

EECE.3170: Microprocessor Systems Design I

Summer 2016

Homework 4 Solution

1. (40 points) Write the following subroutine in x86 assembly:

```
int f(int v1, int v2, int v3) {  
    int x = v1 + v2;  
    return (x + v3) * (x - v3);  
}
```

Recall that:

- Subroutine arguments are passed on the stack, and can be accessed within the body of the subroutine starting at address EBP+8.
- At the start of each subroutine:
 - i. Save EBP on the stack
 - ii. Copy the current value of the stack pointer (ESP) to EBP
 - iii. Create space within the stack for each local variable by subtracting the appropriate value from ESP. For example, if your function uses four integer local variables, each of which contains four bytes, subtract 16 from ESP.
 - iv. Local variables can then be accessed starting at the address EBP-4.
- A subroutine's return value is typically stored in EAX.

See Lectures 14 and 16-18 for more details on subroutines, the x86 architecture, and the conversion from high-level concepts to low-level assembly.

Solution: Solution is shown on the next page; note that many different solutions are possible. The key points are:

- Setting up the stack frame appropriately (save base pointer; point base pointer to appropriate location; create space for local variable(s); save any overwritten registers except eax).
- Adding $v1 + v2$ while appropriately accessing different memory locations (only one memory operand per instruction; accessing arguments at right addresses relative to ebp)
- Computing return value while appropriately accessing different memory locations
- "Cleaning up" stack frame (restoring saved registers; clearing space for local variable(s); restoring base pointer)

```
f PROC                                ; Start of function f
    push    ebp                       ; Save ebp
    mov     ebp, esp                 ; Copy ebp to esp

    sub     esp, 4                    ; Create space on the
                                    ; stack for x

    push    ebx                       ; Save ebx on the stack
    push    edx                       ; Save edx on the stack

    mov     ebx, DWORD PTR 8[ebp]     ; ebx = v1
    add     ebx, DWORD PTR 12[ebp]    ; ebx = v1 + v2
    mov     DWORD PTR -4[ebp], ebx    ; x = ebx = v1 + v2
    mov     eax, ebx                  ; eax = ebx = x
    add     eax, DWORD PTR 16[ebp]    ; eax = eax + v3 = x + v3
    sub     ebx, DWORD PTR 16[ebp]    ; ebx = ebx - v3 = x - v3
    imul    ebx                       ; (edx, eax) = eax * ebx
                                    ; = (x + v3) * (x - v3)

    pop     edx                       ; Restore edx
    pop     ebx                       ; Restore ebx

    mov     esp, ebp                  ; Clear x
    pop     ebp                       ; Restore ebp

    ret                                ; Return from subroutine
f ENDP
```

2. (60 points) Write the following subroutine in x86 assembly:

```
int fib(int n)
```

Given a single integer argument, n , return the n th value of the Fibonacci sequence—a sequence in which each value is the sum of the previous two values. The first 15 values are shown below—note that the first value is returned if n is 0, not 1.

n	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
$fib(n)$	0	1	1	2	3	5	8	13	21	34	55	89	144	233	377

Solution: How you implement the low-level code for this version of the Fibonacci function depends on the algorithm you use. What follows is both C code and assembly for the algorithm implemented either with or without recursion.

```
int fib(int n) {    // FIBONACCI WITHOUT RECURSION
    int i;           // Loop index
    int first, sec;   // Two previous Fibonacci values
    int cur;         // Value from current iteration

    // For n == 0 or n == 1, fib(n) == n
    if (n <= 1)
        return n;

    // Use loop to calculate fib(n)--at each step,
    // current value is sum of previous two values
    else {
        first = 0;
        sec = 1;
        for (i = 2; i <= n; i++) {
            cur = first + sec;
            first = sec;
            sec = cur;
        }
        return cur;
    }
}
```

```

fib          PROC                                ; Start of subroutine
    push     ebp                                ; Save ebp
    mov      ebp, esp                            ; Copy ebp to esp
    sub      esp, 8                             ; Create space for first,
                                                ; sec (cur, if needed,
                                                ; will be in eax)

    push     ebx                                ; Save ebx and ecx (both
    push     ecx                                ; (overwritten in fn)

; CODE FOR: if (n <= 1) return n
    cmp      DWORD PTR 8[ebp], 1                ; Compare n to 1
    jg       L1                                 ; If n isn't <= 1, jump
                                                ; to else case
    mov      eax, DWORD PTR 8[ebp]              ; eax = n (eax holds
                                                ; return value)
    jmp      L3                                 ; Jump to end of function

; CODE FOR: first = 0; sec = 1
L1:
    mov      DWORD PTR -4[ebp], 0                ; first = 0
    mov      DWORD PTR -8[ebp], 1                ; sec = 1

; CODE FOR: loop initialization
; Note that the loop will execute n - 1 iterations, so we
; can initialize ECX to n - 1 and use loop instructions
    mov      ecx, DWORD PTR 8[ebp]              ; cx = n
    dec      ecx                                ; cx = cx - 1 = n - 1

; CODE FOR: cur = first + sec; first = sec; sec = cur
L2:
    mov      eax, DWORD PTR -4[ebp]              ; cur = eax = first
    add      eax, DWORD PTR -8[ebp]              ; cur = first + sec
    mov      ebx, DWORD PTR -8[ebp]              ; ebx = sec
    mov      DWORD PTR -4[ebp], ebx              ; first = ebx = sec
    mov      DWORD PTR -8[ebp], eax              ; sec = eax = cur

; CODE FOR: decrement loop counter & go to start of loop
    loop     L2

; CLEANUP (NOTE: No additional code needed for return cur
; in else case, since cur is already stored in eax)
L3:
    pop      ecx                                ; Restore ecx
    pop      ebx                                ; Restore ebx
    mov      esp, ebp                            ; Clear first, sec
    pop      ebp                                ; Restore ebp
    ret                                           ; Return from subroutine
fib          ENDP

```

```

int fib(int n) {          // FIBONACCI WITH RECURSION
    // For n == 0 or n == 1, fib(n) == n
    if (n <= 1) return n;

    // Otherwise, value is sum of two previous steps
    else return fib(n-1) + fib(n-2);
}

fib                PROC                ; Start of subroutine
    push          ebp                ; Save ebp
    mov           ebp, esp           ; Copy ebp to esp
    push          ebx                ; Save ebx (overwritten
                                    ; in function)

; CODE FOR: if (n <= 1) return n
    cmp           DWORD PTR 8[ebp], 1 ; Compare n to 1
    jg            L1                 ; If n isn't <= 1, jump
                                    ; to else case
    mov           eax, DWORD PTR 8[ebp] ; eax = n (eax holds
                                    ; return value)
    jmp           L2                 ; Jump to end of function

; CODE FOR: calling fib(n-1)
L1:
    mov           ebx, DWORD PTR 8[ebp] ; Copy n to ebx
    dec           ebx                 ; ebx = n - 1
    push          ebx                 ; Push n - 1 to pass it
                                    ; as argument
    call          fib                 ; Call fib(n-1)
                                    ; Return value in eax

; CODE FOR: calling fib(n-2)
; NOTE: We can take advantage of the fact that n-1 is still
; on the stack--decrement that value, and we'll have the
; value n-2 to pass to our next function call
    mov           ebx, eax             ; ebx = eax = fib(n-1)
    dec           DWORD PTR [esp]      ; Value at top of stack =
                                    ; (n-1) - 1 = n-2
    call          fib                 ; Call fib(n-2)
                                    ; Return value in eax

; CODE FOR: return fib(n-1) + fib(n-2)
    add           eax, ebx             ; eax = fib(n-1)+fib(n-2)

; CLEANUP
L2:
    add           esp, 4               ; Clear argument passed to
                                    ; fib(n-2)
    pop           ebx                 ; Restore ebx
    pop           ebp                 ; Restore ebp
    ret                                ; Return from subroutine
fib                ENDP

```