates a hash table

Used because of strong average case performance

HASH TABLES



HASH FUNCTIONS

Converts a variable-size
input to a fixed-size
output

John Doe → **512340**

Jane Smith → **408749**

Same input → same output

Lorem Ipsu → **943275**

Input can be anything -
string, pointer, custom
class

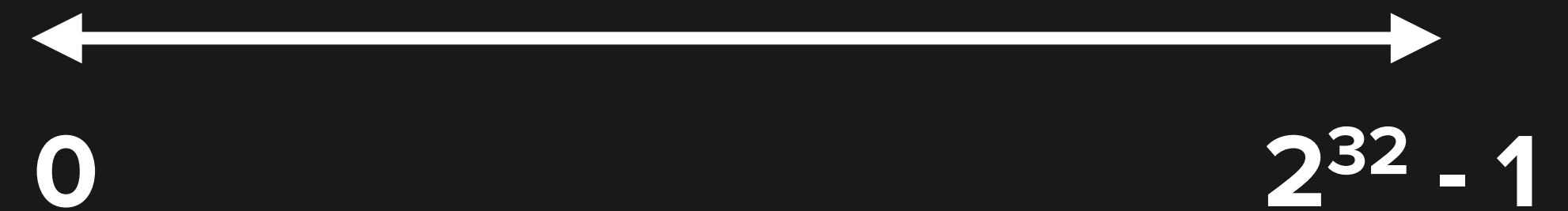
John Doe → **512340**

IDEAL HASH*

Repeatable

Fast

Output is unsigned integer



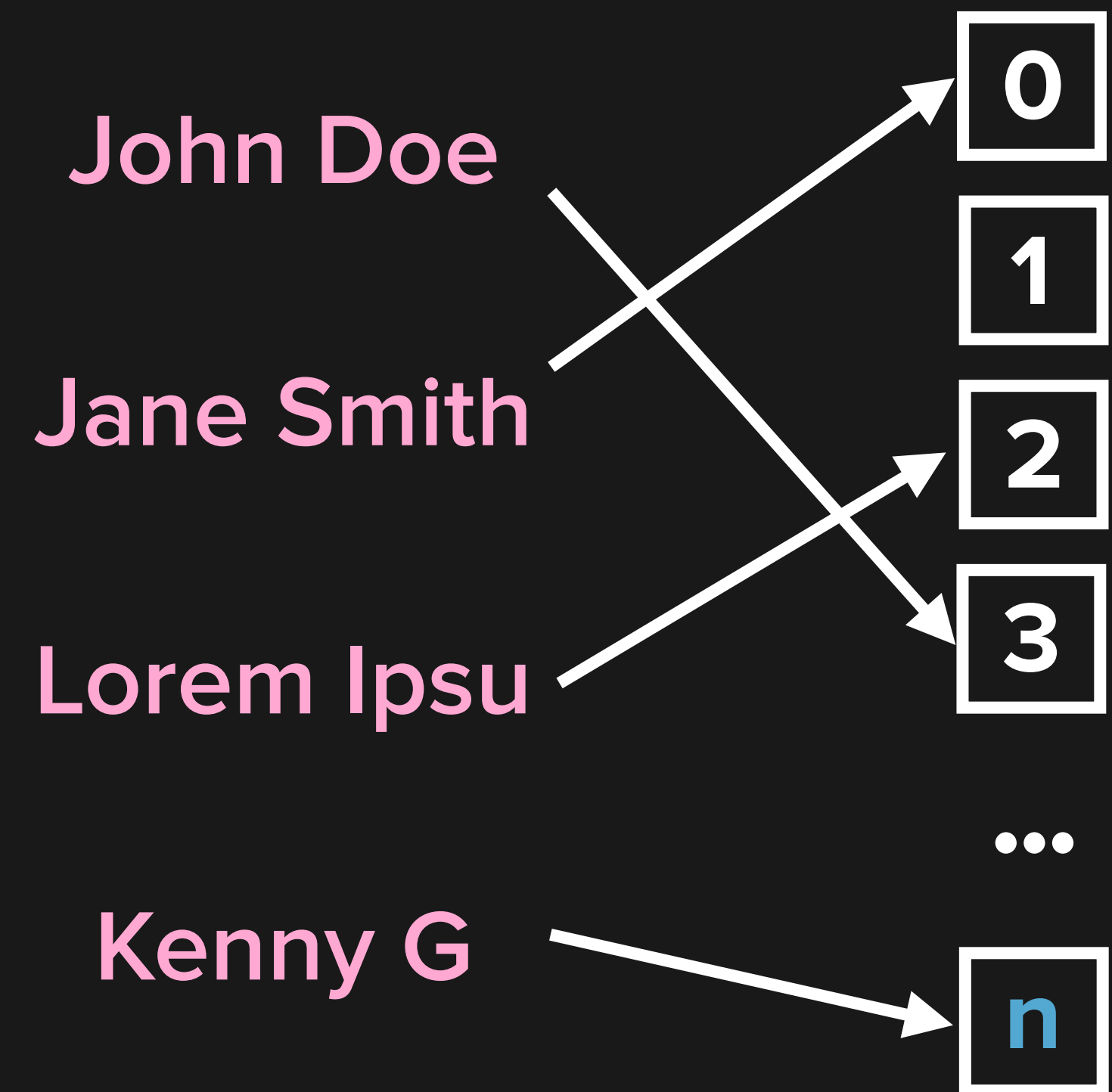
Randomly distributes keys
among output space

Small differences in input result
in large differences in output

**Different for cryptographic hash functions*

WHICH BUCKET?

bucket = hash(key) % n



COLLISIONS

It is impossible to map all possible input to a fixed output space without some inputs generating the same output

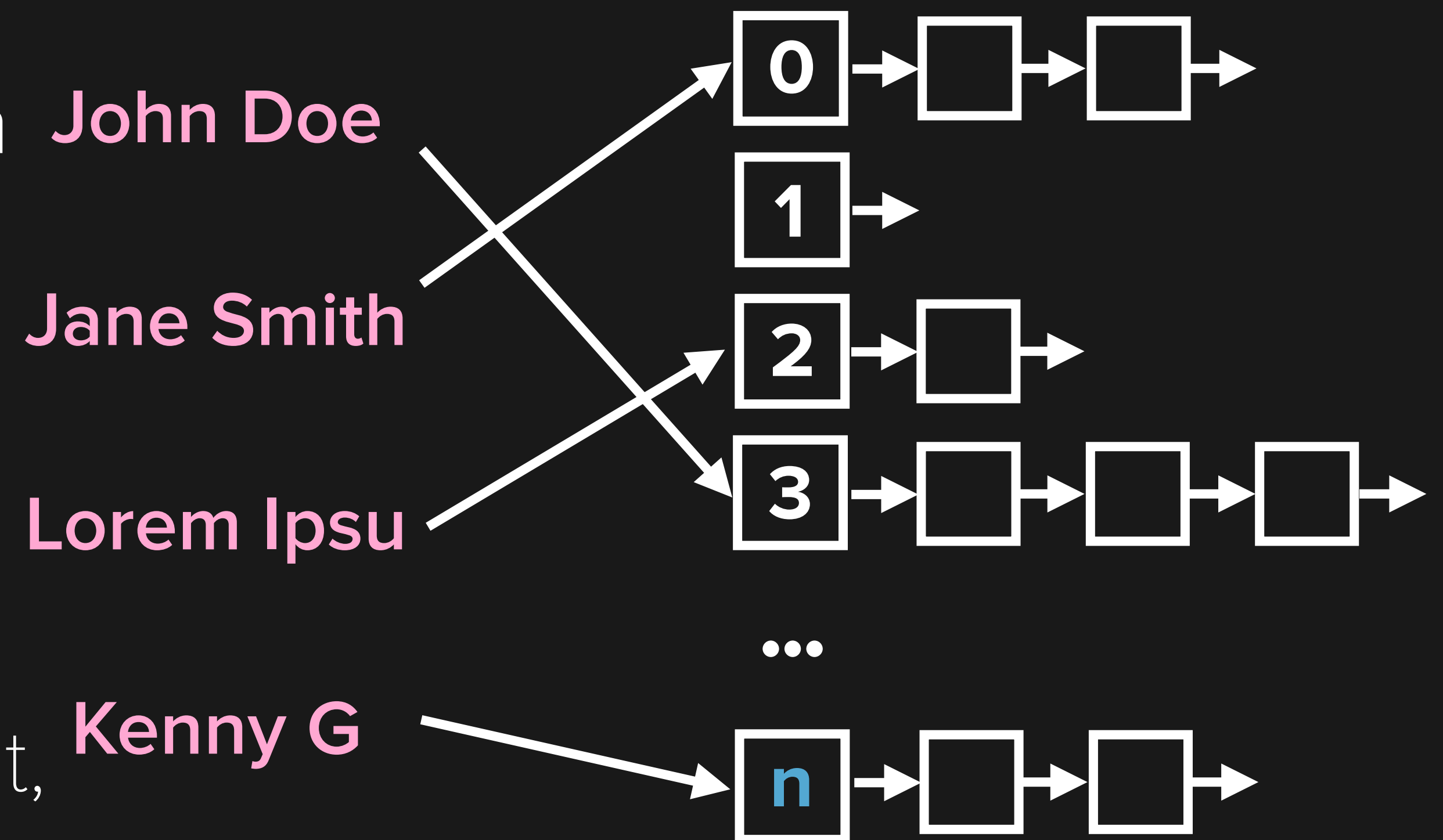
Differing input generating the same output is a collision

CHAINING

Each bucket contains a **John Doe** linked list

On collision - add to end of the linked list

To retrieve - find bucket, **Kenny G** find in linked list



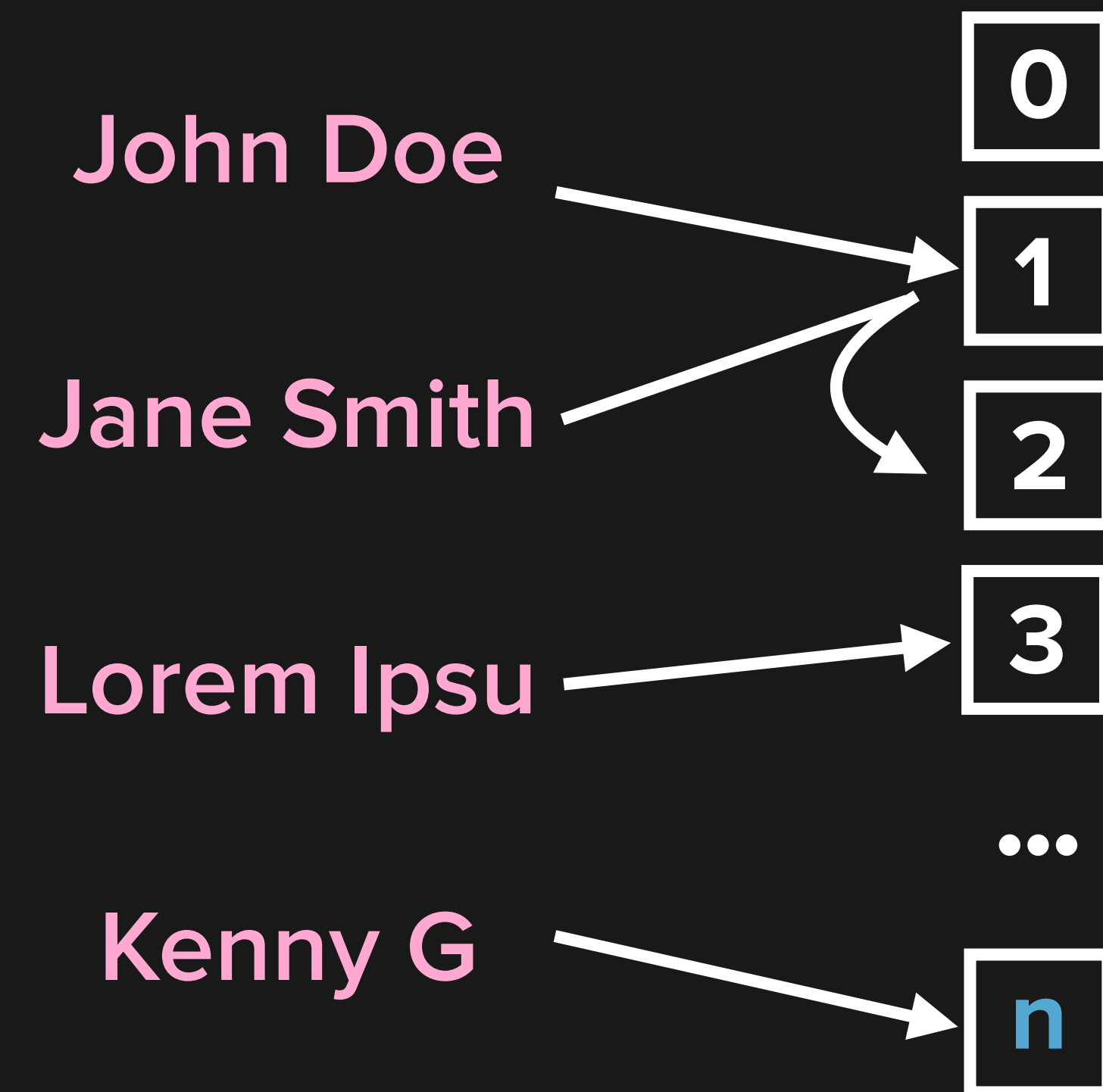
LINEAR PROBING

Each bucket contains one object

On collision - go to next open bucket, add object there

To retrieve - find bucket, if that's not object, iterate buckets until you find it

`dict` does it this way



LOAD FACTOR

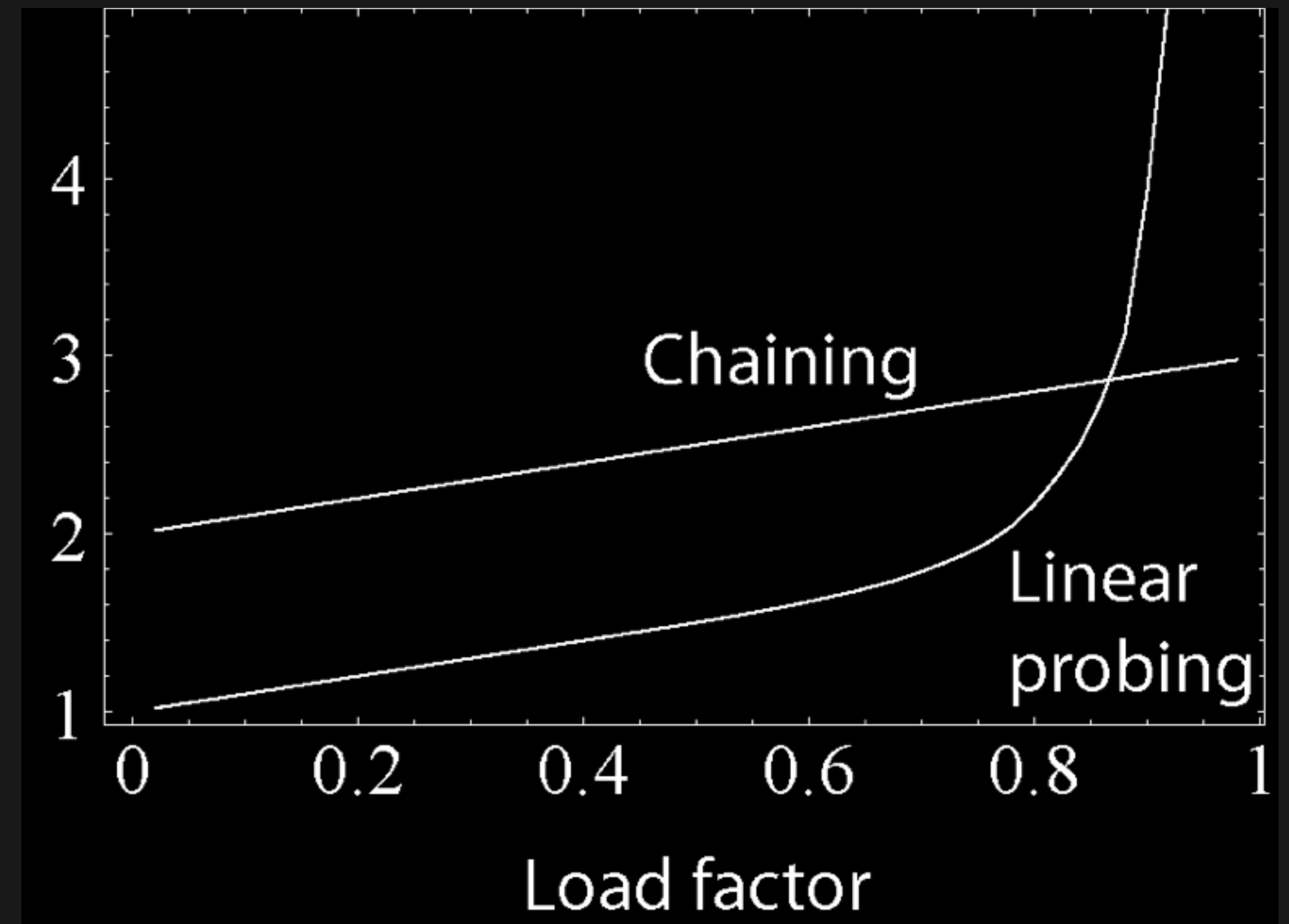
Load Factor = $\text{entries} / \text{buckets}$

For 76 entries in a 128 bucket hash table, that's $76 / 128 = 0.59375$

LOAD FACTOR

Load factor affects performance

Collision resolution affects performance



COMPLEXITY ANALYSIS

	Average Case	Worst Case
Space	$O(n)$	$O(n)$
Search	$O(1)$	$O(n)$
Insert	$O(1)$	$O(n)$
Delete	$O(1)$	$O(n)$

STRING HASHING

Strings are lists of chars

Chars have numerical values

Add up the chars - there's your hash! (Lose Lose algorithm)

But `hash("dog") == hash("god")`